

# Monoidal-Annex

Preston Keel

Dr. Harley Eades

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## 1 Monoidal Categories

**Definition 1.** A monoidal category  $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho)$  is a category  $\mathbb{C}$ , a bifunctor  $\otimes : \mathbb{C} \times \mathbb{C} \rightarrow \mathbb{C}$ , an object  $I \in \mathbb{C}$ , and three natural isomorphisms  $\alpha, \lambda, \rho$ . Where,

$$\alpha = \alpha_{A,B,C} : A \otimes (B \otimes C) \cong (A \otimes B) \otimes C$$

is natural for all  $A, B, C \in \mathbb{C}$  and the diagram

$$\begin{array}{ccccc}
 A \otimes (B \otimes (C \otimes D)) & \xrightarrow{\alpha} & (A \otimes B) \otimes (C \otimes D) & \xrightarrow{\alpha} & ((A \otimes B) \otimes C) \otimes D \\
 \downarrow 1 \otimes \alpha & & & & \uparrow \alpha \otimes 1 \\
 A \otimes ((B \otimes C) \otimes D) & \xrightarrow{\alpha} & & & (A \otimes (B \otimes C)) \otimes D
 \end{array}$$

commutes for all  $A, B, C, D \in \mathbb{C}$ .  $\gamma$  and  $\rho$  are natural

$$\gamma_A : I \otimes A \cong A \quad \rho_A : A \otimes I \cong A$$

for all objects  $A \in \mathbb{C}$ , the diagram

$$\begin{array}{ccc}
 A \otimes (I \otimes C) & \xrightarrow{\alpha} & (A \otimes I) \otimes C \\
 \downarrow 1 \otimes \lambda & & \downarrow \rho_A \otimes 1 \\
 A \otimes C & \xlongequal{\quad} & A \otimes C
 \end{array}$$

[3]

**Definition 2.** A monoidal functor or lax monoidal functor  $(F, m)$  between monoidal categories  $(\mathbb{C}, \otimes, I)$  and  $(\mathbb{D}, \otimes', I')$  is a functor  $F : \mathbb{C} \rightarrow \mathbb{D}$  equipped with a natural transformation

$$m_{A,B} : FA \otimes' FB \rightarrow F(A \otimes B)$$

and an isomorphism

$$m_I : I' \rightarrow FI$$

where the following diagrams commute in the category  $\mathbb{D}$  for all objects  $A, B, C \in \mathbb{C}$

$$\begin{array}{ccc}
(FA \otimes' FB) \otimes' FC & \xrightarrow{\alpha'} & FA \otimes' (FB \otimes' FC) \\
\downarrow m \otimes FC & & \downarrow FA \otimes m \\
F(A \otimes B) \otimes' FC & & FA \otimes' F(B \otimes C) \\
\downarrow m & & \downarrow m \\
F((A \otimes B) \otimes C) & \xrightarrow{F\alpha} & F(A \otimes (B \otimes C))
\end{array}$$
  

$$\begin{array}{ccc}
FA \otimes' I' & \xrightarrow{\rho'} & FA \\
\downarrow FA \otimes' m & & \uparrow F\rho \\
FA \otimes' FI & \xrightarrow{m} & F(A \otimes I)
\end{array}
\quad
\begin{array}{ccc}
I' \otimes' FB & \xrightarrow{\gamma'} & FB \\
\downarrow m \otimes' FB & & \uparrow F\gamma \\
FI \otimes' FB & \xrightarrow{m} & F(I \otimes B)
\end{array}$$

[4]

**Definition 3.** An oplax(colax/comonoidal) monoidal functor  $(F, m)$  between monoidal categories  $(\mathbb{C}, \otimes, I)$  and  $(\mathbb{D}, \otimes', I')$  is a functor  $F : \mathbb{C} \rightarrow \mathbb{D}$  with a natural transformation

$$m_{A,B} : F(A \otimes B) \rightarrow FA \otimes' FB$$

and an isomorphism

$$m_I : FI \rightarrow I'$$

where the following diagrams commute in the category  $\mathbb{D}$ , for all objects  $A, B, C \in \mathbb{C}$

