

Monoidal-Annex

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1 Monoidal Categories

Definition 1. A monoidal category $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho)$ is a category \mathbb{C} , a bifunctor $\otimes : \mathbb{C} \times \mathbb{C} \rightarrow \mathbb{C}$, an object $I \in \mathbb{C}$, and three natural isomorphisms α, λ, ρ . Where,

$$\alpha = \alpha_{A,B,C} : A \otimes (B \otimes C) \cong (A \otimes B) \otimes C$$

is natural for all $A, B, C \in \mathbb{C}$ and the diagram

$$\begin{array}{ccccc}
 A \otimes (B \otimes (C \otimes D)) & \xrightarrow{\alpha} & (A \otimes B) \otimes (C \otimes D) & \xrightarrow{\alpha} & ((A \otimes B) \otimes C) \otimes D \\
 \downarrow 1 \otimes \alpha & & & & \uparrow \alpha \otimes 1 \\
 A \otimes ((B \otimes C) \otimes D) & \xrightarrow{\alpha} & & & (A \otimes (B \otimes C)) \otimes D
 \end{array}$$

commutes for all $A, B, C, D \in \mathbb{C}$. γ and ρ are natural

$$\gamma_A : I \otimes A \cong A \quad \rho_A : A \otimes I \cong A$$

for all objects $A \in \mathbb{C}$, the diagram

$$\begin{array}{ccc}
 A \otimes (I \otimes C) & \xrightarrow{\alpha} & (A \otimes I) \otimes C \\
 \downarrow 1 \otimes \lambda & & \downarrow \rho_A \otimes 1 \\
 A \otimes C & \xlongequal{\quad} & A \otimes C
 \end{array}$$

[2]

2 Symmetric Monoidal Categories

Definition 2. A symmetric monoidal category (SMC), $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho, \gamma)$, is a category \mathbb{C} equipped with a bifunctor $\otimes : \mathbb{C} \times \mathbb{C} \rightarrow \mathbb{C}$ with a neutral element I and natural isomorphisms α, λ, ρ , and γ :

1. $\alpha_{A,B,C} : A \otimes (B \otimes C) \xrightarrow{\sim} (A \otimes B) \otimes C$
2. $\lambda_A : I \otimes A \xrightarrow{\sim} A$
3. $\rho_A : A \otimes I \xrightarrow{\sim} A$
4. $\gamma_{A,B} : A \otimes B \xrightarrow{\sim} B \otimes A$

which make the following 'coherence' diagrams commute.

$$\begin{array}{ccccc}
 A \otimes (B \otimes (C \otimes D)) & \xrightarrow{\alpha_{A,B,C \otimes D}} & (A \otimes B) \otimes (C \otimes D) & \xrightarrow{\alpha_{A \otimes B,C,D}} & ((A \otimes B) \otimes C) \otimes D \\
 \downarrow id_A \otimes \alpha_{B,C,D} & & & & \uparrow \alpha_{A,B,C} \otimes id_D \\
 A \otimes ((B \otimes C) \otimes D) & \xrightarrow{\alpha_{A,B \otimes C,D}} & & & (A \otimes (B \otimes C)) \otimes D
 \end{array}$$

$$\begin{array}{ccccc}
 (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\gamma_{A,B \otimes C}} & (B \otimes C) \otimes A \\
 \downarrow \gamma_{A,B} \otimes id_C & & & & \downarrow \alpha_{B,C,A} \\
 (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{id_B \otimes \gamma_{A,C}} & B \otimes (C \otimes A)
 \end{array}$$

$$\begin{array}{ccccc}
 A \otimes (I \otimes B) & \xrightarrow{\alpha_{A,I,B}} & (A \otimes I) \otimes B & & A \otimes B \xrightarrow{\gamma_{A,B}} B \otimes A \\
 \downarrow id_A \otimes \lambda_B & & \downarrow \rho_A \otimes id_B & & \downarrow \gamma_{B,A} \\
 A \otimes B & \xlongequal{\quad} & A \otimes B & & A \otimes B
 \end{array}$$

$$\begin{array}{ccc}
 A \otimes I & \xrightarrow{\gamma_{A,I}} & I \otimes A \\
 \downarrow \rho_A & & \downarrow \lambda_A \\
 A & \xlongequal{\quad} & A
 \end{array}$$

The following equality is also require to hold:

$$\lambda_I = \rho_I : I \otimes I \rightarrow I$$

[1]

Definition 3. A symmetric monoidal closed category (SMCC), $(\mathbb{C}, \otimes, \multimap, I, \alpha, \lambda, \rho, \gamma)$, is a SMC such that for all objects A in \mathbb{C} , the functor $- \otimes A$ has a specified right adjoint $A \multimap -$.

