

MATH240: Discrete Structures

Matthew He

December 3, 2024

Abstract

1 Predicate Logic

Definition 1 (Predicate). A predicate is a statement containing some number of variable coming from a universe u .

1.1 Tutorial

1. $p \Rightarrow p$: tautology, since $\neg p \vee p = 1$
2. $(p \Rightarrow q) \Rightarrow p$: contingency, true if $p = q = 1$, false if $p = 0$.
3. $(\neg(p \wedge \neg q)) \vee p \equiv 1$: tautology.
4. $(p \Leftrightarrow q) \wedge (p \Leftrightarrow \neg q)$: contradictory.
5. to be filled

When $p = 0$, $p \Rightarrow q$ is always true, and a true \Rightarrow false(p) is always false.

Example 1

Prove that a logical formula is satisfiable iff. its negation is falsifiable.

Proof. A logical formula is satisfiable iff there is an assignment of all the variables which makes the formula true. By definition of the negation, this assignment makes the negation of our formula false, which means the negation is falsifiable. Proof of the converse is analogous.

Prove or disprove:

1. $\forall n \in \mathbb{N}, \exists m \in \mathbb{N}, n + m = 0$
False. Let's prove its negation: $\exists n \in \mathbb{N}, \forall m \in \mathbb{N}, n + m \neq 0$, it's true, e.g., we can let $n = 2$.
2. $\forall n \in \mathbb{N}, \exists m \in \mathbb{Z}, n + m = 0$
True, Let $n \in \mathbb{N}$, then choose $m = -n \in \mathbb{Z}$, and we get $m + n = 0$.
3. $\forall n \in \mathbb{N}, \exists k \in \mathbb{N}, (k \geq m) \Rightarrow (k \geq 5n)$:
The statement is equivalent to $(k < m) \vee (k \geq 5n)$.
For any $n \in \mathbb{N}$, choose $m = 5n$, then for all $k \in \mathbb{N}$, we have that statement is true.

1.2 Proving statement of the form $p \Rightarrow q$

We assume p is true, prove q , which is showing the case that p is true but q is false can't happen.

Example 2

prop. If n is an odd integer, then n^2 is an odd integer.

Write the proposition in predicate formula:

$$u = \mathbb{Z}, \forall n : ((\exists k : n = 2k + 1) \Rightarrow (\exists l : n^2 = 2l + 1))$$

proof. Let n be an integer. Assume that n is odd, that is, there exist k such that $n = 2k + 1$.

Then, $n^2 = (2k + 1)^2 = 2(2k^2 + 2k) + 1$. Let $l = 2k^2 + 2k$, then $n^2 = 2l + 1$, thus, it's odd.

1.3 *To disprove a statement: prove its negation is true***Example 3**

disprove $\exists x \forall y : x + y \neq 0$.

prove $\forall x \exists y : x + y = 0 \equiv \neg(\exists x \forall y : x + y \neq 0)$

proof. Let $x \in \mathbb{R}$ be given. Pick $y = -x$, then $x + y = 0$. Q.E.D

Since $p \wedge \neg p \equiv 0$

1.4 *Converse and Contrapositive***Definition 2.**

1. the **converse** of $p \Rightarrow q$ is $q \Rightarrow p$

NB. $p \Rightarrow q \not\equiv q \Rightarrow p$

2. the **contrapositive** of $p \Rightarrow q$ is $\neg q \Rightarrow \neg p$

$p \Rightarrow q \equiv \neg p \vee q \equiv \neg q \Rightarrow \neg p$

1.5 *Proofs by contradiction*

You assume something is true, and get something nonsense.

$$\neg p \Rightarrow 0 \equiv \neg(\neg p) \vee 0 \equiv p$$

Example 4

prop. there is no least positive rational number.

$$u = \mathbb{Q} : \neg(\exists x : x > 0 \wedge (\forall y : y > 0 \Rightarrow x \leq y))$$

proof. Suppose, for a contradiction that the proposition is false, that is, there exist $x \in \mathbb{Q}$ such that $x > 0$ and for all $y \in \mathbb{Q}$ with $y > 0$, $x \leq y$.

Let $y = \frac{x}{2}$, we have $\frac{x}{2} > 0$ since $x > 0$. Then $x \leq y$, so $x \leq \frac{x}{2}$.

Divide through by x (because $x > 0$) to get $1 \leq \frac{1}{2}$.

the contradiction completes the proof. Q.E.D

1.6 *Case Analysis***Example 5**

prop. There exists irrational numbers a, b such that a is rational.

