1 Grammar

$$
\begin{array}{rcll}
e & ::= & x & \textit{exprs.} \\
& | & r & \\
& | & \lambda x : \tau . e & \\
& | & e\ e & \\
& | & e.\pi & \\
\\
\hat{e} & ::= & x & \textit{labelled exprs.} \\
& | & r & \\
& | & \lambda x : \hat{\tau} . \hat{e} & \\
& | & \hat{e}\ \hat{e} & \\
& | & \hat{e}.\pi & \\
& | & \texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e & \\
\\
v & ::= & r & \textit{values.} \\
& | & \lambda x : \tau . e & \\
\\
\hat{v} & ::= & r & \textit{labelled values} \\
& | & \lambda x : \hat{\tau} . \hat{e} &
\end{array}
$$

$$
\begin{array}{rcll}
\varepsilon & ::= & \{\overline{r.\pi}\} & \textit{effects} \\
\\
\tau & ::= & \{\bar{r}\} & \textit{types} \\
& | & \tau \to \tau & \\
\\
\hat{\tau} & ::= & \{\bar{r}\} & \textit{labelled types} \\
& | & \hat{\tau} \to_{\varepsilon} \hat{\tau} & \\
\\
\Gamma & ::= & \varnothing & \textit{type ctx.} \\
& | & \Gamma, x : \tau & \\
\\
\hat{\Gamma} & ::= & \varnothing & \textit{labelled type ctx.} \\
& | & \hat{\Gamma}, x : \hat{\tau} &
\end{array}
$$

2 Functions

Definition ($\texttt{annot} :: \boldsymbol{\tau} \times \boldsymbol{\varepsilon} \to \boldsymbol{\hat{\tau}}$)

1. $\texttt{annot}(\{\bar{r}\}, \_) = \{\bar{r}\}$
2. $\texttt{annot}(\tau_1 \to \tau_2, \varepsilon) = \texttt{annot}(\tau_1, \varepsilon) \to_{\varepsilon} \texttt{annot}(\tau_2, \varepsilon)$

Definition ($\texttt{annot} :: \texttt{e} \times \boldsymbol{\varepsilon} \to \hat{\texttt{e}}$)

1. $\texttt{annot}(x, \_) = e$
2. $\texttt{annot}(r, \_) = r$
3. $\texttt{annot}(e_1 e_2, \varepsilon) = \texttt{annot}(e_1)\texttt{annot}(e_2)$
4. $\texttt{annot}(e.\pi, \varepsilon) = \texttt{annot}(e).\pi$
5. $\texttt{annot}(\lambda x : \tau.e, \varepsilon) = \lambda x : \texttt{annot}(\tau, \varepsilon).\texttt{annot}(e, \varepsilon)$

Definition ($\texttt{annot} :: \Gamma \times \boldsymbol{\varepsilon} \to \hat{\Gamma}$)

1. $\texttt{annot}(\varnothing, \_) = \varnothing$
2. $\texttt{annot}((\Gamma, x : \tau), \varepsilon) = \texttt{annot}(\Gamma, \varepsilon), x : \texttt{annot}(\tau, \varepsilon)$

Definition ($\texttt{erase} :: \boldsymbol{\hat{\tau}} \to \boldsymbol{\tau}$)

1. $\texttt{erase}(\{\bar{r}\}, \_) = \{\bar{r}\}$
2. $\texttt{erase}(\hat{\tau}_1 \to_{\varepsilon} \hat{\tau}_2) = \texttt{erase}(\hat{\tau}_1) \to \texttt{erase}(\hat{\tau}_2)$

Definition ($\texttt{erase} :: \hat{\texttt{e}} \to \texttt{e}$)

1. $\texttt{erase}(x) = x$
2. $\texttt{erase}(r) = r$
3. $\texttt{erase}(e_1 e_2) = \texttt{erase}(e_1)\texttt{erase}(e_2)$
4. $\texttt{erase}(e.\pi) = \texttt{erase}(e).\pi$
5. $\texttt{erase}(\lambda x : \hat{\tau}.\hat{e}) = \lambda x : \texttt{erase}(\hat{\tau}).\texttt{erase}(\hat{e})$

Definition ($\texttt{effects} :: \hat{\boldsymbol{\tau}} \to \boldsymbol{\varepsilon}$)

1. $\texttt{effects}(\{\bar{r}\}) = \{r.\pi \mid r \in \bar{r}, \pi \in \Pi\}$
2. $\texttt{effects}(\hat{\tau}_1 \to_\varepsilon \hat{\tau}_2) = \texttt{ho-effects}(\hat{\tau}_1) \cup \varepsilon \cup \texttt{effects}(\hat{\tau}_2)$

Definition ($\texttt{ho-effects} :: \hat{\boldsymbol{\tau}} \to \boldsymbol{\varepsilon}$)

1. $\texttt{ho-effects}(\{\bar{r}\}) = \varnothing$
2. $\texttt{ho-effects}(\hat{\tau}_1 \to_\varepsilon \hat{\tau}_2) = \texttt{effects}(\hat{\tau}_1) \cup \texttt{ho-effects}(\hat{\tau}_2)$

Definition ($\texttt{substitution} :: \hat{\texttt{e}} \times \hat{\texttt{v}} \times \hat{\texttt{v}} \to \hat{\texttt{e}}$)

The notation $[\hat{v}/x]\hat{e}$ is short-hand for $\texttt{substitution}(\hat{e}, \hat{v}, x)$. This function is partial, because the third-input must be a variable. We adopt the usual renaming conventions to avoid accidental capture.

