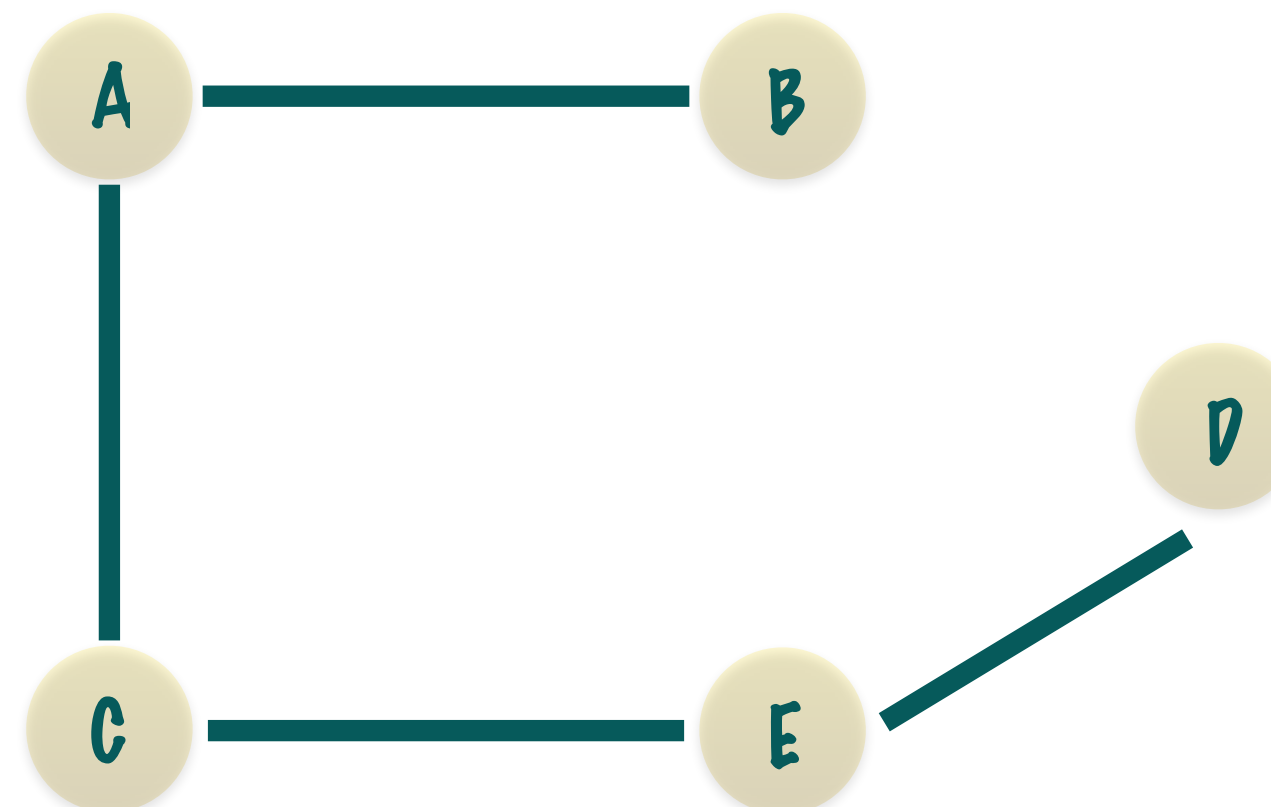
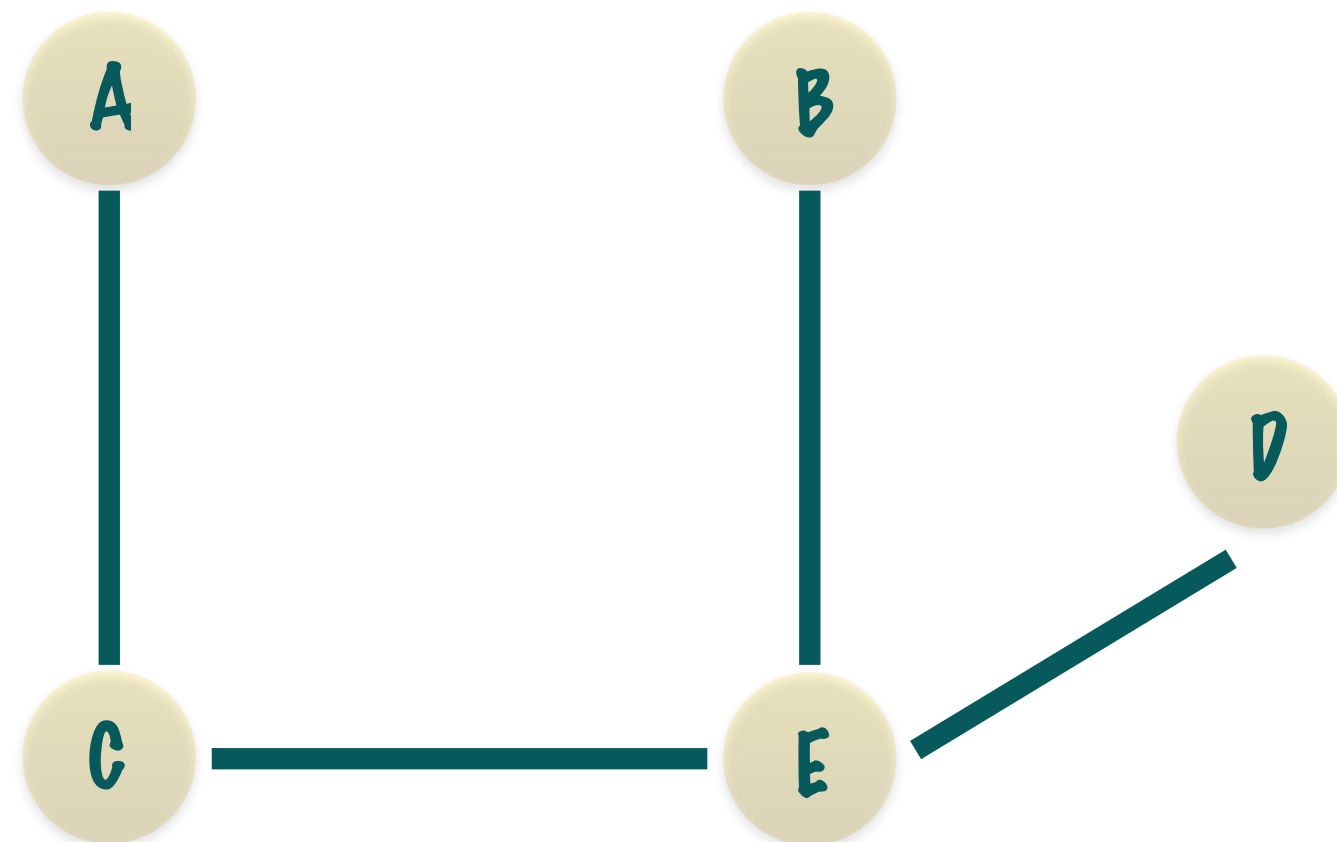
