

## 1 Grammar

$e ::=$	<b>exprs.</b>	$\tau ::=$	<b>types</b>
$x$	<i>variable</i>	$t$	<i>type variable</i>
$v$	<i>value</i>	$\{\bar{r}\}$	<i>effect set</i>
$e e$	<i>application</i>	$\tau \rightarrow \tau$	<i>arrow</i>
$e.\pi$	<i>operation call</i>		
$\forall t : \tau. e$	<i>type polymorphism</i>	$\hat{\tau} ::=$	<b>annotated types</b>
		$t$	<i>type variable</i>
$v ::=$	<b>values</b>	$\{\bar{r}\}$	<i>resource set</i>
$r$	<i>resource literal</i>	$\hat{\tau} \rightarrow_{\varepsilon} \hat{\tau}$	<i>annotated arrow</i>
$\lambda x : \tau. e$	<i>abstraction</i>		
		$\varepsilon ::=$	<b>effects</b>
$\hat{e} ::=$	<b>annotated exprs.</b>	$\epsilon$	<i>effect variable</i>
$x$	<i>variable</i>	$\{\bar{r}.\pi\}$	<i>effect set</i>
$\hat{v}$	<i>value</i>		
$\hat{e} \hat{e}$	<i>application</i>	$\Gamma ::=$	<b>contexts</b>
$\hat{e}.\pi$	<i>operation call</i>	$\emptyset$	<i>empty ctx.</i>
$\forall t : \hat{\tau}.\hat{e}$	<i>type polymorphism</i>	$\Gamma, x : \tau$	<i>var. binding</i>
$\forall \epsilon : \varepsilon.\hat{e}$	<i>effect polymorphism</i>	$\Gamma, t : \tau$	<i>type var. binding</i>
$\text{import}(\varepsilon_s) x = \hat{e} \text{ in } e$	<i>import</i>		
		$\hat{\Gamma} ::=$	<b>annotated contexts</b>
$\hat{v} ::=$	<b>annotated values</b>	$\emptyset$	<i>empty ctx.</i>
$r$	<i>resource literal</i>	$\hat{\Gamma}, x : \hat{\tau}$	<i>var. binding</i>
$\lambda x : \hat{\tau}.\hat{e}$	<i>abstraction</i>	$\hat{\Gamma}, \epsilon : \varepsilon$	<i>effect var. binding</i>
		$\hat{\Gamma}, t : \hat{\tau}$	<i>type var. binding</i>

## 2 Functions

**Definition** ( $\text{annot} :: \tau \times \varepsilon \rightarrow \hat{\tau}$ )

1.  $\text{annot}(t, \_) = t$
2.  $\text{annot}(\{\bar{r}\}, \_) = \{\bar{r}\}$
3.  $\text{annot}(\tau_1 \rightarrow \tau_2, \varepsilon) = \text{annot}(\tau_1, \varepsilon) \rightarrow_{\varepsilon} \text{annot}(\tau_2, \varepsilon)$

**Definition** ( $\text{annot} :: e \times \varepsilon \rightarrow \hat{e}$ )

1.  $\text{annot}(x, \_) = x$
2.  $\text{annot}(r, \_) = r$
3.  $\text{annot}(\lambda x : \tau. e, \varepsilon) = \lambda x : \text{annot}(\tau, \varepsilon). \text{annot}(e, \varepsilon)$
4.  $\text{annot}(e_1 e_2, \varepsilon) = \text{annot}(e_1, \varepsilon) \text{annot}(e_2, \varepsilon)$
5.  $\text{annot}(e.\pi, \varepsilon) = \text{annot}(e, \varepsilon).\pi$
6.  $\text{annot}(\forall t : \tau. e, \varepsilon) = \forall t : \text{annot}(\tau, \varepsilon). \text{annot}(e, \varepsilon)$
7.  $\text{annot}(\forall \epsilon : \varepsilon'. e, \varepsilon) = \forall \epsilon : \varepsilon'. \text{annot}(e, \varepsilon)$

