An Accelerated Procedure for Hypergraph Coarsening on the GPU

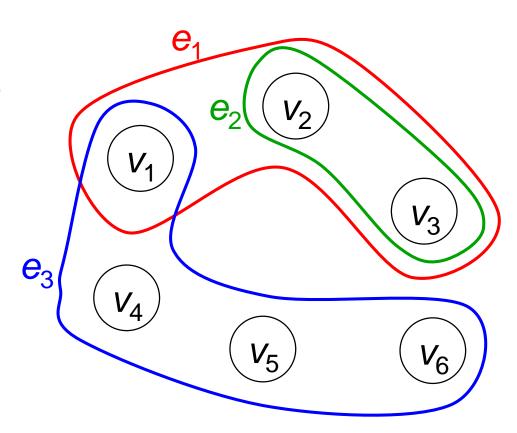
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Outline

- Hypergraph coarsening
- Implementation challenges
- Runtime task planning
- Results

Hypergraph

- Nodes
- Hyperedges (nets)
 - Subsets of nodes



Hypergraph

Hypergraph partitioning

Minimize edge cut

Balance constraint

NP-complete

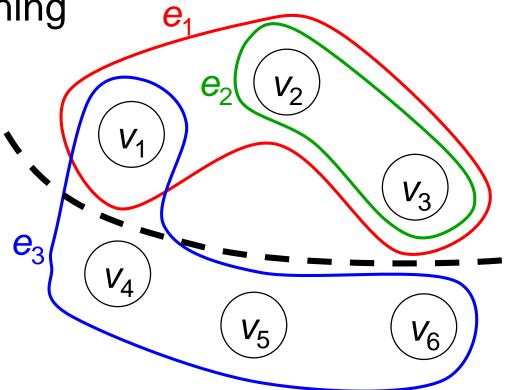
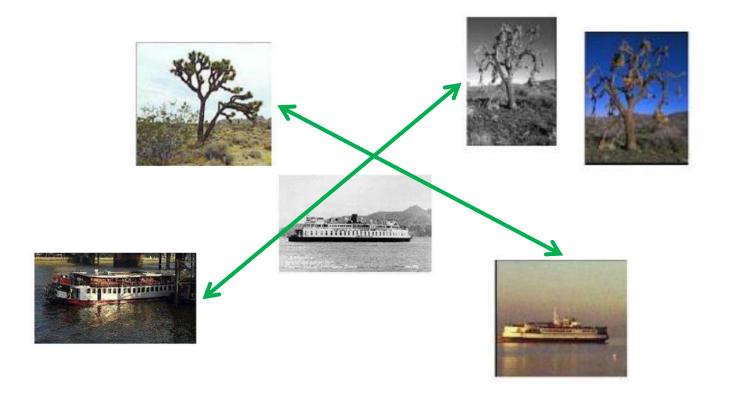


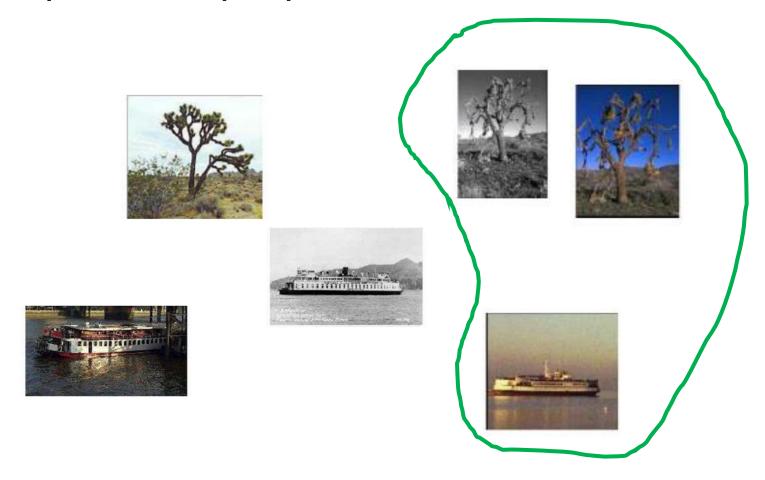
Image classification

- Similar images should go to same category
- Need to compare images to one another