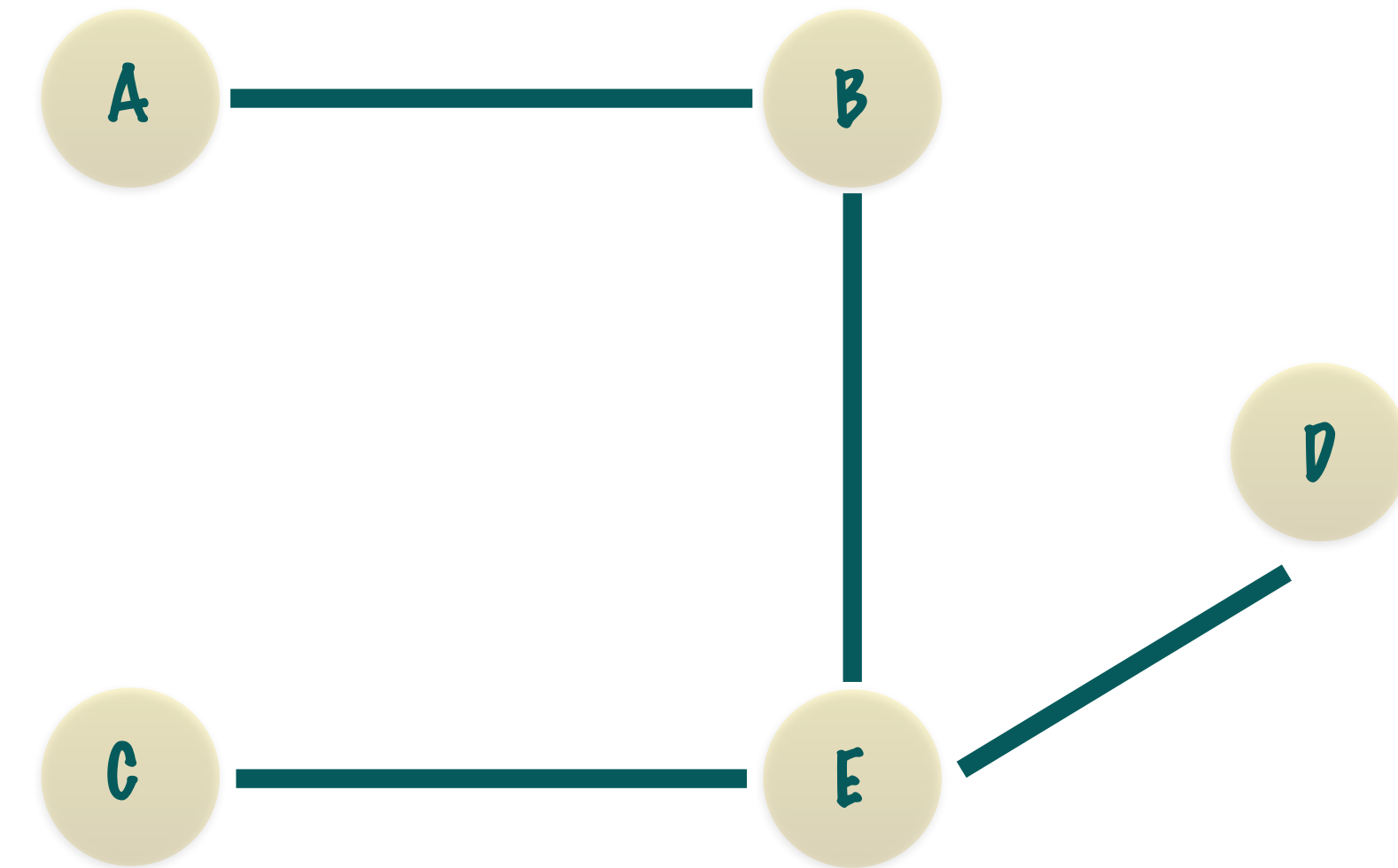
