# Just Enough Category Theory for Haskell (part 1)

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#### Tabel of Contents

- Purpose and Focus
- Category
- Functor

## Purpose and Focus

#### Purpose

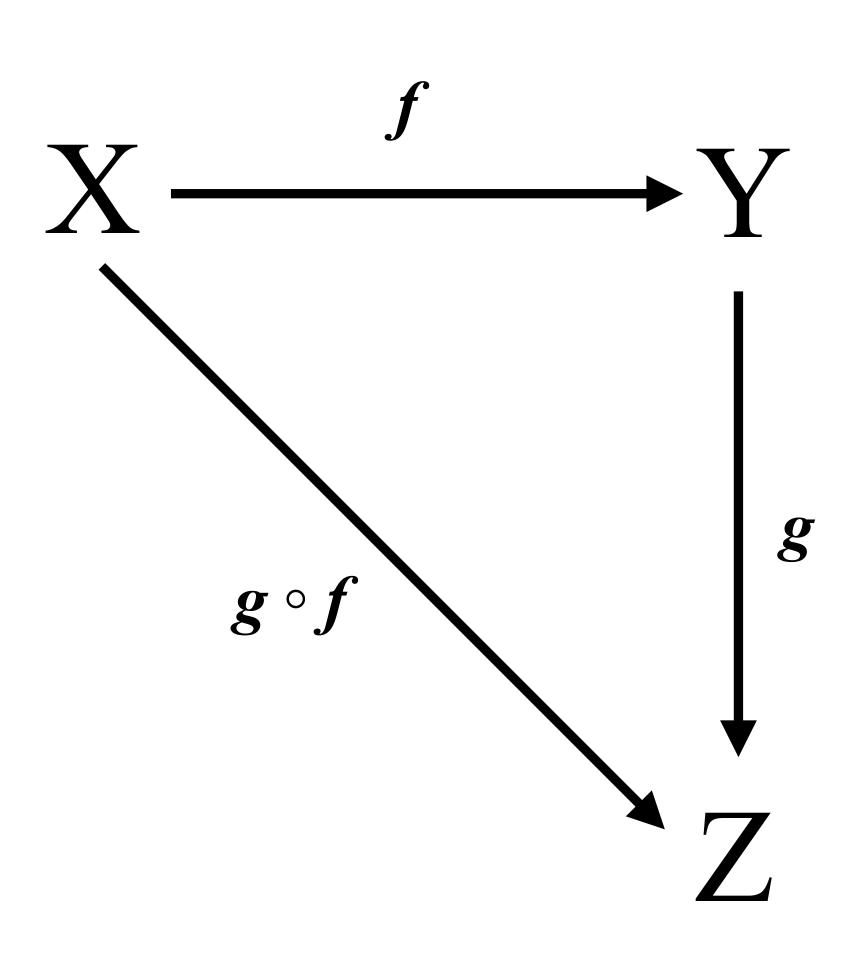
- "Do I have to read this thick math book with stubborn, dull cover thoroughly, just to write a program in Haskell?"
- You don't.
- We've already gone through 5 weeks of exercises without any book!
- Yet, as Haskell borrowed a lot of its core concepts from category theory, some knowledge on it might help you a lot.

