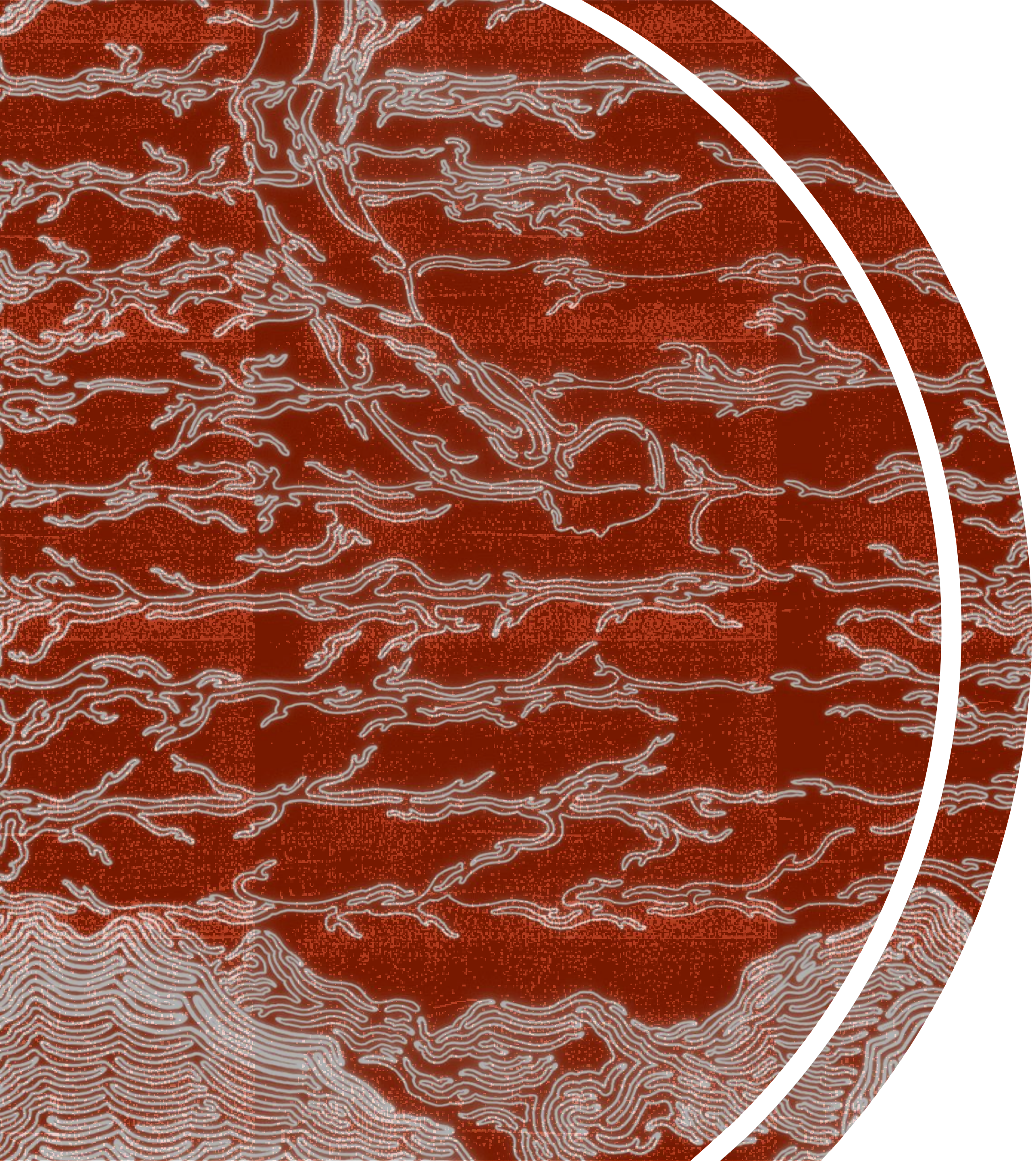

Introduction to Computational Topology

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Lecture 16, November 4, 2021

Discrete Morse Theory

TODAY'S GOAL

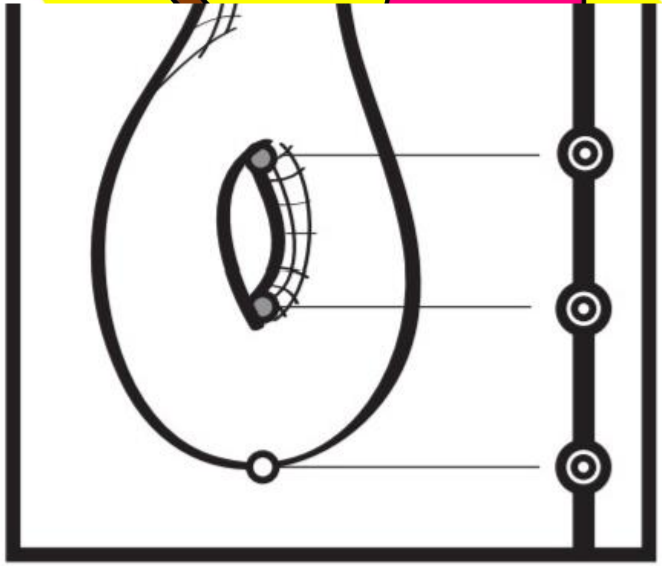
- Introduce a discrete version of the Morse theory that works for complexes

Definitions

- Height function h: $M \rightarrow R$

- Sub-level set $M_{\leq a}$: $h^{-1}(-\infty, a] = \{x : h(x) \leq a\}$

- Critical points: where the topology changes



Definitions

- Intuition: Morse function h is not important, only gradient field ∇h

- Discrete gradient $f: K \to R$
  - For k-simple α and (k+1)-simplex β: $f(\alpha) \geq f(\beta)$

