```
termvar, name, x, y, z, xs
 indecies,\ i,\ j,\ k
 exp, e
                                                                                     Expressions
                                             num[n]
                                             str[s]
                                             \mathsf{plus}\,(\mathit{e}_1;\mathit{e}_2)
                                             \mathsf{mult}\left(e_1;e_2\right)
                                             cat(e_1; e_2)
                                             len(e)
                                             let (e_1; x.e_2)
 type, T
                                                                                    Types
                                             Str
                                             Num
                                                                            S
                                             (T)
 Γ
                                                                                     Typing Contexts
                                             \emptyset
                                             x:T
                                             \Gamma_1, \Gamma_2
 {\cal E}
                                                                                     Evaluation Contexts
                                 ::=
                                             \mathsf{plus}\left(\mathcal{E};e_{2}
ight)
                                             \mathsf{plus}\left(e_1;\mathcal{E}\right)
                                             \mathsf{mult}\left(\mathcal{E}; \mathit{e}_{2}\right)
                                             \mathsf{mult}\left(e_1;\mathcal{E}\right)
                                             \mathsf{cat}\left(\mathcal{E};e_{2}
ight)
                                             \mathsf{cat}\left(e_1;\mathcal{E}
ight)
                                             \mathsf{len}\left(\mathcal{E}\right)
                                             let (\mathcal{E}; x.e_2)
                                             let (e_1; x.\mathcal{E})
                                                                            S
                                             \mathcal{E}[e]
 nat, n
                                                                                    Nats
                                 ::=
                                                                                     Character
 char, c
                                 ::=
 strings, s
                                 ::=
\Gamma \vdash e : T
                            Typing
                                                                                                                        Var
                                                                                \overline{\Gamma,x:T\vdash x:T}
                                                                                                                       \operatorname{Str}
                                                                                  \overline{\Gamma \vdash \mathsf{str}[s] : \mathsf{Str}}
                                                                              \overline{\Gamma \vdash \mathsf{num}[n] : \mathsf{Num}}
                                                                 \frac{\Gamma \vdash e_1 : \mathsf{Num} \quad \Gamma \vdash e_2 : \mathsf{Num}}{\Gamma \vdash \mathsf{plus}\left(e_1; e_2\right) : \mathsf{Num}}
```

$$\begin{array}{c} \Gamma \vdash e_1 : \mathsf{Num} \quad \Gamma \vdash e_2 : \mathsf{Num} \\ \hline \Gamma \vdash \mathsf{mult} \left(e_1; e_2\right) : \mathsf{Num} \\ \hline \\ \frac{\Gamma \vdash e : \mathsf{Str}}{\Gamma \vdash \mathsf{len} \left(e\right) : \mathsf{Num}} \quad \mathsf{LENGTH} \\ \hline \\ \frac{\Gamma \vdash e_1 : \mathsf{Str} \quad \Gamma \vdash e_2 : \mathsf{Str}}{\Gamma \vdash \mathsf{cat} \left(e_1; e_2\right) : \mathsf{Str}} \quad \mathsf{CAT} \\ \hline \\ \frac{\Gamma \vdash e_1 : T_1 \quad \Gamma, x : T_1 \vdash e_2 : T_2}{\Gamma \vdash \mathsf{let} \left(e_1; x . e_2\right) : T_2} \quad \mathsf{LET} \end{array}$$

e val Values

 $e_1 \mapsto e_2$ Evaluation

$$\begin{array}{c} n_1 + n_2 = n \text{ nat} \\ \hline \text{plus} (\text{num}[n_1]; \text{num}[n_2]) \mapsto \text{num}[n] \\ \hline \\ \frac{e_4 \mapsto e_4''}{\text{plus} (e_4; e_5) \mapsto \text{plus} (e_4''; e_5)} \\ \hline \\ \frac{e_1 \text{ val} \quad e_2 \mapsto e_2'}{\text{plus} (e_1; e_2) \mapsto \text{plus} (e_1; e_2')} \\ \hline \\ \frac{e_1 \mapsto e_1'}{\text{mult} (e_1; e_2) \mapsto \text{mult} (e_1'; e_2)} \\ \hline \\ \frac{e_1 \mapsto e_1'}{\text{mult} (e_1; e_2) \mapsto \text{mult} (e_1; e_2')} \\ \hline \\ \frac{n_1 * n_2 = n \text{ nat}}{\text{mult} (\text{num}[n_1]; \text{num}[n_2]) \mapsto \text{num}[n]} \\ \hline \\ \frac{s_1 s_2 = s \text{ str}}{\text{cat} (\text{str}[s_1]; \text{str}[s_2]) \mapsto \text{str}[s]} \\ \hline \\ \frac{e_1 \mapsto e_1'}{\text{cat} (e_1; e_2) \mapsto \text{cat} (e_1'; e_2)} \\ \hline \\ \frac{e_1 \mapsto e_1'}{\text{cat} (e_1; e_2) \mapsto \text{cat} (e_1; e_2')} \\ \hline \\ \frac{|s| = n \text{ num}}{\text{len} (\text{str}[s]) \mapsto \text{num}[n]} \\ \hline \\ \frac{e \mapsto e'}{\text{len} (e) \mapsto \text{len} (e')} \\ \hline \\ \frac{e_1 \text{ val}}{\text{let} (e_1; x. e_2) \mapsto \text{let} (e_1'; x. e_2)} \\ \hline \\ \text{LETVAL} \\ \hline \\ \frac{e_1 \mapsto e_1'}{\text{let} (e_1; x. e_2) \mapsto \text{let} (e_1'; x. e_2)} \\ \hline \\ \text{LETVAL} \\ \hline \\ \text{LETVAL} \\ \hline \\ \hline \\ \text{LETVAL} \\ \hline \\ \text{LETVAL} \\ \hline \\ \hline \\ \text{LETVAL} \\ \hline \\ \hline \\ \text{LETVAL} \\ \hline \\ \text{LETVAL} \\ \hline \\ \hline \\ \end{array}$$