1. $[\hat{v}/y]x = \hat{v}$, if $x = y$
2. $[\hat{v}/y]x = x$, if $x \neq y$
3. $[\hat{v}/y](\lambda x : \hat{\tau}.\hat{e}) = \lambda x : \hat{\tau}.[\hat{v}/y]\hat{e}$, if $y \neq x$ and y does not occur free in $\hat{e}$
4. $[\hat{v}/y](\hat{e}.\pi) = ([\hat{v}/y]\hat{e}).\pi$
5. $[\hat{v}/y](\hat{e}_1\hat{e}_2) = ([\hat{v}/y]\hat{e}_1)([\hat{v}/y]\hat{e}_2)$
6. $[\hat{v}/y](\texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e) = \texttt{import}(\varepsilon)\ x = [\hat{v}/y]\hat{e}\ \texttt{in}\ e$

When performing multiple substitutions we use the notation $[\hat{v}_1/x_1, \hat{v}_2/x_2]\hat{e}$ as shorthand for $[\hat{v}_2/x_2]([\hat{v}_1/x_1]\hat{e})$ (note the order of the variables has been flipped; the substitutions occur as they are written, left-to-right).

3 Static Rules

$\boxed{\Gamma \vdash e : \tau}$

$$\frac{}{\Gamma, x : \tau \vdash x : \tau}\ (\text{T-Var}) \qquad \frac{}{\Gamma, r : \{r\} \vdash r : \{r\}}\ (\text{T-Resource}) \qquad \frac{\Gamma, x : \tau_1 \vdash e : \tau_2}{\Gamma \vdash \lambda x : \tau_1.e : \tau_1 \to \tau_2}\ (\text{T-Abs})$$

$$\frac{\Gamma \vdash e_1 : \tau_2 \to \tau_3 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash e_1\ e_2 : \tau_3}\ (\text{T-App}) \qquad \frac{\Gamma \vdash e : \{\bar{r}\} \quad \forall r \in \bar{r} \mid r \in R \quad \pi \in \Pi}{\Gamma \vdash e.\pi : \texttt{Unit}}\ (\text{T-OperCall})$$

$\boxed{\hat{\Gamma} \vdash \hat{e} : \hat{\tau}\ \texttt{with}\ \varepsilon}$

$$\frac{}{\hat{\Gamma}, x : \tau \vdash x : \tau\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-Var}) \qquad \frac{}{\hat{\Gamma}, r : \{r\} \vdash r : \{r\}\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-Resource})$$

$$\frac{\hat{\Gamma}, x : \hat{\tau}_2 \vdash \hat{e} : \hat{\tau}_3\ \texttt{with}\ \varepsilon_3}{\hat{\Gamma} \vdash \lambda x : \tau_2.\hat{e} : \hat{\tau}_2 \to_{\varepsilon_3} \hat{\tau}_3\ \texttt{with}\ \varnothing}\ (\varepsilon\text{-Abs}) \qquad \frac{\hat{\Gamma} \vdash \hat{e}_1 : \hat{\tau}_2 \to_\varepsilon \hat{\tau}_3\ \texttt{with}\ \varepsilon_1 \quad \hat{\Gamma} \vdash \hat{e}_2 : \hat{\tau}_2\ \texttt{with}\ \varepsilon_2}{\hat{\Gamma} \vdash \hat{e}_1\hat{e}_2 : \hat{\tau}_3\ \texttt{with}\ \varepsilon_1 \cup \varepsilon_2 \cup \varepsilon}\ (\varepsilon\text{-App})$$

$$\frac{\hat{\Gamma} \vdash \hat{e} : \{\bar{r}\} \quad \forall r \in \bar{r} \mid r : \{r\} \in \Gamma \quad \pi \in \Pi}{\hat{\Gamma} \vdash \hat{e}.\pi : \texttt{Unit}\ \texttt{with}\ \{\bar{r}.\pi\}}\ (\varepsilon\text{-OperCall}) \qquad \frac{\hat{\Gamma} \vdash e : \tau\ \texttt{with}\ \varepsilon \quad \tau <: \tau' \quad \varepsilon \subseteq \varepsilon'}{\hat{\Gamma} \vdash e : \tau'\ \texttt{with}\ \varepsilon'}\ (\varepsilon\text{-Subsume})$$

$$\frac{\begin{array}{c}\hat{\Gamma} \vdash \hat{e} : \hat{\tau}\ \texttt{with}\ \varepsilon_1 \qquad \varepsilon = \texttt{effects}(\hat{\tau}) \\ \texttt{ho-safe}(\hat{\tau}, \varepsilon) \qquad x : \texttt{erase}(\hat{\tau}) \vdash e : \tau\end{array}}{\hat{\Gamma} \vdash \texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varepsilon \cup \varepsilon_1}\ (\varepsilon\text{-Module})$$

$\boxed{\texttt{safe}(\tau, \varepsilon)}$

$$\frac{}{\texttt{safe}(\{\bar{r}\},\varepsilon)}\ (\textsc{Safe-Resource}) \qquad \frac{}{\texttt{safe}(\texttt{Unit},\varepsilon)}\ (\textsc{Safe-Unit})$$

$$\frac{\varepsilon \subseteq \varepsilon' \quad \texttt{ho-safe}(\hat{\tau}_1,\varepsilon) \quad \texttt{safe}(\hat{\tau}_2,\varepsilon)}{\texttt{safe}(\hat{\tau}_1 \rightarrow_{\varepsilon'} \hat{\tau}_2,\varepsilon)}\ (\textsc{Safe-Arrow})$$

