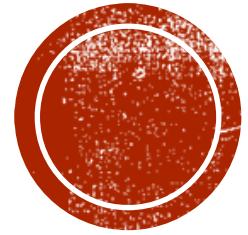




INTRODUCTION TO COMPUTATIONAL TOPOLOGY

HSIEN-CHIH CHANG
SEPTEMBER 14, 2021

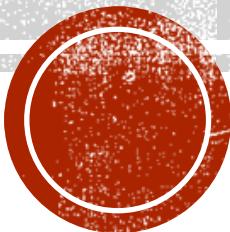


INTRODUCE YOURSELF!

WHO ARE YOU?

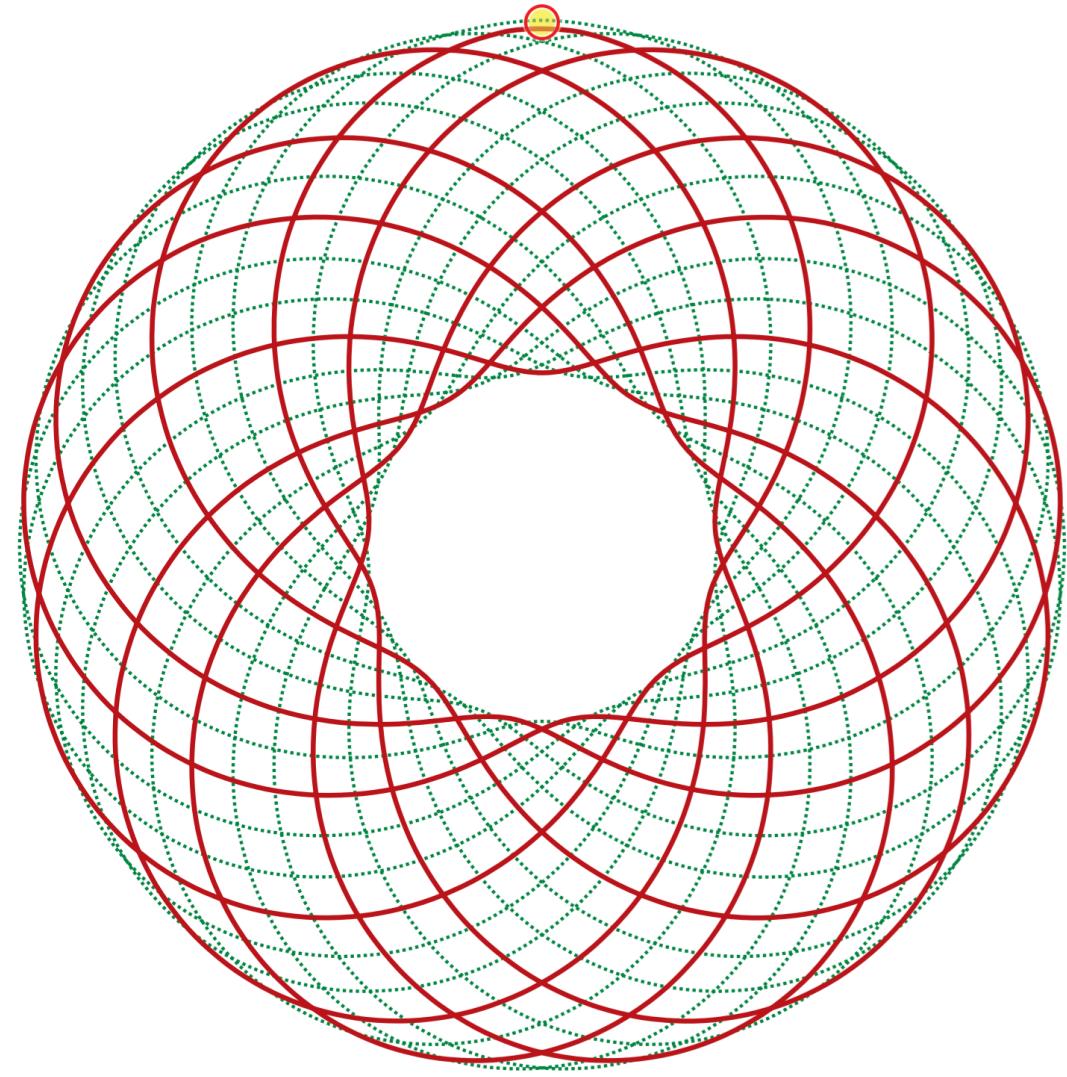
WHY DO YOU CARE ABOUT TOPOLOGY?

SHAPE OF THE SPACE How DO WE COMPUTE IT?



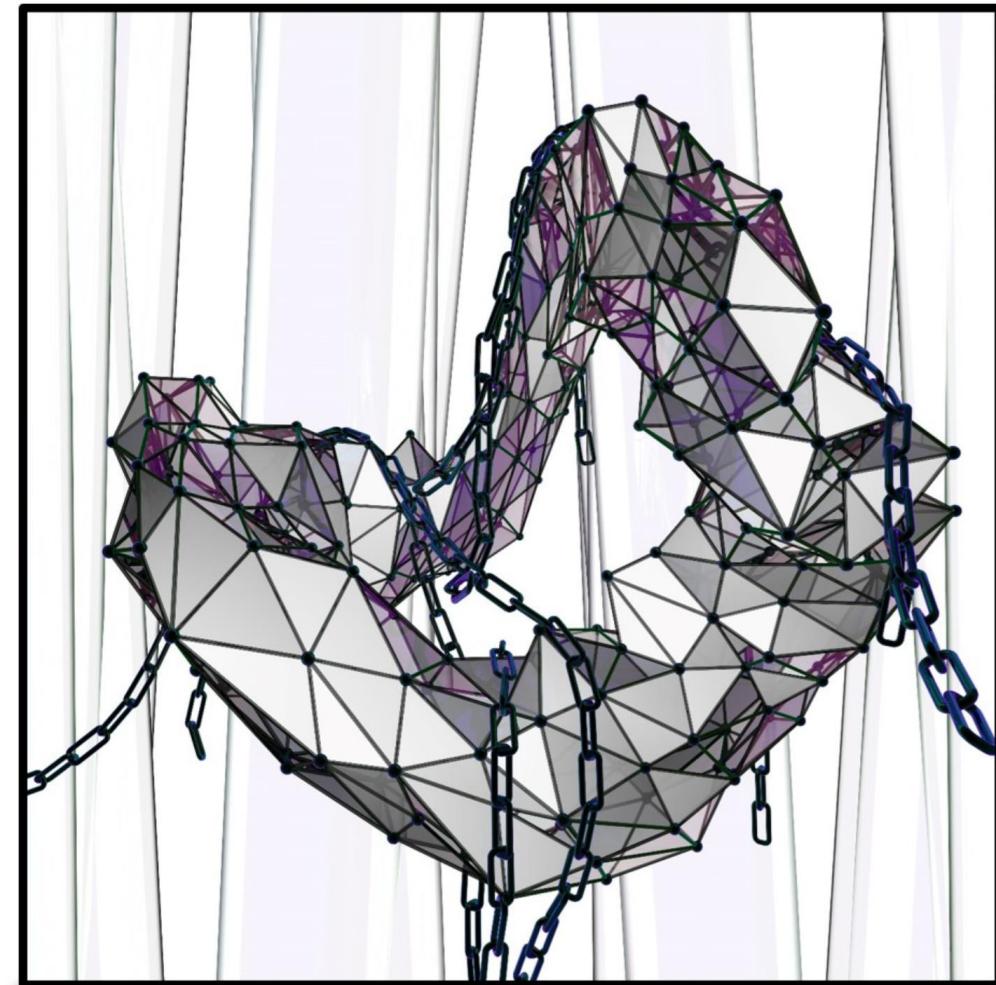
COURSE OUTLINE

- Theory
 - Curves



COURSE OUTLINE

- Theory
 - Curves
 - Surfaces/Complexes



COURSE OUTLINE

- Theory

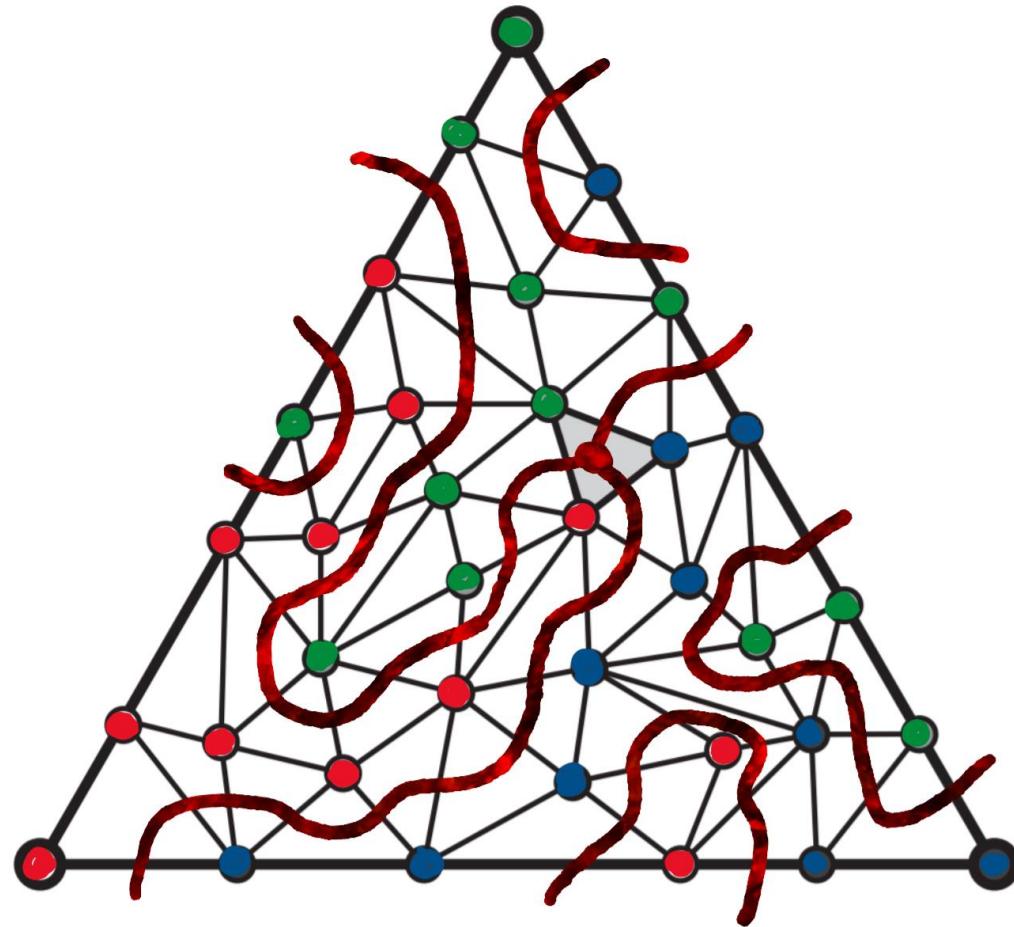
- Curves
- Surfaces/Complexes
- Homotopy



