





### **Optimal Control**

## Introduction to Convex Optimization

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## **About This Course**

Grading System:

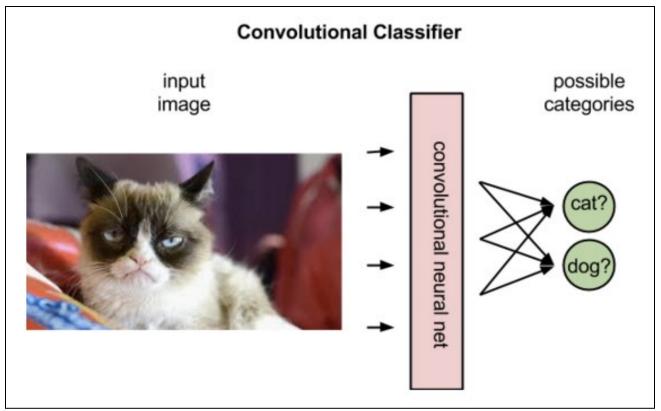
In-class Quizzes: 10%

Homework: 20%

Final Exam: 70%

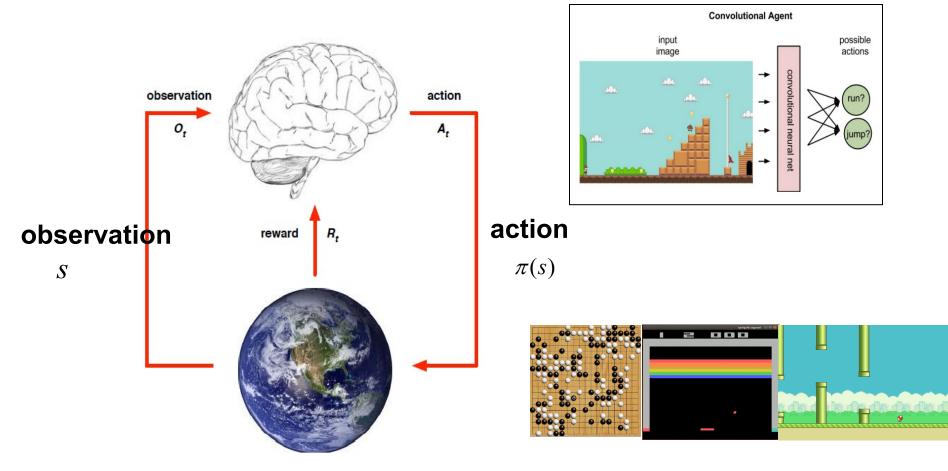
**Daily** performance

# Artificial intelligence (Deep learning)



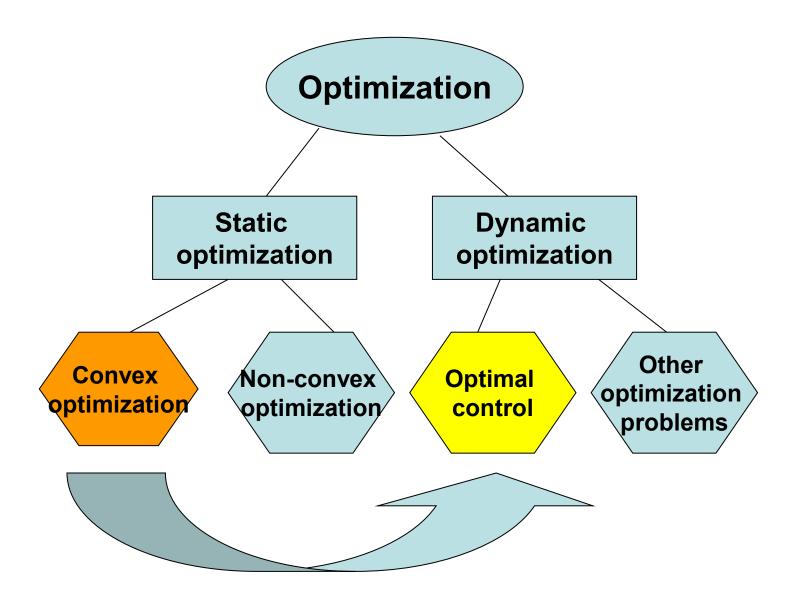
$$\mathbf{Min} \quad \sum_{i} \left\| f(x_i, \omega) - y_i \right\|$$

