

Lecture #4

Sort (1)

Algorithm

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Jinhong Jung

In Previous Lecture

Recursive complexity

- Complexity of a recursive algorithm is also represented recursively as recurrence relation

How to analyze a recursive complexity function

- **Repeated substitution**
 - Repeatedly substitute the complexity function whose input size decreases toward a base case
- **Mathematical induction**
 - Estimate the closed form of a recursive complexity, and then prove it by induction
- **Master theorem**
 - Can solve any function in form of $T(n) = aT(n/b) + f(n)$
- For some cases, changing variables makes an equation simple

In This Lecture

Sorting problem

Basic sorting algorithms

- Selection Sort
- Bubble Sort
- Insertion Sort

Outline

Sorting problem

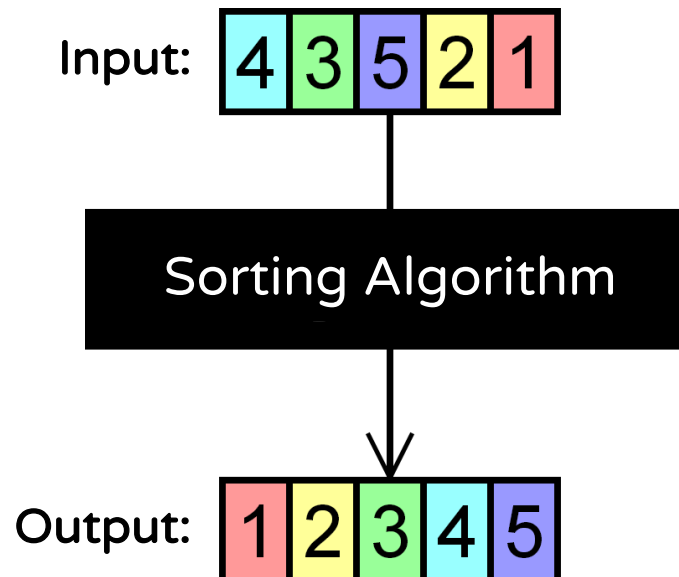
Basic sorting algorithms

- Selection Sort
- Bubble Sort
- Insertion Sort

Sorting Problem

What is sort in CS?

- To rearrange elements in an array in an order
 - If it contains numbers, an order is numerical order
 - If it contains characters, an order is alphabetical order
- **Sorting algorithms** aim to efficiently sort an arbitrary array in a certain order



Classical Sorting Algorithms

Comparison based sorting

- **Basic sorting algorithms**

- Show $\Theta(n^2)$ time complexity for a worst case
- Selection sort, Bubble sort, Insertion sort

- **Advanced sorting algorithms**

- Show $\Theta(n \log n)$ time complexity for a worst case
- Merge sort, Quick sort, Heap sort

Non-comparison based sorting

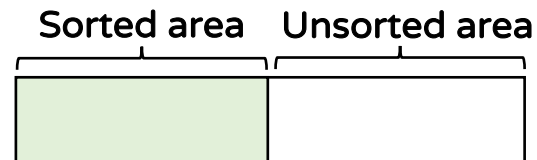
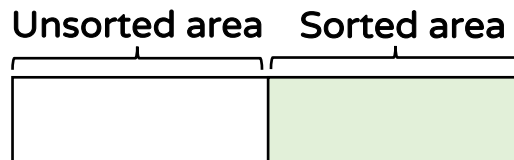
- **Special sorting algorithms**

- Show $\Theta(n)$ time complexity for special cases (not general)
- Radix sort, Counting sort

Before Going Further

Notes

- The index of an array A starts from 1
- An array A will be sorted in the ascending order
 - $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow \dots$
- $A[1 \dots i]$ indicates elements from index 1 to i
- Sorted area (colored by light green)
 - Consecutive sorted partial part in an array while it is sorted by a sorting algorithm
- Unsorted area (colored by white)
 - Remaining part except the sorted area



