

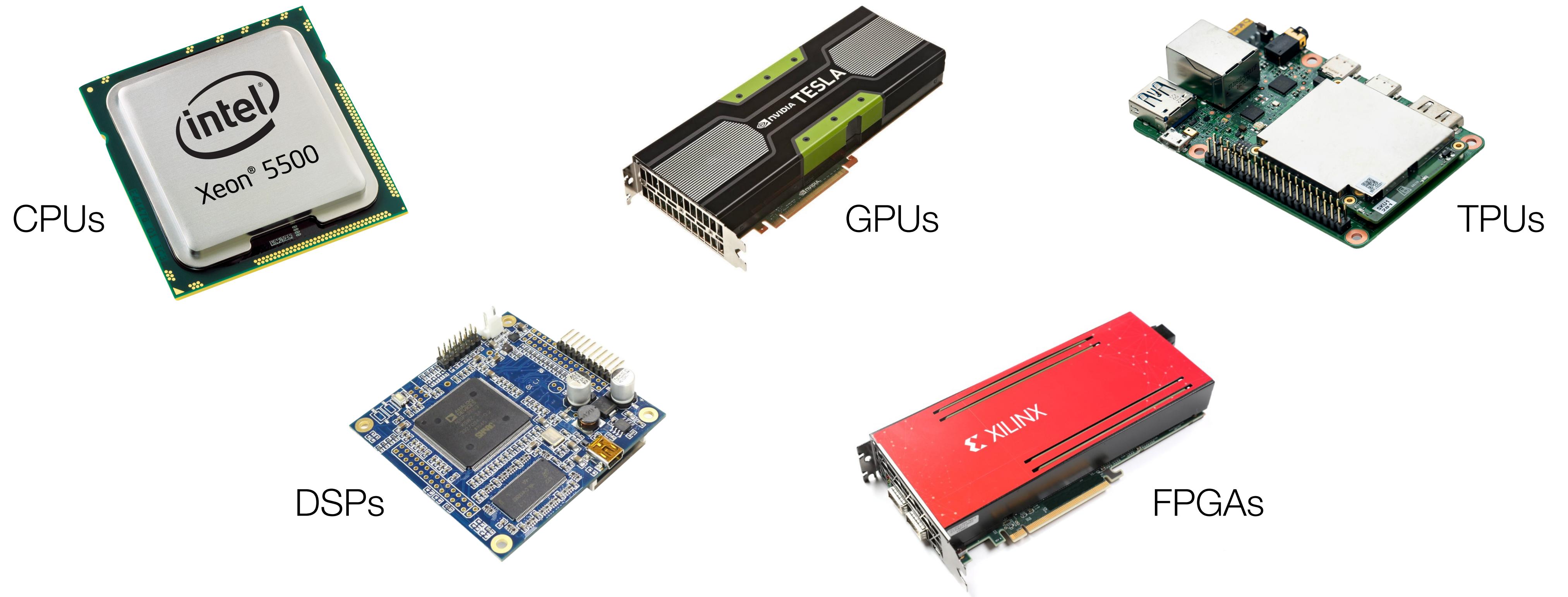
# Query Processing on Heterogeneous Systems

Viktor Rosenfeld  
8 December 2023



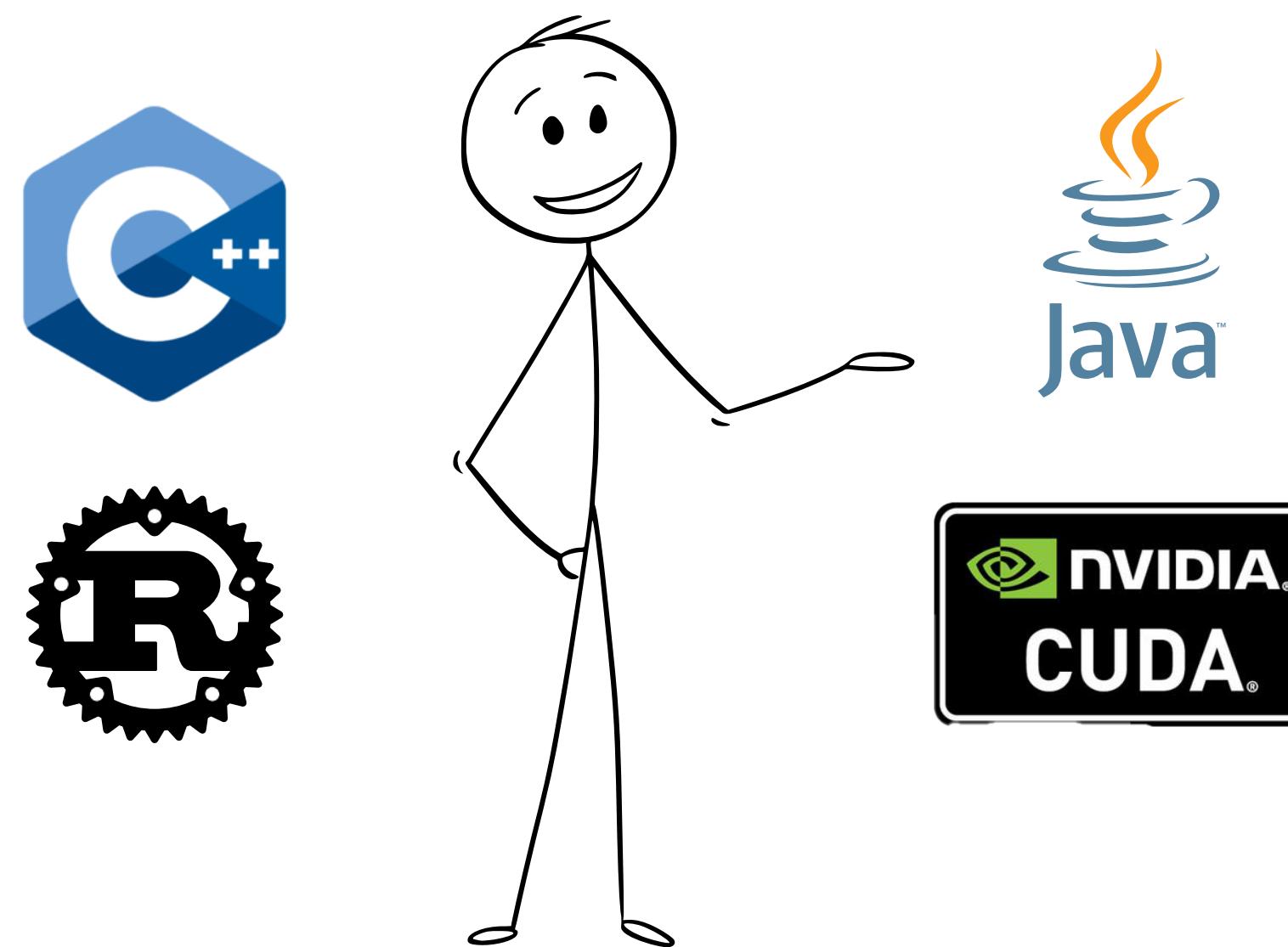
Today's computing systems are highly heterogeneous.

# Specialized processors



Processors are optimized for different application scenarios to overcome the power wall.

# Diverse users requirements



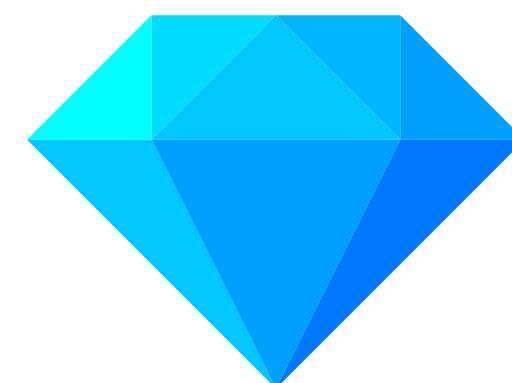
Professional programmers



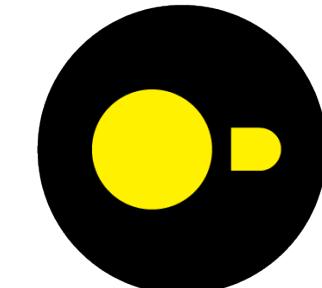
Casual programmers

Programming languages have to meet users' skill set and performance requirements.

# Complex data processing pipelines



IBM Db2



DuckDB



Complex software ecosystems integrate many specialized tools to solve data processing tasks.

# Thesis statement and goals

**Heterogeneity is a benefit.**

It enables the tools that allow different users to process large data sets efficiently.

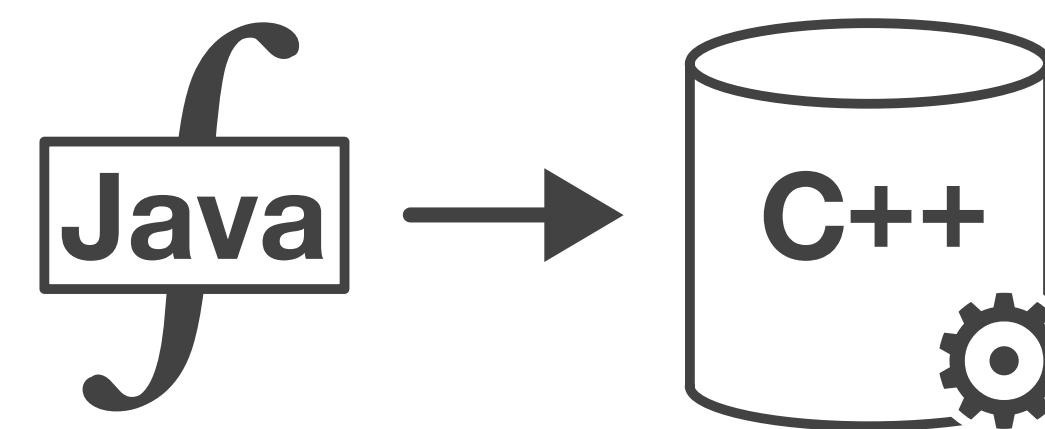
**Heterogeneity is a challenge.**

It increases the complexity of the computing infrastructure.

**Investigate how heterogeneous hardware and software impact query processing systems.**

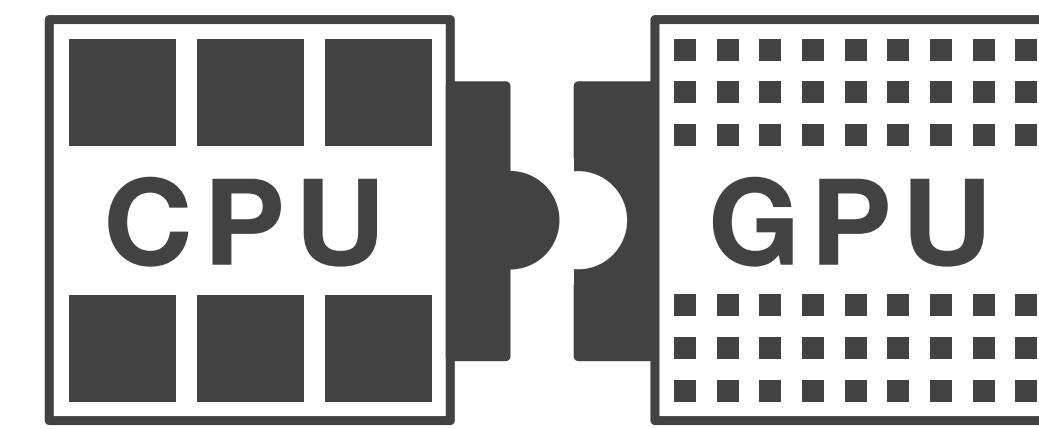
**Optimize query processing performance on heterogeneous hardware and software systems.**

# Three scenarios



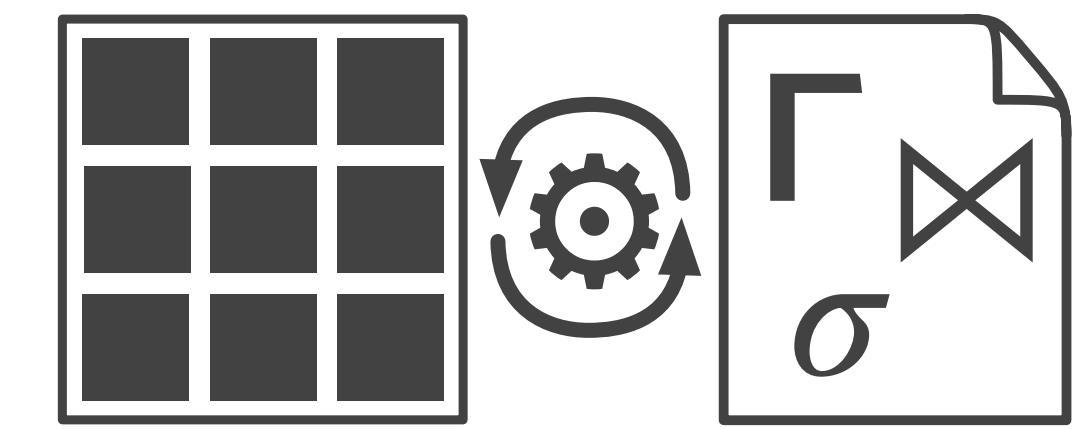
Processing Java UDFs  
in a C++ Environment

ACM SoCC 2019



Query Processing on  
Heterogeneous CPU/GPU Systems

ACM CSUR 2022

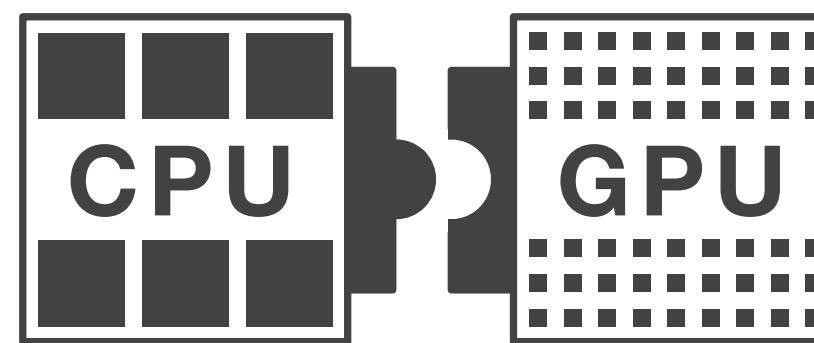


Operator Variant Tuning on  
Heterogeneous Processors

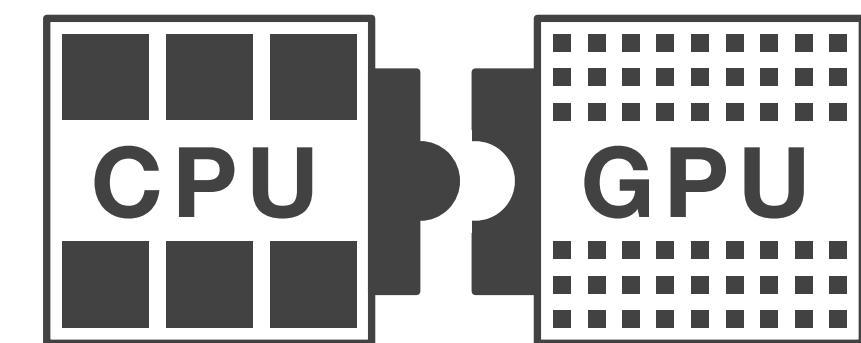
ADMS 2015, DaMoN 2019

# Common challenges posed by heterogeneous hardware and software

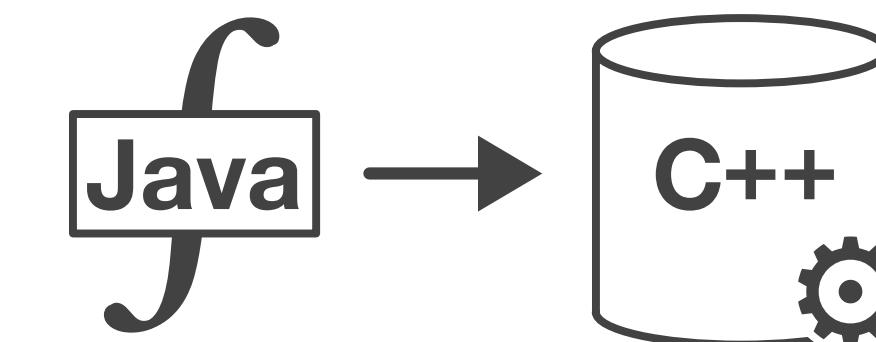
Distributing computation



Reducing data movement

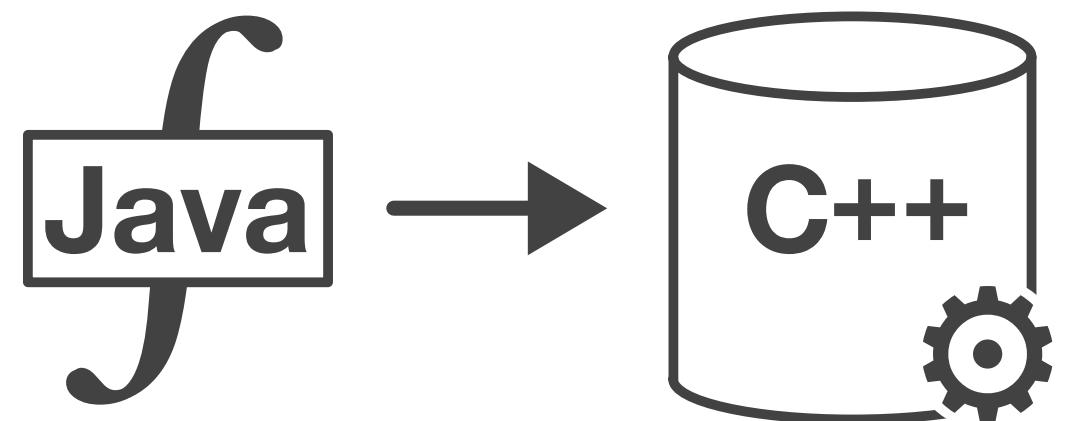


Cross-platform data processing



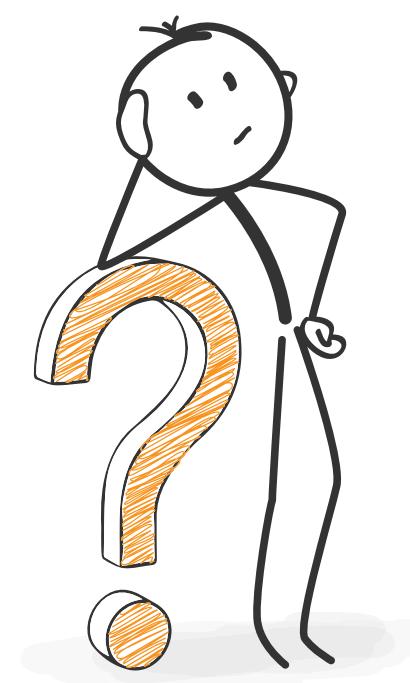
Cross-platform data processing

Techniques to address the heterogeneity of today's computing systems are applicable to a wide range of research areas and engineering tasks in query processing.

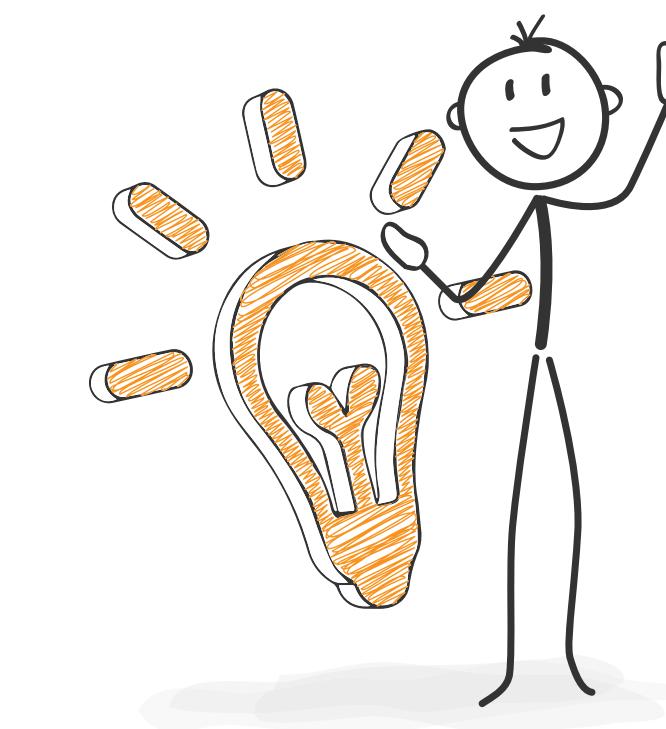


# Processing Java UDFs in a C++ Environment

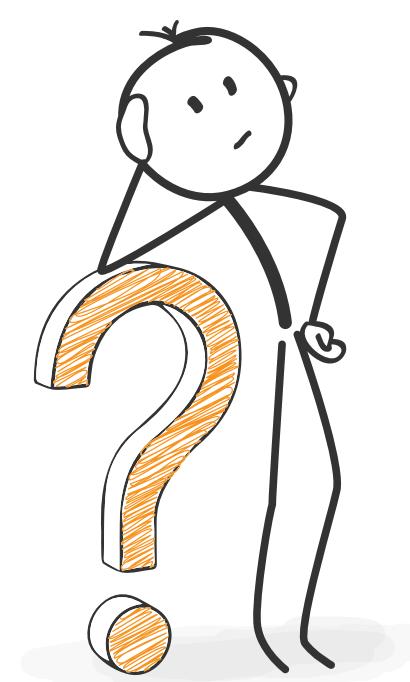
# Extending the Apache Spark ecosystem



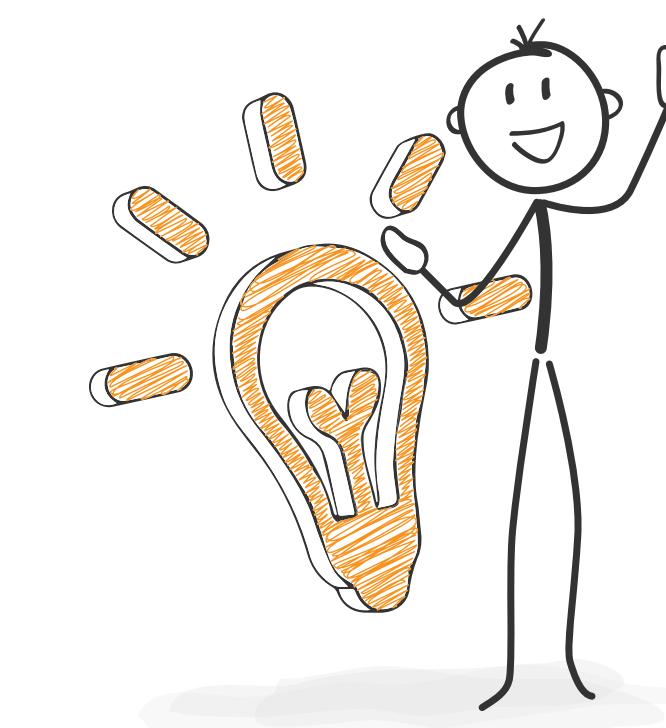
Add transactional processing  
Speed up analytical processing



- Keep Spark SQL as user frontend
- Replace Spark processing engine with a C++ engine
- IBM Wildfire

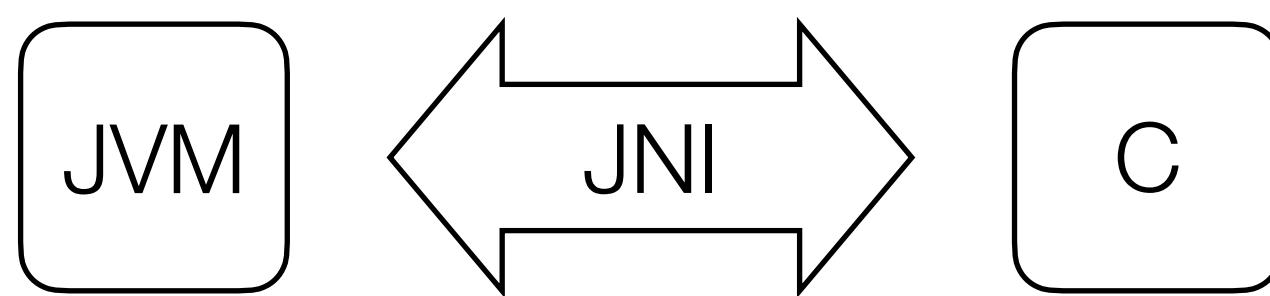


Spark SQL queries can contain arbitrary Java/Scala UDFs



- (1) Execute Java UDFs inside embedded JVM
- (2) Compile Java UDFs to machine code

