# Inductive Definitions

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A set of rules that are used to derive facts about judgments.

#### Inductive Definitions: Judgments

A judgment is an assertion about a syntactic object.

It states that one or more syntactic objects have a property or is related to another.

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#### Examples:

n nat

n is a natural number

 $n = n_1 + n_2$ 

n is the sum of  $n_1$  and  $n_2$ 

t: type

t is a type

e:t

the expression e has type t

### Inductive Definitions: Judgments

We denote an arbitrary judgment as J.



$$\frac{J_1}{J}$$
 ...  $J_n$  Name

$$J_1 \cdots J_n$$
Name

Premises 
$$\longrightarrow I_1 \cdots I_n$$
Name

Premises 
$$\longrightarrow$$
  $J_1$   $\cdots$   $J_n$ 
Conclusion  $\longrightarrow$   $J$ 

true bool false bool 
$$\frac{b_2 \operatorname{bool}}{b_3 \operatorname{bool}}$$
 
$$\frac{b_1 \operatorname{bool}}{\operatorname{if}(b_1, b_2, b_3) \operatorname{bool}}$$

$$\frac{n \operatorname{nat}}{0 \operatorname{nat}}$$
 Succ  $\frac{n \operatorname{nat}}{\operatorname{succ}(n) \operatorname{nat}}$ 

$$\frac{n \text{ nat}}{\text{empty list}} \frac{l \text{ list}}{\cos(n, l) \text{ list}}$$

#### Inductive Definitions: Derivations

To prove that an inductively defined judgment, J, holds it is enough to exhibit a <u>derivation</u> of it.

#### Inductive Definitions: Derivations

Derivations are goal directed proofs written by stacking inference rules to form a derivation tree.

