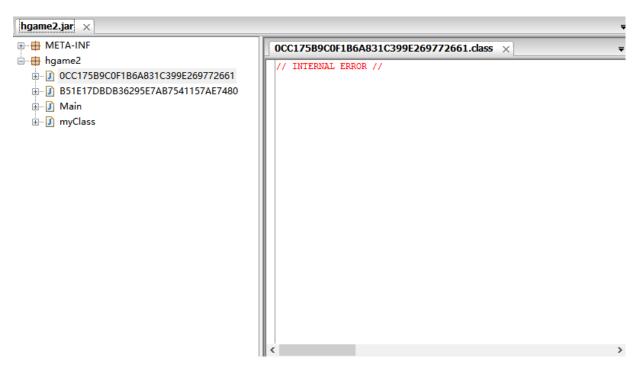
# Week4 WP - Veritas501

## RE

## explorer的奇怪番外4

做完以后我其实想说一句,这题真的不难,是我太蠢。

首先掏出jd-gui,看到这个jar由两个加密的class和两个没加密的class组成:



jar里完整的源码我就不放了,看到两段:

```
{
    myClass mc = new myClass();
    Class clazz = mc.loadClass("hgame2.checkFlag");
    Method c = clazz.getMethod("trueMain", (Class[])null);
    c.invoke(null, new Object[0]);
}
```

```
try {
    res = res.substring(0, l + 1) + md5(className.getBytes()) + ".class";
} catch (NoSuchAlgorithmException e) {
    e.printStackTrace();
}
```

得知,以上加密的两个class的名字是原来名字的md5值,由代码知一个为 checkFlag ,MD5破解得另一个为 a ,当然这不是重点。

分析myClass知, class中的内容使用AES加密的。

key的生成:

```
private static String code = "explorer";
...
MessageDigest md = null;
try {
    md = MessageDigest.getInstance("MD5");
} catch (NoSuchAlgorithmException e) {
    e.printStackTrace();
}
assert (md != null);
md.update(code.getBytes());
byte[] key = md.digest();
```

iv : String ivStr = "\*\*\*\*\*\*\*\*\*\*\*;

我们稍微修改一下代码,即可解密两个class:

```
package test;
import java.io.IOException;
import java.io.File;
import java.io.FileInputStream;
import java.io.FileOutputStream;
import java.io.InputStream;
import java.security.InvalidAlgorithmParameterException;
import java.security.InvalidKeyException;
import java.security.MessageDigest;
import java.security.NoSuchAlgorithmException;
import javax.crypto.BadPaddingException;
import javax.crypto.Cipher;
import javax.crypto.IllegalBlockSizeException;
import javax.crypto.NoSuchPaddingException;
import javax.crypto.spec.IvParameterSpec;
import javax.crypto.spec.SecretKeySpec;
    public static void main (String[] args) throws java.lang.Exception
        String code = "explorer";
        String ivStr = "***********;
        MessageDigest md = null;
            md = MessageDigest.getInstance("MD5");
        } catch (NoSuchAlgorithmException e) {
            e.printStackTrace();
        assert (md != null);
        md.update(code.getBytes());
        byte[] key = md.digest();
        File reader = new File("C:\\Users\\veritas501\\Desktop\\hgame2\\hgame
2\\0CC175B9C0F1B6A831C399E269772661.class");
        InputStream is = new FileInputStream(reader);
        long len = OL;
          len = is.available();
        } catch (IOException e) {
         e.printStackTrace();
        byte[] raw = new byte[(int)len];
          int r = 0;
          int off = 0;
            r = is.read(raw, off, (int)len);
            if (r == len) break;
            len -= r;
            off += r;
```

```
catch (IOException e)
         e.printStackTrace();
       Cipher cipher = null;
         cipher = Cipher.getInstance("AES/CBC/PKCS5Padding");
       } catch (NoSuchAlgorithmException|NoSuchPaddingException e) {
         e.printStackTrace();
       SecretKeySpec skey = new SecretKeySpec(key, "AES");
       IvParameterSpec iv = new IvParameterSpec(ivStr.getBytes());
         assert (cipher != null);
         cipher.init(2, skey, iv);
       } catch (InvalidKeyException|InvalidAlgorithmParameterException e) {
         e.printStackTrace();
       byte[] en = null;
         en = cipher.doFinal(raw);
       } catch (IllegalBlockSizeException|BadPaddingException e) {
         e.printStackTrace();
       File f=new File("C:\\Users\\veritas501\\Desktop\\hgame2\\tes
t.class");
       if (f.exists()==false) {
           f.createNewFile();//create file if not exist
       FileOutputStream fos=new FileOutputStream(f);
       fos.write(en);
       fos.flush();
       fos.close();
```

解密后用压缩软件在塞到jar里,再次用jd-qui打开,得到class的代码:

a.class:

```
package hgame2;
import java.util.Arrays;
public class a
  public static boolean check(String flag)
    byte[] var100 = flag.getBytes();
    byte[] var72 = new byte[var100.length + 2];
    System.arraycopy(var100, 0, var72, 0, var100.length);
    byte[] var140 = new byte[var72.length / 3 * 4];
    byte[] var82 = var140;
    int var17 = 0;
    int var18 = 0;
    while (var17 < var100.length) {</pre>
      var82[var18] = (byte)(var72[var17] >>> 2 & 0x3F);
      var82[(var18 + 1)] = (byte)(var72[(var17 + 1)] >>> 4 & 0xF | var72[var1
7] << 4 & 0x3F);
      var82[(var18 + 2)] = (byte)(var72[(var17 + 2)] >>> 6 & 0x3 | var72[(var1
7 + 1)] << 2 & 0x3F);
      var82[(var18 + 3)] = (byte)(var72[(var17 + 2)] & 0x3F);
      var17 += 3;
      var18 += 4;
    var17 = 0;
    while (var17 < var82.length) {</pre>
      int var10000 = var82[var17];
      if (var10000 < 26) {
        var82[var17] = (byte)(var82[var17] + 65);
        var10000 = var82[var17];
        byte var10001 = 52;
        if (var10000 < var10001) {</pre>
          var82[var17] = (byte)(var82[var17] + 97 - 26);
          var10000 = var82[var17];
          var10001 = 62;
          if (var10000 < var10001) {</pre>
            var82[var17] = (byte)(var82[var17] + 48 - 52);
            var10000 = var82[var17];
            var10001 = 63;
            if (var10000 < var10001) {</pre>
              var82[var17] = 43;
```

```
var140 = var82;
              int var136 = var17;
              var140[var136] = 47;
      var17++;
    int var10000 = var82.length;
    byte var10001 = 1;
    var17 = var10000 - var10001;
    while (var17 > var100.length * 4 / 3) {
      var82[var17] = 61;
      var17--;
    for (int i = 0; i < var82.length; i++) {</pre>
     var82[i] = (byte)(var82[i] ^ 0xCC);
    byte[] f = { -107, -74, -118, -92, -81, -1, -126, -127, -127, -117, -118,
-108, -123, -74, -81, -1, -98, -122, -82, -95, -81, -15 };
    return Arrays.equals(var82, f);
```