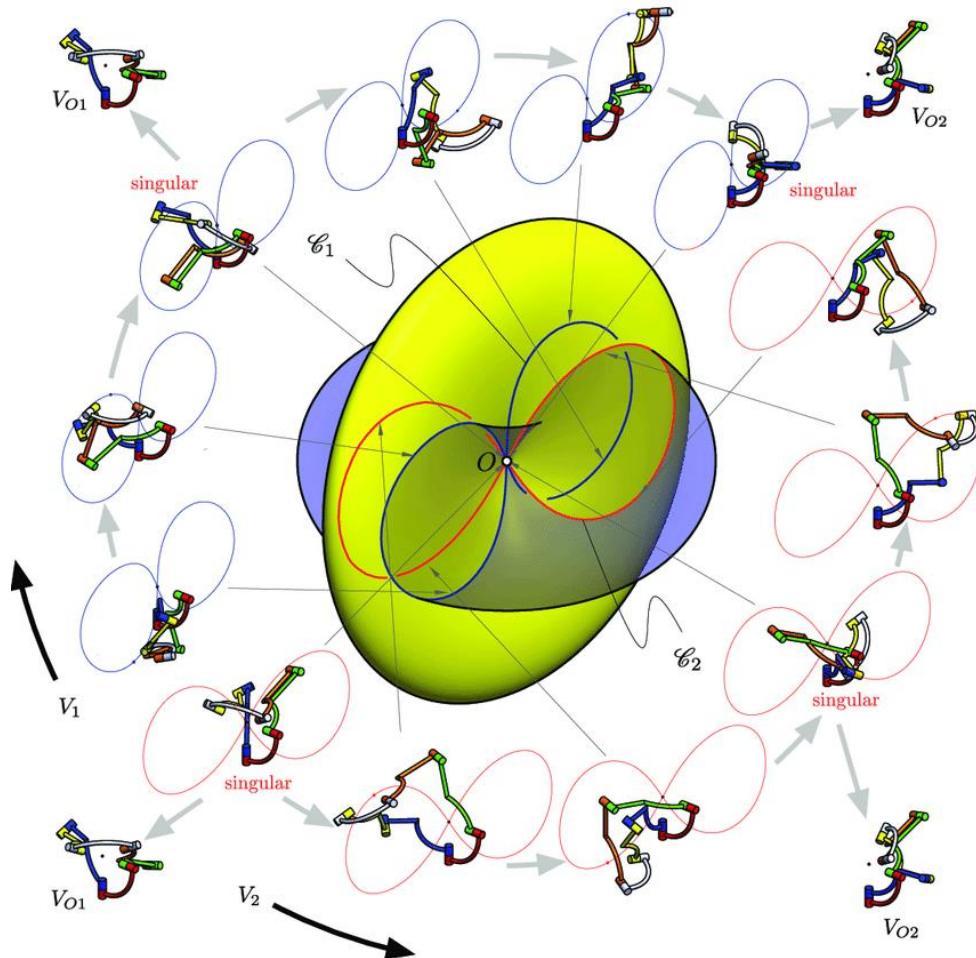
COURSE OUTLINE

- Theory

- Curves
- Surfaces/Complexes
- Homotopy
- Homology



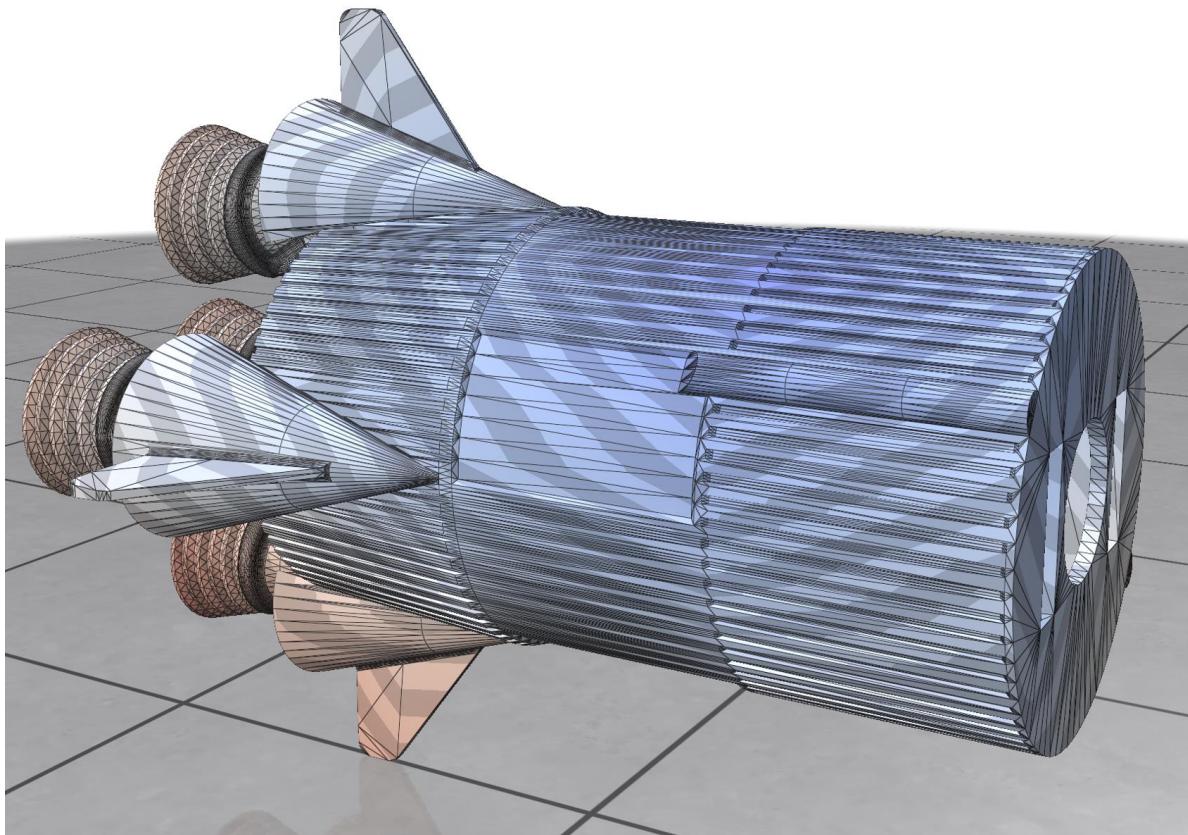
COURSE OUTLINE



- Applications
 - Linkage and Folding



COURSE OUTLINE

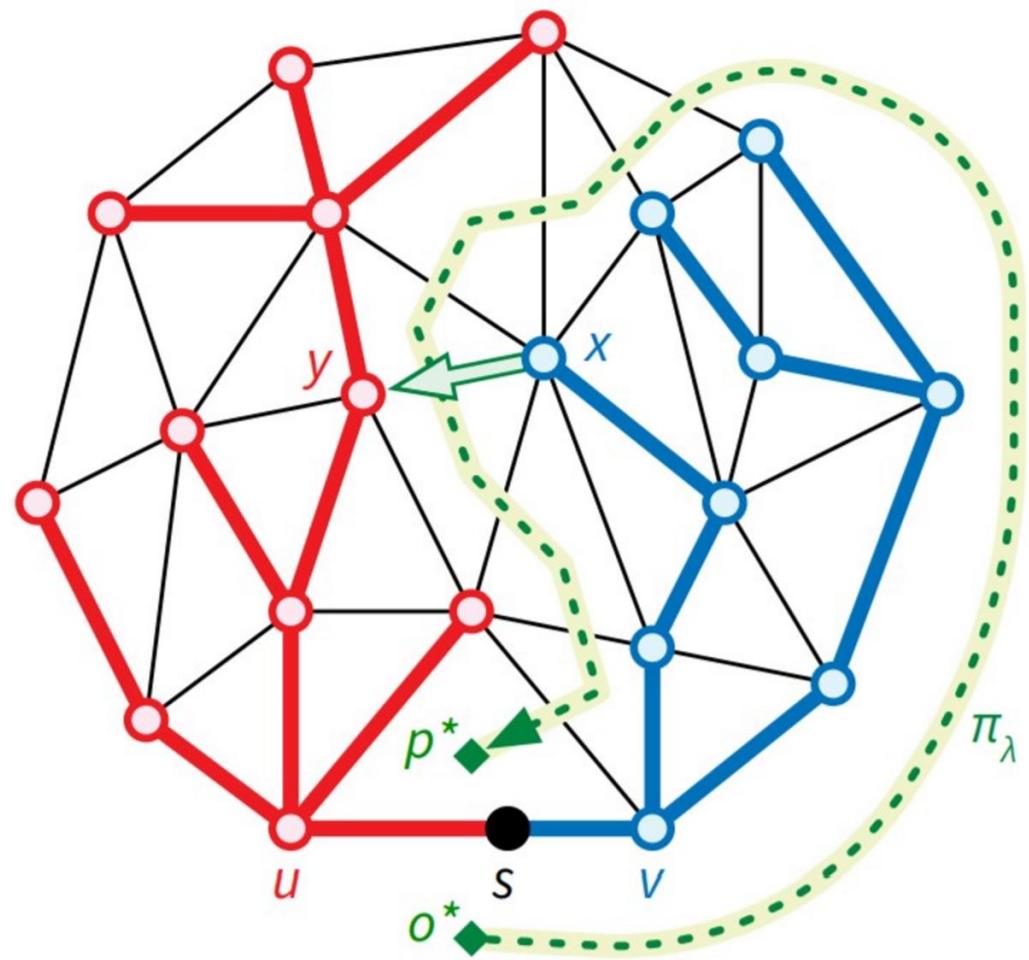


- Applications

- Linkage and Folding
- Mesh Generation



COURSE OUTLINE

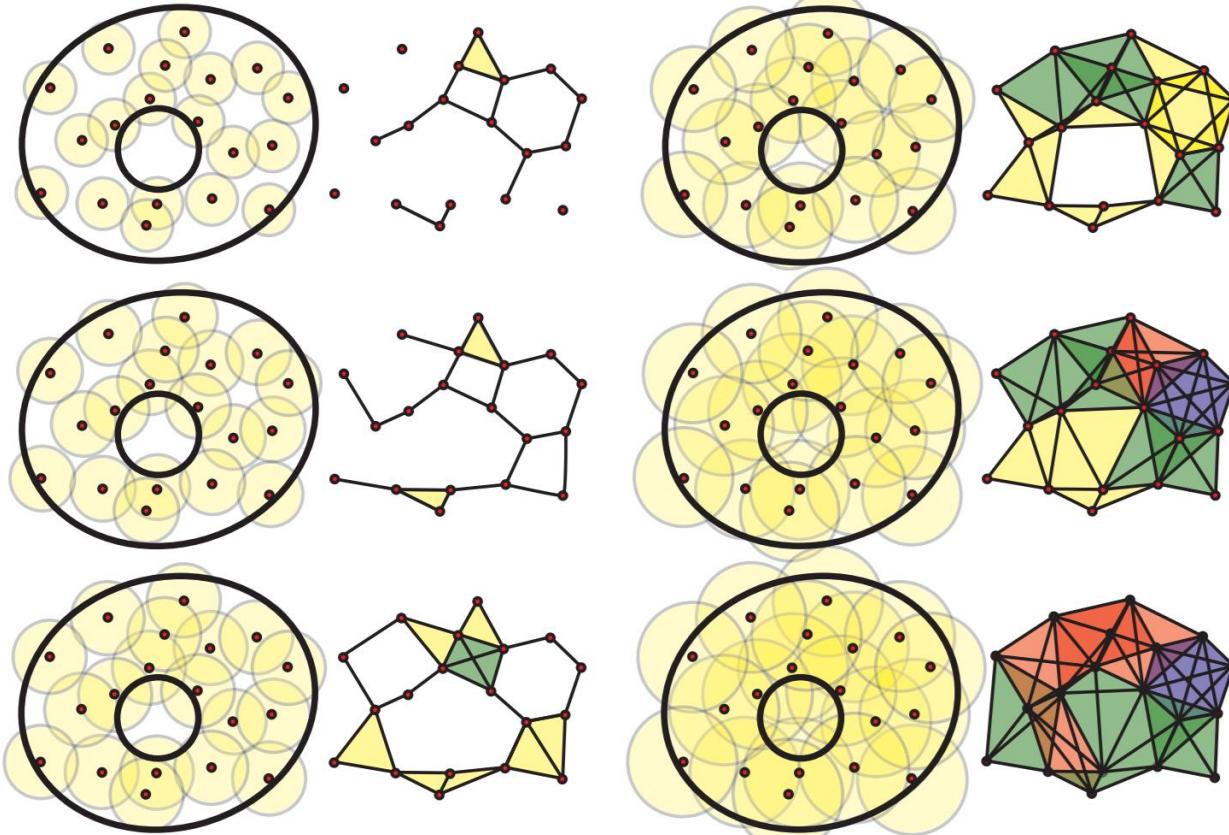


■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization



COURSE OUTLINE

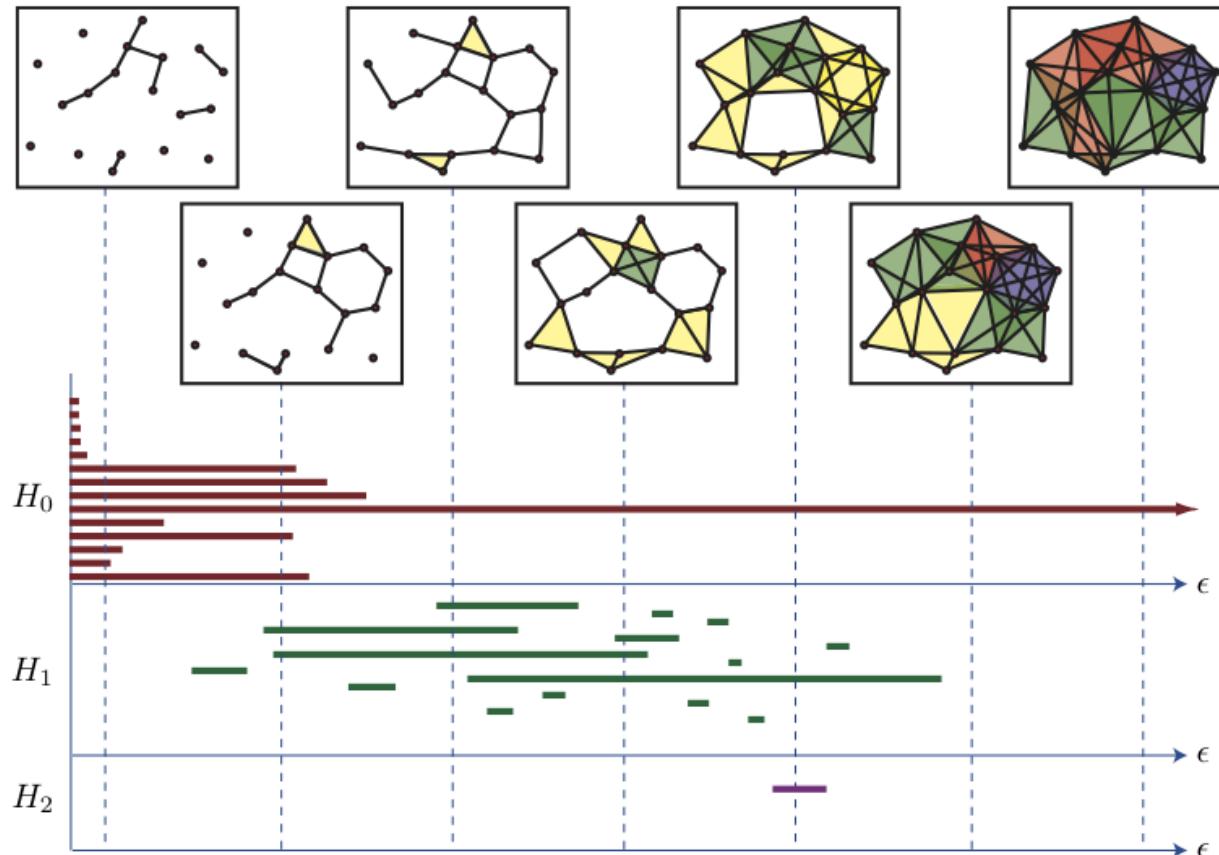


■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- Topological Data Analysis



COURSE OUTLINE



■ Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- Topological Data Analysis



COURSE OUTLINE

- Theory

- Curves
- Surfaces/Complexes
- Homotopy
- Homology

- Applications

- Linkage and Folding