# Java Native Interface (JNI)

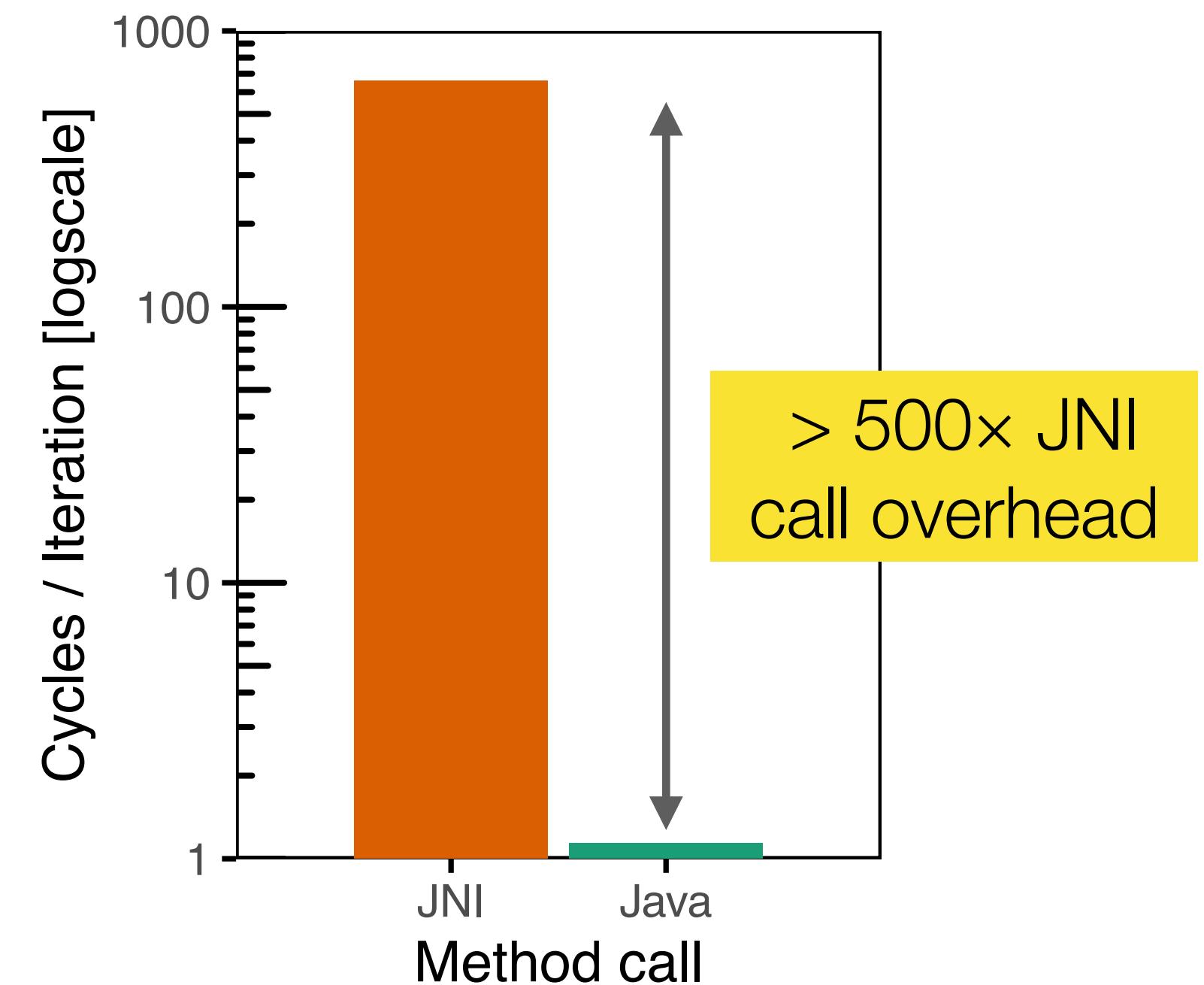


Java methods can be implemented in C

C programs can instantiate a JVM

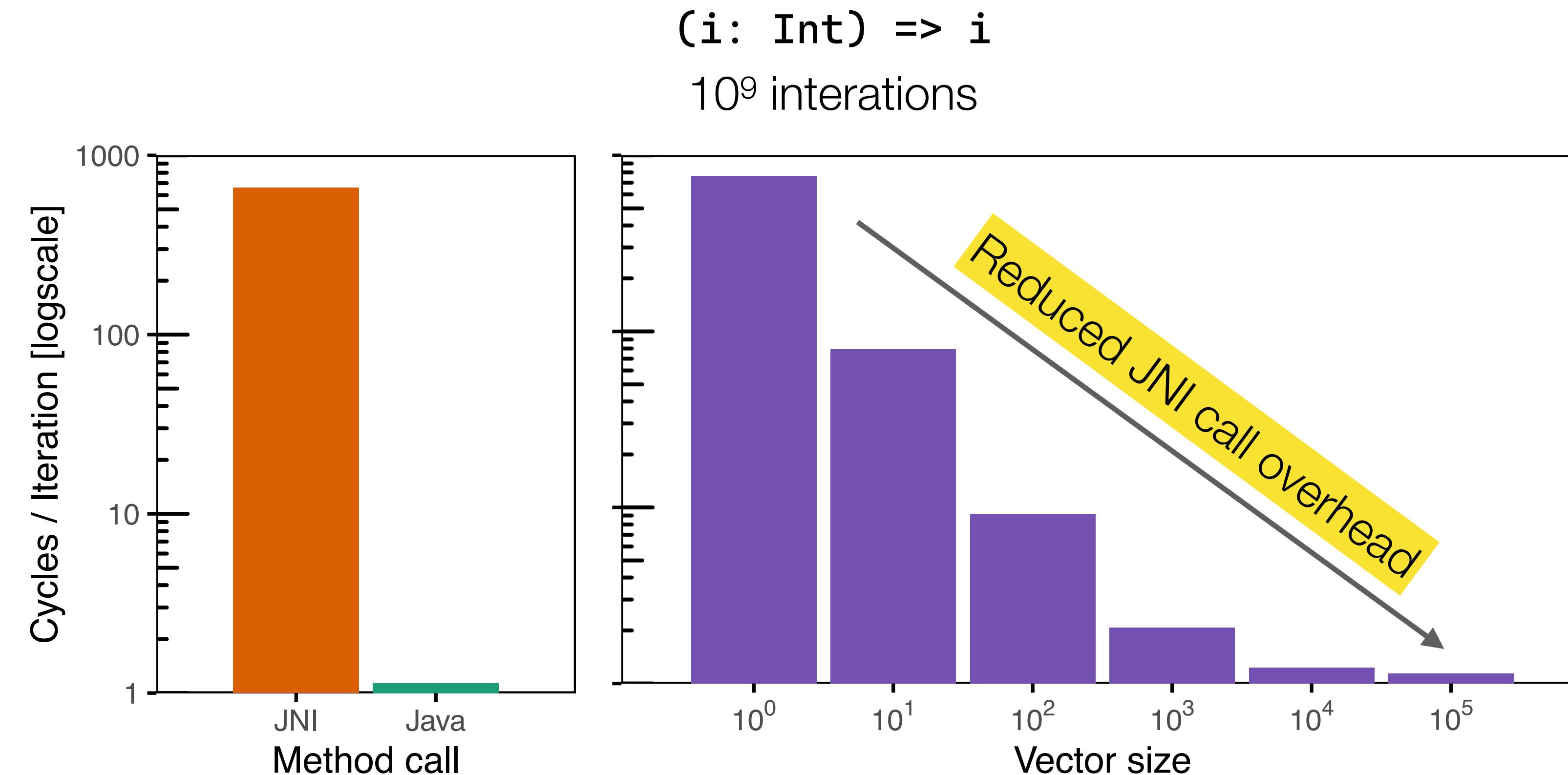
APIs to manipulate Java objects  
and to call Java methods

$(i: \text{Int}) \Rightarrow i$   
 $10^9$  iterations



**JNI call for every tuple has significant overhead.**

# Vectorized execution



Moving part of the loop inside the JVM eliminates JNI call overhead.

# Vectorized execution in embedded JVM

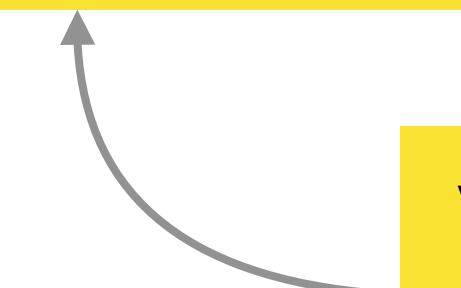
```
output = udf.apply(input0, ..., inputN);
```

# Vectorized execution in embedded JVM

```
public class StridedExecutionWrapper {  
  
    public static void executeUdf(UdfClass udf,  
                                int numRows,  
                                ByteBuffer output,  
                                ByteBuffer[] inputs) {  
  
        for (int i = 0; i < numRows; ++i) {  
            output.putX(udf.apply(inputs[0].getX(), ..., inputs[N].getX()));  
        }  
    }  
}
```

# Vectorized execution in embedded JVM

```
public class StridedExecutionWrapper {  
  
    public static void executeUdf(UdfClass udf,  
                                int numRows,  
                                ByteBuffer output,  
                                ByteBuffer[] inputs) {  
  
        for (int i = 0; i < numRows; ++i) {  
            output.putX(udf.apply(inputs[0].getX(), ..., inputs[N].getX()));  
        }  
    }  
}
```



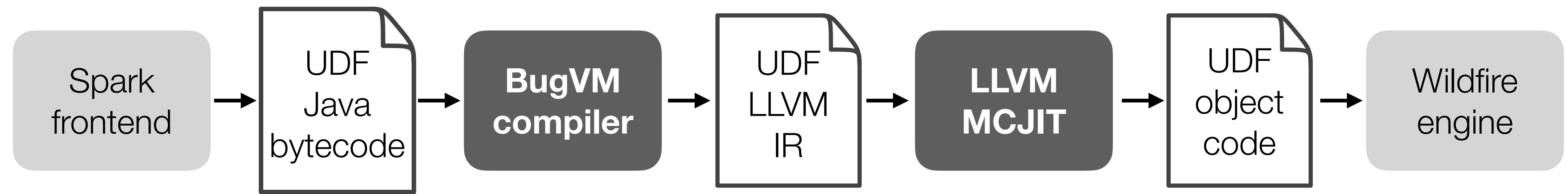
Vectorized  
execution

# Vectorized execution in embedded JVM

```
public class StridedExecutionWrapper {  
  
    public static void executeUdf(UdfClass udf,  
                                  int numRows,  
                                  ByteBuffer output,  
                                  ByteBuffer[] inputs) {  
  
        for (int i = 0; i < numRows; ++i) {  
            output.putX(udf.apply(inputs[0].getX(), ..., inputs[N].getX()));  
        }  
    }  
}
```

The diagram illustrates the code flow. A yellow box highlights the parameters of the `executeUdf` method: `output` and `inputs`. An arrow points from this box to another yellow box labeled "Java Direct ByteBuffers reduce data copies". Another arrow points from the `output` line in the `for` loop to a third yellow box labeled "Vectorized execution".

# JIT compilation to machine code

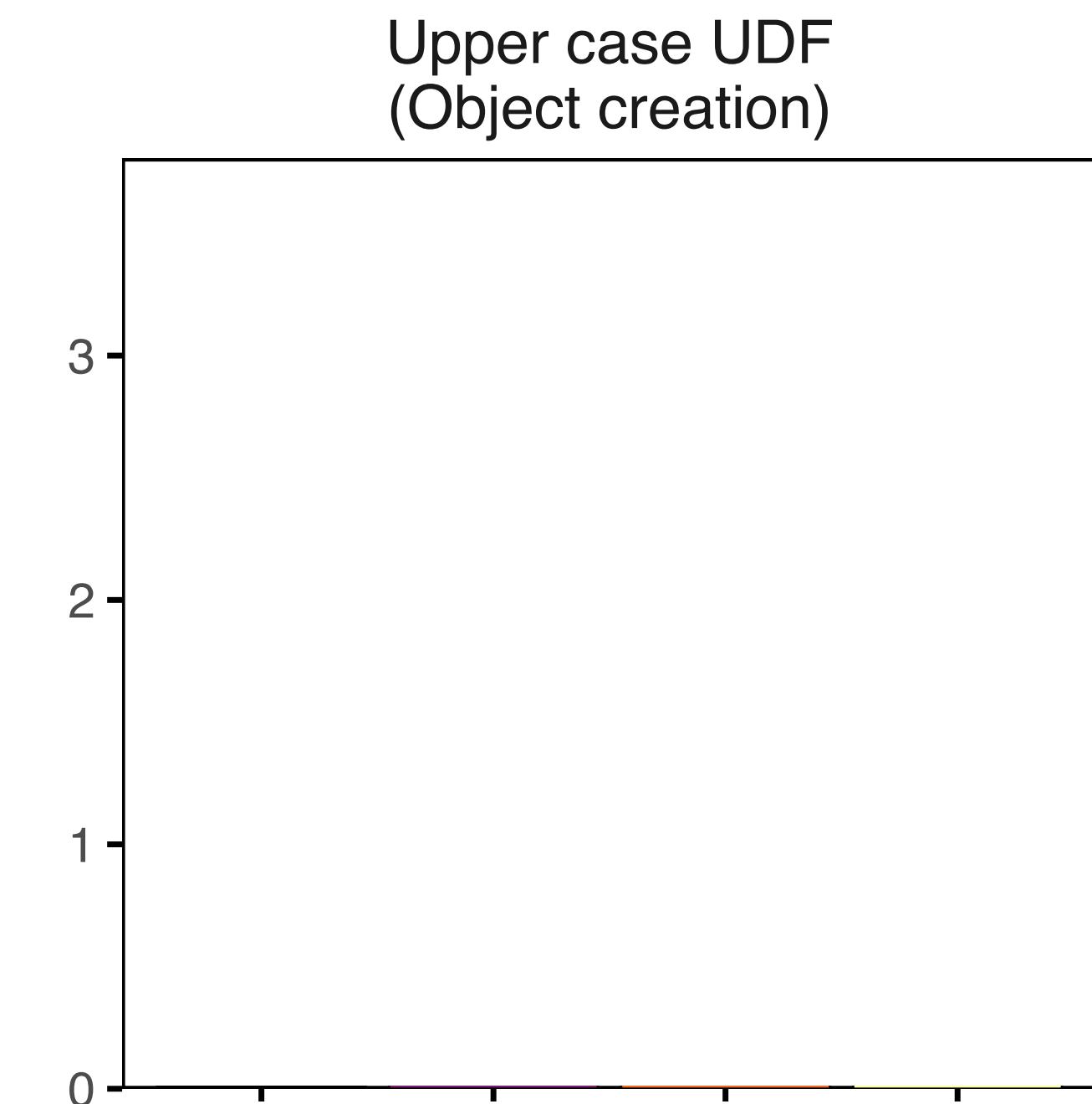
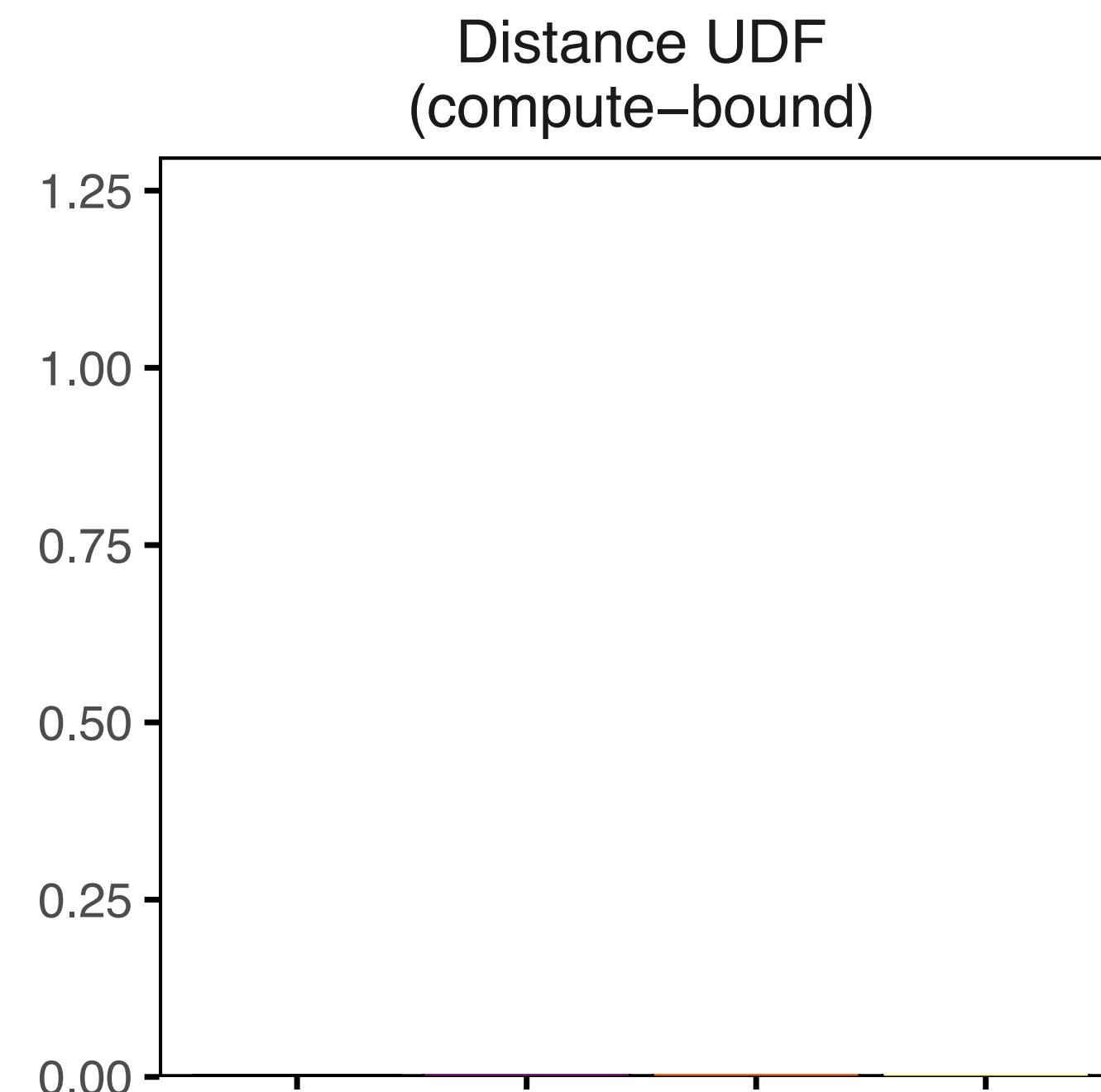
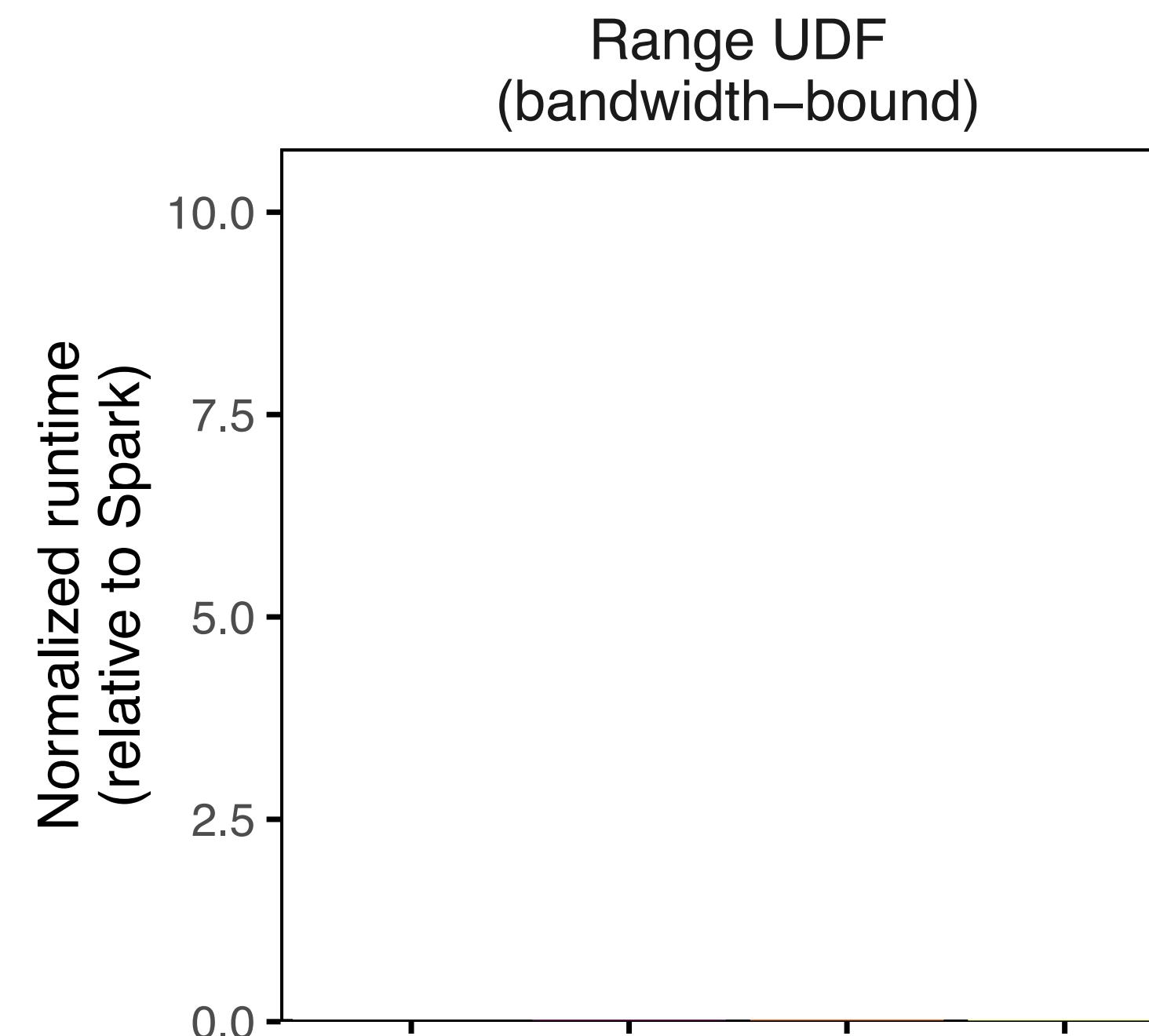


✓ No JNI call overhead

✗ No HotSpot VM optimizations

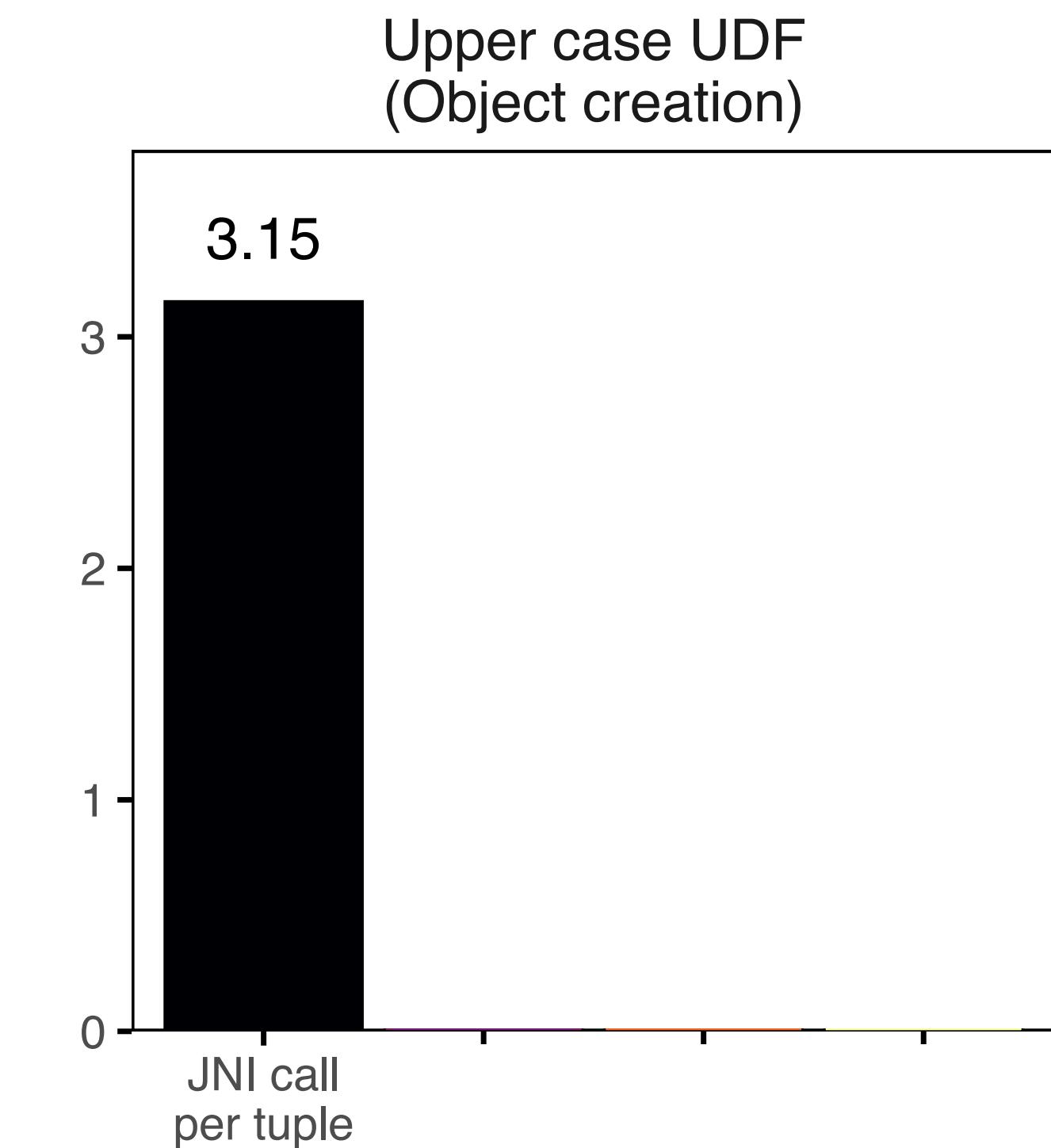
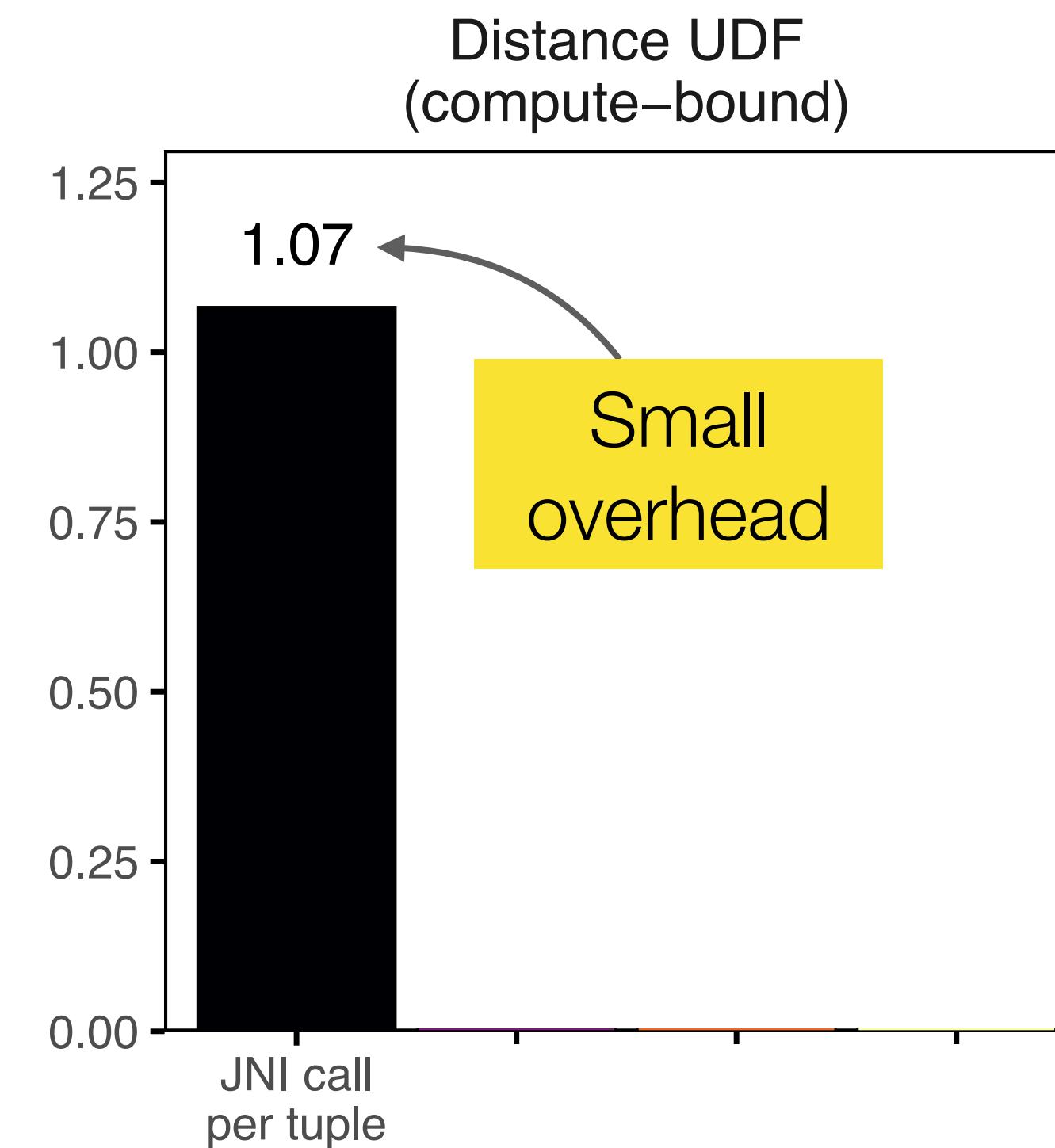
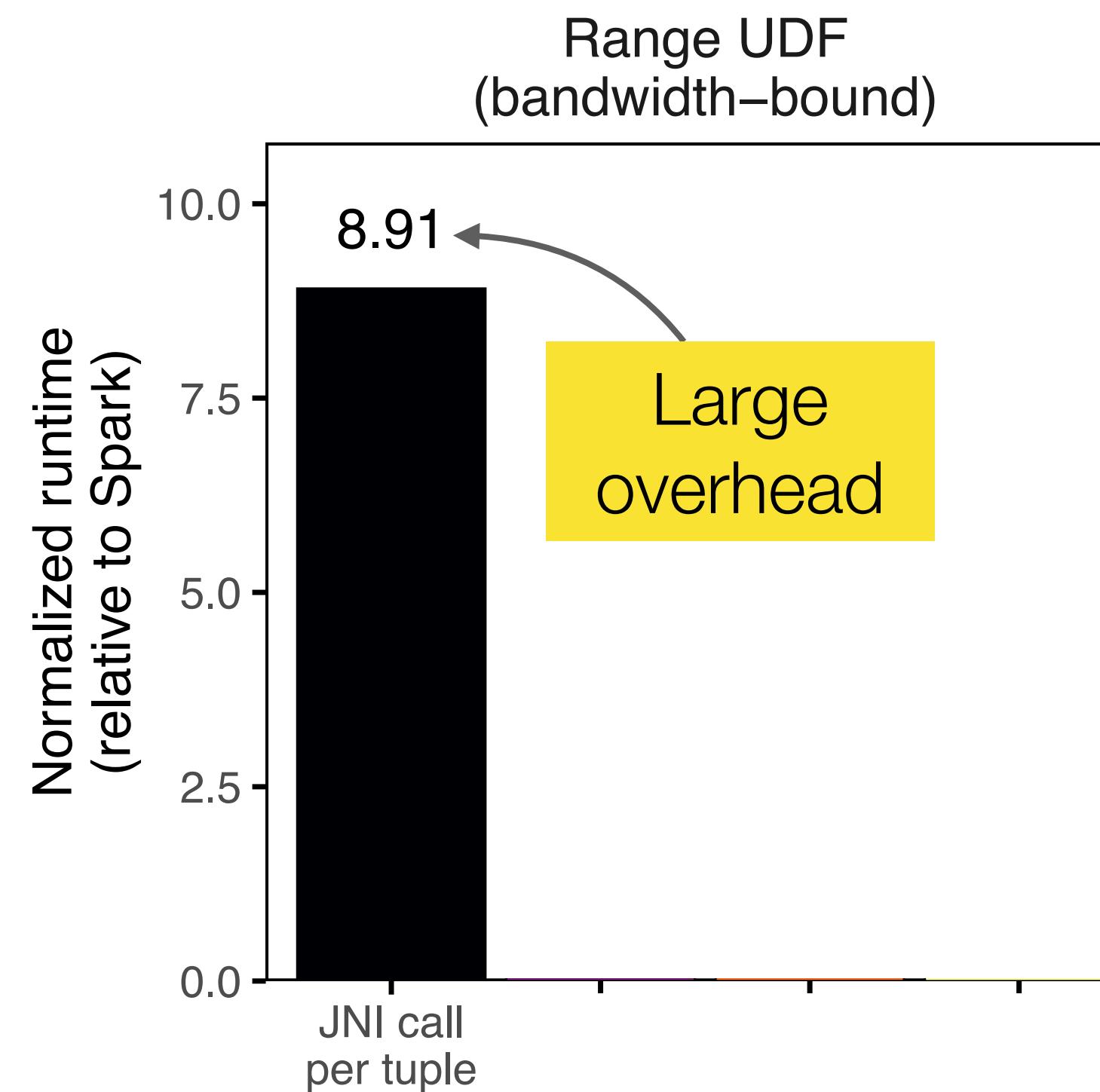
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