$$\frac{\Delta_1 \cdots \Delta_n}{J}$$
 Name

1 nat

cons(2,cons(3,empty)) list

Cons

0,nat

1 nat

cons(2,cons(3,empty)) list

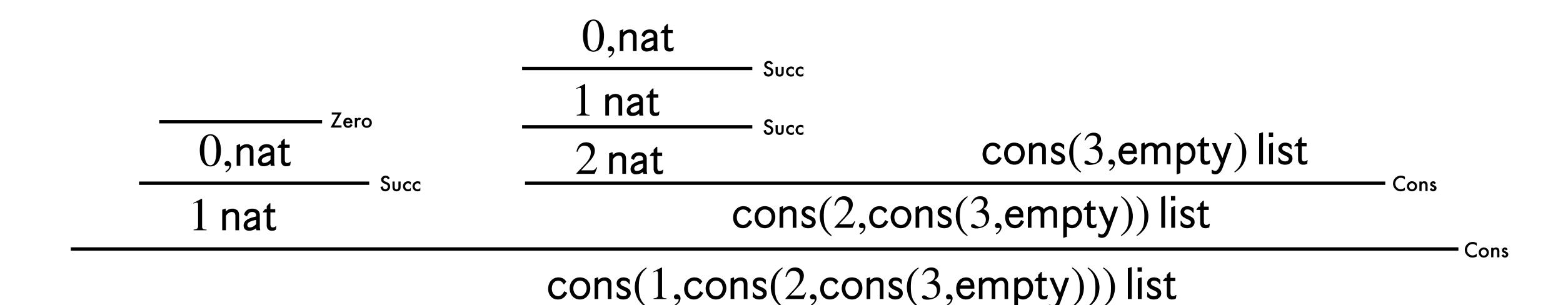
Cons

cons(2,cons(3,empty)) list

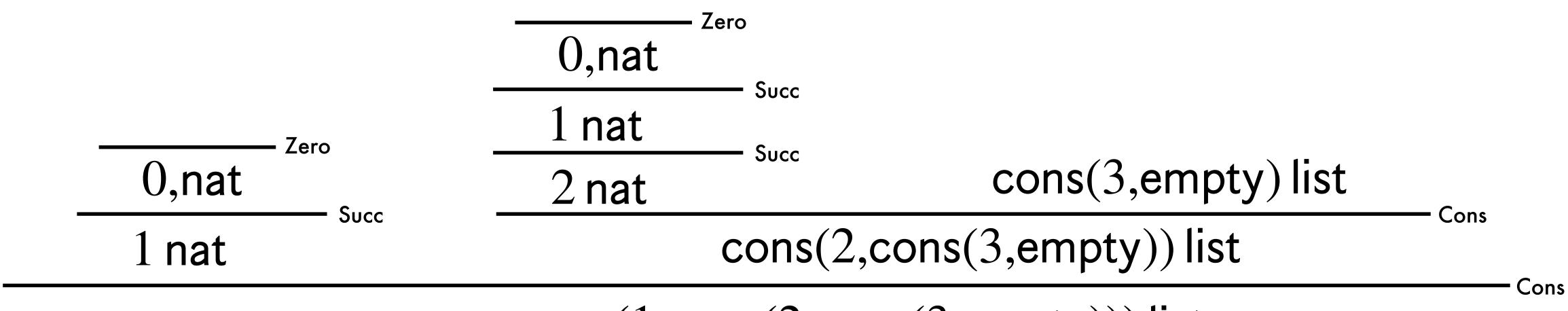
Cons

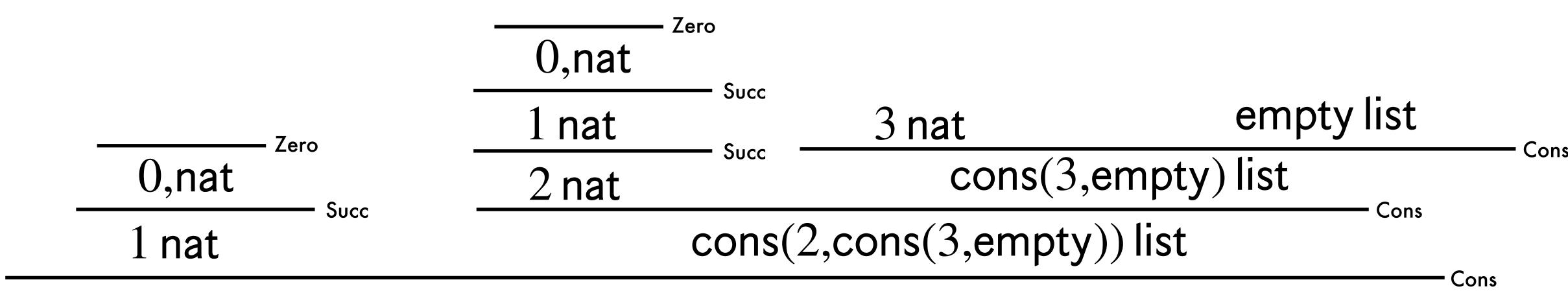
O,nat	2 nat	cons(3,empty) list	
1 nat	cons(2,cons(3,empty)) list		—— Cons
	cons(1,cons(2,c))	cons(3,empty))) list	

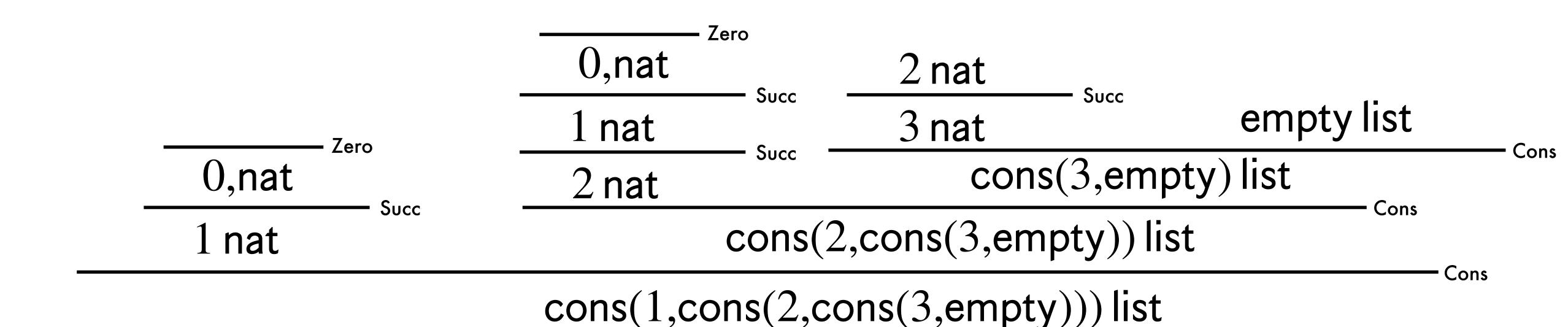
$$\begin{array}{c|c} \hline 0, \mathsf{nat} & \hline 1 & \mathsf{nat} \\ \hline 1 & \mathsf{nat} \\ \hline \\ \hline 1 & \mathsf{nat} \\ \hline \end{array} \begin{array}{c} 1 & \mathsf{nat} \\ \hline \\ & \mathsf{cons}(2, \mathsf{cons}(3, \mathsf{empty})) \ \mathsf{list} \\ \hline \\ & \mathsf{cons}(1, \mathsf{cons}(2, \mathsf{cons}(3, \mathsf{empty}))) \ \mathsf{list} \\ \hline \end{array}$$