$$\begin{array}{ccc}
F((A \otimes B) \otimes C) & \xrightarrow{F\alpha} & F(A \otimes (B \otimes C)) \\
\downarrow m & & \downarrow m \\
F(A \otimes B) \otimes' FC & & FA \otimes' F(B \otimes C) \\
\downarrow m \otimes' FC & & \downarrow FA \otimes' m \\
(FA \otimes' FB) \otimes' FC & \xrightarrow{\alpha'} & FA \otimes' (FB \otimes' FC)
\end{array}$$
  

$$\begin{array}{ccc}
F(A \otimes I) & \xrightarrow{F\rho} & FA \\
\downarrow m & & \uparrow \rho' \\
FA \otimes' FI & \xrightarrow{FA \otimes' m} & FA \otimes' I'
\end{array}
\quad
\begin{array}{ccc}
F(I \otimes B) & \xrightarrow{F\gamma} & FB \\
\downarrow m & & \uparrow \gamma' \\
FI \otimes' FB & \xrightarrow{m \otimes' FB} & I' \otimes' FB
\end{array}$$

[4]

**Definition 4.** Suppose  $(F, m)$  and  $(G, n)$  are (lax) monoidal functors between the monoidal categories:

$$(\mathbb{C}, \otimes, I) \rightarrow (\mathbb{D}, \otimes', I')$$

A monoidal natural transformation

$$\theta : (F, m) \Rightarrow (G, n) : (\mathbb{C}, \otimes, I) \rightarrow (\mathbb{D}, \otimes', I')$$

between the monoidal functors  $(F, m)$  and  $(G, n)$  is a natural transformation

$$\theta : F \Rightarrow G : \mathbb{C} \rightarrow \mathbb{D}$$

between the underlying functors, where the following diagrams commute, for all objects  $A, B \in \mathbb{C}$

$$\begin{array}{ccc}
FA \otimes' FB & \xrightarrow{\theta_A \otimes' \theta_B} & GA \otimes' GB \\
\downarrow m & & \downarrow n \\
F(A \otimes B) & \xrightarrow{\theta_{A \otimes B}} & G(A \otimes B)
\end{array}
\quad
\begin{array}{ccc}
& I' & \\
m \swarrow & & \searrow n \\
FI & \xrightarrow{\theta_I} & GI
\end{array}$$

[4]

**Definition 5.** A monoidal natural transformation

$$\theta : (F, m) \Rightarrow (G, n) : (\mathbb{C}, \otimes, I) \rightarrow (\mathbb{D}, \otimes', I')$$

between two oplax(colax/comonoidal) monoidal functors  $(F, m)$  and  $(G, n)$  is a natural transformation

$$\theta : F \Rightarrow G : \mathbb{C} \rightarrow \mathbb{D}$$

between the underlying functors, where the following diagrams commute, for all objects  $A, B \in \mathbb{C}$

$$\begin{array}{ccc}
 F(A \otimes B) & \xrightarrow{\theta_{A \otimes B}} & G(A \otimes B) \\
 \downarrow m & & \downarrow n \\
 FA \otimes' FB & \xrightarrow{\theta_A \otimes' \theta_B} & GA \otimes' GB
 \end{array}
 \quad
 \begin{array}{ccc}
 FI & \xrightarrow{\theta_I} & GI \\
 \downarrow m & & \downarrow n \\
 & I' &
 \end{array}$$

[4]

**Definition 6.** A monoidal category  $\mathbb{C}$  is said to be biclosed if every  $- \otimes Y$  has a right adjoint  $[Y, -]$  and every  $X \otimes -$  has a right adjoint  $[X, -]$  [2]

**Definition 7.** Given a pair of (lax) monoidal functors:

$$(F_*, m) : (\mathbb{C}, \otimes, I) \rightarrow (\mathbb{D}, \otimes', I') \quad (F^*, n) : (\mathbb{D}, \otimes', I') \rightarrow (\mathbb{C}, \otimes, I).$$

A monoidal adjunction

$$(F_*, m) \dashv (F^*, n)$$

between the monoidal functors is defined as an adjunction  $(F_*, F^*, \eta, \epsilon)$  between the underlying functors

$$F_* : \mathbb{C} \rightarrow \mathbb{D} \quad F^* : \mathbb{D} \rightarrow \mathbb{C}$$

whose natural transformations

$$\eta : id_{\mathbb{C}} \Rightarrow F^* \circ F_* \quad \epsilon : F_* \circ F^* \Rightarrow id_{\mathbb{D}}$$

are monoidal [4]