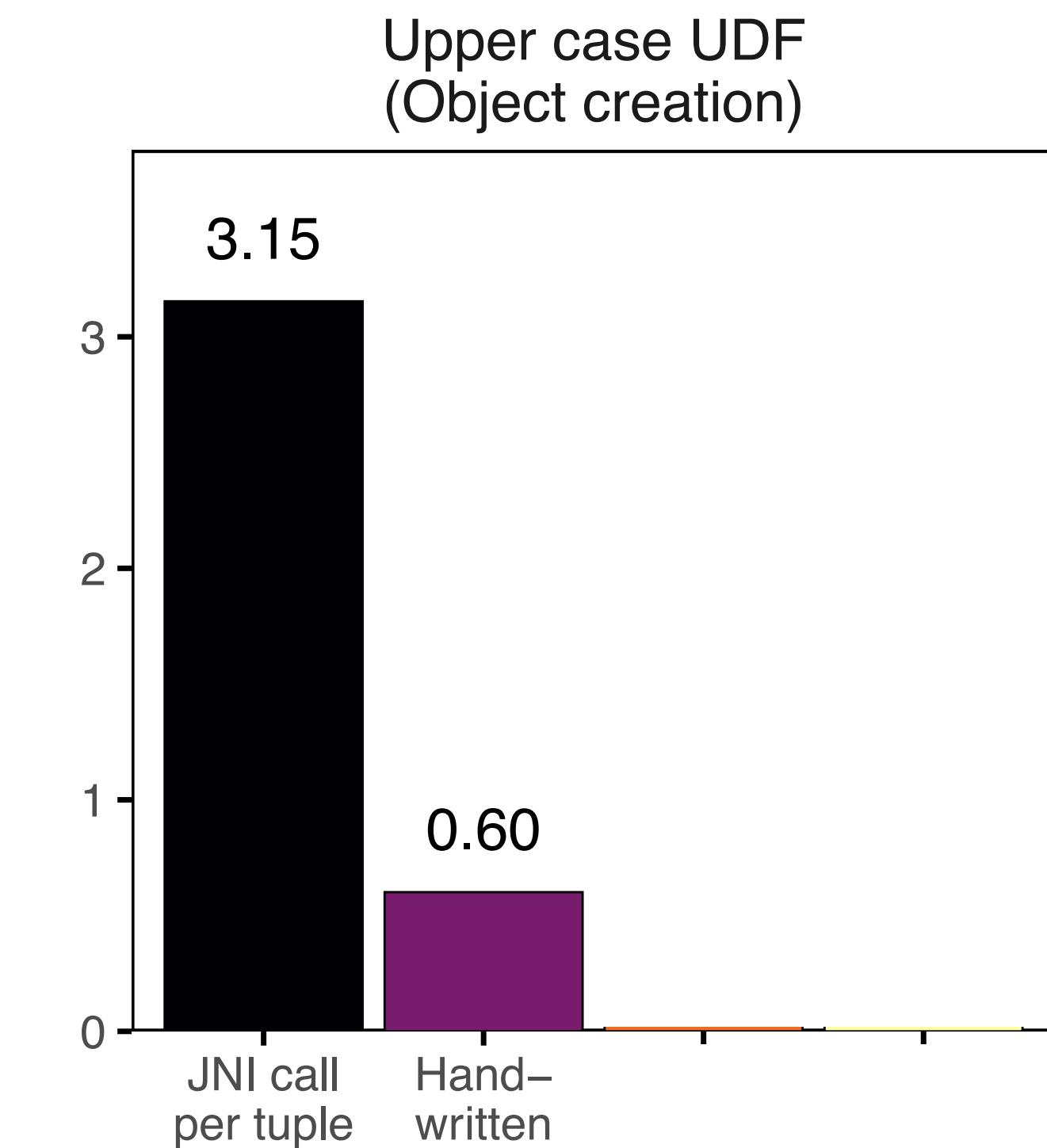
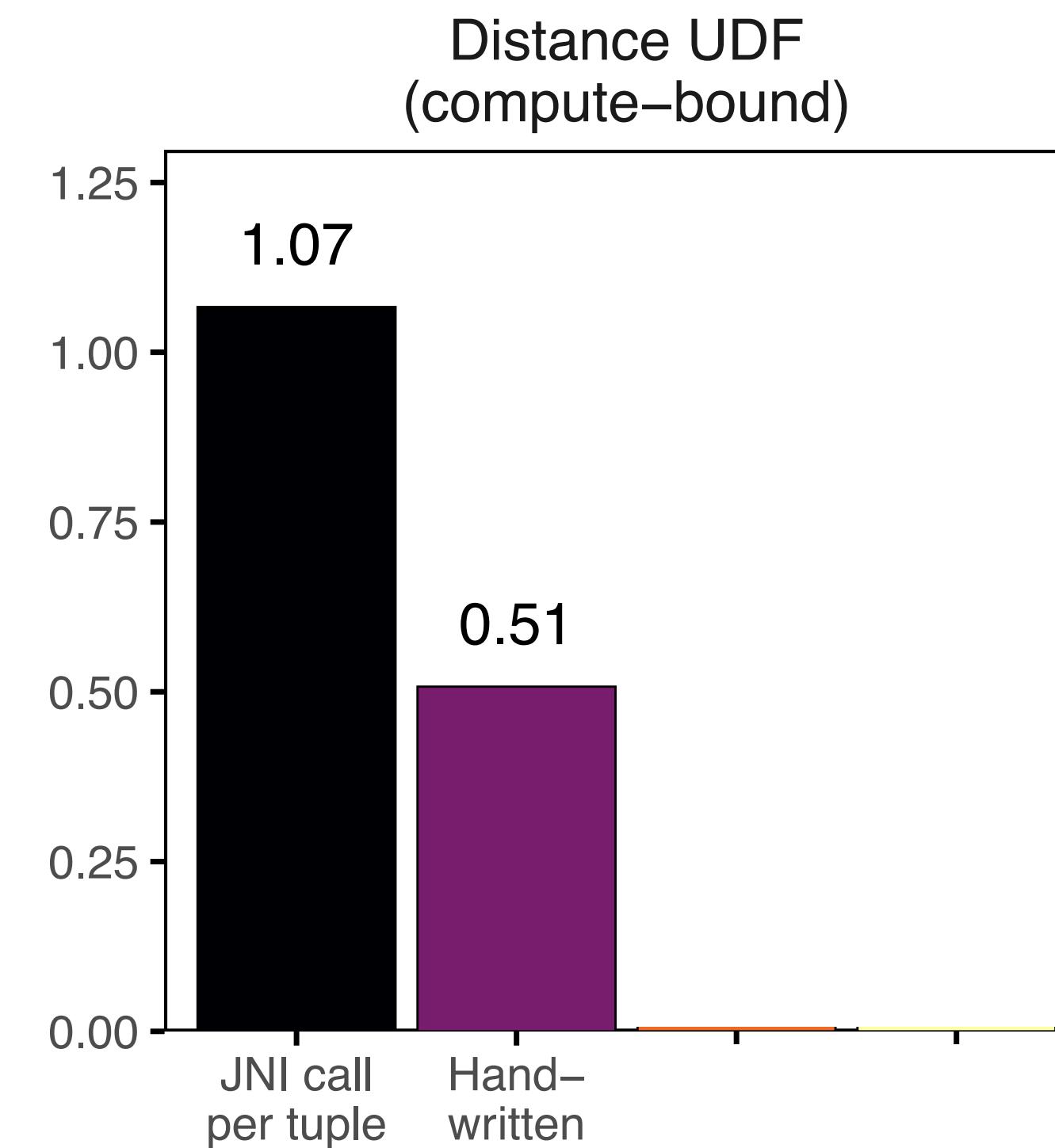
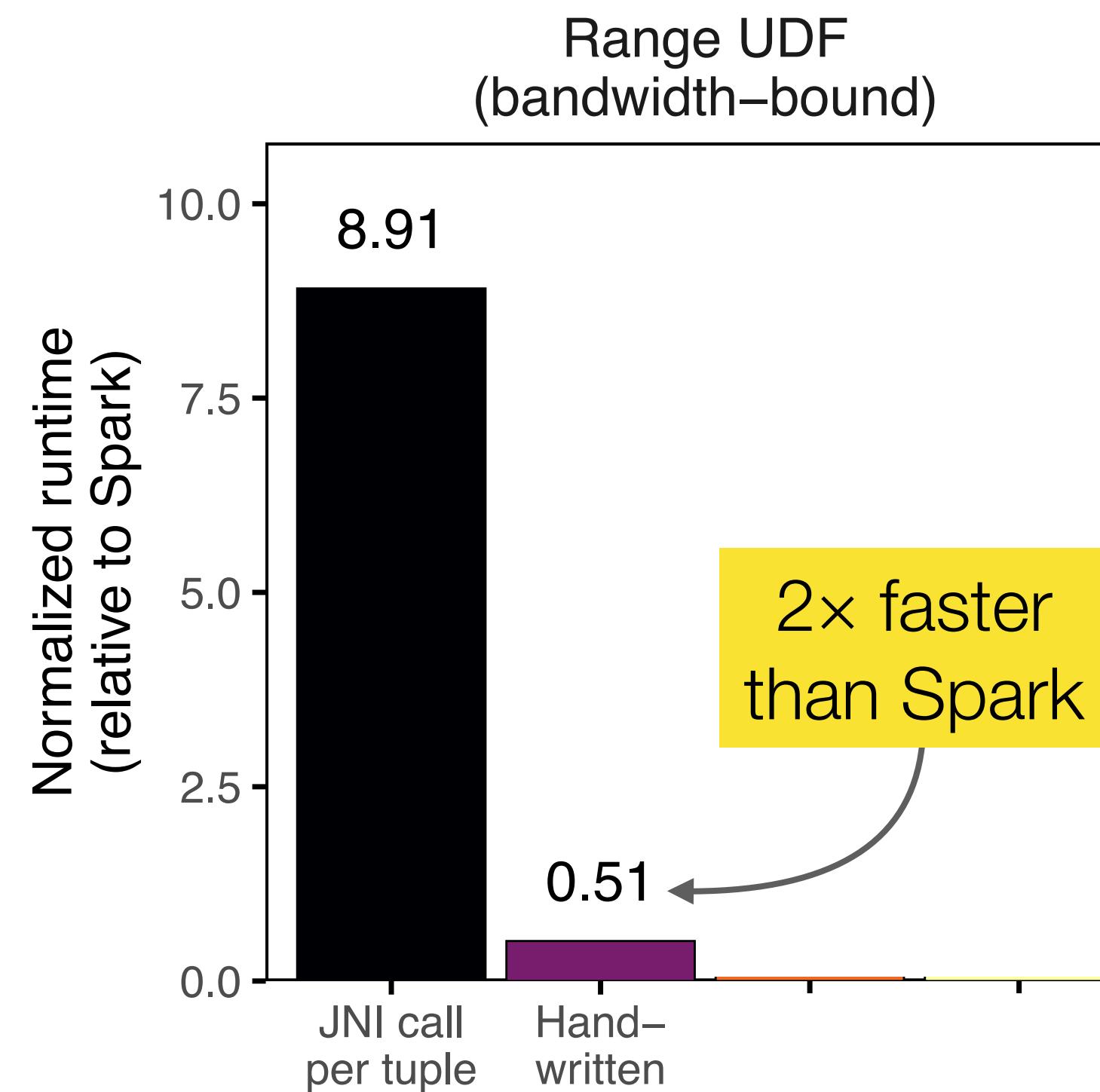
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



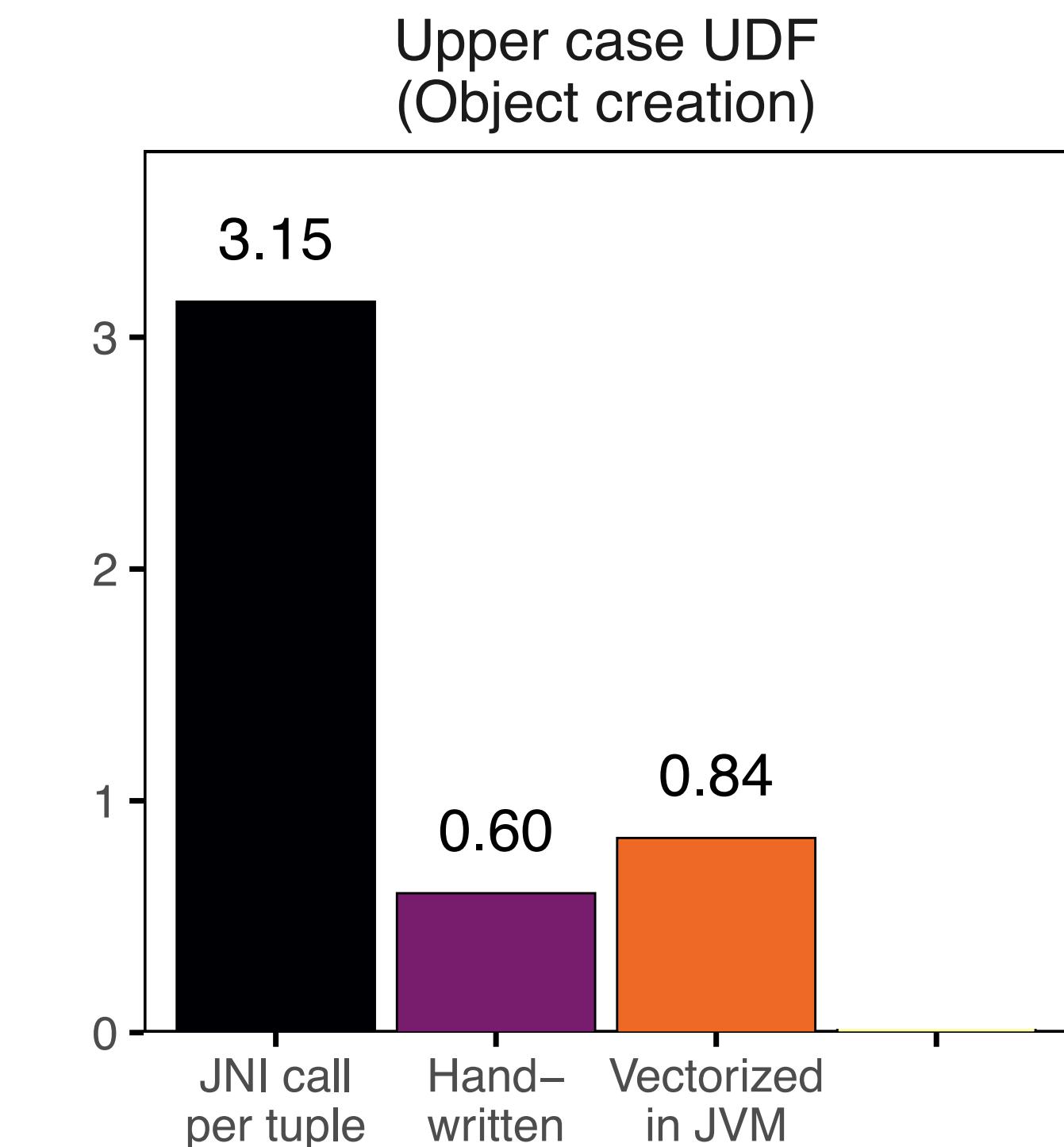
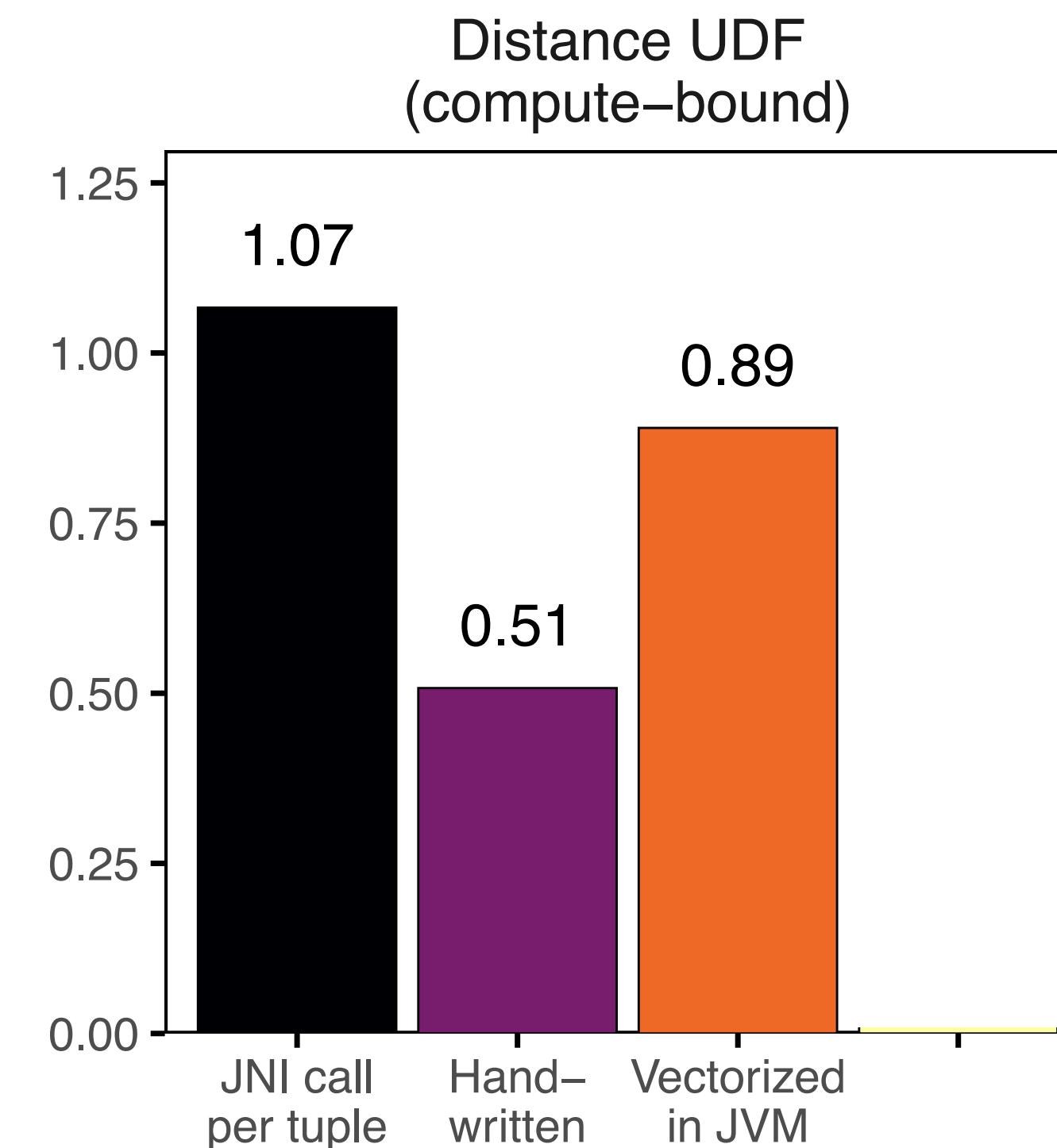
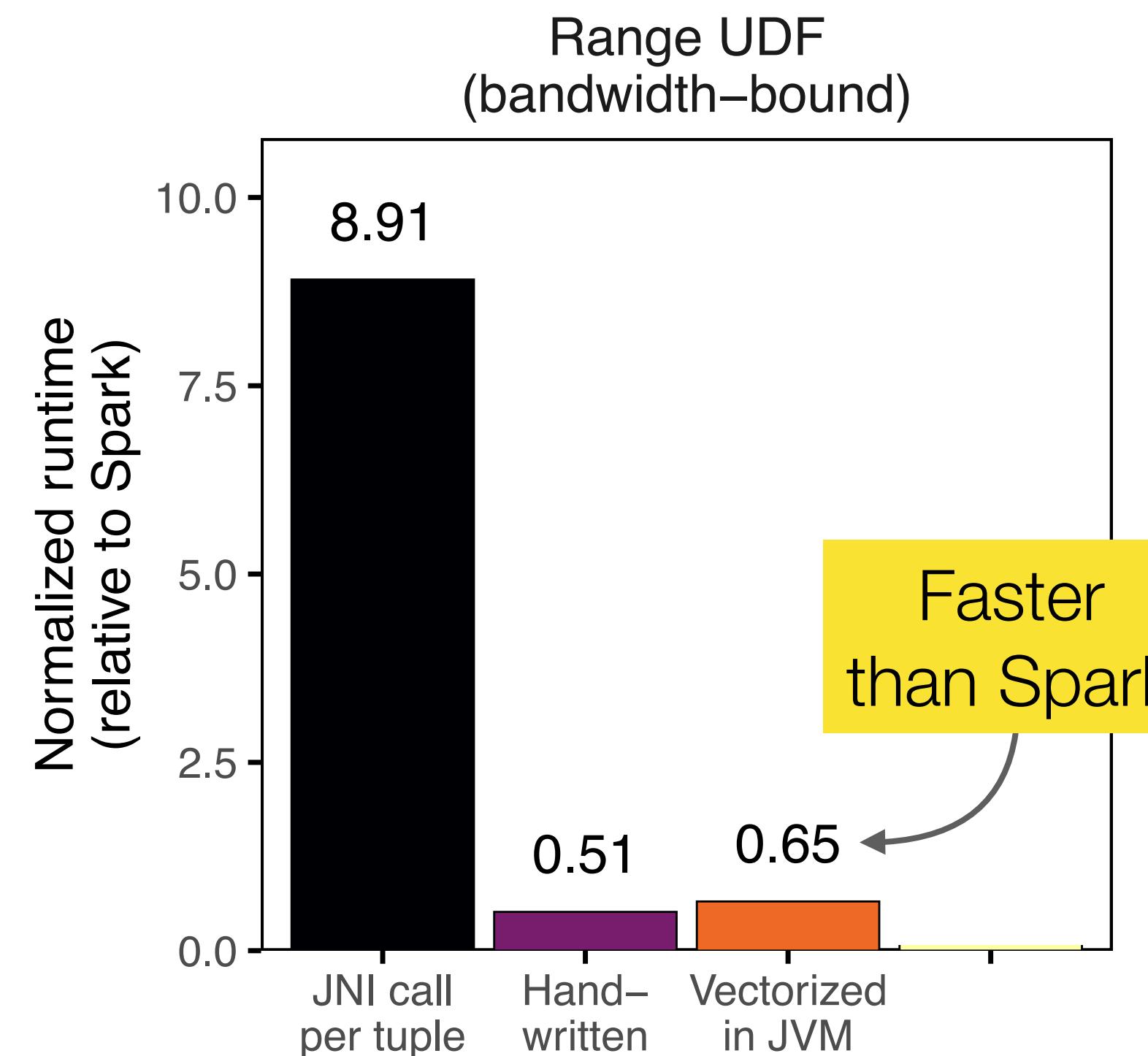
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