Outline

Sorting problem

Basic sorting algorithms

- Selection Sort 

- Bubble Sort

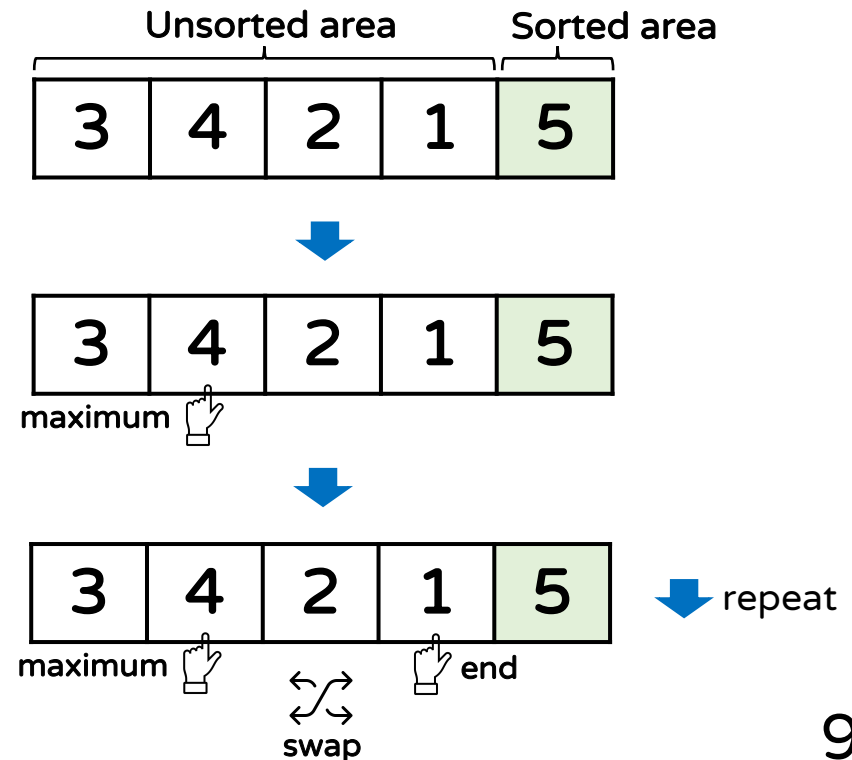
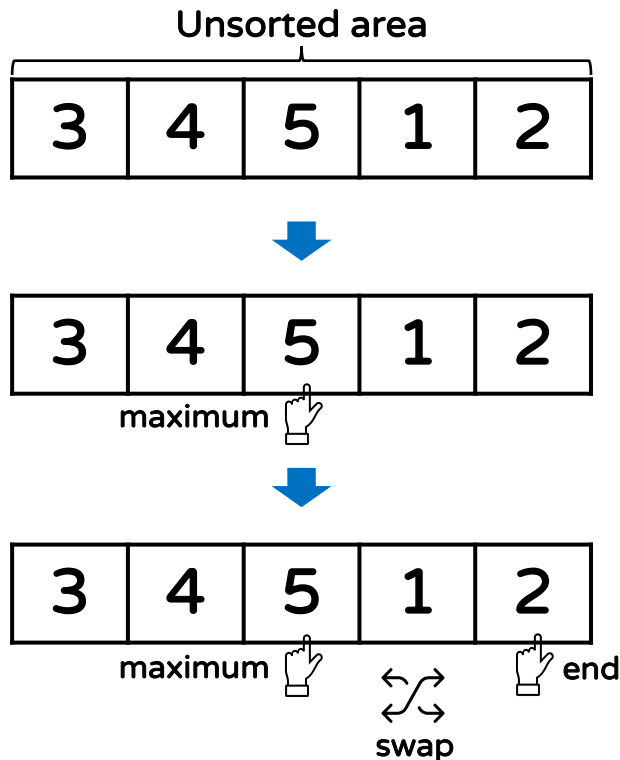
- Insertion Sort

Selection Sort (1)

Idea of selection sort

Or move the minimum
to the front

- Move the maximum to the end of unsorted area
 - Find the maximum & swap it with the end of unsorted area
 - Repeat these until the unsorted area becomes empty



Selection Sort (2)

