## Harvard University Extension School Computer Science E-121

## Problem Set 5

Due October 30, 2015 at 11:59 PM.

Submit your solutions electronically on the course website, located at https://canvas.harvard.edu/courses/4896/assignments. On the site, use the assignments tab to find the correct problem set, then with the "submit assignment" button, upload the PDF file of your solution.

## LATE PROBLEM SETS WILL NOT BE ACCEPTED.

See the syllabus for the collaboration policy.

PROBLEM 1 (4+2 points, suggested length of 1/2 page)

Consider the Turing machine  $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$  where

- $Q = \{q_0, q_1, q_2, q_{\text{accept}}, q_{\text{reject}}\},$
- $\Sigma = \{a, b\}$  and  $\Gamma = \{a, b, \sqcup\},\$
- The start, accept and reject states are  $q_0$ ,  $q_{\text{accept}}$ , and  $q_{\text{reject}}$  respectively.
- The function  $\delta$  is given by:

q	$\sigma$	$\delta(q,\sigma)$
$\overline{q_0}$	a	$(q_0, b, R)$
$q_0$	b	$(q_1, a, R)$
$q_0$	$\sqcup$	$(q_{\text{reject}}, \sqcup, R)$
$q_1$	a	$(q_2, b, R)$
$q_1$	b	$(q_1, a, R)$
$q_1$	$\sqcup$	$(q_{\text{reject}}, \sqcup, R)$
$q_2$	a	$(q_2, b, R)$
$q_2$	b	$(q_2, a, R)$
$q_2$		$(q_{\text{accept}}, \sqcup, R)$

- (A) Give the sequence of configurations describing M's computation on the string aabba
- (B) Describe L(M).

## PROBLEM 2 (4 points, suggested length of 1/4 page)

A Modern Turing Machine (MTM), instead of  $\{L, R\}$ , has  $\{L_k, R_k\}$  (move the head left or right k spaces on the tape, for any integer k encoded in the rule in the MTM). Prove that the MTMs are equivalent in power to the TMs. In other words, any MTM can be converted to a TM with an equivalent language, and vice versa.

PROBLEM 3 
$$(2 + 2 + 1)$$
 points, suggested length of  $1/2$  page)

A Democratic rule of a CFG  $(V, \Sigma, R, S)$  is one whose right-hand side is a member of  $V\Sigma^* \cup \Sigma^*$ , that is, only the leftmost symbol can be a nonterminal. Republican rules are defined analogously where only the rightmost symbol can be a nonterminal. Show that a grammar with only Democratic rules, or only Republican rules, generates a regular language, but a grammar with a mixture of Democratic and Republican rules may generate a non-regular language.

Let us define EvenPalindrome(L) =  $\{ww^R : w \in L\}$ . Show that if L is regular, then EvenPalindome(L) is context-free. Note that this is not the same as the language of all even palindromes, which is actually just EvenPalindrome( $\Sigma^*$ ). Explain briefly how your solution works. You need not provide a formal proof of correctness.

A Boring Turing Machine (BTM) can only write # on the tape (assume  $\# \notin \Sigma$ ). Prove that the BTMs are equivalent in power to the TMs.

Show that every infinite Turing-recognizable language has an infinite decidable sublanguage.

Let  $G = (V, \Sigma, R, S)$  where  $V = \{S, V\}, \Sigma = \{a, b\}$ , and R is the set of rules:

$$S \rightarrow bSS \mid aS \mid aV$$
 
$$V \rightarrow aVb \mid bVa \mid VV \mid \varepsilon$$

- (A) Transform G into an equivalent grammar G' in Chomsky normal form. Show your work for each step of this conversion, but long justifications for each step are not necessary.
- (B) Explain in English what language the grammar G generates. (One clearly worded sentence is fine; thorough justification is not necessary, but it might get you some partial credit if your answer is wrong.)
- (C) Check whether the strings abaab and bbabaa are generated by G, using the recognition algorithm for grammars in Chomsky normal form given in class. Show the complete filled-in matrix for each of the two strings.