

Administrivia

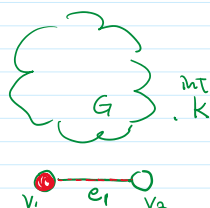
- WS 5/13 merged w/ HW6 (to come), WS 5/15 reductions.
- Quiz next Mon on reductions.
- Guest lectures next week! (No formal work sessions)



SUBSET SUM

input: set X of ^{positive} integers, target $T > 0$
output: Is there subset $Y \subseteq X$ s.t. $\sum_{y \in Y} y = T$?

$\text{MIN-VERTEX COVER.} \leq \text{SUBSET SUM}$
 $(G, k) \quad (X, T)$



$$A[i, j] = 1$$

	#vertices chosen	e_1, \dots, e_m
n_1	1	1 ...) v_1
n_2	1	1 ...) v_2
T	$(k)_2$	1 1 1 ... 1
	2^{m+1}	$2^m 2^{m-1} \dots$ picked 2^j choice = number.

Vertex $v_i \mapsto$ binary number.

$n_i[j] = 1$ if edge e_j incident to v_i

$n_i[m+1] = 1$ always.

$$T = k \cdot 2^{m+1} + \sum_{j=1}^m 2^j$$

X3DM

EXACT 3D MATCHING

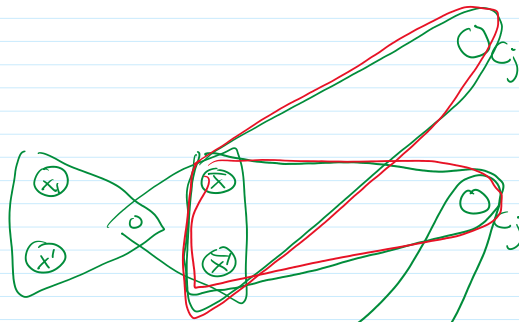
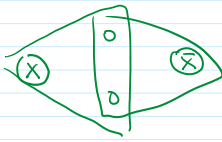
input: set S , collection \mathcal{U} of triples in S .

output: Is there disjoint subcollection of triples covering S ?

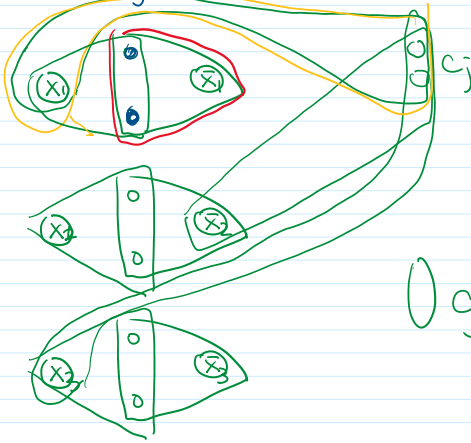
$$3SAT \equiv X3DM$$

$$3\text{-CNF } \Phi \quad (S, U)$$

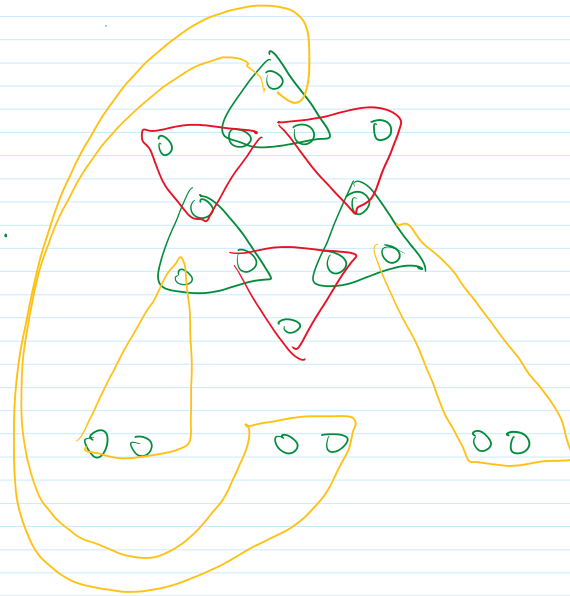
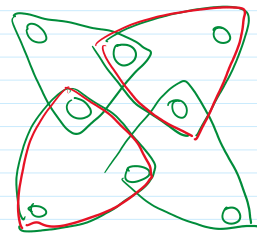
• variable



• clause $C_j = (x_1 \vee \bar{x}_2 \vee x_3)$



O_{C_j}



MAX 2SAT

input: CNF Φ , parameter k
output: Is #satisfied clauses $\geq k$?