Pseudocode of selection sort

```
def selection_sort(A, n):  
    for end ← n downto 2:  
        # at this step, unsorted area is [1, end]  
  
        # selection stage  
        k ← find the maximum's index among A[1...end]  
        swap A[k] and A[end]  
  
        # at this step, unsorted area is [1, end-1]
```

▪ Correctness

- Obviously true since for each step, the sorted area is expanded correctly, and finally its size becomes n at the end

▪ T.M.I. : $k \leftarrow \text{find the maximum's index among } A[1 \dots \text{end}]$

$$k = \underset{i \in [1, \text{end}]}{\operatorname{argmax}} A[i]$$

Selection Sort (3)

Time complexity of selection sort

```
def selection_sort(A, n):  
    for end ← n downto 2:  
        # selection stage  
        k ← find the maximum's index among A[1...end]  
        swap A[k] and A[end]
```

} selection stage

■ Time complexity analysis

- The for-loop repeats the selection process $n - 1$ times
- In the selection stage,
 - Main operations are comparisons for linearly searching the maximum's index
 - $i - 1$ comparisons are required for i -th selection process

$$T(n) = (n - 1) + (n - 2) + \dots + 2 + 1 = \sum_{i=1}^{n-1} i = \frac{n(n - 1)}{2} \in \Theta(n^2)$$

Outline

Sorting problem

Basic sorting algorithms

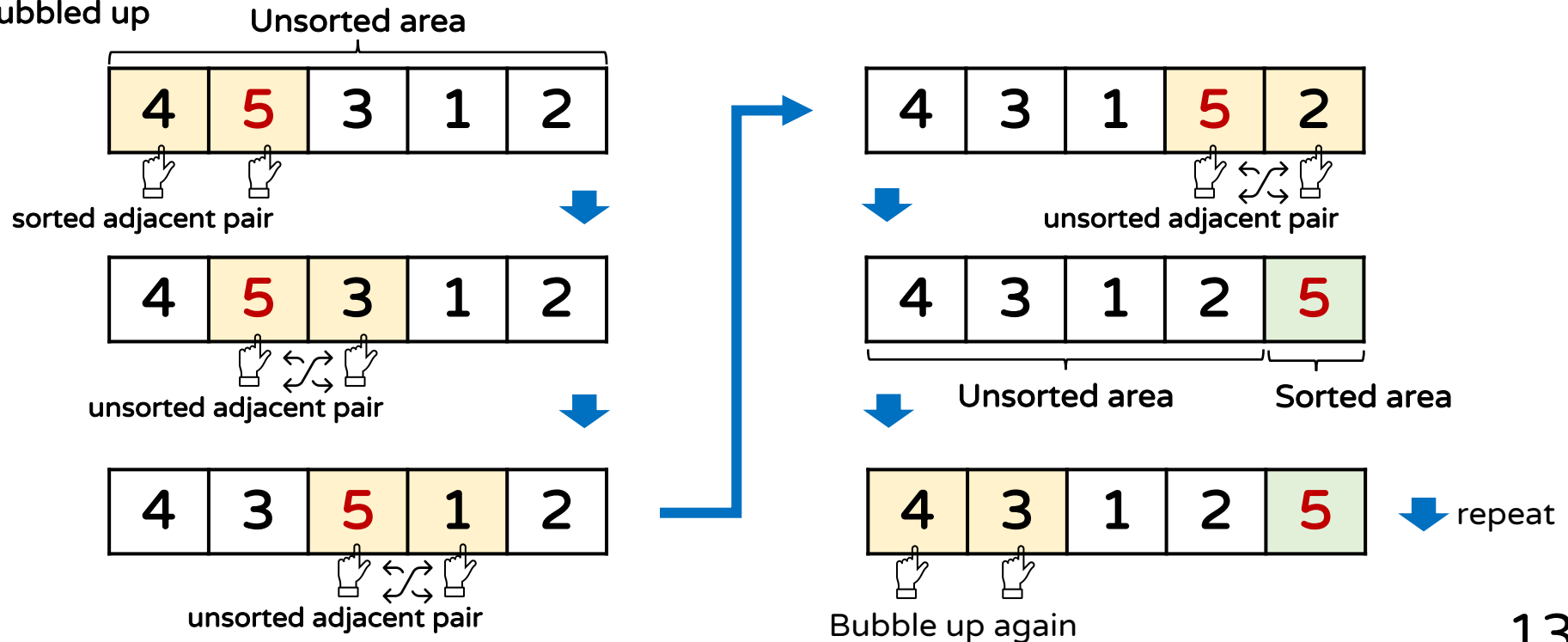
- Selection Sort
- Bubble Sort 
- Insertion Sort