- Mesh Generation
- Optimization
- Topological Data Analysis



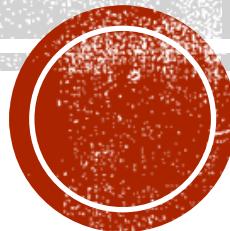
LOGISTICS

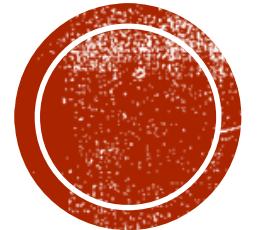
- Course webpage
- Lectures
- Office Hours
- Discussion forum
- Homework
- Project





STOP ME IF YOU ARE LOST





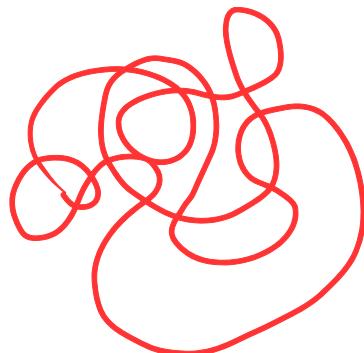
SIMPLE PLANAR CURVES



FORMAL DEFINITIONS

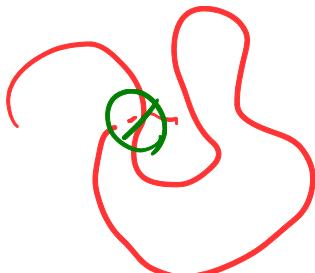
- Planar curve

$$\gamma: S^1 \rightarrow \mathbb{R}^2$$



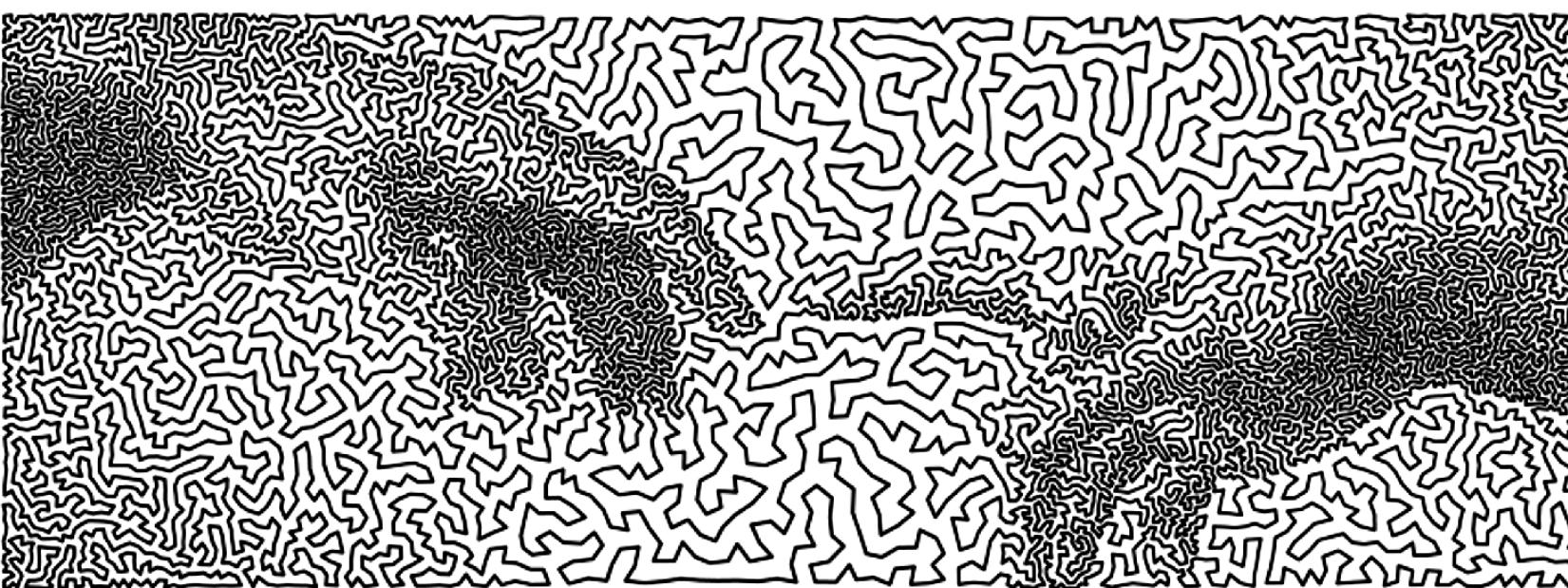
- Simple

γ embedding
injective



no self-crossings





SIMPLE PLANAR CURVES

Jordan curve

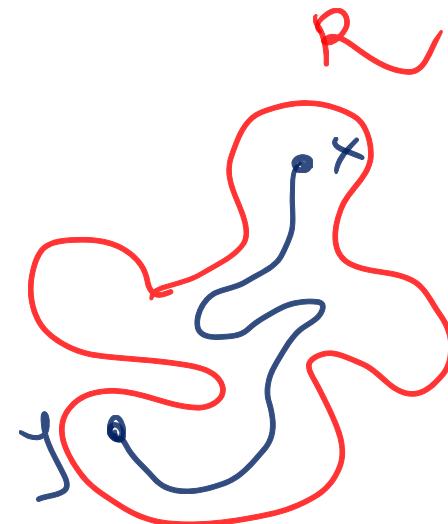
[Jordan 1887]



FORMAL DEFINITIONS

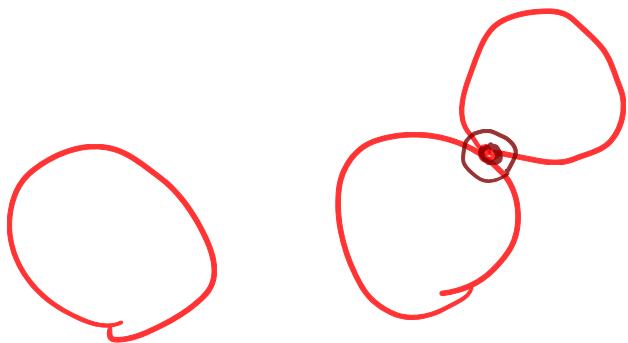
- **Connected** $x \sim y \text{ if } \exists \pi : [0, 1] \rightarrow R, \pi(0) = x, \pi(1) = y$.

