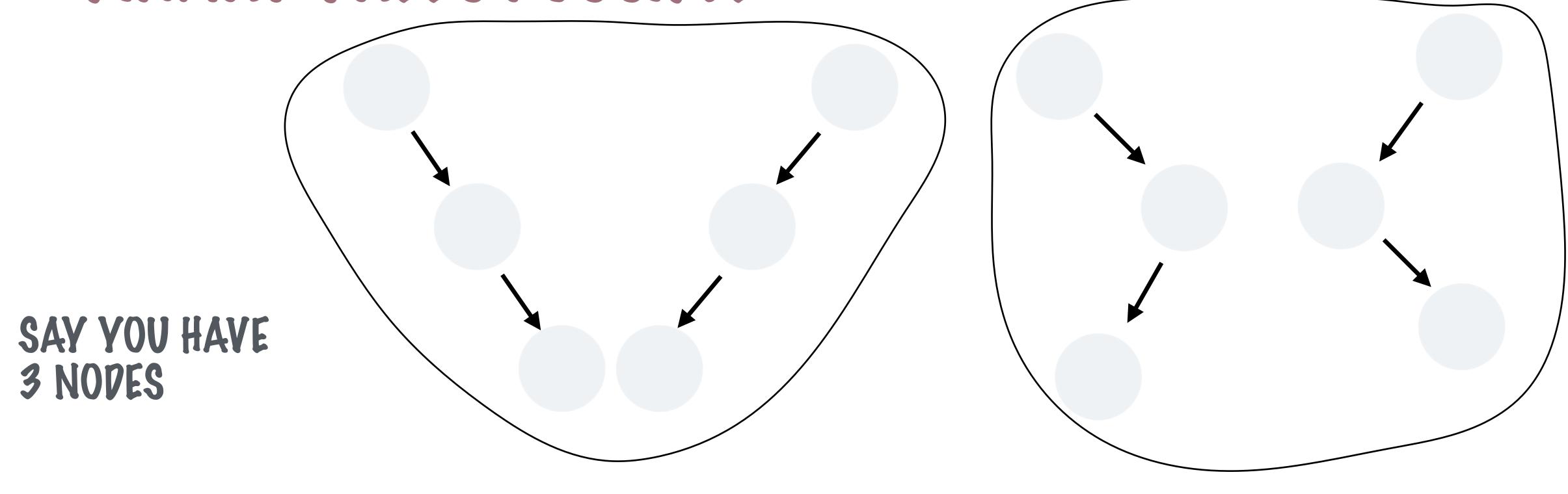
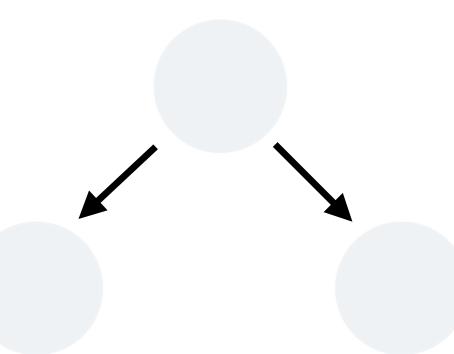
# COUNT THE NUMBER OF STRUCTURALLY UNIQUE BINARY TREES POSSIBLE

COUNT THE NUMBER OF STRUCTURALLY UNIQUE BINARY TREES POSSIBLE





THESE ARE THE VARIOUS POSSIBLE SHAPES OF TREES

NOTE THAT SOME OF THESE ARE MIRRORS OF EACH OTHER

#### COUNT TREES

```
public static int countTrees(int numNodes) {
    if (numNodes <= 1) {</pre>
        return 1;
    int sum = 0;
    for (int i = 1; i <= numNodes; i++) {</pre>
        int countLeftTrees = countTrees(i - 1);
        int countRightTrees = countTrees(numNodes - i);
        sum = sum + (countLeftTrees * countRightTrees):
    return sum;
```

