### The New Assembler-Simulator (nas)

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```
http://www.cs.hku.hk/~fcmlau/nas.zip

$ flex nas.l
$ bison -d nas.y
$ gcc -o nas lex.yy.c nas.tab.c
$ nas fact.as
$ ...
```

#### nas

- A stack machine: all operations use push/pop
- Variables
  - In sas (previous tutorial), there are 26 of
  - them. and functions can have local variables whose names may clash with the global ones
  - In nas, variables are <u>unnamed</u>, stored inside the stack, and there can be as many as you want (or as the stack can hold)

```
push "Enter 5 numbers: "; puts_
geti // = fp[0]
geti
geti
geti
geti
geti
fy and these as variables

Treat these as variables

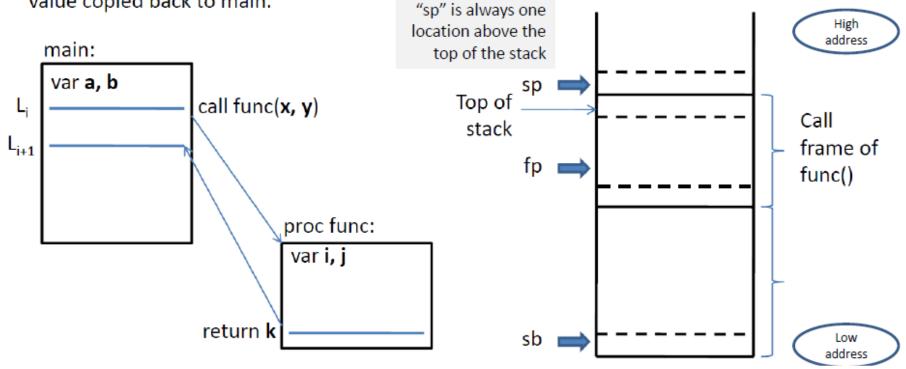
The "index register" (registers are named)

geti // = fp[4]
push 4; pop in // in = 4
```

Note: This "main" is different from C's main() which is a function; here it is simply the outermost scope. Hence, **a**, **b** are global variables; **i**, **j** are local variables of func(). Assuming pass-by-value, **x**, **y** are copied to func() and treated as local variables. **k** is the return value copied back to main.

### **Function Call**

When func() is executing, the stack should look like:



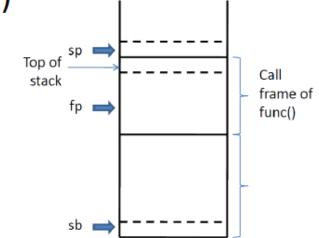
<u>Special registers of the machine</u>: **sp** (stack pointer), **fp** (frame pointer) which points **near** the bottom of the current frame, and **sb** (stack base) ... and **in** (index register) for implementing arrays

Also called **ac** (accumulator), and can be used as a general register, if you need one.

### Using Variables

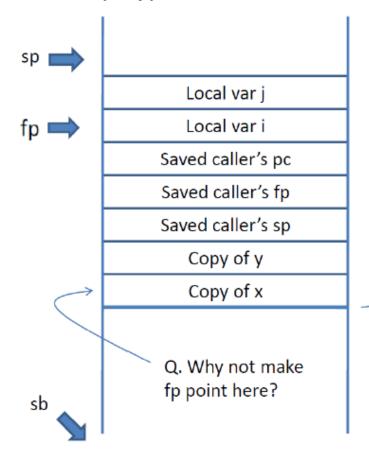
Everything is on the stack, except strings (their addresses are pushed instead)

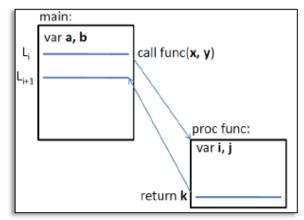
- To access local variables inside a function:
  - Relative to fp  $\rightarrow$  e.g. "fp[-1]"
- To access global variables:
  - Relative to sb  $\rightarrow$  e.g. "sb[3]"
- Only can access own frame and main, but not other frames in between



### A Call Frame

" call func(x, y)"





Call frame of function/callee

i is referred to as fp[0]

j ... fp[1]

x ... fp[-5]

y ... fp[-4]

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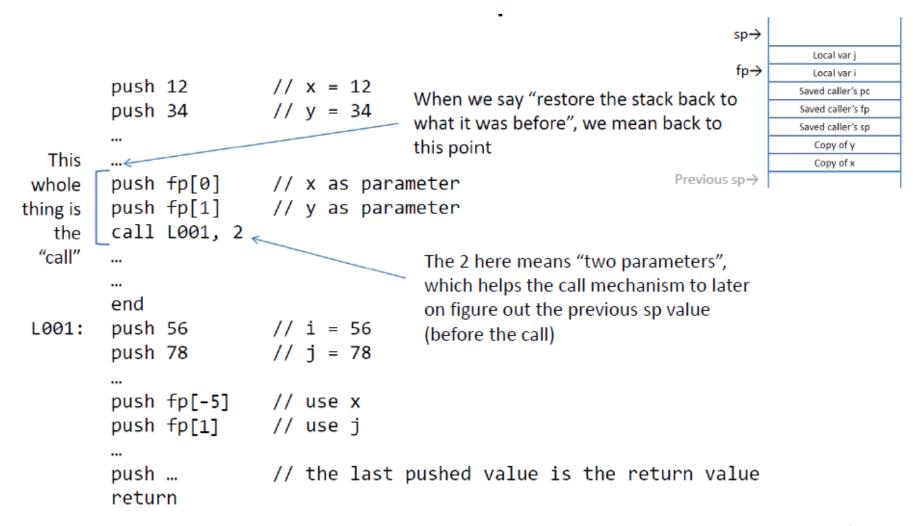
a ... sb[0]

b ... sb[1]

...

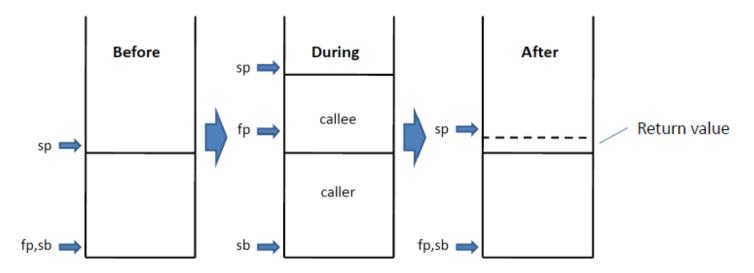
k, the <u>return value</u> will be left on top of the stack after the restoration

### The General Call Pattern



#### Stack Frames

- When a program begins, fp = sp = sb = 0
  - sp rises and drops as actions in Main unfolds
- After a call, the stack must be restored to what it was before the call (+ the return value if there's one)
- At the moment of calling, caller's sp, fp, and pc (program counter) which points at the caller's next instruction  $(L_{i+1})$  are saved in the callee's frame



### Push & Pop