**Definition** ( $\text{annot} :: \Gamma \times \varepsilon \rightarrow \hat{\Gamma}$ )

1.  $\text{annot}(\emptyset, \_) = \emptyset$
2.  $\text{annot}((\Gamma, x : \tau), \varepsilon) = \text{annot}(\Gamma, \varepsilon), x : \text{annot}(\tau, \varepsilon)$
3.  $\text{annot}((\Gamma, t : \tau), \varepsilon) = \text{annot}(\Gamma, \varepsilon), x : \text{annot}(\tau, \varepsilon)$

**Definition** ( $\text{erase} :: \hat{\tau} \rightarrow \tau$ )

1.  $\text{erase}(t) = t$
2.  $\text{erase}(\{\bar{r}\}, \_) = \{\bar{r}\}$
3.  $\text{erase}(\hat{\tau}_1 \rightarrow_{\varepsilon} \hat{\tau}_2) = \text{erase}(\hat{\tau}_1) \rightarrow \text{erase}(\hat{\tau}_2)$

**Definition** ( $\text{erase} :: \hat{e} \rightarrow e$ )

1.  $\text{erase}(x) = x$
2.  $\text{erase}(r) = r$
3.  $\text{erase}(\lambda x : \hat{\tau}. \hat{e}) = \lambda x : \text{erase}(\hat{\tau}). \text{erase}(\hat{e})$
4.  $\text{erase}(e_1 \ e_2) = \text{erase}(e_1) \text{erase}(e_2)$
5.  $\text{erase}(e. \pi) = \text{erase}(e). \pi$
6.  $\text{erase}(\forall t : \hat{\tau}. \hat{e}) = \forall t : \text{erase}(\hat{\tau}). \text{erase}(\hat{e})$
7.  $\text{erase}(\forall \epsilon : \varepsilon. \hat{e}) = \forall \epsilon : \varepsilon. \text{erase}(\hat{e})$

**Definition** ( $\text{erase} :: \hat{\Gamma} \rightarrow \Gamma$ )

1.  $\text{erase}(\emptyset) = \emptyset$
2.  $\text{erase}(\hat{\Gamma}, x : \hat{\tau}) = \text{erase}(\hat{\Gamma}), x : \text{erase}(\hat{\tau})$
3.  $\text{erase}(\Gamma, t : \hat{\tau}) = \text{erase}(\Gamma), x : \text{erase}(\hat{\tau})$
4.  $\text{erase}(\Gamma, \epsilon : \varepsilon) = \text{erase}(\Gamma)$

**Definition** ( $\text{effects} :: \hat{\tau} \rightarrow \varepsilon$ )

1.  $\text{effects}(t) = \emptyset$
2.  $\text{effects}(\{\bar{r}\}) = \{r. \pi \mid r \in \bar{r}, \pi \in \Pi\}$
3.  $\text{effects}(\hat{\tau}_1 \rightarrow_{\varepsilon} \hat{\tau}_2) = \text{ho-effects}(\hat{\tau}_1) \cup \varepsilon \cup \text{effects}(\hat{\tau}_2)$

**Definition** ( $\text{ho-effects} :: \hat{\tau} \rightarrow \varepsilon$ )

1.  $\text{ho-effects}(t) = \emptyset$
2.  $\text{ho-effects}(\{\bar{r}\}) = \emptyset$
3.  $\text{ho-effects}(\hat{\tau}_1 \rightarrow_{\varepsilon} \hat{\tau}_2) = \text{effects}(\hat{\tau}_1) \cup \text{ho-effects}(\hat{\tau}_2)$

**Definition** ( $\text{substitution} :: \hat{e} \times \hat{\nu} \times \hat{\nu} \rightarrow \hat{e}$ )

The notation  $[\hat{v}/x]\hat{e}$  is short-hand for  $\text{substitution}(\hat{e}, \hat{v}, x)$ . This function is partial, because the third input must be a variable. We adopt the usual renaming conventions to avoid accidental capture.