- Discrete Morse function
  - All discrete gradients are unique

Definitions

- Critical cell
  - Cell with no discrete gradient

- Sub-level set $K_{\leq C}$

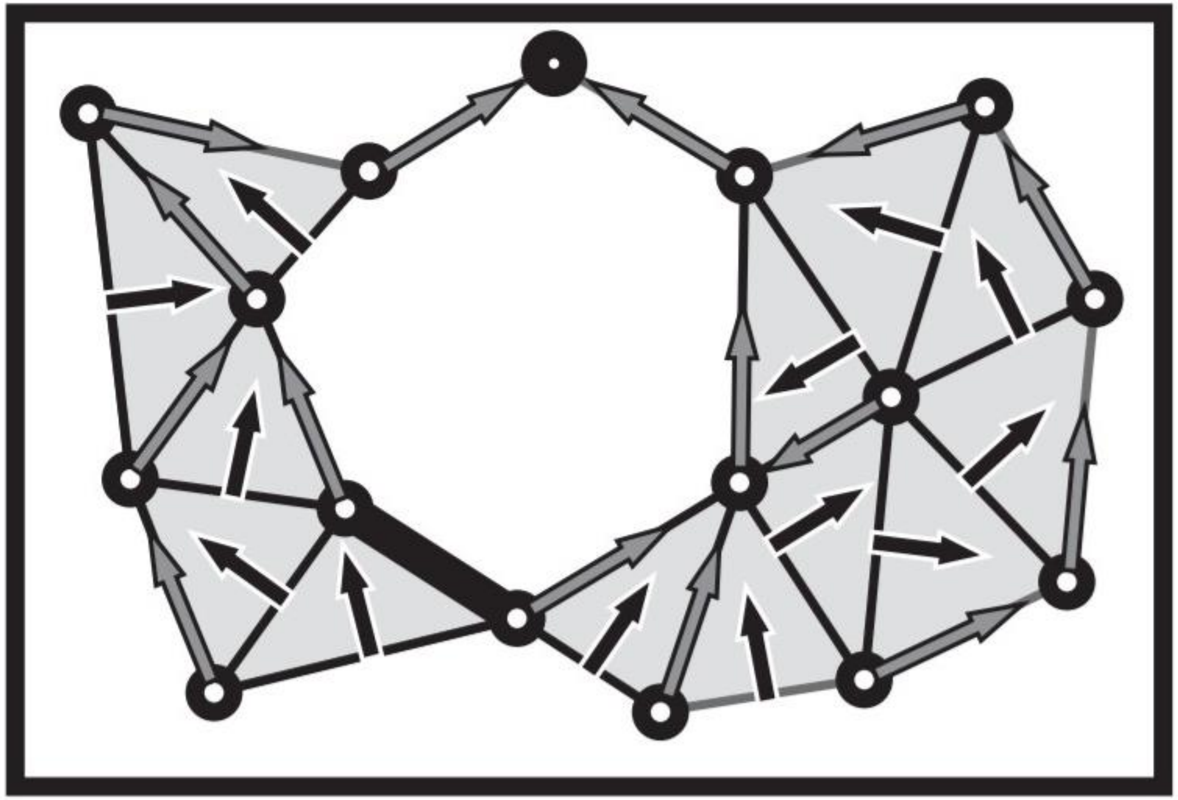
$$\{\beta : \beta \text{ in those } \alpha \text{ that } f(\alpha) \leq C\}$$

Example

- Which one is Morse?