Bubble Sort (1)

Idea of bubble sort

- Move the maximum to the end of unsorted area
 - Bubble it up by repeatedly doing swap unsorted adjacent pairs
 - Repeat these until the unsorted area becomes empty

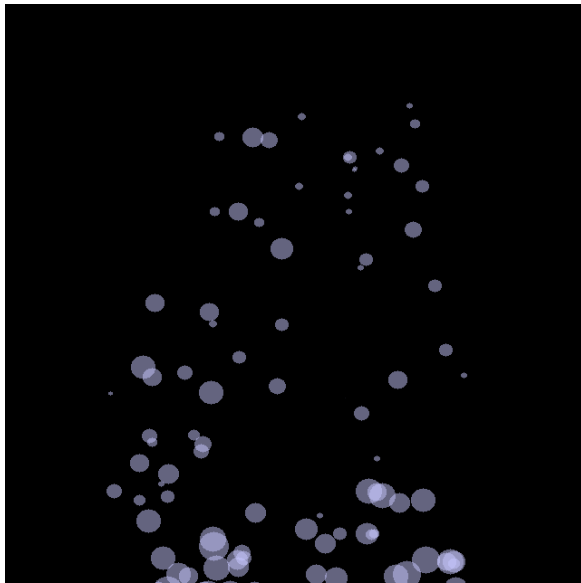
How 5 is
bubbled up



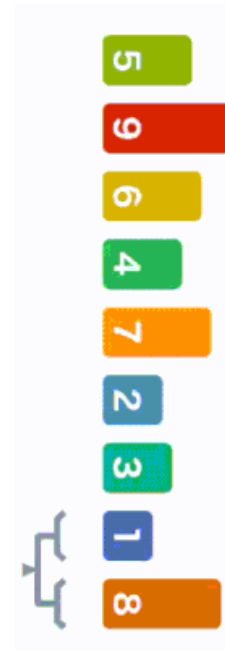
Bubble Sort (2)

Idea of bubble sort

- Move the maximum to the end of unsorted area
 - Bubble it up by repeatedly doing swap unsorted adjacent pairs
 - Repeat these until the unsorted area becomes empty



Bubble up



This sort bubbles
the **minimum** up each stage

Bubble Sort (3)

Pseudocode of bubble sort

```
def bubble_sort(A, n):  
    for end ← n downto 2:  
        # at this step, unsorted area is [1, end]  
  
        # bubble up stage (push the maximum to the end)  
        for i ← 1 to end - 1:  
            if A[i] > A[i+1]:  
                swap A[i] and A[i+1]  
  
        # at this step, unsorted area is [1, end-1]
```

▪ Correctness

- Like selection sort, the sorted area is expanded correctly, and finally its size becomes n at the end

Bubble Sort (4)

Time complexity of bubble sort

```
def bubble_sort(A, n):  
    for end ← n downto 2:  
        # bubble up stage  
        for i ← 1 to end - 1:  
            if A[i] > A[i+1]:  
                swap A[i] and A[i+1]
```

} Bubble up stage

▪ Time complexity analysis


- The outer for-loop repeats the bubble up stage $n - 1$ times
- In the bubble up stage
 - Main operations are comparison and swap having the same # of operations
 - $i - 1$ comparisons are required for the i -th bubble up stage

$$T(n) = (n - 1) + (n - 2) + \dots + 2 + 1 = \sum_{i=1}^{n-1} i = \frac{n(n - 1)}{2} \in \Theta(n^2)$$

