Monoidal-Annex

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1 Monoidal Categories

Definition 1. A monoidal category $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho)$ is a category \mathbb{C} , a bifunctor $\otimes : \mathbb{C} \times \mathbb{C} \to \mathbb{C}$, an object $I \in \mathbb{C}$, and three natural isomorphisms α, λ, ρ . Where,

$$\alpha = \alpha_{A,B,C} : A \otimes (B \otimes C) \cong (A \otimes B) \otimes C$$

is natural for all $A, B, C \in \mathbb{C}$ and the diagram

$$A \otimes (B \otimes (C \otimes D)) \xrightarrow{\alpha} (A \otimes B) \otimes (C \otimes D) \xrightarrow{\alpha} ((A \otimes B) \otimes C) \otimes D$$

$$\uparrow_{0 \otimes 1}$$

$$A \otimes ((B \otimes C) \otimes D) \xrightarrow{\alpha} (A \otimes (B \otimes C)) \otimes D$$

commutes for all $A, B, C, D \in \mathbb{C}$. γ and ρ are natural

$$\gamma_A: I \otimes A \cong A \quad \rho_A: A \otimes I \cong A$$

for all objects $A \in \mathbb{C}$, the diagram

$$\begin{array}{c|c} A\otimes (I\otimes C) & \xrightarrow{\alpha} & (A\otimes I)\otimes C \\ & \downarrow^{\rho_A\otimes 1} & & \downarrow^{\rho_A\otimes 1} \\ & A\otimes C & & & & A\otimes C \end{array}$$

[2]

Definition 2. A monoidal functor or lax monoidal functor (F, m) between monoidal categories (\mathbb{C}, \otimes, I) and $(\mathbb{D}, \otimes', I')$ is a functor $F : \mathbb{C} \to \mathbb{D}$ equipped with nathral transformations

$$m'_{A,B}:FA\otimes'FB o F(A\otimes B)\quad m'':I' o FI$$

where the following diagrams commute in the category \mathbb{D} for all objects $A, B, C \in \mathbb{C}$

$$(FA \otimes' FB) \otimes' FC \xrightarrow{\alpha'} FA \otimes' (FB \otimes' FC)$$

$$\downarrow_{FA \otimes m}$$

$$F(A \otimes B) \otimes' FC \qquad FA \otimes' F(B \otimes C)$$

$$\downarrow_{m}$$

$$\downarrow_{m}$$

$$F((A \otimes B) \otimes C) \xrightarrow{F\alpha} F(A \otimes (B \otimes C))$$

$$FA \otimes' I' \xrightarrow{\rho'} FA \qquad I' \otimes' FB \xrightarrow{\gamma'} FB$$

$$\downarrow_{FA \otimes' m}$$

$$\downarrow_{FA \otimes' m}$$

$$\downarrow_{FA \otimes' FB}$$

$$\downarrow_{$$

[3]

Definition 3. An oplax(colax/comonoidal) monoidal functor (F, n) between monoidal categories (\mathbb{C}, \otimes, I) and $(\mathbb{D}, \otimes', I')$ is a functor $F : \mathbb{C} \to \mathbb{D}$ and natural transformations

$$n'_{AB}: F(A \otimes B) \to FA \otimes' FB \quad n'': FI \to I'$$

where the following diagrams commute in the category \mathbb{D} , for all objects $A, B, C \in \mathbb{C}$

$$F((A \otimes B) \otimes C) \xrightarrow{F\alpha} F(A \otimes (B \otimes C))$$

$$\downarrow n$$

$$\downarrow r$$

$$F(A \otimes B) \otimes' FC \qquad FA \otimes' F(B \otimes C)$$

$$\downarrow rA \otimes' FC \qquad \downarrow rA \otimes' n$$

$$(FA \otimes' FB) \otimes' FC \xrightarrow{\alpha'} FA \otimes' (FB \otimes' FC)$$

$$F(A \otimes I) \xrightarrow{F\rho} FA \qquad F(I \otimes B) \xrightarrow{F\gamma} FB$$

$$\uparrow n \qquad \uparrow n \qquad \uparrow \gamma'$$

$$\downarrow rA \otimes' FI \xrightarrow{FA \otimes' n} FA \otimes' I' \qquad FI \otimes' FB \xrightarrow{n \otimes' FB} I' \otimes' FB$$

[3]

Definition 4. (F, m) and (G, n) are (lax) monoidal functors between the monoidal categories:

$$(\mathbb{C}, \otimes, I) \to (\mathbb{D}, \otimes', I')$$

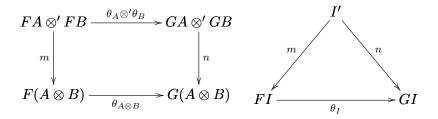
A monoidal natural transformation

$$\theta: (F,m) \Rightarrow (G,n): (\mathbb{C},\otimes,I) \to (\mathbb{D},\otimes',I')$$

between the monoidal functors (F, m) and (G, n) is a naural transformation

$$\theta: F \Rightarrow G: \mathbb{C} \to \mathbb{D}$$

between the underlying functors, where the following diamgrams commute, for all objects $A, B \in \mathbb{C}$



[3]

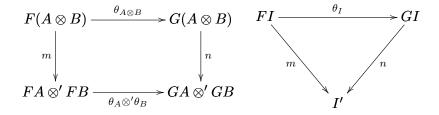
Definition 5. A monoidal natural transformation

$$\theta: (F,m) \Rightarrow (G,n): (\mathbb{C},\otimes,I) \to (\mathbb{D},\otimes',I')$$

between two oplax(colax/comonoidal) monoidal functors (F, m) and (G, n) is a natural transformation

$$\theta: F \Rightarrow G: \mathbb{C} \to \mathbb{D}$$

between the underlying functors, where the following diagrams commute, for all objects $A, B \in \mathbb{C}$