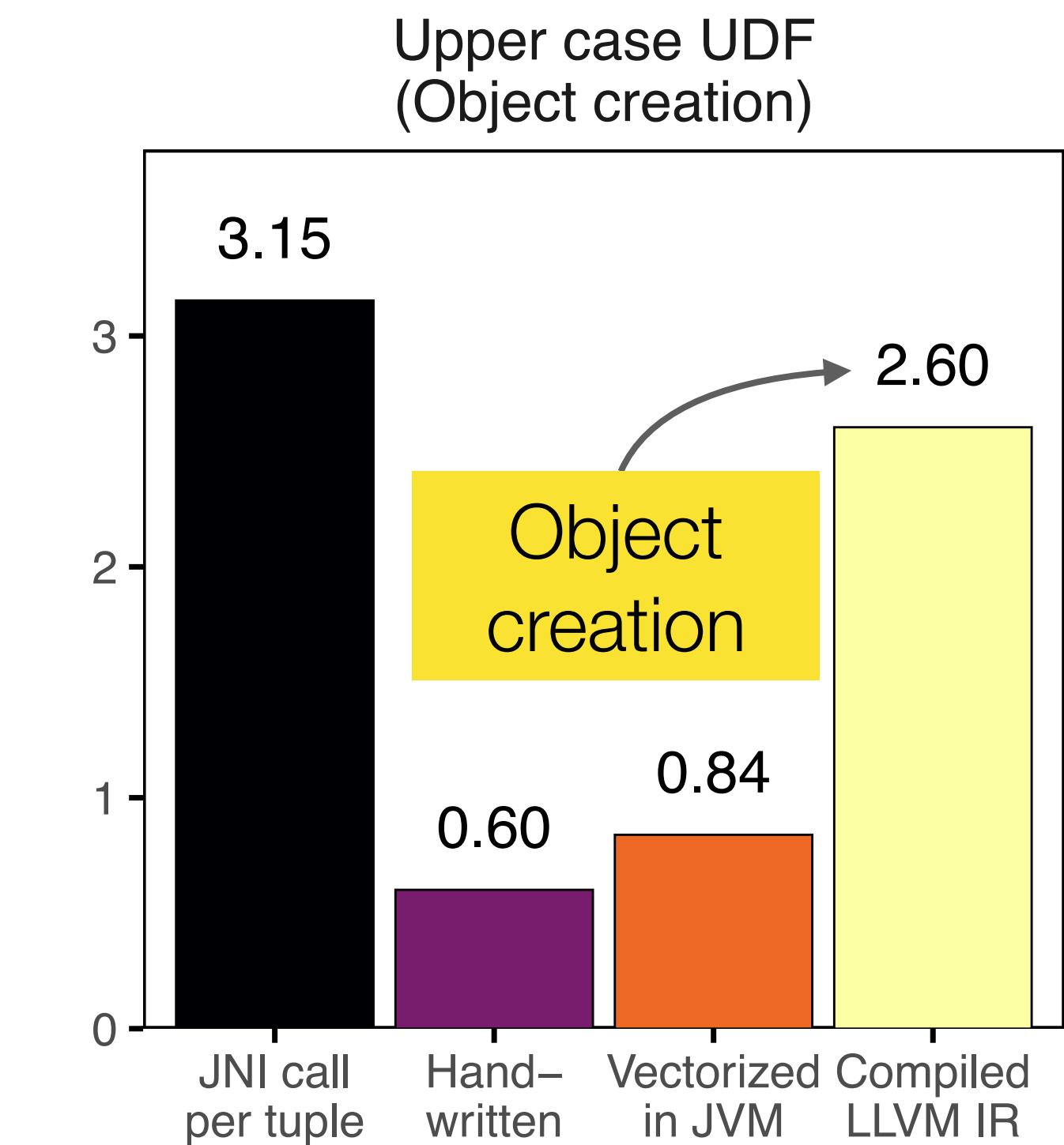
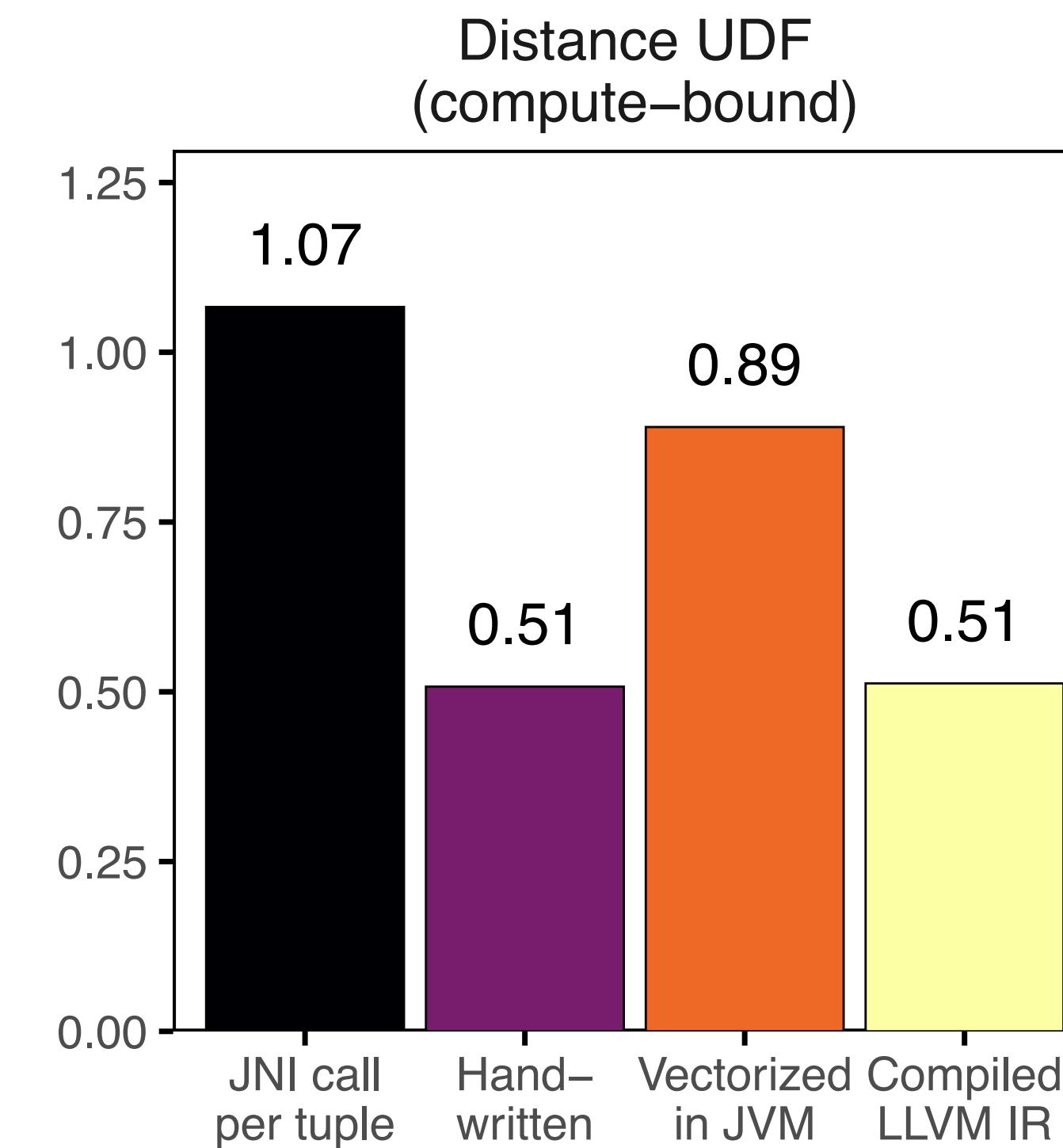
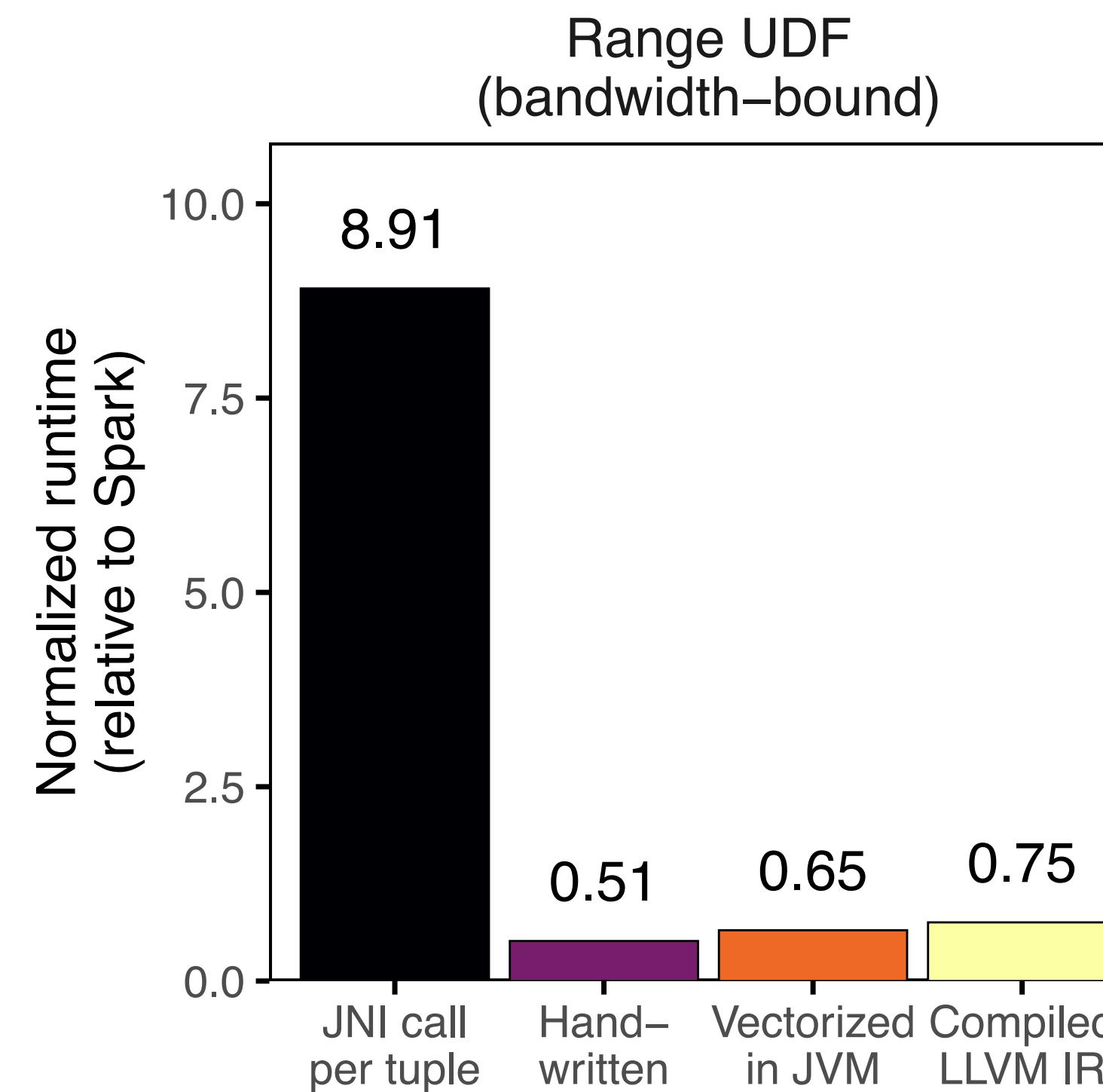
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



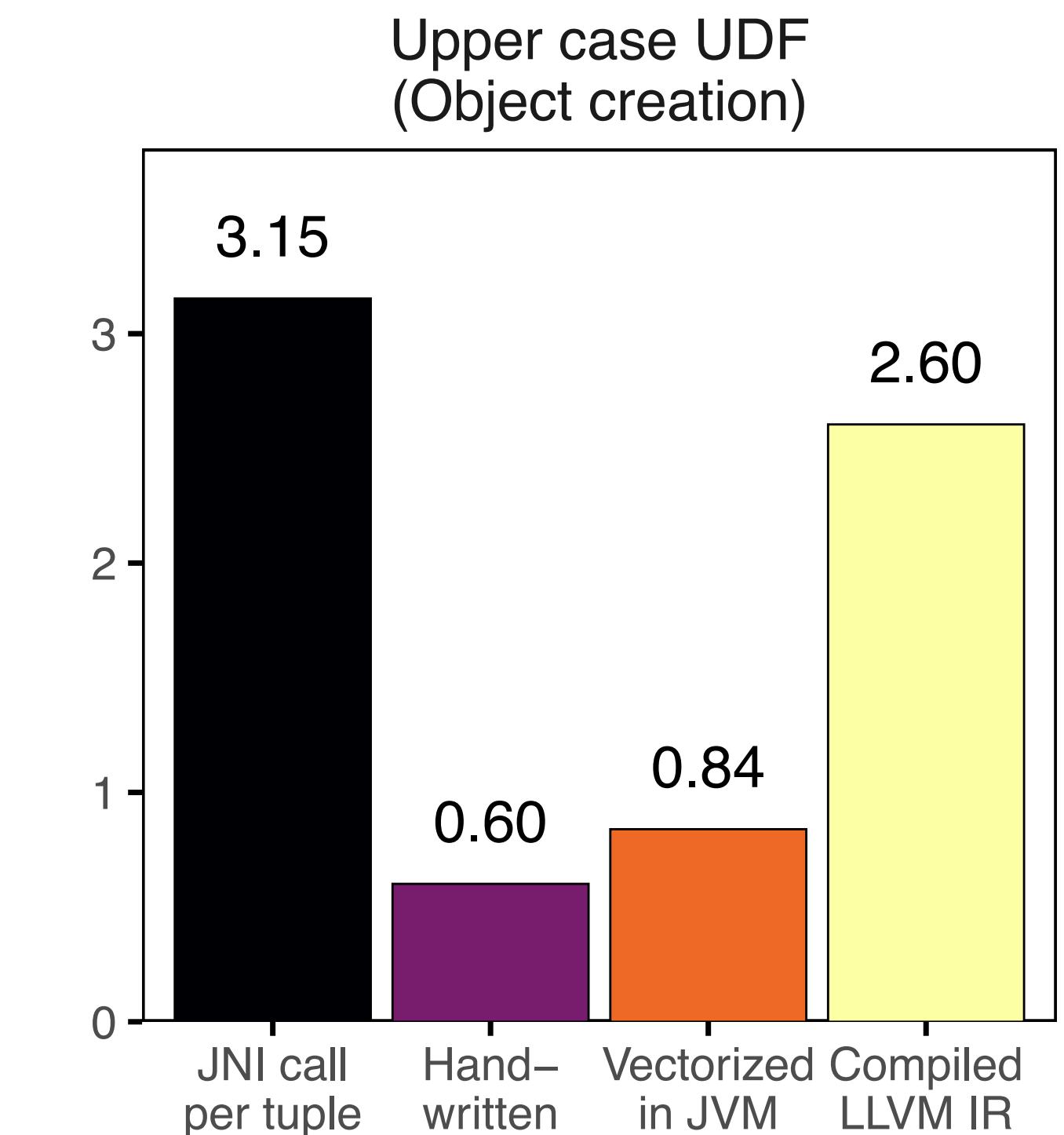
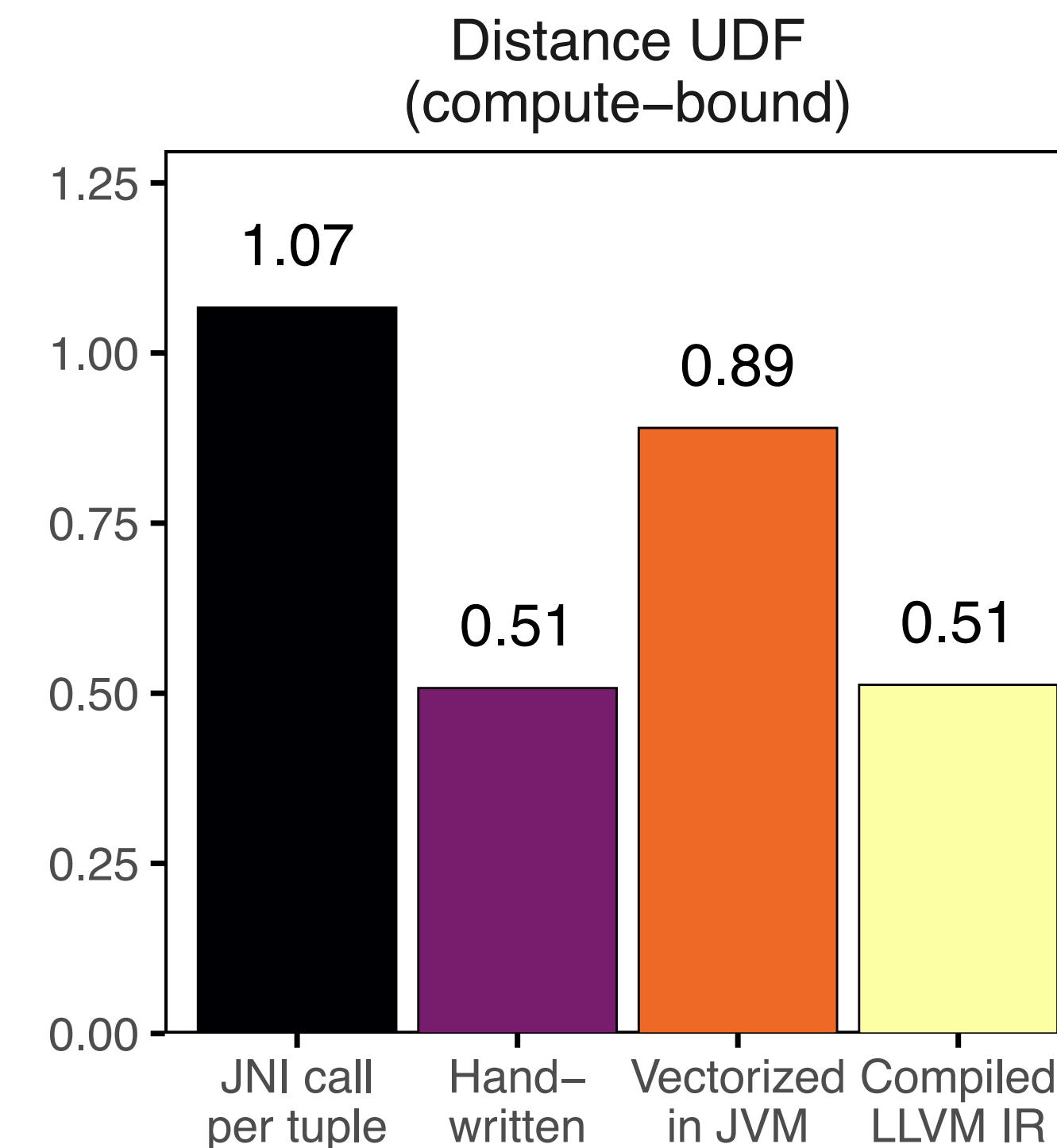
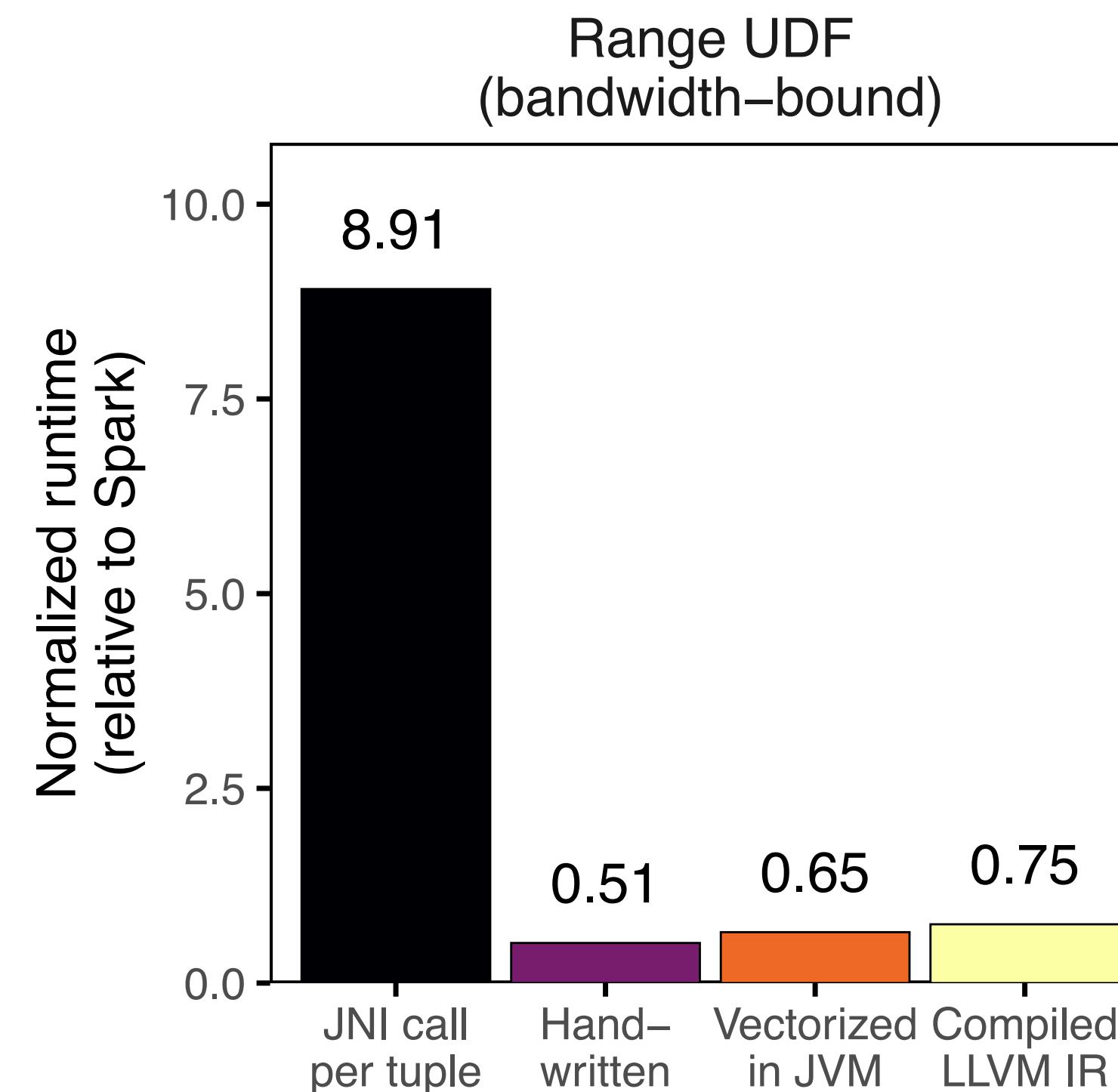
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



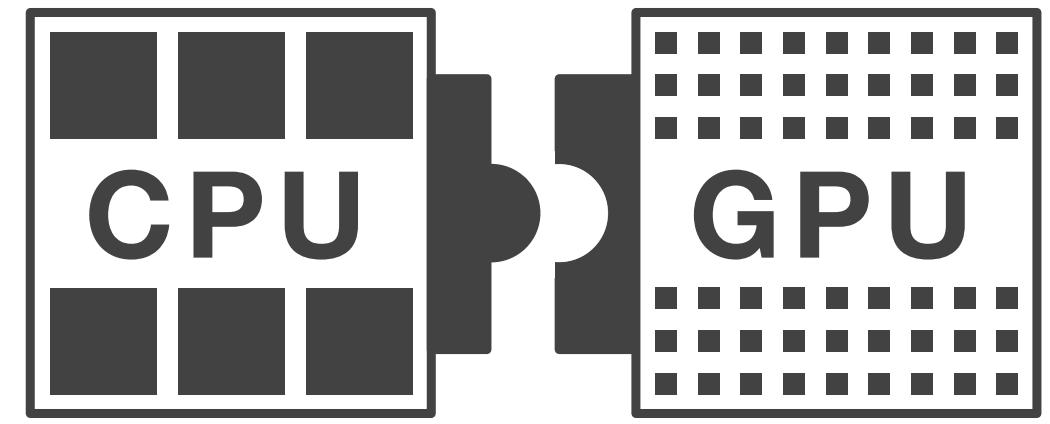
# Evaluation

Runtimes relative to execution in Spark, different scales on y axes!



Vectorized execution in Wildfire is faster than execution in Spark.

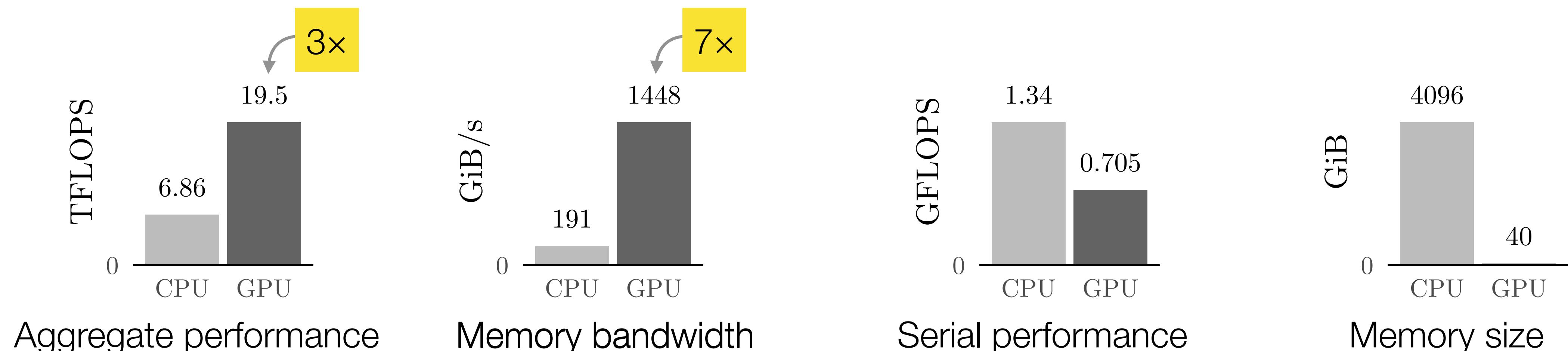
Compilation to machine code is fast if UDF is compute-heavy and does not create objects.



# Query Processing on Heterogeneous CPU/GPU Systems

# Performance characteristics of CPUs and GPUs

AMD EPYC 7702P vs. NVIDIA A100



## CPUs

Optimized for **single thread performance**

Extract implicit instruction parallelism

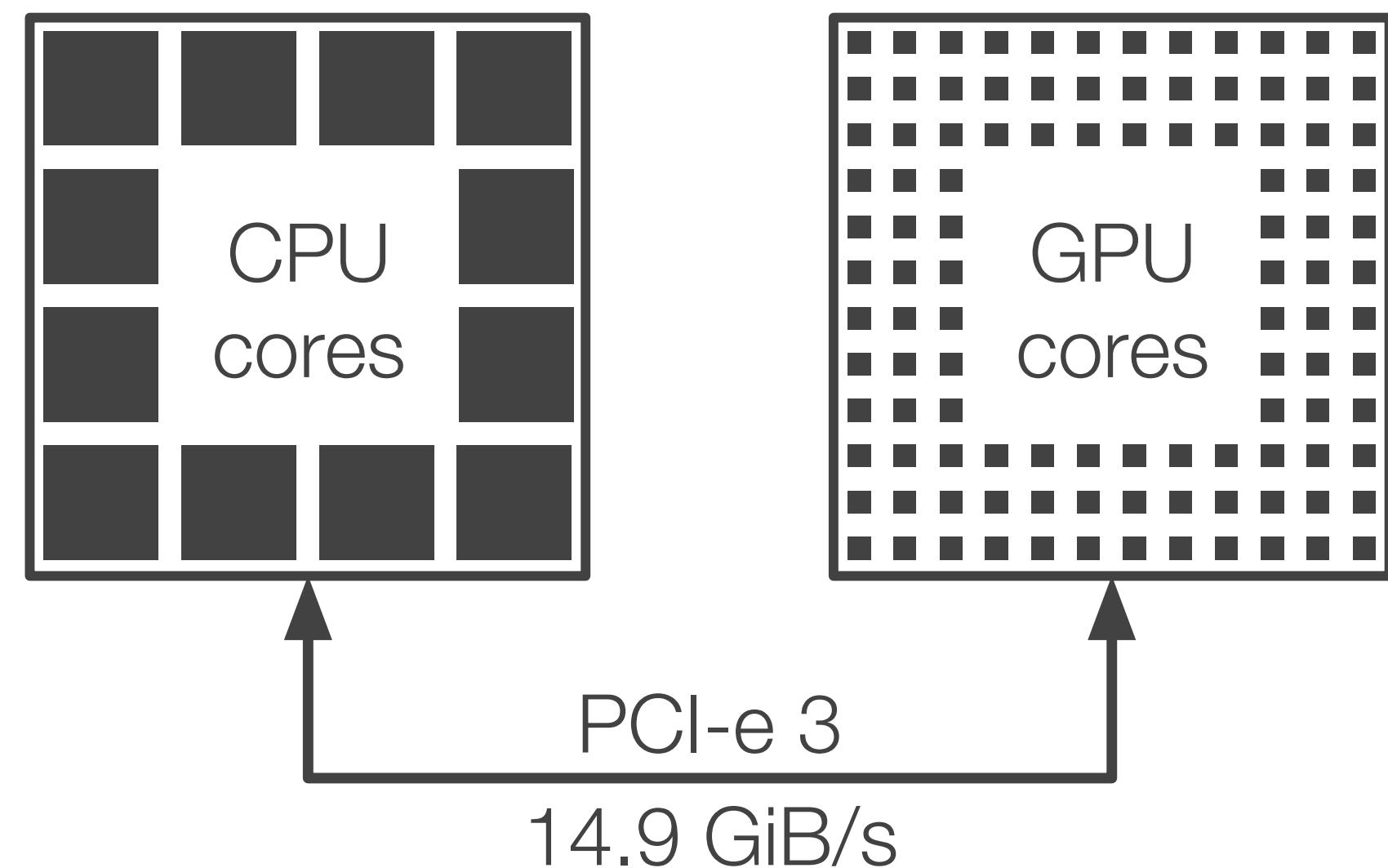
## GPUs

Optimized for **throughput applications**

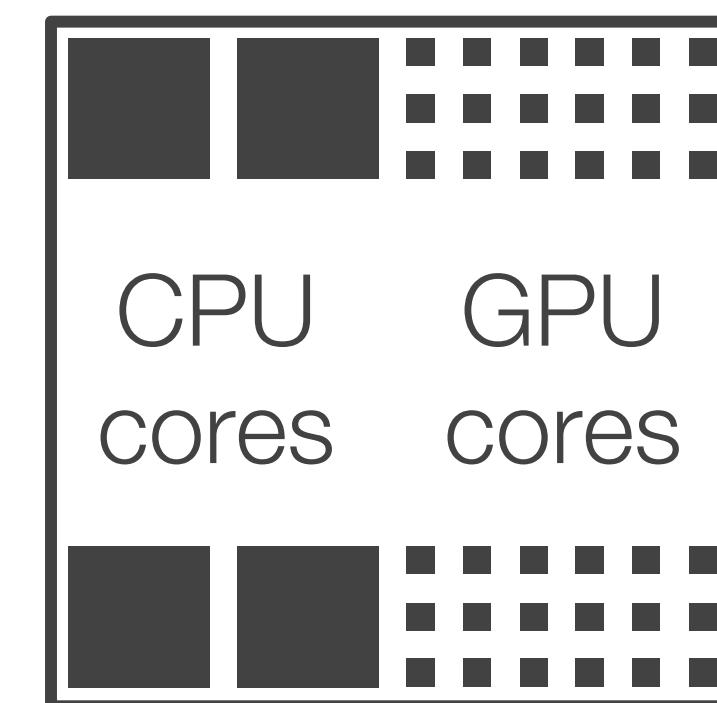
Exploit explicit data parallelism

# GPU integration

Multi-core CPU and dedicated GPU



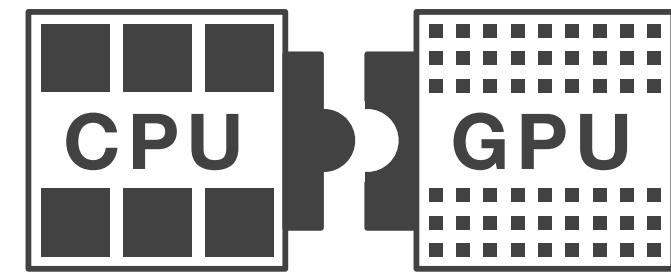
CPU and GPU cores on single die



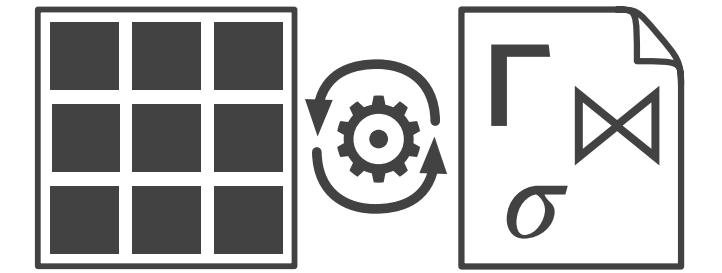
# Challenges of a heterogeneous CPU/GPU query processing system



