LINKEP LISTS

ARRAYS

LIST SIZES ARE UNLIMITED, YOU DON'T NEED TO KNOW UP FRONT HOW MANY ELEMENTS WILL BE ADDED TO THE LINKED LIST, THE SIZE OF THE LINKED LIST CAN GROW DYNAMICALLY

ARRAYS HAVE TO BE DECLARED UPFRONT WITH THE NUMBER OF ELEMENTS IT WILL HOLD, THE SIZE OF THE ARRAY CANNOT BE INCREASED DYNAMICALLY

LINKEP LISTS

ARRAYS

INSERTING A NEW ELEMENT IN A LINKED LIST IS A VERY <u>EASY</u>
OPERATION, THE LOGICAL NEXT POINTERS HAVE TO BE REASSIGNED BUT NOT MUCH ELSE NEEDS TO BE DONE

ARRAYS ARE LOCATED IN CONTIGUOUS MEMORY LOCATIONS, IN ORDER TO INSERT AN ELEMENT, THE ELEMENTS TO IT'S RIGHT HAVE TO MOVE OVER TO MAKE ROOM FOR IT. IT IS A MORE HEAVY WEIGHT OPERATION

SIMILARLY PELETING AN ELEMENT IN A LIST IS VERY EASY AND EFFICIENT

ARRAY ELEMENTS HAVE TO BE MOVED IN THE CASE OF DELETION AS WELL

LINKEP LISTS

ARRAYS

RANDOM ACCESS TO AN ELEMENT AT A SPECIFIC INDEX IN THE LINKED LIST IS NOT POSSIBLE. THE ENTIRE LIST UP-TO THAT ELEMENT HAS TO BE TRAVERSED

ARRAYS PROVIDE VERY QUICK LOOKUP FOR ELEMENTS AT SPECIFIC INDICES, SINCE THEY ARE IN CONTIGUOUS MEMORY LOCATIONS, WE KNOW EXACTLY WHERE IN MEMORY THE ELEMENT IS

LINKEP LISTS

ARRAYS

EACH ELEMENT REQUIRES APPITIONAL SPACE TO STORE A POINTER TO THE NEXT ELEMENT

NO EXTRA SPACE IS REQUIRED OTHER THAN FOR THE ACTUAL ELEMENTS WHICH MAKE UP THE ARRAY

LINKEP LISTS

ARRAYS

CANNOT TAKE ADVANTAGE OF SPATIAL LOCALITY (FOR CACHING) WHEN ACCESSING THE ELEMENTS

EACH ELEMENT CAN LIVE PRETTY MUCH ANYWHERE IN MEMORY AND BE POINTED TO AS ELEMENTS ARE IN CONTIGUOUS MEMORY LOCATIONS ACCESS TO ARRAYS TAKES SIGNIFICANT ADVANTAGE OF SPATIAL LOCALITY IN CACHES

THIS CAN BE A SIGNIFICANT PERFORMANCE IMPROVEMENT

USE LINKEP LISTS WHEN:

YOU HAVE A LARGE NUMBER OF INSERT OR PELETE OPERATIONS TO PERFORM

YOUR LIST MIGHT BE

USE ARRAYS WHEN:

READ OPERATIONS NEED TO BE EXTREMELY FAST AND YOU HAVE RELATIVELY FEW UPDATES TO THE ARRAY

YOU REQUIRE RANDOM ACCESS TO ARRAY ELEMENTS