## Artificial intelligence (Reinforcement learning)



**Total reward:**  $V^{\pi}(s) = \mathbb{E}\left[R(s_0) + \gamma R(s_1) + \gamma^2 R(s_2) + \cdots \mid s_0 = s, \pi\right]$ . Maximization or minimization

## Background



## Content of this Course

- Convex optimization (Static optimization) (70%)
  - Covex set/ Convex function
  - Duality/Algorithms
- Calculus of variations
- Pontryagin maximum principle
- Bellman dynamic programming

(Dynamic optimization) (30%)

## Content

- Static optimization
  - A simple example
  - General mathematical formulation
- Least squares problem
  - An example
- Linear programming
  - An example
- Convex optimization

## Portfolio optimization

#### Portfolio optimization

- X: Capital

 $-X_i$ : Investment in the *i* assets

- Constraints:

$$x_1 + x_2 + \dots + x_n \le X$$

$$x_i \ge 0, i = 1, 2, 3, \dots, n$$

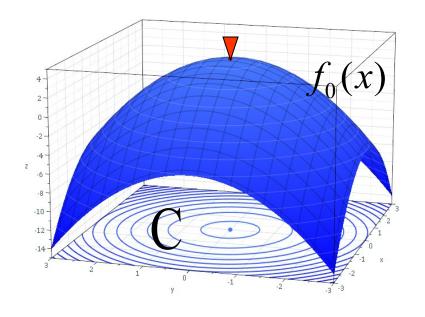


$$\sum_{i=1}^{n} f_i(x_i)$$



## Static optimization

Choosing the best element from some set of available alternatives



$$\max_{x} f_0(x)$$
s.t.  $x \in C \subset \mathbb{R}^n$ 

#### General mathematical formulation

#### **Optimization problem**

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \leq b_i, \quad i = 1, \dots, m$ 

- $x = (x_1, \ldots, x_n)$ : optimization variables
- $f_0: \mathbf{R}^n \to \mathbf{R}$ : objective function
- $f_i: \mathbf{R}^n \to \mathbf{R}, i = 1, \dots, m$ : constraint functions

Optimal Solution:  $x^*$  has smallest value of  $f_0$  among all vectors that satisfy the constraints

### Solving optimization problem

#### General optimization problem

- very difficult to solve
- methods involve some compromise, e.g.
  - very long computation time
  - not always finding the solution

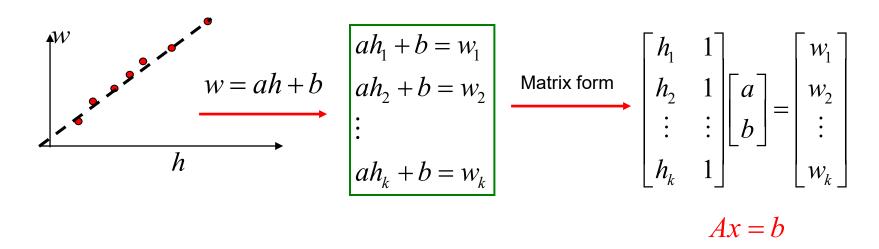
#### What problems are easy to solve?

- least-squares problems
- linear programming problems



### Least-squares problems (LS)

**Example:** weight & height (Assumption: Linear)



LS: minimize 
$$||Ax - b||_2^2$$

#### Solving least-squares problems

- analytical solution:  $x^* = (A^T A)^{-1} A^T b$
- reliable and efficient algorithms and software

### Linear programming problems (LP)

#### **Example**

- **diet problem:** choose quantities  $x_1, \ldots, x_n$  of n foods
  - one unit of food j costs  $c_i$ , contains amount  $a_{ij}$  of nutrient i
  - healthy diet requires nutrient i in quantity at least  $b_i$



LP:

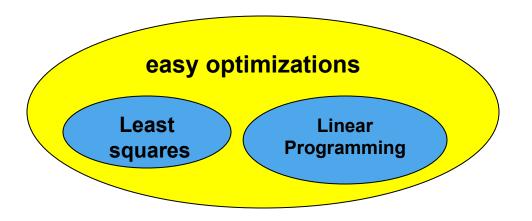
$$\begin{array}{ll} \text{minimize} & c^T x \\ \text{subject to} & Ax \succeq b, \quad x \succeq 0 \\ \end{array}$$