#### checkFlag.class:

```
package hgame2;
import java.io.BufferedReader;
import java.io.IOException;
import java.io.InputStreamReader;
import java.io.PrintStream;

public class checkFlag
{
    public static void trueMain()
        throws IOException
    {
        System.out.print("Now give me flag: ");
        BufferedReader strin = new BufferedReader(new InputStreamReader(System.in));
        String flag = strin.readLine();
        if (a.check(flag))
            System.out.println("hctf{" + flag + "}");
        else
            System.out.println("try again");
        }
    }
}
```

a.class是先对flag做base64编码,然后每一位和0xcc异或,python写出反代码:

```
import base64
c=[-107, -74, -118, -92, -81, -1, -126, -127, -127, -117, -118, -89, -106, -10
8, -122, -86, -127, -102, -126, -86, -127, -101, -7, -4, -106, -108, -123, -7
4, -81, -1, -98, -122, -82, -95, -81, -15 ]

for i in range(len(c)):
    if c[i]<0:
        c[i]+=256
    c[i] ^= 0xCC
    c[i] = chr(c[i])
c = ''.join(c)

print 'hctf{'+base64.b64decode(c)+'}'</pre>
```

得到flag: hctf{classL0ader\_1S\_1nter3stIng}

# explorer**的奇怪番外**7

工具用的jeb。安卓逆向,首先想到的是找找有没有字符串做参考,来到 values\strings.xml ,发现字段:

```
<string name="check_flag">
          check_flag</string>
<string name="enter_password">
          enter password</string>
```

挺好的,再来到 values\public.xml,找到字段:

```
<public id="0x7f04001a" name="check_flag" type="layout" />
...
<public id="0x7f060015" name="check_flag" type="string" />
...
<public id="0x7f0c0054" name="check_flag" type="id" />
```

最后通过type为"id"的public id (2131492948)找到了相关的代码:

```
protected void onCreate(Bundle arg3) {
        super.onCreate(arg3);
        this.setContentView(2130968602);
        this.editText = this.findViewById(2131492949);
        this.button = this.findViewById(2131492948);
        this.textView = this.findViewById(2131492950);
        this.button.setOnClickListener(new View$OnClickListener() {
            public void onClick(View arg11) {
                String v6 = checkFlag.this.editText.getText().toString();
                    MessageDigest v4 = MessageDigest.getInstance("MD5");
                    v4.update(v6.getBytes());
                    if(!Arrays.equals(v4.digest(), new byte[]{-73, 14, 42, 13,
-123, 91, 77, -57, -79, -22, 52, -88, -87, -47, 3, 5})) {
                    MessageDigest v7 = MessageDigest.getInstance("sha-256");
                    v7.update(v6.getBytes());
                    checkFlag.this.textView.setText("hctf{" + checkFlag.bytes2
Hex(v7.digest()) + "}");
                catch(NoSuchAlgorithmException v0) {
                    v0.printStackTrace();
        });
```

代码的意思就是获取输入,对输入进行md5加密,加密结果和已知数据比较,如果想等,则对输入进行sha256加密,结果加上'hctf{}'就是flag了。

得到flag\_md5: {-73, 14, 42, 13, -123, 91, 77, -57, -79, -22, 52, -88, -87, -47, 3, 5} 把signed转换成unsigned再hex:

```
#include <stdio.h>
#include <Windows.h>

int main(void)
{
    unsigned char b[] = { -73, 14, 42, 13, -123, 91, 77, -57, -79, -22, 52, -8
8, -87, -47, 3, 5 };
    for (int i = 0; i < 16; i++)
        {
            printf("%.2x", b[i]);
        }
        system("pause");
        return 0;
}//b70e2a0d855b4dc7blea34a8a9d10305</pre>
```

得到flag\_md5: b70e2a0d855b4dc7b1ea34a8a9d10305

md5在线解密: http://www.md5online.org/

结果为: Gabriel

sha256加密得: 0c030df5a4e7477d218012c0121ebce6d61bb8dc46e0a6c4f8e1cc8091b946a5

最后flag: hctf{0c030df5a4e7477d218012c0121ebce6d61bb8dc46e0a6c4f8e1cc8091b946a5}

#### coder

一开始,我是打算用正常的做re题的方法来解的,先跑跑,逆代码,分析,写反函数之类的来解,但我失败了,大概是水平不够。但我们依然有方法解题。

惯例扔到ida里,在main函数的加密函数中有如下一段:

```
printf("encrypt ok, your key is ", buf);
    for (k = 0; k \le 4; ++k)
     printf("%02x", *((_BYTE *)&pt_key1 + k));
    for ( l = 0; l <= 9; ++l )
     printf("%02x", *((_BYTE *)&pt_key2 + l));
    for ( m = 0; m < size; ++m )</pre>
      *((_BYTE *)&pt_key1 + m % 5) = sub_401766(*((_BYTE *)&pt_key1 + m % 5),
2u);
      *((_BYTE *)&pt_key2 + m % 10) = sub_401766(*((_BYTE *)&pt_key2 + m % 1
0), 4u);
      *((_BYTE *)buf + m) ^= *((_BYTE *)&pt_key1 + m % 5) ^ *((_BYTE *)&pt_key
2 + m % 10);// 将原始文件读入内存,在内存中加密,然后写出
    v8 = open(*(const char **)(v12 + 24), 0x41, 0x1B6LL);
    write(v8, buf, size);
    close(v8);
    putchar(10);
    free(buf);
```