## 2 Symmetric Monoidal Categories

**Definition 8.** A symmetric monoidal category (SMC), is a monoidal category  $\mathbb{C}$  with an additional natural isomorphism  $\gamma$

$$\gamma_{A,B} : A \otimes B \xrightarrow{\sim} B \otimes A$$

and three additional 'coherence' diagrams that commute.

$$\begin{array}{ccccc}
 (A \otimes B) \otimes C & \xrightarrow{\alpha_{A,B,C}} & A \otimes (B \otimes C) & \xrightarrow{\gamma_{A,B \otimes C}} & (B \otimes C) \otimes A \\
 \downarrow \gamma_{A,B} \otimes id_C & & & & \downarrow \alpha_{B,C,A} \\
 (B \otimes A) \otimes C & \xrightarrow{\alpha_{B,A,C}} & B \otimes (A \otimes C) & \xrightarrow{id_B \otimes \gamma_{A,C}} & B \otimes (C \otimes A)
 \end{array}$$

$$\begin{array}{ccc}
 A \otimes B & \xrightarrow{\gamma_{A,B}} & B \otimes A \\
 & \searrow & \downarrow \gamma_{B,A} \\
 & & A \otimes B
 \end{array}
 \quad
 \begin{array}{ccc}
 A \otimes I & \xrightarrow{\gamma_{A,I}} & I \otimes A \\
 \downarrow \rho_A & & \downarrow \lambda_A \\
 A & \xlongequal{\quad} & A
 \end{array}$$

The following equality is also require to hold:

$$\lambda_I = \rho_I : I \otimes I \rightarrow I$$

[1]

**Definition 9.** A symmetric monoidal closed category (SMCC),  $(\mathbb{C}, \otimes, \multimap, I, \alpha, \lambda, \rho, \gamma)$ , is a SMC such that for all objects  $A$  in  $\mathbb{C}$ , the functor  $- \otimes A$  has a specified right adjoint  $A \multimap -$ .

**Definition 10.** A symmetric monoidal functor between SMCs  $(\mathbb{C}, \otimes, I, \alpha, \lambda, \rho, \gamma)$  and  $(\mathbb{C}', \otimes', I', \alpha', \lambda', \rho', \gamma')$  is a functor  $F : \mathbb{C} \rightarrow \mathbb{C}'$  equipped with

1. A morphism  $m_{I'} : I' \rightarrow FI$ .
2. For any two objects  $A$  and  $B$  in  $\mathbb{C}$ , a natural transformation  $m_{A,B} : F(A) \otimes' F(B) \rightarrow F(A \otimes B)$

These must satisfy the following diagrams:

$$\begin{array}{ccc}
 FI \otimes' FA & \xrightarrow{m_{I,A}} & F(I \otimes A) \\
 \uparrow m_I \otimes' id_{FA} & & \downarrow F(\lambda_A) \\
 I' \otimes' FA & \xrightarrow{\lambda'_{FA}} & FA \\
 \end{array}
 \quad
 \begin{array}{ccc}
 FA \otimes' FI & \xrightarrow{m_{A,I}} & F(A \otimes I) \\
 \uparrow id_{FA} \otimes' m_{I'} & & \downarrow F(\rho_A) \\
 FA \otimes' I' & \xrightarrow{\rho'_{FA}} & FA \\
 \end{array}$$
  

$$\begin{array}{ccc}
 (FA \otimes' FB) \otimes' FC & \xrightarrow{m_{A,B} \otimes' id_{FC}} & F(A \otimes B) \otimes' FC \xrightarrow{m_{A \otimes B, C}} F((A \otimes B) \otimes C) \\
 \uparrow \alpha'_{FA, FB, FC} & & \uparrow F(\alpha_{A, B, C}) \\
 FA \otimes' (FB \otimes' FC) & \xrightarrow{id_{FA} \otimes' m_{B, C}} & FA \otimes' F(B \otimes C) \xrightarrow{m_{A, B \otimes C}} F(A \otimes (B \otimes C)) \\
 \end{array}$$
  

$$\begin{array}{ccc}
 FA \otimes' FB & \xrightarrow{m_{A, B}} & F(A \otimes B) \\
 \downarrow \gamma'_{A, B} & & \downarrow F(\gamma_{A, B}) \\
 FB \otimes' FA & \xrightarrow{m_{B, A}} & F(B \otimes A) \\
 \end{array}$$