$\boxed{\texttt{ho-safe}(\hat{\tau},\varepsilon)}$

$$\frac{}{\texttt{ho-safe}(\{\bar{r}\},\varepsilon)}\ (\textsc{HOSafe-Resource}) \qquad \frac{}{\texttt{ho-safe}(\texttt{Unit},\varepsilon)}\ (\textsc{HOSafe-Unit})$$

$$\frac{\texttt{safe}(\hat{\tau}_1,\varepsilon) \quad \texttt{ho-safe}(\hat{\tau}_2,\varepsilon)}{\texttt{ho-safe}(\hat{\tau}_1 \rightarrow_{\varepsilon'} \hat{\tau}_2,\varepsilon)}\ (\textsc{HOSafe-Arrow})$$

$\boxed{\hat{\tau} <: \hat{\tau}}$

$$\frac{\varepsilon \subseteq \varepsilon' \quad \hat{\tau}_2 <: \hat{\tau}_2' \quad \hat{\tau}_1' <: \hat{\tau}_1}{\hat{\tau}_1 \rightarrow_{\varepsilon} \hat{\tau}_2 <: \hat{\tau}_1' \rightarrow_{\varepsilon'} \hat{\tau}_2'}\ (\textsc{S-Effects})$$

4 Dynamic Rules

$\boxed{\hat{e} \longrightarrow \hat{e} \mid \varepsilon}$

$$\frac{\hat{e}_1 \longrightarrow \hat{e}_1' \mid \varepsilon}{\hat{e}_1\hat{e}_2 \longrightarrow \hat{e}_1'\hat{e}_2 \mid \varepsilon}\ (\textsc{E-App1}) \qquad \frac{\hat{e}_2 \longrightarrow \hat{e}_2' \mid \varepsilon}{\hat{v}_1\hat{e}_2 \longrightarrow \hat{v}_1\hat{e}_2' \mid \varepsilon}\ (\textsc{E-App2}) \qquad \frac{}{(\lambda x:\hat{\tau}.\hat{e})\hat{v}_2 \longrightarrow [\hat{v}_2/x]\hat{e} \mid \varnothing}\ (\textsc{E-App3})$$

$$\frac{\hat{e} \rightarrow \hat{e}' \mid \varepsilon}{\hat{e}.\pi \longrightarrow \hat{e}'.\pi \mid \varepsilon}\ (\textsc{E-OperCall1}) \qquad \frac{r \in R \quad \pi \in \Pi}{r.\pi \longrightarrow \texttt{unit} \mid \{r.\pi\}}\ (\textsc{E-OperCall2})$$

$$\frac{\hat{e} \longrightarrow \hat{e}' \mid \varepsilon'}{\texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e \longrightarrow \texttt{import}(\varepsilon)\ x = \hat{e}'\ \texttt{in}\ e \mid \varepsilon'}\ (\textsc{E-Module1})$$

$$\frac{}{\texttt{import}(\varepsilon)\ x = \hat{v}\ \texttt{in}\ e \longrightarrow [\hat{v}/x]\texttt{annot}(e,\varepsilon) \mid \varnothing}\ (\textsc{E-Module2})$$

5 Encodings

5.1 $\perp$

We can define the bottom type as $\perp = \{\}$, because there is no empty-set literal.

5.2 `unit`, `Unit`

Define $\texttt{unit} = \lambda\texttt{x} : \{\}.\texttt{x}$, i.e. the function which takes an empty set of resources and returns it. We shall refer to its type, which is $\{\} \rightarrow_{\varnothing} \{\}$, as `Unit`. It has various properties befitting `unit`.

1. `unit` cannot be invoked, as $\{\}$ is uninhabited.
2. `unit` is a value.
3. The only term with type `Unit` is `unit`.
4. $\vdash \texttt{unit} : \texttt{Unit}$, by using ε-Abs and ε-Var.
5. $\texttt{effects}(\texttt{Unit}) = \texttt{ho-effects}(\texttt{Unit}) = \varnothing$
6. $\texttt{safe}(\texttt{Unit},\varepsilon)$ and $\texttt{ho-safe}(\texttt{Unit},\varepsilon)$

6 Proofs

Theorem 1 (Progress). *If $\hat{\Gamma} \vdash \hat{e}_A : \hat{\tau}_A$ `with` ε_A then $\hat{e}_A$ is a value or $\hat{e}_A \longrightarrow \hat{e}_B \mid \varepsilon$.*

Proof. By induction on $\hat{\Gamma} \vdash \hat{e}_A : \hat{\tau}_A$ `with` ε_A.

$\boxed{\text{Case: } \varepsilon\text{-Resource, } \varepsilon\text{-Unit, } \varepsilon\text{-Abs}}$ Then $\hat{e}_A$ is a value.

$\boxed{\text{Case: } \varepsilon\text{-Subsume}}$ Then $\hat{\Gamma} \vdash e : \tau'$ `with` ε', and $\hat{\Gamma} \vdash e : \tau$ `with` ε, where $\tau' <: \tau$ and $\varepsilon' \subseteq \varepsilon$ are subderivations. The theorem conclusion holds by inductive assumption applied to $\hat{\Gamma} \vdash e : \tau$ `with` ε.