由此我们知道, key1为5位, key2为10位, 加密过程可以大致表示为:

```
enc[i] ^= key3[i%10] ,其中key3为10位, key3[i] = key1[i%5] ^ key2[i]。
```

但我们现在不知道sub\_401766函数是做什么的,起初我怀疑这个函数会对key做某种变换,导致每一轮加密用的key都不一样(满足某种函数关系),我试着分析了一下,但关系实在不好找,以为要GG,但我们打开flag.mp4文件,看到文件尾部为:

```
ZM.EDEUH. OI OR TO DE DO OU DO RO RO IZ OI OR TO DE DO UU HIDUN
                                                           ~~ B4Šü>-.~~ B4Š
2A:EDF0h: 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 A0 42 34 8A
2A:EE00h: FC 3E 96 06 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 | U>-."" B4ŠU>-.""
                                                           B4Šü>-. " B4Šü>
2A:EE10h: A0 42 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E
2A:EE20h: 96 06 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 A0 42 -. " B4Šü>-. " B
2A:EE30h: 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E 96 06
                                                           4Šü>-.~" B4Šü>-.
2A:EE40h: 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 A0 42 34 8A "" B4Šü>-."" B4Š
                                                           ü>-. " B4Šü>-. "
2A:EE50h: FC 3E 96 06 98 A8 A0 42 34 8A FC BE 96 06 98 A8
2A:EE60h: A0 42 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E B4Šü>-. " B4Šü>
                                                          -. " B4Šü>-. " B
2A:EE70h: 96 06 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 A0 42
2A:EE80h: 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E 96 06 45u>-. " B45u>-.
2A:EE90h: 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 CF BD 34 8A
                                                           ~~ B4Šü>-.~~Ͼ4Š
2A:EEAOh: FC 3E 96 06 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 | u>-. " B4Šu>-. "
                                                           B4Šü>-.~" B4Šü>
2A:EEBOh: A0 42 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E
2A:EECOh: 96 06 98 A8 A0 42 34 8A FC 3E 96 06 98 A8 A0 42 -. " B4Šü>-. " B
2A:EEDOh: 34 8A FC 3E 96 06 98 A8 A0 42 34 8A FC 3E 96 06 45u>-.~ B45u>-.
```

非常整齐,十个一组,所以我的顾虑打消了,key3应该是不会改变的。

flag的格式为mp4,这个条件我们不能漏下,我随便打开了我硬盘上的几个mp4文件,发现前八个字节都是相同的,为

```
dec1 = [0x00, 0x00, 0x00, 0x20, 0x66, 0x74, 0x79, 0x70]:
```

```
0 1 2 3 4 5 6 7 8 9 A B C D E F 0123456789ABCDEF

0000h: 00 00 00 20 66 74 79 70 6D 70 34 32 00 00 00 00 ... ftypmp42...

0010h: 6D 70 34 32 6D 70 34 31 69 73 6F 6D 61 76 63 31 mp42mp41isomavc1

0020h: 00 01 A2 3B 6D 6F 6F 76 00 00 00 6C 6D 76 68 64 ..¢;moov...lmvhd
```

而加密后的文件头8个字节为: enc1 = [0xFC,0x3E,0x96,0x1E,0xFE,0xDC,0xD9,0x32]

我们将两者进行异或 , 得到 xor1 = [0xFC,0x3E,0x96,0x3E,0x98,0xA8,0xA0,0x42]

如果没有猜错,文件的结尾出应该是一串相同的字符,就像我手头的这个文件一样:

```
Startup 🔞 flag.mp4 🔞 五月天 貫徹快樂 — MAYDAY Don'ts Don'ts on Vimeo.MP4 🔞 u_e 🗟 🕝 🕦
▼ | Edit As: Hex ▼ | Run Script ▼ | Run Template ▼
    2 3 4 5 6 7
           9 A B C D E F
                 0123456789ABCDEF
.....Ȑ")iii
6F3:F200h: B4 B4 B4 B4 B4 B4 B4 B4 BB EA 84 29 69 69 69
11111111111111111
11111111111111111
```

那么我们用python脚本:

```
file_end = [0x3e,0x96,0x06,0x98,0xa8,0xa0,0x42,0x34,0x8a,0xfc,0x3e,0x96,0x06,0
    x98,0xa8,0xa0,0x42,0x34,0x8a,0xfc]
    xor1 = [0xFC,0x3E,0x96,0x3E,0x98,0xA8,0xA0,0x42]

for i in range(10):
    print 'i=',i
    for j in range(8):
        print xor1[j]^file_end[i+j]
    print '===='
```

观察输出有如下一段:

```
====
i= 9
0
0
0
0
56
0
0
0
0
0
0
====
```

这就验证了我们的猜想,只是flag.mp4的文件头的第4字节和我手头的MP4文件不同。根据file\_end,从 而我们得到了真正的 key3 = [0xFC,0x3E,0x96,0x06,0x98,0xA8,0xA0,0x42,0x34,0x8a]

我们利用key3对flag.mp4进行解密:

```
key3 = [0xFC,0x3E,0x96,0x06,0x98,0xA8,0xA0,0x42,0x34,0x8a]
fp = open(r'D:\flag.mp4','rb')
stream = list(fp.read())
fp.close()
out = []
for ch in stream:
    out.append(chr(ord(ch)^key3[i%10]))
out = ''.join(out)
fp = open(r'D:\flag_dec.mp4','wb')
fp.write(out)
fp.close()
```