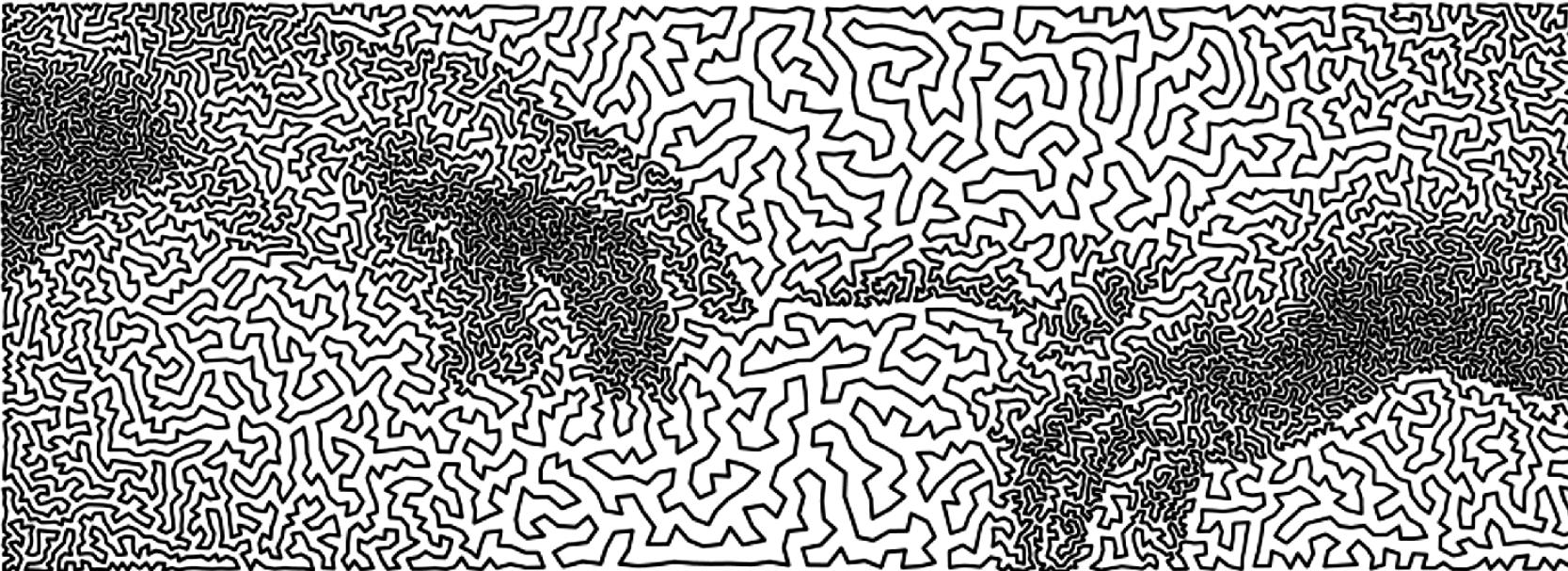
$R \subseteq \mathbb{R}^2$
 R connected if $\forall x, y \in R, x \sim y$



- **Connected component**

maximal subregions that are connected.

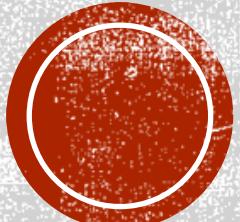


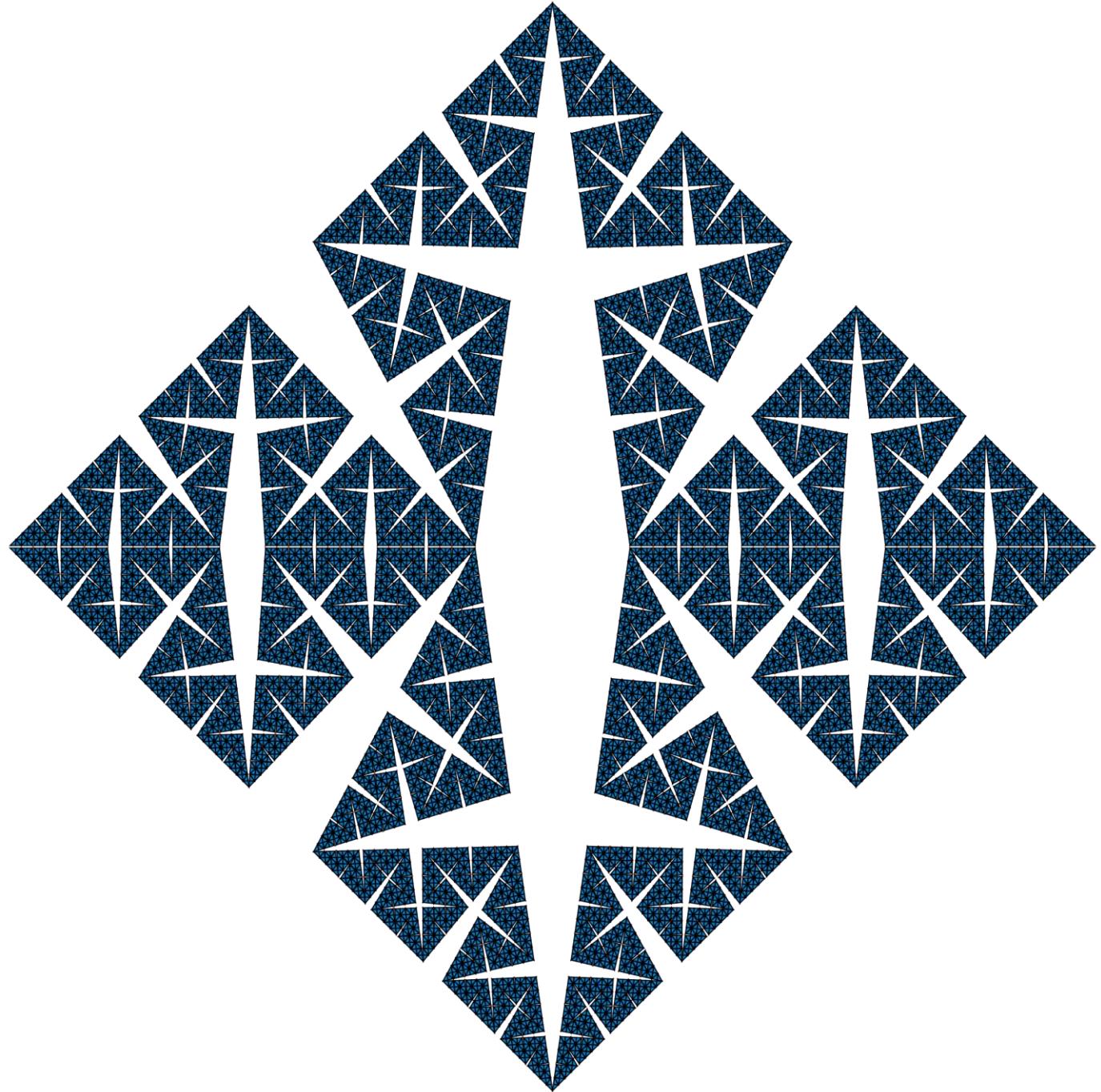


JORDAN CURVE THEOREM

[Bolzano ~1800s] [Jordan 1887]

Any simple closed curve C separates $\mathbb{R}^2 \setminus C$ into exactly
two connected components





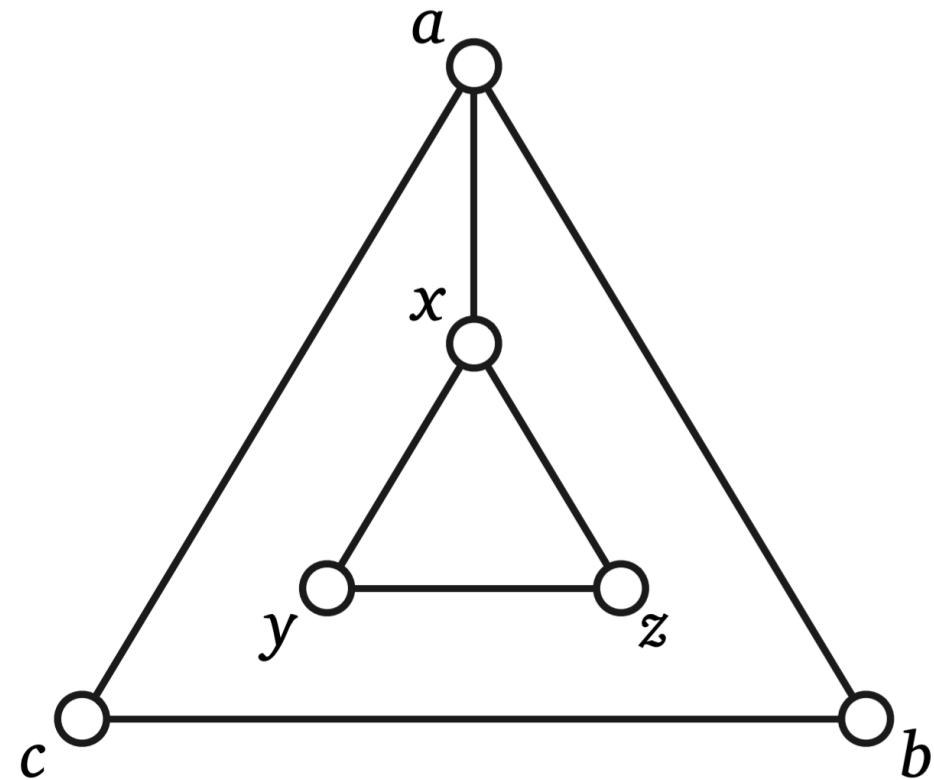
SIMPLE PLANAR CURVES

Osgood curve
[Osgood 1903]