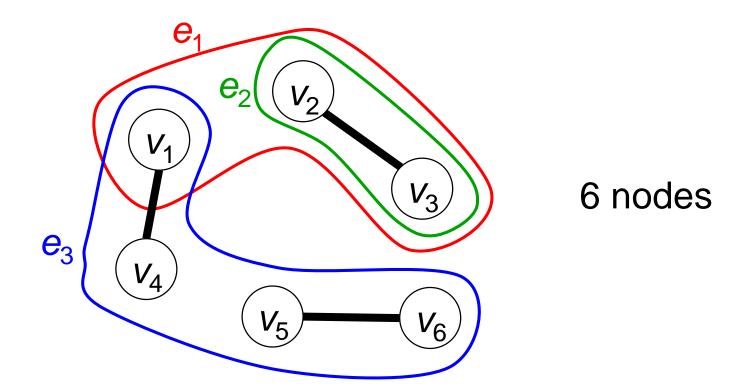
Hypergraph construction

Compare multiple pictures at a time



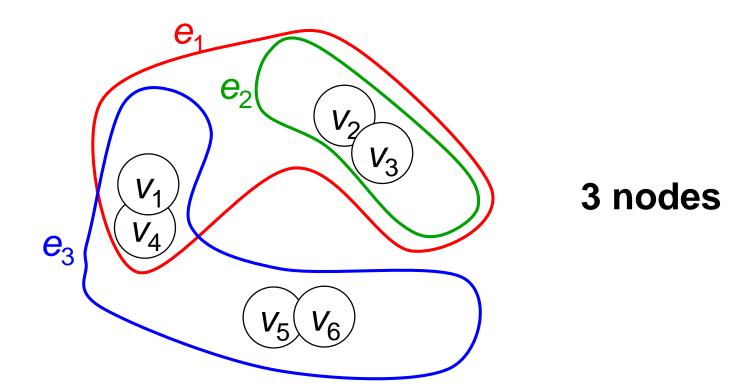
Hypergraph coarsening

Heuristic: reduce # nodes by fusing

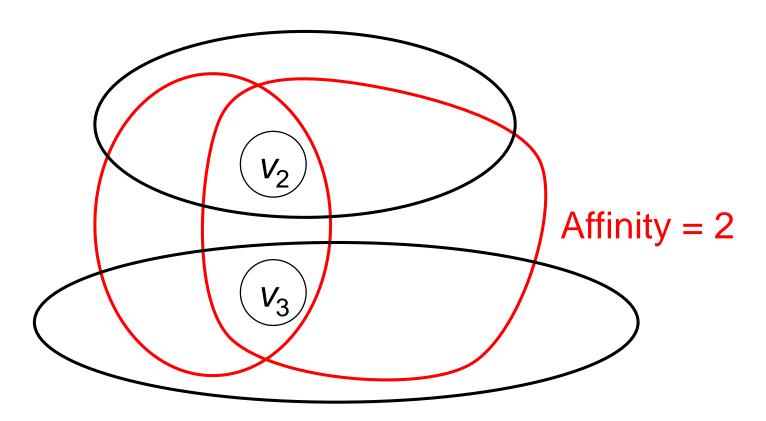


Hypergraph coarsening

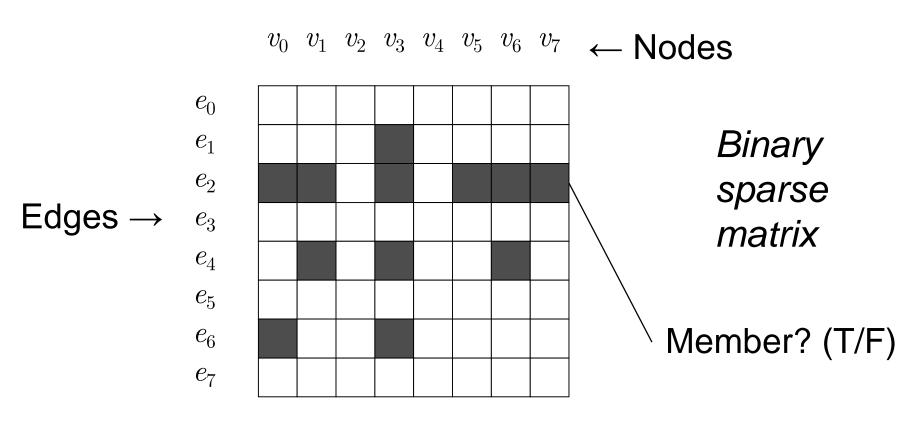
Heuristic: reduce # nodes by fusing



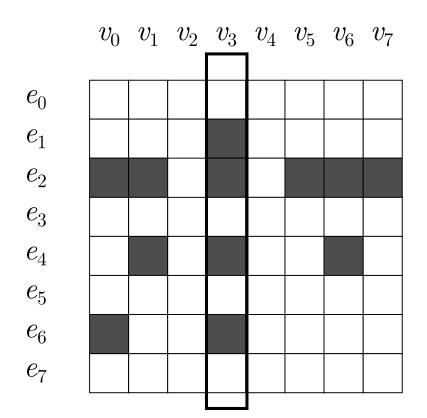
- Given a node, find most "similar" neighbor
- Similarity = # hyperedges containing both