#### Focus

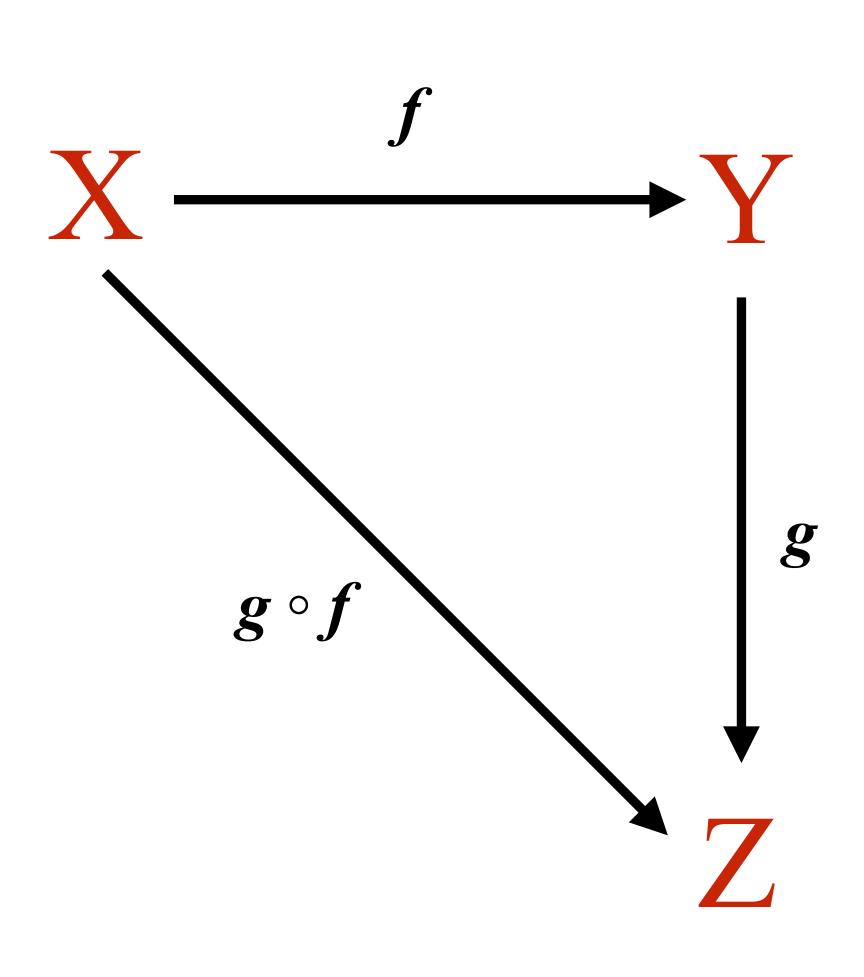
- · I'll try to avoid analogues, which is often used to give an intuition.
- That's because sometimes those keep you from getting actual, clear idea about a concept.
- Instead, I'd rather stick to mathematical terms and laws.

#### Disclaimer

- · I've never been excellent in math.
- Haven't been really passionate about it, neither.
- Though I've put a huge amount of time, there will be lots of unclear, even wrong spots throughout the slides.
- If you find any, please feel free to interrupt and correct me @

