



# **DISCRETE MATHEMATICS IN COMPUTER SCIENCE**

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JANUARY 28, 2022**

# ADMINISTRIVIA

- Midterm 1
  - Jan 31 (Mon) 6–9PM
  - Moore Hall B13 Filene Auditorium
- Conflict Midterm 1
  - Feb 1 (Tue) 5–8PM
  - Moore Hall B13 Filene Auditorium
- SAS/Conflict Conflict/COVID
  - Come talk to me
- Closed-book written exam
- Scope: Module P on proofs
- One-page two-sided cheatsheet
  - Must be hand-written

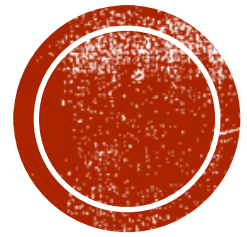


# ADMINISTRIVIA

- Homework 3 due today
- Homework 4 out today

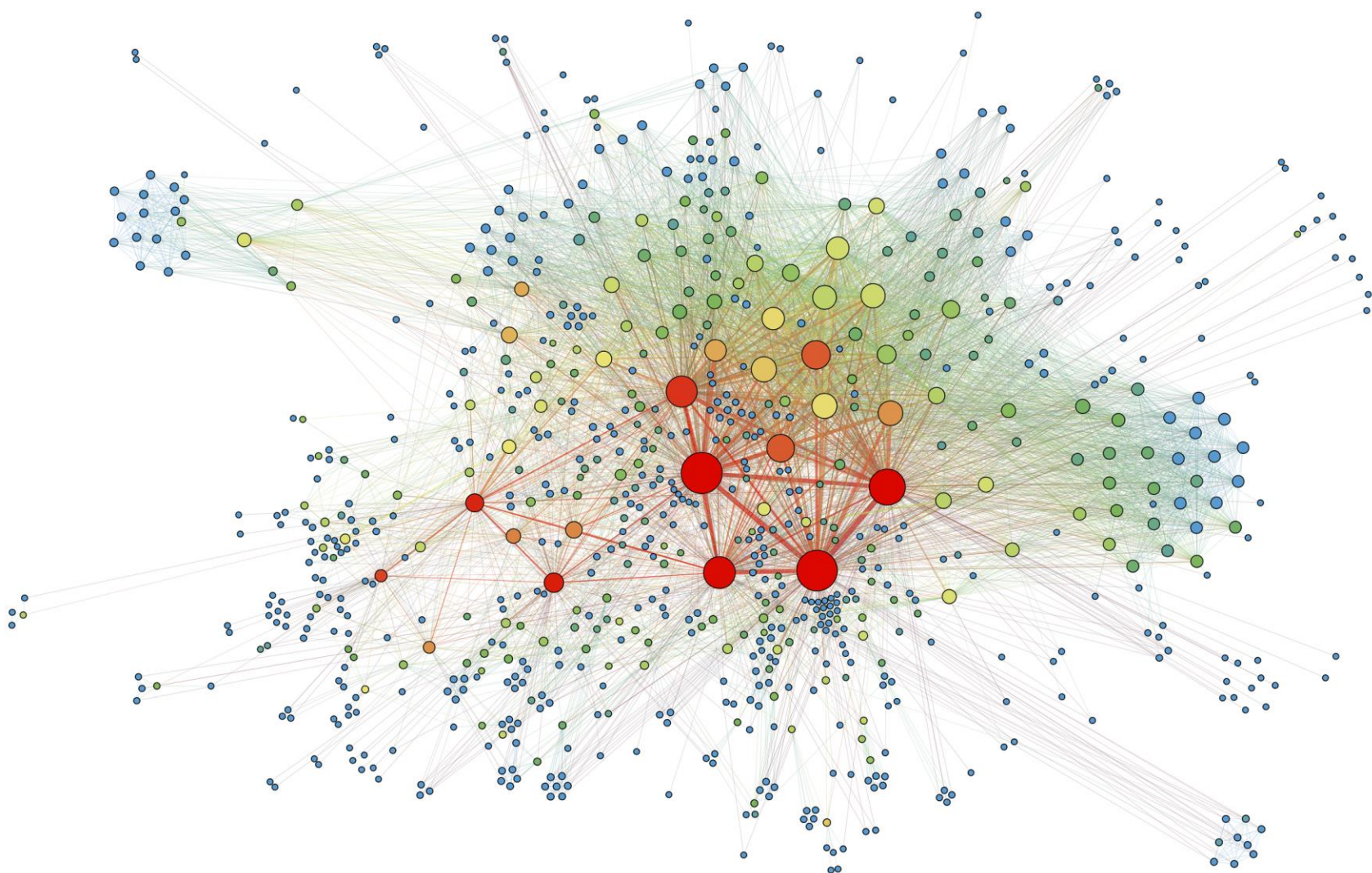






# GRAPH TRAVERSAL





# Jargon

walk

path

closed walk

cycle

length

distance

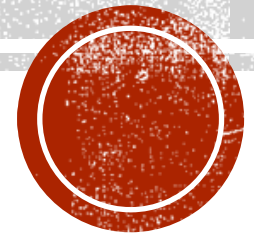
reachable

connected

components



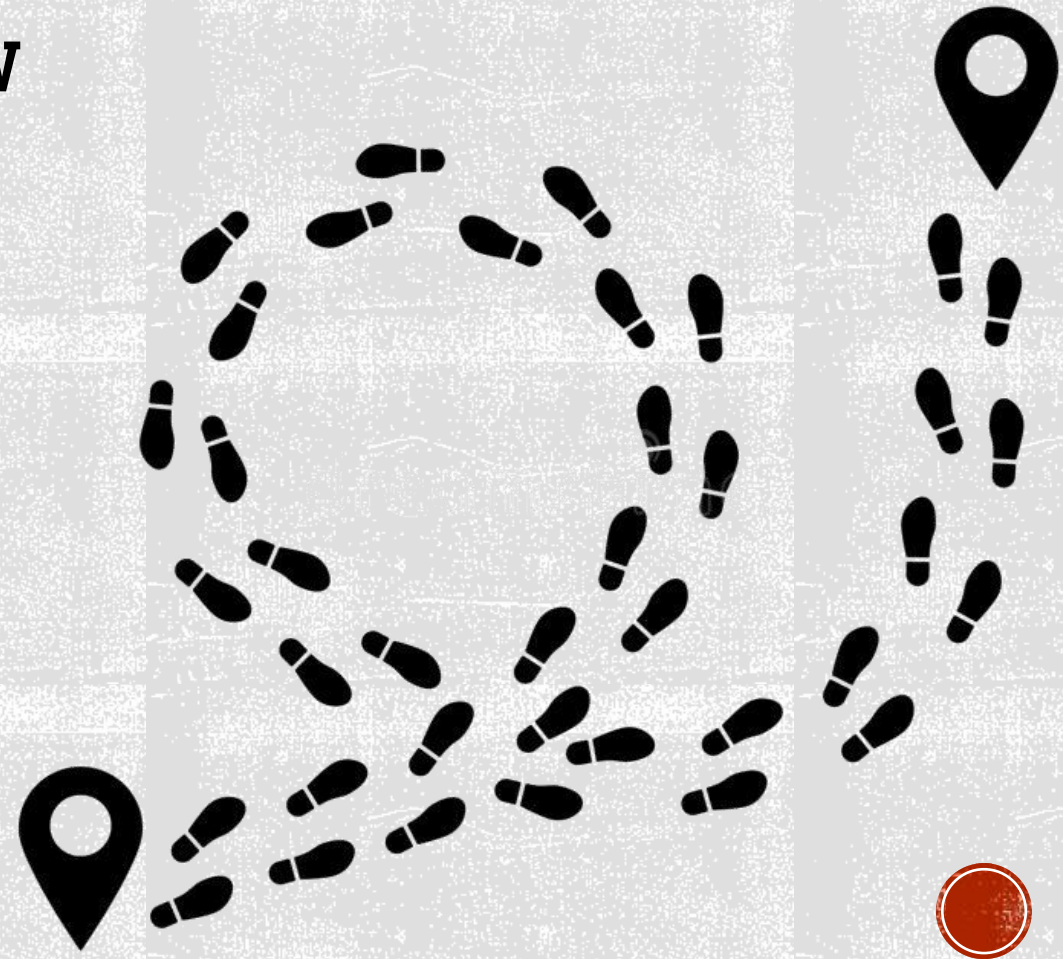
**IS THERE A PATH FROM  $u$  TO  $v$ ?**





# WANDERING LEMMA

- For any graph  $G$  and two vertices  $u$  and  $v$ 
  - if there is a walk from  $u$  to  $v$ , then there is a path from  $u$  to  $v$ .



**ANY SHORTEST WALK BETWEEN TWO VERTICES MUST  
BE A PATH.**

## **COROLLARY**

**ANY SHORTEST PATH IN AN  $n$ -VERTEX GRAPH MUST  
HAVE LENGTH AT MOST  $n-1$ .**





## WHATEVERFIRSTSEARCH(s, G):

```
put s in an empty bag
while the bag is not empty:
    take v from the bag
    if v is unmarked:
        mark v
        for each edge  $v \rightarrow w$  in G:
            put w in the bag
```

**THEOREM.** **WHATEVERFIRSTSEARCH(s, G)** MARKS ALL  
VERTICES REACHABLE FROM s IN G.

## WHATEVER-FIRST SEARCH

Implement the bag

- stack: **DFS**
- queue: **BFS**
- priority queue:
  - **MST** (weight)
  - **Dijkstra** (distance)



**WHATEVERFIRSTSEARCH(s, G):**

put  $\epsilon \rightarrow s$  in an empty bag

while the bag is not empty:

take  $p \rightarrow v$  from the bag

if  $v$  is unmarked:

mark  $v$

$\text{parent}(v) \leftarrow p$

for each edge  $v \rightarrow w$  in  $G$ :

put  $v \rightarrow w$  in the bag

## WHATEVER-FIRST SEARCH

**THEOREM. WHATEVERFIRSTSEARCH(s, G) MARKS  
EXACTLY THE VERTICES REACHABLE FROM  $s$  IN  $G$ .**



## BREADTHFIRSTSEARCH(s, G):

```
put  $\epsilon \rightarrow s$  in an empty queue
while the queue is not empty:
    take  $p \rightarrow v$  from the queue
    if  $v$  is unmarked:
        mark  $v$ 
         $\text{parent}(v) \leftarrow p$ 
        for each edge  $v \rightarrow w$  in  $G$ :
            put  $v \rightarrow w$  in the queue
```

**THEOREM.** BREADTHFIRSTSEARCH(s, G) COMPUTES  
THE SHORTEST PATH FROM  $s$  TO ANY VERTEX  $v$  IN  $G$ .

## BREADTH-FIRST SEARCH





**CAN THE SET OF PARENT EDGES FORM A CYCLE?**

**NEXT TIME.  
FORESTS AND TREES**