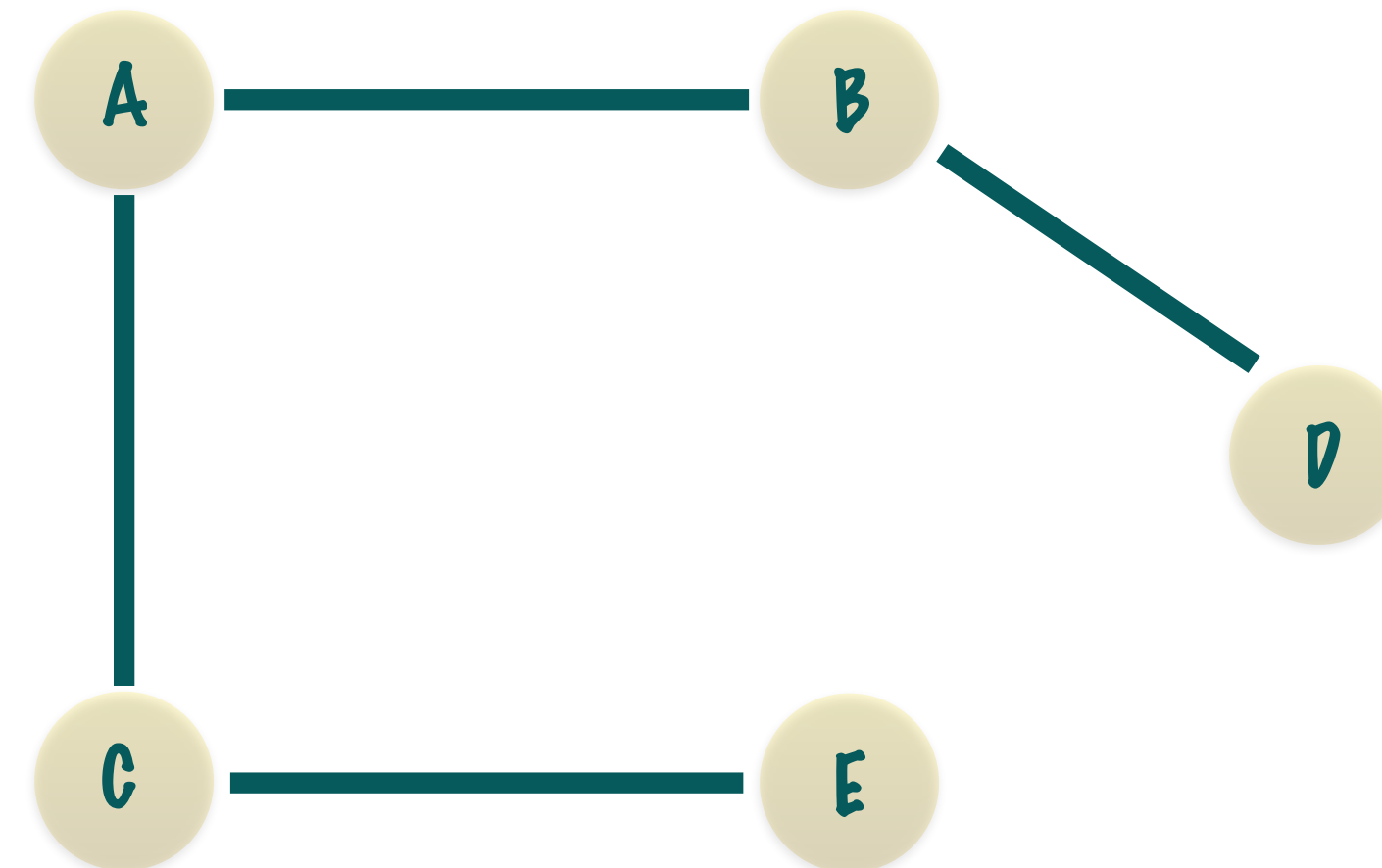
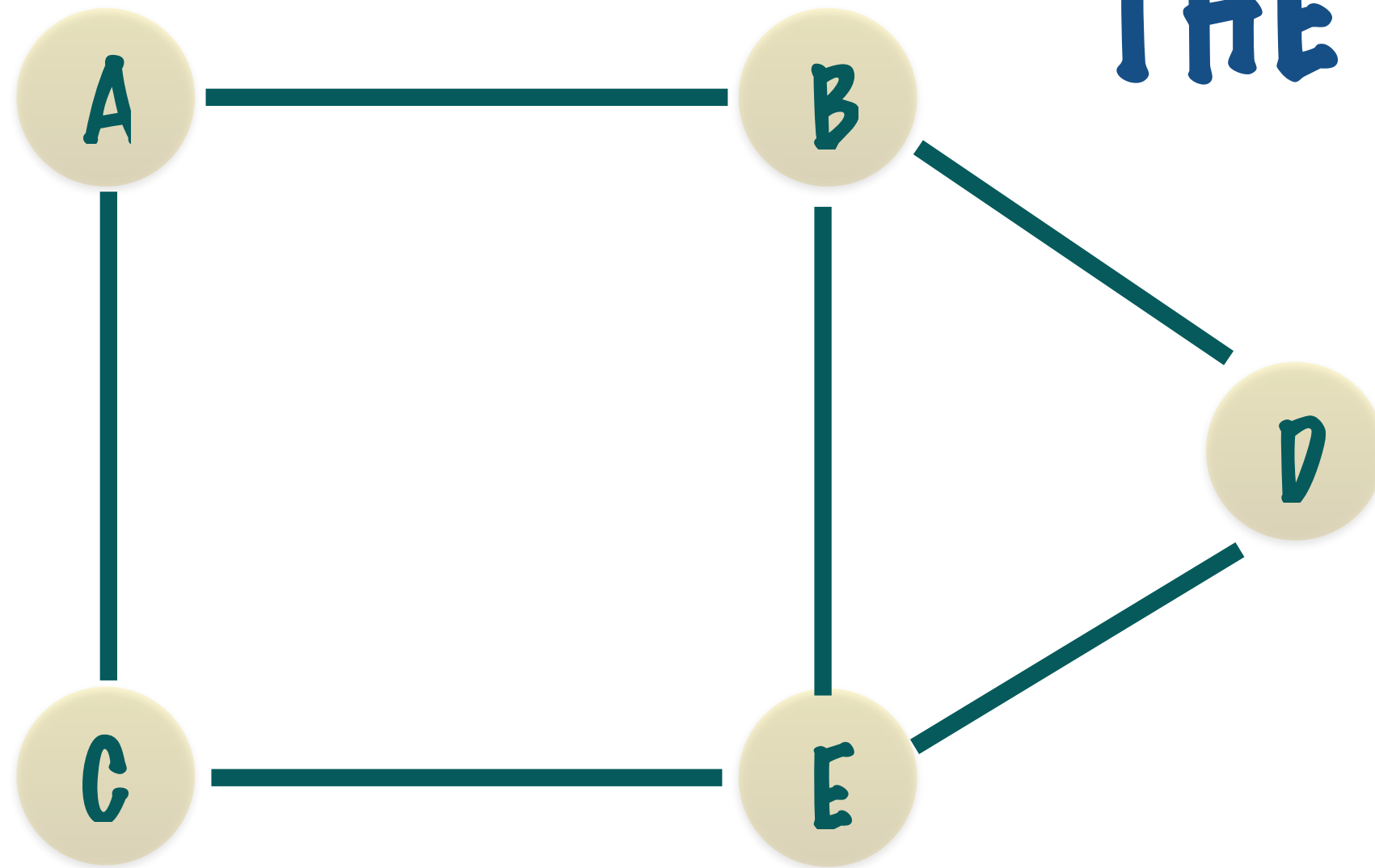


