

Consider the game-tree above. The maximising player controls red nodes and the minimising player controls the green ones.

For each of the following assignments of values to  $x, w, y, z$  decide if under those assignments Alpha-Beta search explores all states of the game-tree.

Vælg en svarmulighed på hver linje

$x := 8, w := 7, y := 9, z := 8$

Yes

☐

No

☒

$x := 6, w := 8, y := 5, z := 6$

☐
☒

$x := 6, w := 5, y := 6, z := 8$

☐
☒

It is known that  $MINIMAX(s_1) = 7$ . Decide which of the following assignment of values to  $x, w, y, z$  are possible in such case.

Vælg en svarmulighed på hver linje

$x := 8, w := 7, y := 6, z := 8$

Yes

☒

No

☐

$x := 6, w := 5, y := 6, z := 8$

☒
☐

$x := 6, w := 8, y := 6, z := 5$

☒
☐

$p, q$	$p, \bar{q}$	$\bar{p}, q$	$\bar{p}, \bar{q}$
$x$	$y$	$z$	$w$

↓ *more plausible*

Which state is most plausible after the revision with  $\neg p$ ?

Vælg én svarmulighed

☒  $z$

☐  $x$

☐  $w$

☐  $y$

Which states are most plausible after contraction with  $p$ ?

Vælg én svarmulighed

☐  $z$  and  $w$

☐  $x$  and  $y$

☒  $x$  and  $z$

☐  $w$  and  $x$

Propositional Logic

Consider this numbered list of formulas:

- 1.  $(p \vee q) \wedge (\neg p \vee q)$
- 2.  $(p \vee \neg q) \vee (\neg p \vee q) \vee (p \vee \neg p)$
- 3.  $p \vee p \vee q$
- 4.  $q \wedge \neg q \wedge p$

Which of the above formulas match the following descriptions?

Vælg en svarmulighed på hver linje

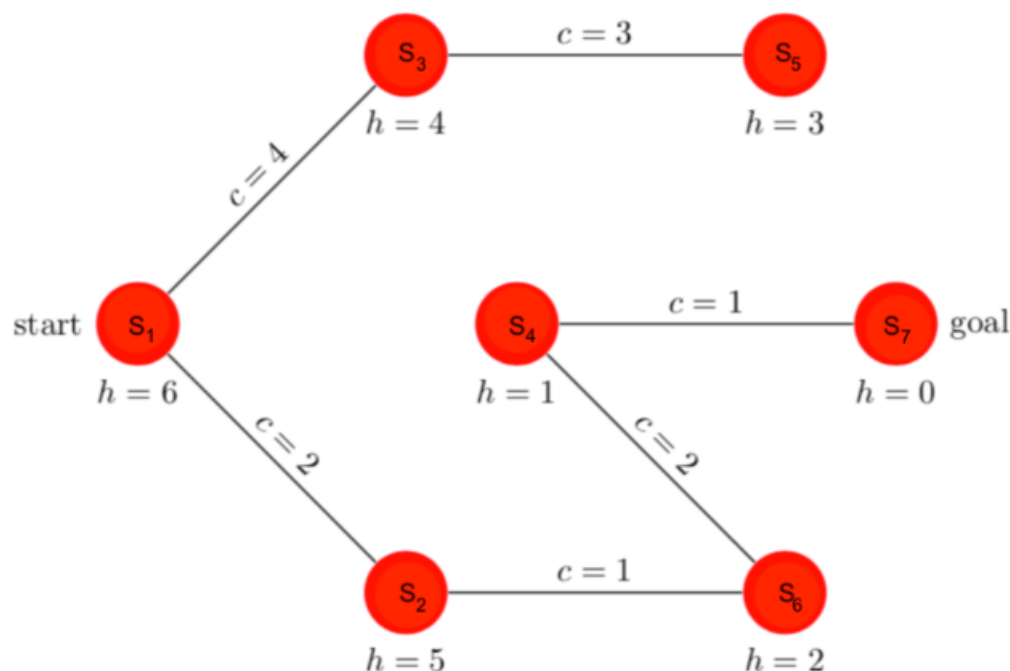
valid disjunction of definite clauses

result of resolving  $\neg r \vee p$  and  $p \vee q \vee r$

equivalent to  $q$

unsatisfiable conjunction of Horn clauses

1	2	3	4
<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>
<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>	<input type="radio"/>
<input checked="" type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input type="radio"/>
<input type="radio"/>	<input type="radio"/>	<input type="radio"/>	<input checked="" type="radio"/>



Consider the state space in figure above. Heuristic function for each node is given by  $h$  and the step-costs are given by  $c$ . All transitions are bidirectional. You are asked to answer questions about the run of various tree and graph search algorithms on this space.

Evaluate the following statement:

Uniform Cost Graph Search expands node  $s_5$ .

Vælg én svarmulighed

☐ True

☒ False

Evaluate the following statement:

$A^*$  Graph Search expands node  $s_3$ .