Schedule query workload on different processors



Adapt query processing code to different processors



Mitigate the data transfer bottleneck

# Workload scheduling classification scheme

**Processor usage** – generic compute resource or specialized for specific tasks

**Scheduling time** – before or during query execution

**Scheduling strategy** – heuristics, cost models, work stealing

**Workload distribution**

**Task granularity**

**Data partitioning**



– what kind of tasks are scheduled

# Workload scheduling classification scheme

**Processor usage** – generic compute resource or specialized for specific tasks

**Scheduling time** – before or during query execution

**Scheduling strategy** – heuristics, cost models, work stealing, ...

**Workload distribution**

**Task granularity**

**Data partitioning**



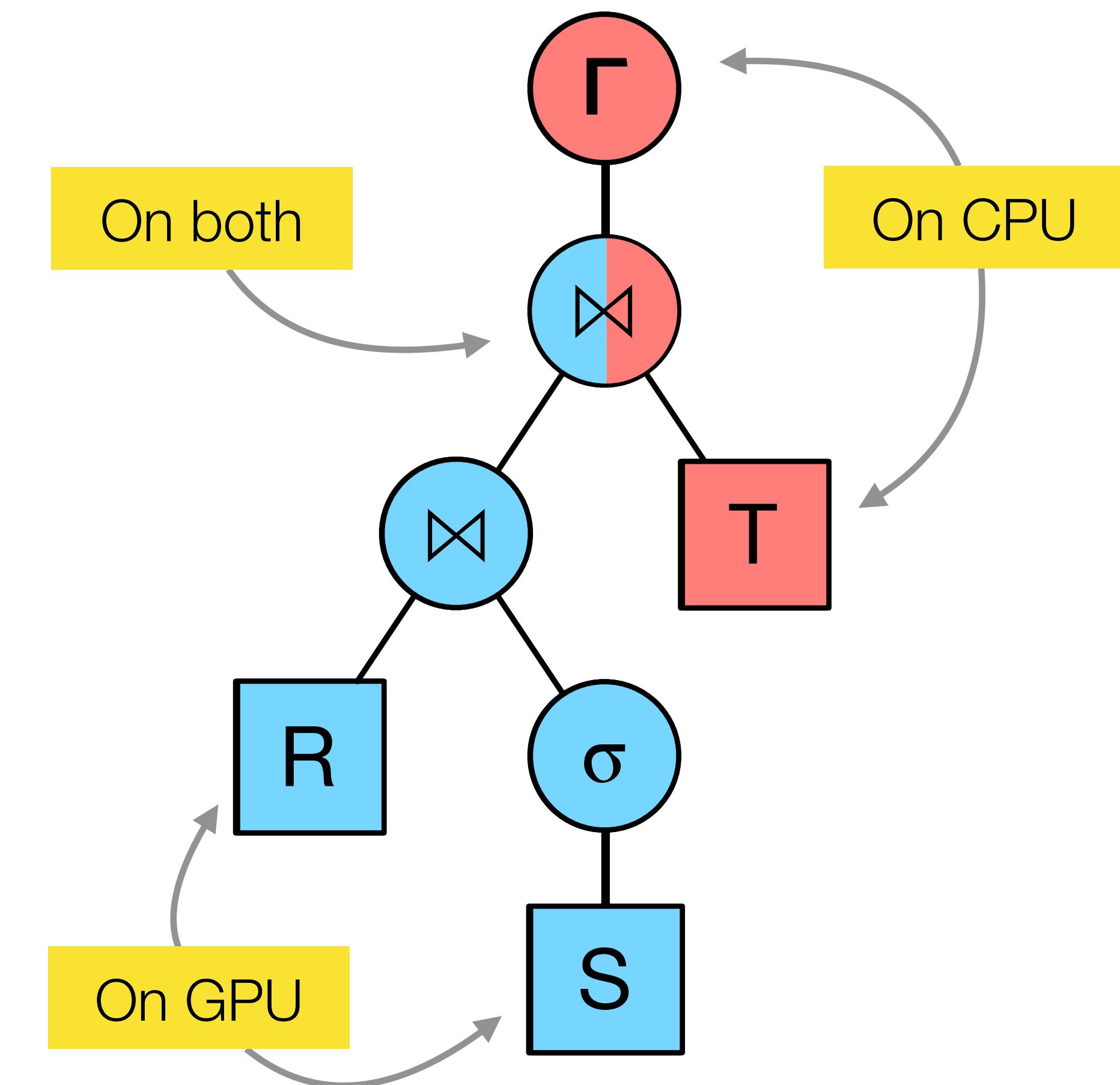
– what kind of tasks are scheduled

# Processors as generic compute resources

Tasks can be scheduled on any processor

Heuristics, cost models, work stealing

Processors are distinguished by their relative throughput



# Processors as specialized resources

Developer analyzes tasks and assigns them to suitable processor

Example: Approximate & Refine  
[Pirk et al., ICDE 2014]

Bitwise partition of data

```
001000110101|10110100111101110101
```

# Processors as specialized resources

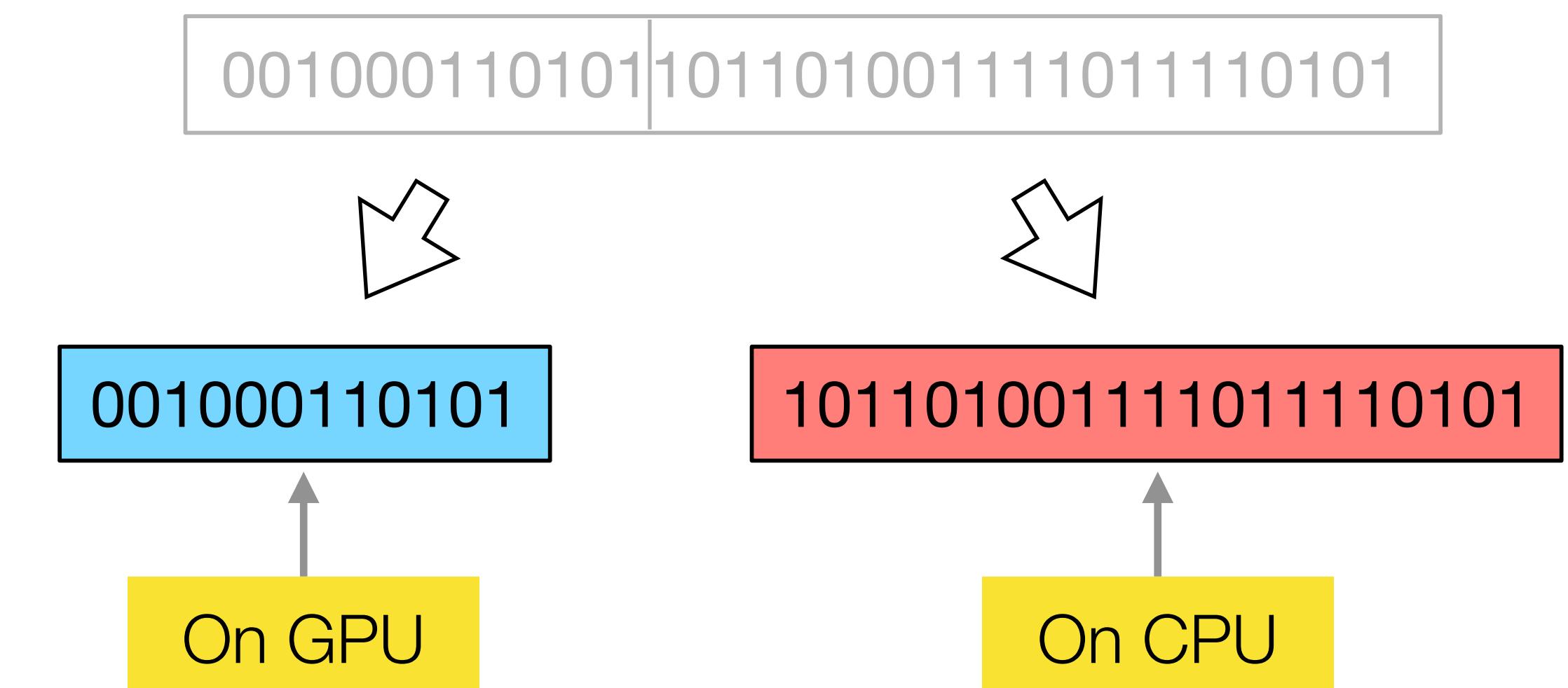
Developer analyzes tasks and assigns them to suitable processor

Example: Approximate & Refine  
[Pirk et al., ICDE 2014]

Bitwise partition of data

Place higher bits on GPU, lower bits on CPU

Compute approximate result on GPU & refine into exact result on CPU



Publication		GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<i>Full query processing systems</i>								
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	Bits
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
GDB	He et al., 2009	Dedicated	Generic	Static	Cost model	Task partitions	Operator	Tuples
CoGaDB	Breß, 2014	Dedicated	Generic	Both	Data locality, cost model	Operator placement	Operator	Columns
SABER	Koliouisis et al., 2016	Dedicated	Generic	Dynamic	Load balancing, cost model	Single partition	Query	Data batch
DB2 BLU	Meraji et al., 2016	Dedicated	Generic	Dynamic	Task nature, load balancing	Task partitions	Operator	Tuples
HetExchange	Chrysogelos et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Pipeline	Data batch
He et al.	He et al., 2014	Integrated	Hybrid	Static	Task nature, cost model	Task partitions	Primitive	Tuples
DIDO	Zhang et al., 2017	Integrated	Hybrid	Hybrid	Task nature, load balancing, cost model	Operator placement	Operator	Query batch
FineStream	Zhang et al., 2020	Integrated	Generic	Dynamic	Cost model	Operator placement	Operator	—
HERO	Karnagel et al., 2017	Both	Generic	Dynamic	Data locality, cost model	Operator placement	Primitive	—
<i>Individual query processing tasks</i>								
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	Algorithm-specific	—	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
G-Grid	Li et al., 2018	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
GAT	Zhang et al., 2018	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Sioulas et al.	Sioulas et al., 2019	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Gubner et al.	Gubner et al., 2019	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Operator	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Tuple batch
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	—	Query batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	—	Cuboids, points
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Primitive	Tuples
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	Algorithm-specific	—	—

Publication	GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<b>Full query processing systems</b>							
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type
GDB	He et al., 2009	Dedicated	Specialized	Static	Task partitions	Task partitions	Operator
CoGaDB	Breß, 2014	Dedicated	Specialized	Dynamic	Task partitions	Task partitions	Tuples
SABER	Koliousis et al., 2016	Dedicated	Specialized	Hybrid	Load balancing, data locality	Task partitions	Operator
DB2 BLU	Meraji et al., 2016	Dedicated	Generic	Dynamic	Task nature, load balancing	Task partitions	Columns
HetExchange	Chrysogelos et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Query
He et al.	He et al., 2014	Integrated	Hybrid	Static	Task nature, cost model	Task partitions	Data batch
DIDO	Zhang et al., 2017	Integrated	Hybrid	Hybrid	Task nature, load balancing, cost model	Operator placement	Tuples
FineStream	Zhang et al., 2020	Integrated	Generic	Dynamic	Cost model	Operator placement	Pipeline
HERO	Karnagel et al., 2017	Both	Generic	Dynamic	Data locality, cost model	Operator placement	Data batch
<b>Individual query processing tasks</b>							
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	Algorithm-specific	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
G-Grid	Li et al., 2018	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
GAT	Zhang et al., 2018	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
Sioulas et al.	Sioulas et al., 2019	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—
Gubner et al.	Gubner et al., 2019	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	Tuple batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	Query batch
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Cuboids, points
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	Algorithm-specific	Tuples

## Full query processing systems

Publication		GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<i>Full query processing systems</i>								
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	Bits
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	Query type	—
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	Query type	—
GDB	He et al., 2009	Dedicated	Generic	Static	Cost model	Task partitions	Operator	Tuples
CoGaDB	Breß, 2014	Dedicated	Generic	Both	Data locality, cost model	Operator placement	Operator	Columns
SABER	Koliouisis et al., 2016	Dedicated	Generic	Dynamic	Load balancing, cost model	Single partition	Query	Data batch
DB2 BLU	Meraji et al., 2016	Dedicated	Generic	Dynamic	Task nature, load balancing	Task partitions	Operator	Tuples
HetExchange	Chrysogelos et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Pipeline	Data batch
He et al.	He et al., 2014	Integrated	Hybrid	Static	Task nature, cost model	Task partitions	Primitive	Tuples
DIDO	Zhang et al., 2017	Integrated	Hybrid	Hybrid	Task nature, load balancing, cost model	Operator placement	Operator	Query batch
FineStream	Zhang et al., 2020	Integrated	Generic	Dynamic	Cost model	Operator placement	Operator	—
HERO	Karnagel et al., 2017	Both	Generic	Dynamic	Data locality, cost model	Operator placement	Primitive	—
<i>Individual query processing tasks</i>								
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	<i>Algorithm-specific</i>	—	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
G-Grid	Li et al., 2018	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
GAT	Zhang et al., 2018	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Sioulas et al.	Sioulas et al., 2019	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Gubner et al.	Gubner et al., 2019	Dedicated	Specialized	Static	Task nature	Task partitions	Operator	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Operator	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Tuple batch
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	—	Query batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	—	Cuboids, points
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Primitive	Tuples
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—

## Individual query processing tasks

Publication		GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<i>Full query processing systems</i>								
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	Bits
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	Query type	—
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	Query type	—
GDB	He et al., 2009	Dedicated	Generic	Static	Cost model	Task partitions	Operator	Tuples
CoGaDB	Breß, 2014	Dedicated	Generic	Both	Data locality, cost model	Operator placement	Operator	Columns
SABER	Koliouisis et al., 2016	Dedicated	Generic	Dynamic	Load balancing, cost model	Single partition	Query	Data batch
DB2 BLU	Meraji et al., 2016	Dedicated	Generic	Dynamic	Task nature, load balancing	Task partitions	Operator	Tuples
HetExchange	Chrysogelos et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Pipeline	Data batch
He et al.	He et al., 2019	Dedicated	Generic	Hybrid	Task nature	Task partitions	Primitive	Tuples
DIDO	Zhang et al., 2019	Dedicated	Generic	Hybrid	Task nature	Task partitions	Operator	Query batch
FineStream	Zhang et al., 2019	Dedicated	Generic	Hybrid	Task nature	Task partitions	Operator	—
HERO	Karnagel et al., 2013	Dedicated	Specialized	Hybrid	Task nature	Task partitions	Primitive	—
<i>Individual query processing tasks</i>								
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	<i>Algorithm-specific</i>	—	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
G-Grid	Li et al., 2018	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
GAT	Zhang et al., 2018	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Sioulas et al.	Sioulas et al., 2019	Dedicated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—
Gubner et al.	Gubner et al., 2019	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Operator	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Tuple batch
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	—	Query batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	—	Cuboids, points
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Primitive	Tuples
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	<i>Algorithm-specific</i>	—	—

Dedicated GPUs are used as specialized compute resources

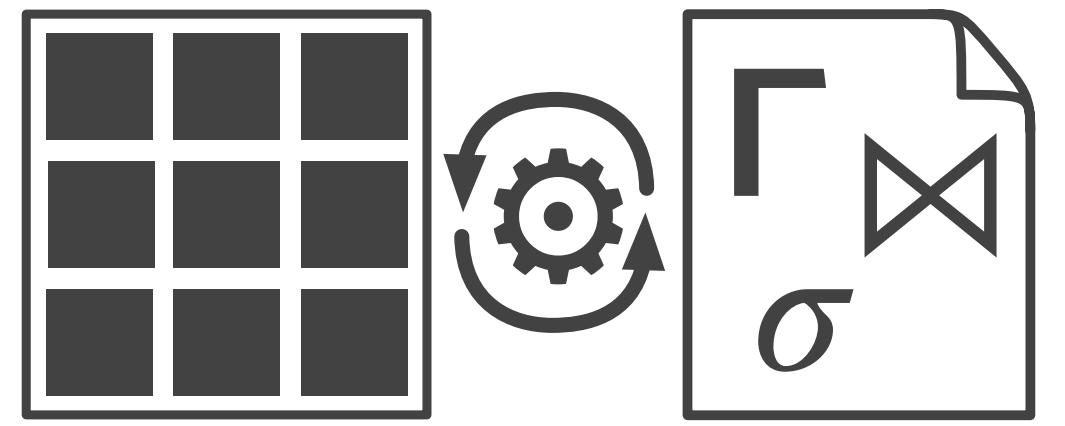
Publication		GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<i>Full query processing systems</i>								
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	Bits
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
GDB	He et al., 2009	Dedicated	Generic	Static	Cost model	Task partitions	Operator	Tuples
CoGaDB	Breß, 2014	Dedicated	Generic	Both	Data locality, cost model	Operator placement	Operator	Columns
SABER	Koliouisis et al., 2016	Dedicated	Generic	Dynamic	Load balancing, cost model	Single partition	Query	Data batch
DB2 BLU	Meraji et al., 2016	Dedicated	Generic	Dynamic	Task nature, load balancing	Task partitions	Operator	Tuples
HetExchange	Chrysogelos et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Pipeline	Data batch
He et al.	He et al., 2014	Integrated	Hybrid	Static	Task nature, cost model	Task partitions	Primitive	Tuples
DIDO	Zhang et al., 2017	Integrated	Hybrid	Hybrid	Task nature, load balancing, cost model	Operator placement	Operator	Query batch
FineStream	Zhang et al., 2020	Integrated	Generic	Dynamic	Cost model	Operator placement	Operator	—
HERO	Karnagel et al., 2017	Both	Generic	Dynamic	Data locality, cost model	Operator placement	Primitive	—
<i>Individual query processing tasks</i>								
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	Algorithm-specific	—	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
G-Grid	Li et al., 2013	—	—	—	—	—	—	—
GAT	Zhang et al., 2017	—	—	—	—	—	—	—
Sioulas et al.	Sioulas et al., 2017	—	—	—	—	—	—	—
Gubner et al.	Gubner et al., 2017	—	—	—	—	—	—	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Operator	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Tuple batch
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	—	Query batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	—	Cuboids, points
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Primitive	Tuples
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	Algorithm-specific	—	—

Integrated GPUs are used as generic compute resources

Publication		GPU integration	Processor usage	Scheduling time	Scheduling strategy	Workload distribution	Task granularity	Data partitioning
<i>Full query processing systems</i>								
Approx. & Refine	Pirk et al., 2014	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	Bits
Stat. coproc.	Heimel et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Mega-KV	Zhang et al., 2015	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
Caldera	Appuswamy et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
Raza et al.	Raza et al., 2020	Dedicated	Specialized	Static	Task nature	Algorithm-specific	Query type	—
GDB	He et al., 2009	Dedicated	Generic	Static	Cost model	Task partitions	Operator	Tuples
CoG								—
SAE								—
DB2								—
HetExchange	Chrysogios et al., 2019	Dedicated	Generic	Hybrid	Load balancing, data locality	Task partitions	Pipeline	Data batch
He et al.	He et al., 2014	Integrated	Hybrid	Static	Task nature, cost model	Task partitions	Primitive	Tuples
DIDO	Zhang et al., 2017	Integrated	Hybrid	Hybrid	Task nature, load balancing, cost model	Operator placement	Operator	Query batch
FineStream	Zhang et al., 2020	Integrated	Generic	Dynamic	Cost model	Operator placement	Operator	—
HERO	Karnagel et al., 2017	Both	Generic	Dynamic	Data locality, cost model	Operator placement	Primitive	—
<i>Individual query processing tasks</i>								
GSS	Bøgh et al., 2013	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
STIG	Doraiswamy et al., 2016	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
HB <sup>+</sup> -tree	Shahvarani et al., 2016	Dedicated	Specialized	Static	Task nature, cost model	Algorithm-specific	—	—
Stehle et al.	Stehle et al., 2017	Dedicated	Specialized	Static	Task nature	Algorithm-specific	—	—
G-G								—
GAT								—
Siou								—
Gubler et al.	Gubler et al., 2019	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Key batch
SCCG	Wang et al., 2012	Dedicated	Hybrid	Dynamic	Task nature, load balancing	Task partitions	Operator	Polygon pairs
Lutz et al.	Lutz et al., 2020	Dedicated	Hybrid	Hybrid	Task nature, load balancing	Task partitions	Operator	Tuple batch
Beier et al.	Beier et al., 2012	Dedicated	Generic	Dynamic	Cost model	Single partition	—	Query batch
Bøgh et al.	Bøgh et al., 2017	Dedicated	Generic	Dynamic	Load balancing	Single partition	—	Cuboids, points
He et al.	He et al., 2013	Integrated	Generic	Static	Cost model	Task partitions	Primitive	Tuples
HELLS join	Karnagel et al., 2013	Integrated	Specialized	Static	Task nature	Algorithm-specific	—	—

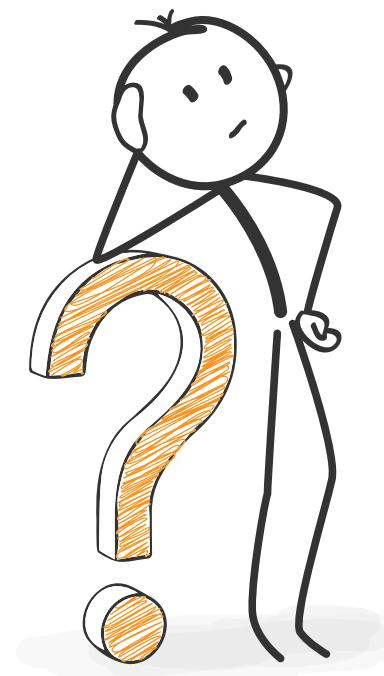
**Dedicated GPUs should perform specialized coarse-grained tasks.**

**Integrated GPUs should cooperate closely with CPU.**

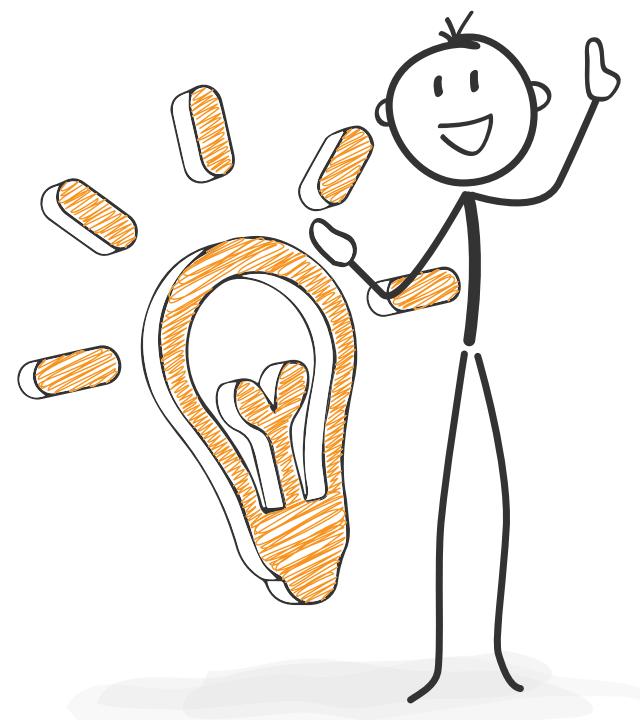


# Operator Variant Tuning on Heterogeneous Processors

# Processor sensitivity



How sensitive are processors to operator implementation details?