Discrete Flowlines

- Pairing of neighboring k- and (k+1)-simplex are canceled

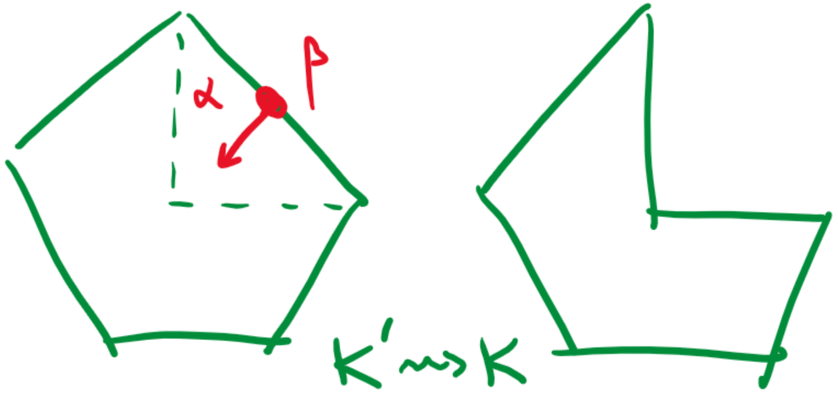


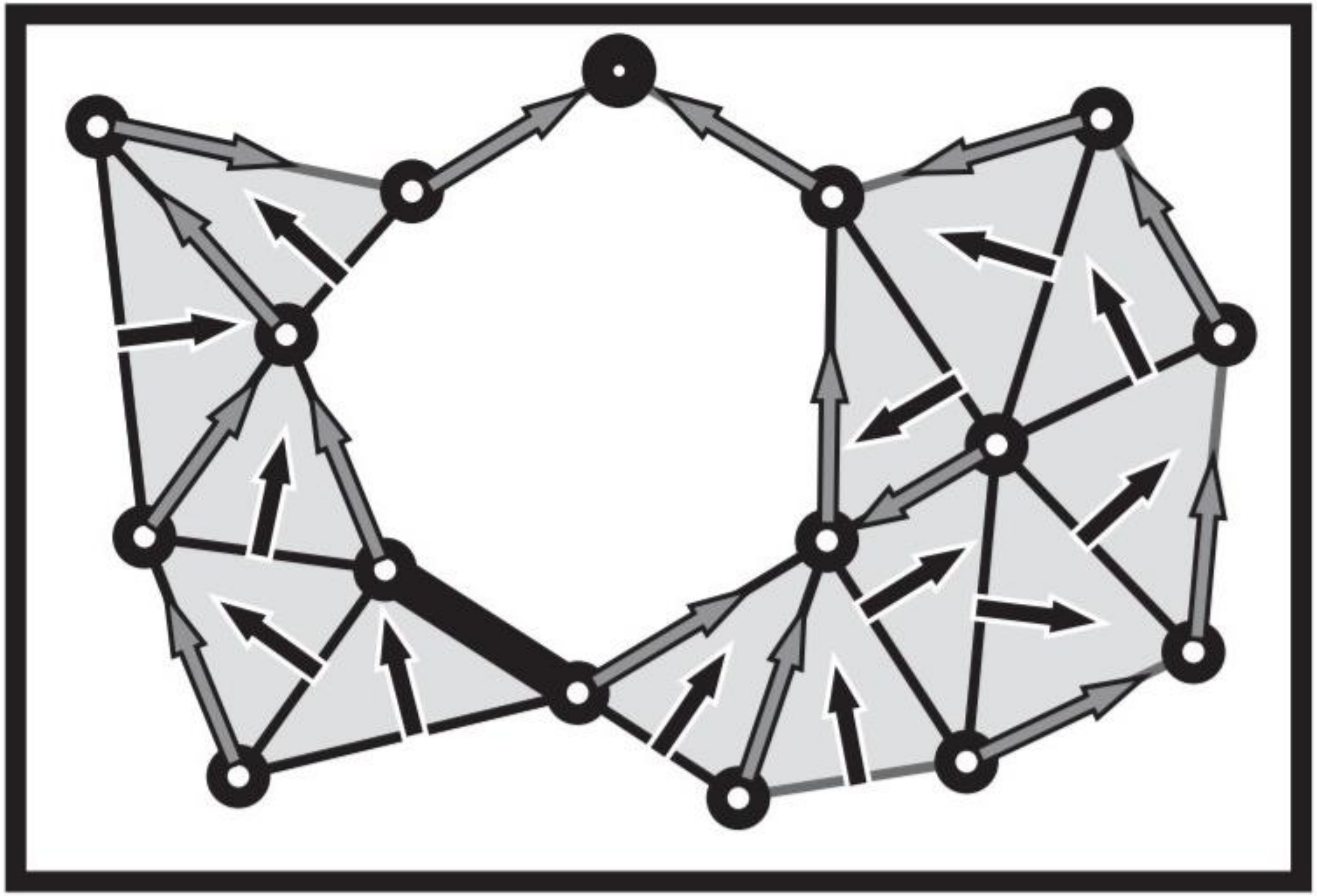
Proposition. A vector field is the gradient field of a discrete Morse function iff it is acyclic.

Properties

- $K_{\leq b} \simeq K_{\leq a}$ if no critical points in $(a, b]$
- $K_{\leq b} \simeq K_{\leq a} \cup \{\text{k-handle}\}$ if $(a, b]$ has k-dim critical point p

- Collapse (discrete homotopy)
  - If $K' = K \cup \{\alpha, \beta\}$ where β is the face to only α, then K' can be collapsed to K