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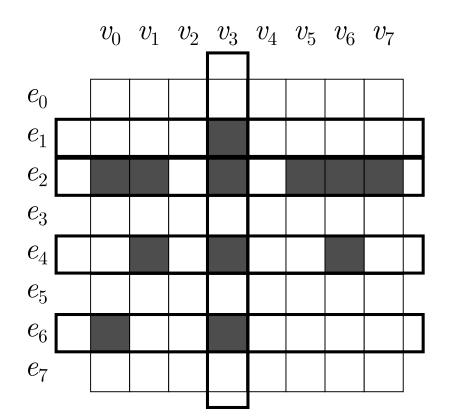


- Given a node, find most "similar" neighbor
- Similarity = # hyperedges containing both



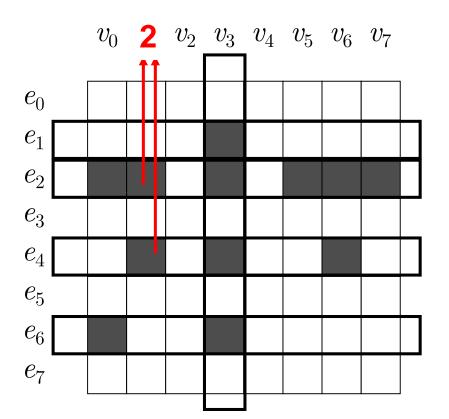
1. Find edges containing v_3

- Given a node, find most "similar" neighbor
- Similarity = # hyperedges containing both



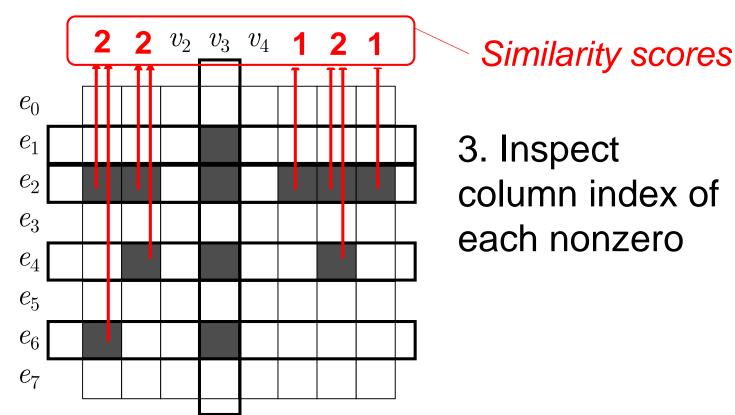
- 1. Find edges containing v_3
- 2. Collect nonzeros

- Given a node, find most "similar" neighbor
- Similarity = # hyperedges containing both

