

Ternary Search

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Computer systems use different methods to find specific data. There are various search algorithms, each better suited for certain situations. For instance, a binary search divides information into two parts, while a ternary search does the same but into three equal parts. It's worth noting that ternary search is only effective for sorted data. In this article, we're going to uncover the secrets of **Ternary Search** – how it works, why it's faster in some situations.



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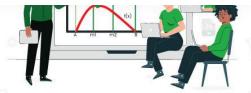


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What is the Ternary Search?

Ternary search is a search algorithm that is used to find the position of a target value within a sorted array. It operates on the principle of dividing the array into **three parts** instead of two, as in <u>binary search</u>. The basic idea is to narrow down the search space by comparing the target value with elements at two points that divide the array into **three equal parts**.

- mid1 = l + (r-l)/3
- mid2 = r (r-1)/3

When to use Ternary Search:

- When you have a large ordered array or list and need to find the position of a specific value.
- When you need to find the maximum or minimum value of a function.
- When you need to find bitonic point in a bitonic sequence.
- When you have to evaluate a quadratic expression

Working of Ternary Search:

The concept involves dividing the array into **three equal segments** and determining in which segment the **key** element is located. It works similarly to a binary search, with the distinction of reducing time complexity by dividing the array into three parts instead of two.

Below are the step-by-step explanation of working of Ternary Search:

1. Initialization:

• Set two pointers, **left** and **right**, initially pointing to the first and last elements of our search space.

2. Divide the search space:

- Calculate two midpoints, **mid1** and **mid2**, dividing the current search space into three roughly equal parts:
- mid1 = left + (right left) / 3
- mid2 = right (right left) / 3
- The array is now effectively divided into [left, mid1], (mid1, mid2), and [mid2, right].

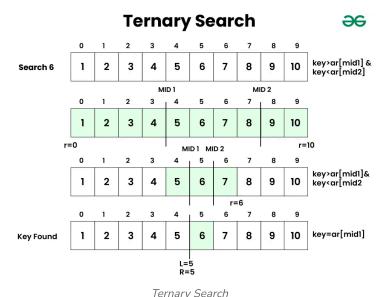
3. Comparison with Target:.

- If the **target** is equal to the element at **mid1** or **mid2**, the search is successful, and the index is returned
- If the target is less than the element at mid1, update the right pointer to
 mid1 1.
- If the **target** is greater than the element at **mid2**, update the **left** pointer to **mid2 + 1**.
- If the **target** is between the elements at **mid1** and **mid2**, update the **left** pointer to **mid1 + 1** and the **right** pointer to **mid2 1**.

4. Repeat or Conclude:

- Repeat the process with the reduced search space until the **target** is found or the search space becomes empty.
- If the search space is empty and the **target** is not found, return a value indicating that the **target** is not present in the array.

Illustration:



Implementation:

Below is the implementation of Ternary Search Approach:



```
import math as mt
4
5 # Function to perform Ternary Search
6 def ternarySearch(1, r, key, ar):
7
        if (r >= 1):
8
9
            # Find the mid1 and mid2
10
            mid1 = 1 + (r - 1) //3
11
            mid2 = r - (r - 1) //3
12
13
            # Check if key is present at any mid
14
            if (ar[mid1] == key):
15
                return mid1
16
17
            if (ar[mid2] == key):
18
                return mid2
19
20
            # Since key is not present at mid,
21
22
            # check in which region it is present
23
            # then repeat the Search operation
            # in that region
24
            if (key < ar[mid1]):</pre>
25
26
                # The key lies in between 1 and mid1
27
                return ternarySearch(l, mid1 - 1, key, ar)
28
29
            elif (key > ar[mid2]):
30
31
                # The key lies in between mid2 and r
32
                return ternarySearch(mid2 + 1, r, key, ar)
33
34
            else:
35
36
                # The key lies in between mid1 and mid2
37
38
                return ternarySearch(mid1 + 1,
39
                                      mid2 - 1, key, ar)
40
        # Key not found
41
        return -1
42
43
   # Driver code
44
   1, r, p = 0, 9, 5
```

```
47 # Get the array
48 # Sort the array if not sorted
   ar = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
49
50
   # Starting index
51
   1 = 0
52
53
  # end element index
54
55 r = 9
56
   # Checking for 5
57
58
   # Key to be searched in the array
59
   key = 5
60
61
   # Search the key using ternarySearch
62
   p = ternarySearch(1, r, key, ar)
63
64
   # Print the result
   print("Index of", key, "is", p)
66
67
   # Checking for 50
68
69
   # Key to be searched in the array
70
   key = 50
71
72
   # Search the key using ternarySearch
   p = ternarySearch(1, r, key, ar)
74
75
   # Print the result
76
  print("Index of", key, "is", p)
77
79 # This code is contributed by
80 # Mohit kumar 29
```

