

Container With Most Water

🔍 Difficulty	Medium
☰ Category	Two Pointers
🔗 Question	https://leetcode.com/problems/container-with-most-water/
🔗 Solution	https://youtu.be/UuiTKBwPgAo
🌟 Status	Done

Question

You are given an integer array `height` of length `n`. There are `n` vertical lines drawn such that the two endpoints of the `i`th line are `(i, 0)` and `(i, height[i])`.

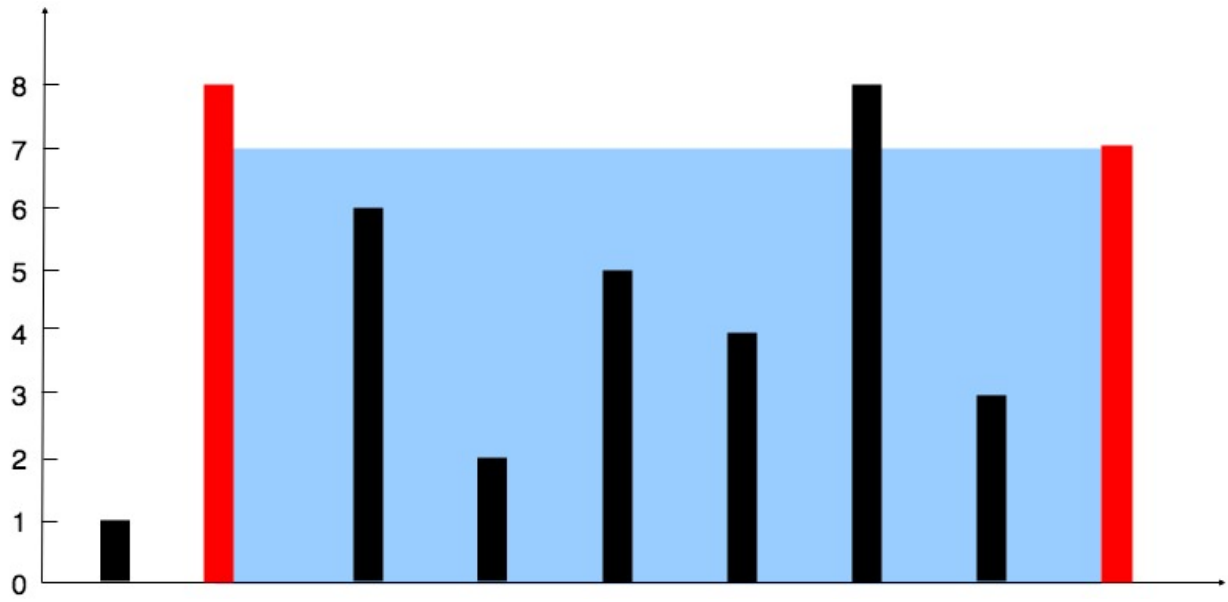
Find two lines that together with the x-axis form a container, such that the container contains the most water.

Return *the maximum amount of water a container can store*.

Notice that you may not slant the container.

Example

Example 1:



Input: height = [1,8,6,2,5,4,8,3,7]

Output: 49

Explanation: The above vertical lines are represented by array [1,8,6,2,5,4,8,3,7]. In this case, the max area of water (blue section) the container can contain is 49.

Example 2:

Input: height = [1,1]

Output: 1

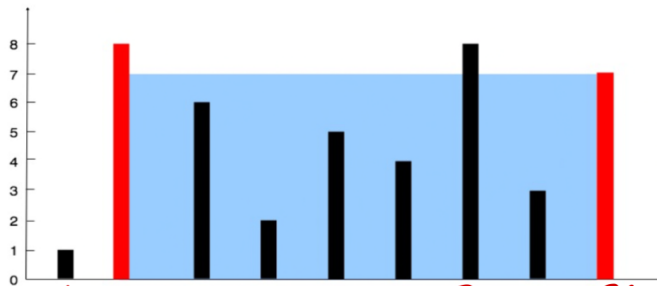
Idea



Shrinking window, left/right initially at endpoints, shift the pointer with min height;

④ Container with Most Water.

Example 1:

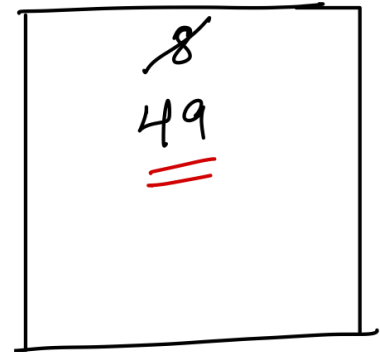


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Example 2:

Input: height = [1,1]
Output: 1

Max Area



$$\text{area} = (r - l) * \min(h[l], h[r])$$

$$= (8 - 0) * 1 = 8$$

Solution

```
class Solution:
    def maxArea(self, height: List[int]) -> int:
        res = 0 # Initialize the result (maximum area) to 0
        l, r = 0, len(height) - 1 # Initialize two pointers, 'l' and 'r', to the two ends of the array

        while l < r: # Loop until 'l' is less than 'r'
            # Calculate the area between the two lines (height[l] and height[r])
            area = (r - l) * min(height[l], height[r])

            # Update the result by taking the maximum of the current result and the calculated area
            res = max(res, area)

            if height[l] < height[r]:
                l += 1 # Move the left pointer to the right if the left line is shorter
            else:
                r -= 1 # Move the right pointer to the left if the right line is shorter

        return res # Return the maximum area found
```

Explanation

Time Complexity:

The time complexity of this code is $O(n)$, where 'n' is the length of the `height` list. This is because we use a two-pointer approach that moves towards each other, and each step of the loop either increments the left pointer or decrements the right pointer.

Space Complexity:

The space complexity is $O(1)$ because we only use a fixed amount of extra space to store the variables `res`, `l`, `r`, and `area`, regardless of the size of the input list.