If

2 Functions

Definition 3 (Surjective and injective). A function is surjective if

$$\forall b \in B, \exists a \in A, f(a) = b$$

A function is injective if

$$\forall a_1, a_2 \in A, a_1 \neq a_2 \rightarrow f(a_1) \neq f(a_2)$$

or $a_1 = a_2 \rightarrow f(a_1) = f(a_2)$.

3 Modular Arithmetic

3.1 Division

Proposition 1. Let a and b be nonzero integers. The set

$$X = \{s'a + t'b : s', t' \in \mathbb{Z}\}$$

is exactly the set of multiples of $d = \gcd(a, b)$.

Proof. By Bezout's identity, there exist integers s, t such that $d = sa + tb$. First let $n \in \mathbb{Z}$ be a multiple of d , so $n = kd$ for some $k \in \mathbb{Z}$. Then we have

$$n = kd = d(sa + tb) = (ds)a + (dt)b,$$

which means that $n \in X$, since $ds, dt \in \mathbb{Z}$.

Conversely, suppose that $n \in X$. Then $n = s'a + t'b$ for some $s', t' \in \mathbb{Z}$. Then since d divides a and b , we can write $a = ld$ and $b = md$ for some integers $l, m \in \mathbb{Z}$. So we have

$$n = s'a + t'b = s'ld + t'md = (s'l + t'm)d,$$

which shows that $d|n$, since $s'l + t'm \in \mathbb{Z}$. \square

4 Graph Theory

4.1 Definitions

A graph is a pair $G = (V, E)$, where V is a nonempty set and

$$E \subseteq \{u, v : u, v \in V, u \neq v\}.$$

Degrees and k-regularity The neighbors of a vertex v are all $u \in V$ such that $uv \in E$. The degree of a vertex v is the number of neighbors

The graph G is said to be *finite* if both V are finite sets. A simple graph is a graph that does not have more than one edge between any two vertices and no edge starts and ends at the same vertex.

of v , denoted by $\deg(v)$. A graph is said to be k -regular for some $k \in \mathbb{N}$ if every $v \in V$ has degree k .

The following theorem relates vertex degree to the number of edges.

Theorem 1. *Let $G = (V, E)$ be a finite graph. Then*

$$\sum_{v \in V} \deg(v) = 2|E|.$$

Corollary 1 (Handshaking Lemma). *In every finite simple graph, the number of vertices having odd degree is even.*

From Theorem above, we can derive a corollary that counts the number of edges in k -regular graphs.

Corollary 2. *Let $G = (V, E)$ be k -regular. Then*

$$|E| = \frac{k|V|}{2}.$$

Walks, paths, and cycles. A *walk* in a graph $G = (V, E)$ is a sequence of vertices $\sigma = (v_0, v_1, \dots, v_k)$ such that $v_i v_{i+1} \in E$ for all i with $0 \leq i < k$. The *endpoints* of the walk are v_0 and v_k , and the *length* of the walk σ is $|\sigma|$. The walk is said to be *closed* if $v_0 = v_k$ and *open* otherwise.

A walk is a *path* if no vertices are repeated.

Theorem 2. *Let $G = (V, E)$ be a graph. If u and v are vertices such that there exists a walk from u to v , then there exists a path from u to v .*

Proof. Let $\sigma = (v_0, v_1, \dots, v_n)$ be a walk from u to v of shortest length. We claim that σ is a path. Indeed, suppose for a contradiction that it is not a path; then there is some repeated vertex, so there exist $i, j \in 0, 1, \dots, n$ such that $i < j$ and $v_i = v_j$. Hence there is no need to visit any of the vertices between v_i and v_j in the walk, since $v_i = v_j$ is connected to v_{j+1} . Concretely, consider the walk

$$\sigma' = (v_0, v_1, \dots, v_i, v_{j+1}, \dots, v_n).$$

Note that $|\sigma'| = |\sigma| - (j - i)$, and $j - i > 0$, so σ' is a shorter walk from u to v . But this contradicts our choice of σ as a walk of shortest length. We conclude that σ is a path. \square

A *cycle* is a walk of length at least 3 and no vertices repeated except for $v_0 = v_k$.

Proposition 2. *Let $G = (V, E)$. If G contains a closed walk of odd length, then G contains a cycle of odd length.*

There we perform a minimality argument. Note that this theorem is about existence of path, a walk is not necessary a path.

Proof. Let $\sigma = (v_0, v_1, \dots, v_n)$ be an odd-length closed walk in G , and choose this walk to have minimal odd length (i.e, any shorter closed walk has even length). We shall prove that σ is a cycle.