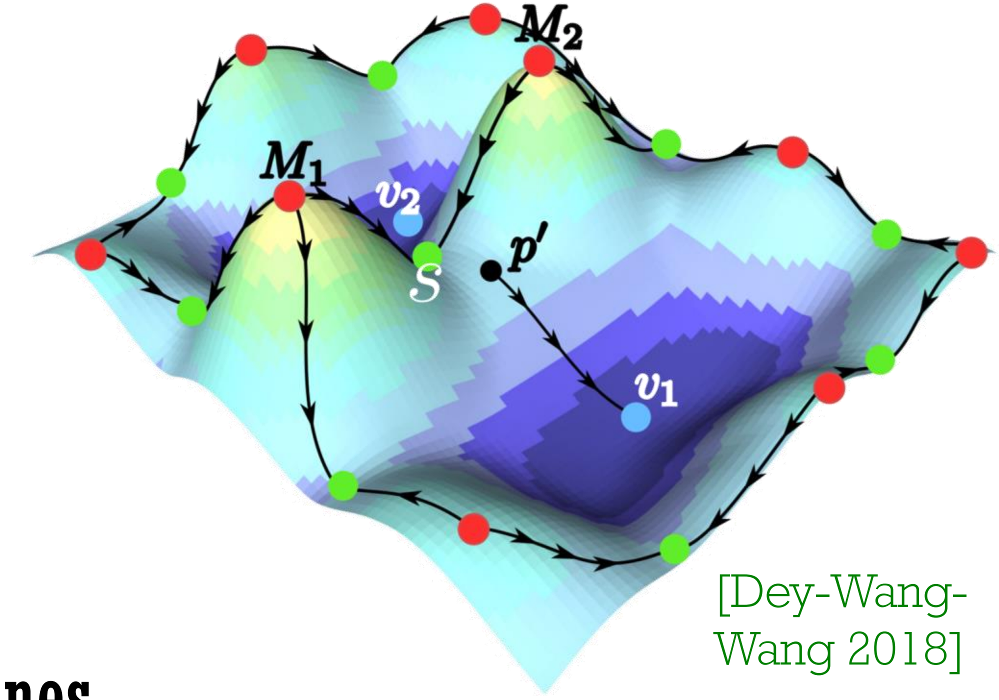




Example

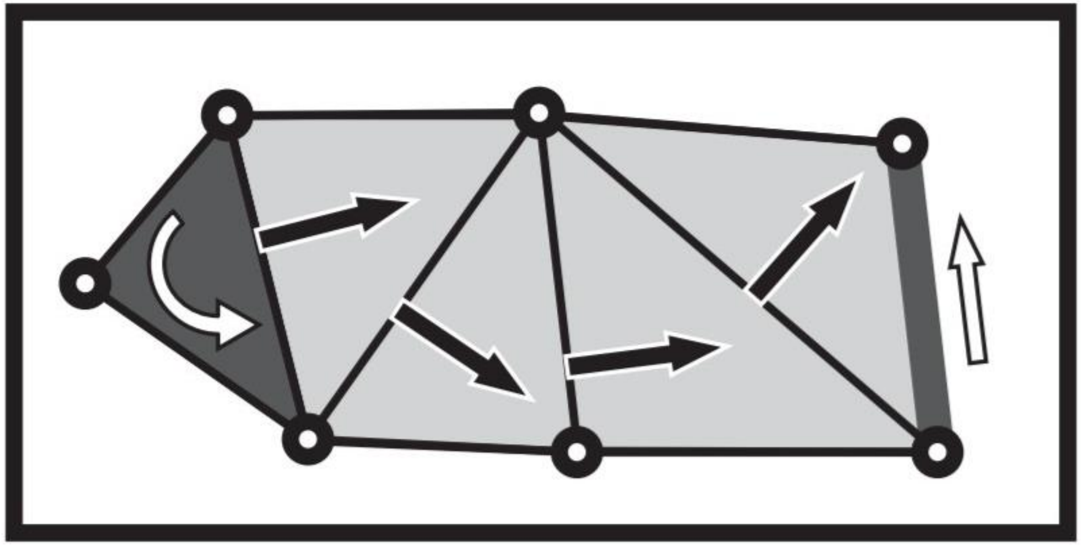
- Collapsing the flowlines

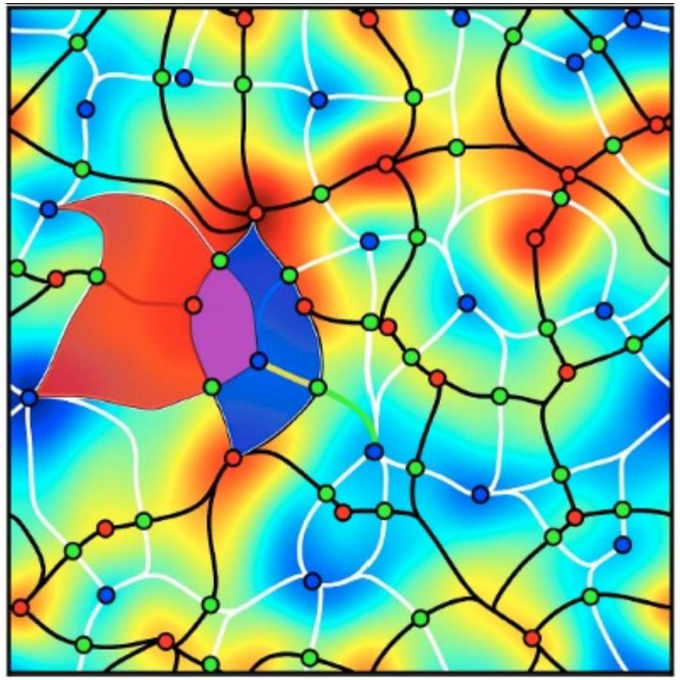
DISCRETE MORSE COMPLEX



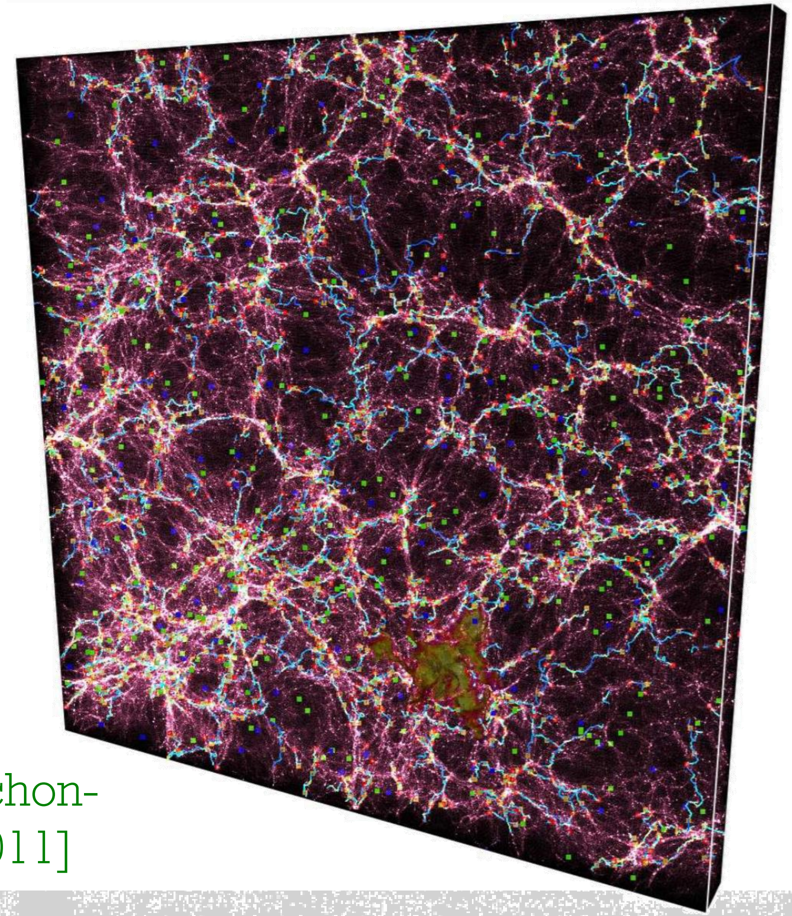
[Dey-Wang-Wang 2018]