# THE GRAPH

MINIMAL SPANNING TREE FOR FORESTS

# THE GRAPH SPANNING TREE

SPANNING TREE IS A SUBGRAPH THAT CONTAINS ALL THE VERTICES AND IS ALSO A TREE



# THE GRAPH MINIMAL SPANNING TREE FOR FORESTS

A FOREST IS AN UNCONNECTED GRAPH

## KRUSKAL'S ALGORITHM

THIS IS A 'GREEDY ALGORITHM' IT TRIES TO  
FIND THE OPTIMAL NEXT STEP AT EVERY STEP -  
A LOCAL OPTIMUM NOT THE GLOBAL OPTIMUM

AT EVERY STEP WE CHOOSE THE SMALLEST  
WEIGHTED EDGE FROM THE ENTIRE GRAPH

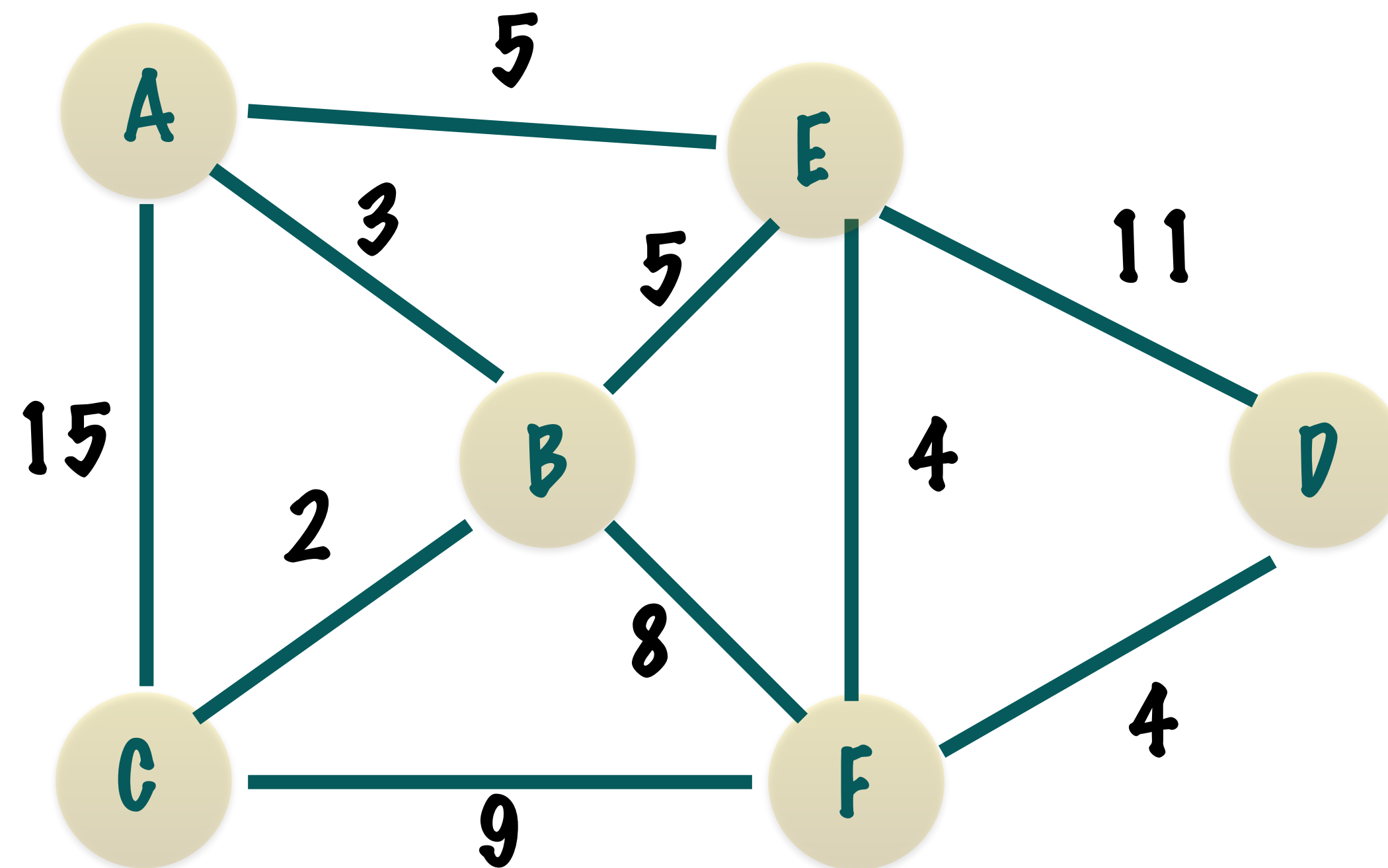
2 THINGS TO KEEP IN MIND WHILE  
IMPLEMENTING KRUSKAL'S ALGORITHM

# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS

1

USE A PRIORITY QUEUE OF EDGES WHERE THE WEIGHTS OF THE EDGES DETERMINE THE PRIORITY OF THE EDGE



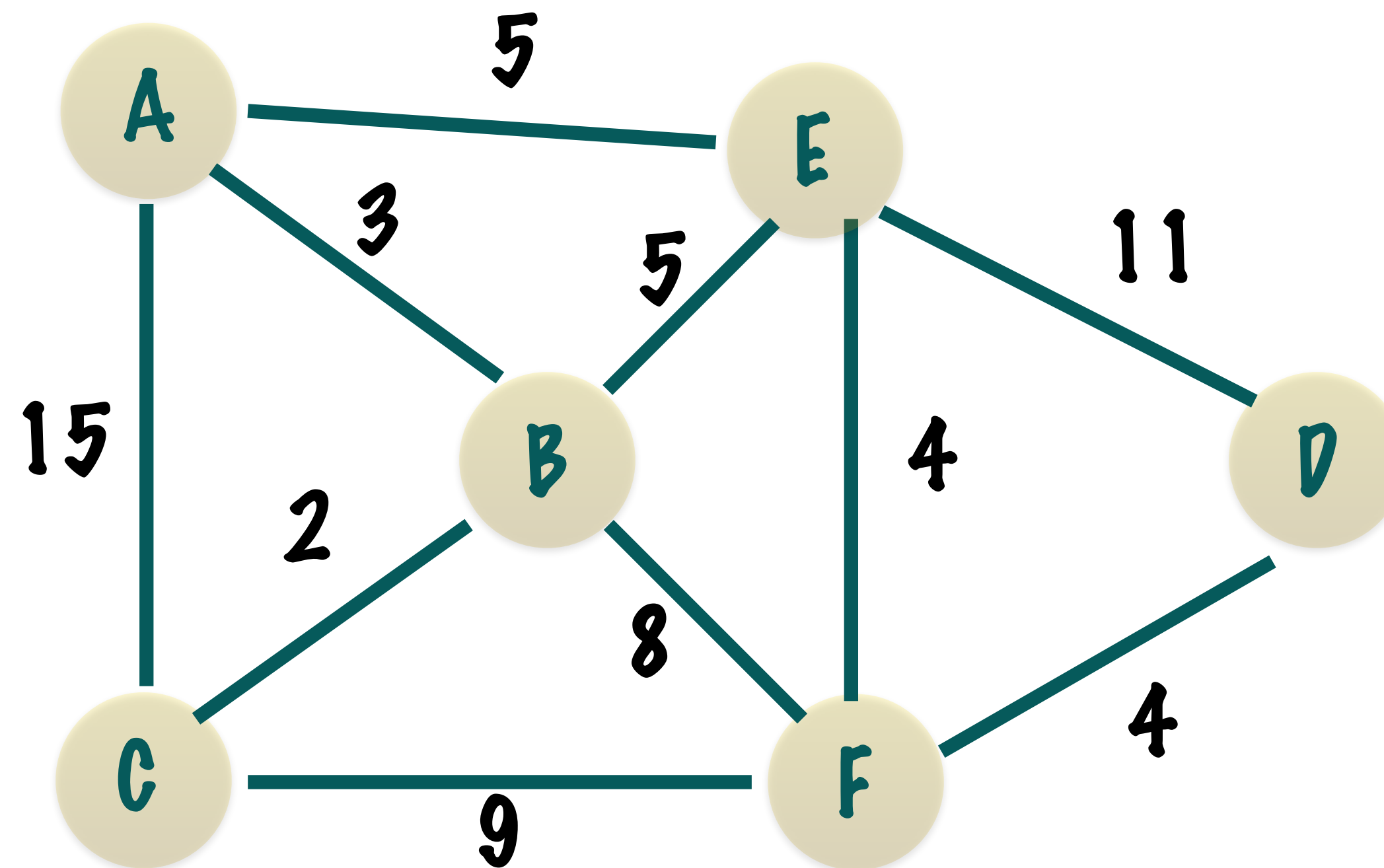
# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS

2

WHILE ADDING A NEW EDGE, ALWAYS MAKE SURE THAT THE NEW EDGE DOES NOT CREATE A CYCLE IN THE SPANNING TREE

CONTINUE ADDING EDGES TILL WE GET  $V - 1$  EDGES SO THE GRAPH IS CONNECTED I.E. IT'S A TREE



# THE GRAPH

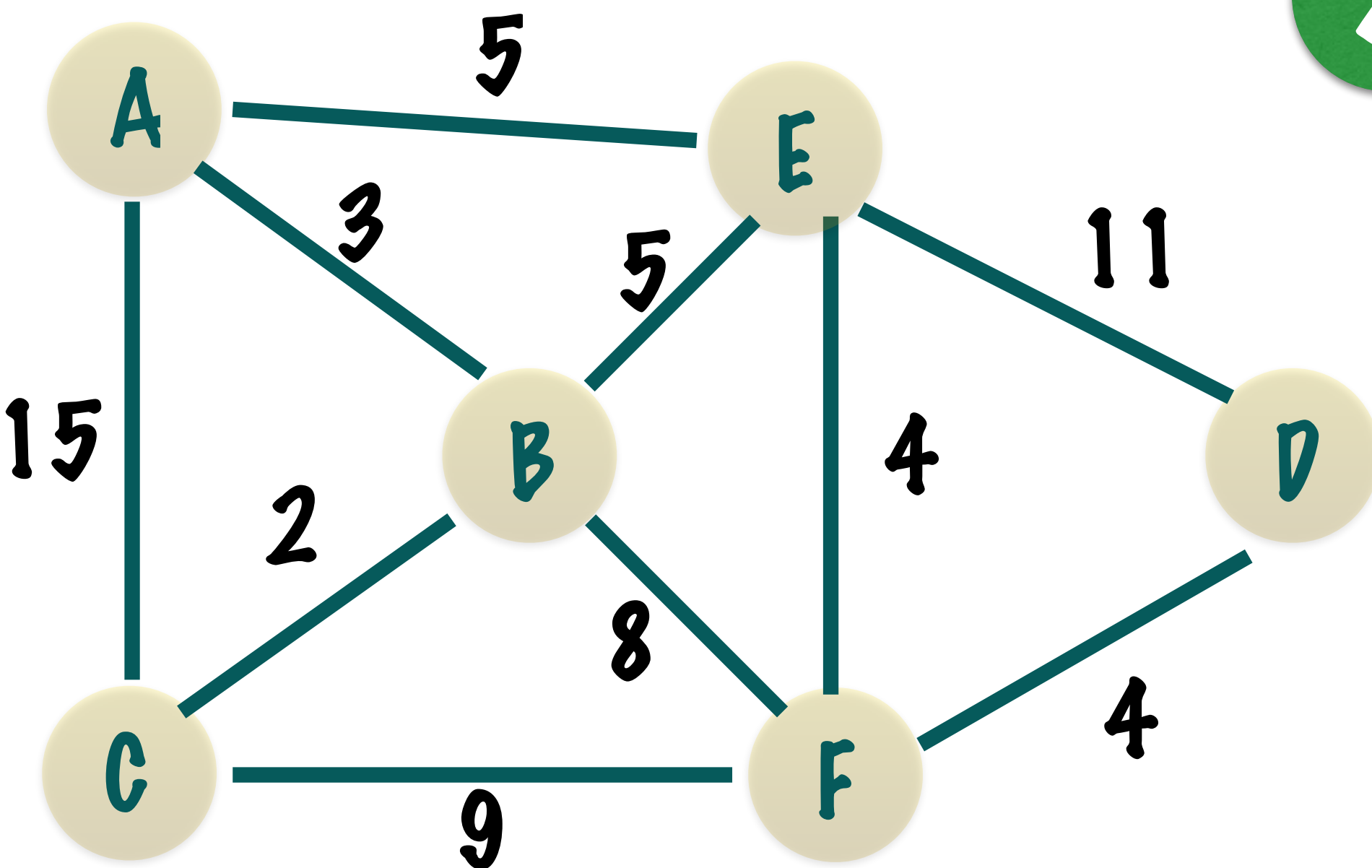
## MINIMAL SPANNING TREE FOR FORESTS

1

USE A PRIORITY QUEUE OF EDGES WHERE THE WEIGHTS OF THE EDGES DETERMINE THE PRIORITY OF THE EDGE

2