For a contradiction, suppose that σ is not a cycle, so that there exist indices i, j with $0 \leq i < j \leq n$ such that $v_i = v_j$. Consider the two closed walks

$$\sigma_1 = (v_0, v_1, \dots, v_i, v_{j+1}, \dots, v_n)$$

and

$$\sigma_2 = (v_i, v_{j+1}, \dots, v_j).$$

Both are shorter than σ , so by the minimality of σ , they must have even length. But this implies that $|\sigma| = |\sigma_1| + |\sigma_2|$ is even. This leads to a contradiction.

We conclude that σ is a cycle. \square

Connectedness. We say a graph $G = (V, E)$ is *connected* if for all $u, v \in V$, there exists a walk from u to v .

Proposition 3. For all $n \geq 1$, the graph K_n and Q_n are connected.

Proof. Any two vertices u and v in K_n are connected by an edge, so we have a path (u, v) of length 1 between u and v . This shows that K_n is connected.

Now let u and v be any two vertices in Q_n . Suppose there are m bits that differ between u and v . Then we can flip them one by one to change u to v . This gives us a walk of length m between u and v , since there is an edge in Q_n between any two strings in Q_n that differ at exactly one bit. \square

Here is an example:

Example 6 (Disconnected Graph and Modular Arithmetic)

Let $G = (V, E)$, where for $i, j \in \mathbb{Z}$ with $i < j$, we have $ij \in E$ if and only if $j - i \in 9, 15$. Then G is disconnected.

Proof. Starting at $n \in \mathbb{Z}$, we can reach any vertex m that is of the form

$$m = n + 15s + 9t$$

for some $s, t \in \mathbb{Z}$. By proposition in division section, the integers representable as $n + 15s + 9t$ are exactly the multiples of $\gcd(15, 9) = 3$. So in fact, from n one can reach any integer m of the form $n + 3k$ for some $k \in \mathbb{Z}$. That is, one can reach any integer m with $m \equiv n \pmod{3}$. Hence the three connected components of G are $[0]_3, [1]_3, [2]_3$, the equivalence classes of integers modulo 3. \square

The idea is similar to the proof of the previous theorem.

This is a disconnected graph:



4.2 Triangles and bipartite graphs

A *subgraph* of a graph $G = (V, E)$ is a graph $G' = (V', E')$ such that $V' \subseteq V$ and $E' \subseteq E$ where for all $e = uv \in E'$, we have $u, v \in V'$.

The extremal question An extremal question asks for the external (maximum or minimum) number of objects we can have, subject to some restrictions.

We now want to derive the maximum number of edges in a triangle-free graph on n vertices.

Theorem 3 (Cauchy-Schwarz Inequality).

For all $u_1, \dots, u_n, v_1, \dots, v_n \in \mathbb{R}$, we have

$$\left(\sum_{i=1}^n u_i v_i \right)^2 \leq \left(\sum_{i=1}^n u_i^2 \right) \left(\sum_{i=1}^n v_i^2 \right).$$

Theorem 4 (Mantel's theorem). Let $G = (V, E)$ be a graph not containing a triangle as a subgraph. Then

$$|E| \leq \lfloor \frac{|V|^2}{4} \rfloor.$$

Proof. Consider the sum

$$\sum_{uv \in E} (\deg(u) + \deg(v)).$$

The term $\deg(u)$ appears in the sum exactly once for every edge incident on u ; that is, it appears $\deg(u)$ times. This is true for all $u \in V$, so we conclude that

$$\sum_{uv \in E} (\deg(u) + \deg(v)) = \sum_{u \in V} \deg(u)^2.$$

On the other hand, since G contains no triangle, for every pair of vertices u and v , the set of neighbors of u is disjoint from the set of neighbors of v . So $\deg(u) + \deg(v) \leq |V|$, and we have

$$\sum_{u \in V} \deg(u)^2 = \sum_{uv \in E} (\deg(u) + \deg(v)) \leq |E| |V|.$$

By the Cauchy-Schwarz inequality, we have

$$(2|E|)^2 \leq |V| \left(\sum_{u \in V} \deg(u)^2 \right)$$

Hence

$$4|E|^2 \leq |V| \left(\sum_{u \in V} \deg(u)^2 \right) \leq |V|^2 |E|.$$

This implies that $|E| \leq \frac{|V|^2}{4}$. We can take the floor function on the R.H.S, since $|E|$ must be an integer. \square

Derivation using Cauchy-Schwarz inequality:

$$(2|E|)^2 = \left(\sum_{u \in V} (\deg(u)) \right)^2 \quad (1)$$

$$= \left(\sum_{u \in V} \deg(u) \cdot 1 \right)^2 \quad (2)$$

$$\leq \left(\sum_{u \in V} \deg(u)^2 \right) \left(\sum_{u \in V} 1^2 \right) \quad (3)$$

$$= |V| \left(\sum_{u \in V} \deg(u)^2 \right). \quad (4)$$

So if a graph has $|E| > \lfloor |V|^2/4 \rfloor$, there must be a triangle sub-graph in G .