- MC_k: ⟨k-dim critical cells⟩
- Boundary map ∂_k:
  all (k-1)-dim critical cells reachable by flowlines

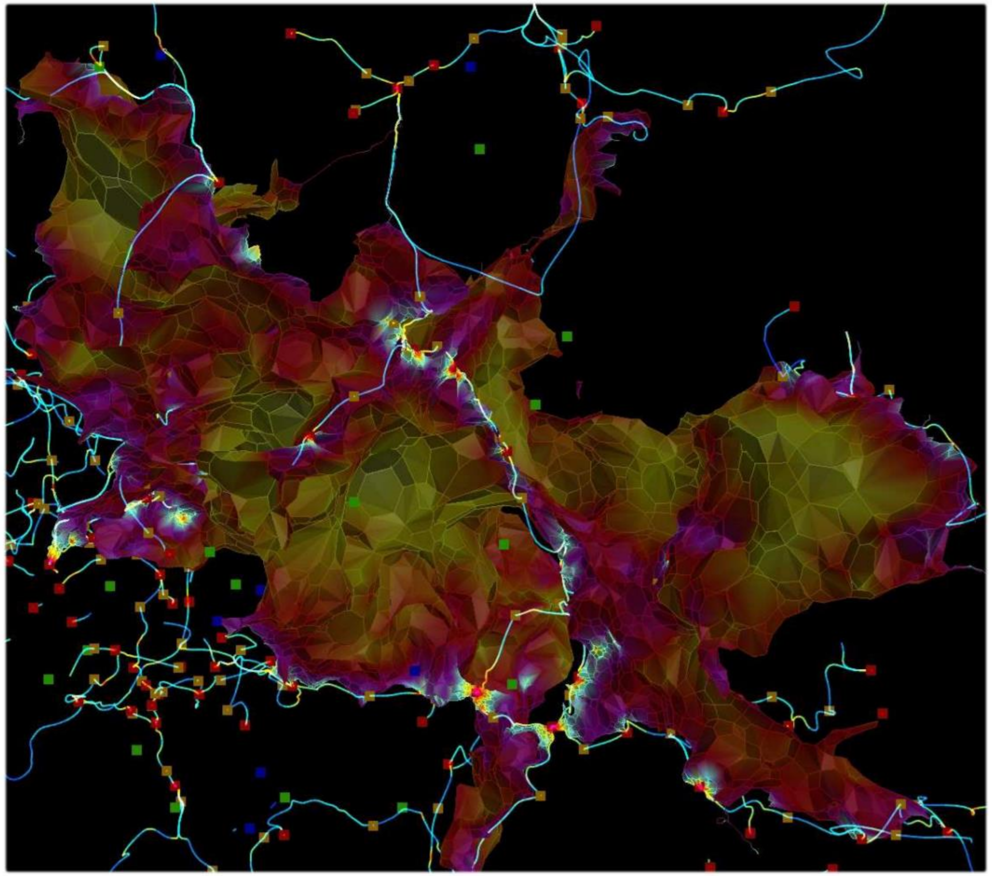




[Sousbie 2011]



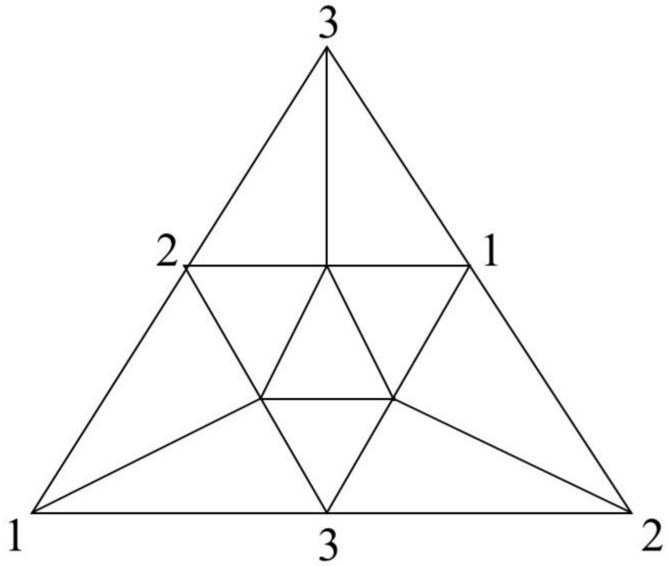
[Sousbie-Pichon-Kawahara 2011]

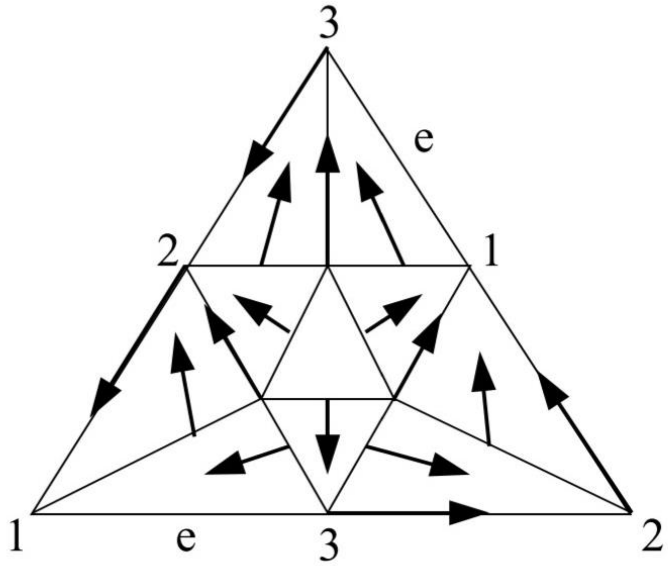


Morse Homology Theorem [Forman 1998]

$MH_n(K) \cong H_n(K)$ for any complex K

Exercise





Intermission

Food for Thought.

