

# Data Mining: Itemsets

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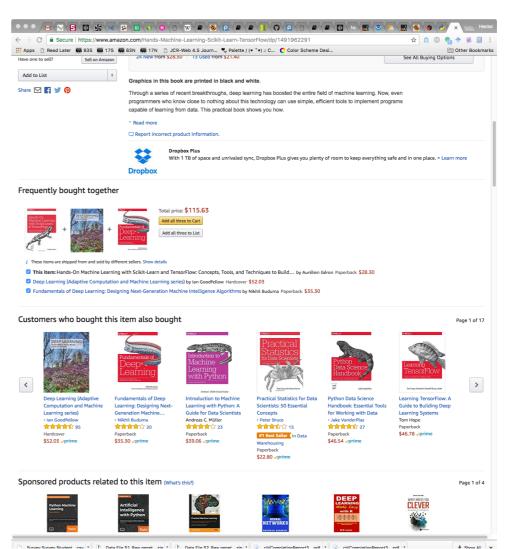
## Data Mining

For today: the analysis of itemsets

- Similar itemsets
- Frequent itemsets

#### Collaborative Filtering

- Items: customers
- Itemsets: customers that bought a specific book
- Similar itemsets: books purchased by same customers



#### Similar Documents

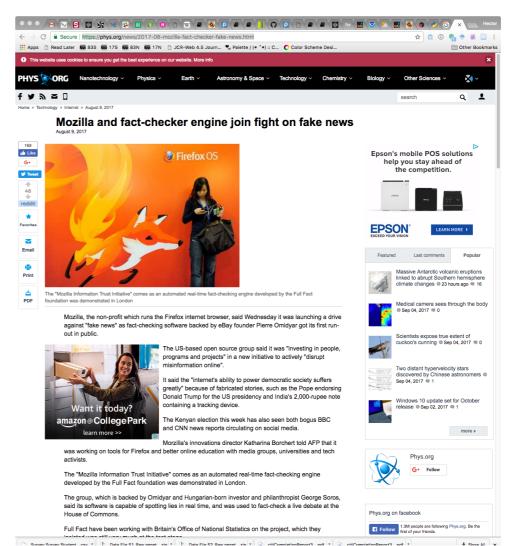
- Items: words
- Itemsets: documents
- Similar itemsets: documents using many of the same words





#### Similar News Articles

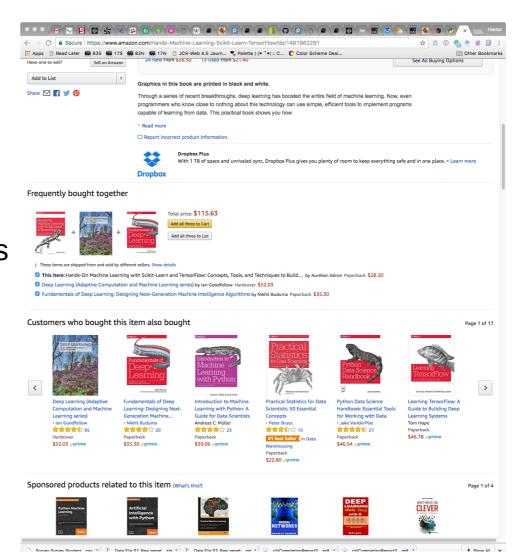
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### Frequent Itemsets

#### Online Purchasing

- Items: books
- Itemsets: orders (sets of books)
- Frequent itemsets: sets of books that are purchased together frequently

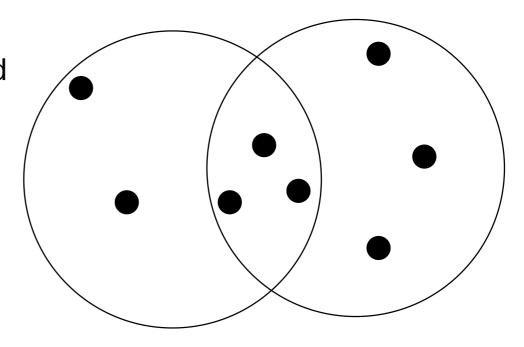


- Describing set similarity (Jaccard Similarity)
- Representing documents as sets (Shingling)
- Similarity-preserving set summaries (Minhash)
- Search for similar itemsets using Locality-Sensitive Hashing (LSH)

# **Jaccard Similarity**

The *Jaccard Similarity* of sets s and t is

$$\frac{S \bigcap T}{S \bigcup T}$$



#### Exercises

- Compute the Jaccard bag similarity of each pair of sets: {1,1,1,2}, {1,1,2,3}, {1,2,3,4}
- Suppose we have a universal set v of n elements. We chose two subsets s and r, each with m of the n elements. What is the expected value of the JS of s and r?

