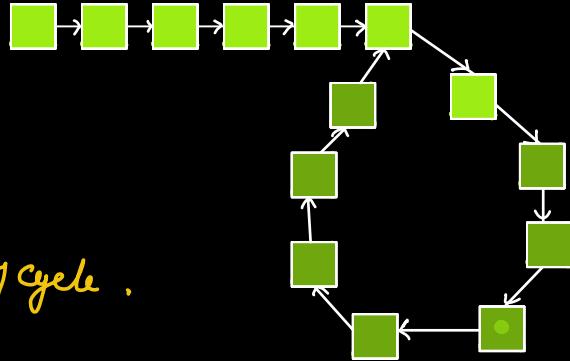


Cycle detection

1. Detect if there is a cycle.
2. Find the first node /
start node of cycle .



Approach 1

Use Set / Map

HashSet < List Node > set ;

- Iterate over the LL
for every node :

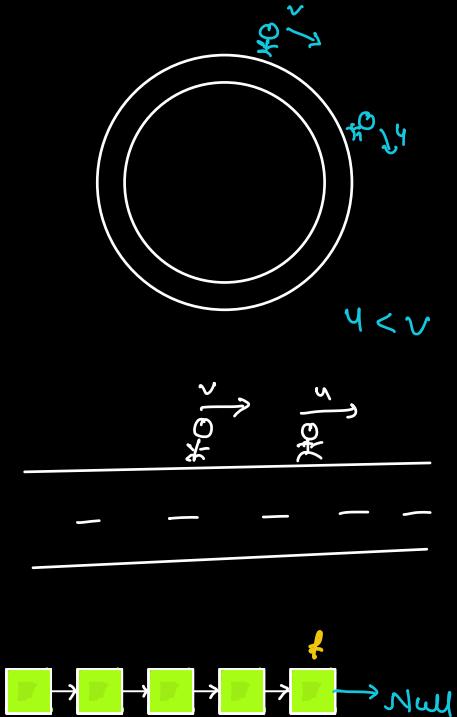
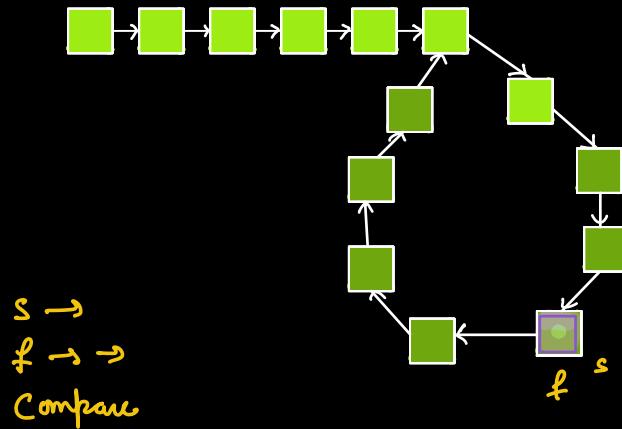
Check if it present in set

if present → Cycle is present
else add in set start of it

TC : O(N)

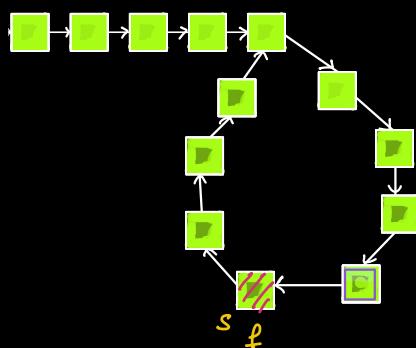
SC : O(N)