However in this particular case, assuming that  $!$  is a symmetric monoidal (endo) functor means that  $!$  comes equipped with a natural transformation

$$m_{A,B} : !A \otimes !B \rightarrow !(A \otimes B)$$

and a morphism

$$m_I : I \rightarrow !I$$

(where  $m_I$  is just the nullary version of the natural transformation.) The diagrams given in the above definition become the following:

$$\begin{array}{ccc}
!I \otimes !A & \xrightarrow{m_{I,A}} & !(I \otimes A) \\
\uparrow m_I \otimes id_{!A} & & \downarrow !(\lambda_A) \\
I \otimes !A & \xrightarrow{\lambda_{!A}} & !A
\end{array}
\quad
\begin{array}{ccc}
!A \otimes !I & \xrightarrow{m_{A,I}} & !(A \otimes I) \\
\uparrow id_{!A} \otimes m_I & & \downarrow !(\rho_A) \\
!A \otimes I & \xrightarrow{\rho_{!A}} & !A
\end{array}$$
  

$$\begin{array}{ccccc}
(!A \otimes !B) \otimes !C & \xrightarrow{m_{A,B} \otimes id_{!C}} & !(A \otimes B) \otimes !C & \xrightarrow{m_{A \otimes B, C}} & !((A \otimes B) \otimes C) \\
\uparrow \alpha_{!A, !B, !C} & & & & \uparrow !(\alpha_{A, B, C}) \\
!A \otimes (!B \otimes !C) & \xrightarrow{id_{!A} \otimes m_{B,C}} & !A \otimes !(B \otimes C) & \xrightarrow{m_{A, B \otimes C}} & !(A \otimes (B \otimes C))
\end{array}$$
  

$$\begin{array}{ccc}
!A \otimes !B & \xrightarrow{m_{A,B}} & !(A \otimes B) \\
\downarrow \gamma'_{A,B} & & \downarrow !(\gamma_{A,B}) \\
!B \otimes !A & \xrightarrow{m_{B,A}} & !(B \otimes A)
\end{array}$$

[1]

**Definition 11.** A symmetric monoidal functor,  $(F, m_{A,B}, m_{I'}) : \mathbb{C} \rightarrow \mathbb{C}'$ , is said to be

1. Strict if  $m_{A,B}$  and  $m_{I'}$  are identities.
2. Strong if  $m_{A,B}$  and  $m_{I'}$  are natural isomorphisms.

[1]

**Definition 12.** An oplax monoidal functor

$$(F, m) : (\mathbb{C}, \otimes, I) \rightarrow (\mathbb{D}, \otimes', I')$$

is symmetric when the following diagram commutes in the category  $\mathbb{D}$  for all objects  $A, B \in \mathbb{C}$

$$\begin{array}{ccc}
F(A \otimes B) & \xrightarrow{F\gamma} & F(B \otimes A) \\
\downarrow m & & \downarrow m \\
FA \otimes' FB & \xrightarrow{\gamma'} & FB \otimes' FA
\end{array}$$

[4]

### 3 Monoidal Double Category

**Definition 13.** The strict double category  $\mathbb{D}$  has a full double subcategory  $\mathbb{M}$  of monoidal categories. These are viewed as:

- vertical double categories on a formal object  $*$
- vertical arrows  $A : * \rightarrow *$
- cells  $a : A \rightarrow A'$

The horizontal arrows of  $\mathbb{M}$  are monoidal functors (lax with respect to tensor product) and the vertical arrows are comonoidal functors (colax). A cell  $\alpha : (FRSG)$  associates to every object  $A$  in  $\mathbb{A}$  with an arrow  $\alpha A : GRA \rightarrow SFA$  in  $\mathbb{D}$  which satisfies the naturality condition

$$(GRA|\alpha v) = (\alpha u|SFA) \quad (\text{for } a : (ufgv) \text{ in } \mathbb{A})$$

and the coherence conditions

$$\begin{aligned} (C\rho A|\alpha I|S\varphi A) &= (\gamma RA|I|\sigma FA) & (\text{for } a \text{ in } \mathbb{A}) \\ (G\rho(u,v)|\alpha v|S\varphi(u,v)) &= (\gamma(Ru,Rv)|(\alpha u \otimes \alpha v)|\sigma(Fu,Fv)) & (\text{for } w = u \otimes v \text{ in } \mathbb{A}) \end{aligned}$$

and have the following diagrams

$$\begin{array}{ccc} GRA & \xrightarrow{\alpha A} & SFA \\ \downarrow GRA & & \downarrow SFA \\ GRA' & \xrightarrow{\alpha A'} & SFA' \end{array} \quad \begin{array}{ccccc} & & GRI & \xrightarrow{\alpha I} & SFI \\ & \nearrow G\rho & & & \searrow S\varphi \\ GI & \xrightarrow{\gamma} & I & \xrightarrow{\sigma} & SI \end{array}$$