#### 打开解密后的视频:



## easy-shell

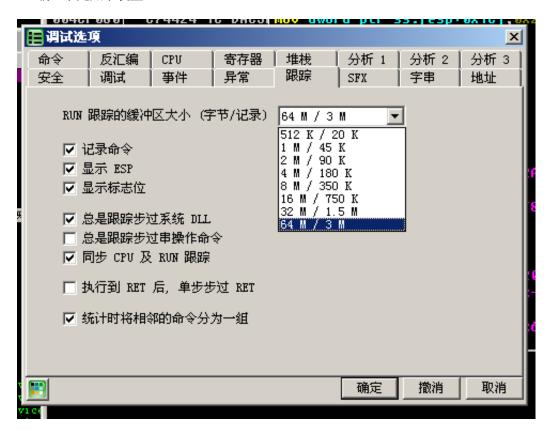
首先扫下壳:



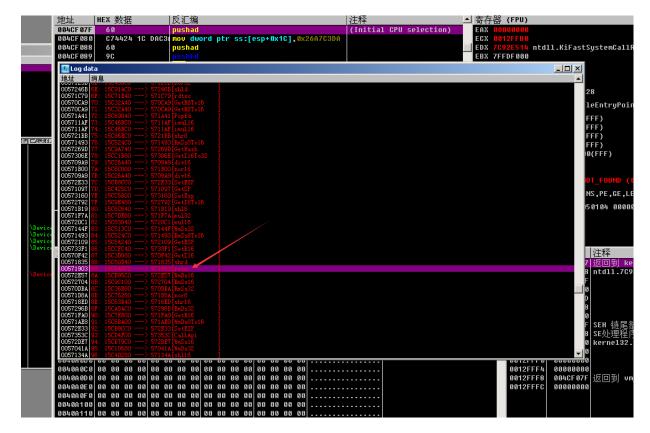
正如文件的名字,是个vmp壳,说真的我挺害怕这个壳的。

前排提醒:此处调试用的OD不能是原版OD,你可以使用网上各大论坛改的OD,比如52pojie,学破解,飘云阁等等。调试时需要Strong OD, FKVMP,忽略异常等。(你应该需要在XP上调试)

OD载入, 先如下设置:



在此时的代码处右键选择FKVMP>>start,点击OD上方的L按钮,找到retn:



记录retn的地址: 0x00571903

接着在OD上下断点: bp VirtualProtect ,按f9运行,观察堆栈区:

```
地址
         |数值
                    |注释
0012F6C4
           00572FE8
                     rCALL 到
                     Address = vm_vmp.00401000
0012F6C8
           00401000
0012F6CC
           0000830A
                      Size = 830A (33546.)
                    NewProtect = PAGE_EXECUTE_READWRITE
LpOldProtect = 0012FF98
0012F6D0
           00000040
0012F6D4
           0012FF98
0012F6D8
           7C801AD4 kernel32.VirtualProtect
0012F6DC
           00000170
0012F6E0
           00000000
0012F6E4
           00000206
0012F6E8
           00000202
0012F6EC
           0012FF98 UNICODE "X"
0012F6F0
          58531C00
```

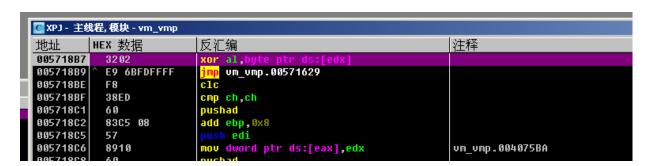
直到NewProtect = READONLY:

```
212 EMPCA 0.0
   ST4 empty 0.0
地址
         数值
                   注释
0012F6C4
          00572FE8 CALL 到
0012F6C8
          0040A000
                    Address = vm_vmp.0040A000
0012F6CC
          00002386
                    Size = 2386 (9094.)
0012F6D0
          00000002
                    NewProtect = PAGE READONLY
0012F6D4
          0012FF98
                   Lp01dProtect = 0012FF98
0012F6D8
          00000206
0012F6DC
          00000202
0012F6E0
          00000001
0012F6E4
          00000000
0012F6E8
          00187168
0012F6EC
          689F232A
```

此时alt+B, 断点界面取消或禁用断点, 然后alt+M, 对text段下内存访问断点:



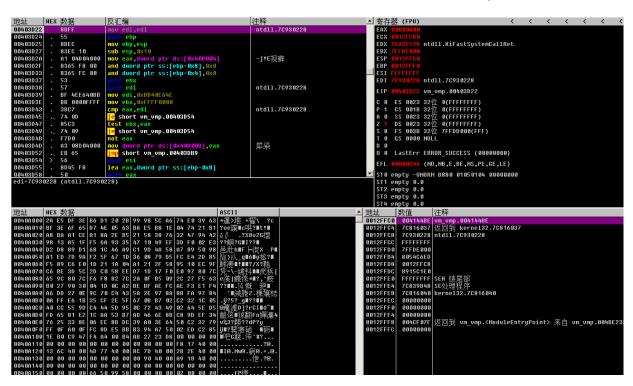
f9一下,取消text段的访问断点,来到这里:



掏出我们之前记录下的retn地址:ctrl+G转到然后f2下断:

```
☑ XPJ - 主线程, 模块 - vm_vmp
         HEX 数据
                                                                      注释
地址
                           反汇编
            F6C4 C9
005718E0
                            test ah,0xC9
test sp,0xBB35
005718E3
            66:F7C4 35BB
                                vm_vmp.00570863
005718E8
           E9 76EFFFFF
            66:0FBCC3
                                ax,bx
                                                                       shr16
005718F1
           66:8B45 00
                            mov ax,word ptr ss:[ebp]
005718F5
           F6D9
                            neg cl
005718F7
           FEC9
                            dec
005718F9
                            add cl,bh
            aafg
005718FB
            8A4D 02
                            mov cl,byte ptr ss:[ebp+0x2]
                                vm_vmp.00570542
005718FE
           E8 3FECFFFF
            F7C1 A21663EE
                                                                       retn
00571909
           89EC
                            mov esp,ebp
                            clc
0057190B
           F8
0057190C
           80EA 29
                            sub d1,0x29
0057190F
                            cld
           F.C
00571910
           58
                            pop eax
00571911
           66:29D2
                            sub dx,dx
00571914
            5A
                            pop edx
                                                                       vm vmp.004075BA
00571915
            80E9 A8
                                cl,0xA8
```

f9一下,取消retn的断点,再对text段下内存访问断点,f9一下,来到了我们的oep:



(此时我们可以对oep下硬件执行断点,方便下次调试,不用再重复之前那些动作)

我们现在可以用lordpe dump一下镜像,虽然IAT没有修复,但是IDA还是能分析部分的。

ida载入,找到关键函数:

```
int sub_401000()
 int v0; // eax@1
 int v1; // esi@1
 signed int v2; // eax@2
 unsigned int v3; // eax@4
 char *v4; // ecx@4
 char *v5; // edx@4
 __int16 v7; // [sp+8h] [bp-50h]@1
 char v8; // [sp+Ah] [bp-4Eh]@1
 char v9; // [sp+Ch] [bp-4Ch]@1
 __int16 v10; // [sp+24h] [bp-34h]@1
 char v11[28]; // [sp+28h] [bp-30h]@2
 int v12; // [sp+44h] [bp-14h]@1
 __int16 v13; // [sp+48h] [bp-10h]@1
 int v14; // [sp+4Ch] [bp-Ch]@1
 char v15; // [sp+50h] [bp-8h]@1
 v8 = 0;
 qmemcpy(&v9, word_40B90C, 24u);
 LOBYTE(v13) = 0;
 v15 = 0;
 v12 = 0x4030201;
 LOBYTE(v12) = 'f';
 BYTE1(v12) ^= 'n';
 HIWORD(v12) = 'ga';
 v14 = 0x4030201;
 v10 = word_40B90C[12];
 v13 = 0;
 v0 = sub_401311((int)&v12, (int)&unk_40B928); // 'r'
 v1 = v0;
 if ( v0 )
   sub_40111D(v11, 26, v0);
   sub_4014E4(v1);
   LOBYTE(v7) = v7 ^ 'f';
   HIBYTE(v7) ^= 'e';
   LOBYTE(v14) = v14 ^{\prime} 'b';
   BYTE1(v14) ^= 'm';
   BYTE2(v14) ^= 'l';
   BYTE3(v14) ^= 'h';
   v2 = 0;
     v11[v2] = (v11[v2] - 3) ^ '3';
     ++v2;
   while ( v2 < 25 );
   v3 = 25;
   v4 = v11;
   v5 = &v9;
   while (*(_DWORD *)v5 == *(_DWORD *)v4)
     v3 -= 4;
     v4 += 4;
```

```
v5 += 4;
if ( v3 < 4 )
{
    if ( *v4 != *v5 )
        break;
    sub_401328(&v14);
    return 0;
    }
}
sub_401328(&v7);
return 0;
}</pre>
```

忽略中间的n多细节, 我们只看两段:

结合OD动态调试, 我们发现, 函数是将写在0x0012ff48处的数据做 byte = (byte - 3) ^ '3' 变换:

```
004010A9 | .
             8DA424 00000 lea esp,dword ptr ss:[esp]
004010B0
             8A4C 05
                   D Ø
                          mov cl,byte ptr ss:[ebp+eax-0x30]
             80E9 03
                          sub cl,0x3
004010B4
004010B7
             80F1 33
                           xor cl,0x33
                           mov byte ptr ss:[ebp+eax-0x30],cl
004010BA
             884C05 D0
004010BE
             40
                          inc eax
004010BF
             83F8 19
                           cmp eax,0x19
004010C2
             7C EC
                             short vm_vmp.004010B0
00401004
             B8 19000000
                          mov eax,0x19
             8D4D D0
00401009
                          lea ecx,[local.12]
c1 = 06
跳转来自 004010C2
         HEX 数据
批批
                                                        ASCII
0012FF28 67 67 00 00 56 53 42 50 4B 5F 56 6F 65 7F 61 6F
                                                        qq..VSBPK Voe∎ao
0012FF38 0D 7C 71 6F 63 7F 6F 63 79 0D 7C 62 49 00 40 00 .|qoc∎ocy.|bI.@.
0012FF48 34 FF 12
                  00 2C A1 40 00 B0 FF 12
                                          00 C0 1F
                                                  40
                                                     00 4ij∎., .?∎.?@.
0012FF58 BC D3 0C
                    FE FF FF
                             FF
                                79 6C 40 00 66 6C 61 67 加.w?jjjy1@.flag
                                                        ..@.cool./@.膽
0012FF68 00 00 40 00 63 6F 6F
                             6C 00 2F 40 00 C4 91 5E 77
                                                  BA 00 ?■.?@.■... /?
0012FF78 C0 FF 12
                  00 8C 16 40
                             00 01 00
                                      00
                                         00 20 2F
0012FF88 68 2F BA
                    7C 91 5E
                             77 28 02 93
                                         7C FF FF FF
                                                     FF
                 00
                                                        h/?|慯w(■搢ÿÿÿÿ
0012FF98 00 E0 FD 7F
                     06 02 00 00 D9 18 41 00 00 00 00 00
                                                        .帻■■■..?A....
0012FFA8 8C FF 12
                 00 14 E5 92 7C E0 FF 12 00 C0 1F 40 00 ?■.■鰕|?■.?@.
0012FFB8 3C D4 0C
                    00 00 00
                             00 F0 FF 12
                                         00 37 60
                                                  81 7C <?w....?■.7` ]
                 77
                                         7F ED C6 54 80 (■搢ÿÿÿÿ.帳■砥T■
0012FFC8 28 02 93
                 7C FF FF FF
                             FF 00 E0 FD
                                                        ?■.嗔■?ijijH泝[
                             89 FF FF FF
                                         FF 48 9B 83 7C
0012FFD8 C8 FF 12
                 00 E0 C1 15
                                                        0012FFE8 40 60 81
                 7C 00 00 00
                             00 00 00 00 00 00 00 00 00
0012FFF8
         7F F0 4C 00
                     00 00 00 00
                                                        ■餖.....
```