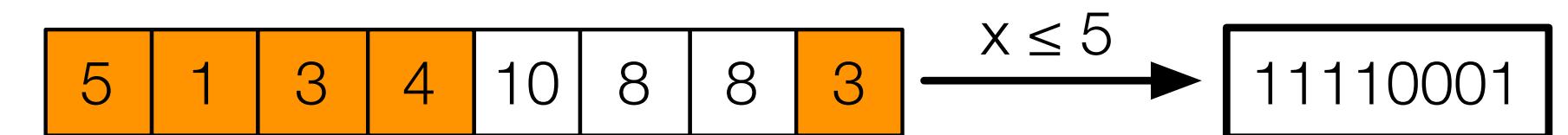


Performance analysis of selection and hash aggregation operators

Multiple CPUs, GPUs + Intel Xeon Phi coprocessor

# Selection variants

Operation



Basic selection kernels

SEQUENTIAL, GLOBALATOMIC, LOCALATOMIC,  
REDUCE, COLLECT, TRANSPOSE

Variants

Thread configuration

Low-level implementation parameters

Processors

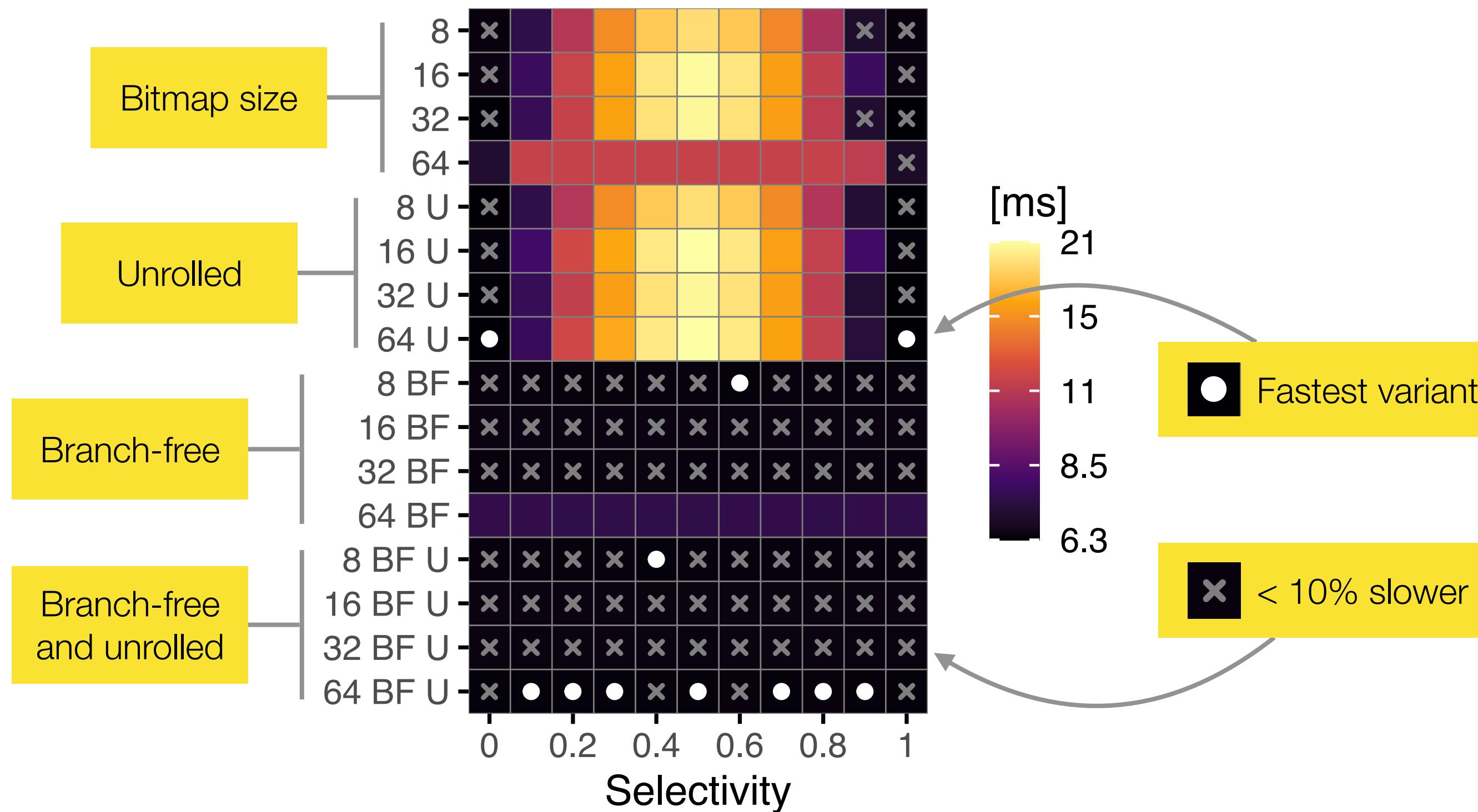
7 CPUs, 5 GPUs, Xeon Phi

AMD, IBM, Intel, Nvidia

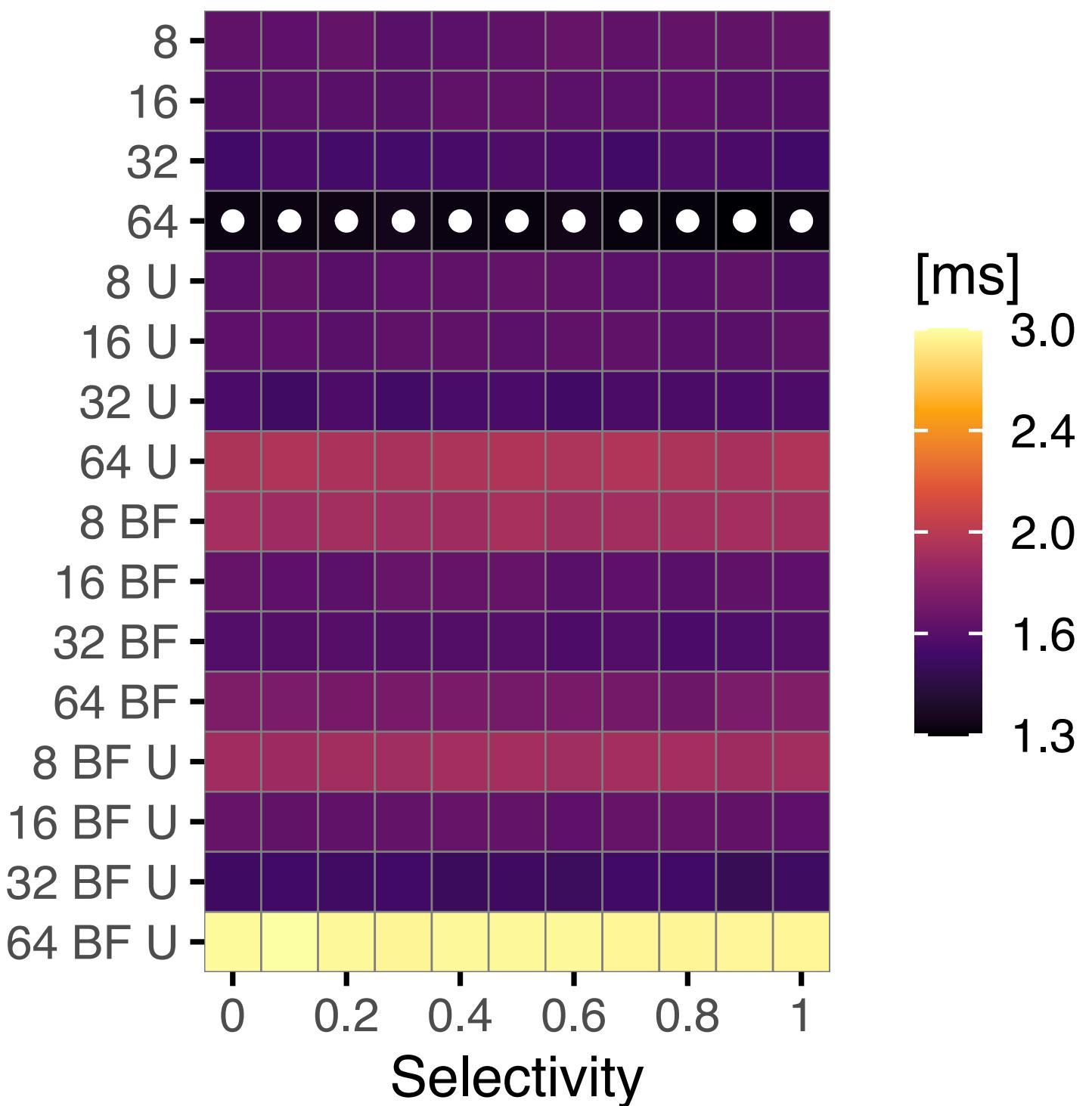
# Selection variants performance

Selection on 32 million random integers, uniform distribution, median of 10 repetitions

**SEQUENTIAL** on Intel Core i7-4900MQ



**TRANSPOSE** on Intel Xeon Phi 7120



No single variant performs best on every processor.

# Competitive variants

At most 10% slower than the fastest variant

Processor	Number of variants	Competitive variants	Percentage	Maximum slowdown
Intel Xeon E5620	5880	1370	23 %	32
Nvidia Tesla K40m	4696	129	2.7 %	136
Intel Xeon Phi 7120	3886	6	0.15 %	147

Some processors are very sensitive to implementation details.

# Hash aggregation variants

Parallelization strategy

SHARED, INDEPENDENT, WORKGROUPLOCAL

Variants

Thread configuration

Processors

6 GPUs

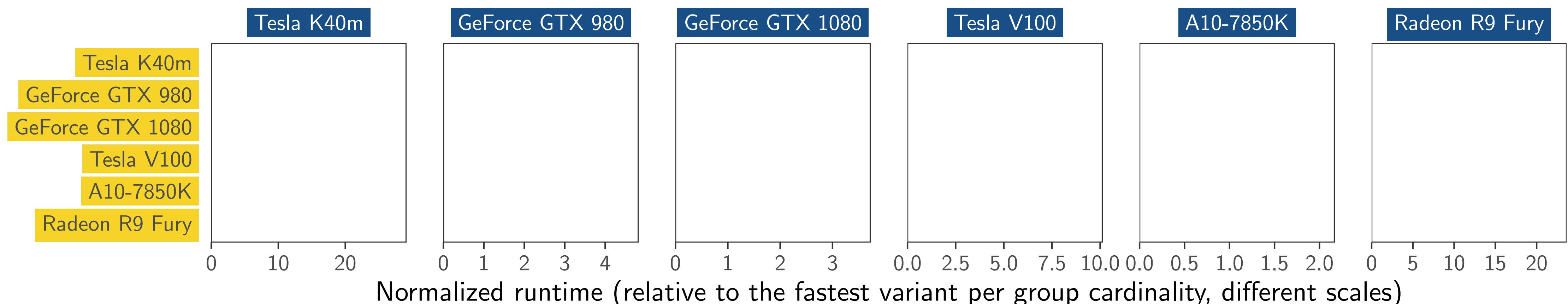
AMD, Nvidia

Different micro architectures

# Influence of thread configuration

128 million input rows, 32-bit keys and values, Sum aggregation, group cardinality between 1 and  $2^{28}$

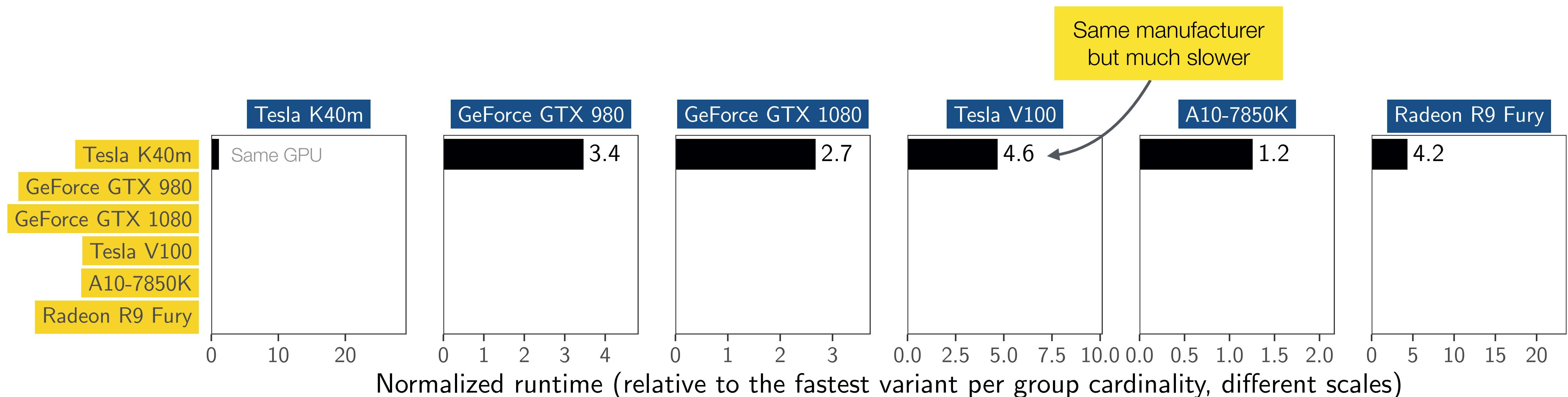
Performance penalty when a thread configuration optimized for a specific GPU (rows) is executed on another GPU (columns).



# Influence of thread configuration

128 million input rows, 32-bit keys and values, Sum aggregation, group cardinality between 1 and 2<sup>28</sup>

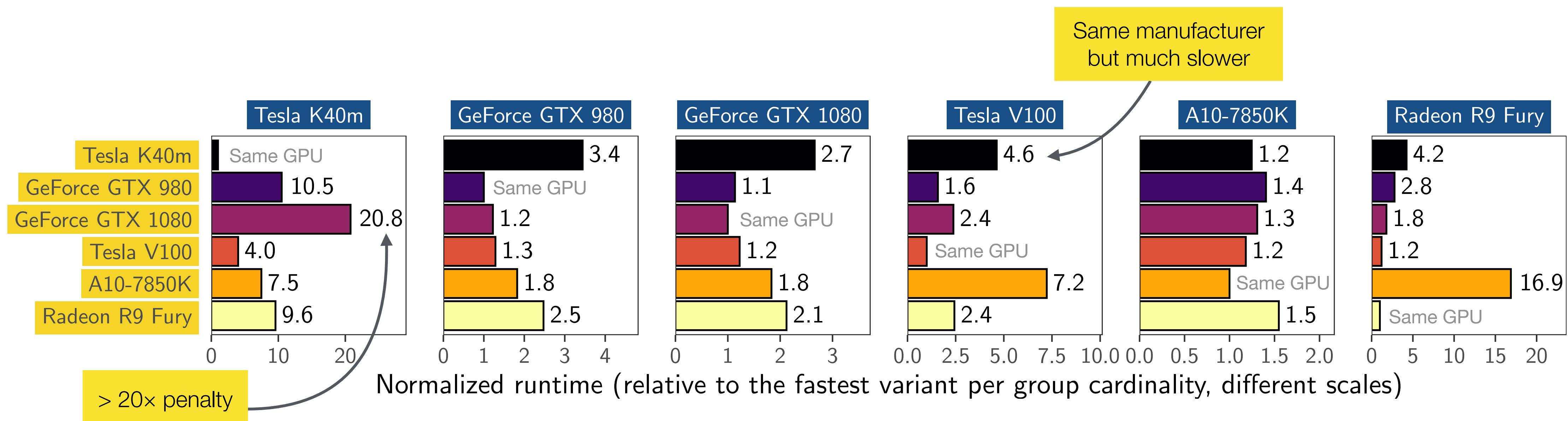
Performance penalty when a thread configuration optimized for a specific GPU (rows) is executed on another GPU (columns).



# Influence of thread configuration

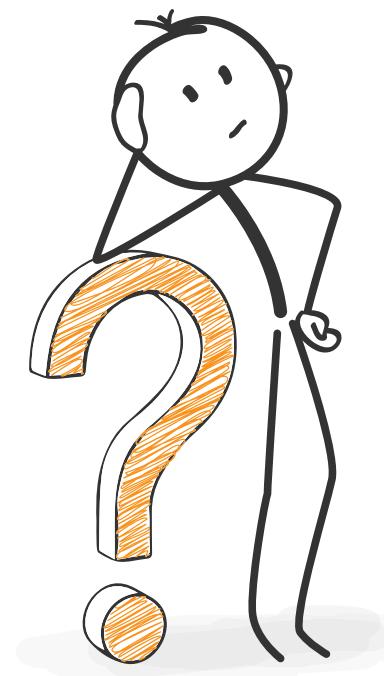
128 million input rows, 32-bit keys and values, Sum aggregation, group cardinality between 1 and 2<sup>28</sup>

Performance penalty when a thread configuration optimized for a specific GPU (rows) is executed on another GPU (columns).

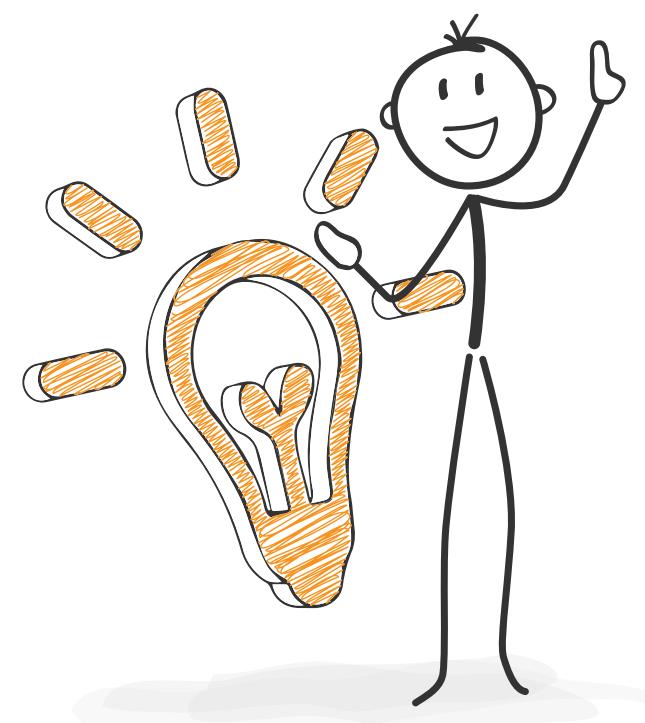


The fastest thread configuration is highly GPU-specific.

# Variant tuning



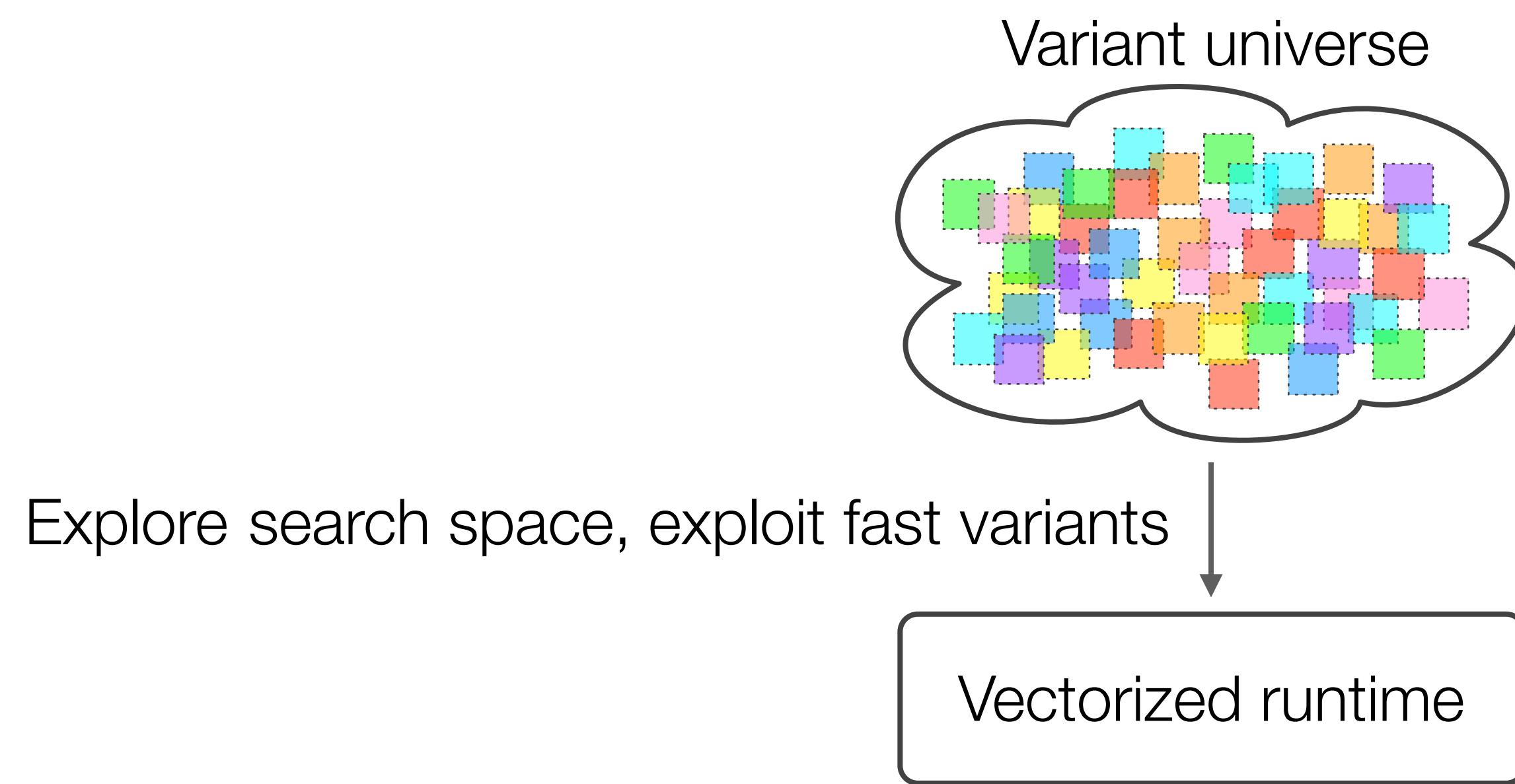
Let the database find a fast operator implementation automatically



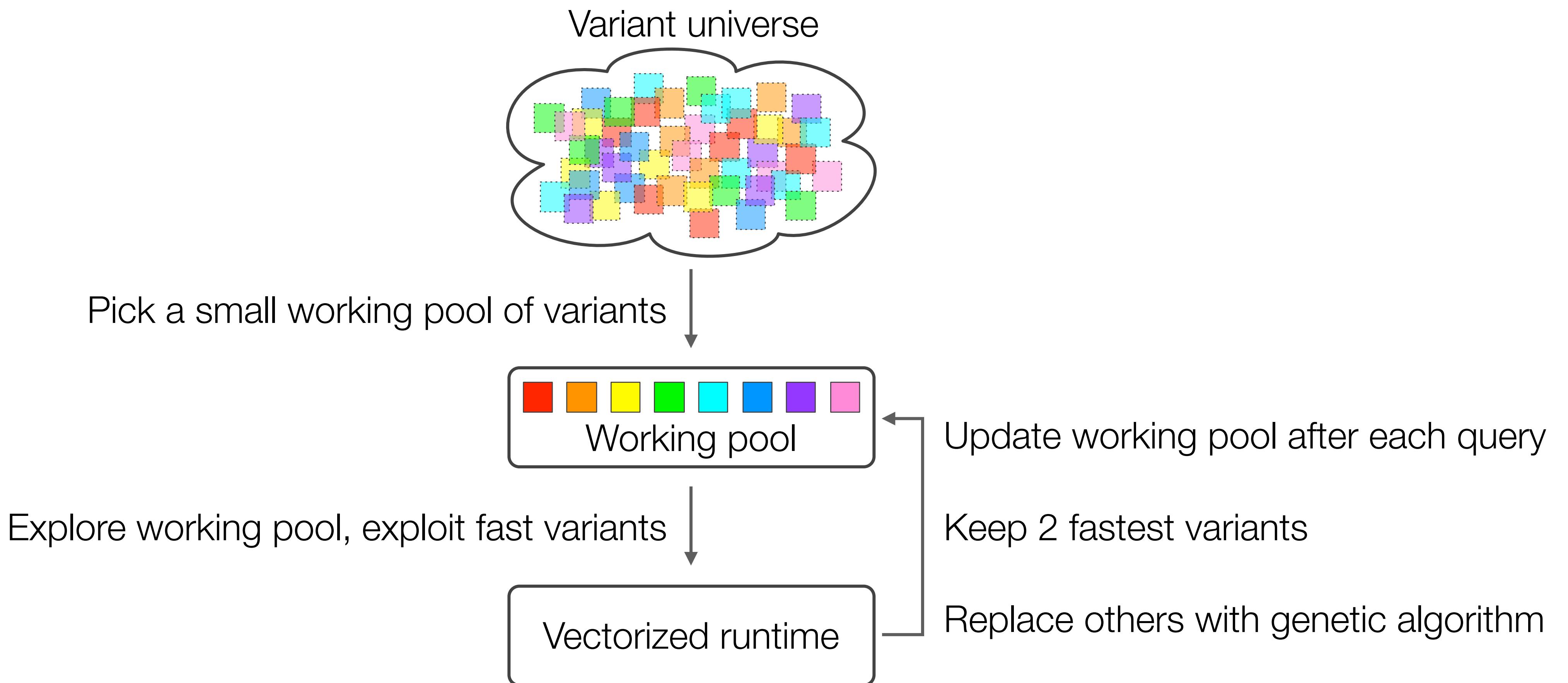
Extend micro adaptivity to handle large search spaces

