1. Why Wyvern?

Wyvern is a new programming language exploring how to help software engineers build software better at scale. Our primary aim is to do research that discovers and validates new principles for designing engineering-focused programming languages. Part of the "validation" bit is actually writing interesting programs in Wyvern, and so we hope to also make it a great language to write code in—for us, and for you!

Wyvern's design incorporates a lot of great ideas from prior languages: it is a statically-typed, garbage-collected general-purpose applications programming language with excellent support for both object-oriented and functional programming. However, Wyvern is most interesting because of the new ideas it explores. Most of these are best illustrated through the examples below, but here's a brief overview of Wyvern's most interesting design features and what motivates them.

- Large programs must be composed from parts, so Wyvern has an advanced module system with features echoing those of Standard ML's module system. However, modern programs often load and (re-)compose modules at run time, so Wyvern modules and functors are first-class objects and functions, respectively, providing programmers with the power and flexibility they need.
- No one language can be good at everything, so Wyvern is extensible: libraries can define new syntax for the abstractions that they provide, and that syntax can be embedded seamlessly in Wyvern expressions.
- Engineers need to control the access that untrusted code has to resources such as the network or file system, and so Wyvern's module system is the first to be designed from the ground up to be capability-safe. We are designing an effect system that leverages these capabilities to provide lightweight and automatically-checked control of resources.
- More broadly, understanding and controlling the software architecture of a program is critical to understanding its properties and evolving it over time. Thus we are building a way of expressing software architecture as an integrated part of Wyvern programs, including a static view building on the module system and a dynamic view showing run-time components and the connections between them. Wyvern architectures are "live" in that changing the architecture specification affects the program semantics, and they are "trusted" in that they are guaranteed to be an accurate abstraction of what the program does.

2. Hello, World! in Wyvern

Here is a "Hello, World!" program in Wyvern (examples/rosetta/hello.wyv): require stdout

```
stdout.print("Hello, World!")
```

This program already illustrates a couple of basic aspects of Wyvern. First, Wyvern is object-oriented: stdout is an object, and we are invoking the print method on it. For expressions, much of Wyvern's syntax is similar to Java's.

Second, system resources such as the standard output object, $\tt stdout$, are not ambiently available to programs, but must be explicitly required from the operating system. A primary goal of Wyvern's module system is helping developers to reason about the use of resources. Thus even a simple script such as Hello World must declare the resources it requires in order to execute. This allows engineers to determine at a glance what kind of I/O a program might do, and provides a basis for making a decision about whether to run this program in a particular situation. In this case, even without looking at the actual code, we know that this program may write to the standard output stream, but will not access the file system or access the network.

3. Anonymous Functions

Wyvern has good support for functional programming, and anonymous functions can be defined in Wyvern using the syntax:

```
(x:Int) \Rightarrow x + 1
```

We can bind the expression above to a variable and invoke it:

```
val addOne = (x:Int) \Rightarrow x + 1 addOne(1)
```

and the result will be 2.

Anonymous functions can also have multiple parameters:

```
(x:Int,y:Int) => x + y
```

or no parameters:

```
() => 7
```

Function types can be denoted with an arrow, and we can annotate a variable with this type. If we annotate the type of the variable we are binding to the function, we can leave out the type annotation (and even the parentheses) on the function's argument:

```
val annotatedAddOne : Int \rightarrow Int = x \Rightarrow x + 1
```

This also works if we pass an anonymous function to a higher-order function:

```
val invokeIt = (f:Int -> Int, x:Int) => f(x)
invokeIt(x => x+1, 5)
```

The code above can be found in examples/introductory/functions.wyv

4. Functions in Wyvern

Consider the definition of the factorial function in Wyvern (examples/rosetta/factorial.wyv): require stdout

```
def factorial(n:Int):Int
    (n < 2).ifTrue(
          () => 1,
          () => n * factorial(n-1)
    )
stdout.print("factorial(15) = ")
stdout.printInt(factorial(15))
```

A function is defined with the def keyword, and its argument and return types are given in Algol-like syntax. Functions defined with def are recursive, so we can call factorial in the body. The example illustrates how an integer comparison n-2 is a boolean object, on which we can invoke the ifTrue method. This method takes two functions, one of which is evaluated in the true case and one of which is evaluated in the false case.

Note that factorial(15) would overflow in languages such as Java in which the default integer types is represented using only 32 bits. In Wyvern, Int means an arbitrary precision integer.

Wyvern provides two nicer ways to write the if statement above:

In the first case, we put the then computation on a new, indented line, followed by else (also indented, but not as much), and then the else computation. In the second case, we put it all on one line of code, but put the then and else computations in curly braces to set them off.

It turns out that both versions of the if function can be user-defined. In the first case, Wyvern takes all the indented code and passes it as an argument to the if function. In the second case, Wyvern passes the code in each set of curly braces as a separate argument to the if...else... function.