Vælg én svarmulighed

☐ True

☒ False

Evaluate the following statement:

Greedy Best-First Tree Search expands node  $s_6$ .

Vælg én svarmulighed

☒ False

☐ True

Evaluate the following statement:

Greedy Best-First Graph Search expands node  $s_2$ .

Vælg én svarmulighed

☐ False

☒ True

Evaluate the following statement:

Greedy Best-First Tree Search expands node  $s_4$ .

Vælg én svarmulighed

☒ False

☐ True

Evaluate the following statement:

Uniform Cost Tree Search expands node  $s_5$ .

Vælg én svarmulighed

☐ True

☒ False

Evaluate the following statement:

$A^*$  Tree Search expands node  $s_3$ .

Vælg én svarmulighed

☒ False

☐ True

Partially observable domains

Consider the following partially observable domain. An agent is in a room that contains a light bulb and a coin. There are four possible physical states:

- *light\_on\_heads\_up* : the light bulb is **on** and the coin is laying **heads up**.
- *light\_on\_tails\_up* : the light bulb is **on** and the coin is laying **tails up**.
- *light\_off\_heads\_up* : the light bulb is **off** and the coin is laying **heads up**.
- *light\_off\_tails\_up* : the light bulb is **off** and the coin is laying **tails up**.

There is one action available:

- *turn\_light\_on* : the agent can turn the light bulb on, if it is currently off

The ACTIONS and RESULTS functions are defined as follows:

- $ACTIONS(light\_off\_heads\_up) = ACTIONS(light\_off\_tails\_up) = \{turn\_light\_on\}$
- $ACTIONS(light\_on\_heads\_up) = ACTIONS(light\_on\_tails\_up) = \{\}$
- $RESULTS(light\_off\_heads\_up, turn\_light\_on) = \{light\_on\_heads\_up\}$
- $RESULTS(light\_off\_tails\_up, turn\_light\_on) = \{light\_on\_tails\_up\}$

The agent can only see whether the coin is laying heads up or tails up, if the light is currently on. Formally, the PERCEPT function is as follows:

- $PERCEPT(light\_on\_heads\_up) = heads\_up\_visible$
- $PERCEPT(light\_on\_tails\_up) = tails\_up\_visible$
- $PERCEPT(light\_off\_heads\_up) = PERCEPT(light\_off\_tails\_up) = null$

Suppose that the agent's initial belief state is  $b = \{light\_off\_heads\_up, light\_off\_tails\_up\}$ .

What is  $POSSIBLE-PERCEPTS(\bigcup_{s \in b} RESULTS(s, turn\_light\_on))$ ?

Vælg én svarmulighed

- ☐  $\{null, heads\_up\_visible, tails\_up\_visible\}$
- ☐  $\{null\}$
- ☐  $\{\{null\}, \{heads\_up\_visible, tails\_up\_visible\}\}$

☒  $\{heads\_up\_visible, tails\_up\_visible\}$

What is  $\bigcup_{s \in b} RESULTS(s, turn\_light\_on)$ ?

Vælg én svarmulighed

- ☐  $\{light\_off\_heads\_up, light\_off\_tails\_up, light\_on\_heads\_up, light\_on\_tails\_up\}$
- ☒  $\{light\_on\_heads\_up, light\_on\_tails\_up\}$
- ☐  $\{\{light\_on\_heads\_up\}, \{light\_on\_tails\_up\}\}$
- ☐  $\{light\_on\_heads\_up\}$

What is  $RESULTS'(b, turn\_light\_on)$ ?

Vælg én svarmulighed

- ☐  $\{\{light\_on\_heads\_up, light\_on\_tails\_up\}\}$
- ☐  $\{\{light\_on\_heads\_up\}, \{light\_on\_tails\_up\}\}$
- ☐  $light\_on\_heads\_up$

☒  $\{light\_on\_heads\_up, light\_on\_tails\_up\}$

☐  $\{light\_off\_heads\_up, light\_off\_tails\_up, light\_on\_heads\_up, light\_on\_tails\_up\}$

Belief Revision: AGM  
Answer the following questions about belief revision.

Let  $A = \{p, q, \neg(\neg p \wedge \neg q), p \rightarrow q, \neg p \rightarrow q\}$  be a belief base. Which of the following sets are in  $A \perp q$ ?

Vælg en svarmulighed på hver linje

$\{p, \neg q, p \vee q, \neg p \rightarrow q\}$

Yes

☐

No

☒

$\{p, \neg(\neg p \wedge \neg q)\}$

☐☒

$\{p, \neg(\neg p \wedge \neg q), p \rightarrow q, \neg p \rightarrow q\}$

☐☒

$\{\neg(\neg p \wedge \neg q), p \rightarrow q, \neg p \rightarrow q\}$

☐☒

$\{p, \neg(\neg p \wedge \neg q), \neg p \rightarrow q\}$

☒☐

Is it the case that  $\neg p \vee p \in Cn(\{q, p \wedge \neg q\})$ ?

Vælg én svarmulighed

☐ No

☒ Yes