```
push 123
               push "123" onto the stack
push -456
               push "-456" onto the stack
push fp[2]
               push the content of "where fp is pointing + 2"
               push the "... - 7"
push fp[-7]
pop sb[4]
               pop the stack and store the value in "the stack bottom + 4"
               push ... "where fp is pointing + the value of in"
push fp[in]
               Illegal; instead, you can make the value of in negative
push fp[-in]
push in
               push the value of in
               pop the stack and store the value in in
pop in
push fp
               push the value of fp
push fp[0]
                push the content of where in the stack fp is pointing
```

### Example: max.as

```
// max.as
                                                   Do not print \n
        push "Enter 2 numbers: "; puts_
        geti
               Reads inputs and passes them as arguments to function
        geti
        call L001, 2
                                                   Which is at the stack's top
        puti_ // print the return value
        push " is larger"; puts
        end
L001: push fp[-4] push fp[-5]
                       - Retrieves and pushes the two arguments
        compgt
        j1 L002
                                Return value
        push fp[-5]
        ret
L002: push fp[-4]
        ret
```

# Example: fact.as

```
// recursive fact.as
        push "Please enter a +ve int < 13: "; puts_</pre>
        geti
        call L001, 1
                                 Read n
        puti -
                                  Call fact(n)
        end
                                  Print return value
// factorial():
L001: push fp[-4]
        i0 L002
                                                 n
        push fp[-4]; push 1; sub -
                                                 n = n - 1
        call L001, 1 // recursive call
                                                 Return n x fact(n - 1)
        push fp[-4]
        mu1
        ret
L002: push 1
        ret
```

## Example: rev-c.as

```
// rev-c.as
        push "Please enter a line:"; puts
        push 0; pop in
                                                    // in = 0
L001:
        getc; // NO pop fp[in] here !!
                                                   // fp[in] = getc
        push fp[in]; push 10; compeq; j1 L002  // if newline goto L002
        push in; push 1; add; pop in
                                                    // in++
        jmp L001
L002:
        push in; push 1; sub; pop in
                                                    // in--
        push fp[in]; putc_
        push in; j0 L003; jmp L002
        push ''; putc
L003:
        end
```

Print a newline