WHILE ADDING A NEW EDGE, ALWAYS MAKE SURE THAT THE NEW EDGE DOES NOT CREATE A CYCLE IN THE SPANNING TREE



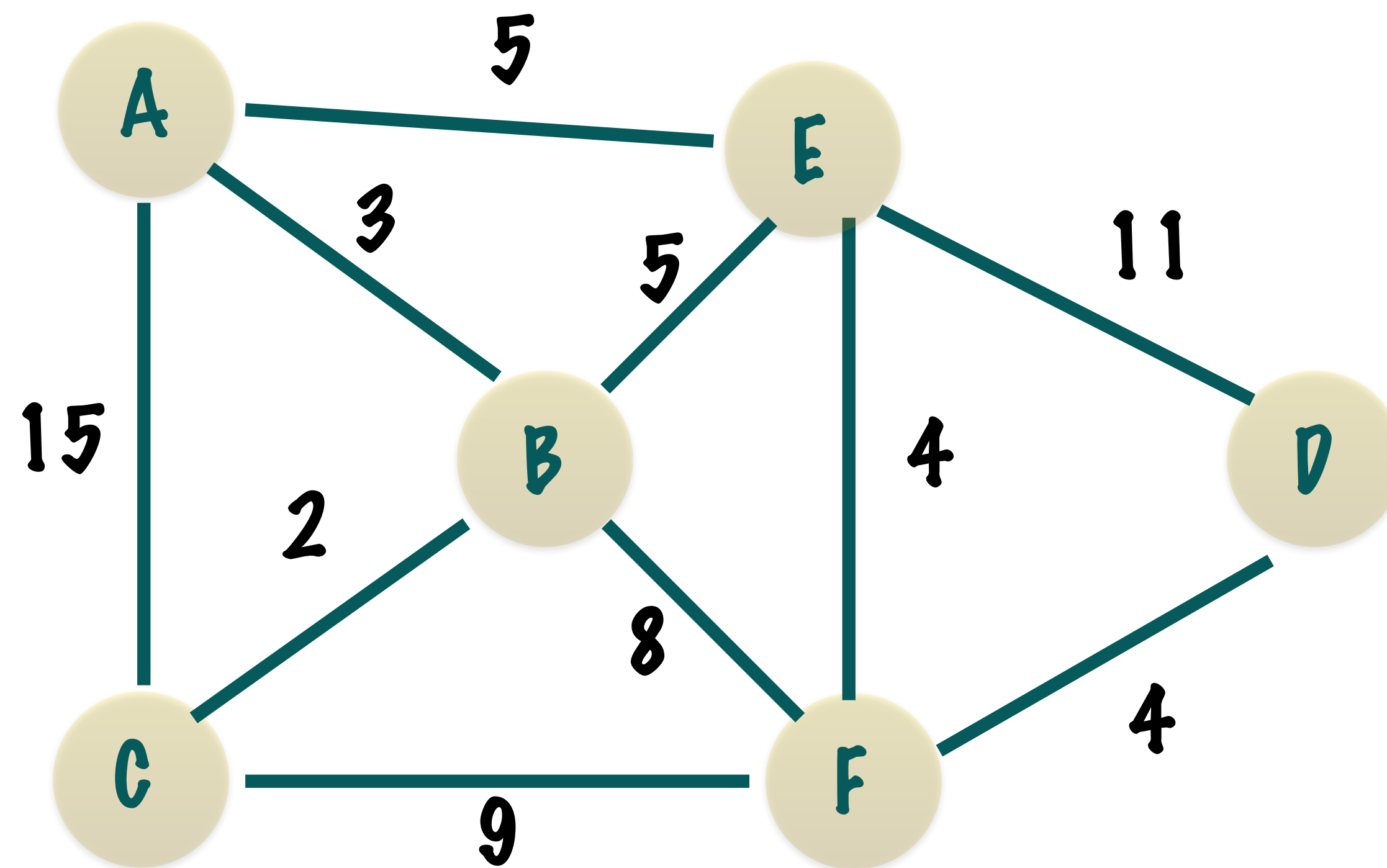


# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS

PRIORITY QUEUE OF EDGES

PRIORITY = WEIGHT OF EDGE



BC 2

AB 3

DF 4

EF 4

BE 5

AE 5

BF 8

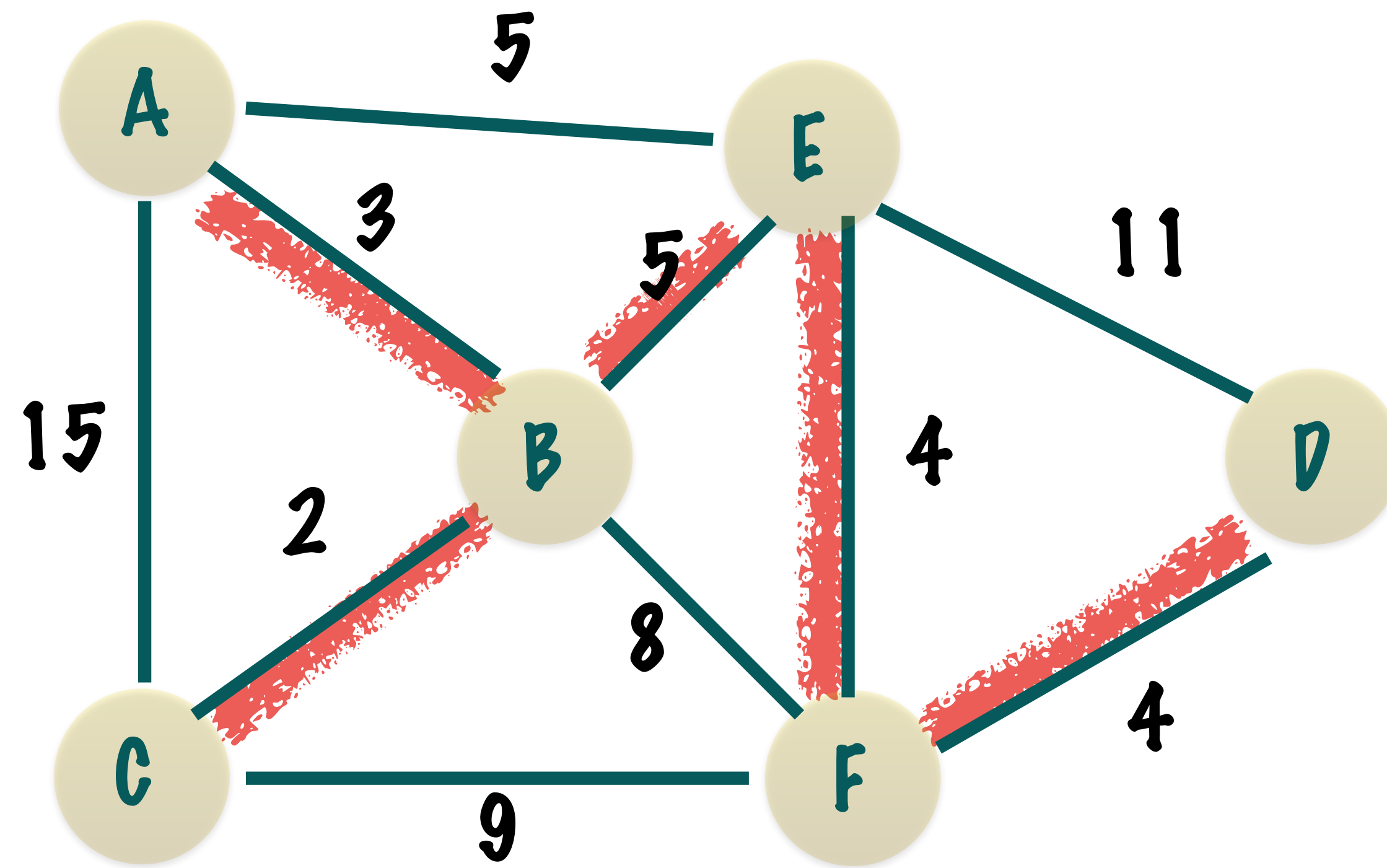
CF 9

ED 11

AC 15

# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS



WE HAVE COVERED ALL THE VERTICES!