1.  $[\hat{v}/y]x = \hat{v}$ , if  $x = y$
2.  $[\hat{v}/y]x = x$ , if  $x \neq y$
3.  $[\hat{v}/y](\lambda x : \hat{\tau}. \hat{e}) = \lambda x : \hat{\tau}. [\hat{v}/y]\hat{e}$ , if  $y \neq x$  and  $y$  does not occur free in  $\hat{e}$
4.  $[\hat{v}/y](\hat{e}_1 \ \hat{e}_2) = ([\hat{v}/y]\hat{e}_1) ([\hat{v}/y]\hat{e}_2)$
5.  $[\hat{v}/y](\hat{e}. \pi) = ([\hat{v}/y]\hat{e}). \pi$
6.  $[\hat{v}/y](\forall t : \hat{\tau}. \hat{e}) = \forall t : \hat{\tau}. [\hat{v}/y]\hat{e}$ , if  $y \neq t$  and  $y$  does not occur free in  $\hat{e}$
7.  $[\hat{v}/y](\forall \epsilon : \varepsilon. \hat{e}) = \forall \epsilon : \varepsilon. [\hat{v}/y]\hat{e}$ , if  $y \neq \epsilon$  and  $y$  does not occur free in  $\hat{e}$
8.  $[\hat{v}/y](\text{import}(\varepsilon_s) \ x = \hat{e} \ \text{in} \ e) = \text{import}(\varepsilon_s) \ x = [\hat{v}/y]\hat{e} \ \text{in} \ e$

When performing multiple substitutions the notation  $[\hat{v}_1/x_1, \hat{v}_2/x_2]\hat{e}$  is used as shorthand for  $[\hat{v}_2/x_2]([\hat{v}_1/x_1]\hat{e})$  (note the order of the variables has been flipped; the substitutions occur as they are written, left-to-right).

### 3 Static Rules

$$\boxed{\Gamma \vdash e : \tau}$$

$$\begin{array}{c} \frac{}{\Gamma, x : \tau \vdash x : \tau} \text{ (T-VAR)} \quad \frac{}{\Gamma, r : \{r\} \vdash r : \{r\}} \text{ (T-RESOURCE)} \quad \frac{\Gamma, x : \tau_1 \vdash e : \tau_2}{\Gamma \vdash \lambda x : \tau_1. e : \tau_1 \rightarrow \tau_2} \text{ (T-ABS)} \\[10pt] \frac{\Gamma \vdash e_1 : \tau_2 \rightarrow \tau_3 \quad \Gamma \vdash e_2 : \tau_2}{\Gamma \vdash e_1 \ e_2 : \tau_3} \text{ (T-APP)} \quad \frac{\Gamma \vdash e : \{\bar{r}\}}{\Gamma \vdash e. \pi : \text{Unit}} \text{ (T-OPERCALL)} \end{array}$$