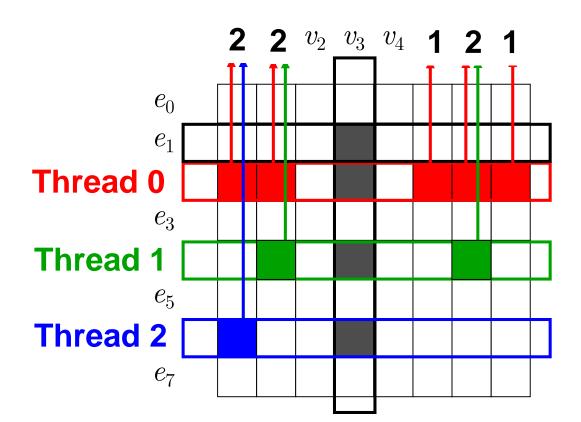


3. Inspect column index of each nonzero

- Given a node, find most "similar" neighbor
- Similarity = # hyperedges containing both

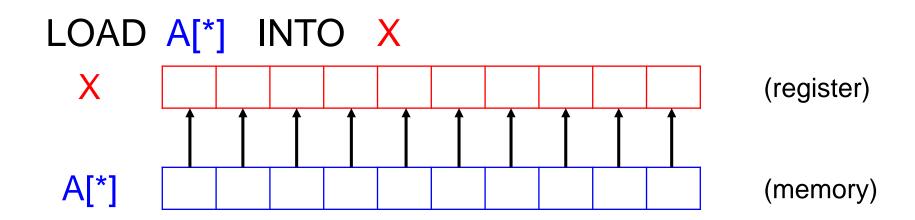


 Parallel algorithm: Inspect edges in parallel

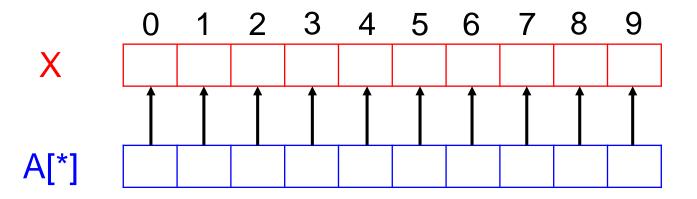


GPU as parallel accelerator

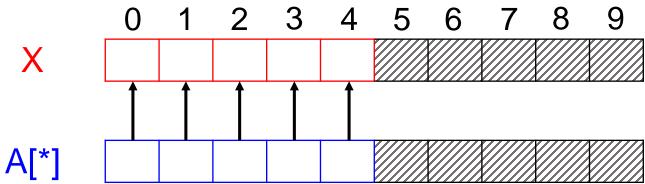
NVIDIA GPUs : organized in warps
 32 threads share one instruction counter