Extend local search to handle runtime variation

# Candidate selection in large variant spaces

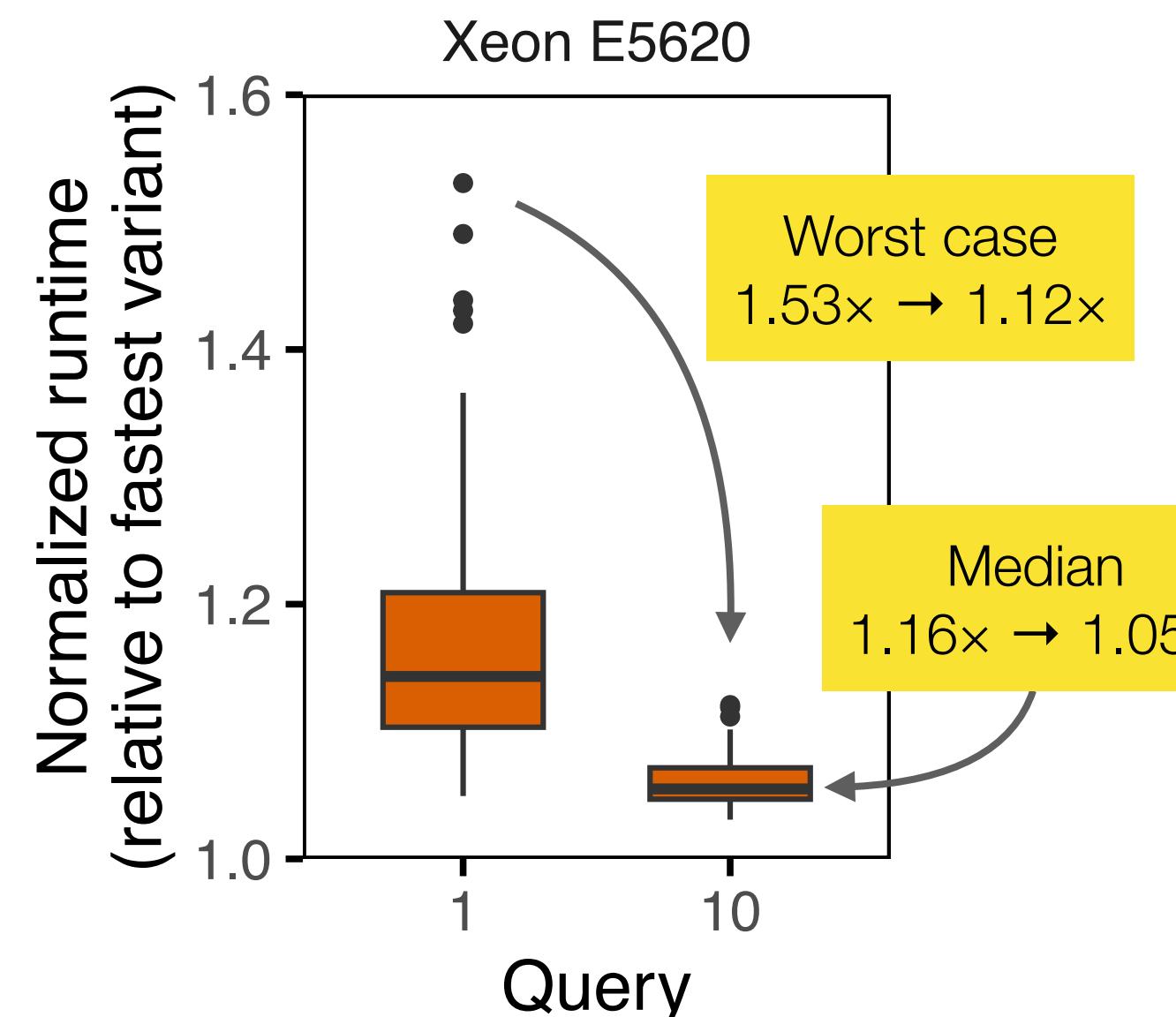


# Candidate selection in large variant spaces

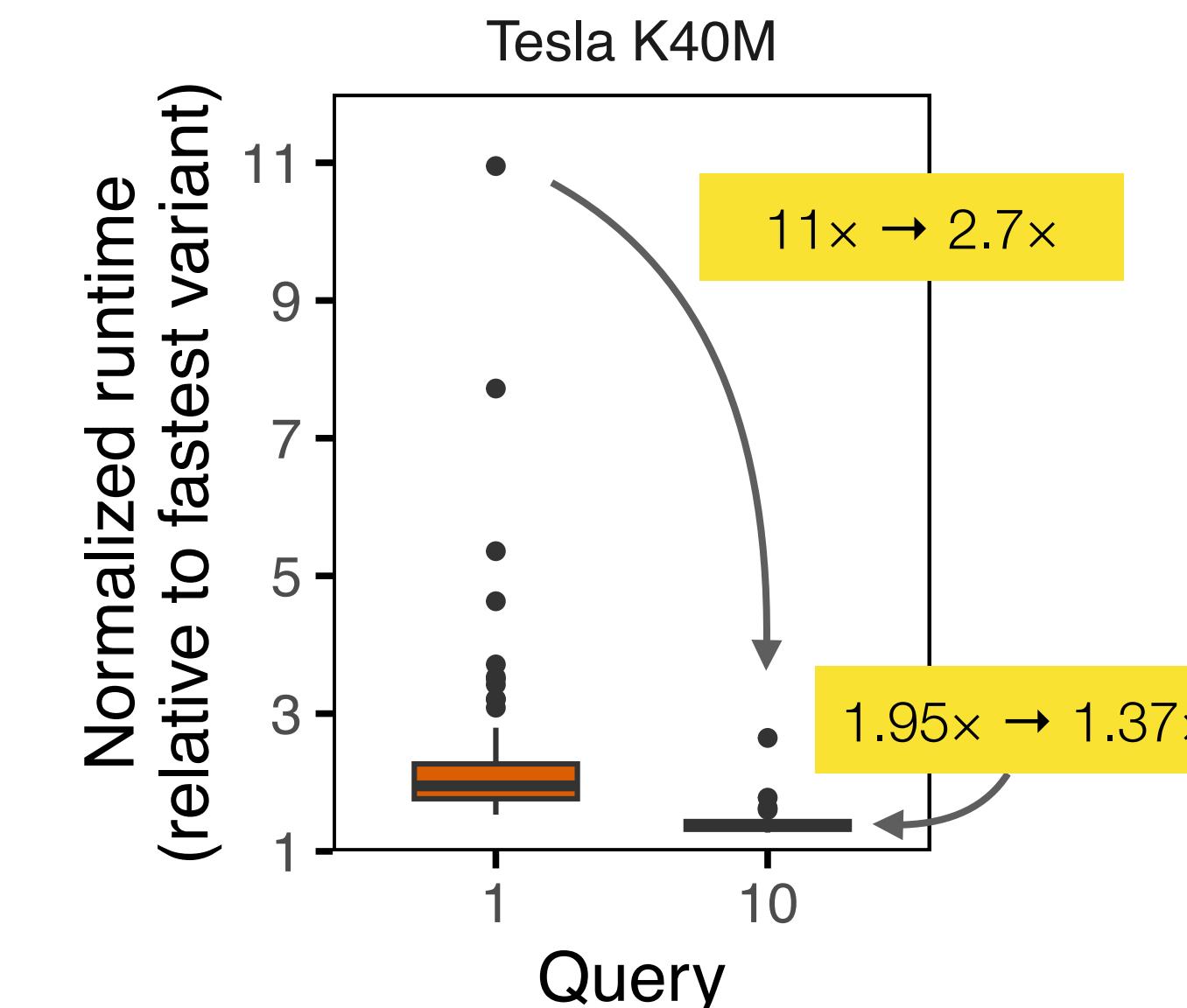


# Evaluation

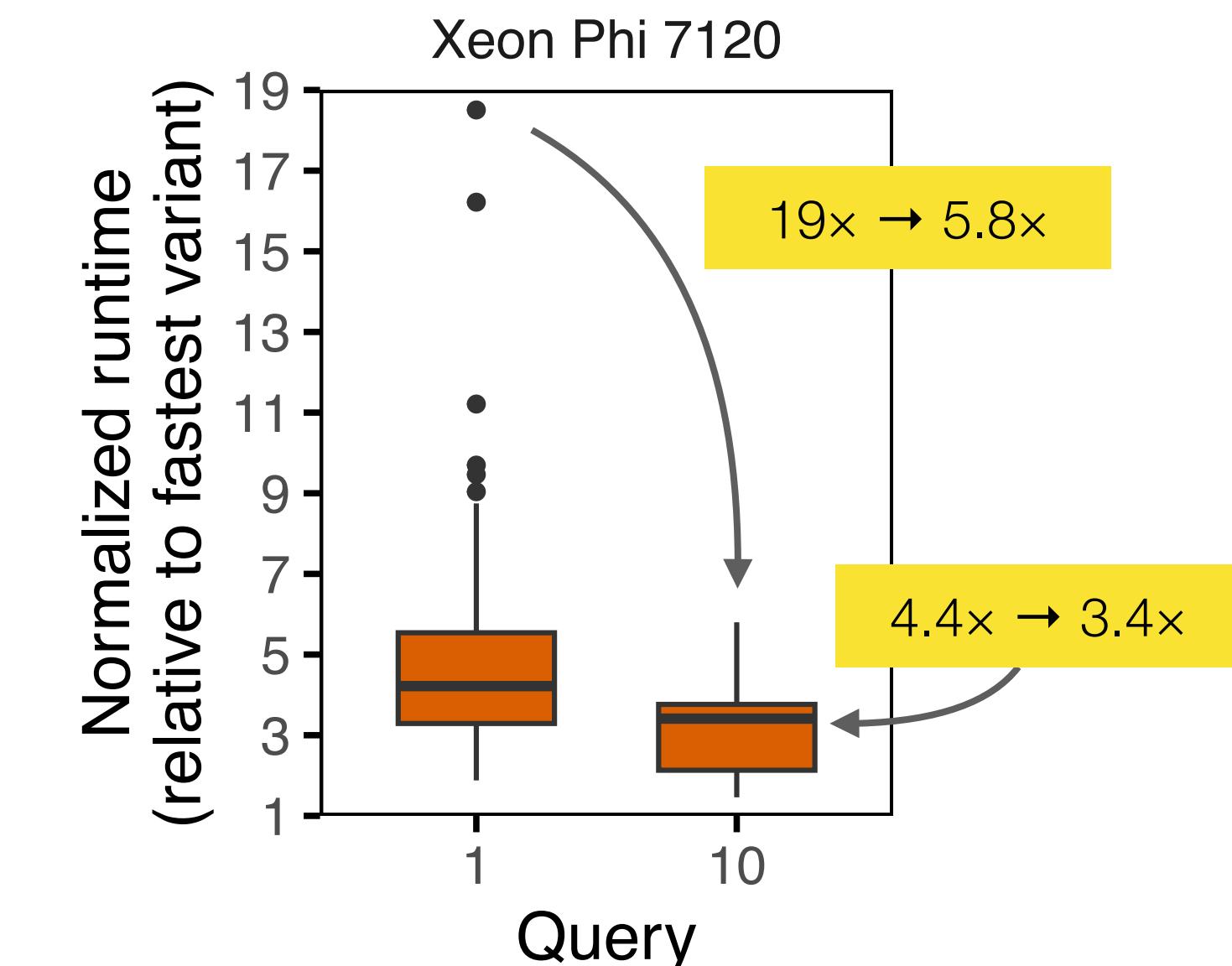
Selection variants, selectivity 0.5, 1024 blocks/query, 8 variants in pool, genetic algorithm after each query, 100 reps



23% competitive variants  
 $\rightarrow$  88% chance that initial pool contains a competitive variant



2.7% competitive variants  
 $\rightarrow$  20% chance

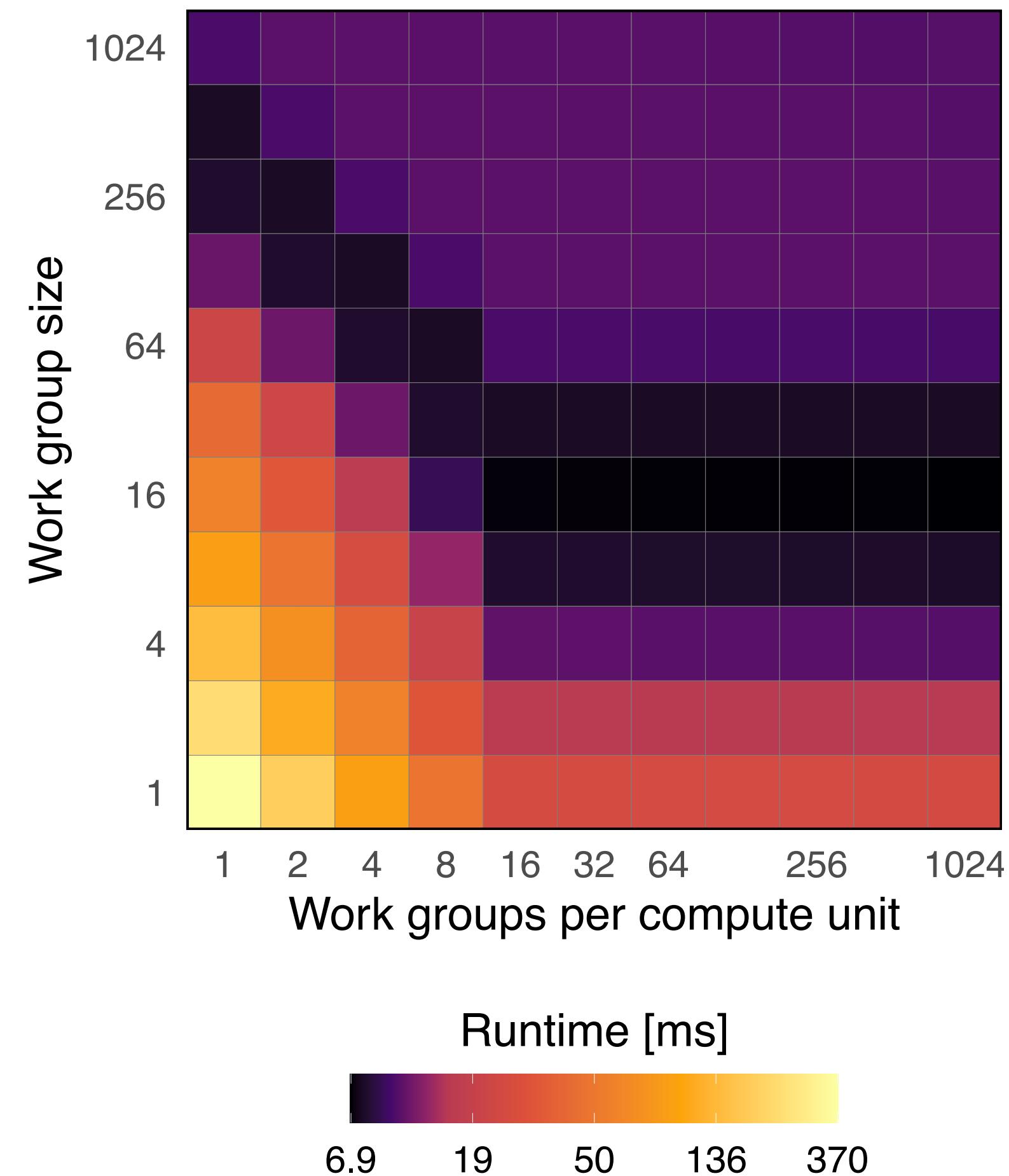


0.15% competitive variants  
 $\rightarrow$  1.2% chance

**Variant tuning performance depends on initial working pool and number of competitive variants.**

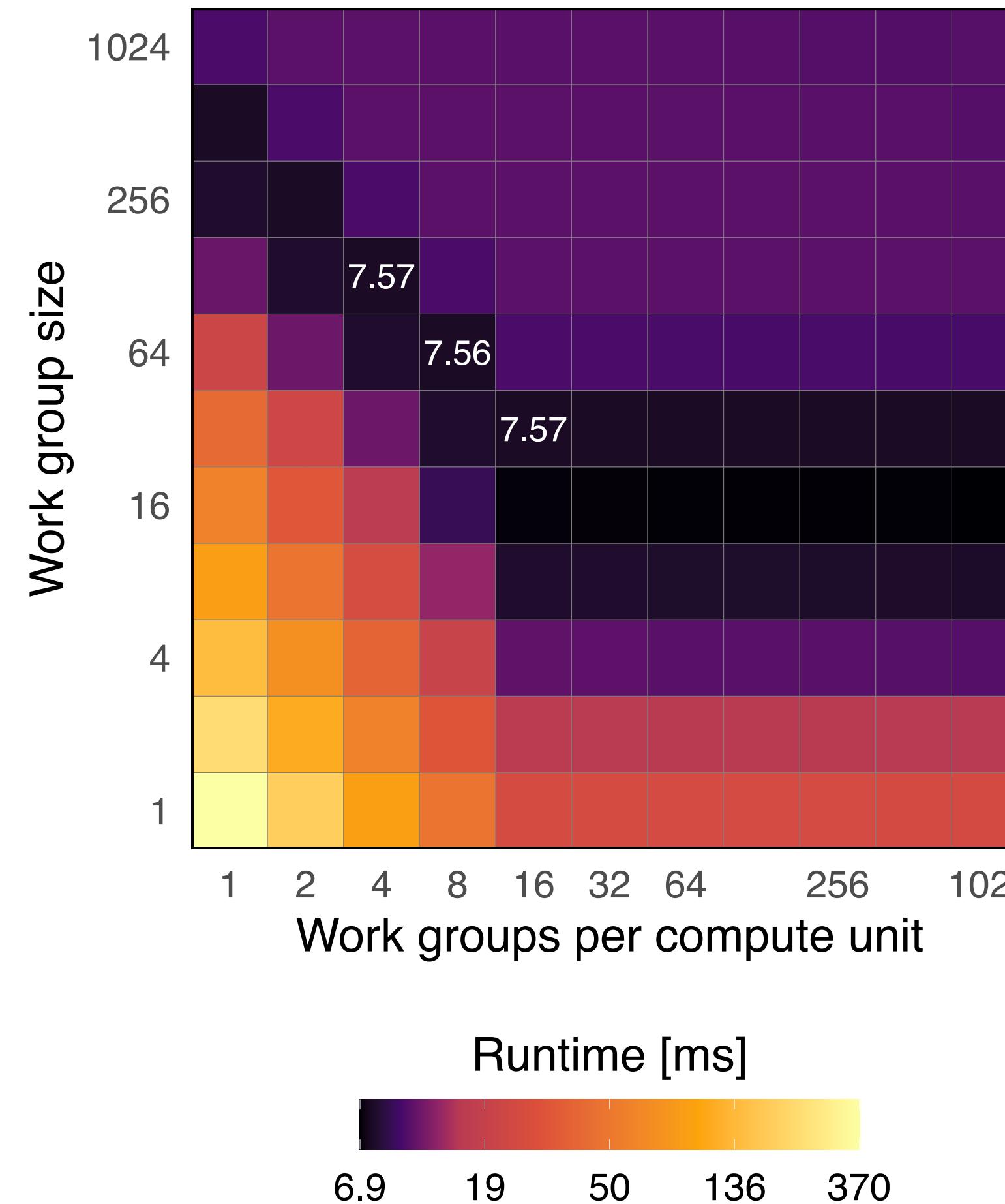
# Performance plateaus

Shared aggregation, Tesla K40m, 8 million groups, 128 million input rows, 32-bit keys and values



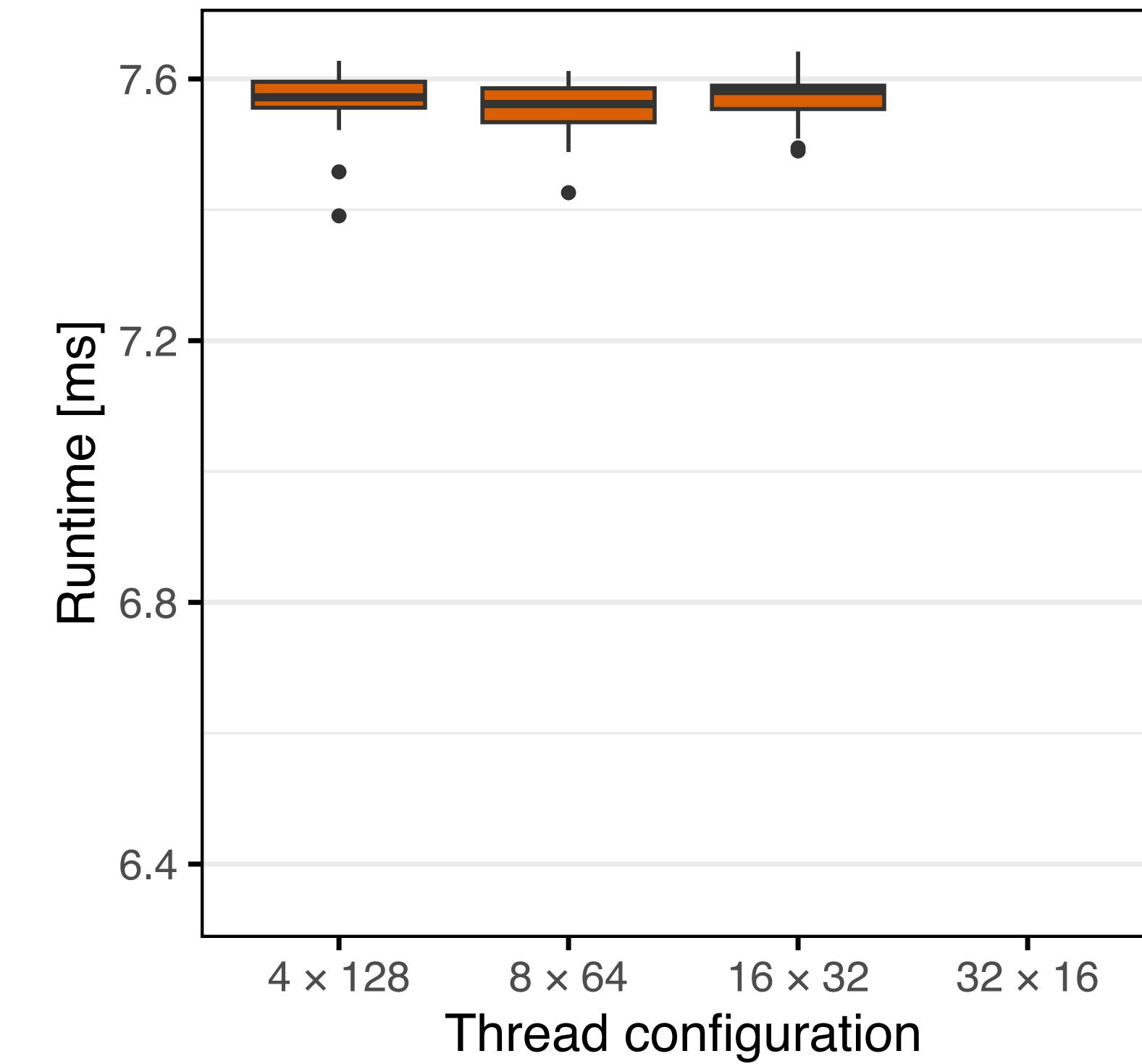
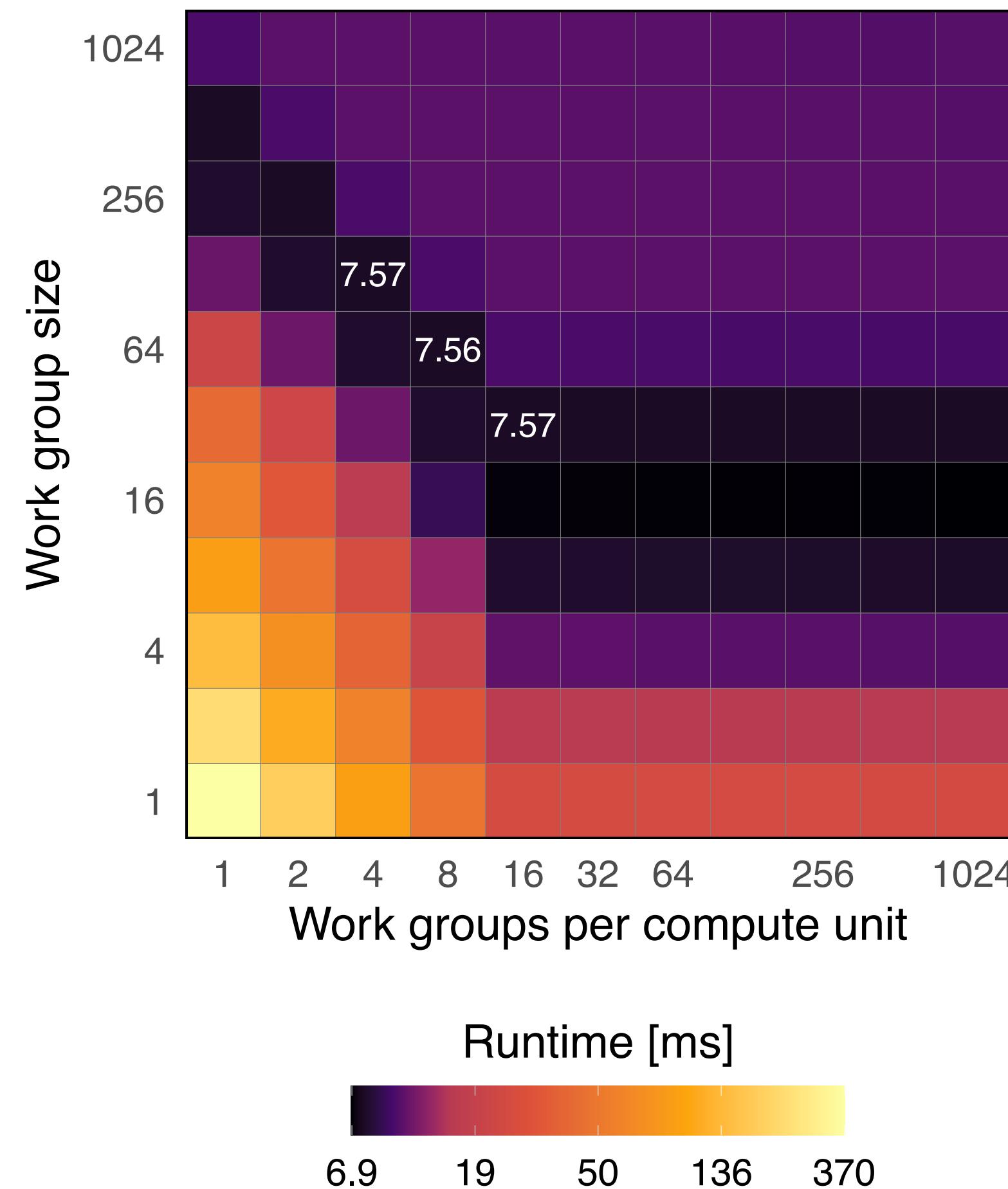
# Performance plateaus

Shared aggregation, Tesla K40m, 8 million groups, 128 million input rows, 32-bit keys and values



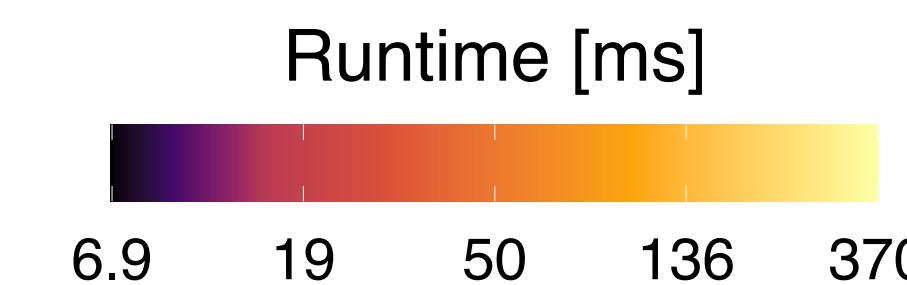
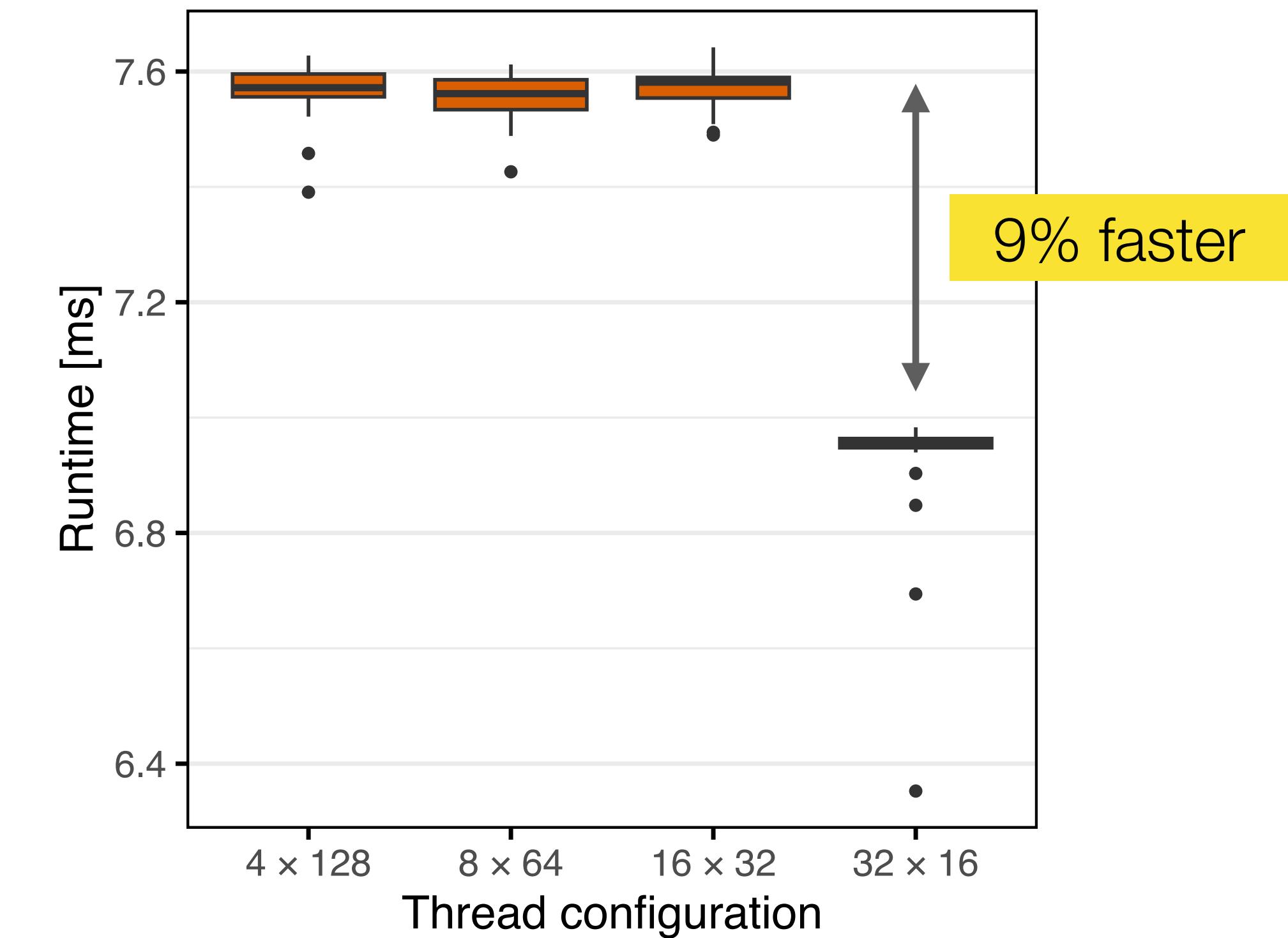
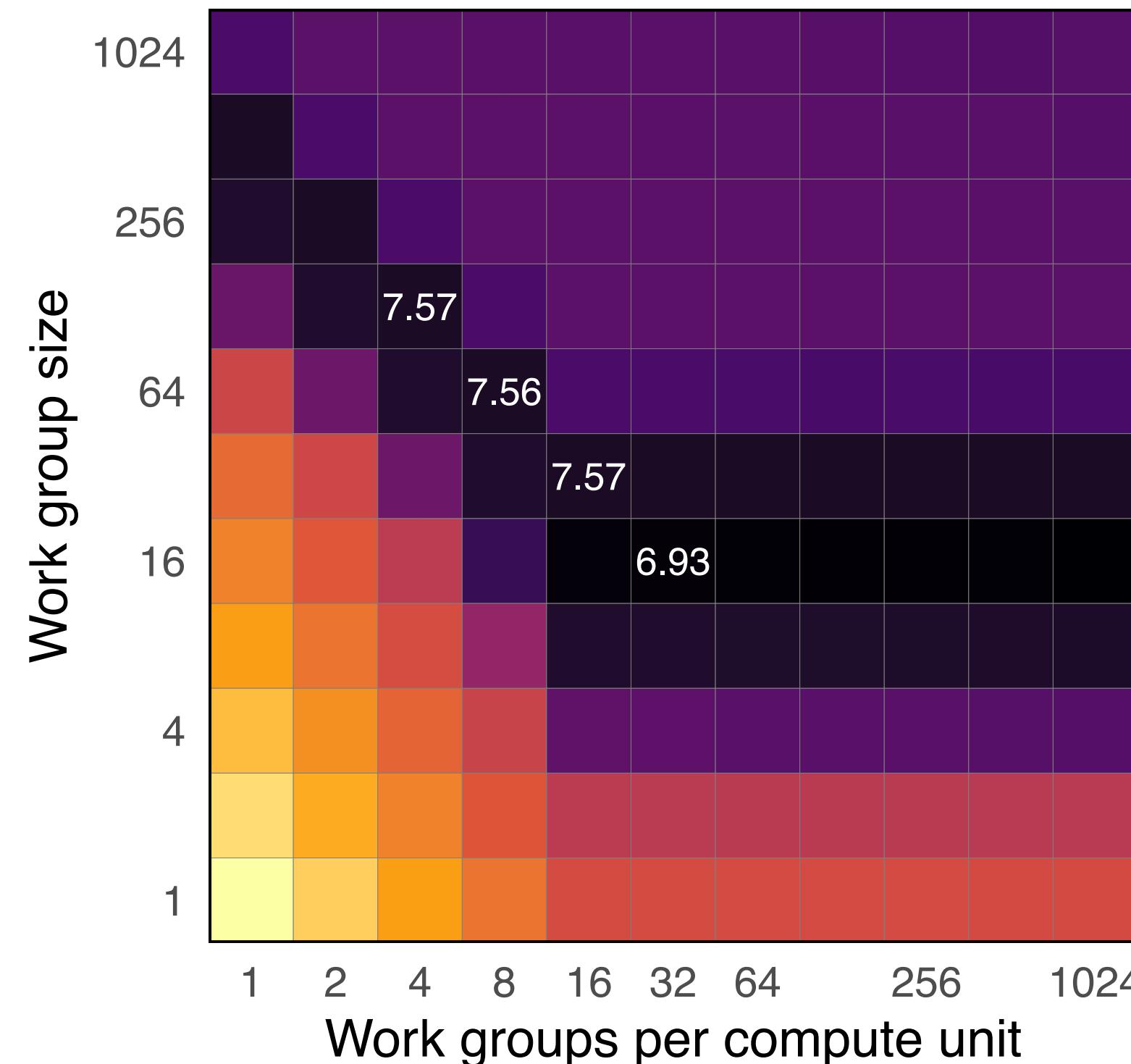
# Performance plateaus