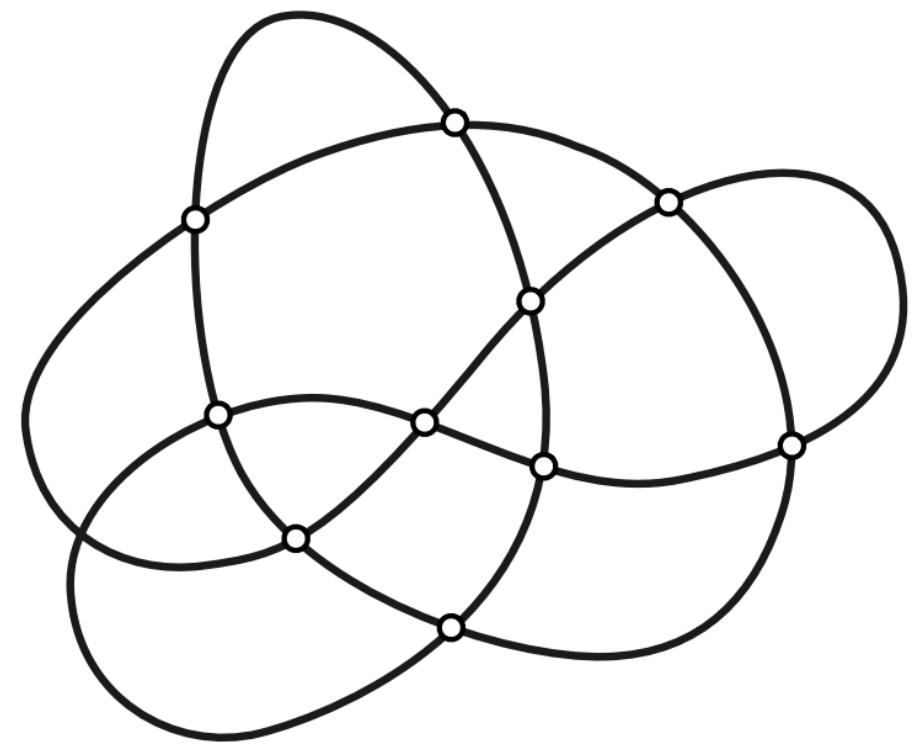


REPRESENTATION OF CURVES

- Polygonal

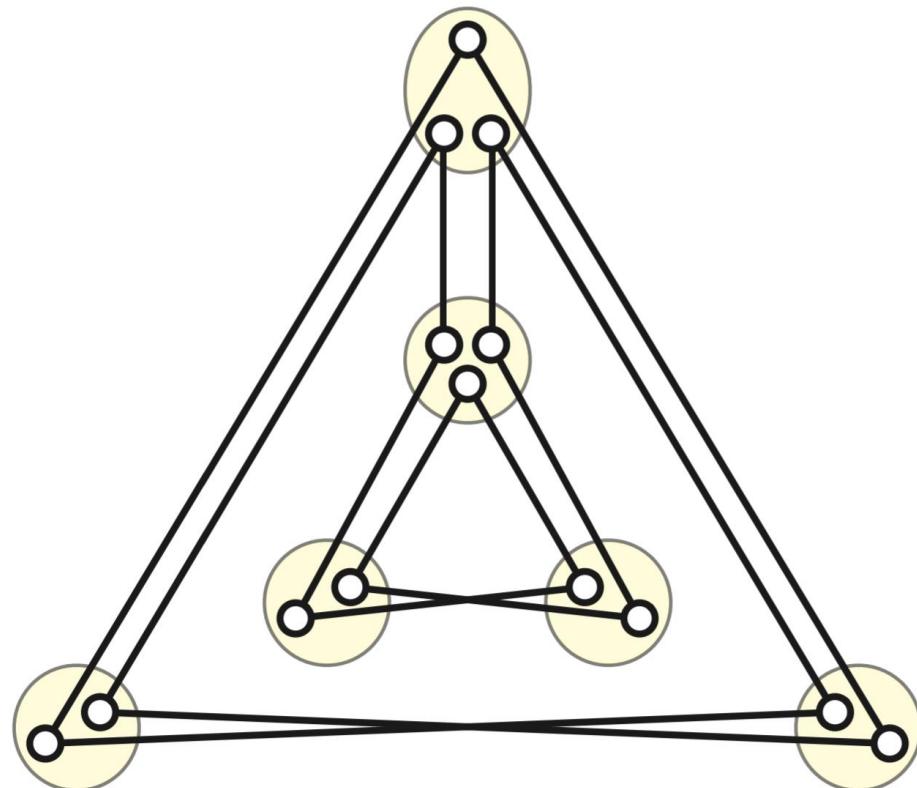


- Generic

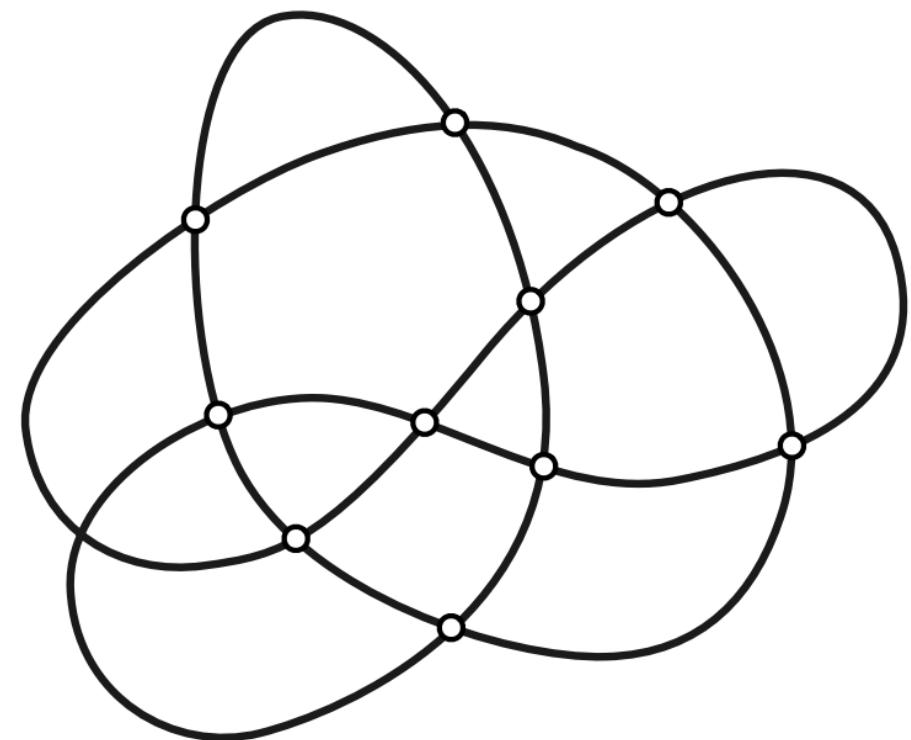


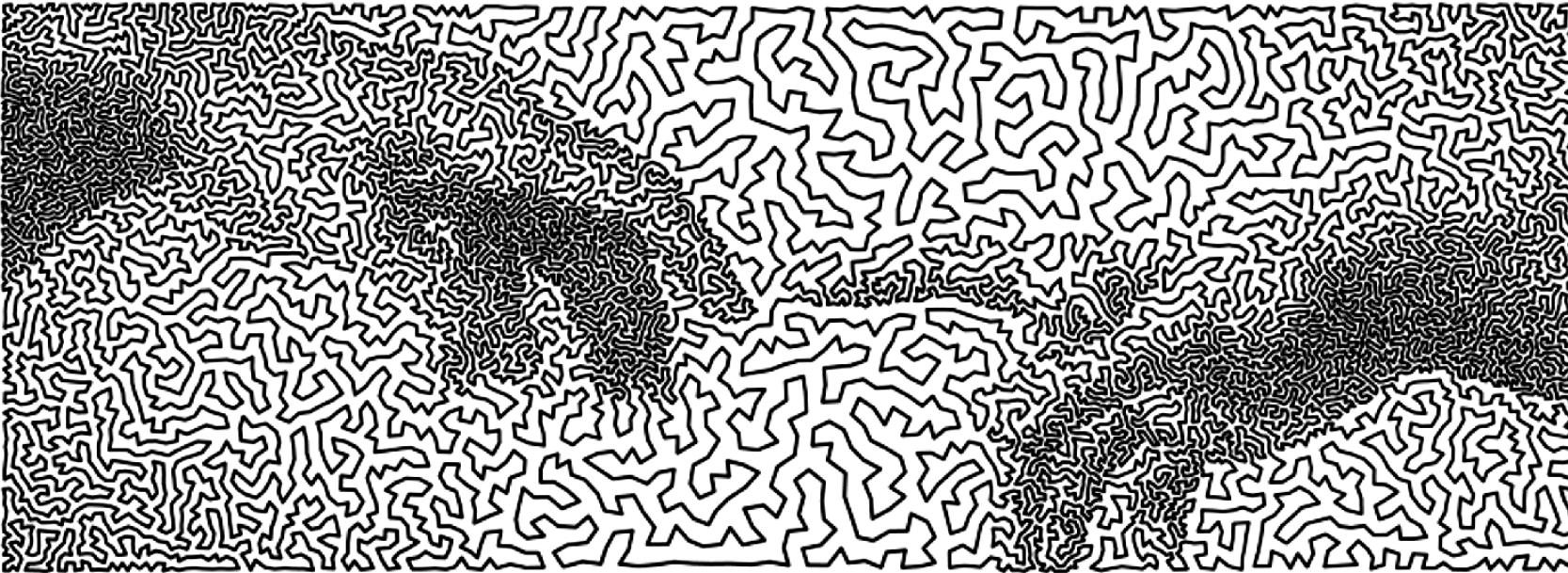
REPRESENTATION OF CURVES

- Polygonal



- Generic

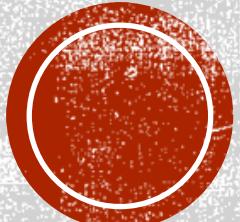




JORDAN POLYGON THEOREM

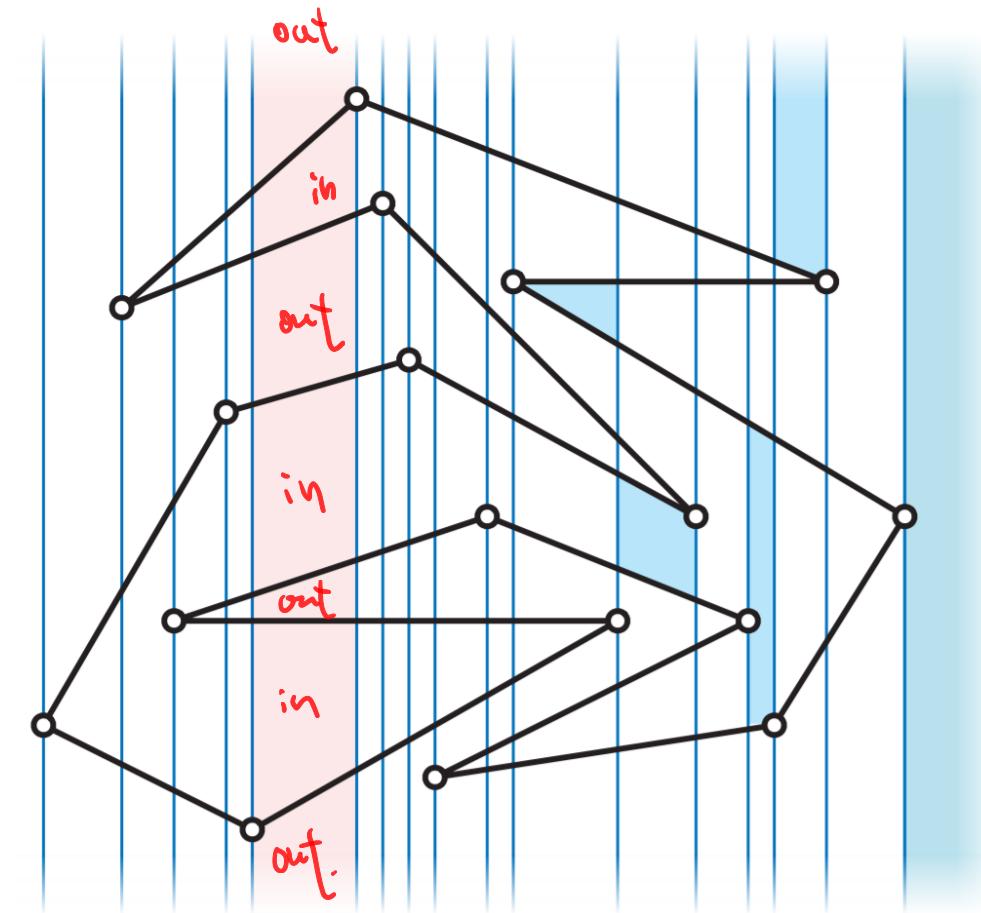
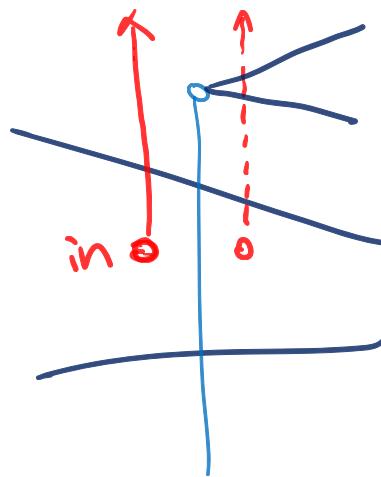
[Dehn 1899]

Any simple **polygon** P separates $\mathbb{R}^2 \setminus P$ into exactly two connected components