How about the case a graph has exactly $|V|^2/4$ edges? The Mantel's theorem does not assert that it must have triangles. This brings us to the definition of a bipartite graph.

Bipartite graphs A graph $G = (V, E)$ is *bipartite* if there exists a partition of $V = A \cup B$ of the vertex set ($A \cap B = \emptyset$) called the *bipartition* such that each edge has one endpoint in A and the other in B . For example, hypercubes Q_n are bipartite.

Proposition 4. For all $n \geq 1$, the graph Q_n is bipartite.

Proof. Let $Q_n = (V, E)$. Every element $s \in V$ corresponds to a binary string of length n , $S = (s_1, s_2, \dots, s_n)$ where each s_i is either 0 or 1. Define

$$V_0 = \{s \in V : s_1 + \dots + s_n \equiv 0 \pmod{2}\}$$

and

$$V_1 = \{s \in V : s_1 + \dots + s_n \equiv 1 \pmod{2}\}.$$

It's clear that $V_0 \cup V_1 = V$ and $V_0 \cap V_1 = \emptyset$, so this is a bipartition of the vertex set. For every $e = s_1 s_2 \in E$, the strings s_1 and s_2 differ in exactly one bit, so if $s_1 \in V_0$, then $s_2 \in V_1$ and vice versa. Hence Q_n is bipartite. \square

The *complete bipartite graph* $K_{m,n}$ is a bipartite graph with bipartition $V = V_m \cup V_n$, where $|V_m| = m$ and $|V_n| = n$, and E is the set $uv : u \in V_m, v \in V_n$ of all possible edges between the two sets.

When n is even, the graph $K_{n/2, n/2}$ has n vertices and exactly $n^2/4$ edges. When n is odd, the graph $K_{(n-1)/2, (n+1)/2}$ has n vertices and

$$\frac{n+1}{2} \cdot \frac{n-1}{2} = \frac{n^2-1}{4} = \frac{n^2}{4} - \frac{1}{4} = \lfloor \frac{n^2}{4} \rfloor$$

edges. (The last equality here follows from the fact that any odd n is congruent to either 1 or 3 modulo 4, which means its square is congruent to 1 modulo 4.)

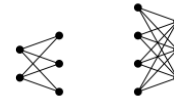
Lemma 1. A graph $G = (V, E)$ is bipartite if and only if its connected components are bipartite.

Proof. Let $G_1 = (V_1, E_1), \dots, G_k = (V_k, E_k)$ be the connected components of G .

If each connected component is bipartite, then we can bipartition each V_i into $A_i \cup B_i$. Then $A = \bigcup_{i=1}^n A_i$ and $B = \bigcup_{i=1}^n B_i$ is a bipartition of V . And every edge in G has one endpoint in A and the other in B , so G is bipartite.

Now suppose that there is some connected component G_i that is not bipartite. Now let $V = A \cup B$ be any partition of V into disjoint

This is the converse of the proposition we just proved. Pay attention here, because that means the theorem does not assert any graph has edges less or equal to $\lfloor |V|^2/4 \rfloor$ is triangle-free.



The complete bipartite graphs $K_{2,3}$ and $K_{4,3}$.

$K_{n/2, n/2}$ has the largest possible number of edges such that the Mantel's theorem does not apply. But adding a single edge to $K_{n/2, n/2}$ results in a graph that must contain a triangle by Mantel's theorem.

nonempty sets. Then $A_i = A \cap V$ and $B_i = B \cap V$ form a partition of V_i into disjoint nonempty sets. But since G_i is not bipartite, there must be an edge $uv \in E_i$ with both endpoints in A_i or both in B_i . This means that $A \cup B$ is not a bipartition of V , and since A and B were arbitrary, we conclude that G is not bipartite either. \square

As a matter of fact, we shall prove something much stronger than just the fact that $K_{n/2, n/2}$ does not contain any triangles.

Theorem 5. *A graph is bipartite if and only if it does not contain any cycles of odd length.*

Proof. First, assume that $G = (V, E)$ is bipartite with bipartition $V = A \cup B$. Let σ be a cycle in G . Each edge changes sides between A and B , so in order for starting and ending vertices of this cycle to be the same, $|\sigma|$ must be even.

Now suppose that G does not contain any cycles of odd length. To show that G is bipartite, it suffices to show that each of its connected components is bipartite by the lemma above. So without loss of generality we may assume that G is connected. That means $\text{dist}(u, v) < \infty$ for all $u, v \in V$.