BC 2

AB 3

DF 4

EF 4

BE 5

AE 5

BF 8

CF 9

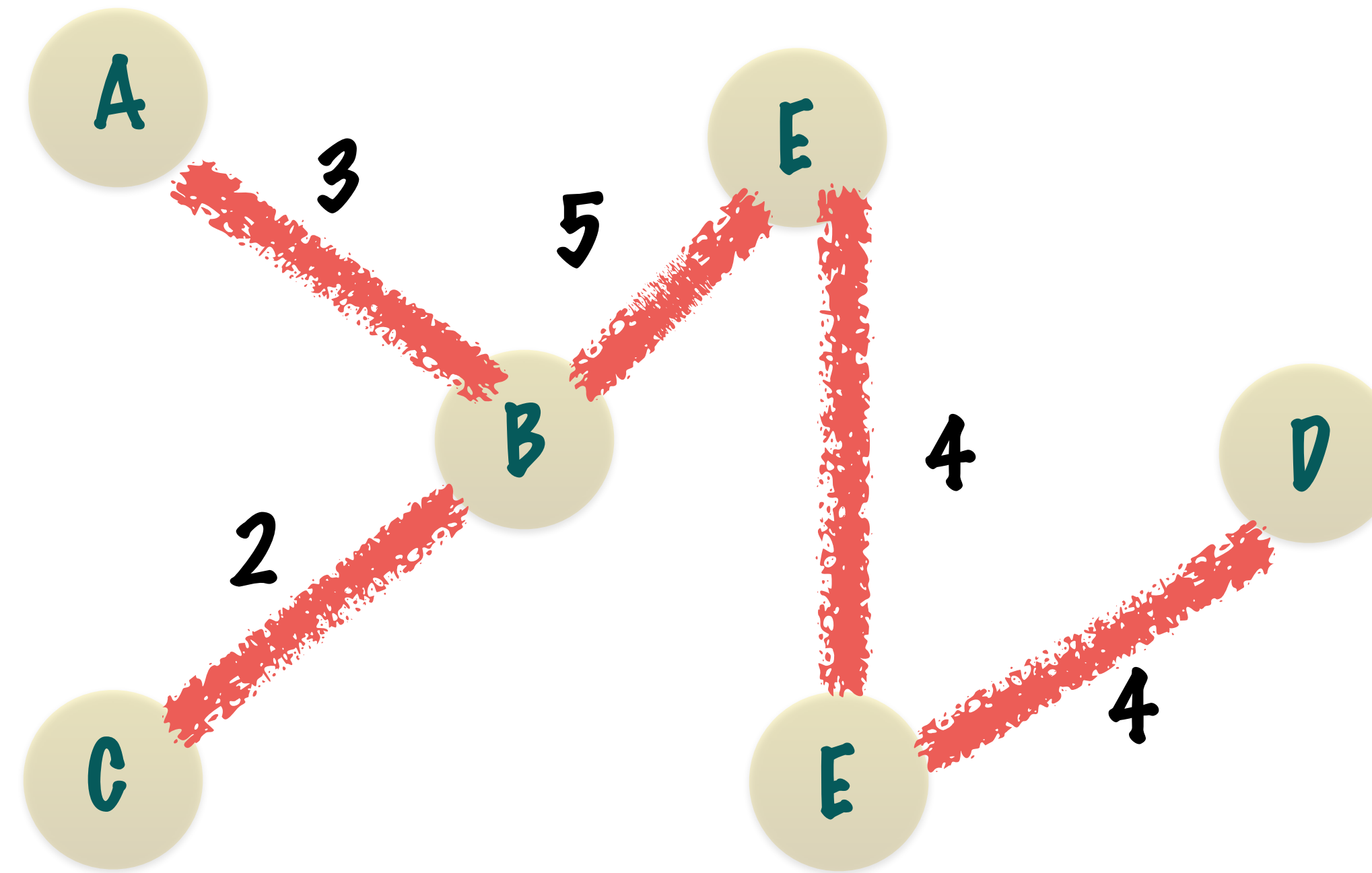
ED 11

AC 15



# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS



THIS ALGORITHM WORKS FOR BOTH  
CONNECTED AND UNCONNECTED GRAPHS  
I.E FORESTS

# THE GRAPH

## MINIMAL SPANNING TREE FOR FORESTS

THE ALGORITHM'S RUNNING  
TIME IS  $E \lg E$ !

THE MAIN PROCESSING TIME INVOLVES  
SORTING THE EDGES BY WEIGHT AND  
THIS IS THE RUNNING TIME OF THE BEST  
SORTING ALGORITHMS

# EDGE INFO DATA STRUCTURE

```
/**
 * A class which represents an edge in an undirected weighted graph.
 */
public static class EdgeInfo {

    private Integer vertex1;
    private Integer vertex2;
    private Integer weight;

    public EdgeInfo(Integer vertex1, Integer vertex2, Integer weight) {
        this.vertex1 = vertex1;
        this.vertex2 = vertex2;
        this.weight = weight;
    }

    public Integer getVertex1() {
        return vertex1;
    }

    public Integer getVertex2() {
        return vertex2;
    }

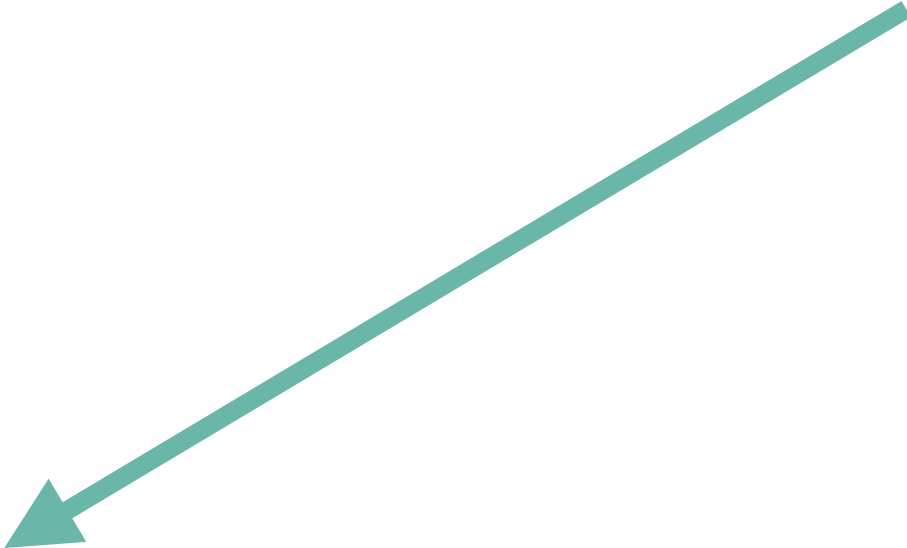
    public Integer getWeight() {
        return weight;
    }

    @Override
    public String toString() {
        return String.valueOf(vertex1) + String.valueOf(vertex2);
    }
}
```

