

Predicate Logic (Pt.2)

CS 5012



Predicate Sentences and Formulae (1)

- **User supplies symbols for constants, functions and predicates**
- **Constant symbols**, which generally represent names and fixed numerical values:
 - Mary, 3, Green
- **Function symbols**, which map individuals to individuals
 - father-of(Mary) = John, color-of(Sky) = Blue
- **Predicate symbols**, which map individuals to truth values
 - greater(5,3), green(Grass) , color(Grass, Green)

Predicate Sentences and Formulae (2)

- The input provided by the user is combined with the symbols of the predicate logic to complete a formulae
- Variable symbols
 - E.g., x , y , foo
- **Connectives**
 - Same as in PL: not (\neg), and (\wedge), or (\vee), implies (\rightarrow), if and only if (*biconditional* \leftrightarrow)
- **Quantifiers**
 - Universal $\forall x$ or (Ax)
 - Existential $\exists x$ or (Ex)

Building Sentences (1)

- Predicate Logic sentences are built from Terms and Atoms
- A **term** (*denoting a real-world individual*) is a constant symbol, a variable symbol, or an n-place function of n terms.
 - x and $f(x_1, \dots, x_n)$ are terms, where each x_i is a term
 - A term with no variables is a **ground term**
- An **atomic sentence** (*which has value true or false*) is an n-place predicate of n terms

Building Sentences (2)

- A **complex sentence** is formed from atomic sentences connected by the logical connectives:
 - $\neg P$, $P \vee Q$, $P \wedge Q$, $P \rightarrow Q$, $P \leftrightarrow Q$ where P and Q are sentences
- A **quantified sentence** adds quantifiers \forall and \exists
- A **well-formed formula (wff)** is a sentence containing no “free” variables. That is, **all variables are “bound”** by universal or existential quantifiers
 - $(\forall x)P(x,y)$ has x bound as a universally quantified variable, but y is free

Syntax of Predicate Logic (1)

Primitive vocabulary

(a) A set of terms:

- A set of individual **constants**: a, b, c, d, \dots
 - John, Mary, Steve, Loren
- A set of individual **variables**: x, y, z, x_1, x_2, \dots
 - he, she, it

Syntax of Predicate Logic (2)

Primitive vocabulary (cont'd)

(b) A set of predicates: P, Q, R, \dots

- One-place predicates: *is happy, is boring*
- Two-place predicates: *like, hate, love, meet*
- Three-place predicates: *introduce, give*

(c) A binary identity predicate: $=$

(d) The connectives of propositional logic:

$\neg, \wedge, \vee, \rightarrow, \leftrightarrow$

(e) Quantifiers: \forall, \exists

(f) brackets: $() []$

Syntax of Predicate Logic (3)

Syntactic rules

(a) If P is an n -place predicate and t_1, \dots, t_n are all terms, then $P(t_1, \dots, t_n)$ is an atomic formula

- John is happy: $H(j)$
- John loves Mary: $L(j,m)$
- John introduced Mary to Sue: $I(j,m,s)$

(b) If t_1 and t_2 are individual constants or variables, then $t_1 = t_2$ is a formula.

- John is Bill: $j=b$

Syntax of Predicate Logic (4)

NOTE: Φ is Greek letter Phi and Ψ is Greek letter Psi, being used as symbols for 2 different formulae below

- (c) If Φ is a formula, then $\neg \Phi$ is a formula.
- (d) If Φ and Ψ are formulas, then $(\Phi \wedge \Psi)$, $(\Phi \vee \Psi)$, $(\Phi \rightarrow \Psi)$, and $(\Phi \leftrightarrow \Psi)$ are formulas too
- (e) If Φ is a formula and x is a variable, then $\forall x\Phi$, and $\exists x\Phi$ are formulas too
 - Everyone is happy: $\forall x H(x)$
 - Someone loves John: $\exists x L(x, j)$
 - Everyone loves someone: $\forall x \exists y L(x, y)$, $\exists y \forall x L(x, y)$
- (f) Nothing else is a formula in predicate logic

Translate from English Sentences to Predicate Logic Formulae

NOTE: \forall = *For all, for each, every, ...*
 \exists = *There exists, there is at least one, ...*

(1) Every student is happy:

Right: $\forall x [\text{student}(x) \rightarrow \text{happy}(x)]$

Wrong: $\forall x [\text{student}(x) \wedge \text{happy}(x)]$

(2) Some students are happy:

Right: $\exists x [\text{student}(x) \wedge \text{happy}(x)]$

Wrong: $\exists x [\text{student}(x) \rightarrow \text{happy}(x)]$

1. *For all x , if x is a student, then x is happy*
2. *There exists an x , where x is a student and x is happy*

Translate from English Sentences to Predicate Logic Formulae

(3) No student complained:

- $\forall x [\text{student}(x) \rightarrow \neg \text{complained}(x)]$
- $\neg \exists x [\text{student}(x) \wedge \text{complained}(x)]$

(4) Not every student complained:

- $\neg \forall x [\text{student}(x) \rightarrow \text{complained}(x)]$
- $\exists x [\text{student}(x) \wedge \neg \text{complained}(x)]$

3.
 - For all x , if member of set 1 then not member of set 2
 - There doesn't exist an x , member of set 1 and member of set 2
4.
 - For all x , it's false, if member of set 1 then member of set 2
 - There exists an x , member of set 1 and not member of set 2

Are CS 5012 Class students good?

NOTE: $S = \text{Class of CS 5012}$ (set of students)
 $x = \text{is a student}$ – an **Element** of the **Set S**

(5) Every student is very good:

Right: $\forall x [S(x) \rightarrow \text{very good}(x)]$

Wrong: $\forall x [S(x) \wedge \text{very good}(x)]$

(6) Some students are excellent:

Right: $\exists x [S(x) \wedge \text{excellent}(x)]$

Wrong: $\exists x [S(x) \rightarrow \text{excellent}(x)]$

5. *For every x , if x belongs to S , then x is very good*
6. *There exists an x , where x belongs to S and x is excellent*

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