Select on vertex $h \in V$ and set

$$V_0 = \{v \in V : \text{dist}(h, v) \equiv 0 \pmod{2}\}$$

and

$$V_1 = \{v \in V : \text{dist}(h, v) \equiv 1 \pmod{2}\}.$$

Let $e = uv \in E$. Consider a closed walk σ that follows a shortest path from u to h then a shortest path from h to v , and finally the edge uv . The length of σ is

$$|\sigma| = \text{dist}(u, h) + \text{dist}(h, v) + 1$$

But since G contains no cycles of odd length, $|\sigma|$ must be even, by (the Contrapositive of) Proposition 2, any closed walk in G must have even length. This means that $\text{dist}(u, h), \text{dist}(h, v)$ are not the same modulo 2. That is, either $u \in V_0$ and $v \in V_1$ or vice versa. \square

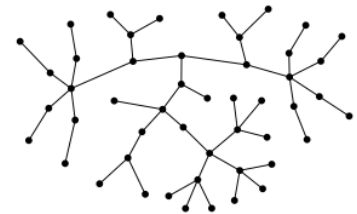
Proof of this theorem requires the metric of distance imposed on a graph G . For all $u, v \in V$, the distance $\text{dist}(u, v)$ is the length of the shortest path from u to v . If no such path exists, then $\text{dist}(u, v) = \infty$.

4.3 Trees

A graph G is called a *forest* if it contains no cycles. A *tree* is a connected forest. A vertex of degree 1 in a forest is called a *leaf*.

Proposition 5. *Let F be a forest with at least one edge (and hence at least two vertices). There are at least two leaves in F .*

Proof. Let $\sigma = (v_0, v_1, \dots, v_k)$ be a path of maximum length in F . Since there is at least one edge in F , we have $|\sigma| \geq 1$. The claim is that



A tree with 25 leaves.

$\deg(v_0) = \deg(v_k) = 1$. Suppose for a contradiction that $\deg(v_0) \geq 2$. So there exists $u \in V$ such that uv_0 is an edge. If $u = v_i$ for some i , then there is a cycle v_i, v_0, \dots, v_i contradicting the fact that F is a forest. Hence for all $1 \leq i \leq k$, we have $u \neq v_i$. But this a path of length $k + 1$ exists. This contradicts the maximality of σ . \square

The intuition is that, a tree should have at least two ends, if not, either there is a cycle or there is an extending part (a new end) that can be added to the tree.

Proposition 6 (A characterization of trees). *A graph $G = (V, E)$ is a tree if and only if for all $u, v \in V$, there exist a unique path from u to v .*

Proof. First, suppose that G is a tree. Let $u, v \in V$, since G is a tree, it is connected, so there must exist some path σ from u to v . Now we prove that this path is unique. For a contradiction, let σ' be another path from u to v . Both σ and σ' starts at u , but eventually they must diverge. Let v_1 be the last vertex where σ and σ' have in common (this should be u). Eventually they must meet again, since they both end at v . Let v_2 be the first vertex where they meet again. Now let σ'' be the path obtained by following σ from v_1 to v_2 , then following σ' from v_2 to v . This is a cycle, by choice of v_1 and v_2 . But this give us a contradiction, since G is a tree and contains no cycles.

Now suppose that for all $u, v \in V$, there exists a unique path from u to v . The fact that there exists a path at all between every two vertices immediately implies that G is connected. It remains to prove that there is no cycle in G . But this is clear, since if σ is a cycle in G , then picking any two distinct vertices u, v on σ , we have two distinct paths from u to v (clockwise and counterclockwise along σ). \square

We can prove another characterization of trees, but we need a lemma first.

Lemma 2 (Detour Lemma). *Let $G = (V, E)$ be a connected graph and let σ be a cycle in G . If G' is the graph obtained by deleting one edge from σ , then G' is still connected.*

Lemma 3. *A graph $G = (V, E)$ is a tree if and only if it is connected and $|E| = |V| - 1$.*

Proof. We prove the forward implication by induction on $n = |V|$. Concretely, the statement we shall prove is, for all $n \geq 1$, for all $G = (V, E)$ with $|V| = n$, if G is a tree, then $|E| = n - 1$ and G is connected.

The base case in $n = 1$, where the only possible graph G is simply a single vertex with no edges. This is a tree, since it is connected and contains no cycles. We have $|E| = 0 = 1 - 1$.