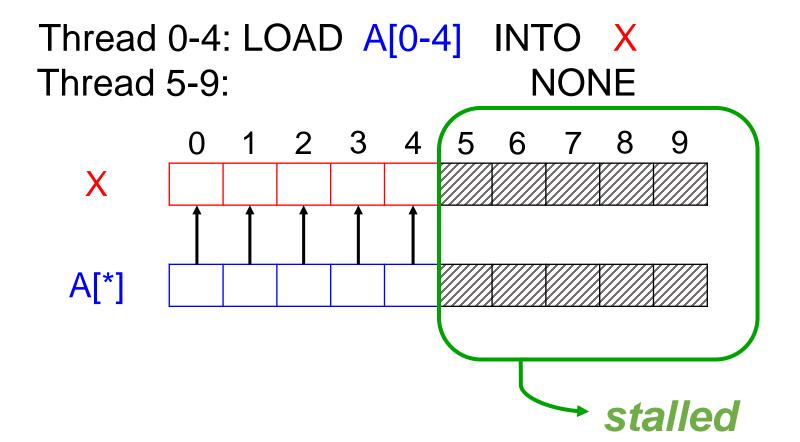


Thread 0-9: LOAD A[0-9] INTO X



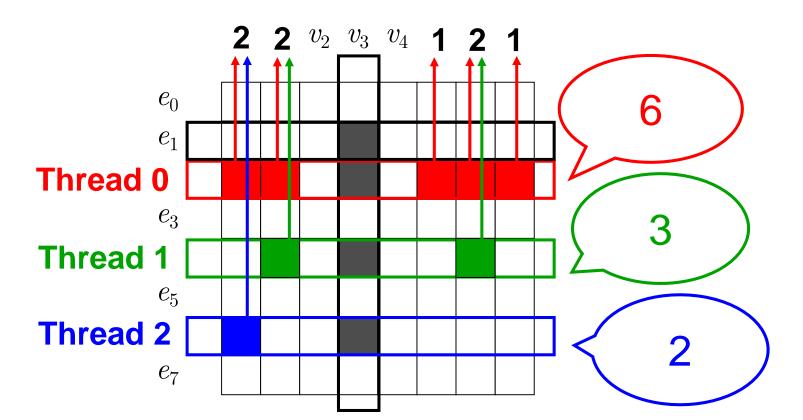
Thread 0-4: LOAD A[0-4] INTO X Thread 5-9: NONE





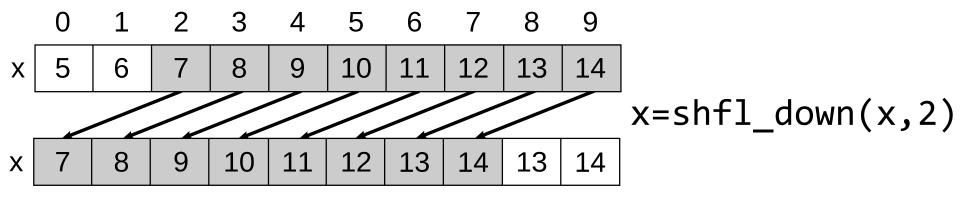
- Serializes execution
- Caused by load imbalance
 - Sparse/irregular data

- A naïve strategy results in load imbalance
- Nonzero entry = workload



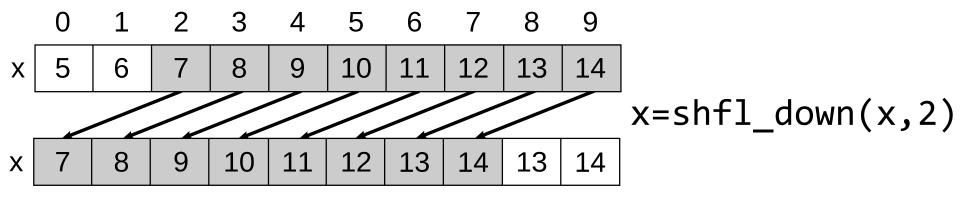
SHFL to the rescue

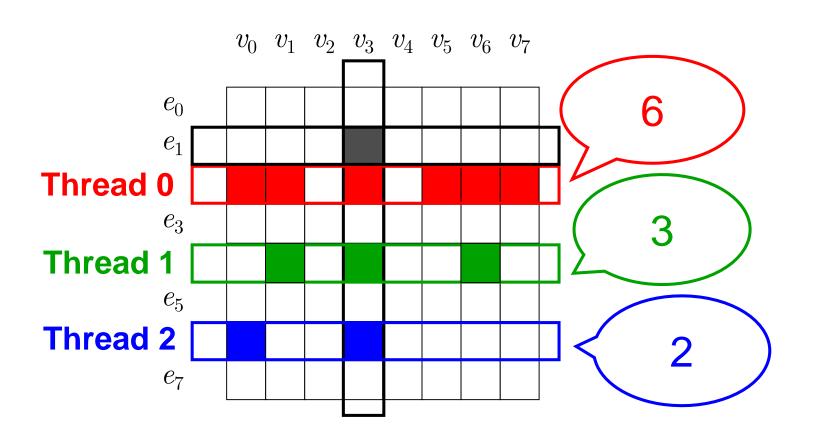
- Compiler primitive
- Shuffles content of adjacent registers