然后和写在0x0012FF2C处的数据一位一位比较,如果相同则输出'cool':

```
004010CF
             90
                           nop
004010D0
             8B32
                           -mov esi,dword pt
             3B31
                            cmp esi,dword ptr ds:[ecx]
004010D2
004010D4
             75 1A
                               short vm_vmp.004010F0
                            sub eax,0x4
004010D6
             83E8 04
004010D9
             8301 04
                            add ecx,0x4
004010DC
             83C2 04
                            add edx,0x4
004010DF
             83F8 04
                            cmp eax,0x4
004010E2
             73 EC
                               short vm vmp.004010D0
                           mov al,byte ptr ds:[ecx]
004010E4
             8A 01
004010E6
             3A 02
                           cmp al,byte ptr ds:[edx]
0.01.01.00
                               chart um uma AALA1AEA
堆栈 ds:[0012FF2C]=5<u>0425356</u>
esi=0040D070 (vm vmp.0040D070)
跳转来自 004010E2
地址
         HEX 数据
                                                           ASCII
0012FF2C
         56 53 42
                  50 4B 5F 56 6F
                                                           USBPK_Voe∎ao.|qo
                                  65 7F 61
                                           6F
                                               0D 7C 71
                                                        6F
0012FF3C 63 7F 6F
                  63
                     79
                        0D 7C
                               62
                                  49 00 40
                                           99
                                              02 CF 3C CE clocy.|bI.@.■?■
                                                           ■?螢?螏/■螉?G
0012FF4C 1A AD 0E CE 9E CF 3C CE 8E 2F 0E CE
                                              8A E3 3A 47
                                                       00 认痃Eí@.fíag..@.
0012FF5C C8 CF CF
                  CF 45 6C 40
                               99
                                  66 6C 61
                                           67
                                              00 00 40
0012FF6C 63 6F 6F
                                                        00 cool./@.膽^w?■.
                  6C
                     00 2F 40
                               00
                                  C4 91 5E
                                           77
                                              CO FF 12
0012FF7C 8C 16 40
                  00
                     01 00 00
                                  20 2F BA
                                              68 2F BA
                                                        99
                                                           ?@.■... /?h/?
                               00
                                           99
                                                           |慯w(■搢ÿÿÿÿ-帱■
                                           FF
0012FF8C
         7C 91 5E
                  77
                     28
                        02 93
                               7C
                                  FF FF FF
                                              00 E0 FD
                                                        7F
0012FF9C 06
            02 00 00 D9
                         18
                                           ១១ ខេត
                                                 FF 12
                            Д1
                                  aa
                                     00
                                        aa
```

由此用py写出反函数:

```
enc = [0x56,0x53,0x42,0x50,0x4B,0x5F,0x56,0x6F,0x65,0x7F,0x61,0x6F,0x0D,0x7C,0
x71,0x6F,0x63,0x7F,0x6F,0x63,0x79,0x0D,0x7C,0x62,0x49]
dec=[]
for i in range(len(enc)):
    dec.append(chr(((enc[i])^0x33)+3))
print ''.join(dec)

#hctf{oh_YOU_ARE_SO_SMART}
```

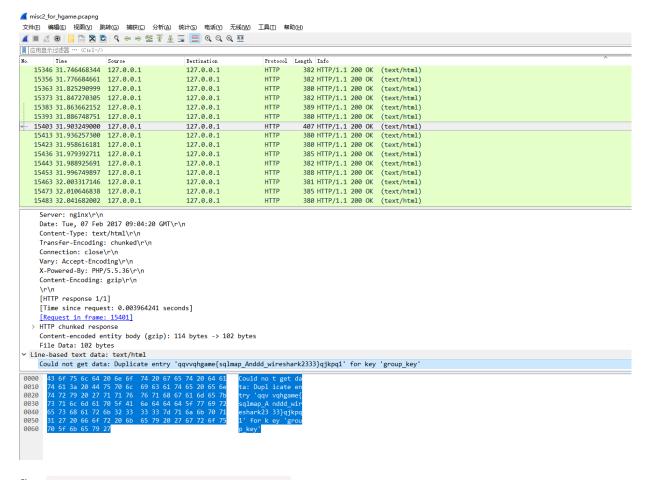
得flag: hctf{oh\_YOU\_ARE\_SO\_SMART}

后话:如果你知道vs编译出来的程序的用户函数一般都从data段最前面开始,而且你有成功猜到这道题的主函数在0x401000处的话,这道题会简单很多:首先下一个退出断点:BP ExitProcess,运行后断下,ctrl+G来到401000,发现代码完整,接着在数据窗口ctrl+G来到401000,对401000下硬件执行断点,重新载入,f9运行,程序成功断在401000处,然后和上面一样,很快就解得了flag。

### Misc

### 来看看自己是怎么日自己的

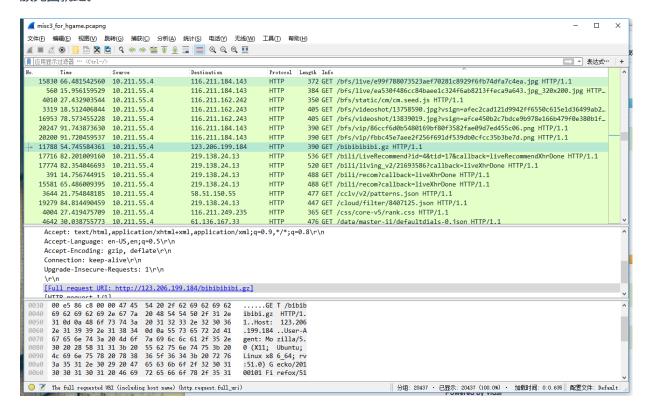
放完图就跑。



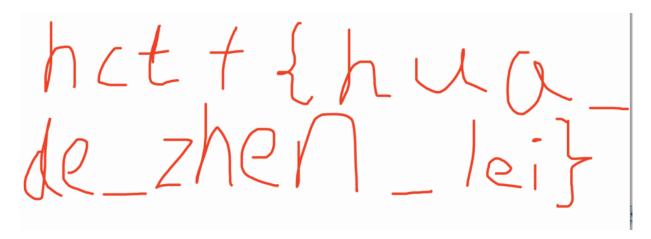
flag: hgame{sqlmap\_Anddd\_wireshark2333}