If you'd like to see how these if functions are actually user-defined, look in the Wyvern standard prelude (stdlib/prelude.wyv):

```
import metadata wyvern.IfTSL
val ifelseARG = (cond:Boolean, tt:IfTSL.FnExpr, ff:IfTSL.FnExpr) => IfTSL.doif(cond, tt, ff)
val if = (cond:Boolean, blocks:IfTSL.Blocks) => IfTSL.doifblk(cond, blocks)
```

The import statement loads the IfTSL module from the wyvern package. The metadata keyword indicates that the library defines new syntax, in this case for then and else clauses of an if statement. We define an if shorthand for calling the doifblk function from that library. The ifelseARG function is how if...else... is written: keywords separated by a single argument are concatenated, and whenever more than one argument separates keywords (or comes after them all) ARG is added to the name. Parsers associated with the IfTSL.FnExpr and IfTSL.Blocks are defined in stdlib/wyvern/IfTSL.wyv.

5. Objects and Object Types in Wyvern

We can define a sumable integer list type as follows (examples/introductory/objects.wyv):

```
type IntList
    def sum():Int
```

The type keyword declares a new object type, called IntList in this case. The public methods available in the type are listed below, but no method bodies may be given as we are defining a type, not an implementation.

We can implement a constant representing the empty list and a constructor for creating a larger list out of a smaller one as follows:

```
val empty:IntList = new
    def sum():Int = 0

def cons(elem:Int,rest:IntList):IntList = new
    def sum():Int = elem + rest.sum()

cons(3,cons(4,empty)).sum() // evalutes to 7
```

The new expression creates an object with the methods given. In the example above, we just have one method, sum(), which evaluates to 0 in the case of the empty list and sums up the integers in the list otherwise.

6. Strings and Characters

String literals can be written in quotes, using the same escapes as in Java. Strings support several operations, including ==, <, >, length(), and charAt(Int). The last of these returns a Character, which supports ==, <, and > operations. A simple program illustrating these is in examples/introductory/strings.wyv Wyvern will support character literals but doesn't yet.

7. Anonymous Functions as Objects

The anonymous function syntax described above is actually a shorthand for

creating an object with an apply method that has the same arguments and body:

```
new
    def apply(x:Int):Int = x + 1
which is an instance of the following type:
type IntToIntFn
    def apply(x:Int):Int
```

As mentioned earlier, the type above can be abbreviated Int -> Int, as in many other languages with good support for functional programming.

8. Mutable State and Resource Types

Types with mutable state can be defined, but need to be marked as resource types (examples/introductory/cell.wyv):

```
resource type Cell
   def set(newValue:Int):Unit
  def get():Int

def makeCell(initVal:Int):Cell = new
  var value : Int = initVal
  def set(newValue:Int):Unit
       this.value = newValue
  def get():Int = this.value

val c = makeCell(5)
c.get() // evalutes to 5
c.set(3)
c.get() // evalutes to 3
```

Here makeCell uses a new statement to create an object with a var field value. var fields are assignable, so the set funtion is implemented to assign the value field of the receiver object this to the passed-in argument. Note that we must initialize a var field with an initial value. If we had not declared Cell to be a resource type, we would get an error because the new expression creates a stateful object that is a resource.

In the example above, Unit is used as the return type of functions that do not return any interesting value.

9. Modules

We can define the Cell abstraction above in a module (examples/modules/cell.wyv): module cell

```
resource type Cell
   def set(newValue:Int):Unit
   def get():Int

def make(initVal:Int):Cell = new
   var value : Int = initVal
   def set(newValue:Int):Unit
        this.value = newValue
   def get():Int = this.value
```

In Wyvern, analogously to Java, a module named m should be stored in a file m.wyv (we expect that the implementation will enforce this in the near future). The file system forms a hierarchical namespace with one name per directory that allows us to find modules by their qualified name. In this case, within the examples directory of the Wyvern distribution we have the directory modules that contains cell.wyv, so we can use it in a program as follows (see examples/modules/cellClient.wyv):

```
import modules.cell
val myCell : cell.Cell = cell.make(3)
myCell.set(7)
myCell.get() // evalutes to 7
```

Here the import statement takes a fully qualified name and uses this to find the file defining module cell. The module is actually an object that gets bound to the name cell. We can invoke make() on cell just as if it were a method. Types such as Cell defined in the cell module can be referred to by their qualified names, i.e. cell.Cell. In fact, types can be defined as members of an object as well, and the same qualified syntax can be used to refer to them. So modules are not special semantically: they are just a convenient syntax for defining an object. Consider what is the type of cell? The answer could be written as follows:

```
type TCell
    resource type Cell
        def set(newValue:Int):Unit
        def get():Int

def make(initVal:Int):this.Cell
```

Wyvern files that define a type use a .wyt extension (for Wyvern Type), and you can find the above definition at examples/modules/TCell.wyt.