For the inductive step, let $n \geq 1$ and suppose that the statement holds n . Let $G = (V, E)$ be a tree with $|V| = n + 1$. Since $n \geq 1$, $|V| \geq 2$, there are at least two vertices, and the connectedness of G means that there is at least one edge in G . Hence G is a forest with

at least one edge, and by Proposition 5, there are at least two leaves in G . Pick one of these leaves and call it u . Form a new graph $G' = (V', E')$ by removing the leaf u and the edge incident on it. We have $|V'| = n$ and $|E'| = |E| - 1$. The graph G' is a tree, because for each pair of vertices $v, w \in V'$, there is a unique path from v to w in G , and this path does not pass through u (since u is a leaf). Hence

$$|E| = |E'| + 1 = |V'| - 1 + 1 = |V| - 1,$$

where in the second equality we have to use the inductive hypothesis.

Now assume that G is connected and $|E| = |V| - 1$. We want to show that G has no cycles. For a contradiction, suppose that σ is a cycle in G . Remove an edge from σ , by detour lemma, the resulting graph is still connected. If this graph still has a cycle, remove an edge from it again, and repeat this process until no cycles remain. This graph G' still has $|V|$ vertices, but it has a new edge set E' with $|E'| < |E|$. But G' is now a tree, since it contains no cycles but is still connected. So by the previous paragraph, we have $|E'| = |V| - 1 = |E|$. The contradiction completes the proof. \square

Now we are able to answer the following extremal question: How many edges can a graph on n vertices have if it does not contain a cycle (i.e., is a forest)?

Corollary 3. *Let $F = (V, E)$ be a forest, then $|E| \leq |V| - 1$.*

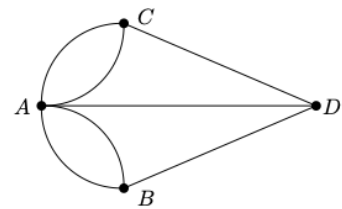
Proof. Let $k \geq 1$ denote the number of connected components of F , and number the connected components as C_1, \dots, C_k . Join C_1 to C_2 by an edge, then join C_2 to C_3 by an edge, and so on. This creates a connected graph $G = (V, E')$ with $|E'| = |E| + k - 1$, and G must be a tree, since adding these edges does not introduce a cycle. By the previous theorem, $|E'| = |V| - 1$, so $|E| \leq |E'| = |V| - 1$. \square

4.4 Eulerian trails and circuits

Recall that a walk that does not repeat any vertex is called a path. Now we introduce the concept of *trail*: this is a walk that does not repeat any *edge*. A walk is called an *Eulerian trail* if it uses every edge of G exactly once, and an *Eulerian circuit* if it is an Eulerian trail and it is closed.

Theorem 6. *Let $G = (V, E)$ be a connected graph with a finite number of vertices and edges, where multiple edges between the same pair of vertices are allowed. Then G has an Eulerian circuit if and only if every vertex in G has even degree.*

Proof. Suppose G has an Eulerian circuit $\sigma = (v_0, v_1, \dots, v_k)$ where $v_0 = v_n$. For each $v_i \neq v_0$, each visit uses two edges, one edge to enter



The graph representing the seven bridges of Königsberg

v_i and one edge to leave v_i . So $\deg(v_i)$ is even for all $0 \leq i < n$. For $v_0 = n$, there is one edge that leaves it at the beginning, two edges for each visit during the circuit, and then one returning to it at the end. So $\deg(v_0)$ is even as well.

We prove the reverse implication by induction on $m = |E|$. That is, we prove that for all integers $m \geq 0$, if $G = (V, E)$ with $|E| = m$ and all vertices in G have even degree, then G has an Eulerian circuit.

For the base case $m = 0$, the graph G is a single vertex with no edges, and it has an Eulerian circuit.

Now let $m \geq 1$ and suppose the statement holds for all integers less than or equal to m . Let $G = (V, E)$ with $|E| = m + 1$. Since there is at least one edge in G , G cannot be a tree. This is because every tree with at least one edge contains at least two leaves, and leaves have (odd) degree 1. Let σ be a cycle in G and form a new graph $G' = (V, E')$ by deleting every edge in σ . This graph may no longer be connected, but since we remove a cycle, every vertex in G' either has the same degree as in G , or degree two less than in G . So all vertices in G' have even degree as well. Let k be the number of connected components of G' ; number them as H_1, \dots, H_k . Each has at most m edges, so by the inductive hypothesis, each has an Eulerian circuit.

