

1 Bytecode Abstract Syntax

$b ::= v \ P \ \bar{i} \ \overline{M}$ *bytecode file*

$v ::= \text{magic major.minor}$ *magic+version number*

$P ::= \text{fully qualified path}$ *path to module*

$i ::= \text{import } \mu \ \text{URI} : \tau \text{ as } x$ *module import*

$\mu ::= [\text{metadata}] \ [\text{type}]$

$M ::= \begin{array}{l} \text{module } P : \tau = e \\ \text{type } P = T \ [\delta] \end{array}$ *top level modules*

$e ::= \begin{array}{l} x \\ \text{new } \tau \ \{x \Rightarrow \bar{d}\} \\ e.m(\bar{e}) \\ e.f \\ e.f = e \\ \mathcal{L} \\ \overline{e.\text{match } x : p.L \Rightarrow e \ [\text{else } e]} \\ \overline{e \mid d} \end{array}$ *expressions*

$\mathcal{L} ::= \begin{array}{l} \text{string} \\ \text{integer} \end{array}$ *literals*

$d ::= \begin{array}{l} \text{val } f : \tau = e \\ \text{var } f : \tau = e \\ \text{def } m(\overline{x : \tau}) : \tau = e \\ \text{type } L = T \ [\delta] \end{array}$ *declarations*

$T ::= \begin{array}{l} c \\ \text{extag } c \\ \text{datatag } \overline{p.L} \ c \end{array}$ *type desc.*

$c ::= \begin{array}{l} \tau \\ \text{extends } p.L \ \tau \end{array}$ *case desc.*

$\tau ::= \begin{array}{l} \tau \ \{\mathbf{x} \Rightarrow \bar{\sigma}\}_s \\ p.L \\ \top \\ \perp \\ ? \end{array}$ *type*

$p ::= \begin{array}{l} x \\ p.f \end{array}$ *paths*

$s ::= \text{stateful} \mid \text{pure}$

$\sigma ::= \begin{array}{l} \text{val } f : \tau \\ \text{var } f : \tau \\ \text{def } m : \Pi \overline{x : \tau}. \tau \\ \text{type } L = T \ [\delta] \\ \text{type}_s L \ [\delta] \end{array}$ *decl type*

$\delta ::= \text{metadata } e$ *metadata*

Notation: overbar means a list of elements, as in Java