Forget about homology.

We can use it to simplify complexex!

Evasiveness (Why Lowerbound is Hard)

Motivating Puzzle

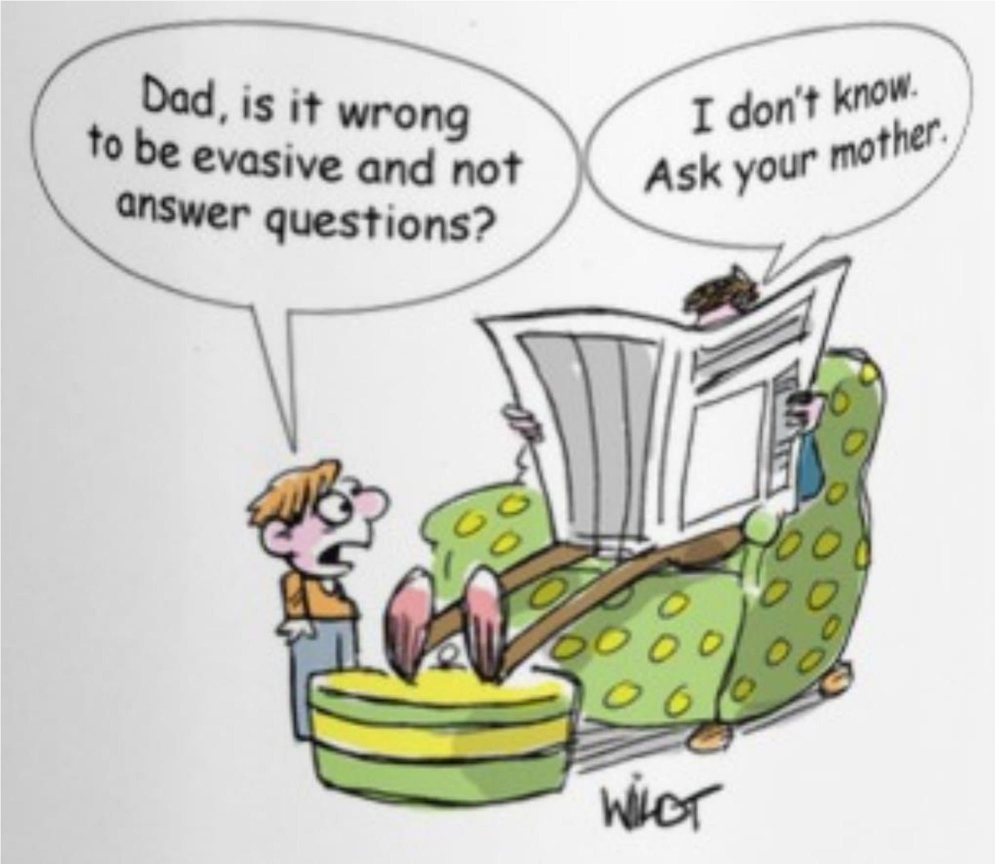
- Question allowed:
  "Is edge (i,j) in G?"
- Goal:
  Does G have a cycle?

Formulation

- Let $g(x_1, \ldots, x_E)$ be Boolean function
- Property T
  - g(X) = 0 iff graph X has property T
- Monotone property
  - If graph X has T, subgraph Y of X must be in T

- Determine if graph G satisfies T

How many Qs do you need?



Evasiveness Conjecture

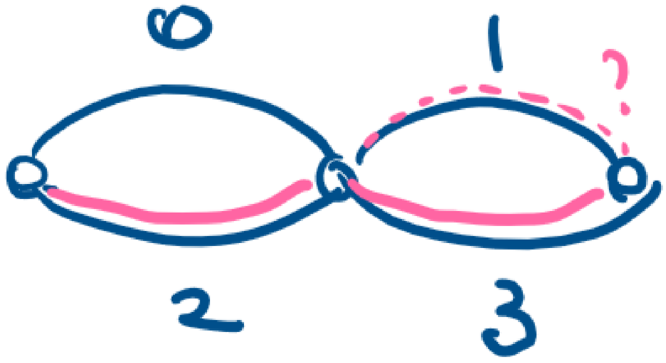
[Aanderaa-Rosenberg 1973]

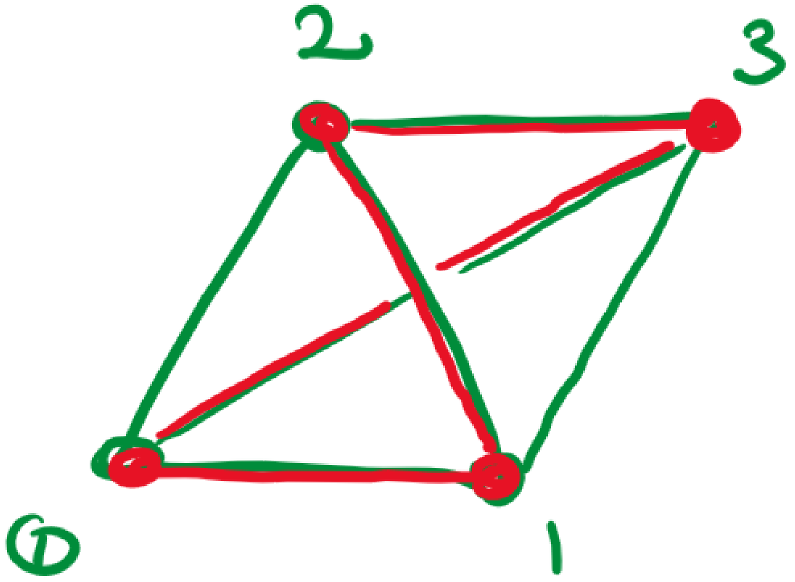
If property T is monotone, nontrivial, and symmetric, then T is evasive, i.e. requires $\binom{n}{2}$ questions