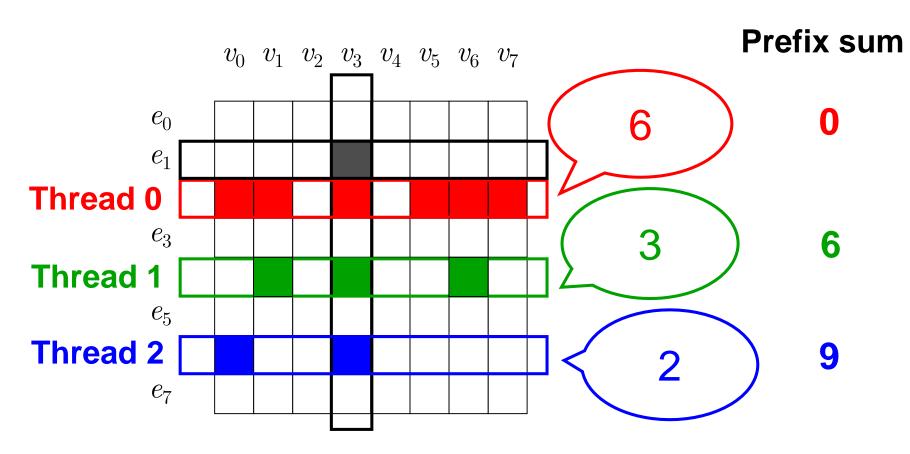


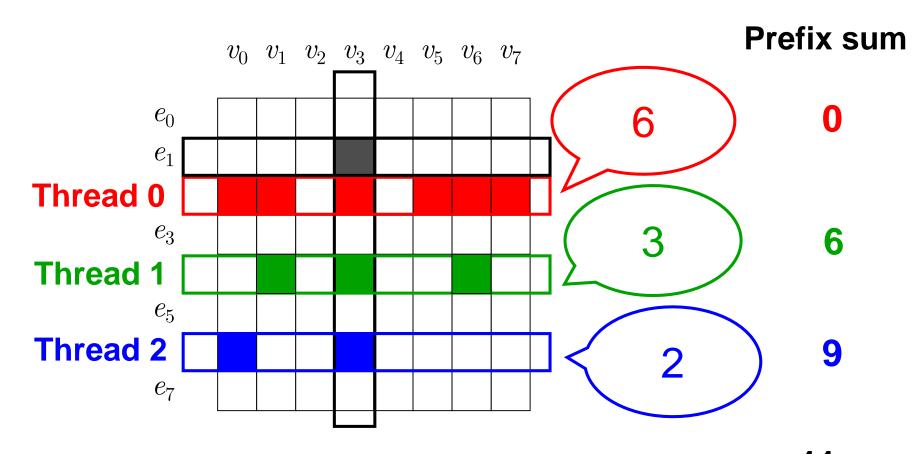
SHFL to the rescue

- Compiler primitive
- Shuffles content of adjacent registers
- Single machine instruction
- Warp-synchronous; no sync. needed after