$$\boxed{\hat{I} \vdash \hat{e} : \hat{\tau} \text{ with } \varepsilon}$$

$$\begin{array}{c}
\frac{}{\hat{I}, x : \tau \vdash x : \tau \text{ with } \emptyset} (\varepsilon\text{-VAR}) \quad \frac{}{\hat{I}, r : \{r\} \vdash r : \{r\} \text{ with } \emptyset} (\varepsilon\text{-RESOURCE}) \\
\\
\frac{\hat{I}, x : \hat{\tau}_2 \vdash \hat{e} : \hat{\tau}_3 \text{ with } \varepsilon_3}{\hat{I} \vdash \lambda x : \tau_2. \hat{e} : \hat{\tau}_2 \rightarrow_{\varepsilon_3} \hat{\tau}_3 \text{ with } \emptyset} (\varepsilon\text{-ABS}) \\
\\
\frac{\hat{I} \vdash \hat{e}_1 : \hat{\tau}_2 \rightarrow_{\varepsilon} \hat{\tau}_3 \text{ with } \varepsilon_1 \quad \hat{I} \vdash \hat{e}_2 : \hat{\tau}_2 \text{ with } \varepsilon_2}{\hat{I} \vdash \hat{e}_1 \hat{e}_2 : \hat{\tau}_3 \text{ with } \varepsilon_1 \cup \varepsilon_2 \cup \varepsilon} (\varepsilon\text{-APP}) \\
\\
\frac{\hat{I} \vdash \hat{e} : \{\bar{r}\}}{\hat{I} \vdash \hat{e}. \pi : \text{Unit} \text{ with } \{r. \pi \mid r \in \bar{r}\}} (\varepsilon\text{-OPERCALL}) \quad \frac{\hat{I} \vdash e : \tau \text{ with } \varepsilon \quad \tau <: \tau' \quad \varepsilon \subseteq \varepsilon'}{\hat{I} \vdash e : \tau' \text{ with } \varepsilon'} (\varepsilon\text{-SUBSUME}) \\
\\
\frac{\text{effects}(\hat{\tau}) \cup \text{ho-effects}(\text{annot}(\tau, \emptyset)) \subseteq \varepsilon \quad \hat{I} \vdash \hat{e} : \hat{\tau} \text{ with } \varepsilon_1 \quad \text{ho-safe}(\hat{\tau}, \varepsilon) \quad x : \text{erase}(\hat{\tau}) \vdash e : \tau}{\hat{I} \vdash \text{import}(\varepsilon) x = \hat{e} \text{ in } e : \text{annot}(\tau, \varepsilon) \text{ with } \varepsilon \cup \varepsilon_1} (\varepsilon\text{-IMPORT})
\end{array}$$

$$\boxed{\text{safe}(\tau, \varepsilon)}$$

$$\begin{array}{c}
\frac{}{\text{safe}(\{\bar{r}\}, \varepsilon)} (\text{SAFE-RESOURCE}) \quad \frac{}{\text{safe}(\text{Unit}, \varepsilon)} (\text{SAFE-UNIT}) \\
\\
\frac{\varepsilon \subseteq \varepsilon' \quad \text{ho-safe}(\hat{\tau}_1, \varepsilon) \quad \text{safe}(\hat{\tau}_2, \varepsilon)}{\text{safe}(\hat{\tau}_1 \rightarrow_{\varepsilon'} \hat{\tau}_2, \varepsilon)} (\text{SAFE-ARROW})
\end{array}$$

$$\boxed{\text{ho-safe}(\hat{\tau}, \varepsilon)}$$

$$\begin{array}{c}
\frac{}{\text{ho-safe}(\{\bar{r}\}, \varepsilon)} (\text{HOSAFE-RESOURCE}) \quad \frac{}{\text{ho-safe}(\text{Unit}, \varepsilon)} (\text{HOSAFE-UNIT}) \\
\\
\frac{\text{safe}(\hat{\tau}_1, \varepsilon) \quad \text{ho-safe}(\hat{\tau}_2, \varepsilon)}{\text{ho-safe}(\hat{\tau}_1 \rightarrow_{\varepsilon'} \hat{\tau}_2, \varepsilon)} (\text{HOSAFE-ARROW})
\end{array}$$

$$\boxed{\hat{\tau} <: \hat{\tau}}$$

$$\frac{\varepsilon \subseteq \varepsilon' \quad \hat{\tau}_2 <: \hat{\tau}'_2 \quad \hat{\tau}'_1 <: \hat{\tau}_1}{\hat{\tau}_1 \rightarrow_{\varepsilon} \hat{\tau}_2 <: \hat{\tau}'_1 \rightarrow_{\varepsilon'} \hat{\tau}'_2} (\text{S-EFFECTS}) \quad \frac{r \in \bar{r}_2 \implies r \in \bar{r}_1}{\{\bar{r}_2\} <: \{\bar{r}_1\}} (\text{S-RESOURCESET})$$