Topological Approach

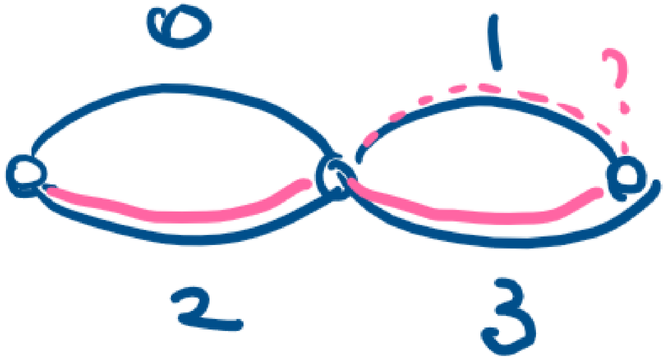
[Kahn-Saks-Sturtevant 1984]

- Construct complex K_P
  - Add cell σ if σ satisfies P

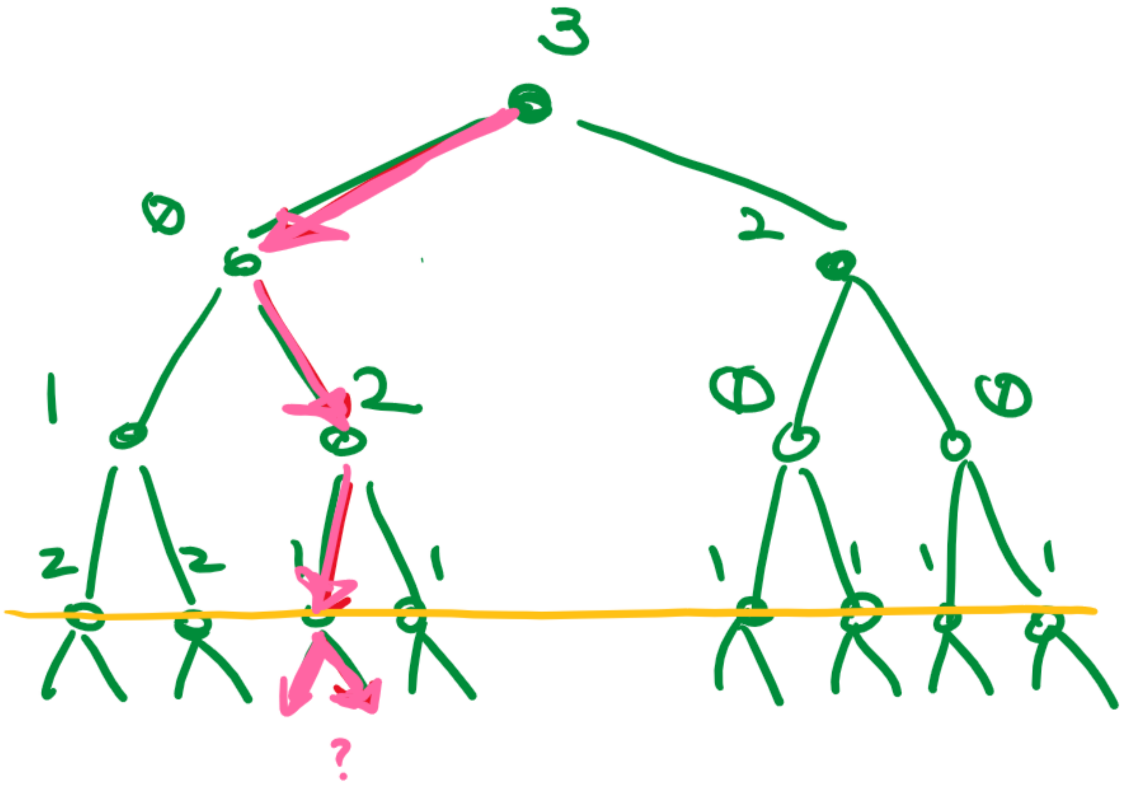


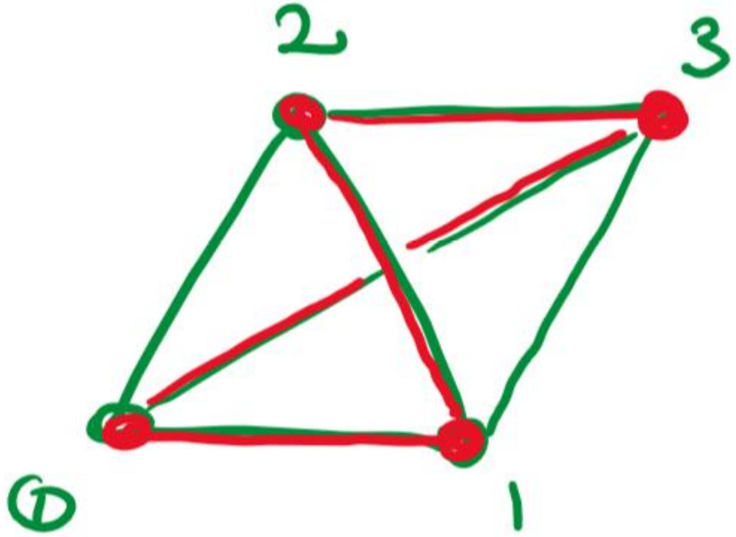


Observation

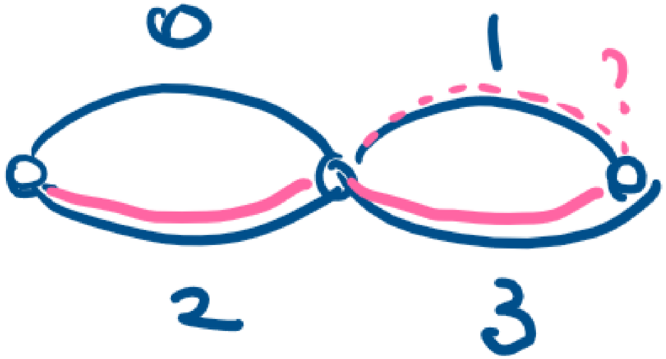


- Guessing algorithm induces discrete gradient field