$\boxed{\text{Case: } \varepsilon\text{-App}}$ Then $\hat{e}_A = \hat{e}_1\ \hat{e}_2$. We consider the cases in which $\hat{e}_1$ and $\hat{e}_2$ are values.

If $\hat{e}_1$ is not a value then by inductive assumption there is a reduction $\hat{e}_1 \longrightarrow \hat{e}_1' \mid \varepsilon$. Then $\hat{e}_1\ \hat{e}_2$ reduces by the rule E-App1, giving $\hat{e}_1\ \hat{e}_2 \longrightarrow \hat{e}_1'\ \hat{e}_2 \mid \varepsilon$.

If $\hat{e}_2$ is not a value then WLOG $\hat{e}_1$ is a value. By inductive assumption $\hat{e}_2 \longrightarrow \hat{e}_2' \mid \varepsilon$. Then $\hat{v}_1\ \hat{e}_2$ reduces by the rule E-App2, giving $\hat{v}_1\ \hat{e}_2 \longrightarrow \hat{v}_1\ \hat{e}_2' \mid \varepsilon$.

If $\hat{e}_1$ and $\hat{e}_2$ are both values then by canonical forms $\hat{e}_1 = \hat{v}_1 = \lambda x : \tau_2.e$. Then $\hat{v}_1\ \hat{v}_2$ reduces by the rule E-App3, giving $\hat{v}_1\ \hat{v}_2 \longrightarrow [\hat{v}_2/x]\hat{e} \mid \varnothing$.

$\boxed{\text{Case: } \varepsilon\text{-OperCall}}$ Then $\hat{e}_A = \hat{e}_1.\pi$. We consider whether $\hat{e}_1$ is a value.

If $\hat{e}_1$ is not a value then by inductive assumption there is a reduction $\hat{e}_1 \longrightarrow \hat{e}_1' \mid \varepsilon$. Then $\hat{e}_1.\pi$ reduces by the rule E-OperCall1, giving $\hat{e}_1.\pi \longrightarrow \hat{e}_1'.\pi \mid \varepsilon$.

If $\hat{e}_1$ is a value then $\hat{e}_1 = r$ by canonical forms. By the assumption that $r.\pi$ is closed under Γ, we know $r \in R$ and $\pi \in \Pi$. Then $\hat{e}_1.\pi$ reduces by the rule E-OperCall2, giving $r.\pi \longrightarrow$ `unit` $\mid \varepsilon$.

$\boxed{\text{Case: } \varepsilon\text{-Module}}$ Then $e_A =$ `import`$(\varepsilon)\ x = \hat{e}$ `in` e. If $\hat{e}$ is an expression then it can be reduced, so $\hat{e} \longrightarrow \hat{e}' \mid \varepsilon'$, and so by E-Module1 we get `import`$(\varepsilon)\ x = \hat{e}$ `in` $e \longrightarrow$ `import`$(\varepsilon)\ x = \hat{e}'$ `in` $e \mid \varepsilon'$. Otherwise $\hat{e} = \hat{v}$ is a value. Then by E-Module2 we get `import`$(\varepsilon)\ x = \hat{v} \longrightarrow [\hat{v}/x]$`annot`$(e, \varepsilon) \mid \varnothing$.

Lemma 1 (Substitution). *If $\hat{\Gamma}, x : \hat{\tau}' \vdash e : \hat{\tau}$ `with` ε and $\hat{\Gamma} \vdash \hat{v} : \hat{\tau}'$ `with` $\varnothing$ then $\hat{\Gamma} \vdash [\hat{v}/x]e : \hat{\tau}$ `with` ε.*

Proof. By induction on $\hat{\Gamma}, x : \hat{\tau}' \vdash e : \hat{\tau}$ `with` ε.

$\boxed{\varepsilon\text{-Var}}$ Then $\hat{e} = y$. Either $y = x$ or $y \neq x$.

Subcase: $y \neq x$. Then $[\hat{v}/x]y = y$ and $\hat{\Gamma} \vdash y : \hat{\tau}$ `with` $\varnothing$. Therefore $\hat{\Gamma} \vdash [\hat{v}/x]y : \hat{\tau}$ `with` $\varnothing$.

Subcase: $y = x$. By inversion on ε-Var, the original typing judgement is $\hat{\Gamma}, x : \hat{\tau}' \vdash x : \hat{\tau}'$ `with` $\varnothing$. Since $[\hat{v}/x]y = \hat{v}$ and by assumption $\hat{\Gamma} \vdash \hat{v} : \hat{\tau}'$ `with` $\varnothing$, then we have $\hat{\Gamma} \vdash [\hat{v}/x]x : \hat{\tau}'$ `with` $\varnothing$.

$\boxed{\varepsilon\text{-Resource}}$ Because $\hat{e} = r$ is a resource literal then $\hat{\Gamma} \vdash r : \hat{\tau}$ `with` $\varnothing$. By definition, $[\hat{v}/x]r = r$, so $\hat{\Gamma} \vdash [\hat{v}/x]r : \hat{\tau}$ `with` $\varnothing$.