### Solving linear programming problem

- no analytical formula for solution
- reliable and efficient algorithms (simplex) and software



#### Extend LS and LP



#### **Linear function:**

#### 2-norm function:

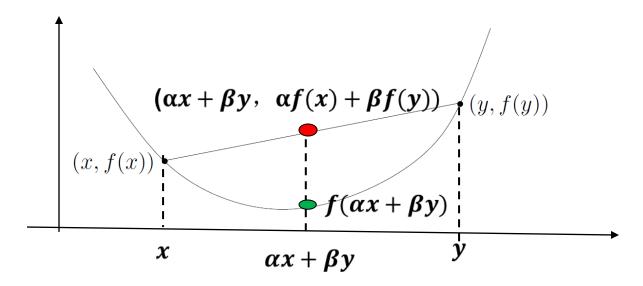
$$f(\alpha x + \beta y) = \alpha f(x) + \beta f(y) \qquad \|\alpha x + \beta y\|_{2} \le \alpha \|x\|_{2} + \beta \|y\|_{2}$$

$$\int \int \int \int \int dx dx + \beta y \le \alpha f(x) + \beta f(y)$$

#### **Convex function**

**Convex function**  $f(\alpha x + \beta y) \le \alpha f(x) + \beta f(y)$ 

if 
$$\alpha + \beta = 1$$
,  $\alpha \ge 0$ ,  $\beta \ge 0$ 



## **Convex optimization**

minimize 
$$f_0(x)$$
  
subject to  $f_i(x) \leq b_i, \quad i = 1, \dots, m$ 

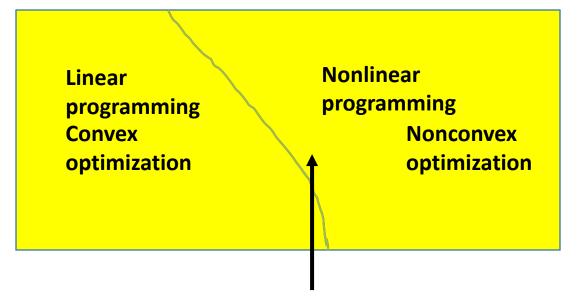
objective and constraint functions are convex

$$f_i(\alpha x + \beta y) \le \alpha f_i(x) + \beta f_i(y)$$

if 
$$\alpha + \beta = 1$$
,  $\alpha \ge 0$ ,  $\beta \ge 0$ 

## **Summary**

#### **Static optimization**

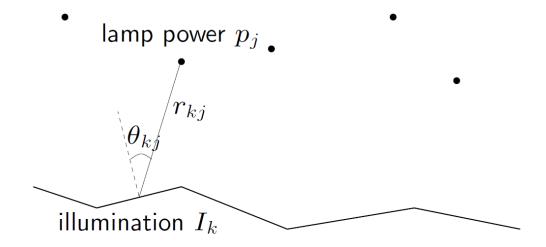




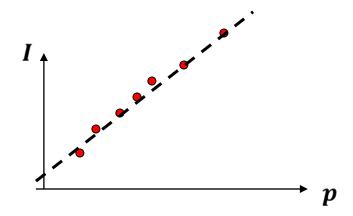
**Convex nonlinear programming** 

## An engineering example

m lamps illuminating n (small, flat) patches



#### Model



intensity  $I_k$  at patch k depends linearly on lamp powers  $p_j$ :

$$I_k = \sum_{j=1}^m a_{kj} p_j, \qquad a_{kj} = r_{kj}^{-2} \max\{\cos \theta_{kj}, 0\}$$

**problem**: achieve desired illumination  $I_{\text{des}}$  with bounded lamp powers

$$0 \leq p_j \leq p_{max}$$

#### **Optimization**

- 1. Use uniform power:  $p_j = p$ , vary p
- 2. Use least-squares:

#### Researching

minimize 
$$\sum_{k=1}^{n} (I_k - I_{\text{des}})^2$$

round 
$$p_j$$
 if  $p_j > p_{\text{max}}$  or  $p_j < 0$ 

3. Use weighted least-squares:

minimize 
$$\sum_{k=1}^{n} (I_k - I_{\text{des}})^2 + \sum_{j=1}^{m} w_j (p_j - p_{\text{max}}/2)^2$$

iteratively adjust weights  $w_j$  until  $0 \le p_j \le p_{\text{max}}$ 

easy to recognize

#### **Linear programming**

$$\begin{array}{ll} \text{minimize} & \max_{k=1,\ldots,n} |I_k - I_{\mathsf{des}}| \\ \text{subject to} & 0 \leq p_j \leq p_{\mathsf{max}}, \quad j = 1,\ldots,m \end{array}$$









$$\begin{cases}
\min \| x \|_1 \\
s \cdot t \cdot Ax = b
\end{cases}$$

$$\begin{cases} \min e^{T} x^{+} + e^{T} x^{-} \\ s \cdot t \cdot A x^{+} - A x^{-} = b \\ x^{+} \geqslant 0, x^{-} \geqslant 0 \end{cases}$$

$$x_{j}^{+} = \frac{1}{2}(|x_{j}| + x_{j}), x_{j}^{-} = \frac{1}{2}(|x_{j}| - x_{j})$$

minimize 
$$t$$
  
subject to  $\left|I_k - I_{des}\right| \le t$   
 $0 \le p_j \le p_{\max}, \ j = 1,...,m$ 

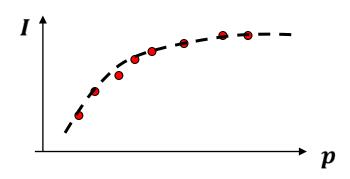
#### Pr oblem

min imize 
$$\sum_{k=1}^{n} |I_k - I_{des}|$$

s.t. 
$$0 \le p_j \le p_{\text{max}}, j = 1, 2, ..., m$$

#### Convex optimization

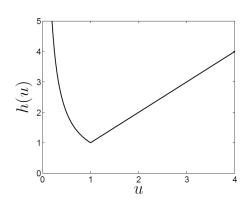




minimize 
$$\max_{k=1,...,n} |\log I_k - \log I_{\text{des}}|$$
  
subject to  $0 \le p_j \le p_{\text{max}}, \quad j = 1,...,m$ 

minimize 
$$f_0(p) = \max_{k=1,...,n} h(I_k/I_{\text{des}})$$
  
subject to  $0 \le p_j \le p_{\text{max}}, \quad j=1,...,m$ 

with 
$$h(u) = \max\{u, 1/u\}$$



 $f_0$  is convex because maximum of convex functions is convex

#### **Problems**

additional constraints: does adding 1 or 2 below complicate the problem?



- 1. no more than half of total power is in any 10 lamps
- 2. no more than half of the lamps are on  $(p_j > 0)$
- answer: with (1), still easy to solve; with (2), extremely difficult
- moral: (untrained) intuition doesn't always work; without the proper background very easy problems can appear quite similar to very difficult problems

$$I(p_1 > 0) + I(p_2 > 0) + \dots + I(p_m > 0) \le m/2$$

Research is going on!