### Documents (Shingles)

• Set all words to lowercase, remove all whitespace and punctuation

```
"Hurricane Irma, they confirmed landfall" -> "hurricaneirmatheyconfirmedlandfall"
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• For some parameter *k*, represent document as the set of *k*-long subsequences of document

```
For k=3 {hur,urr,rri,...,eir,irm,rma,...,fir,irm,rme,...}
```

# Documents (Shingles)

- Choosing k: choose large enough that probability of any given shingle appearing in any given document is low. Depends on collection.
- Hashing: hash k-shingles instead of using them directly in algorithms that follow
- Using words, effective for similarity (more meaning) but sparser, set of possible shingles is huge

Clever idea: let's summarize item(sets) (reduce data size!) but make it easy to find similar item(sets).

#### **Characteristic Matrix**

Element	$S_1$	$S_2$	$S_3$	$S_4$
a	1	0	0	1
b	0	0	1	0
c	0	1	0	1
d	1	0	1	1
e	0	0	1	0

- Permute the rows of the characteristic matrix
- Min-Hash value of set: first non-zero row in corresponding column

#### Permuted characteristic Matrix

Element	$S_1$	$S_2$	$S_3$	$S_4$
b	0	0	1	0
e	0	0	1	0
a	1	0	0	1
d	1	0	1	1
<i>c</i>	0	1	0	1

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Property:  $Pr(h(S_i) = h(S_j)) = JS(S_i, S_j)$ 

Pf: on the board

### Min-Hash Signatures

- Choose n permutations of rows, and set  $h_i(S_j)$  as the Min-Hash given by permutation i of set j
- Represent set j by the *signature* vector of Min-hashes  $[h_1(S_j), \ldots, h_n(S_j)]$
- Collect signature vectors into a *signature matrix*

# Min-Hash Signatures in Practice

Instead of row permutations, use hash functions  $h_i$  over row indices

Let SIG(i,c) be the *i*th hash of *c*th element

Initialize: set  $SIG(i,c) = \infty$  for all i and c Row r:

- Compute  $h_i(r)$  for all i
- For each column *c*:
  - $\circ$  If c has a 0 in row r, do nothing
  - o If c has a 1 in row r, then for each  $i=1,\ldots,n$ : set SIG(i,c) to  $\min(SIG(i,c),h_i(r))$

#### Exercise

Element	$S_1$	$S_2$	$S_3$	$S_4$
0	0	1	0	1
1	0	1	0	0
2	1	0	0	1
3	0	0	1	0
4	0	0	1	1
5	1	0	0	0

 $h_1(x) = 2x + 1 \mod 6 \ h_2(x) = 3x + 2 \mod 6$ 

 $h_3(x) = 5x + 2 \mod 6$ 

# JS and Minhashing

Estimate  $JS(S_i, S_j)$  as the proportion of rows (hashes) of the signature matrix that are equal for columns  $S_i$  and  $S_j$ .

#### Exercise

Prove that if the JS of two sets is 0, then Min-Hash always gives the right answer.

# Locality-Sensitive Hashing

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But to find all pairs of similar item(sets) can still take a lot of time

LSH gives us a way of only comparing likely similar pairs.

Conversely, ignore pairs that are unlikely similar

#### LSH for Minhash

- Divide signature matrix into b ands, each with r rows
- For each column (itemset) and band, hash it's r entries according to some hash function
- Use same hash function in each of the bands, but use different hash arrays

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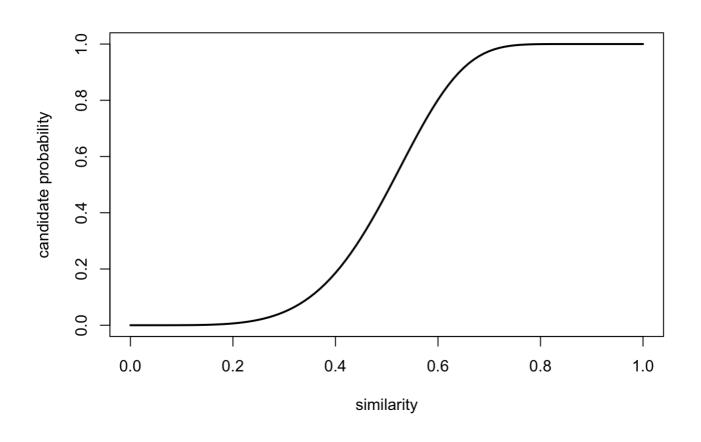
Itemsets without matching signatures will not

# Analysis of LSH

Let JS(S,T)=s

- Probability signatures agree in all rows of one band: s<sup>r</sup>
- Probability do not agree in at least one row of a band:  $1-s^r$
- Probability that signatures do not agree in all rows of any of the bands:  $(1-s^r)^b$
- Probability that signatures agree in all the rows of at least one band (hash to the same bucket at least once):  $1 (1 s^r)^b$ .

# Analysis of LSH



#### Part I: Shingles

- Pick a value of k, construct k-shingles for each document (optionally hashing k-shingles)
- Sort documents by document-shingle pairs by shingle

Part II: Minhash

- Pick a length *n* for minhash signatures
- Compute minhash signatures for all documents

Part III: LSH

- Choose threshold t for how similar documents have to be to consider as a similar pair
- Choose number of bands b and number of rows r such that br = n and threshold t is approximately  $(1/b)^{(1/r)}$
- Construct candidate pairs using LSH

Part IV: Confirm similar pairs

- For each candidate pair, confirm that their signatures match in at least
  t fraction of rows
- Optionally, verify similarity in shingled documents

### Frequent Itemsets

Find items that occur frequently together in sets

#### Examples:

- items frequently bought together in the same transaction
- words that appear frequently together in the same document

### Market-Basket Model

Items: objects we are modeling Baskets: sets of items (transactions)

Frequent itemsets: items that co-occur frequently in baskets