## 4 Dynamic Rules

$$\boxed{\hat{e} \longrightarrow \hat{e} \mid \varepsilon}$$

$$\begin{array}{c}
\frac{\hat{e}_1 \longrightarrow \hat{e}'_1 \mid \varepsilon}{\hat{e}_1 \hat{e}_2 \longrightarrow \hat{e}'_1 \hat{e}_2 \mid \varepsilon} \text{ (E-APP1)} \quad \frac{\hat{e}_2 \longrightarrow \hat{e}'_2 \mid \varepsilon}{\hat{v}_1 \hat{e}_2 \longrightarrow \hat{v}_1 \hat{e}'_2 \mid \varepsilon} \text{ (E-APP2)} \quad \frac{}{(\lambda x : \hat{\tau}. \hat{e}) \hat{v}_2 \longrightarrow [\hat{v}_2/x] \hat{e} \mid \emptyset} \text{ (E-APP3)} \\
\\
\frac{\hat{e} \rightarrow \hat{e}' \mid \varepsilon}{\hat{e}. \pi \longrightarrow \hat{e}'. \pi \mid \varepsilon} \text{ (E-OPERCALL1)} \quad \frac{r \in R \quad \pi \in \Pi}{r. \pi \longrightarrow \mathbf{unit} \mid \{r. \pi\}} \text{ (E-OPERCALL2)} \\
\\
\frac{}{\forall t : \tau. e \longrightarrow [\tau/t] e} \text{ (E-TYPEPOLY)} \quad \frac{}{\forall \epsilon : \varepsilon. e \longrightarrow [\varepsilon/\epsilon] e} \text{ (E-FXPOLY)} \\
\\
\frac{\hat{e} \longrightarrow \hat{e}' \mid \varepsilon'}{\mathbf{import}(\varepsilon_s) \ x = \hat{e} \ \mathbf{in} \ e \longrightarrow \mathbf{import}(\varepsilon_s) \ x = \hat{e}' \ \mathbf{in} \ e \mid \varepsilon'} \text{ (E-IMPORT1)} \\
\\
\frac{}{\mathbf{import}(\varepsilon_s) \ x = \hat{e} \ \mathbf{in} \ e \longrightarrow [\hat{v}/x] \mathbf{annot}(e, \varepsilon_s) \mid \emptyset} \text{ (E-IMPORT2)}
\end{array}$$

## 5 Encodings

### 5.1 $\perp$

The bottom type is defined as  $\perp \stackrel{\text{def}}{=} \emptyset$ , which is the literal for an empty set of resources.

$$\frac{}{\Gamma \vdash \perp : \emptyset} \text{ (T-}\perp\text{)} \quad \frac{}{\Gamma \vdash \perp : \emptyset \ \mathbf{with} \ \emptyset} \text{ (}\varepsilon\text{-}\perp\text{)}$$

### 5.2 **unit**, **Unit**

Define  $\mathbf{unit} = \lambda x : \emptyset. x$ , i.e. the function which takes an empty set of resources and returns it. We shall refer to its type, which is  $\emptyset \rightarrow_{\emptyset} \emptyset$ , as **Unit**. It has various properties befitting **unit**.

1. **unit** cannot be invoked as  $\emptyset$  is uninhabited.
2. **unit** is a value.
3. The only term with type **Unit** is **unit**.
4.  $\vdash \mathbf{unit} : \mathbf{Unit}$  by using  $\varepsilon$ -ABS and  $\varepsilon$ -VAR.
5.  $\mathbf{effects}(\mathbf{Unit}) = \mathbf{ho-effects}(\mathbf{Unit}) = \emptyset$
6.  $\mathbf{safe}(\mathbf{Unit}, \varepsilon)$  and  $\mathbf{ho-safe}(\mathbf{Unit}, \varepsilon)$

$$\frac{}{\Gamma \vdash \mathbf{unit} : \mathbf{Unit}} \text{ (T-UNIT)} \quad \frac{}{\Gamma \vdash \mathbf{unit} : \mathbf{Unit} \ \mathbf{with} \ \emptyset} \text{ (}\varepsilon\text{-UNIT)}$$