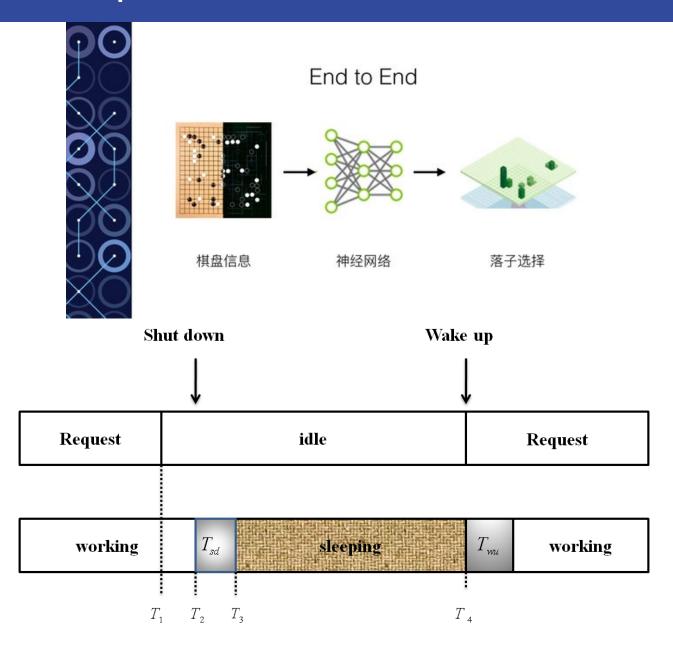
## Dynamic optimization

**\diamondsuit** Dynamic process x(t)

Max performance 
$$g_0(u) = \Phi_0(x(T)) + \int_0^T L_0(t, x(t), u(t)) dt$$
  
S.t. State evolution rule 
$$\begin{cases} \frac{dx}{dt} = f(t, x(t), u(t)), & t \in (0, T] \\ x(0) = x^0 \end{cases}$$

Find a sequence or functional to optimize the performance

## Two Examples



## Standard Optimal Control Problem

Minimize subject to

$$g_0(\mathbf{u}) = \Phi_0(\mathbf{x}(T)) + \int_0^T L_0(t, \mathbf{x}(t), \mathbf{u}(t)) dt$$

State equation 
$$\begin{cases} \frac{d\mathbf{x}}{dt} = \mathbf{f}(t, \mathbf{x}(t), \mathbf{u}(t)), & t \in (0, T] \\ \mathbf{x}(0) = \mathbf{x}^0 \end{cases}$$

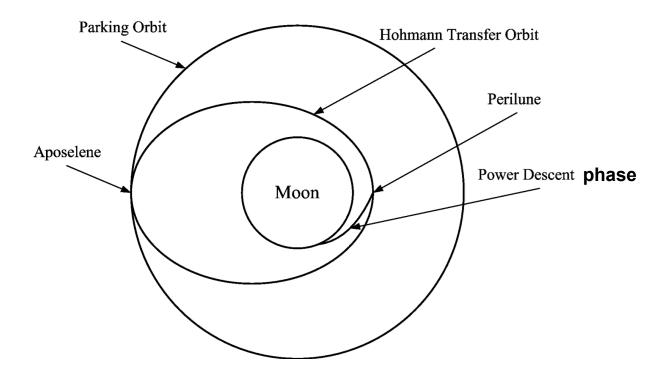
**Terminal constraint** 

$$\psi_i(\mathbf{x}(T)) \le 0$$
, and / or = 0,  $i = 1,..., N_T$   
 $h_i(t, \mathbf{x}(t)) \le 0$ ,  $t \in [0, T]$ ,  $i = 1,..., N$ 

Continuous constraint Control constraint

$$a_i \le u_i(t) \le b_i, \ \forall t \in [0,T], \ i = 1,...,r.$$
  
$$\boldsymbol{x} = [x_1,...,x_n]^T, \boldsymbol{u} = [u_1,...,u_r]^T$$

## Lunar Module Soft Landing

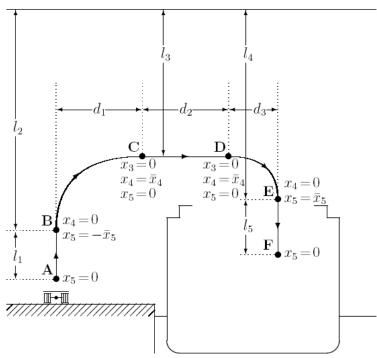


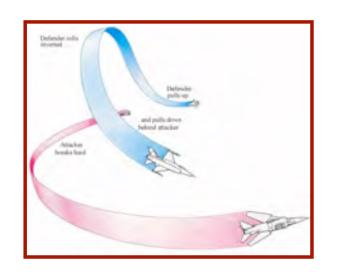
**Aposelene: 100km above the moon surface** 