### 考眼力喽

放完图就跑。



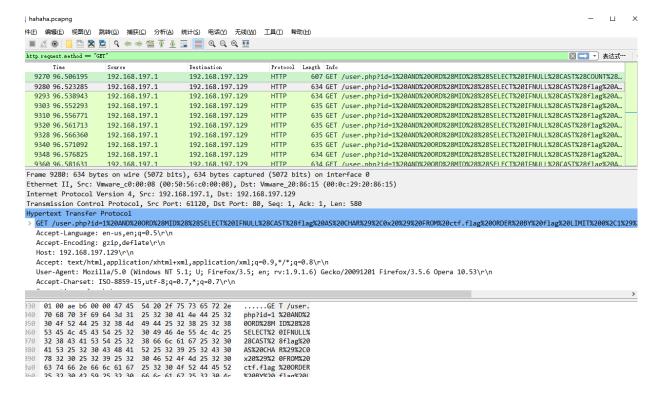
下载压缩包,解压得到一张图片:



flag : hctf{hua\_de\_zhen\_lei}

### 正在前往翻车大道

首先过滤"GET":



发现这些包中含有flag字段,成功引起了我的注意,url解码一下,形式如

下: http://192.168.197.129/user.php?id=1 AND ORD(MID((SELECT IFNULL(CAST(flag AS CHAR),0x20) FROM ctf.flag ORDER BY flag LIMIT 0,1),1,1))>96,一条条向下观察,发现最后一位是ascii码,大概用的是二分法。因为是flag,那八成是hctf开头,看了一下数据包,是 gcte,第一位和第四位要加一。又因为是大于号,所以最终的字符应该是ord(i)或ord(i+1)。先一位位提出来再说,得到 gcte{ekoweq\_sqk\_imiecsiom},适当加一后拼出连贯单词,即flag: hctf{flower\_sql\_injection}

# Crypto

# 进击的 Crypto [4]

题目不长, 先粘上:

```
import hashlib
from gmpy2 import mpz, invert
p = 19071602744764679284418973368135837932734306119936934138447490255730976968
292325467422028404595111739386865163837954151118521242161165396546010097988772
558347603718936959999966832729952954249400929004999205715774700715299402518659
1269943960410319186487203043168774653327128061548663247131284489765017
q = 930788704028200015275140127068138499329817310955
g = 22023715606272465708641346383195448856570387583936863915316264582425973547
564841743681268927004583530561320186941832818910141784591238694469574773757041
506397939323959002218076757573665282491833676469401442167351817806524451936760
21245167165713347633549737<u>17129155502536345229391073998560517516716958</u>
k = !???!
x = '???'
def data_to_int(s):
   return int(s.encode('hex'), 16)
def SHA1(data):
    return data_to_int(hashlib.sha1(data).hexdigest())
def encrypt(data, p, q, g, x, k):
    r = pow(g, k, p) % q
    s = (invert(k, q) * (SHA1(data) + x * r)) % q
   return (r, s)
data1 = "guest"
data2 = "admin"
(r1, s1) = encrypt(data1, p, q, g, x, k)
(r2, s2) = encrypt(data2, p, q, g, x, k)
print SHA1(data1)
print SHA1(data2)
print s1
print s2
print r1
print r2
413685540380226864
835940898148680488372488685713345793755099380413493862399556052721366535745667
186387858109315383
618159893787048300752592802884467155388759696698
659836539307844663175437862395252943516139307036
568752653628483014849549142909331362115254788206
568752653628483014849549142909331362115254788206
def getflag(data):
    if data == "getflag":
       (r, s) = encrypt(data, p, q, g, x, k)
       flag = "hctf{" + str(s % r) + "}"
       print flag
```

这里我做一个小小的改动,把gmpy2库改成gmpy,我的linux不知怎的就是装不上gmpy2,就拿gmpy将就一下,反正也能做题,如果有大佬知道怎么装请务必告诉我谢谢。

顺便说一下,我查资料的时候看到一种叫DSA的加密算法和他好像,虽然不知道这是不是DSA。

前排提醒:因本人还未学过离散,以下内容可能有误!!

先观察加密函数:

 $r = (g^k \mod p) \mod q$  $s = (invert(k, q) * (SHA1(data) + x * r)) \mod q$ 其中invert函数是求逆元。

题中给了两组加密,说明只用一组来暴力跑肯定是不科学的。如果从r入手肯定是暴力了,所以我从s入手。

两组的s相减可以发现:

$$s2 - s1 = (invert(k, q) * (sha2 - sha1)) \mod q$$

消掉了其中的x\*r。我们记s = s2 - s1, sha = sha2 - sha1。

得到:  $s = (y * sha) \mod q$ 

其中:  $k * y \mod q = 1$ 

因此我们两边乘k, 化简得到:  $k = (sha \div s) \mod q$ 

在python中利用gmpy的 divm(sha,s,q) 解出k。

对原始的s式子两边乘k,除r得到:

$$\frac{s1*k}{r} = (\frac{sha1}{r} + x) \mod q$$

从而化简得到:  $x = ((\frac{s1*k}{r} \mod q) - (\frac{sha1}{r} \mod q)) \mod q$ 

在python中利用gmpy的 (divm(s1\*k,r,q) - divm(sha1,r,q)) % q 解出k。

以下为完整的python解密代码:

```
import hashlib
```

from gmpy import mpz, invert,divm

 $\begin{array}{lll} \textbf{p} &=& 19071602744764679284418973368135837932734306119936934138447490255730976968\\ 292325467422028404595111739386865163837954151118521242161165396546010097988772\\ 558347603718936959999966832729952954249400929004999205715774700715299402518659\\ 1269943960410319186487203043168774653327128061548663247131284489765017 \end{array}$ 

q = 930788704028200015275140127068138499329817310955

sha1 = 42726297627322808322187199831313194501002956120959170626211891393748957
7133576413685540380226864

sha2 = 83594089814868048837248868571334579375509938041349386239955605272136653
5745667186387858109315383

s1 = 618159893787048300752592802884467155388759696698

s2 = 659836539307844663175437862395252943516139307036

r = 568752653628483014849549142909331362115254788206

```
s = s2 - s1
sha = sha2 - sha1
k = divm(sha, s, q)
x = (divm(s1*k,r,q) - divm(shal,r,q)) % q
def data_to_int(s):
    return int(s.encode('hex'), 16)
def SHA1(data):
    return data_to_int(hashlib.sha1(data).hexdigest())
def encrypt(data, p, q, g, x, k):
    r = pow(g, k, p) % q
    s = (invert(k, q) * (SHA1(data) + x * r)) % q
    return (r, s)
def getflag(data):
    if data == "getflag":
        (r, s) = encrypt(data, p, q, g, x, k)
        flag = "hctf{" + str(s % r) + "}"
        print flag
getflag("getflag")
```