PROOF OF JORDAN POLYGON THEOREM

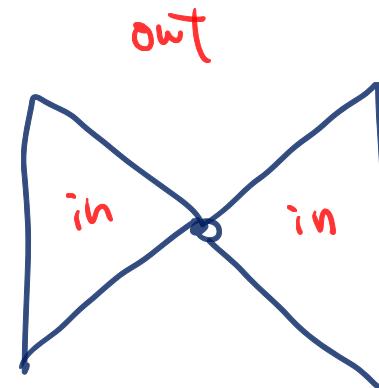
- Intuition: Parity argument



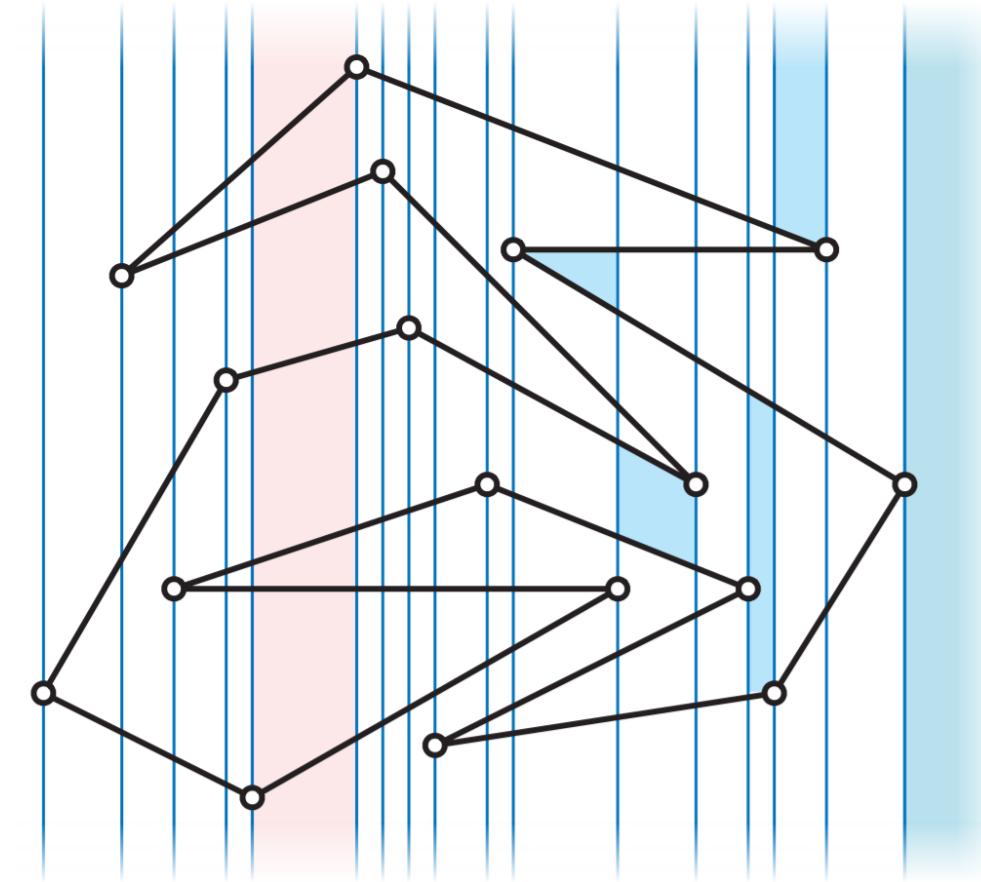
PROOF OF JORDAN POLYGON THEOREM

- Lemma ≥ 2

- Parity argument

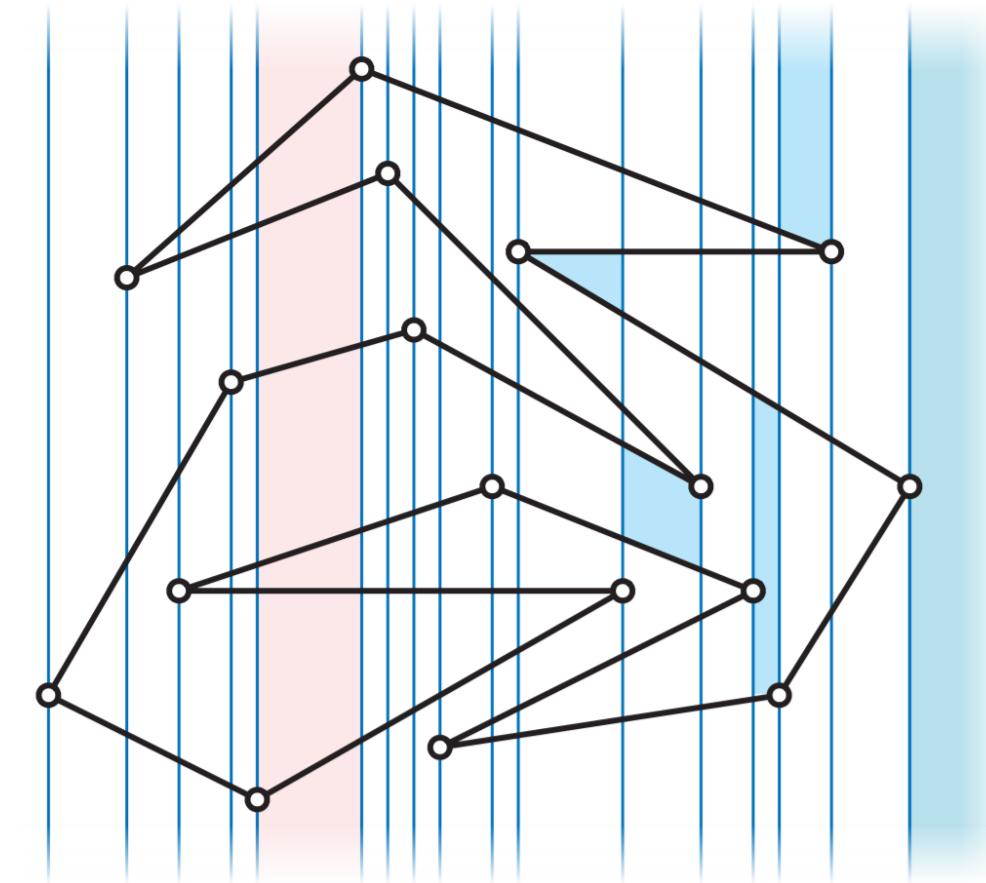
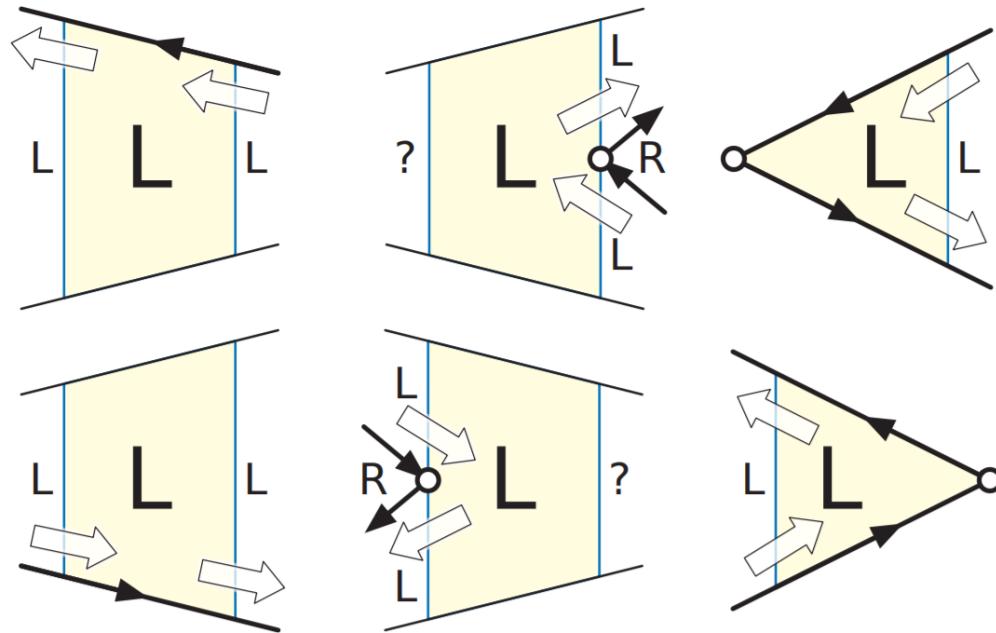


- Lemma $\leq 2?$



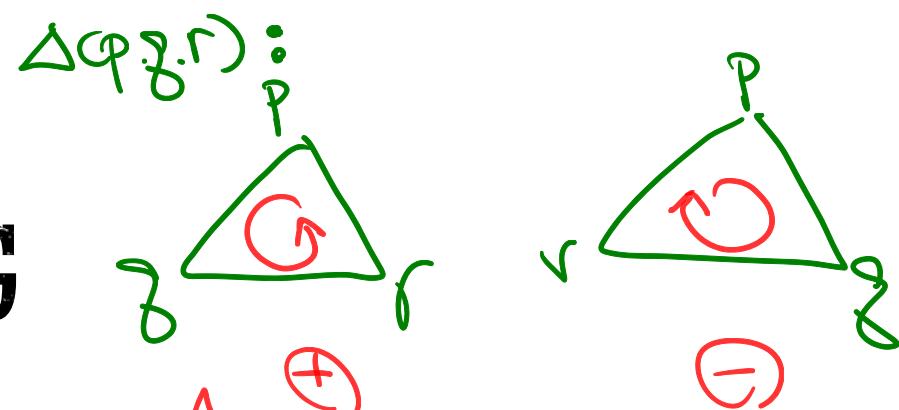
PROOF OF JORDAN POLYGON THEOREM

- Lemma ≤ 2



$$\Delta(p,g,r) = \begin{vmatrix} 1 & p.x & p.y \\ 1 & g.x & g.y \\ 1 & r.x & r.y \end{vmatrix}$$

INSIDE-POLYGON TESTING



Inside Polygon? ($P-g$):

sign $\leftarrow 0$

for each segment pr :

$\Delta \leftarrow \Delta(p,g,r)$

sign of triangle (p,g,r)

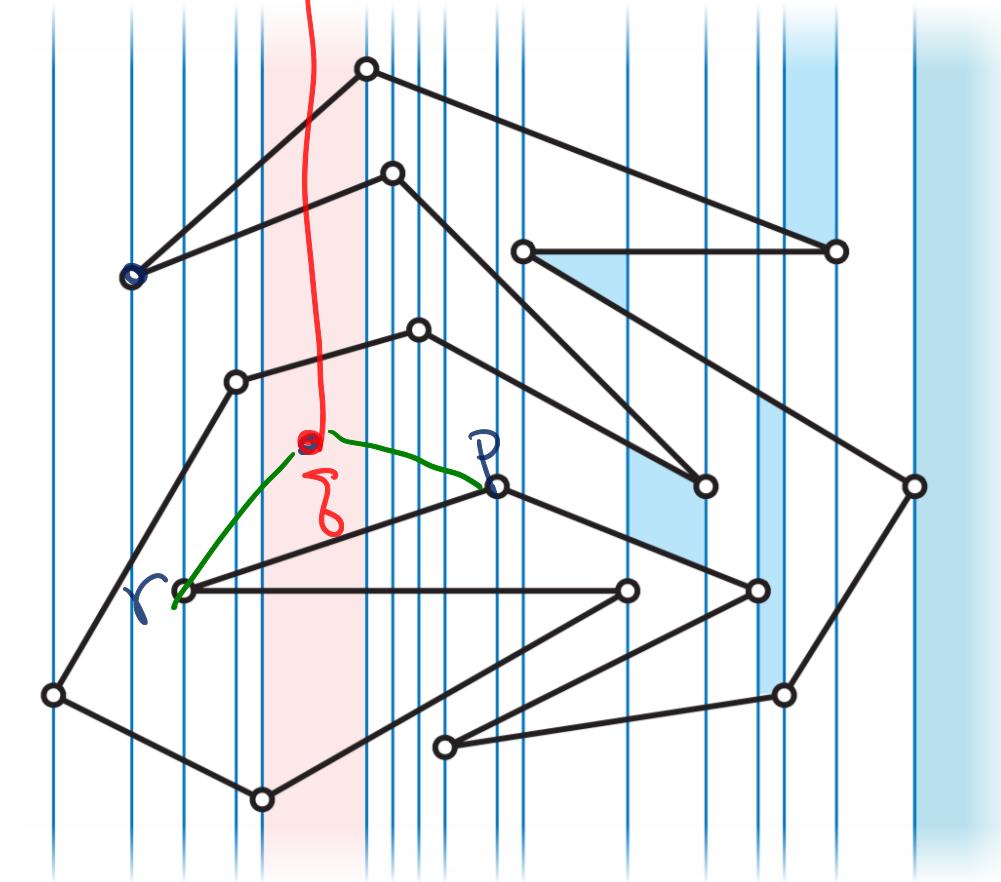
if $p.x \leq g.x < r.x$:

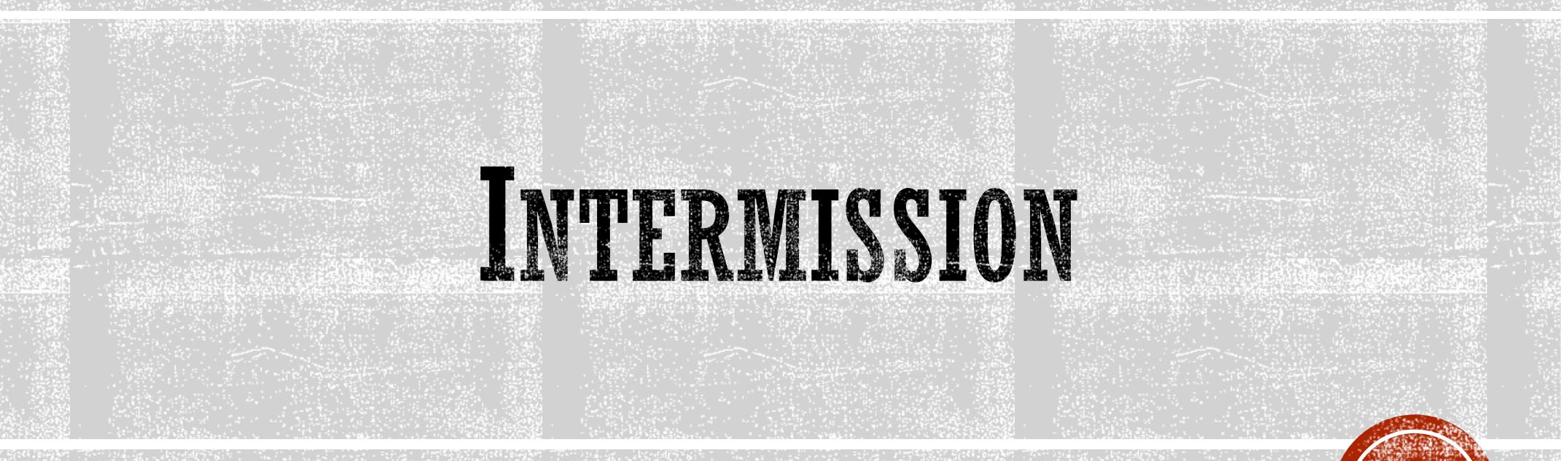
sign $\leftarrow -\Delta \cdot \text{sign}$ crossing above red

if $r.x \leq g.x < p.x$:

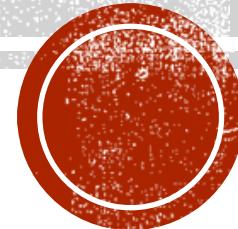
sign $\leftarrow \Delta \cdot \text{sign}$ crossing v below red

return sign



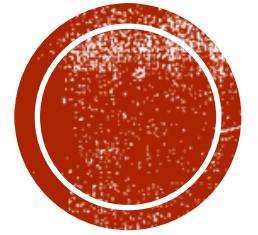


INTERMISSION



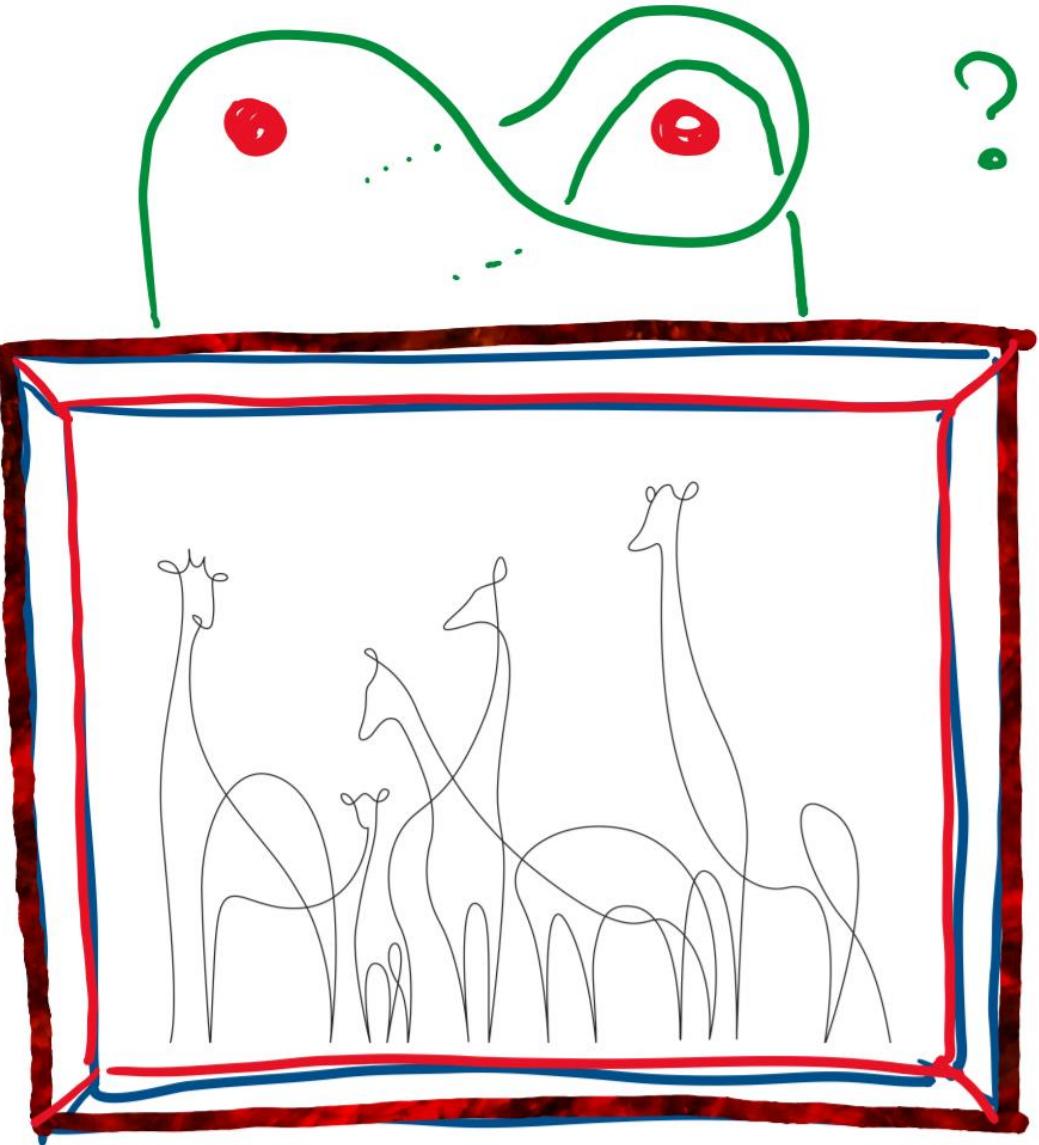
FOOD FOR THOUGHT.

How to compute the area of a simple polygon?



WINDING NUMBER

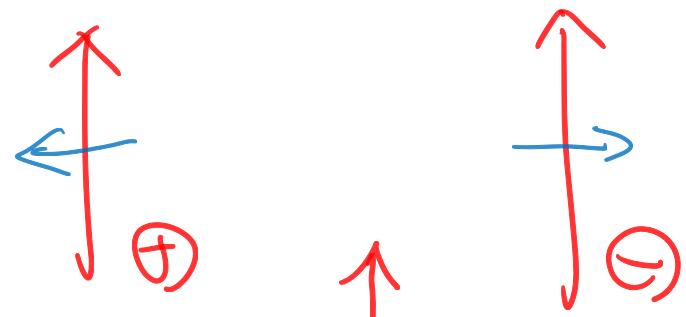
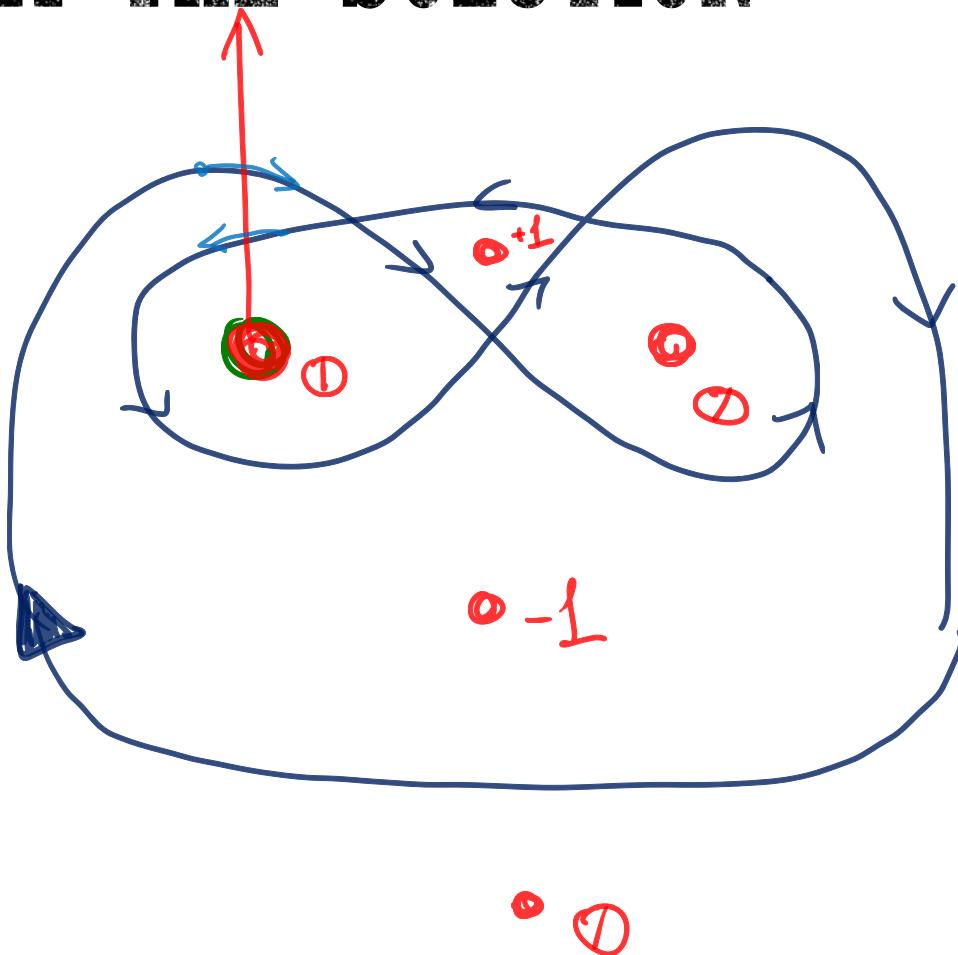




PICTURE-HANGING PUZZLE

Can you hang a picture with two nails, such that the picture falls if either nail is pulled?