Perilune: 15km above the moon surface

## Optimal Control of Container Cranes



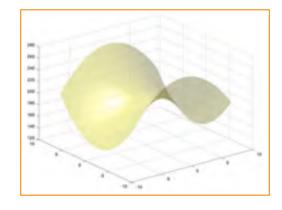




# Example: Pursuit-Evasion: A Competitive Optimization Problem

- Pursuer's goal: minimize final miss distance
- · Evader's goal: maximize final miss distance
- "Minimax" (saddle-point) cost function
- Optimal control laws for pursuer and evader

$$\mathbf{u}(t) = \begin{bmatrix} \mathbf{u}_{P}(t) \\ \mathbf{u}_{E}(t) \end{bmatrix} = - \begin{bmatrix} \mathbf{C}_{P}(t) & \mathbf{C}_{PE}(t) \\ \mathbf{C}_{EP}(t) & \mathbf{C}_{E}(t) \end{bmatrix} \begin{bmatrix} \hat{\mathbf{x}}_{P}(t) \\ \hat{\mathbf{x}}_{E}(t) \end{bmatrix}$$

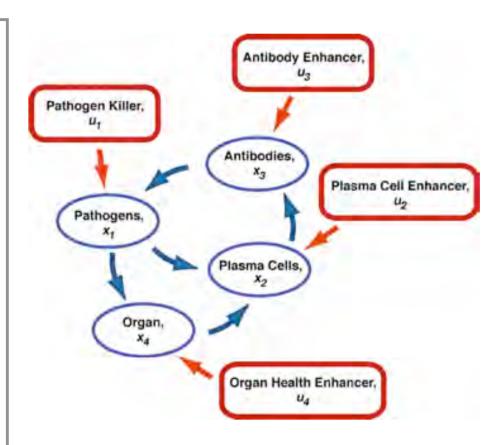


Example of a differential game, Isaacs (1965), Bryson & Ho (1969)



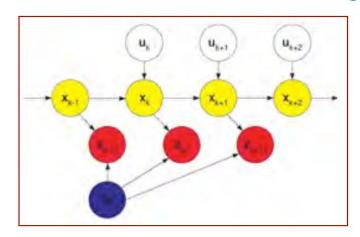
## **Example:** Minimize Concentrations of Virus, Infected Cells, and Drug Usage

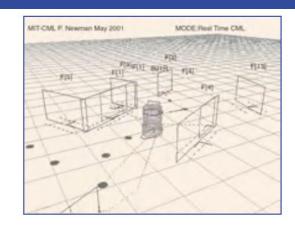
- x<sub>1</sub> = Concentration of a pathogen, which displays antigen
- x<sub>2</sub> = Concentration of plasma cells, which are carriers and producers of antibodies
- x<sub>3</sub> = Concentration of antibodies, which recognize antigen and kill pathogen
- x<sub>4</sub> = Relative characteristic of a damaged organ [0 = healthy, 1 = dead]

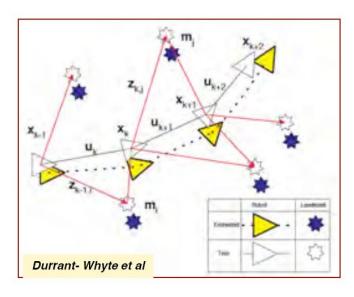


## Example: Simultaneous Location and Mapping (SLAM)

- Build or update a local map within an unknown environment
  - Stochastic map, defined by mean and covariance
  - SLAM Algorithm = State estimation with extended Kalman filter
  - Landmark and terrain tracking







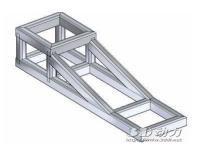
#### Conclusion

#### Static

- Optimal state, x\*, and control, u\*, are fixed,
   i.e., they do not change over time
  - $J^* = J(x^*, u^*)$
  - Functional minimization (or maximization)
  - Parameter optimization

#### Dynamic

- Optimal state and control vary over time
  - $J^* = J[x^*(t), u^*(t)]$
  - Optimal trajectory
  - Optimal feedback strategy
- Optimized cost function, J\*, is a scalar, real number in both cases



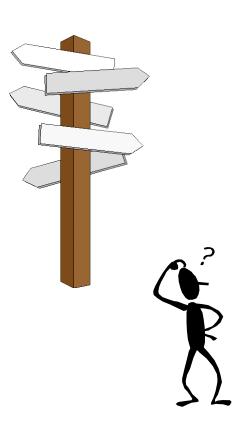




## What is optimization

- Choice of best parameter
- Choice of best strategy
- Choice of best control
- Choice of best estimate

Choosing the best element from some set of available alternatives



## Thank You