解得flag: hctf{88169191231439818447681393510021281730269252095}

# 进击的 Crypto [5]

这道题给个十组RSA(别说你看不出来), e=10, 如果e比较大的话应该是用同素因子分解n的方法做,但这题e很小,且给了10组,所以应该是低加密指数广播攻击了。

破解这种用中国剩余定理,不知道的百度查一下。

以下为py脚本:

```
import gmpy

def my_parse_number(number):
    string = "%x" % number
    erg = []
    while string != '':
        erg = erg + [chr(int(string[:2], 16))]
        string = string[2:]
    return ''.join(erg)
```

```
def e_gcd(a, b):
    x,y = 0, 1
    lastx, lasty = 1, 0
    while b:
        a, (q, b) = b, divmod(a,b)
        x, lastx = lastx-q*x, x
        y, lasty = lasty-q*y, y
    return (lastx, lasty, a)
def chinese_remainder_theorem(items):
  N = 1
  for a, n in items:
    N *= n
  result = 0
  for a, n in items:
    m = N/n
    r, s, d = e_gcd(n, m)
    if d != 1:
      raise "Input not pairwise co-prime"
    result += a*s*m
  return result % N, N
e=10
```

409329951990653329877191831971994483364358510155402721554414867460273151078213 409691056234515758725801709783686505938802740760255204539563785397662284304291 421179946169853189631257390768266354592151149461593486783982680993895254144316 185177438893140857143904309727832232709856889889952577703631172254725900844891 209583971891309584627845326374534821827635057917206414226867496162165217240121 08378680177131928438795893866425040489551258902610047574579899692834984999209528110849342042529825594061894238371591910354584305136154355365191215400149,218 404372844226015841776018573558452964203001577673391095723776404083627266745612 462104004007604741871212468937124801093268936141627794707684152829007138009838 155706022500686736950447499325326898513107700875525246208576234420194280444827 879654280759849519821049209015624564266721115909091007362091539058530351272477 202769740504046638475470903627807290778295946849959794573264149474490832194616 173026433929045227338129031413013322273576895178808018455105795188874729280837 188268903690169445491604614287106157470012735212536387662671433418191702496483 85341127263707417642925464809788975715332938666417316695941347477577,300159141 339867581331050150829224609104717268190004798728168128067941408872092943939630 632733778912030698647114667762001081736724287792933083204601164930405728269150 206542939292418436867282963874001100620990745730096455635208053018999623446808 213962002560424785395406758504273779585533245813713621126527524448249196017678216220384883655010316365582361260526542416627430706512133695668989988986344412 618116550537893845887115476859399277211164044299143290814329909133212969649572 606911487047105816817745629049240808132749848099400089686295153048595186294732 84018332651831587916402539386242421630250537030688162675751761,180097184358254 456493726296348677722470351382294931083627136309476803383542277355720708823903 786324403651631170942334135201070705819082242392109691720944027609240553345842 596782765059194181916237629982497242483793027866216890491075007603483039401128 406649262313456463259641332815507654547196281616004751437925673099473301307478 601255575470518436208996292176369180029298464630978415394460140195798303473560840804885619174403568484658129372468090181685826351916196220410790124505318112

 $\begin{array}{l} \mathbf{c} = [467518260554971134765329951477782623855904020052774781084690199143312702572\\ 336180721167220805286287025804728528594921251566427841049960382966326621396851\\ 149397476667407182522553666526704904461934037843853161598369640630584194575639\\ 657937179482626554992950304313466836005502727652236887919362093745788148650931\\ 171190572248318773465012779651080863789018995284097133076891394688452559462765\\ 245874593053221923825891160874411200179813639480600472647289083848188102671660 \end{array}$ 

```
574381267435095525800379459074825852392946044698602079444547807655963514518190
9093786617716473870604463315696183050495143502762614767485881979217577416,4586
033930015814975553487614321341290358072960539564849589758593801587823414394351
862074093343255479635303199078561905260793645869942622136164615776822867275172
888812989071016687658023005838776181437470204528335444413852561895603377955959
308318399246963213964044984235974352824060625970587119586612774322385216810264
819426458641852166009899106655036755913322283219739741020656812034348404092018
340277026844114272699170323014050087289715403989671918767712613414346320222745
328800096615947138667139053339483278319521201983885423491905171167675501563556
563890821756892474999549635650486944010346554641603343238490572689,12012575342
366442210994368605032582129674485327006093552902983877957202172783938071684841
031902396156899396879136388606996807608296316795628987877357547213450360953742
947781700549506469226564975927361748596520878542651668963131021269851928174681
007347771519569914690980744941965086762285630082781013781879511681353119117280
354102271439057204066405072328422057721895983540492150402309552609035008272366
176657384174081229798739099558539083074489500359062010164358300087124447037922
844709870430218019673914305742939927888838599616520300011913022792060392260618
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```
data=[]
for i in range(len(c):
    data += [(c[i],n[i])]
x, n = chinese_remainder_theorem(data)
realnum = gmpy.mpz(x).root(e)[0].digits()
print my_parse_number(int(realnum))
#When e are small and same,it can be Hastad's broadcast attack.Maybe we won't
have topic aboout RSA,but I wish you can explore it Non-stop.hctf{Hastad's_broadcast_attack_is_interesting}
```

解得flag: hctf{Hastad's\_broadcast\_attack\_is\_interesting}