LOOK AT THE SOLUTION



COMPUTING WINDING NUMBER

Winding Number (P, g) :

wind $\leftarrow \emptyset$

for each segment pr :

$\Delta \leftarrow \Delta(p, g, r)$

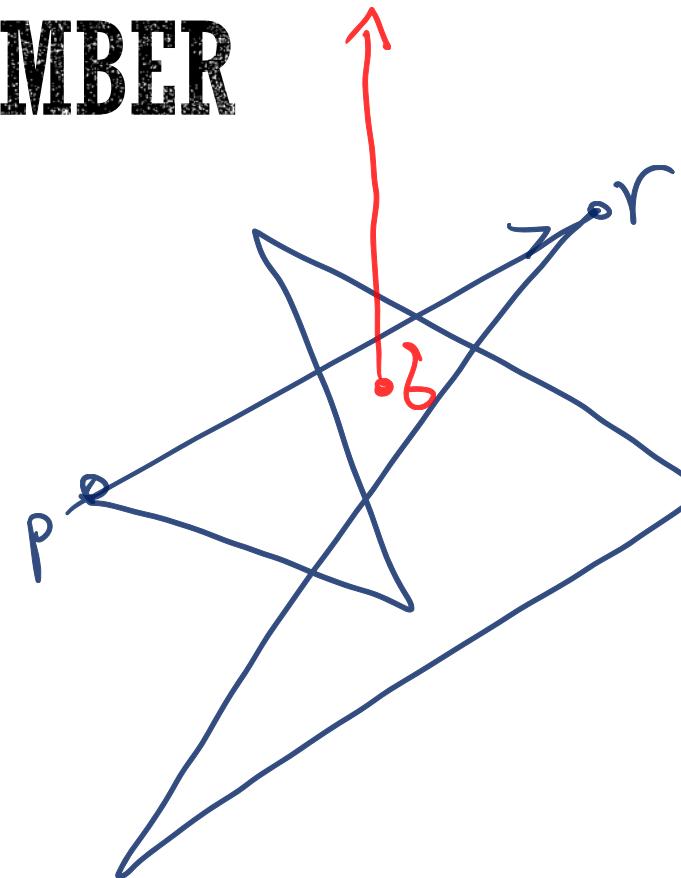
if $p.x \leq g.x < r.x$
and $\Delta = +1$:

wind --

if $r.x \leq g.x < p.x$
and $\Delta = -1$:

wind ++

return wind



$WIND_q(P)$ INVARIANT UNDER MORPHING

- Homotopy





CURVE MORPHING





CURVE MORPHING

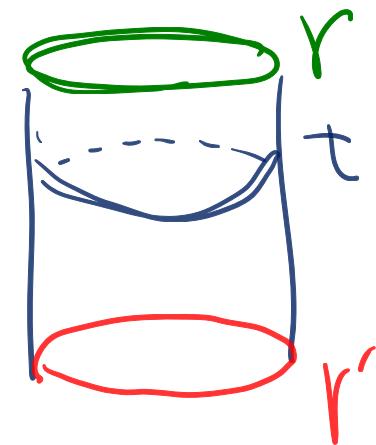
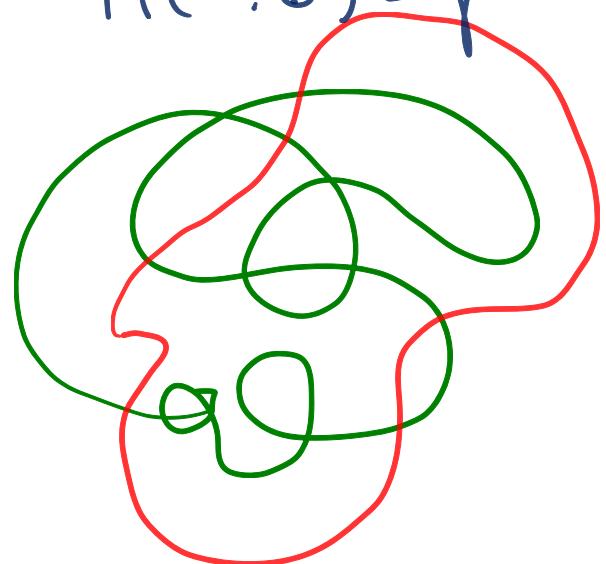


$\text{WIND}_q(P)$ INVARIANT UNDER MORPHING

■ Homotopy

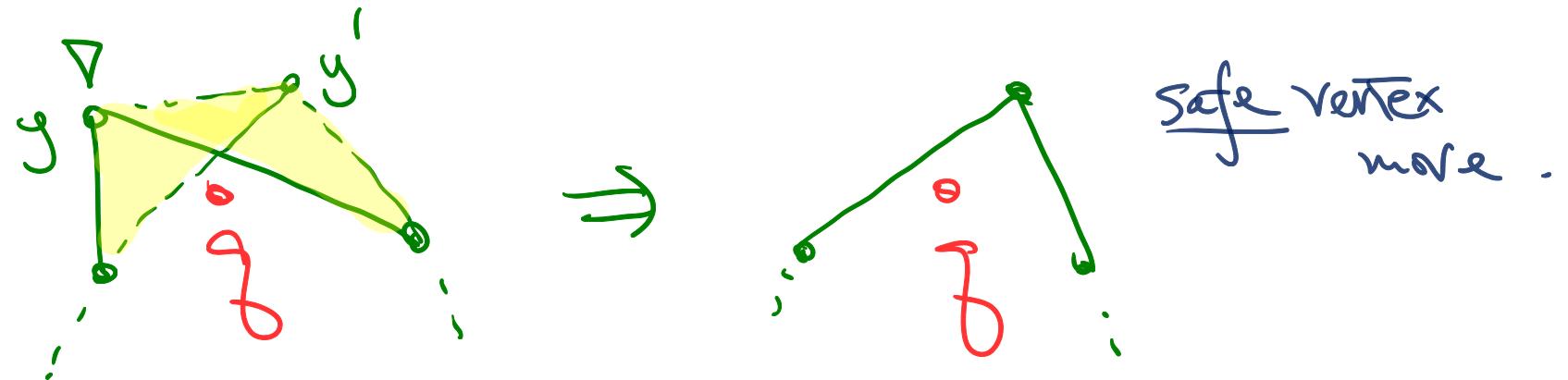
cont. change from $\gamma \rightarrow \gamma'$

$$H: S^1 \times [0,1] \rightarrow \mathbb{R}^2$$
$$H(\cdot, 0) = \gamma \quad H(\cdot, 1) = \gamma'$$

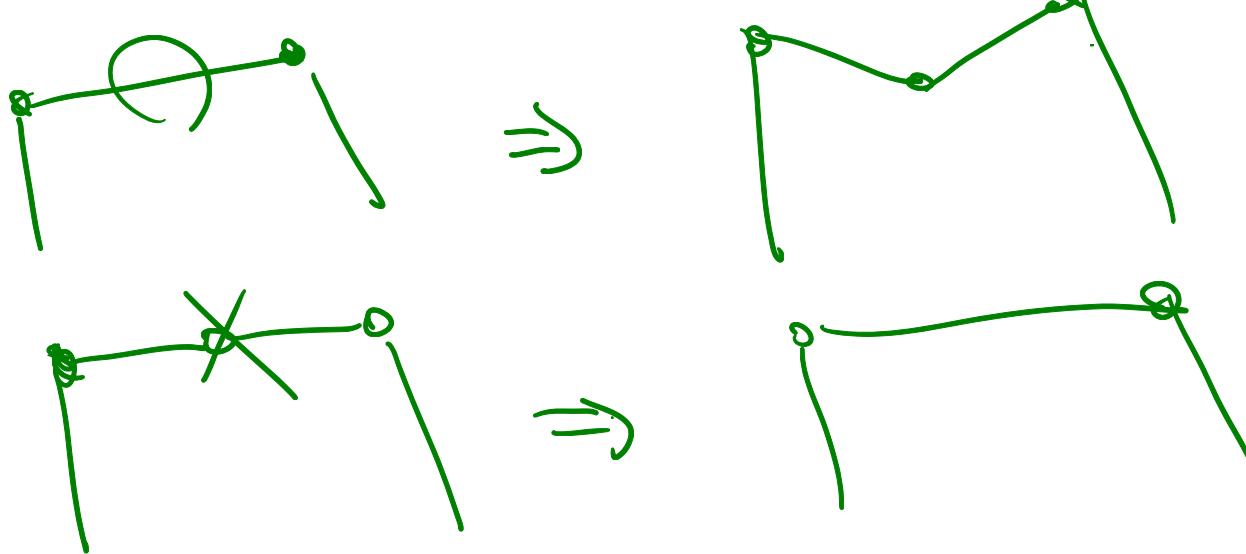


$\text{WIND}_q(P)$ INVARIANT UNDER MORPHING

- Homotopy

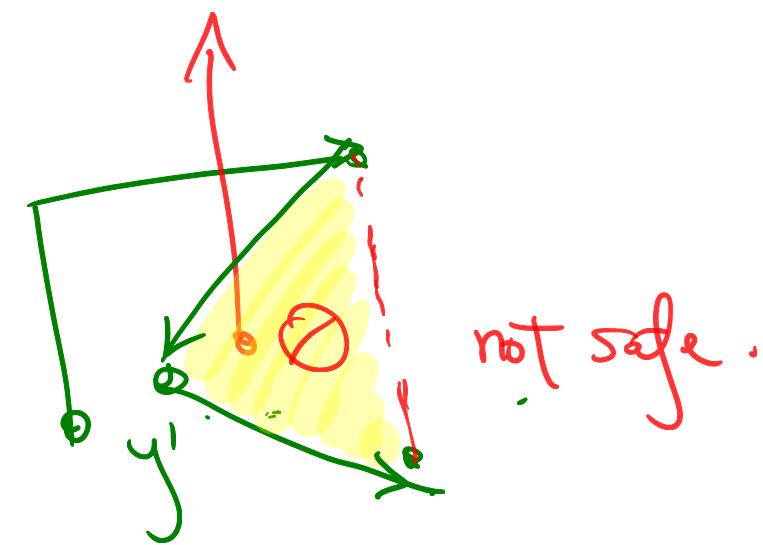
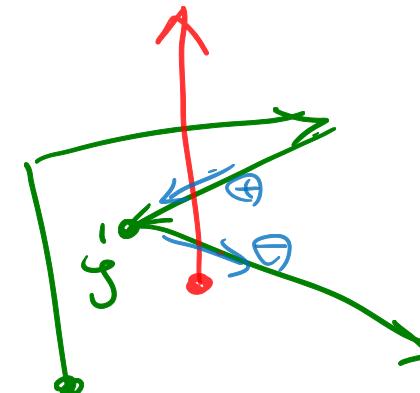
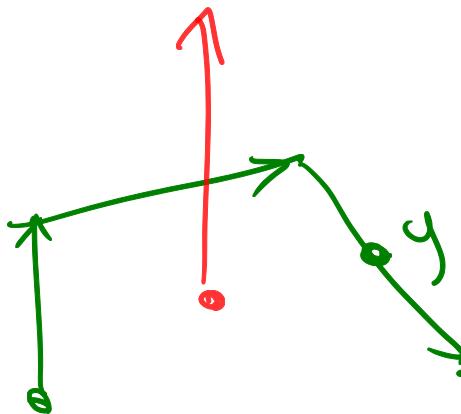


- Vertex move



THEOREM. $\text{Wind}_q(P)$ is invariant under safe vertex moves.

Pf sketch.



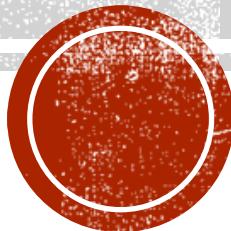
not safe.



THEOREM. Two polygons P and Q are homotopic in $\mathbb{R}^2 \setminus q$ if and only if they have the same Wind_q . [Hopf 1935]



**WIND_q IS A COMPLETE
HOMOTOPIC INVARIANT!**



TAKEAWAY.

Planar curve can be described by how many times it goes around reference points.