10. Functors

Although modules are objects, they cannot contain any state and they cannot encapsulate system resources. How can we express designs similar to those in other languages, in which modules do these things? We can do so by defining functors. A functor is a module that is defined as a function: it takes zero or more arguments and returns an object. For example, we can define a functor that returns a stateful Cell object (examples/modules/cellasModule.wyv):

```
module def cellAsModule()
var value : Int = 0
def set(newValue:Int):Unit
   value = newValue
def get():Int = value
```

Wyvern does not allow implicitly shared global state, because this often causes problems in software development. So cellasModule does not evaluate to an object, but rather a function that, when invoked, yields a fresh object with its own copy of the internal state defined by the module. The module def syntax indicates this; it is reminiscent of the def syntax for defining functions. The word functor that we used from this is based on a similar construct in Standard ML (and inspired by category theory, for the mathematically inclined). We can use cellasModule in a program as follows (examples/modules/cellModuleClient.wyv):

import modules.cellAsModule

```
val m1 = cellAsModule()
val m2 = cellAsModule()
m1.set(1)
m2.set(2)
m1.get() // evalutes to 1
m2.get() // evalutes to 2
```

In this example you can see that we have called the functor cellAsModule twice, and each resulting object instance has its own internal state.

11. Module Parameters

If objects are produced by functors (module functions), we expect to be able to pass parameters—and so we can. First let's define the type that cellAsModule returns. For convenience, we will put this type in a file TCellAsModule.wyt (here .wyt stands for Wyvern Type):

```
resource type TCellAsModule
   def set(newValue:Int):Unit
   def get():Int
```

```
Here is a client of the cellAsModule (examples/modules/cellClientFunctor.wyv):
module def cellClientFunctor(cell : modules.TCellAsModule)

def addOne():Unit
    cell.set(cell.get()+1)

def getValue():Int = cell.get()
```

We can put module parameters in between the parentheses in the definition of the functor, specifying the type in the usual way.

Now we can use cellAsModule together with cellClientFunctor in a program (examples/modules/cellClientMain.wyv):

```
import myPackage.cellAsModule
import myPackage.cellClientFunctor

val client = cellClientFunctor(cellAsModule())
client.addOne()
client.getValue() // evalutes to 1
```

12. Dynamic Types

Wyvern is intended to be a mostly statically typed language. However, while getting parameterized types to work, we implemented a Dyn type that partially implements dynamic types. Specifically, Dyn is a subtype of any type, and any type is a subtype of Dyn. (Note that subtyping is not transitive where Dyn is involved, as this would effectively collapse the type system to a single type.)

We recommend avoiding Dyn where possible, and now that parameterized types work well, we are gradually transitioning existing Dyn code to remove use of this construct. If we do keep this in the long term, we need to think about how it interacts with resource types.

13. Declaration Sequences and Mutual Recursion

Programs are made up of four kinds of core declarations: val, var, def, and type, as well as expressions. The val and var declarations and the expressions in a program are evaluated in sequence, and the variables defined earlier in the sequence are in scope in later declarations and expressions. In contrast, def and type declarations do not evaluate, and therefore these declaration forms can be safely used to define mutually recursive functions and types. Each sequence of declarations that consists exclusively of def and type is therefore treated as a mutually recursive block, so that the definition or type defined in each of the declarations is in scope in all the other declarations.

To understand why we allow recursive def and type declarations but do not allow this for val and var declarations, consider the following example:

```
// NOTE: this example does not typecheck
type IntCell
    def get():Int

def foo():IntCell = baz()
val bar:IntCell = foo()
def baz():IntCell = bar
bar.get()
```

When we try to initialize the bar value, we call foo(), which in turn invokes baz(). However, baz() reads the bar variable, which is what we are defining, so there is no well-defined result. Languages such as Java handle this by initializing bar to null at first and then writing a permanent value to it after the initializer executes. However, in order to avoid null pointer errors, Wyvern does not allow null as a value. Languages such as Haskell would use a special "black hole" value and signal a run-time error if the black hole is ever used, as in the bar.get statement at the end. We avoid this semantics as it adds complexity and means the program can fail at run time. Of course, infinite loops can still exist in Wyvern, but they come from recursive functions, never recursively defined values.

14. Some More Examples

The Wyvern standard library files are in subdirectories of stdlib. For example, earlier we used the IfTSL library, defined in stdlib/wyvern/IfTSL.wyv. When importing files from the standard library, leave out stdlib; i.e. in the example above, just use import wyvern.IfTSL

Platform-specific definitions are in the platform subdirectory of stdlib, in a sub-subdirectory named after the platform (e.g. java). For example, stdout is defined in stdlib/platform/java/stdout.wyv for the java platform; there is an analogous definition for the python platform too. The definition of stdout for java uses some Java helper code defined in the Java class wyvern.stdlib.support.Stdio. These libraries should be imported without using stdlib.platform.java.

An example of a utility library that provides a small part of a regular expression package is in wyvern/util/matching/regex.wyv. The design approximately follows the corresponding Scala library.

An example of a data structure library is wyvern/collections/list.wyv. Also see wyvern/option.wyv.

All of the above examples are tested by the Wyvern regression test suite that is run as part of ant test when building Wyvern.