## 4 The Double Category of (co)Lax Monoidal Functors

**Definition 14.** A lax monoidal multifunctor from  $(\mathcal{C}_1, \dots, \mathcal{C}_n)$  to  $\mathcal{D}$  is a functor:

$$F : \mathcal{C}_1 \times \dots \times \mathcal{C}_n \longrightarrow \mathcal{D}$$

such that:

- Each functor,  $(F^i, m^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}), m_{I_i}^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}))$  defined by  $F^i(X) = F(\overrightarrow{A_1}, X, \overrightarrow{A_{i+1}}) : \mathcal{C}_i \longrightarrow \mathcal{D}$ , for  $1 \leq i \leq n$  is lax monoidal.
- The following equations hold for any  $1 \leq i < j \leq n$ :

$$\begin{aligned} - & m_{I_i}(\overrightarrow{I_1}, -, \overrightarrow{I_{i+1}}) = m_{I_j}(\overrightarrow{I_1}, -, \overrightarrow{I_{i+1}}), \\ - & (m_{I_i}^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, X, \overrightarrow{A_{j+1}}) \otimes m_{I_i}^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, Y, \overrightarrow{A_{j+1}})); m_{X,Y}^j(\overrightarrow{A_1}, I_i, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}}) \\ & = \iota; m_{I_i}^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, X \otimes Y, \overrightarrow{A_{j+1}}) \\ - & (m_{I_j}^j(\overrightarrow{A_1}, X, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}}) \otimes m_{I_j}^j(\overrightarrow{A_1}, Y, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}})); m_{X,Y}^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, I_j, \overrightarrow{A_{j+1}}) \\ & = \iota; m_{I_j}^j(\overrightarrow{A_1}, X \otimes Y, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}}) \end{aligned}$$

In the above definition  $\iota$  is the isomorphism  $\iota : I \otimes I \longrightarrow I$  which holds in any monoidal category.

**Definition 15.** A symmetric lax monoidal multifunctor from  $(\mathcal{C}_1, \dots, \mathcal{C}_n)$  to  $\mathcal{D}$  is a lax monoidal multifunctor

$$F : (\mathcal{C}_1, \dots, \mathcal{C}_n) \longrightarrow \mathcal{D}$$

such that:

- Each functor,  $(F^i, m^i(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}), m^i_{I_i}(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}))$  defined by  $F^i(X) = F(\overrightarrow{A_1}, X, \overrightarrow{A_{i+1}}) : \mathcal{C}_i \longrightarrow \mathcal{D}$ , for  $1 \leq i \leq n$  is symmetric lax monoidal.
- The following additional coherence axiom holds for any  $1 \leq i < j \leq n$ :

$$\begin{aligned}
& - (m^i_{X,Y}(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, P, \overrightarrow{A_{j+1}}) \otimes m^i_{X,Y}(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, Q, \overrightarrow{A_{j+1}})); m^j_{P,Q}(\overrightarrow{A_1}, X \otimes Y, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}}) \\
& = \tau; (m^j_{P,Q}(\overrightarrow{A_1}, X, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}}) \otimes m^j_{P,Q}(\overrightarrow{A_1}, Y, \overrightarrow{A_{i+1}}, -, \overrightarrow{A_{j+1}})); m^i_{X,Y}(\overrightarrow{A_1}, -, \overrightarrow{A_{i+1}}, P \otimes Q, \overrightarrow{A_{j+1}})
\end{aligned}$$

In the above definition  $\tau$  is the isomorphism  $\tau : (A \otimes B) \otimes (C \otimes D) \longrightarrow (A \otimes C) \otimes (B \otimes D)$  which holds in any monoidal category.

## References

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- [4] Paul-André Mellies. Categorical semantics of linear logic. *Panoramas et synthèses*, 27:15–215, 2009.