$\boxed{\varepsilon\text{-Abs}}$ Then $\hat{\Gamma}, x : \hat{\tau}' \vdash \lambda z : \hat{\tau}_2.\hat{e}_{body} : \hat{\tau}_2 \rightarrow_{\varepsilon_3} \hat{\tau}_3$ `with` $\varnothing$. From inversion on ε-Abs we get the judgement $\hat{\Gamma}, x : \hat{\tau}', z : \hat{\tau}_2 \vdash \hat{e}_{body} : \hat{\tau}_3$ `with` ε_3. By applying the inductive assumption to $[\hat{v}/x]e_{body}$, we get $\hat{\Gamma}, z : \hat{\tau}_2 \vdash [\hat{v}/x]\hat{e}_{body} : \hat{\tau}_3$ `with` ε_3. Then applying ε-Abs, we get $\hat{\Gamma} \vdash \lambda z : \hat{\tau}_2.[\hat{v}/x]\hat{e}_{body} : \hat{\tau}_2 \rightarrow_{\varepsilon_3} \hat{\tau}_3$ `with` $\varnothing$. Then we are done, as $\lambda z : \hat{\tau}_2.[\hat{v}/x]\hat{e}_{body} = [\hat{v}/x](\lambda z : \hat{\tau}_2.\hat{e}_{body})$

$\boxed{\varepsilon\text{-App}}$ By inversion we know $\hat{\Gamma}, x : \hat{\tau}' \vdash \hat{e}_1 : \hat{\tau}_2 \rightarrow_{\varepsilon_3} \hat{\tau}_3$ `with` ε_A and $\hat{\Gamma}, x : \hat{\tau}' \vdash \hat{e}_2 : \hat{\tau}_2$ `with` ε_B, where $\varepsilon = \varepsilon_A \cup \varepsilon_B \cup \varepsilon_3$ and $\hat{\tau} = \hat{\tau}_3$. By inductive assumption, $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e}_1 : \hat{\tau}_2 \rightarrow_{\varepsilon_3} \hat{\tau}_3$ `with` ε_A and $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e}_2 : \hat{\tau}_2$ `with` ε_B. By ε-App we have $\hat{\Gamma} \vdash ([\hat{v}/x]\hat{e}_1)([\hat{v}/x]\hat{e}_2) : \hat{\tau}_3$ `with` $\varepsilon_A \cup \varepsilon_B \cup \varepsilon_3$. By simplifying and applying the definition of `substitution`, this is the same as $\hat{\Gamma} \vdash [\hat{v}/x](\hat{e}_1\hat{e}_2) : \hat{\tau}$ `with` ε.

$\boxed{\varepsilon\text{-OperCall}}$ By inversion we know $\hat{\Gamma}, x : \hat{\tau}' \vdash \hat{e}_1 : \{\bar{r}\}$ `with` ε_1, where $\varepsilon = \varepsilon_1 \cup \{r.\pi \mid r.\pi \in \bar{r} \times \Pi\}$ and $\hat{\tau} = \{\bar{r}\}$. By applying the inductive assumption, $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e}_1 : \{\bar{r}\}$ `with` ε_1. Then by ε-OperCall,

$\hat{\Gamma} \vdash ([\hat{v}/x]\hat{e}_1).\pi : \{\bar{r}\}$ `with` $\varepsilon_1 \cup \{r.\pi \mid r.\pi \in \bar{r} \times \Pi\}$. By simplifying and applying the definition of `substitution`, this is the same as $\hat{\Gamma} \vdash [\hat{v}/x](\hat{e}_1.\pi) : \hat{\tau}$ `with` ε.

$\boxed{\varepsilon\text{-Subsume}}$ By inversion we know $\hat{\Gamma}, x : \hat{\tau}' \vdash \hat{e} : \hat{\tau}_2$ `with` ε_2, where $\hat{\tau}_2 <: \hat{\tau}$ and $\varepsilon_2 \subseteq \varepsilon$. By inductive hypothesis, $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e} : \hat{\tau}_2$ `with` ε_2. Then by ε-Subsume we get $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e} : \hat{\tau}$ `with` ε.

$\boxed{\varepsilon\text{-Module}}$ Then $\hat{\Gamma}, x : \hat{\tau}' \vdash \texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e : \texttt{annot}(\tau, \varepsilon)$ `with` $\varepsilon \cup \varepsilon_1$. By inversion we know $\hat{\Gamma}, x : \hat{\tau}' \vdash \hat{e} : \hat{\tau}_1$ `with` ε_1. By inductive assumption, $\hat{\Gamma} \vdash [\hat{v}/x]\hat{e} : \hat{\tau}_1$ `with` ε_1. Then by ε-Module we have $\hat{\Gamma} \vdash \texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e : \texttt{annot}(\tau, \varepsilon)$ `with` $\varepsilon \cup \varepsilon_1$.

Lemma 2. *If* $\texttt{effects}(\hat{\tau}) \subseteq \varepsilon$ *and* $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$ *then* $\hat{\tau} <: \texttt{annot}(\texttt{erase}(\hat{\tau}), \varepsilon)$.

Lemma 3. *If* $\texttt{ho-effects}(\hat{\tau}) \subseteq \varepsilon$ *and* $\texttt{safe}(\tau, \varepsilon)$ *then* $\texttt{annot}(\texttt{erase}(\hat{\tau}), \varepsilon) <: \hat{\tau}$.

Proof. By simultaneous induction.

$\boxed{\text{Case: } \hat{\tau} = \{\bar{r}\}}$ Then $\hat{\tau} = \texttt{annot}(\texttt{erase}(\hat{\tau}), \varepsilon)$ and the results for both lemmas hold immediately.