$$\frac{1}{\text{let}(e_1; x.e_2) \mapsto [e_1/x]e_2}$$
 LetL

 $e_1 \mapsto^* e_2$ Multistep Evaluation

$$\frac{e_1 \mapsto e_2 \quad \text{Refl}}{e_1 \mapsto e_2 \quad e_2 \mapsto^* e_3}$$

$$\frac{e_1 \mapsto e_2 \quad e_2 \mapsto^* e_3}{e_1 \mapsto^* e_3} \quad \text{Step}$$

 $e = \mathcal{E}$ Contextual Equality

 $e_1 \mapsto_{\mathsf{S}} e_2$ Single Step Evaluation

$$\begin{split} \frac{n_1 + n_2 = n \text{ nat}}{\text{plus} \left(\text{num}[n_1]; \text{num}[n_2] \right) \mapsto_{\text{S}} \text{num}[n]} & \text{S_PLUSVAL} \\ \frac{n_1 * n_2 = n \text{ nat}}{\text{mult} \left(\text{num}[n_1]; \text{num}[n_2] \right) \mapsto_{\text{S}} \text{num}[n]} & \text{S_MULTVAL} \\ \frac{s_1 s_2 = s \text{ str}}{\text{cat} \left(\text{str}[s_1]; \text{str}[s_2] \right) \mapsto_{\text{S}} \text{str}[s]} & \text{S_CATVAL} \\ \frac{|s| = n \text{ num}}{\text{len} \left(\text{str}[s] \right) \mapsto_{\text{S}} \text{num}[n]} & \text{S_LENGTHVAL} \\ \frac{e_1 \text{ val}}{\text{let} \left(e_1; x. e_2 \right) \mapsto_{\text{S}} \left[e_1 / x \right] e_2} & \text{S_LETVAL} \end{split}$$

 $e_1 \mapsto e_2$ Evaluation Contexts

$$\frac{e = \mathcal{E}[e_1] \quad e_1 \mapsto_{\mathsf{S}} e_2 \quad e' = \mathcal{E}[e_2]}{e \mapsto e'} \quad \text{C_STEP}$$

$\Gamma \vdash e_1 \equiv e_2 : T$ Equality

$$\begin{array}{c} \Gamma \vdash e \equiv e : T \\ \hline \Gamma \vdash e \equiv e : T \\ \hline \Gamma \vdash e_1 \equiv e_2 : T \\ \hline \Gamma \vdash e_1 \equiv e_2 : T \\ \hline \Gamma \vdash e_1 \equiv e_3 : T \\ \hline \Gamma \vdash e_1 \equiv e_3 : T \\ \hline \Gamma \vdash e_1 \equiv e_3 : T \\ \hline \Gamma \vdash e_1 \equiv e_3 : T \\ \hline \Gamma \vdash \text{plus} \left(\text{num}[n_1]; \text{num}[n_2] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{plus} \left(\text{num}[n_1]; \text{num}[n_2] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{plus} \left(e_1; e_2 \right) \equiv \text{plus} \left(e_1'; e_2' \right) : \text{Num} \\ \hline \Gamma \vdash \text{plus} \left(e_1; e_2 \right) \equiv \text{plus} \left(e_1'; e_2' \right) : \text{Num} \\ \hline \Gamma \vdash \text{mult} \left(\text{num}[n_1]; \text{num}[n_2] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{mult} \left(\text{num}[n_1]; \text{num}[n_2] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{mult} \left(e_1; e_2 \right) \equiv \text{mult} \left(e_1'; e_2' \right) : \text{Num} \\ \hline \Gamma \vdash \text{mult} \left(e_1; e_2 \right) \equiv \text{mult} \left(e_1'; e_2' \right) : \text{Num} \\ \hline \Gamma \vdash \text{cat} \left(\text{str}[s_1]; \text{str}[s_2] \right) \equiv \text{str}[s] : \text{Str} \\ \hline \Gamma \vdash \text{cat} \left(e_1; e_2 \right) \equiv \text{cat} \left(e_1'; e_2' \right) : \text{Str} \\ \hline \Gamma \vdash \text{cat} \left(e_1; e_2 \right) \equiv \text{cat} \left(e_1'; e_2' \right) : \text{Str} \\ \hline \Gamma \vdash \text{len} \left(\text{str}[s] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{len} \left(\text{str}[s] \right) \equiv \text{num}[n] : \text{Num} \\ \hline \Gamma \vdash \text{len} \left(e \right) \equiv \text{len} \left(e' \right) : \text{Num} \\ \hline \Gamma \vdash \text{len} \left(e \right) \equiv \text{len} \left(e' \right) : \text{Num} \\ \hline \Gamma \vdash \text{len} \left(e \right) \equiv \text{len} \left(e' \right) : \text{Num} \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right) : T_2 \\ \hline \Gamma \vdash \text{let} \left(e_1 \right) = \text{let} \left(e_1'; x. e_2 \right) \equiv \text{let} \left(e_1'; x. e_2 \right$$

Definition rules: 54 good 0 bad Definition rule clauses: 98 good 0 bad