We build an Eulerian circuit in G using the process to the right. Eventually, this gives us an Eulerian circuit in G . \square

Theorem 7. *Let $G = (V, E)$ be a connected graph (with multiple edges allowed). Then G contains an Eulerian trail that isn't an Eulerian circuit if and only if exactly two vertices of G have odd degree.*

Proof. Suppose that G contains an Eulerian trail $\sigma = (v_0, v_1, \dots, v_k)$ in which $v_0 \neq v_k$. By a similar reasoning to the first paragraph of the previous proof, every vertex in the walk has even degree, except both v_0 and v_k must have odd degree, since starting at v_0 we must leave for the first time via an edge, and at the end we enter v_k one last time without exiting.

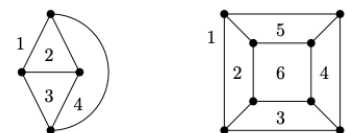
Now assume that G contains exactly two odd-degree vertices. By the previous theorem, there does not exist an Eulerian circuit in G . Let u and v be the two odd-degree vertices. Add a new edge uv , even if there was already an edge between u and v before. We get G' in which every vertex has even degree. So G' must have an Eulerian circuit σ' . It contains the edge uv , but we can remove it to get an Eulerian trail in G . \square

Process of building an Eulerian circuit in G :

- Start at any vertex σ
- Follow σ
- Each time we reach a vertex of a connected component H_i of G' for the first time, follow its Eulerian circuit. (We know that each connected component has an Eulerian circuit by the inductive hypothesis.)
- Once we return to the beginning of the Eulerian circuit, continue along σ .
- Repeat this process until we return to the vertex of σ we started at.

4.5 Planar graphs

A graph G is called *planar* if it can be drawn in the plane without any edges crossing. If we can, such a drawing is called a *planar embedding* of G .



Planar embeddings of K_4 and Q_3

Planarity is a property that passes down to subgraphs.

Proposition 7. *If G is planar then any subgraph H of G is also planar.*

Proof. If G is drawn in the plane without edges crossing, and we erase some edges and vertices to create a drawing of H , we cannot introduce any edge crossings along the way. \square

A planar embedding of a graph determines regions, or *faces* bounded by edges, including on region outside the graph. We define the *Euler characteristic* of a planar embedding of G to be the quantity

$$\chi = |V| - |E| + f,$$

where f is the number of faces in the embedding.

From the figure above, we notice that both K_4 and Q_3 have $\chi = 2$, this is not just a coincidence.

Theorem 8 (Euler's formula). *Let $G = (V, E)$ be a connected planar graph, let f be the number of faces in a planar embedding of G . Then $|V| - |E| + f = 2$.*

To prove this, we need the following theorem.

Theorem 9 (Jordan curve theorem). *Every closed curve in the plane \mathbb{R}^2 that does not intersect itself divides the plane into two regions.*

Proof of Theorem 8. We proceed by induction on $m = |E|$. In the case that $m = 0$, G can only be a single vertex, in which we have $\chi = 1 - 0 + 1 = 2$.

Now let $m \geq 0$ and assume that the theorem holds for all natural numbers at most m . Let $G = (V, E)$ be a connected planar graph with $|E| = m + 1$. Draw G in the plane and let f be the number of faces in the embedding. There are two cases, according to whether G is a tree or not. If G is a tree, then $f = 1$ since G has no cycles. But we also have $|E| = |V| - 1$, so

$$\chi = |V| - |E| + f = |V| - |V| + 1 + 1 = 2.$$

If G is not a tree, then G has a cycle, call it σ . Form G' by removing exactly one edge e from σ , so

$$|E'| = |E| - 1 = m + 1 - 1 = m.$$

By the detour lemma, G' is still connected, and it is planar, since removing an edge doesn't introduce any crossings. By the inductive hypothesis, the Euler characteristic χ' of G' satisfies

$$\chi' = |V| - |E'| + f' = 2,$$

where f' is the number of faces in a planar embedding of G' . Returning to the drawing of G , the cycle σ is closed curve without intersections in the plane, so it divides the plane into two regions (each of which may have multiple faces within them). In particular, the removed edge e forms a border between two distinct faces, and in G' , these two faces combine into one face. Hence $f = f' + 1$. Putting everything together, we calculate the Euler characteristic of G satisfies

$$\chi = |V| - |E| + f = |V| - (|E'| + 1) + (f' + 1) + 1 = 2. \quad \square$$

Euler's formula tells us that the number of faces in any drawing of G depends entirely on the number of vertices and edges in G . We can use this information to give conditions for a graph to be non-planar.