$\boxed{\text{Case: } \hat{\tau} = \hat{\tau}_1 \to_{\varepsilon'} \hat{\tau}_2,\ \texttt{effects}(\hat{\tau}) \subseteq \varepsilon,\ \texttt{ho-safe}(\hat{\tau}, \varepsilon)}$ It is sufficient to show $\hat{\tau}_2 <: \texttt{annot}(\texttt{erase}(\hat{\tau}_2), \varepsilon)$ and $\texttt{annot}(\texttt{erase}(\hat{\tau}_1), \varepsilon) <: \hat{\tau}_1$, because the result will hold by S-Effects. To achieve this we shall inductively apply **lemma 2** to $\hat{\tau}_2$ and **lemma 3** to $\hat{\tau}_1$.

From $\texttt{effects}(\hat{\tau}) \subseteq \varepsilon$ we have $\texttt{ho-effects}(\hat{\tau}_1) \cup \varepsilon' \cup \texttt{effects}(\hat{\tau}_2) \subseteq \varepsilon$ and therefore $\texttt{effects}(\hat{\tau}_2) \subseteq \varepsilon$. From $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$ we have $\texttt{ho-safe}(\hat{\tau}_2, \varepsilon)$. Therefore we can apply **lemma 2** to $\hat{\tau}_2$.

From $\texttt{effects}(\hat{\tau}) \subseteq \varepsilon$ we have $\texttt{ho-effects}(\hat{\tau}_1) \cup \varepsilon' \cup \texttt{effects}(\hat{\tau}_2) \subseteq \varepsilon$ and therefore $\texttt{ho-effects}(\hat{\tau}_1) \subseteq \varepsilon$. From $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$ we have $\texttt{ho-safe}(\hat{\tau}_1, \varepsilon)$. Therefore we can apply **lemma 3** to $\hat{\tau}_1$.

$\boxed{\text{Case: } \hat{\tau} = \hat{\tau}_1 \to_{\varepsilon'} \hat{\tau}_2,\ \texttt{ho-effects}(\hat{\tau}) \subseteq \varepsilon,\ \texttt{safe}(\hat{\tau}, \varepsilon)}$ It is sufficient to show $\texttt{annot}(\texttt{erase}(\hat{\tau}_2), \varepsilon) <: \hat{\tau}_2$ and $\hat{\tau}_1 <: \texttt{annot}(\texttt{erase}(\hat{\tau}_1), \varepsilon)$, because the result will hold by S-Effects. To achieve this we shall inductively apply **lemma 3** to $\hat{\tau}_2$ and **lemma 2** to $\hat{\tau}_1$.

From $\texttt{ho-effects}(\hat{\tau}) \subseteq \varepsilon$ we have $\texttt{effects}(\hat{\tau}_1) \cup \texttt{ho-effects}(\hat{\tau}_2) \subseteq \varepsilon$ and therefore $\texttt{ho-effects}(\hat{\tau}_2) \subseteq \varepsilon$. From $\texttt{safe}(\hat{\tau}, \varepsilon)$ we have $\texttt{safe}(\hat{\tau}_2, \varepsilon)$. Therefore we can apply **lemma 3** to $\hat{\tau}_2$.

From $\texttt{ho-effects}(\hat{\tau}) \subseteq \varepsilon$ we have $\texttt{effects}(\hat{\tau}_1) \cup \texttt{ho-effects}(\hat{\tau}_2) \subseteq \varepsilon$ and therefore $\texttt{effects}(\hat{\tau}_1) \subseteq \varepsilon$. From $\texttt{safe}(\hat{\tau}, \varepsilon)$ we have $\texttt{ho-safe}(\hat{\tau}_1, \varepsilon)$. Therefore we can apply **lemma 2** to $\hat{\tau}_1$.

Theorem 2 (Preservation). *If* $\hat{\Gamma} \vdash \hat{e}_A : \hat{\tau}_A$ `with` ε_A *and* $e_A \longrightarrow e_B \mid \varepsilon_C$, *then* $\hat{\Gamma} \vdash e_B : \tau_B$ `with` ε_B, *where* $e_B <: e_B$ *and* $\varepsilon \cup \varepsilon_B \subseteq \varepsilon_A$.

Proof. By induction on $\hat{\Gamma} \vdash \hat{e}_A : \tau_A$ `with` ε_A, and then on $e_A \longrightarrow e_B \mid \varepsilon$.

$\boxed{\varepsilon\text{-Var}, \varepsilon\text{-Resource}, \varepsilon\text{-Unit}, \varepsilon\text{-Abs}}$ Then e_A cannot be reduced and so the theorem statement vacuously holds.

$\boxed{\varepsilon\text{-App}}$ Then $e_A = \hat{e}_1\hat{e}_2$ and $\hat{e}_1 : \hat{\tau}_2 \to_\varepsilon \hat{\tau}_3$ `with` ε_1 and $\hat{\Gamma} \vdash \hat{e}_2 : \hat{\tau}_2$ `with` ε_2. If the reduction rule used was E-App1 or E-App2, then the result follows by applying the inductive hypothesis to $\hat{e}_1$ and $\hat{e}_2$ respectively.