Output

```
Index of 5 is 4
Index of 50 is -1
```

Time Complexity: O(2 * log₃n)

Auxiliary Space: O(log₃n)

Iterative Approach of Ternary Search:

C Java Python3 C# JavaScript PHP C++

```
Q
      1 # Python 3 program to illustrate iterative
        # approach to ternary search
      3
        # Function to perform Ternary Search
      4
         def ternarySearch(1, r, key, ar):
      5
             while r >= 1:
      6
      7
                  # Find mid1 and mid2
      8
                 mid1 = 1 + (r-1) // 3
      9
                 mid2 = r - (r-1) // 3
     10
     11
                  # Check if key is at any mid
     12
                  if key == ar[mid1]:
     13
                      return mid1
     14
                  if key == ar[mid2]:
     15
                      return mid2
     16
     17
                  # Since key is not present at mid,
     18
                  # Check in which region it is present
     19
                 # Then repeat the search operation in that region
     20
                  if key < ar[mid1]:</pre>
     21
     22
                      # key lies between l and mid1
                      r = mid1 - 1
     23
     24
                  elif key > ar[mid2]:
                      # key lies between mid2 and r
     25
                      1 = mid2 + 1
     26
                  else:
     27
                      # key lies between mid1 and mid2
     28
                      1 = mid1 + 1
     29
                      r = mid2 - 1
     30
     31
             # key not found
     32
             return -1
     33
     34
         # Driver code
     35
```

```
37 # Get the list
38 # Sort the list if not sorted
39 ar = [1, 2, 3, 4, 5, 6, 7, 8, 9, 10]
40
  # Starting index
41
42 1 = 0
43
44 # end element index
  r = 9
45
46
47 # Checking for 5
48 # Key to be searched in the list
  key = 5
49
50
  # Search the key using ternary search
51
   p = ternarySearch(1, r, key, ar)
52
53
54 # Print the result
55 print("Index of", key, "is", p)
56
57 # Checking for 50
58 # Key to be searched in the list
59 key = 50
60
  # Search the key using ternary search
61
  p = ternarySearch(1, r, key, ar)
62
63
  # Print the result
64
  print("Index of", key, "is", p)
65
66
67 # This code has been contributed by Sujal Motagi
```

Output

```
Index of 5 is 4
Index of 50 is -1
```

Time Complexity: $O(2 * log_3 n)$, where n is the size of the array. **Auxiliary Space:** O(1)

Complexity Analysis of Ternary Search:

Time Complexity:

• Worst case: O(log₃N)

• Average case: $\Theta(\log_3 N)$

• Best case: $\Omega(1)$

Auxiliary Space: O(1)

Binary search Vs Ternary Search:

The time complexity of the binary search is less than the ternary search as the number of comparisons in ternary search is much more than binary search. Binary Search is used to find the maxima/minima of <u>monotonic functions</u> where as Ternary Search is used to find the maxima/minima of <u>unimodal functions</u>.

Note: We can also use ternary search for monotonic functions but the time complexity will be slightly higher as compared to binary search.

Advantages:

- Ternary search can find maxima/minima for unimodal functions, where binary search is not applicable.
- Ternary Search has a time complexity of $O(2 * log_3 n)$, which is more efficient than linear search and comparable to binary search.
- Fits well with optimization problems.

Disadvantages:

- Ternary Search is only applicable to ordered lists or arrays, and cannot be used on unordered or non-linear data sets.
- Ternary Search takes more time to find maxima/minima of monotonic functions as compared to Binary Search.

Summary:

- Ternary Search is a divide-and-conquer algorithm that is used to find the position of a specific value in a given array or list.
- It works by dividing the array into three parts and recursively performing the search operation on the appropriate part until the desired element is found.

- The algorithm has a time complexity of $O(2 * log_3 n)$ and is more efficient than a linear search, but less commonly used than other search algorithms like binary search.
- It's important to note that the array to be searched must be sorted for Ternary Search to work correctly.

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