Theorem 10. *Let $G = (V, E)$ be a connected planar graph with $|V| \geq 5$. The*

$$|E| \leq 3|V| - 6.$$

Under the further assumption that G contains no triangles, we have a better bound

$$|E| \leq 2|V| - 4.$$

Proof. The proof is by double counting. Let R be the set of all regions into which G divides the plane, so that $|R| = f$, the number of faces. Consider the set

$$S = \{(e, r) \in E \times R : \text{the edge } e \text{ touches the region } r\}.$$

We will count $|S|$ in two ways. First off, each edge $e \in E$ touches at most two regions, so $|S| \leq 2|E|$. On the other hand, every region is bounded by a cycle, and a cycle has at least three edges, so each region $r \in R$ touches at least 3 edges. In other words, $|S| \geq 3f$. Chaining these into two inequalities, we have

$$3f \leq |S| \leq 2|E|.$$

But by Euler's formula, we have $|V| - |E| + f = 2$, so $f = 2 + |E| - |V|$, and substituting this above, we get

$$3(2 + |E| - |V|) \leq 2|E|.$$

Distributing and rearranging terms yields the desired inequality $|E| \leq 3|V| - 6$.

If we have the further assumption that there are no triangles in G , then every region must touch at least 4 edges, allowing us to conclude the stronger inequality $4f \leq 2|E|$. Then we proceed as above to get

$$4(2 + |E| - |V|) \leq 2|E|,$$

which can be manipulated to get $|E| \leq 2|V| - 4$. \square

The condition that $|V| \geq 5$ is a complete non-issue, since a graph with fewer than 5 vertices is a subgraph of K_4 , and we can apply Proposition 7, since we already know that K_4 is planar. As a corollary of the (contrapositive of the) previous theorem, we can give two important examples of nonplanar graphs.

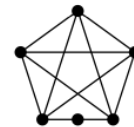
Corollary 4. *The complete graph K_5 and the complete bipartite graph $K_{3,3}$ are nonplanar.*

Proof. If $(V, E) = K_5$ we have $|V| = 5$ so $3|V| - 6 = 9$, but $|E| = 10$. So by the contrapositive of the previous theorem, K_5 is nonplanar.

The graph $K_{3,3}$ does not contain any triangles, so if it were planar, then the stronger bound $|E| \leq 2|V| - 4$ must hold. But in $(V, E) = K_{3,3}$, we have $|V| = 6$ and $2|V| - 4 = 8$, but $|E| = 9$. So $K_{3,3}$ is nonplanar. \square

This corollary, combined with the contrapositive of Proposition 7, shows that K_n is nonplanar for all $n \geq 5$, and that $K_{m,n}$ is nonplanar for all $m, n \geq 3$. More generally, any graph that contains K_5 or $K_{3,3}$ as a subgraph is nonplanar.

However, subgraphs are not exactly the right notion to be considering when talking about planarity. To see why, consider the graph to the right.



The graph K_5 with an extra vertex subdividing an edge.

It is nonplanar, since if it had a planar embedding, then by contracting on of the edges incident on the vertex, and bending other edges accordingly, we would obtain a planar embedding of K_5 . On the other hand, this graph does not contain K_5 as a subgraph. We thus introduce the following notion.

We say that a graph H is a *graph minor* of a graph G if H can be obtained from G repeatedly by either one of the following operations to the right.

Any number of the first two operations simply creates a subgraph. It is the contraction operation that produces interesting examples of graph minors. It is easy to see pictorially that contracting edges preserves planarity, so any graph minor of a planar graph is planar. We can use this to show, for instance, that the Petersen graph, depicted to the right, is nonplanar. Since by contracting each of the five edges connecting the inner star pentagram to the other pentagon, we obtain a K_5 . (Note that Theorem 10 does not apply to the Petersen graph.)

Not only all graphs that have K_5 and $K_{3,3}$ as minors are nonplanar, it turns out that there are the only nonplanar graphs.

Theorem 11 (Wagner's theorem). *A graph G is nonplanar if and only if either K_5 or $K_{3,3}$ is a minor of G .*

It deals with the notion of a *subdivision* of a graph G , subdividing

- deleting an edge;
- deleting a vertex; or
- contracting an edge uv by removing it and merging u and v into a single vertex (and also combining any resulting multiple edges into a single edge).

an edge into two edges each time. (We have shown an example of a subdivision of K_5)

Theorem 12 (Kuratowski's theorem). *A graph G is nonplanar if and only if it contains a subdivision of either K_5 or $K_{3,3}$ as a subgraph.*

It is easy to see that if H is a subdivision of G , then G is a minor of H , since we can reobtain G by contracting all the edges created by the subdivision operations. So the two theorems above are certainly very closely related, though it is not immediately obvious if one should imply the other. The truth is that they are equivalent, because it can be shown that any graph with either K_5 or $K_{3,3}$ as a minor also has a subgraph that is a subdivision of one of them.

References