Outline

Sorting problem

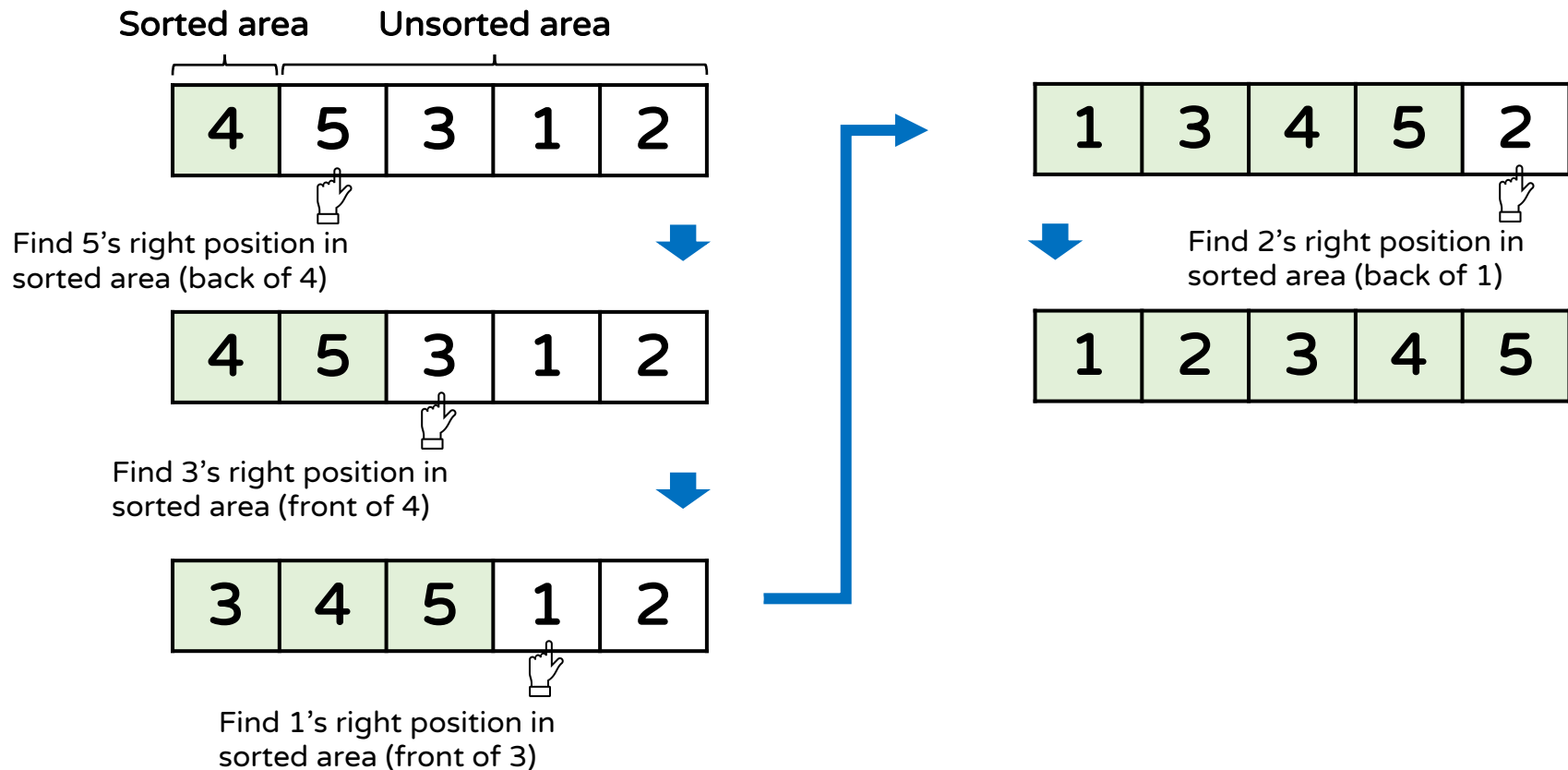
Basic sorting algorithms

- Selection Sort
- Bubble Sort
- Insertion Sort 

Insertion Sort (1)

Idea of insert sort

- Pick the front of unsorted area
- Insert it into its right position in sorted area