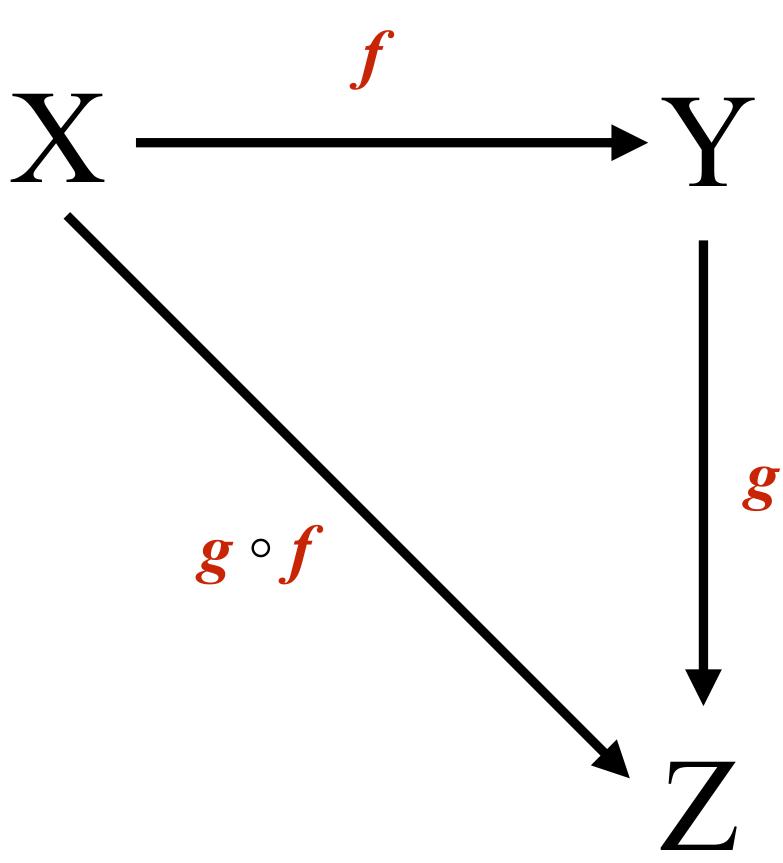


- A collection of objects
- A collection of morphisms
- A notion of composition

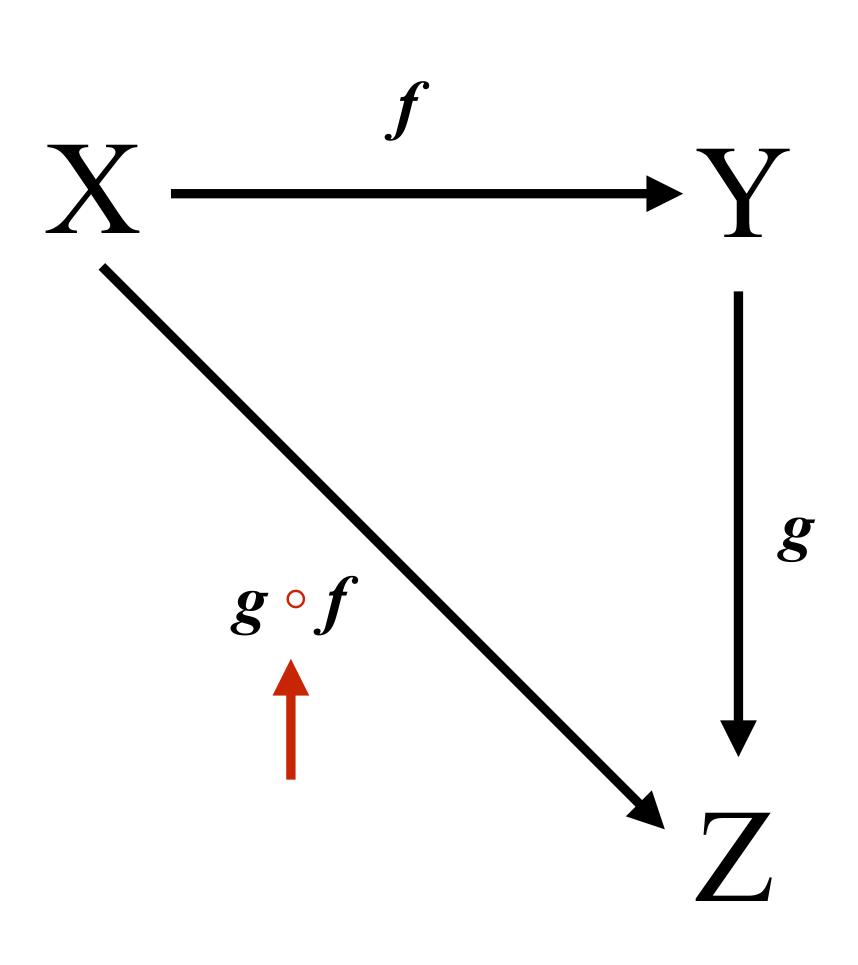


- A collection of objects
- A collection of morphisms
- A notion of composition

source object X target object Y



- A collection of objects
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- A collection of objects
- A collection of morphisms
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#### Category Example

- · Category of sets(Set) is a category.
- Set has a collection of objects, which is, well, all sets.
- **Set** has a collection of morphisms, which is a collection of functions from a set to another set.
- Set has a notion of composition, which is just a function composition.

## Category Example (cont'd)

- · Category of groups(Grp) is also a category.
- Grp has a collection of objects, which is, again, all groups.
- **Grp** has a collection of morphisms, which is a collection of functions which preserves group operation (the group homomorphism)
- Grp has a notion of composition, which is just a function composition.

#### Category Laws

Categories must meet these requirements:

- Composition of mophisms should be associative  $f \circ (g \circ h) = (f \circ g) \circ h$
- For every object in a category, there must be an identity morphism

#### Hask

- · Hask, the category of Haskell types and functions, is our main topic.
- Speaking of category, we've heard that every category must have...
- · Objects, morphisms, and composition!

## Hask (cont'd)

- Objects: Haskell types
- Morphisms: Haskell functions
- · Composition : (.)
- We know haskell functions are associative, and we have id.
- On top of that, we *define* two functions f and g as the same morphism if f x = g x for all x. (Why? Is it really needed?)

#### Functor

#### Functor

- "In mathematics, a functor is a type of mapping between categories which is applied in category theory. Functors can be thought of as homomorphisms between categories." (Wikipedia)
- Sorry, what?

## Functor (cont'd)

"In mathematics, a functor is a type of **mapping between categories** which is applied in category theory.

Functors can be thought of as **homomorphisms between categories**."

- To map a category to another category, we have to map...
- · Objects, Morphisms, and Composition!
- Composition is mostly trivial for this case, so we won't focus on it.

