

# Question 1

Consider the three transactions T1, T2 and T3, and the schedules S1, S2, S3 and S4 given below. Which of the schedules is conflict serializable? The subscript for each database operation in a schedule denotes the transaction number for that operation. For each schedule, show all conflicts, draw the precedence graph, determine and write down if it is serializable or not, and the equivalent serial schedules if exist

## Transactions

T1: R1(X), W1(X)

T2: R2(X)

T3: R3(X) W3(X)

## Schedules

S1: R1(X) R3(X) W1(X) R2(X) W3(X)

S2: R1(X) R3(X) W3(X) W1(X) R2(X)

S3: R3(X) R2(X) W3(X) R1(X) W1(X)

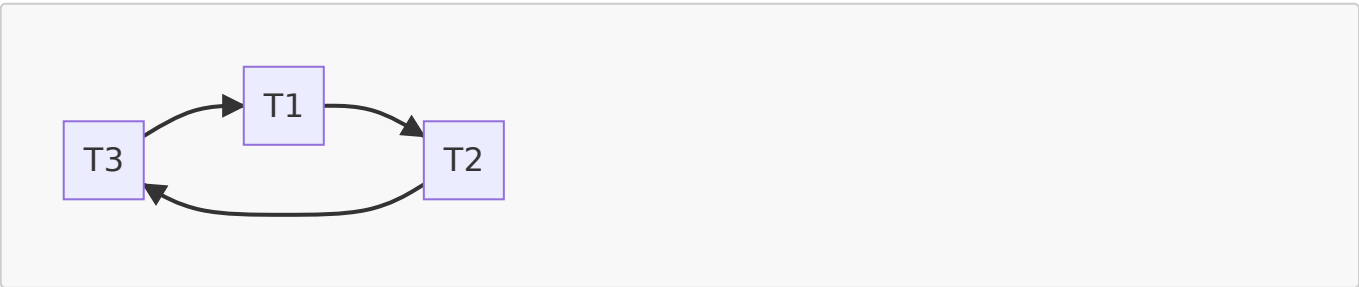
S4: R3(X) R2(X) R1(X) W3(X) W1(X)

## Schedule 1

S1: R1(X) R3(X) W1(X) R2(X) W3(X)

T1	T2	T3
R1(X)		
		R3(X)
W1(X)		
	R2(X)	
		W3(X)

## Schedule 1 Graph



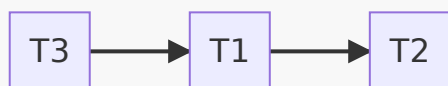
S1 is not conflict serializable, cycle detected in graph

## Schedule 2

S2: R1(X) R3(X) W3(X) W1(X) R2(X)

T1	T2	T3
R1(X)		
		R3(X)
		W3(X)
W1(X)		
	R2(X)	

### Schedule 2 Graph



Schedule 2 is Conflict Serializable in the type T3->T1->T2

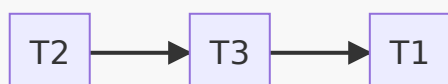
Schedule 2 Conflict Serialized: R3(X) W3(X) R1(X) W1(X) R2(X)

## Schedule 3

S3: R3(X) R2(X) W3(X) R1(X) W1(X)

T1	T2	T3
		R3(X)
	R2(X)	
		W3(X)
R1(X)		
W1(X)		

### Schedule 3 Graph



Schedule 3 is Conflict Serializable in the Type  $T2 \rightarrow T3 \rightarrow T1$

Schedule 3 Conflict Serialized:  $R2(X) R3(X) W3(X) R1(X) W1(X)$

## Schedule 4

S4:  $R3(X) R2(X) R1(X) W3(X) W1(X)$

T1	T2	T3
		R3(X)
	R2(X)	
R1(X)		
		W3(X)
W1(X)		

Schedule 4 Graph



Schedule 4 is NOT Conflict Serializable, Cycle detected on graph between T1 and T3