Shared aggregation, Tesla K40m, 8 million groups, 128 million input rows, 32-bit keys and values

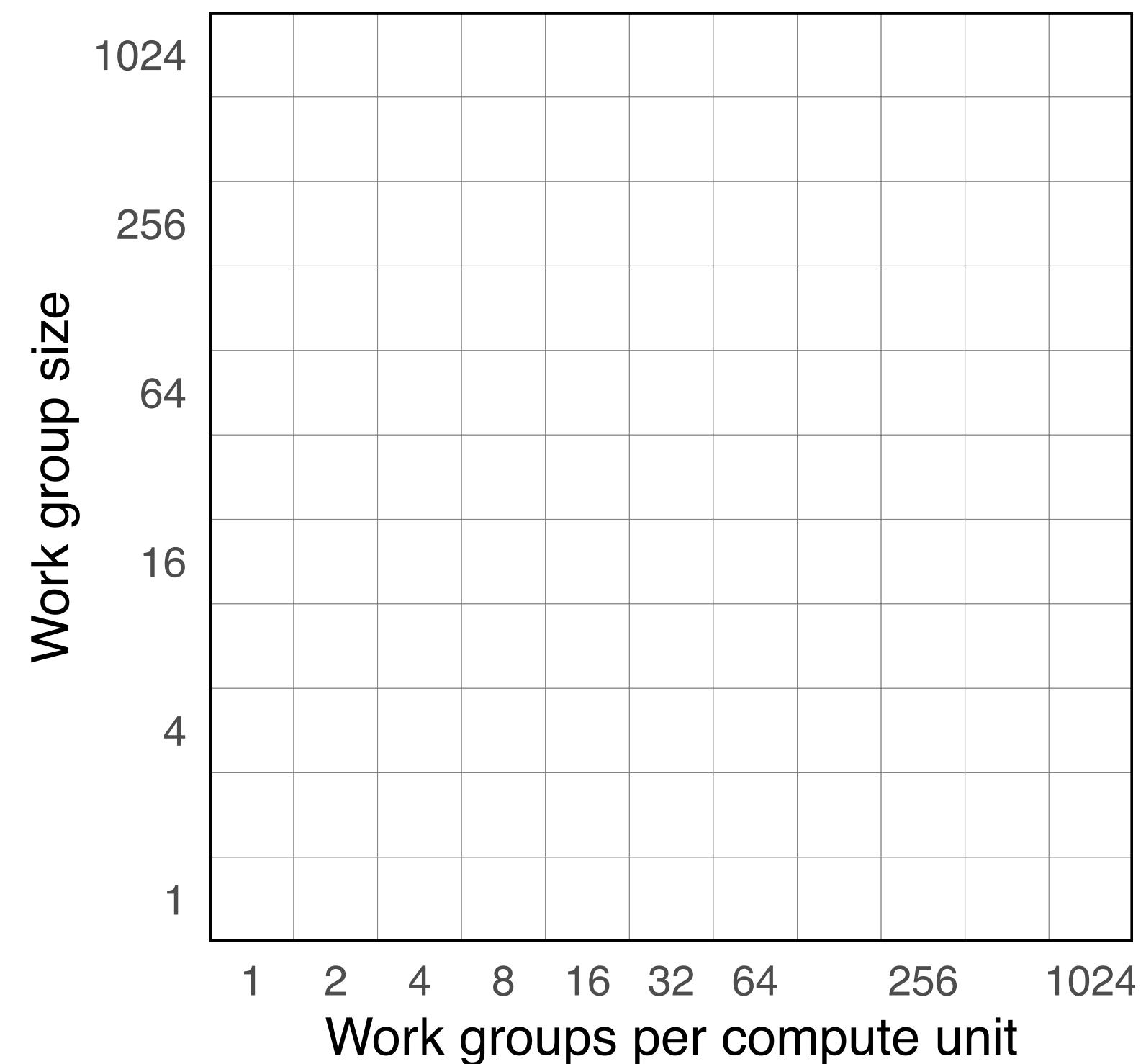


# Performance plateaus

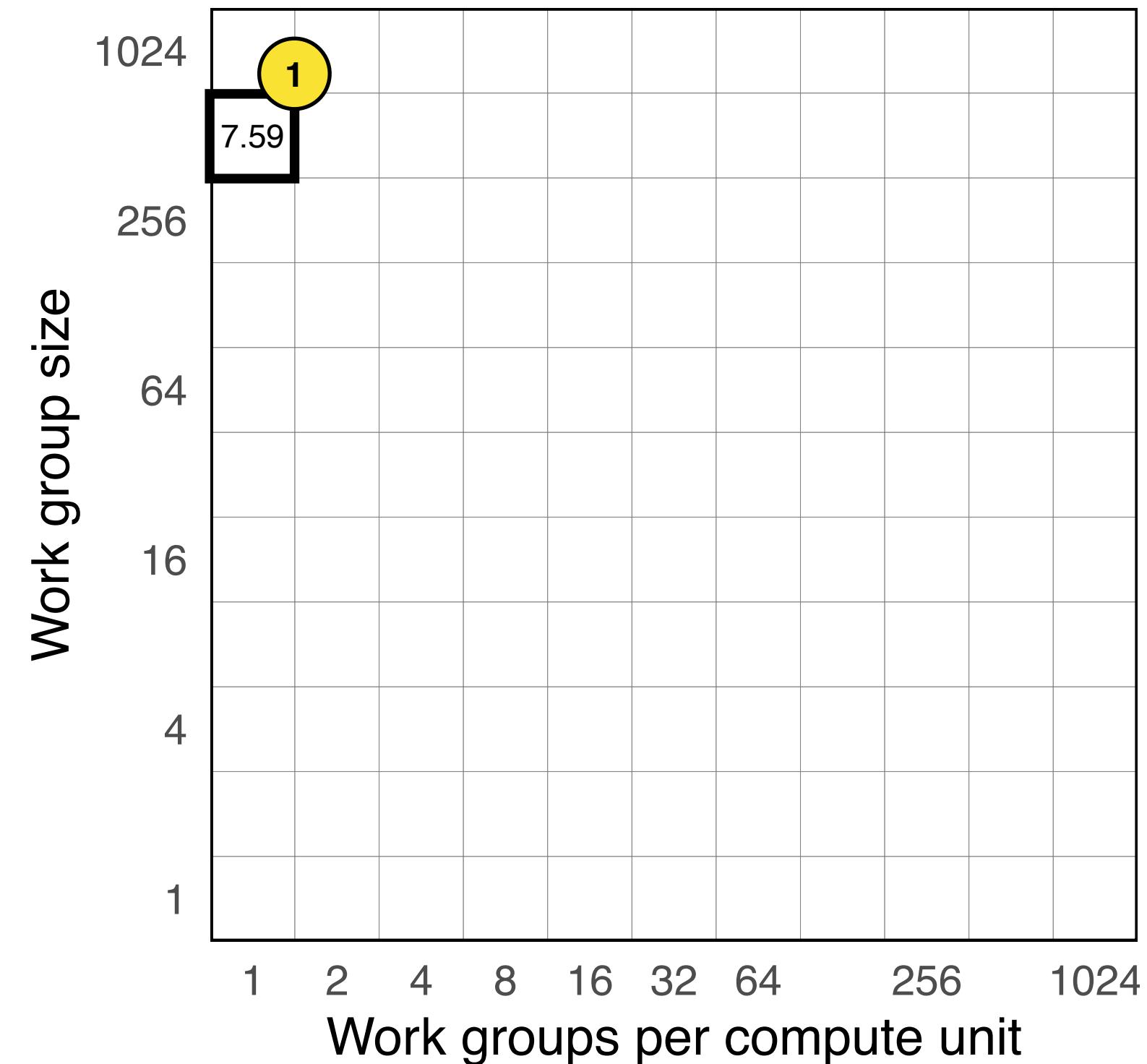
Shared aggregation, Tesla K40m, 8 million groups, 128 million input rows, 32-bit keys and values



# Runtime variation-aware local search

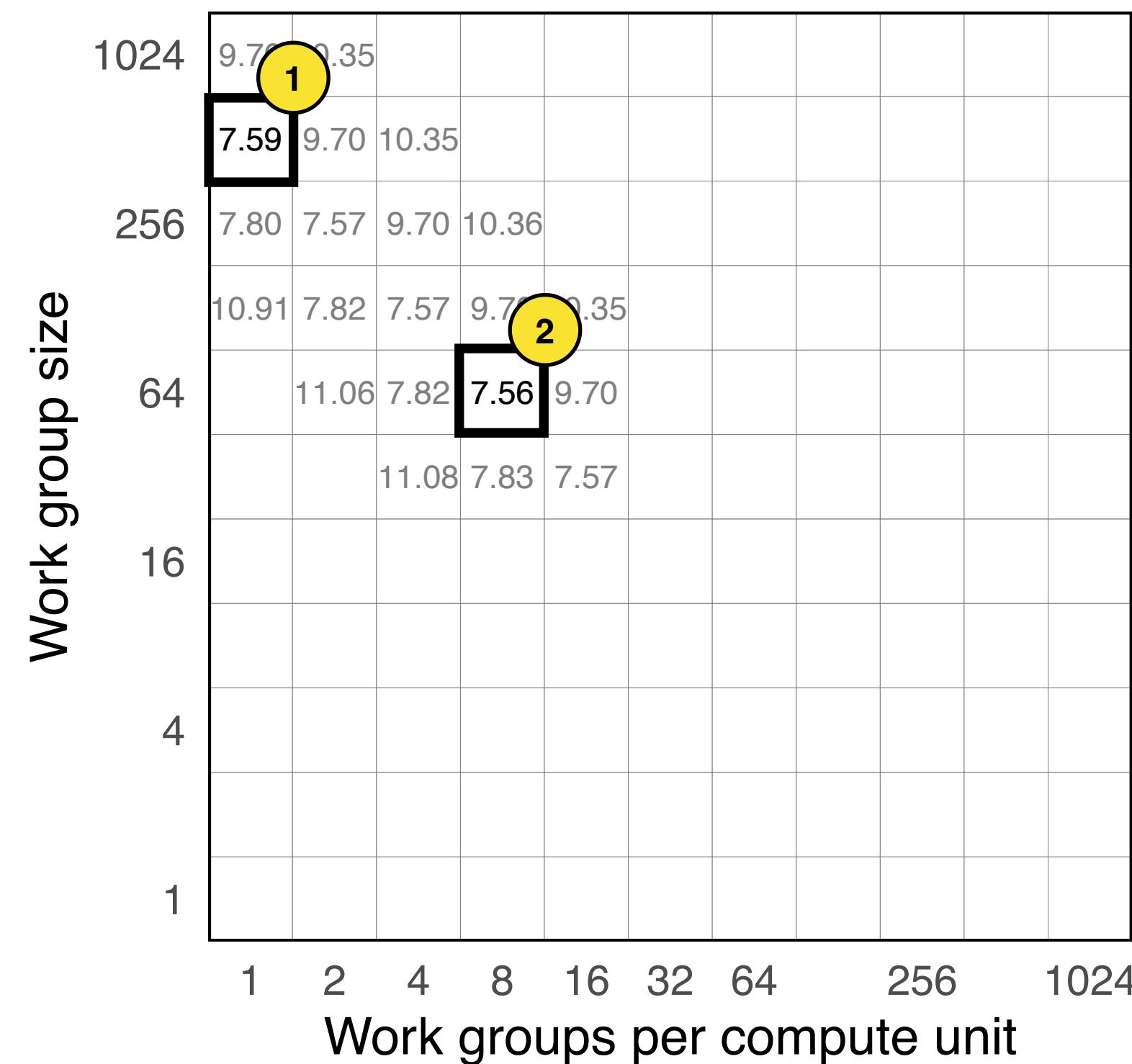


# Runtime variation-aware local search



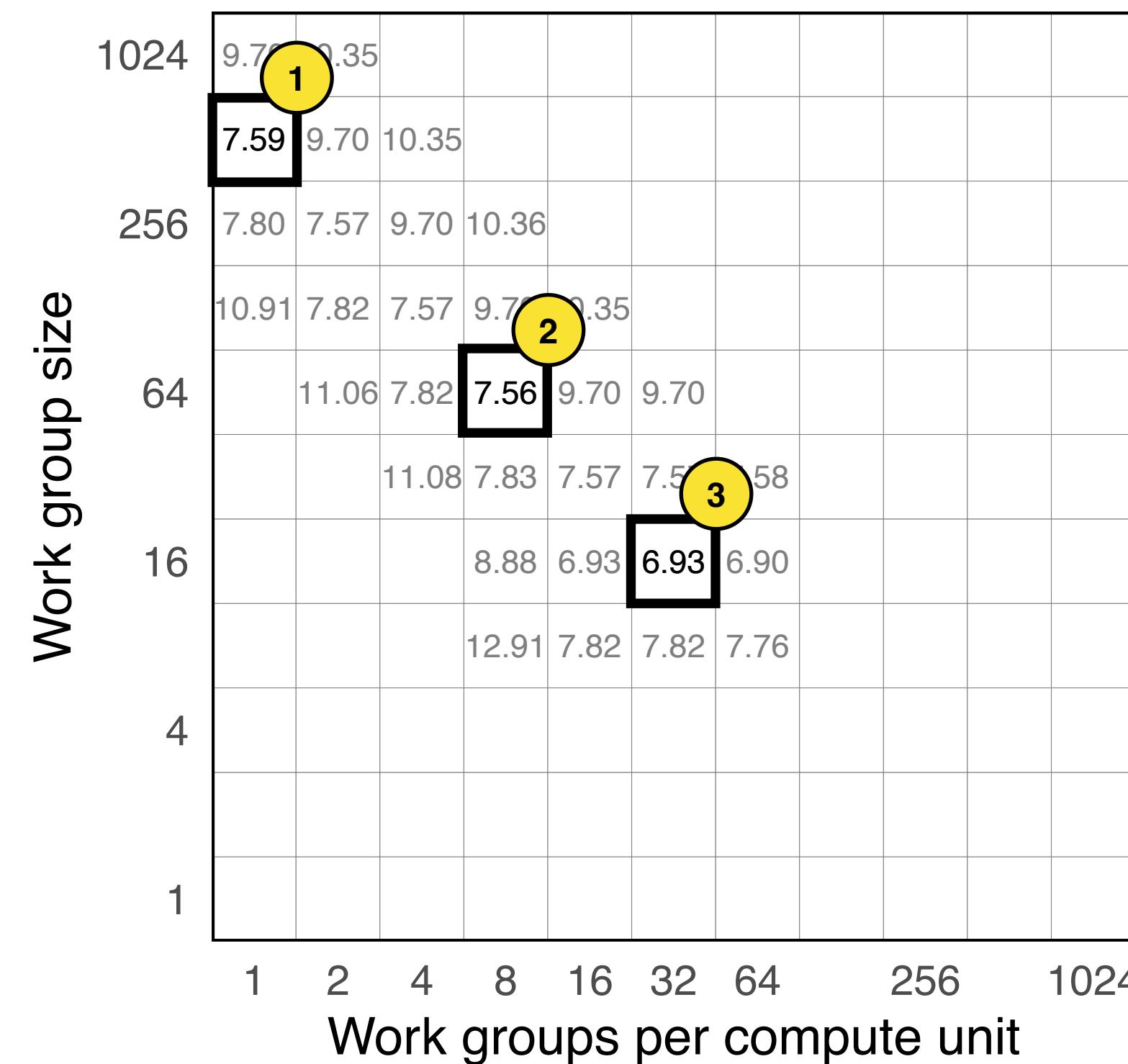
- ① Start with initial thread configuration

# Runtime variation-aware local search



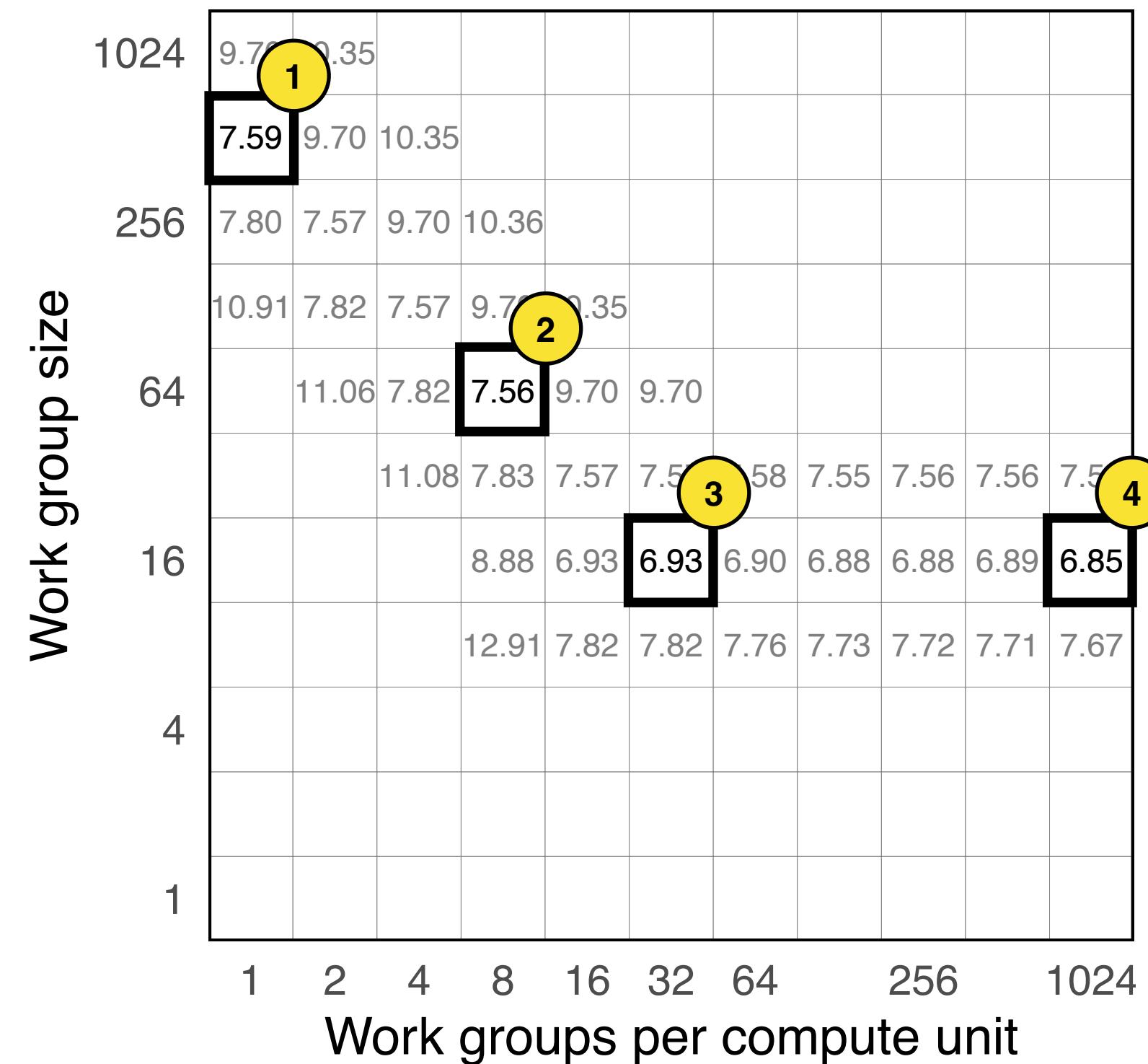
- ① Start with initial thread configuration
- ② Follow gradient in search space

# Runtime variation-aware local search



- ① Start with initial thread configuration
- ② Follow gradient in search space
- ③ Push past performance plateaus

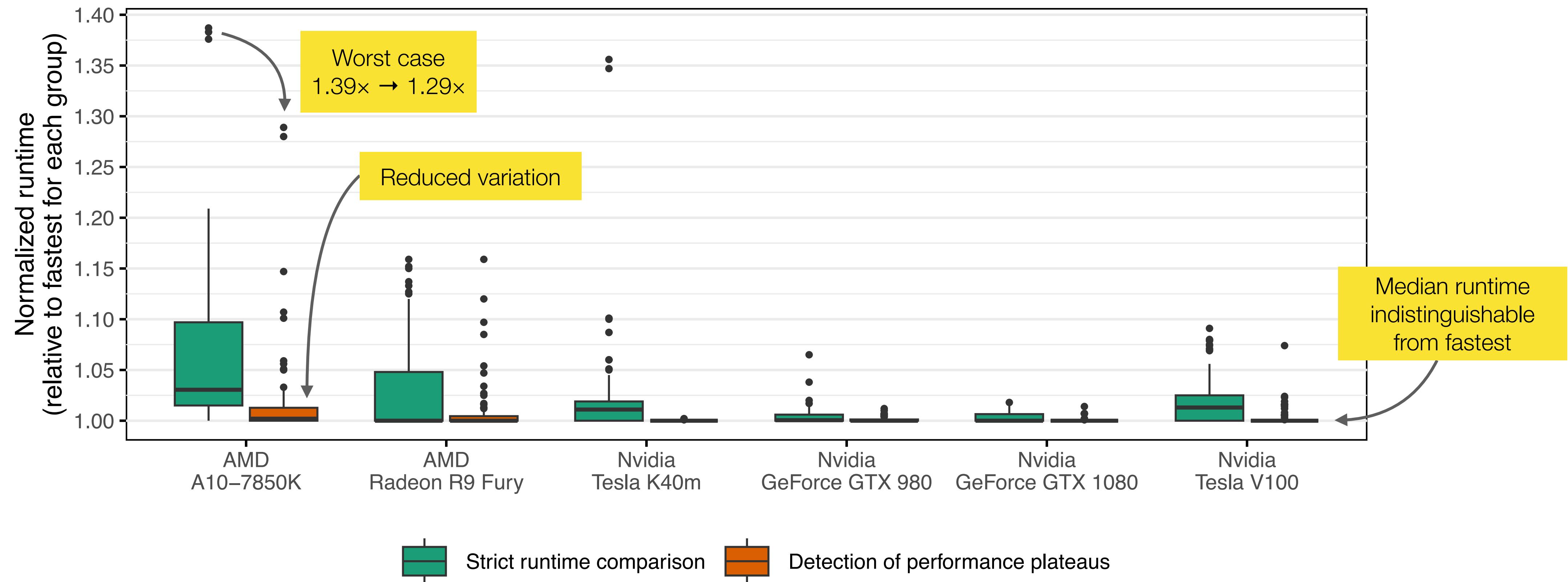
# Runtime variation-aware local search



- ① Start with initial thread configuration
- ② Follow gradient in search space
- ③ Push past performance plateaus
- ④ Stop at minimum

# Evaluation

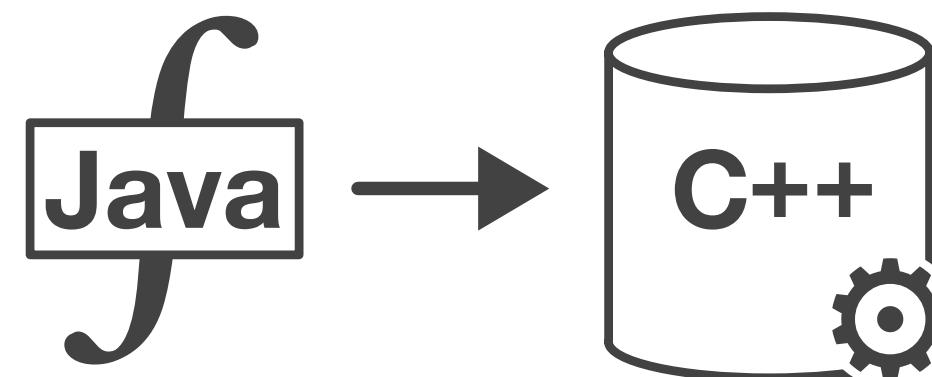
Hash aggregation variants, 128 million input rows, 32-bit keys and values, 3 samples per group cardinality



Exploiting processor and operator characteristics finds fast variants.

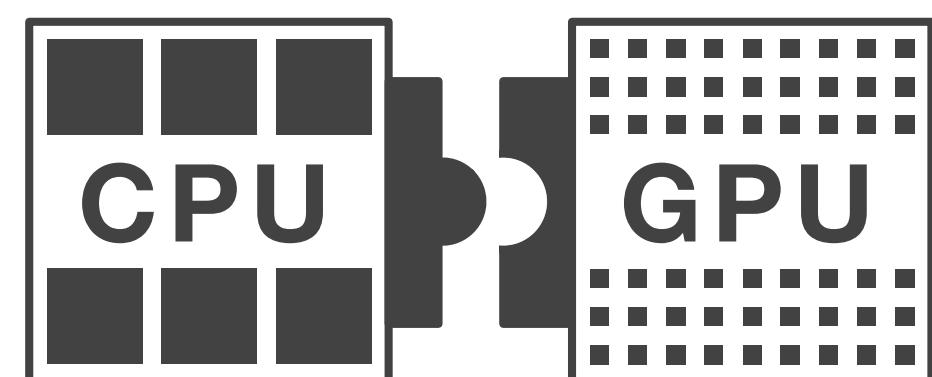
# Conclusion

# Summary of contributions



## Processing Java UDFs in a C++ environment

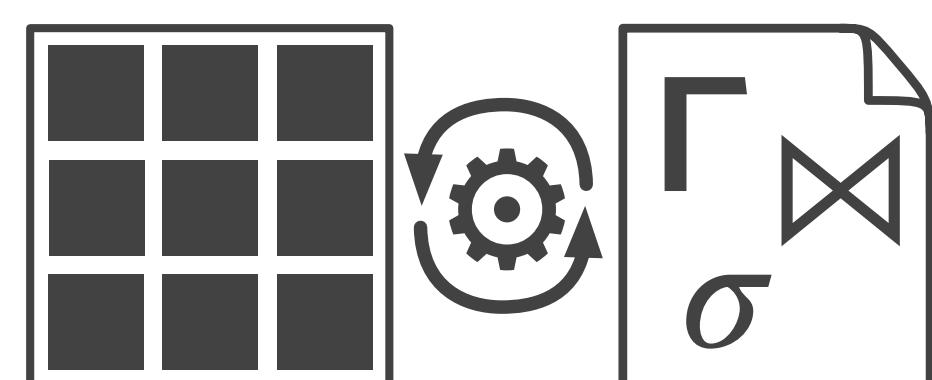
Overcome the JNI overhead when executing Java UDFs



## Query Processing on Heterogeneous CPU/GPU Systems

Survey of query processing systems and individual query processing tasks

Classification scheme for workload distribution on heterogeneous processors



## Operator Variant Tuning on Heterogeneous Processors

Analysis of selection and hash aggregation on heterogeneous processors

Let the database find fast operator implementations automatically