Otherwise the rule used was E-App3. Then $(\lambda x : \hat{\tau}_2.\hat{e})\hat{v}_2 \longrightarrow [\hat{v}_2/x]\hat{e} \mid \varnothing$. By inversion on the typing rule for $\lambda x : \hat{\tau}_2.\hat{e}$ we know $\Gamma, x : \hat{\tau}_2 \vdash \hat{e} : \hat{\tau}_3$ `with` ε_3. By canonical forms, $\varepsilon_2 = \varnothing$ because $\hat{e}_2 = \hat{v}_2$ is a value. Then by the substitution lemma, $\hat{\Gamma} \vdash [\hat{v}_2/x]\hat{e} : \hat{\tau}_3$ `with` ε_3. By canonical forms, $\varepsilon_1 = \varepsilon_2 = \varnothing = \varepsilon_C$. Therefore

$\varepsilon_A = \varepsilon_3 = \varepsilon_B \cup \varepsilon_C$.

$\boxed{\varepsilon\text{-OperCall}}$ Then $e_A = e_1.\pi$ and $\hat{\Gamma} \vdash e_1 : \{\bar{r}\}$ $\texttt{with}$ ε_1. If the reduction rule used was E-OperCall1 then the result follows by applying the inductive hypothesis to $\hat{e}_1$.

Otherwise the reduction rule used was E-OperCall2 and $v_1.\pi \longrightarrow \texttt{unit} \mid \{r.\pi\}$. By canonical forms, $\hat{\Gamma} \vdash v_1 : \texttt{unit with } \{r.\pi\}$. Also, $\hat{\Gamma} \vdash \texttt{unit} : \texttt{Unit with } \varnothing$. Then $\tau_B = \tau_A$. Also, $\varepsilon_C \cup \varepsilon_B = \{r.\pi\} = \varepsilon_A$.

$\boxed{\varepsilon\text{-Module}}$ Then $e_A = \texttt{import}(\varepsilon)\ x = \hat{e}\ \texttt{in}\ e$. If the reduction rule used was E-ModuleCall1 then the result follows by applying the inductive hypothesis to $\hat{e}$.

Otherwise $\hat{e}$ is a value and the reduction used was E-ModuleCall2. The following are true:

1. $e_A = \texttt{import}(\varepsilon)\ x = \hat{v}\ \texttt{in}\ e$
2. $\hat{\Gamma} \vdash e_A : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varepsilon \cup \varepsilon_1$
3. $\texttt{import}(\varepsilon)\ x = \hat{v}\ \texttt{in}\ e \longrightarrow [\hat{v}/x]\texttt{annot}(e, \varepsilon) \mid \varnothing$
4. $\hat{\Gamma} \vdash \hat{v} : \hat{\tau}\ \texttt{with}\ \varnothing$
5. $\varepsilon = \texttt{effects}(\hat{\tau})$
6. $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$
7. $x : \texttt{erase}(\hat{\tau}) \vdash e : \tau$

Apply the **annotation lemma** with $\Gamma = \varnothing$ to get $\hat{\Gamma}, x : \hat{\tau} \vdash \texttt{annot}(e, \varepsilon) : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varepsilon$.

By **4.** we have $\hat{\Gamma} \vdash \hat{v} : \hat{\tau}\ \texttt{with}\ \varnothing$.

By **substitution lemma**, $\hat{\Gamma} \vdash [\hat{v}/x]\texttt{annot}(e, \varepsilon) : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varepsilon$.

By **canonical forms**, $\varepsilon_1 = \varepsilon_C = \varnothing$. Then $\varepsilon_B = \varepsilon = \varepsilon_A \cup \varepsilon_C$. By examination, $\tau_A = \tau_B = \texttt{annot}(\tau, \varepsilon)$.

Lemma 4 (Annotation). *If the following are true:*

- $\hat{\Gamma} \vdash \hat{v} : \hat{\tau}\ \texttt{with}\ \varnothing$
- $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e : \tau$
- $\varepsilon = \texttt{effects}(\hat{\tau})$
- $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$

Then $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e, \varepsilon) : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varepsilon \cup \texttt{effects}(\texttt{annot}(\Gamma, \varepsilon))$.

Proof. By induction on $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e : \tau$.

$\boxed{\text{Case: T-Var}}$ Then $e = x$ and $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash x : \tau$. There are two cases: $x = y$ or $x \neq y$.

Subcase 1: $x = y$. Then by ε-Var we get $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash x : \hat{\tau}\ \texttt{with}\ \varnothing$. First note that $\texttt{annot}(x, \varepsilon) = x$ in this case. Therefore $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash \texttt{annot}(\texttt{erase}(x), \varepsilon) : \hat{\tau}\ \texttt{with}\ \varnothing$. We know by assumption that $\texttt{effects}(\hat{\tau}) = \varepsilon$ and $\texttt{ho-safe}(\hat{\tau}, \varepsilon)$. Applying **Lemma 2** we know $\hat{\tau} <: \texttt{annot}(\texttt{erase}(\hat{\tau}), \varepsilon)$. Lastly, by ε-Subsume we have $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash \texttt{annot}(\texttt{erase}(x), \varepsilon) : \texttt{annot}(\texttt{erase}(x), \varepsilon)\ \texttt{with}\ \varepsilon \cup \texttt{effects}\ (\texttt{annot}(\Gamma, \varepsilon))$.

Subcase 2: $x \neq y$. Then $x : \tau \in \Gamma$. Together with the definition $\texttt{annot}(x, \varepsilon) = x$, we know $x : \texttt{annot}(\tau, \varepsilon) \in \texttt{annot}(\Gamma, \varepsilon)$. By ε-Var we have $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(x, \varepsilon) : \texttt{annot}(\tau, \varepsilon)\ \texttt{with}\ \varnothing$. Lastly, by ε-Subsume we have $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash \texttt{annot}(\texttt{erase}(x), \varepsilon) : \texttt{annot}(\texttt{erase}(x), \varepsilon)\ \texttt{with}\ \varepsilon \cup \texttt{effects}\ (\texttt{annot}(\Gamma, \varepsilon))$.

