# A Model of Co-intuitionistic Linear Logic in Dialectica Categories

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Abstract. TODO

#### 1 Introduction

TODO [?]

# 2 Symmetric Monoidal Co-closed Categories

In this section we recall the definition of a symmetric monoidal category [?], and perhaps the lesser known definition of a symmetric monoidal left-closed category [?].

**Definition 1.** A symmetric monoidal category (SMC) is a category, M, with the following data:

- An object I of M,
- *A bi-functor*  $\oplus$  :  $M \times M$  → M,
- The following natural isomorphisms:

$$\lambda_A : I \oplus A \longrightarrow A$$

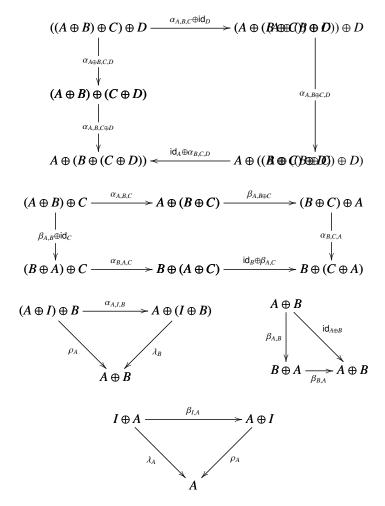
$$\rho_A : A \oplus I \longrightarrow A$$

$$\alpha_{A,B,C} : (A \oplus B) \oplus C \longrightarrow A \oplus (B \oplus C)$$

- A symmetry natural transformation:

$$\beta_{A,B}: A \oplus B \longrightarrow B \oplus A$$

- Subject to the following coherence diagrams:



**Definition 2.** A symmetric monoidal left closed category (SMLCC) is a symmetric monoidal category  $(C, \oplus, I, \alpha_{A,B,C}, \lambda_A, \rho_A, \beta_{A,B})$  such that the for any object B of C the functor  $-\oplus B: C \to C$  has a specified left adjoint  $B \leftarrow -: C \to C$ . This means that for any objects A,B, and C of C, we have the following bijection:

$$\operatorname{\mathsf{Hom}}_{\mathcal{C}}(C,A\oplus B)\cong \operatorname{\mathsf{Hom}}_{\mathcal{C}}(B \longleftarrow C,A)$$

that is natural in all arguments.

As we can see from the definition a SMLCC is the categorical dual to symmetric monoidal closed categories. Just as SMCCs model intuitionistic linear logics [?] SMLCCs model co-intuitionistic linear logics [?]. There are many concrete examples of SMLCCs, in fact, the dual of any interesting concrete SMCCs must be an interesting SMLCC.

In this paper I make use of a specific SMLCC, Set<sup>op</sup>, which is the dual of the category Set of all sets and functions between them. The definition of Set<sup>op</sup> I use here is the well-known one in power-set algebras  $(\mathcal{P}(X), \cup, \cap, \overline{\phantom{A}}, \emptyset, X)$  where X is a set. In fact, it is well-known that we can define the functor  $\mathcal{P}: \text{Set} \to \text{Set}^{\text{op}}$  on morphisms by  $\mathcal{P}(f: X \to Y)(S \in \mathcal{P}(Y)) = \{x \in X \mid f(x) \in S\}$ . Thus, objects of Set<sup>op</sup> are powersets and morphisms are set theoretic functions between powersets.

## 3 The category Dial<sub>L</sub>(Sets<sup>op</sup>)

Dialectica categories originate from de Paiva's thesis [1], and are one of the first models of intuitionistic linear logic. In fact, they are the first categorical model of full intuitionistic linear logic which is intuitionistic linear logic with every logical connective from linear logic; complete with multiple conclusions. In full generality, dialectica categories, denoted  $Dial_L(C)$ , are symmetric monoidal closed categories constructed in terms of a lineale [?], L, and a symmetric monoidal closed category C. However, I take L to be a colineale (Definition ??) and C to be a symmetric monoidal left-closed category (Definition 2), particularly, I take C to be  $Set^{op}$ , but I take care to insure that all constructions lift to the general case of an arbitrary symmetric monoidal left-closed category. This is the first time this construction has been given.

We begin by introducing colineales as the categorical dual to lineales. The following defines when a proset is symmetric monoidal.

**Definition 3.** A monoidal proset is a proset,  $(L, \leq)$ , with a given symmetric monoidal structure  $(L, \bullet, e)$ . That is, a set L with a given binary relation  $\leq$ :  $L \times L \rightarrow L$  satisfying the following:

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- (reflexivity) a \le a for any a \in L
- (transitivity) If a \le b and b \le c, then a \le c
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together with a monoidal structure  $(\bullet, e)$  consisting of a binary operation, called multiplication,  $\bullet: L \times L \to L$  and a distinguished element  $e \in L$  called the unit such that the following hold:

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- (associativity) (a \bullet b) \bullet c = a \bullet (b \bullet c)

- (identity) a \bullet e = a = e \bullet a

- (symmetry) a \bullet b = b \bullet a
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Finally, the structures must be compatible, that is, if  $a \le b$ , then  $a \bullet c \le b \bullet c$  for any  $c \in L$ .

A colineale is essentially a symmetric monoidal left-closed category in the category of prosets.

**Definition 4.** A *colineale* is a monoidal proset,  $(L, \leq, \bullet, e)$ , with a given binary operation, called coimplication,  $\bullet$ :  $L \times L \to L$ , such that the following hold:

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- (relative complement) b \le a \bullet (a \bullet b)
- (adjunction) If b \bullet c \le a, then c \le a \bullet b
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An example of a concrete colineale is the three element set  $3 = \{0, \bot, 1\}$  where  $\bot$  stands for undefined<sup>1</sup>. However, one must be careful when defining the colineale because it is possible to degenerate to classical logic.

### References

1. Valeria de Paiva. Dialectica categories. In J. Gray and A. Scedrov, editors, *Categories in Computer Science and Logic*, volume 92, pages 47–62. Amerian Mathematical Society, 1989.

<sup>&</sup>lt;sup>1</sup> The full definition of the colineale 3 can be found in the formal development: ?