

CSE 522: RTES

Assignment 01

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Description of implementation:

1) Parse and create arrays:

- a) Use task_model.h to create the thread, mutex, semaphore, and timer arrays.

2) The “activate” command:

- a) I have used “conditional variable” to implement functionality to start all threads after getting the “active” command from the user.
- b) All threads will wait for conditional variable initially and upon signal/broadcast, they will start the computation.

3) Execution time:

- a) I have not used traditional k_msleep or k_sleep to implement <wait_for_period> Instead I am using a combination of a timer and semaphore for each thread.
- b) Each thread will start the timer at the start of the computation (i.e. first statement of while(1) with timeout equal to the period of the thread.
- c) The timer callback will give “sem_give” the semaphore.
- d) Each thread will get blocked at “sem_take” upon completion of the task_body.

4) Deadline error message:

- a) I am maintaining a global “compute_flag” array of size equal to NUM_THREADS.
- b) A timer callback will check the flag and throw the error message accordingly if compute_flag[thread_id] is false.
- c) Each thread will set compute_flag[thread_id] to true upon completion of the task.

System View Screenshots:

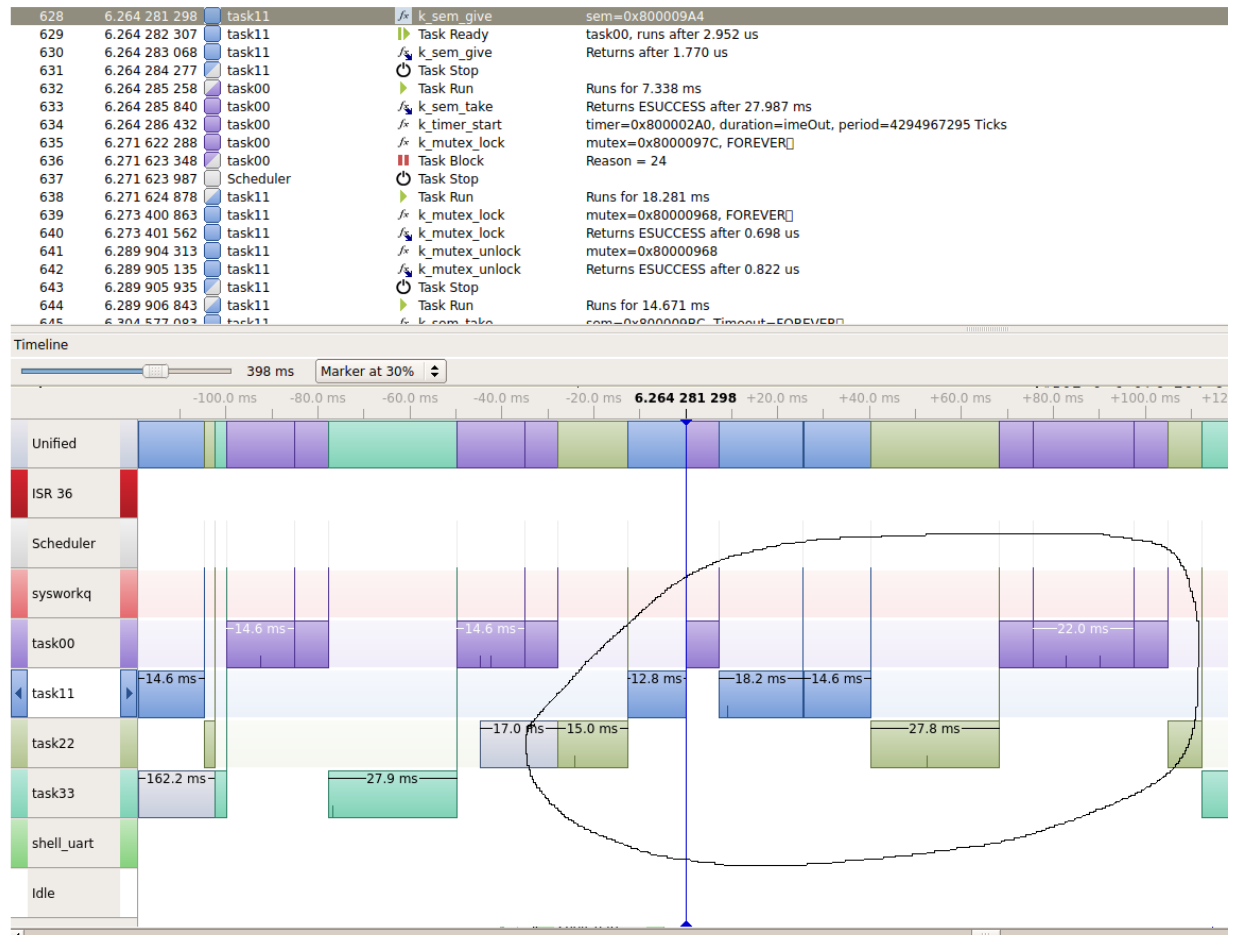
The task set used for both screenshots is given below

```
#define NUM_MUTEXES 3           // number of mutexes
#define NUM_THREADS 4           // number of threads
#define TOTAL_TIME 4000        // total execution time in milliseconds

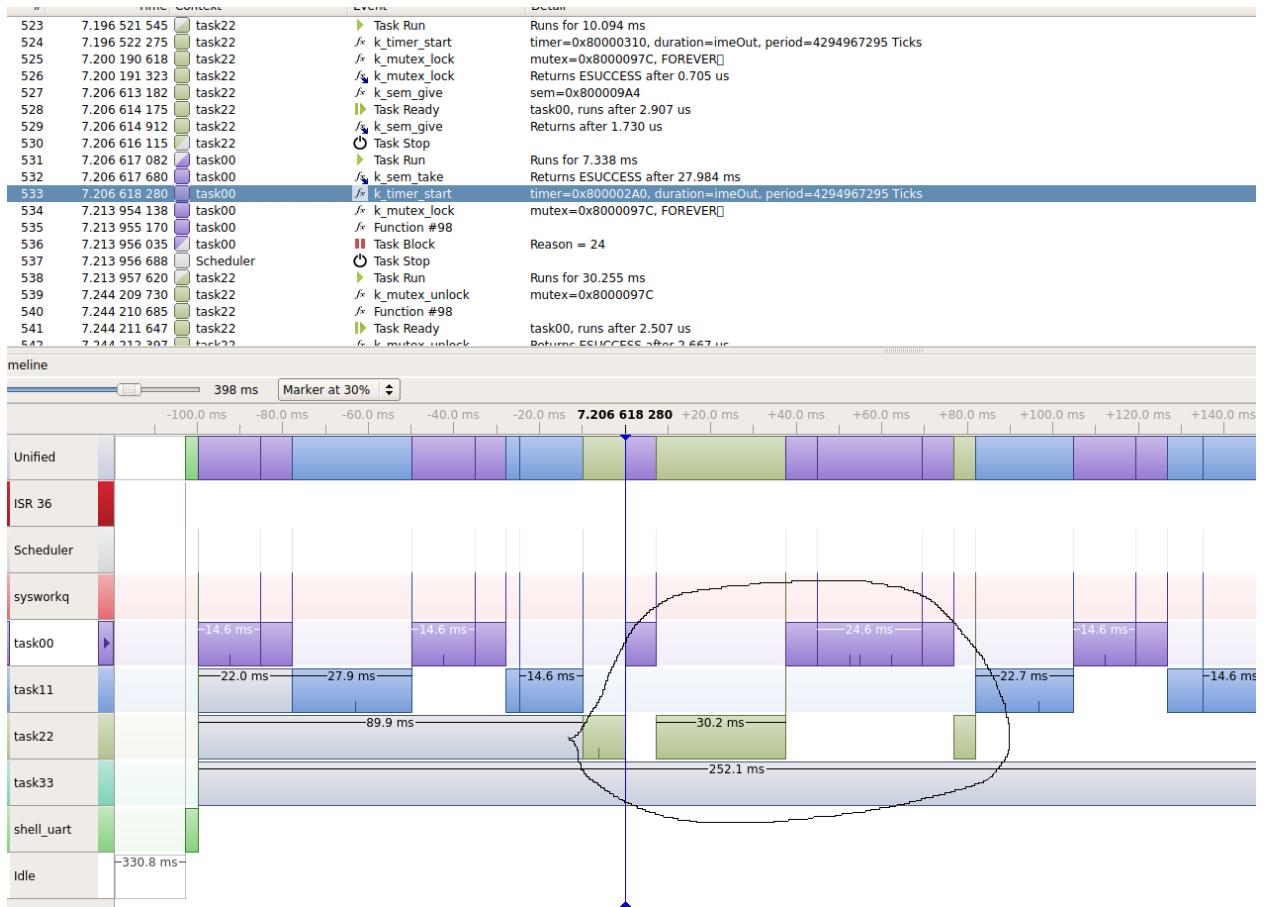
#define THREAD0 {"task00", 2, 50, {400000, 400000, 400000}, 1}
#define THREAD1 {"task11", 3, 160, {800000, 900000, 800000}, 0}
#define THREAD2 {"task22", 4, 220, {200000, 2000000, 400000}, 1}
```

```
#define THREAD3 {"task33", 5, 360, {200000, 2000000, 400000}, 2}
```

1) CONFIG_PRIORITY_CEILING=10



2) CONFIG_PRIORITY_CEILING=0



Explanation of above screenshots.

Scenario:

In the above figures, the priority of task22 < task11 < task00.

Task00 and task22 share the common mutex (resource).

Task22 gets preempted by task11, task11 gets preempted by task00.

Task00 tries to get the mutex lock which is currently acquired by task22 and hence task00 goes in a waiting state.

- 1) When CONFIG_PRIORITY_CEILING is set to 0 the task22 now becomes a task with the highest priority and the scheduler makes task22 in a ready state to release the mutex lock as soon as possible.
- 2) When CONFIG_PRIORITY_CEILING is set to 10, task11 gets scheduled over task22 as the priority of task11 is greater than task22 and hence task00 gets delayed.