Definition 4. A symmetric monoidal functor between SMCs $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho, \gamma)$ and $(\mathbb{C}', \otimes', I', \alpha', \lambda', \rho', \gamma')$ is a functor $F : \mathbb{C} \rightarrow \mathbb{C}'$ equipped with

1. A morphism $m_{I'} : I' \rightarrow FI$.
2. For any two objects A and B in \mathbb{C} , a natural transformation $m_{A,B} : F(A) \otimes' F(B) \rightarrow F(A \otimes B)$

These must satisfy the following diagrams:

$$\begin{array}{ccc}
 FI \otimes' FA & \xrightarrow{m_{I,A}} & F(I \otimes A) \\
 \uparrow m_{I'} \otimes' id_{FA} & & \downarrow F(\lambda_A) \\
 I' \otimes' FA & \xrightarrow{\lambda'_{FA}} & FA \\
 \end{array}
 \qquad
 \begin{array}{ccc}
 FA \otimes' FI & \xrightarrow{m_{A,I}} & F(A \otimes I) \\
 \uparrow id_{FA} \otimes' m_{I'} & & \downarrow F(\rho_A) \\
 FA \otimes' I' & \xrightarrow{\rho'_{FA}} & FA \\
 \end{array}$$

$$\begin{array}{ccccc}
 (FA \otimes' FB) \otimes' FC & \xrightarrow{m_{A,B} \otimes' id_{FC}} & F(A \otimes B) \otimes' FC & \xrightarrow{m_{A \otimes B, C}} & F((A \otimes B) \otimes C) \\
 \uparrow \alpha'_{FA, FB, FC} & & & & \uparrow F(\alpha_{A, B, C}) \\
 FA \otimes' (FB \otimes' FC) & \xrightarrow{id_{FA} \otimes' m_{B, C}} & FA \otimes' F(B \otimes C) & \xrightarrow{m_{A, B \otimes C}} & F(A \otimes (B \otimes C))
 \end{array}$$

$$\begin{array}{ccc}
 FA \otimes' FB & \xrightarrow{m_{A,B}} & F(A \otimes B) \\
 \downarrow \gamma'_{A,B} & & \downarrow F(\gamma_{A,B}) \\
 FB \otimes' FA & \xrightarrow{m_{B,A}} & F(B \otimes A)
 \end{array}$$

However in this particular case, assuming that $!$ is a symmetric monoidal (endo) functor means that $!$ comes equipped with a natural transformation

$$m_{A,B} : !A \otimes !B \rightarrow !(A \otimes B)$$

and a morphism

$$m_I : I \rightarrow !I$$

(where m_I is just the nullary version of the natural transformation.) The diagrams given in the above definition become the following:

$$\begin{array}{ccc}
!I \otimes !A & \xrightarrow{m_{I,A}} & !(I \otimes A) \\
\uparrow m_I \otimes id_{!A} & & \downarrow !(\lambda_A) \\
I \otimes !A & \xrightarrow{\lambda_{!A}} & !A
\end{array}
\quad
\begin{array}{ccc}
!A \otimes !I & \xrightarrow{m_{A,I}} & !(A \otimes I) \\
\uparrow id_{!A} \otimes m_I & & \downarrow !(\rho_A) \\
!A \otimes I & \xrightarrow{\rho_{!A}} & !A
\end{array}$$

$$\begin{array}{ccccc}
(!A \otimes !B) \otimes !C & \xrightarrow{m_{A,B} \otimes id_{!C}} & !(A \otimes B) \otimes !C & \xrightarrow{m_{A \otimes B, C}} & !((A \otimes B) \otimes C) \\
\uparrow \alpha_{!A, !B, !C} & & & & \uparrow !(\alpha_{A, B, C}) \\
!A \otimes (!B \otimes !C) & \xrightarrow{id_{!A} \otimes m_{B,C}} & !A \otimes !(B \otimes C) & \xrightarrow{m_{A, B \otimes C}} & !(A \otimes (B \otimes C))
\end{array}$$

$$\begin{array}{ccc}
!A \otimes !B & \xrightarrow{m_{A,B}} & !(A \otimes B) \\
\downarrow \gamma'_{A,B} & & \downarrow !(\gamma_{A,B}) \\
!B \otimes !A & \xrightarrow{m_{B,A}} & !(B \otimes A)
\end{array}$$

[1]

Definition 5. A symmetric monoidal functor, $(F, m_{A,B}, m_{I'}) : \mathbb{C} \rightarrow \mathbb{C}'$, is said to be

1. Strict if $m_{A,B}$ and $m_{I'}$ are identities.
2. Strong if $m_{A,B}$ and $m_{I'}$ are natural isomorphisms.

[1]

References

- [1] G.M. Bierman. *On Intuitionistic Linear Logic*. PhD thesis, University of Cambridge, 1993.
- [2] Saunders Mac Lane. *Categories for the Working Mathematician*, volume 5 of *Graduate Texts in Mathematics*. Springer-Verlag, New York, 1971.