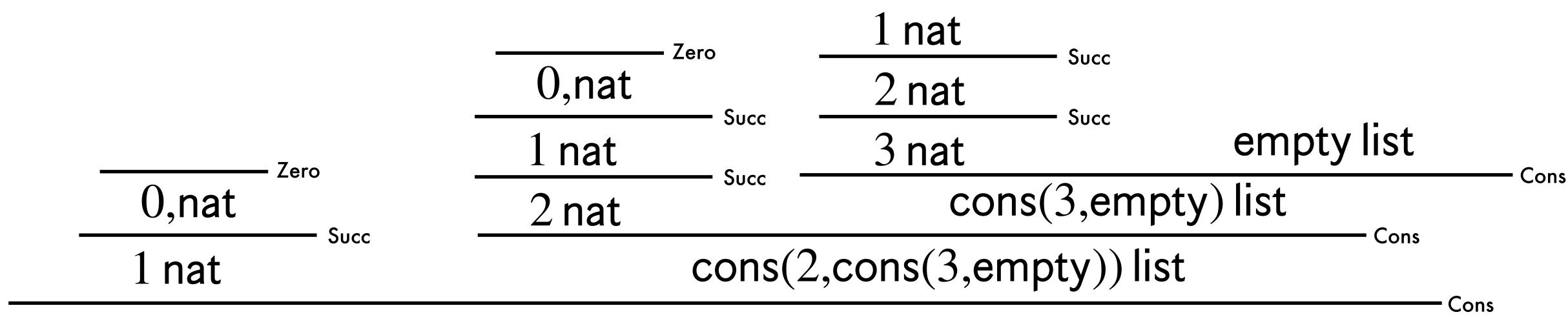


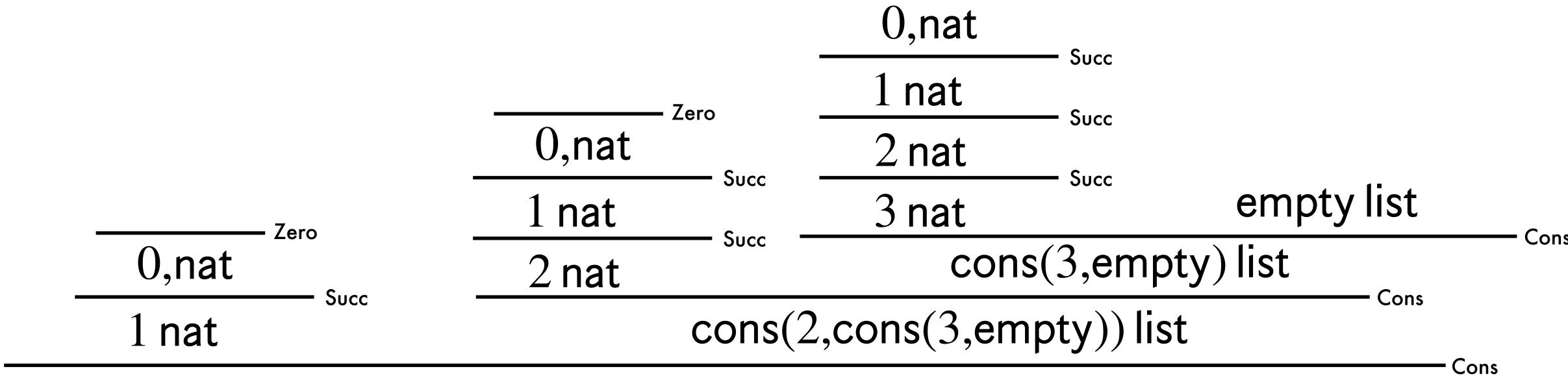
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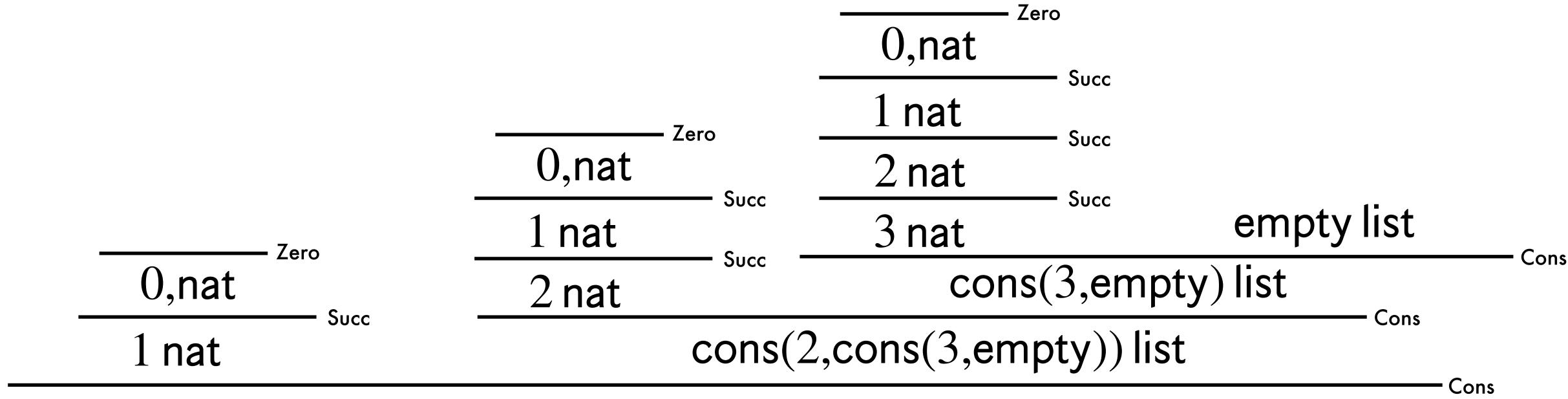