$\boxed{\text{Case: T-Resource}}$ Then $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash r : \{r\}$. By definition, $\texttt{annot}(r, \varepsilon) = r$ and $\texttt{annot}(\{r\}, \varepsilon)$. By ε-Resource $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash r : \{r\}\ \texttt{with}\ \varnothing$. By ε-Subsume, $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash r : \{r\}\ \texttt{with}\ \varepsilon \cup \texttt{effects}(\texttt{annot}(\Gamma, \varepsilon))$.

$\boxed{\text{Case: T-Abs}}$ Then $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash \lambda x : \tau_1.e_{body} : \tau_1 \to \tau_2$.

By inversion, we get the sub-derivation $\Gamma, y : \texttt{erase}(\hat{\tau}), x : \tau_1 \vdash e_2 : \tau_2$. By definition, $\texttt{annot}(e, \varepsilon) = \texttt{annot}(\lambda x : \tau_1.e_2, \varepsilon) = \lambda x : \texttt{annot}(\tau_1, \varepsilon).\texttt{annot}(e_2, \varepsilon)$ and $\texttt{annot}(\tau, \varepsilon) = \texttt{annot}(\tau_1 \to \tau_2, \varepsilon) = \texttt{annot}(\tau_1, \varepsilon) \to_\varepsilon \texttt{annot}(\tau_2, \varepsilon)$.

To apply the inductive assumption to e_2 we use the unlabelled context $\Gamma, x : \tau_1$. The inductive assumption tells us $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau}, x : \texttt{annot}(\tau_1, \varepsilon) \vdash \texttt{annot}(e_2, \varepsilon) : \texttt{annot}(\tau_2, \varepsilon)$ $\texttt{with}$ $\varepsilon \cup \texttt{effects}(\texttt{annot}(\Gamma, \varepsilon)) \cup \texttt{effects}(\texttt{annot}(\tau_1, \varepsilon))$. Call this last effect-set ε'.

By ε-Abs, we get $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \lambda x : \texttt{annot}(\tau_1, \varepsilon).\texttt{annot}(e_2, \varepsilon) : \texttt{annot}(\hat{\tau}_1) \to_{\varepsilon'} \texttt{annot}(\hat{\tau}_2)$ $\texttt{with}$ $\varnothing$.

By ε-Subsume, we get $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e, \varepsilon) : \texttt{annot}(\hat{\tau}_1) \to_\varepsilon \texttt{annot}(\hat{\tau}_2)$ $\texttt{with}$ $\varepsilon \cup \texttt{effects}(\texttt{annot}(\Gamma), \varepsilon)$.

Case: T-App Then $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_1\ e_2 : \tau_3$, where $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_1 : \tau_2 \to \tau_3$ and $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_2 : \tau_2$.

By applying the inductive assumption to e_1 and e_2, we get $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e_1, \varepsilon) : \texttt{annot}(\tau_1, \varepsilon)$ $\texttt{with}$ ε and $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e_2, \varepsilon) : \texttt{annot}(\tau_2, \varepsilon)$ $\texttt{with}$ ε.

By simplifying: $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e_1, \varepsilon) : \texttt{annot}(\tau_2, \varepsilon) \to_\varepsilon \texttt{annot}(\tau_3, \varepsilon)$ $\texttt{with}$ ε.

By ε-App, we get $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash \texttt{annot}(e_1\ e_2, \varepsilon) : \texttt{annot}(\tau_3, \varepsilon)$ $\texttt{with}$ ε.

Case: T-OperCall Then $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_1.\pi : \texttt{Unit}$.

By inversion we get the sub-derivation $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_1 : \{\bar{r}\}$.

By definition, $\texttt{annot}(\{\bar{r}\}, \varepsilon) = \{\bar{r}\}$.

By inductive assumption, $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash e_1 : \{\bar{r}\}$ $\texttt{with}$ $\varepsilon \cup \texttt{effects}(\texttt{annot}(\Gamma, \varepsilon))$.

By ε-OperCall, $\hat{\Gamma}, \texttt{annot}(\Gamma, \varepsilon), y : \hat{\tau} \vdash e_1.\pi : \{\bar{r}\}$ $\texttt{with}$ $\varepsilon \cup \{\bar{r}.\pi\}$.

It remains to show $\{\bar{r}.\pi\} \subseteq \varepsilon$. We shall do this by considering where r must have come from (which subcontext left of the turnstile).

Subcase 1. $r = \hat{\tau}$. As $\varepsilon = \texttt{effects}(\hat{\tau})$, then $r.\pi \in \texttt{effects}(\hat{\tau})$.

Subcase 2. $r : \{r\} \in \Gamma$. As $\texttt{annot}(r, \varepsilon) = r$, then $r.\pi \in \texttt{annot}(\Gamma, \varepsilon)$.

Subcase 3. $r : \{r\} \in \hat{\Gamma}$. Then because $\Gamma, y : \texttt{erase}(\hat{\tau}) \vdash e_1 : \{\bar{r}\}$, then $r \in \Gamma$ or $r = \texttt{erase}(\hat{\tau}) = \hat{\tau}$ and one of the above subcases must also hold.