Honey, I Shrunk the Guests

Page Access Tracking using a Minimal Virtualisation Layer

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Plenty of memory, hardly any knowledge

Many memory technologies

- HBM
- DRAM
- Persistent Memory
- CXL-attached memory
- NUMA

Different characteristics

- Capacity
- Latency
- Bandwidth
- Persistence
- Cost

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- ϟ ”Expert interfaces” are difficult (PMDK, libnuma)
- ϟ Language support does not expand to legacy programs (NV-Heaps)

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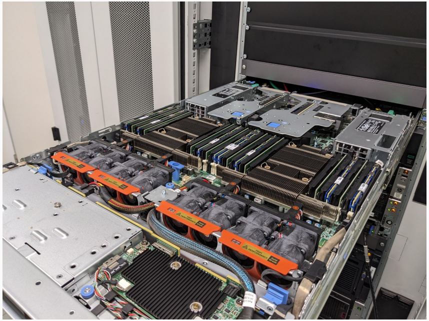
→ Operating system support for transparent memory placement

Problems we want to tackle

- Fast-tier memory is rare
  → it must be distributed efficiently
- Process workloads shift over time
  - in intensity
  - in locality
- Memory placement decisions should be adaptable



→ Detailed runtime information on memory utilisation is necessary

Identifying which memory works best for a given process?

Conventional approaches

- Instrumentalisation
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Page Modification Logging (PML)

- Virtualisation extension for VM checkpointing, VM migration
- Hardware-based logging of write accesses
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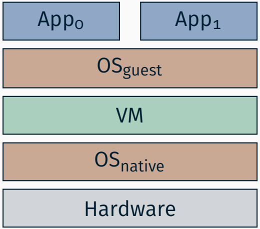
→ Is virtualisation viable for gathering memory access statistics?

vmmload: A minimal virtualisation layer

- Guest OS introduces overhead
- Processes should communicate with the host OS
- Processes should communicate with other processes



PML with VM

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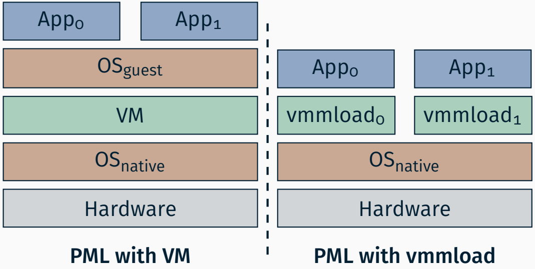
→ **vmmload** as minimal hypervisor



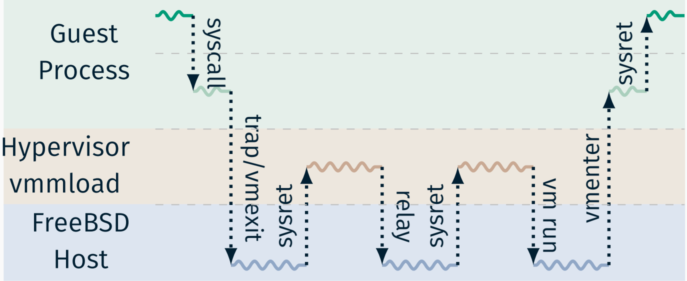
Relaying of System Calls

Requirements

- Isolate hypervisor from guest
- Interface with host OS

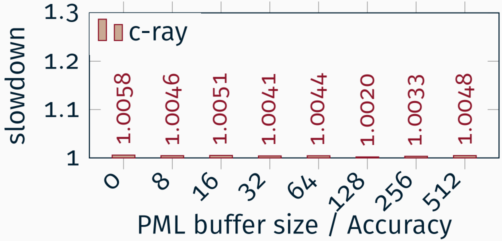
Implications

- Emulate system calls that manipulate the issuers process' state
- Translation of memory addresses



Slowdown

CPU-bound processes



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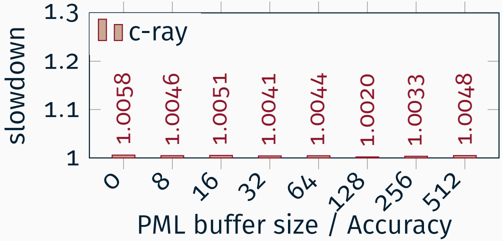

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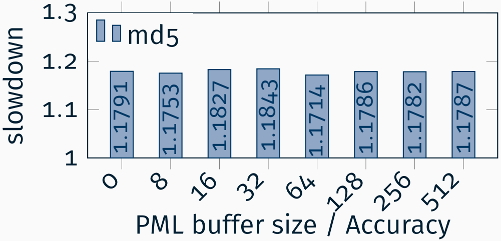
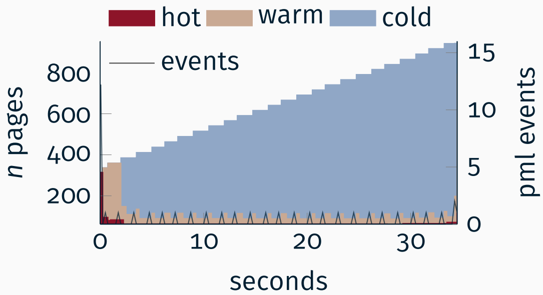
I/O bound processes



- Highly interactive processes
- System call overhead introduces great slowdown

Access Patterns for *c-ray*



Temperature

- Identify total number of pages
- Identify frequently accessed pages
- Frequency of PML events

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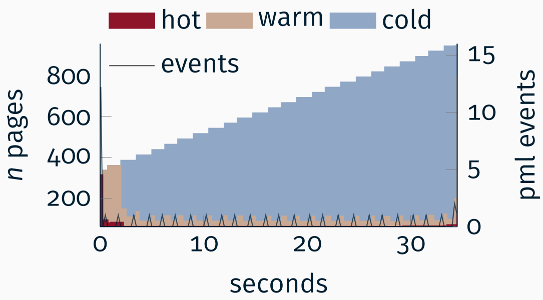

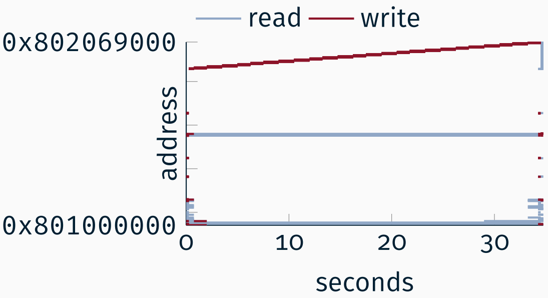

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Distribution

- Distinguish between *read*/*write*
- Alteration of working set
- Sparsity/density of accesses

Summary

We have

- ✓ shown that **vmmload** can collect memory access statistics
- ✓ achieved statistics over read/write accesses

Next, we plan to

- 🕓 Reduce system call overhead to $\frac{1}{4}$ with kernel integration
- 🕓 Derive memory placement/migration decisions

→ Source code is freely available

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→ Thank you for your attention