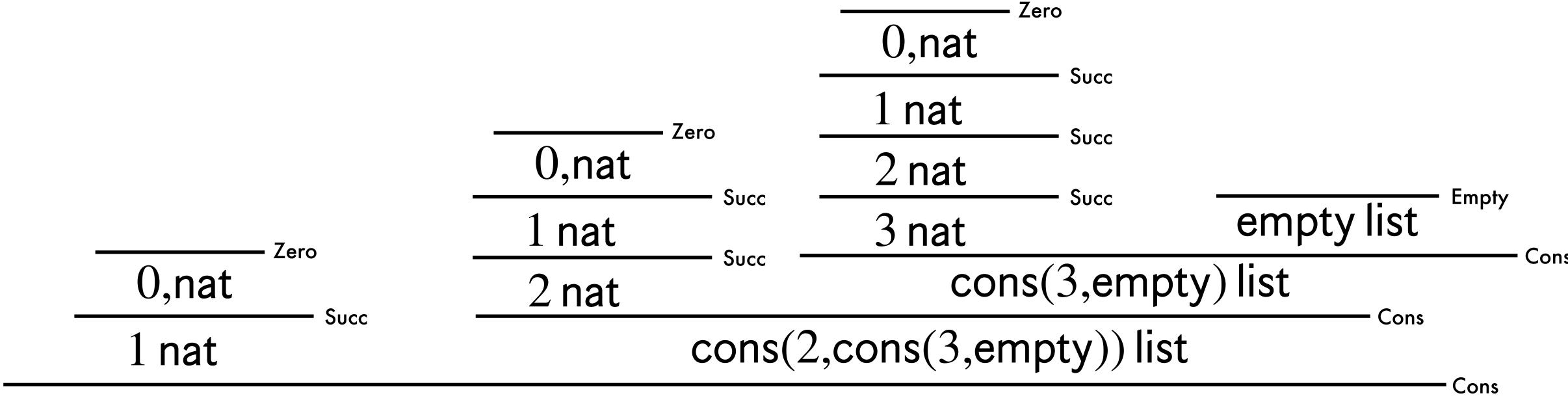












Suppose we want to prove a property P holds for judgment J.

Then for each rule defining J:

$$A_1 \cdots A_n$$
Name

It is enough to assume  $P(A_1), ..., P(A_n)$  hold and show P(C).

$$\frac{a = b \text{ nat}}{0 = 0 \text{ nat}}$$
 
$$\frac{a = b \text{ nat}}{\text{succ}(a) = \text{succ}(b) \text{ nat}}$$

- 1. P(0 = 0 nat)
- 2. if a = b nat and P(a = b nat), then succ(a) = succ(b) nat and P(succ(a) = succ(b) nat)

Lemma. If a nat, then a = a nat

$$\frac{a = b \text{ nat}}{0 = 0 \text{ nat}}$$
 
$$\frac{a = b \text{ nat}}{\text{succ}(a) = \text{succ}(b) \text{ nat}}$$

Lemma. If a nat, then a = a nat

Proof. By rule induction on a nat.

Rule Zero: In this case, we know a=0, and by rule EqZero we know 0=0 nat.

Rule Succ: We know that  $a = \operatorname{succ}(b)$  and b nat by assumption. By the induction hypothesis we know b = b nat, and by applying the EqSucc rule we know  $\operatorname{succ}(b) = \operatorname{succ}(b)$  nat.

Lemma. If succ(a) = succ(b) nat, then a = b nat

Lemma. If succ(a) = succ(b) nat, then a = b nat

Proof. By rule induction on succ(a) = succ(b) nat.

Rule EqZero: In this case we need  $0 = \operatorname{succ}(a)$ , but this is impossible.

Rule EqSucc: We know that a=b nat by the induction hypothesis.