REPRESENTS AN EDGE USING  
THE TWO VERTICES AND THE  
EDGE WEIGHT



STRING REPRESENTATION OF AN  
EDGE WHICH IS "02" FOR AN EDGE  
WHICH CONNECTS VERTEX 0 WITH  
VERTEX 2





# BUILD THE EDGE MAP AND SPANNING TREE- SETUP

```
static void spanningTree(Graph graph) {  
    // A priority queue to store and retrieve the edges on the basis of their  
    // weights.  
    PriorityQueue<EdgeInfo> queue = new PriorityQueue<>(new Comparator<EdgeInfo> () {  
        @Override  
        public int compare(EdgeInfo o1, EdgeInfo o2) {  
            return o1.getWeight().compareTo(o2.getWeight());  
        }  
    });  
  
    // Add all edges to the priority queue.  
    for (int i= 0; i < graph.getNumVertices(); i++) {  
        for (int neighbour : graph.getAdjacentVertices(i)) {  
            queue.add(new EdgeInfo(i, neighbour, graph.getWeightedEdge(i, neighbour)));  
        }  
    }  
  
    Set<Integer> visitedVertices = new HashSet<>();  
    Set<EdgeInfo> spanningTree = new HashSet<>();  
    Map<Integer, Set<Integer>> edgeMap = new HashMap<>();  
    for (int v = 0; v < graph.getNumVertices(); v++) {  
        edgeMap.put(v, new HashSet<>());  
    }  
}
```

SET UP A PRIORITY QUEUE  
WHICH RETURNS EDGES WITH  
THE SMALLEST WEIGHT - "THE  
GREEDY SOLUTION"

ADD EVERY EDGE TO THE  
PRIORITY QUEUE

KEEP TRACK OF THE VERTICES  
ALREADY VISITED, EACH EDGE  
SHOULD ADD A NEW VERTEX TO  
THE SET TILL WE GET NUMBER OF  
VERTICES - 1 EDGES

THE EDGE MAP TRACKS THE EDGES ADDED TO  
THE SPANNING TREE TO SEE IF IT FORMS A  
CYCLE

THE SPANNING TREE IS THE SET OF EDGES  
CONNECTING ALL THE NODES OF THE GRAPH,  
AN EDGE IS REPRESENTED BY "0 1" IF IT  
CONNECTS VERTICES 0 AND 1



# BUILD THE EDGE MAP AND SPANNING TREE - PROCESS

```
while(!queue.isEmpty() && spanningTree.size() < graph.getNumVertices() - 1) {  
    EdgeInfo currentEdge = queue.poll();  
  
    // Add the new edge to the edge map and see if it ends up with a cycle.  
    // If yes then discard this edge and get the next edge from the priority  
    // queue.  
    edgeMap.get(currentEdge.getVertex1()).add(currentEdge.getVertex2());  
    if (hasCycle(edgeMap)) {  
        edgeMap.get(currentEdge.getVertex1()).remove(currentEdge.getVertex2());  
        continue;  
    }  
  
    spanningTree.add(currentEdge);  
  
    // Add both vertices to the visited list, the set will ensure  
    // that only one copy of the vertex exists.  
    visitedVertices.add(currentEdge.getVertex1());  
    visitedVertices.add(currentEdge.getVertex2());  
}  
  
// Check whether all vertices have been covered with the spanning tree.  
if (visitedVertices.size() != graph.getNumVertices()) {  
    System.out.println("Minimum Spanning Tree is not possible");  
} else {  
    System.out.println("Minimum Spanning Tree using Kruskal's Algorithm");  
    for(EdgeInfo edgeInfo : spanningTree) {  
        System.out.println(edgeInfo);  
    }  
}
```

THE SPANNING TREE SHOULD HAVE  
NUMBER OF VERTICES - 1 EDGES

RETRIEVE THE EDGES WITH THE  
SMALLEST WEIGHT FIRST - THE  
GREEDY SOLUTION

ADD THE EDGE TO THE EDGE MAP  
AND SEE IF IT CAUSES A CYCLE - IF  
YES THEN DO NOT USE THE EDGE  
IN THE SPANNING TREE

ADD THE EDGE TO THE SPANNING  
TREE

ADD BOTH VERTICES TO THE VISITED  
VERTEX LIST

IF ALL VERTICES HAVE BEEN COVERED  
THE SPANNING TREE EXISTS!



# CHECK FOR CYCLES IN THE SPANNING TREE

```
private static boolean hasCycle(Map<Integer, Set<Integer>> edgeMap) {  
    for (Integer sourceVertex : edgeMap.keySet()) {  
        LinkedList<Integer> queue = new LinkedList<>();  
        queue.add(sourceVertex);  
        Set<Integer> visitedVertices = new HashSet<>();  
        while (!queue.isEmpty()) {  
            int currentVertex = queue.pollFirst();  
            if (visitedVertices.contains(currentVertex)) {  
                return true;  
            }  
  
            visitedVertices.add(currentVertex);  
            queue.addAll(edgeMap.get(currentVertex));  
        }  
    }  
  
    return false;  
}
```

START FROM EVERY VERTEX IN THE  
EDGE MAP AND EXPLORE ALL  
VERTICES PRESENT IN THE SPANNING  
TREE

IF WE EVER RE-VISIT A VERTEX  
WE'VE ALREADY SEEN IN THE  
SPANNING TREE IT MEANS THERE IS  
A CYCLE IN THE SPANNING TREE

ADD TO QUEUE ALL THE ADJACENT  
VERTICES OF THE CURRENT  
VERTEX WHICH ARE PART OF THE  
SPANNING TREE