THIS IS THE NUMBER OF POSSIBLE TREES WITH THIS ROOT - THE COMBINATION OF RIGHT AND LEFT SUBTREES

WHEN THE NUMBER OF NODES IS 1 THERE IS JUST ONE POSSIBLE TREE - THIS IS THE BASE CASE

CONSIDER THAT EVERY NODE CAN BE THE ROOT - THE NODES BEFORE IT WILL BE ON THE LEFT AND THE NODES AFTER IT ON THE RIGHT

NOPES ON THE LEFT AND RIGHT FORM THEIR OWN SUBTREES

## PRINT ALL NODES WITHIN A RANGE IN A BINARY SEARCH TREE

## PRINT ALL NODES WITHIN A RANGE IN A BINARY SEARCH TREE

A RANGE WILL INCLUDE A SUBSET OF NODES IN THE BINARY SEARCH TREE THIS SUBSET CAN INCLUDE O NODES AS WELL

CHECK EVERY NODE TO SEE IF IT'S VALUE IS WITHIN THE RANGE - PRINT IT TO SCREEN IF THE RANGE CONSTRAINTS ARE MET

#### PRINT NOPES WITHIN A RANGE

```
public static void printRange(Node<Integer> root, int low, int high) {
   if (root == null) {
      return;
   }

   if (low <= root.getData()) {
      printRange(root.getLeftChild(), tow. high);
   }

   if (low <= root.getData() && root.getData() <= high) {
      System.out.println(root.getData());
   }

   if (high > root.getData()) {
      printRange(root.getRightChild(), low, high);
   }
}
```

### PASS IN THE MIN AND MAX INDICATING THE RANGE WE CARE ABOUT

BASE CASE, NOTHING TO PO FOR A NULL ROOT

IF THE RANGE LOW VALUE IS LESS THAN THE CURRENT NODE, RUN THE OPERATION ON THE LEFT SUBTREE

CHECK THE NODE VALUE TO SEE IF IT'S WITHIN THE RANGE - IF YES, PRINT!

IF THE RANGE HIGH VALUE IS GREATER THAN THE CURRENT NODE, RUN THE OPERATION ON THE RIGHT SUBTREE

### CHECK IF A BINARY TREE IS A BINARY \*SEARCH\* TREE

#### CHECK IF A BINARY TREE IS A BINARY \*SEARCH\* TREE

FOR EVERY NODE IN A BINARY SEARCH TREE - THE NODES WITH VALUES <= NODE ARE IN THE LEFT SUBTREE AND NODES WITH VALUES > NODE ARE IN THE RIGHT SUBTREE

CHECK EVERY NODE TO SEE IF THIS CONSTRAINT IS VIOLATED

> THIS CAN BE SOLVED ITERATIVELY AND RECURSIVELY

#### IS BINARY SEARCH TREE?

```
PASS IN THE MIN AND MAX INDICATING THE RANGE FOR THE SUBTREE
```

```
public static boolean isBinarySearchTree(Node<Integer> root, int min, int max) {
   if (root == null) {
      return true;
   }

if (root.getData() <= min || root.getData() > max) {
      return false;
   }

return isBinarySearchTree(root.getLeftChild(), min, root getData())
   && isBinarySearchTree(root.getRightChild(), root.getData(), max);
}
```

A NULL NOPE IS A VALIP BINARY TREE

IF ANY NODE LIES OUTSIDE
THE RANGE THEN THE BST
CONSTRAINT HAS BEEN
VIOLATED AND WE RETURN
FALSE

FOR THE LEFT SUBTREE THE CURRENT NODE'S VALUE SHOULD BE THE MAX

FOR THE RIGHT SUBTREE THE CURRENT NODE'S VALUE SHOULD BE THE MIN

CHECK THE LEFT AND RIGHT SUBTREES TO SEE IF THEY'RE VALID SEARCH TREES - NOTE HOW THE RANGE FOR THE CHECKS CHANGE