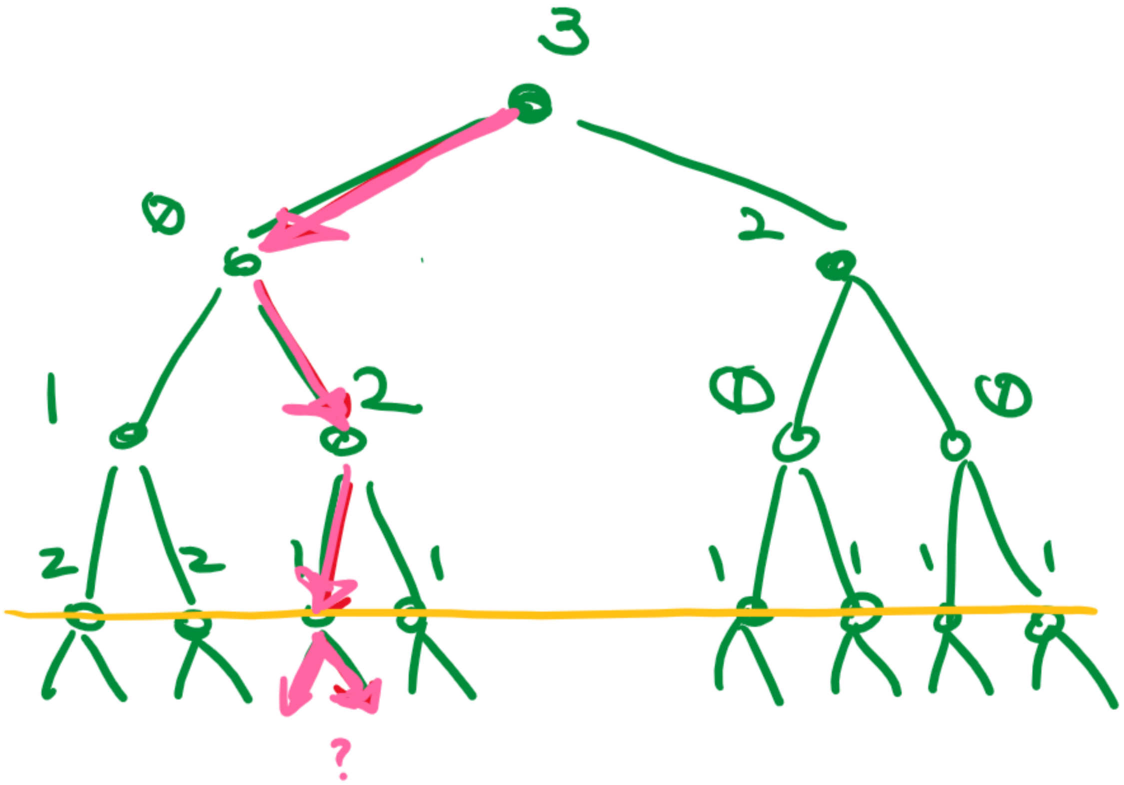


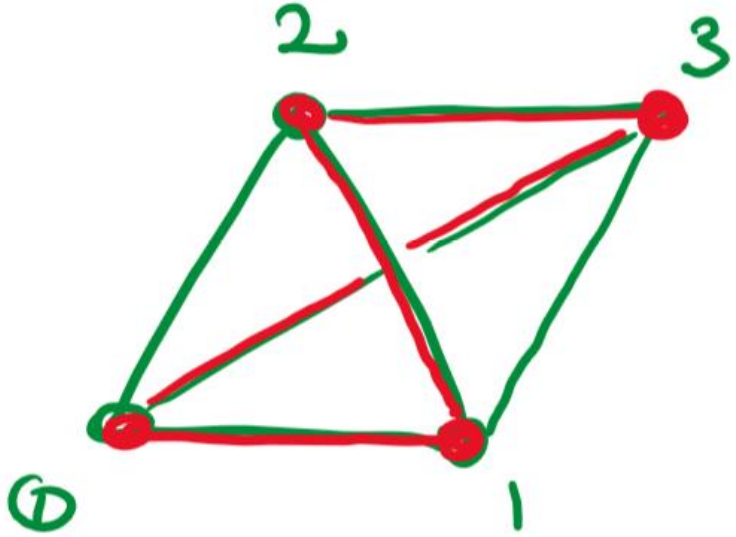


Observation



- Critical cells in K_p correspond to graph pairs that require the last question





Counting Evaders [Forman 2000]

#Evaders under any algorithm is at least

$2\ \Sigma_k \dim H_k(K_P)$

Proving Lowerbound by Showing K_p Non-Trivial

Next Time.

Almost end of the term. We'll see!