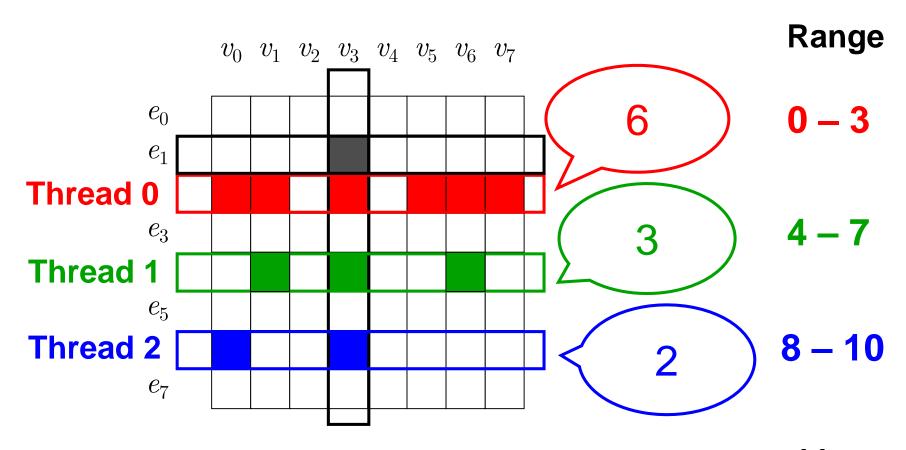




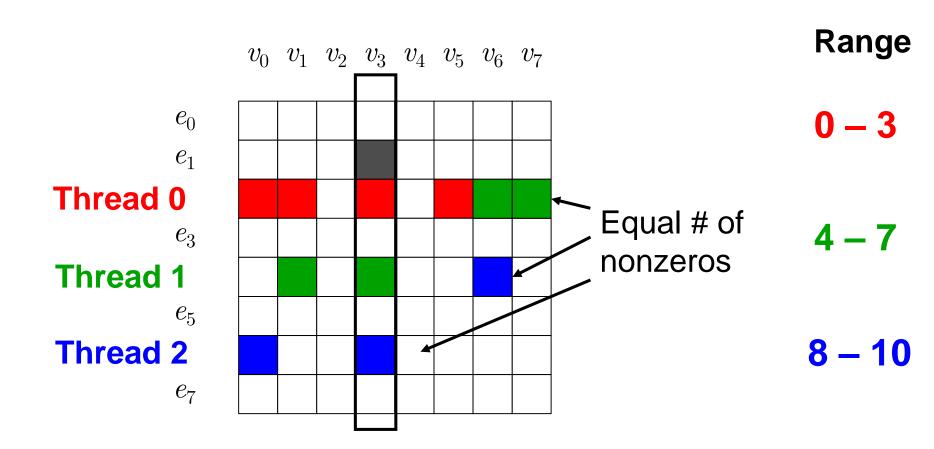


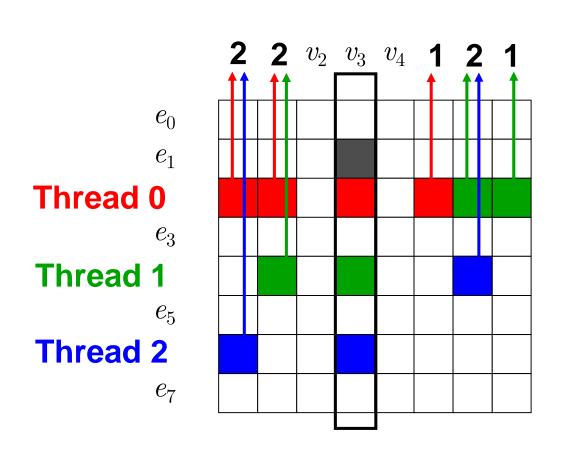


$$ceil(11 / 3) = 4$$



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Range

0 - 3

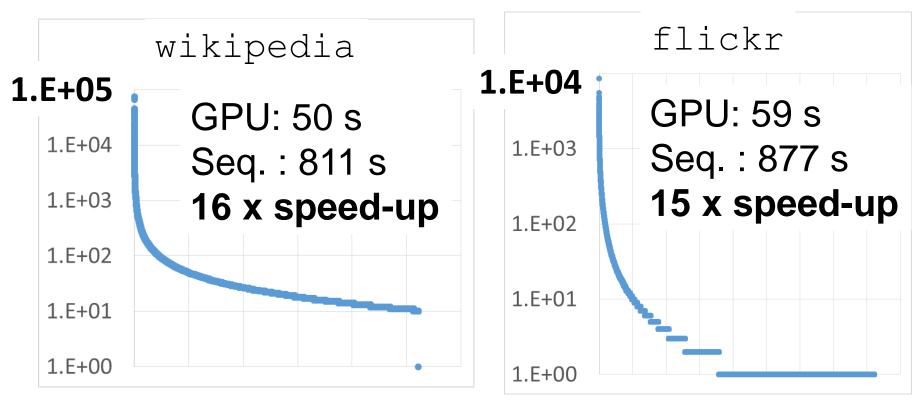
4 - 7

8 - 10

Results

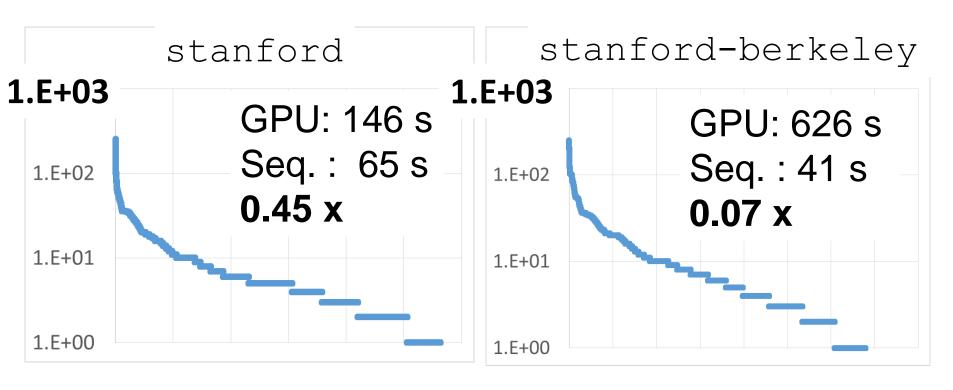
- NVIDIA Tesla K20c (5GB mem)
- Intel Xeon E5-2620
- Reference sequential implementation of Mondriaan