Insertion Sort (2)

Pseudocode of insertion sort

```
def insertion_sort(A, n):  
    for i ← 2 to n:  
        # at this step, sorted area is [1, i-1]  
  
        # insertion stage  
        insert A[i] into its right position in A[1..i]  
  
        # at this step, sorted area is [1, i]
```

■ Correctness

- For base case ($i = 1$), $A[i]$ is already sorted
- Assume it correctly works for $i = k$ (i.e., $A[1 \dots k]$ is sorted)
- Then, the insertion stage will correctly insert $A[k + 1]$ into the sorted area, meaning $A[1 \dots k + 1]$ is also sorted
- Thus, for $i = n$, $A[1 \dots n]$ will be sorted correctly (by induction)

Insertion Sort (3)

Pseudocode of insertion sort

```
def insertion_sort(A, n):  
    for i ← 2 to n:  
        # at this step, sorted area is [1, i-1]
```

---▶ # insertion stage
insert A[i] into its right position in A[1..i]

at this step, sorted area is [1, i]

i-th insertion stage

loc ← i - 1

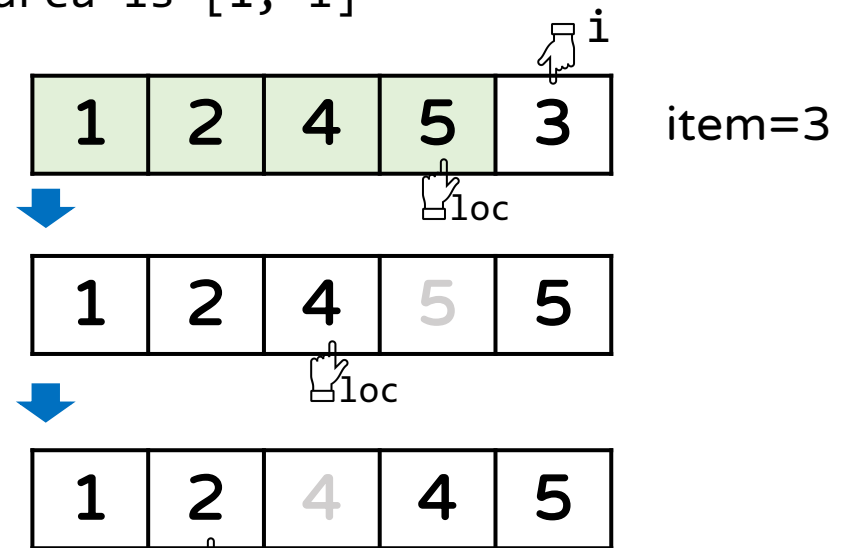
item ← A[i]

while loc ≥ 1 and item < A[loc]:

 A[loc+1] ← A[loc]

 loc ← loc - 1

A[loc+1] ← item



Back of 2 is the right position of 3

insert 3

Insertion Sort (4)

Time complexity of insertion sort

```
def insertion_sort(A, n):  
    for i ← 2 to n:  
        # insertion stage  
        insert A[i] into its right position in A[1..i]
```

▪ Time complexity analysis

- The outer for-loop repeats the insertion stage $n - 1$ times
- The i -th insertion stage performs at most i operations (comparisons)

$$T(n) = 1 + 2 + \dots + (n - 2) + (n - 1) = \sum_{i=1}^{n-1} i = \frac{n(n-1)}{2} \in \Theta(n^2)$$

What You Need To Know

Sorting problem

- To efficiently rearrange elements in an array in an order

Basic sorting algorithms

- **Selection Sort**
 - Move the maximum to the end of unsorted area
 - Find the maximum directly by linearly searching
- **Bubble Sort**
 - Move the maximum to the end of unsorted area
 - Move the maximum by bubbling it up
- **Insertion Sort**
 - Pick the front of unsorted area
 - Insert it into its right position in sorted area

In Next Lecture

Discussion on basic sorting algorithms

- Selection Sort
- Bubble Sort
- Insertion Sort

Advanced sorting algorithms

- Merge Sort
- Quick Sort

Thank You