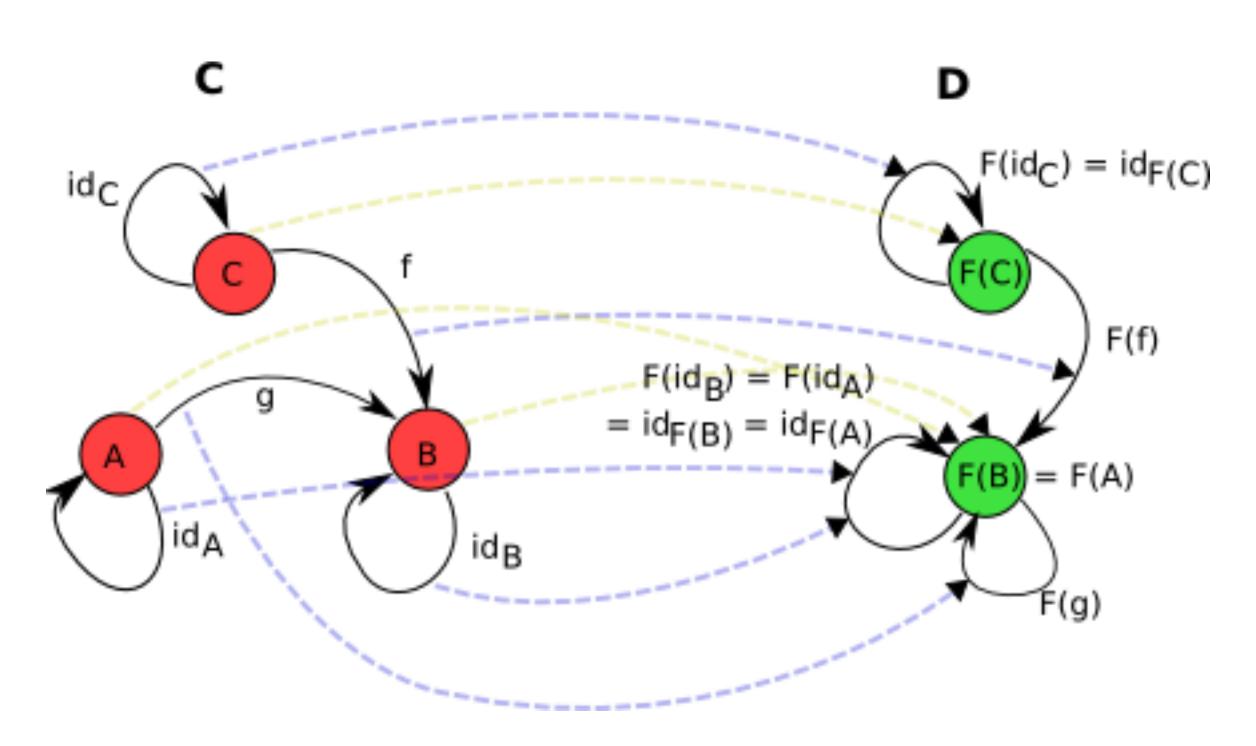
#### Functor Laws

Functors must meet these requirements:

$$F(id_A)=id_{F(A)}$$

- Functor must map identity of an object to that of a mapped object
- Functor must distribute over morphism composition  $F(f \circ g) = F(f) \circ F(g)$

#### Functor Example



- Diagram of a functor from category C to category D.
- You can see:
  - Object mapping (yellow)
  - Morphism mapping (purple)

#### Functor in Hask

```
class Functor (f :: * -> *) where
  fmap :: (a -> b) -> f a -> f b

instance Functor Maybe where
  fmap f (Just x) = Just (f x)
  fmap _ Nothing = Nothing
```

- What does Maybe do?
  - It turns a **Hask** object(type), T, to another object, Maybe T.

## Functor in Hask (cont'd)

```
class Functor (f :: * -> *) where
  fmap :: (a -> b) -> f a -> f b

instance Functor Maybe where
  fmap f (Just x) = Just (f x)
  fmap _ Nothing = Nothing
```

- What does fmap do, when applied to a Maybe instance?
  - It turns a Hask morphism(function), with type (A -> B),
     to another morphism, with type (Maybe A -> Maybe B).

## Functor in Hask (cont'd)

- So as we have
  - Type constructor, which maps objects between categories
  - fmap, which maps morphisms between categories
- Maybe is an instance of Functor it maps Hask to its subcategory, a category of Maybe types and functions defined on Maybe types.

## Functor in Hask (cont'd)

- Same logic applies to any Functor types.
- · Both Functor laws are fulfilled for any functor on Hask.
  - fmap id = id
  - fmap(f.g) = (fmap f).(fmap g)

# Summary

#### Summary

- We've talked about Hask, the category of Haskell types and functions.
- (types, functions, (.) ) corresponds to (objects, morphisms, composition), respectively.
- Functor typeclasses works as Functor in category theory, that is, a mapping between two different categories.
- Next topics: Monoid, Applicative Functor, Monad, ...

## Appendix

#### References

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