Results: long-tailed distribution



Y-axis: # of nonzeros in columns, descending, log-scale

Results: non-long-tailed distribution



Y-axis: # of nonzeros in columns, descending, log-scale

Analysis of results

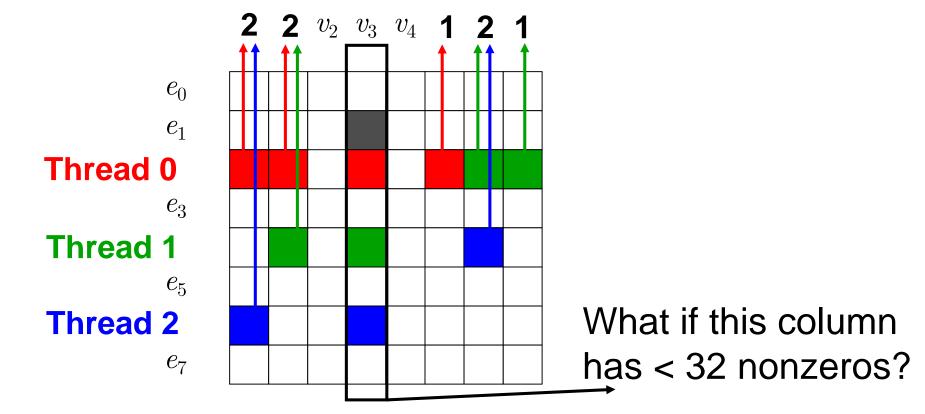
 Good speedup for data with long-tailed distribution of nonzeros

Analysis of results

- Good speedup for data with long-tailed distribution of nonzeros
- Synthetic data
 - First 1,000 columns: 100,000 nonzeros each
 - Next 699,000 columns: all zero
 - Speedup: 123 x

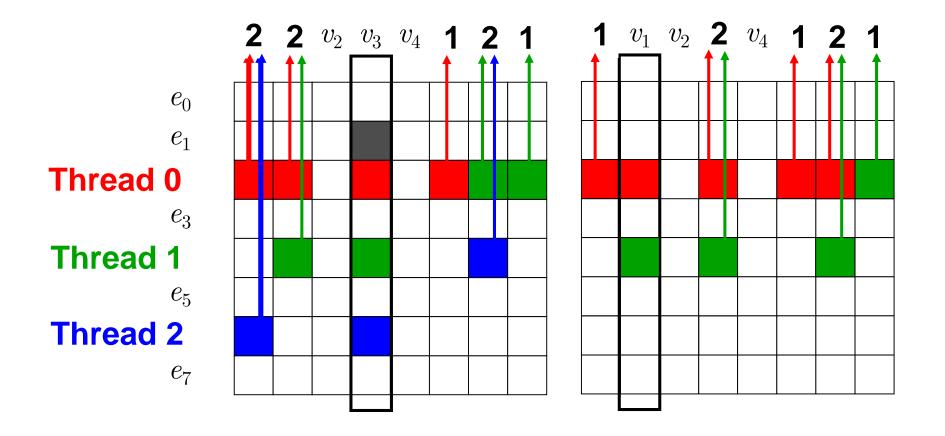
Analysis of results

- Most nodes sparsely connected (<< 32 edges)
- Now: one warp, one instance of Mondriaan



Work in progress

Pool multiple instances of Mondriaan into warp



Conclusion

- Implemented hypergraph algorithm handling arbitrary connectivity patterns
- Explored SHFL for task planning

Future work: more flexible allocation strategy

Acknowledgment

- Trinity College: Student Research Program
- NVIDIA: CUDA Teaching Center