### Frequent Itemsets

Support: define the support of an itemset I as the number of baskets in which itemset I appears

Frequent itemsets: Itemsets I with support at least some support threshold S

## Example

- (1) {Cat, and, dog, bites}
- (2) {Yahoo, news, claims, a, cat, mated, with, a, dog, and, produced, viable, offspring}
- (3) {Cat, killer, likely, is, a, big, dog}
- (4) {Professional, free, advice, on, dog, training, puppy, training}
- (5) {Cat, and, kitten, training, and, behavior}
- (6) {Dog, &, Cat, provides, dog, training, in, Eugene, Oregon}
- (7) {"Dog, and, cat", is, a, slang, term, used, by, police, officers, for, a, male-female, relationship}
- (8) {Shop, for, your, show, dog, grooming, and, pet, supplies}

### **Association Rules**

Rules of the form  $I \rightarrow j$ : if itemset I is in basket, then item I is likely in basket as well

*rule confidence*: ratio of support of  $I \cup \{j\}$  to support of I.

*rule interest*: difference between confidence of rule and fraction of baskets that contain *j* 

### **Association Rules**

Note: once we have itemsets, we can get association rules easily

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Suppose we find all frequent itemsets over some support threshold

Let itemset *J* with *n* items be one of those itemsets, then

- 1. there are only n candidate Association Rules  $J \{j\} \rightarrow j$
- 2. Both  $J \{j\}$  and j are also frequent itemsets, so we have already calculated their support
- 3. We can quickly compute the *confidence* and *interest* of each rule

Suppose there are 100 items, numbered 1 to 100, and also 100 baskets, numbered 1 to 100.

Item i is in basket b iff i divides b with no remainder

- Item 1 is in all baskets, item 2 in the even-numbered baskets
- Basket 24 contains items {1,2,3,4,6,8,12,24}
- a) If support threshold is 5, which items are frequent? b) Which pairs are frequent?

## Itemset monotonicity

If I is a frequent itemset, then every subset of I is a frequent itemset

Why?

Suppose we are given baskets over *n* items

First pass

Count the number of occurrences of each item (array of n values)

After first pass

Identify frequent singletons (above support threshold)

Second pass

Count the number of occurrences of pairs of frequent items

- For each basket:
  - Check which of its items are frequent (first pass)
  - For each pair of items increase occurrence count

After the second pass

Identify frequent pairs (above support threshold)

Third pass

Count the number of occurrences of frequent pairs + a frequent item

- For each basket:
  - Check which item pairs and singletons that are frequent (first and second pass)
  - For each combination of pair and singleton, increase occurrence count

After the third pass

And so on until no more frequent sets are identified

#### Notes:

- The data structure to store pair counts will be important consideration
- The algorithm has a construct-filter structure: at each pass, *construct* the set of candidate itemsets, *filter* to those that are frequent

Apply A-priori algorithm to previous exercise

# Handling large datasets

For large datasets storing occurrences of candidate frequent pairs is problematic

PCY algorithm: hash item pairs and keep count in hash bucket

Define candidate frequent pairs as

- *i* and *j* anre frequent items
- $\{i,j\}$  hashes to a frequent bucket (with count > threshold)

# Handling large datasets

Identify frequent buckets with a bitmap (little memory)

Only count (and verify) candidate pairs as defined above (expected to be much fewer)

Consider baskets over items 1,...,6

```
\{1,2,3\} \{2,3,4\} \{3,4,5\} \{4,5,6\} \{1,3,5\} \{2,4,6\} \{1,3,4\} \{2,4,5\} \{3,5,6\} \{1,2,4\} \{2,3,5\} \{3,4,6\}
```

- Compute support for each item and each pair of items
- Using hash function  $i \times j \mod 11$  (hash table with 11 buckets), which pairs hash to the same buckets?

- Which buckets are frequent?
- Which pairs are counted in the second pass of PCY algorithm?

## Summary

Itemset analysis: applications to collaborative filtering, recommendation engines

### Finding Similar Itemsets

- Jaccard similarity: measure of set similarity based on common items
- Minhashing with LSH: effective way of finding similar itemsets with efficient data structures for large datasets

## Summary

### Finding Frequent Itemsets

- Market-basket data: model of item transactions
- Frequent Itemsets: Sets of items appearing frequently in "baskets"
- Association Rules:  $I \rightarrow j$
- Pair-counting Bottleneck: frequent itemset mining memory space taken mostly in keeping counts of pairs of frequent items
- Monotonicity of frequent itemsets
- A-priori Algorithm
- Hashing for large datasets