

1 Bytecode Abstract Syntax

b	$::= v P \bar{i} M$	<i>bytecode file</i>	d	$::=$	$\text{val } f : \tau = v$ $\text{var } f : \tau = v$ $\text{def } m(\overline{x:\tau}) : \tau = e$ $\text{type } L = T$	<i>declarations</i>
v	$::= \text{major.minor}$	<i>version number</i>	T	$::=$	$c [\text{metadata } e]$ $\text{extag } c [\text{metadata } e]$ $\text{datatag } \overline{p.L} c [\text{metadata } e]$	<i>type desc.</i>
P	$::= \text{fully qualified path}$	<i>path to module</i>	c	$::=$	τ $\text{extends } p.L \tau$	<i>case desc.</i>
i	$::= \text{import } [\mu] \text{ URI} : \tau \text{ as } x$	<i>module import</i>	τ	$::=$	$\tau \{x \Rightarrow \bar{\sigma}\}_s$ $p.L$ \top $?$	<i>type</i>
μ	$::= \text{metadata} \mid \text{type}$		p	$::=$	x $p.f$	<i>paths</i>
M	$::=$ $\text{module } P : \tau = e$ \mid $\text{type } L = T$	<i>top level module</i>	s	$::=$	$\text{stateful} \mid \text{pure}$	
e	$::=$ x \mid $\text{new } \tau \{x \Rightarrow \bar{d}\}$ \mid $e.m(\bar{e})$ \mid $e.f$ \mid $e.f = e$ \mid $\text{let } x = e \text{ in } e$ \mid \mathcal{L} \mid $e.\text{match } \overline{x : p.L \Rightarrow e} [\text{else } e]$	<i>expressions</i>	σ	$::=$	$\text{val } f : \tau$ \mid $\text{var } f : \tau$ \mid $\text{def } m : \Pi \overline{x:\tau}. \tau$ \mid $\text{type } L = T$ \mid $\text{type}_s L$	<i>decl type</i>
\mathcal{L}	$::=$ string \mid integer	<i>literals</i>				
v	$::= x$	<i>values</i>				

Notation: overbar means a list of elements, as in Java