[3]

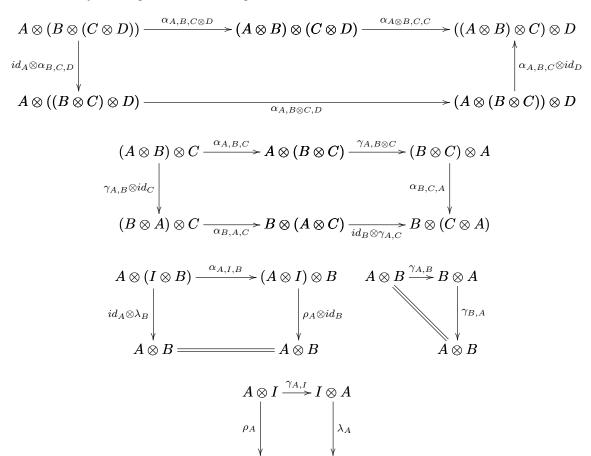
2 Symmetric Monoidal Categories

Definition 6. A symmetric monoidal catergory (SMC), $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho, \gamma)$, is a category \mathbb{C} equippped with a bifunctor $\otimes : \mathbb{C} \times \mathbb{C} \to \mathbb{C}$ with a neutral element I and natural isomorphisms α, λ, ρ , and γ :

- 1. $\alpha_{A,B,C}: A \otimes (B \otimes C) \xrightarrow{\sim} (A \otimes B) \otimes C$
- 2. $\lambda_A: I \otimes A \xrightarrow{\sim} A$
- 3. $\rho_A: A\otimes I \xrightarrow{\sim} A$

4.
$$\gamma_{A,B}: A \otimes B \xrightarrow{\sim} B \otimes A$$

which make the following 'coherence' diagrams commute.



The following equality is also require to hold:

$$\lambda_I = \rho_I : I \otimes I \to I$$

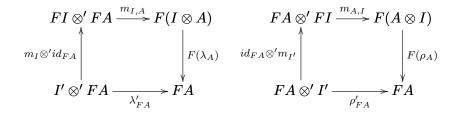
[1]

Definition 7. A symmetric monoidal closed category (SMCC), $(\mathbb{C}, \otimes, \neg, I, \alpha, \lambda, \rho, \gamma)$, is a SMC such that for all objects A in \mathbb{C} , the functor $-\otimes A$ has a specified right adjoint $A \multimap -$.

Definition 8. A symmetric monoidal functor between SMCs $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho, \gamma)$ and $(\mathbb{C}', \otimes', I', \alpha', \lambda', \rho', \gamma')$ is a functor $F : \mathbb{C} \to \mathbb{C}'$ equipped with

- 1. A morphism $m_{I'}: I' \to FI$.
- 2. For any two objects A and B in \mathbb{C} , a natural transformation $m_{A,B}: F(A) \otimes' F(B) \to F(A \otimes B)$

These must satisfy the following diagrams:



$$(FA \otimes' FB) \otimes' FC \xrightarrow{m_{A,B} \otimes' id_{FC}} F(A \otimes B) \otimes' FC \xrightarrow{m_{A \otimes B,C}} F((A \otimes B) \otimes C)$$

$$\alpha'_{FA,FB,FC} \downarrow \qquad \qquad \uparrow F(\alpha_{A,B,C})$$

$$FA \otimes' (FB \otimes' FC)_{id_{FA} \otimes' m_{B,C}} FA \otimes' F(B \otimes C) \xrightarrow{m_{A,B \otimes C}} F(A \otimes (B \otimes C))$$

$$FA \otimes' FB \xrightarrow{m_{A,B}} F(A \otimes B)$$

$$\uparrow_{A,B} \qquad \qquad \downarrow F(\gamma_{A,B})$$

$$FB \otimes' FA \xrightarrow{m_{B,A}} F(B \otimes A)$$

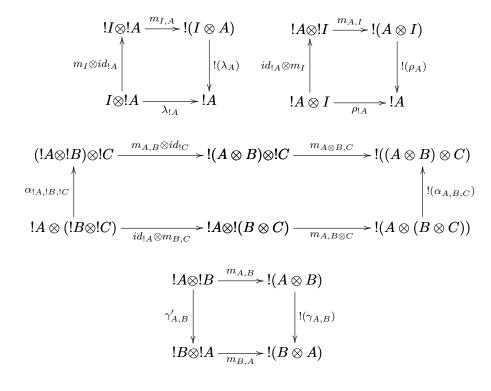
However in this particular case, assuming that ! is a symmetric monoidal (endo) functor means that ! comes equipped with a natural transformation

$$m_{A,B}: !A \otimes !B \rightarrow !(A \otimes B)$$

 $and\ a\ morphism$

$$m_I:I\to !I$$

(where m_I is just the nullary version of the natural transformation.) The diagrams given in the above definition become the following:



[1]

Definition 9. A symmetric monoidal functor, $(F, m_{A,B}, m_{I'}) : \mathbb{C} \to \mathbb{C}'$, is said to be

- 1. Strict if $m_{A,B}$ and $m_{I'}$ are identities.
- 2. Strong if $m_{A,B}$ and $m_{I'}$ are natural isomorphisms.

[1]

References

- [1] G.M. Bierman. On Intuitionistic Linear Logic. PhD thesis, University of Cambridge, 1993.
- [2] Saunders Mac Lane. Categories for the Working Mathematician, volume 5 of Graduate Texts in Mathematics. Springer-Verlag, New York, 1971.
- [3] Paul-André Mellies. Categorical semantics of linear logic. *Panoramas et syntheses*, 27:15–215, 2009.