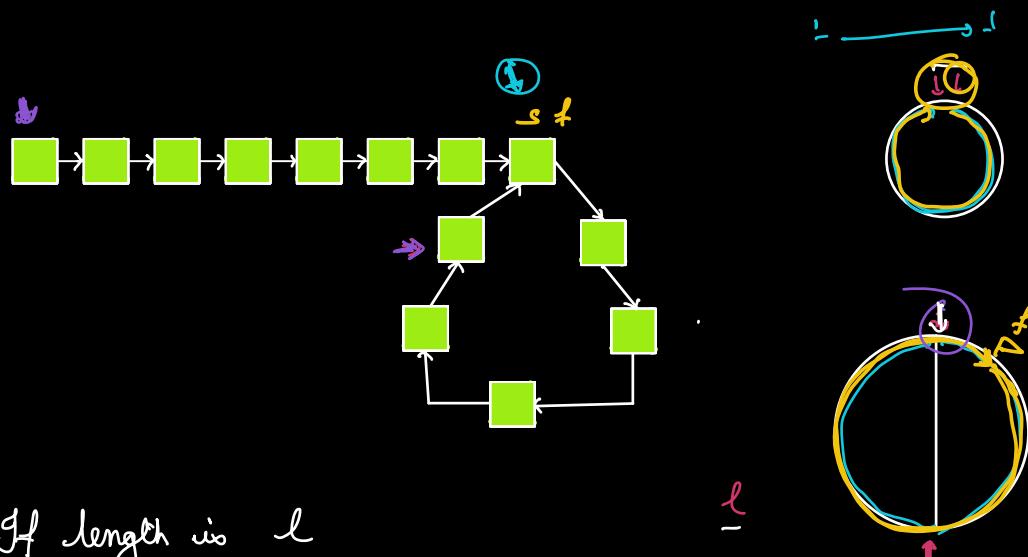
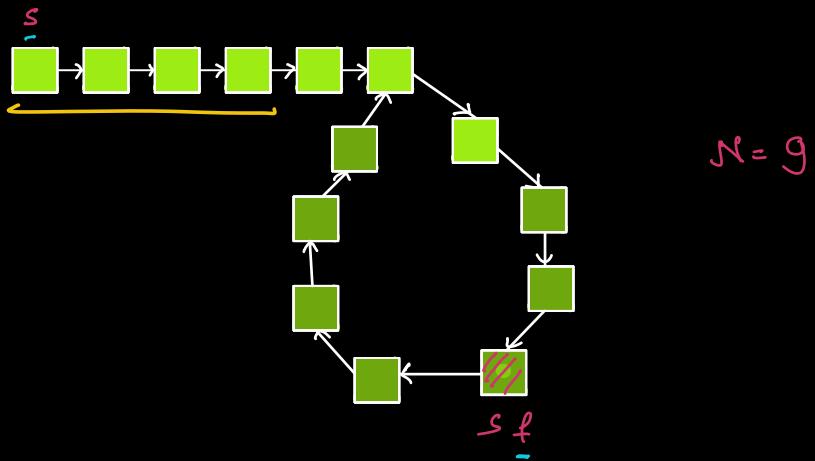
fast & slow pointers



```

if (head == null) ret false;
while (fast != null && fast.next != null) {
  slow = slow.next,
  fast = fast.next.next,
  if (slow == fast)
    ret true,
}
  
```





If length is l

& we start from the start of cycle

l iterations \rightarrow starty pool

$2l$ iterations \rightarrow starty pool

$3l$ iterations \rightarrow starty pool

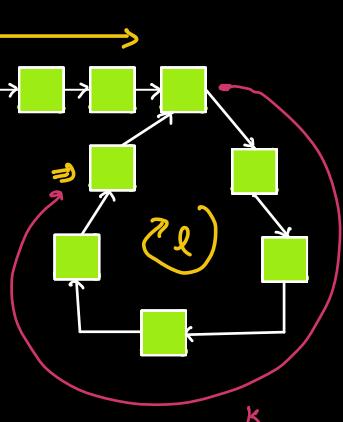
⋮

Nl \rightarrow start



$$\text{dist(fast)} = M + xl + K$$

$$\text{dist(slow)} = M + yl + K$$



$$\text{dist(fast)} = 2 \text{ dist(slow)}$$

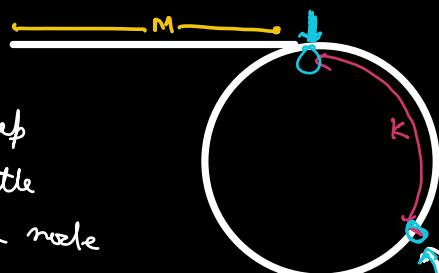
$$M + xl + K = 2M + 2yl + 2K$$

$$xl - 2yl = M + K$$

$$(x - 2y)l = M + K$$

$(M+k)$ \Rightarrow integral multiple of l

If we start iteratively from the 1st node of cycle $\xrightarrow{(M+k)}$ end up at the same node

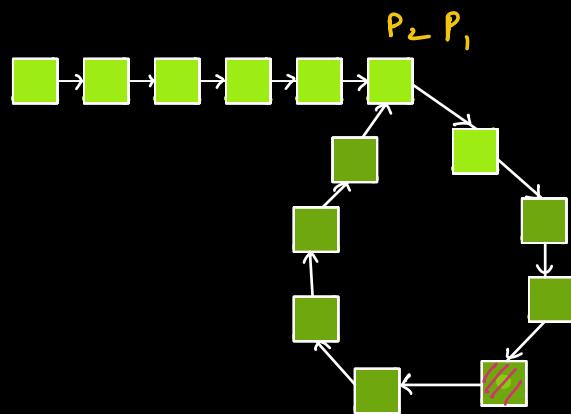
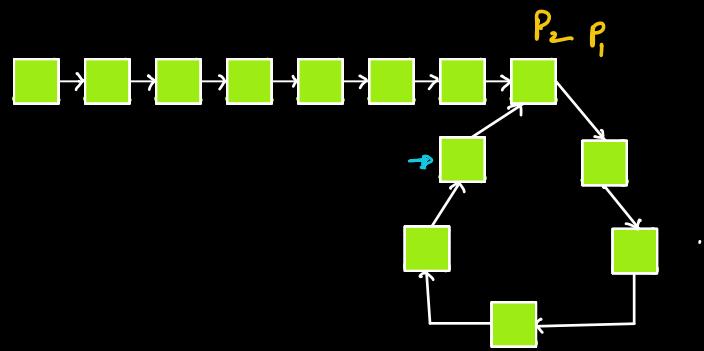


The meeting point is at K dist from the start.

So M iterations from the meeting point \longrightarrow start node

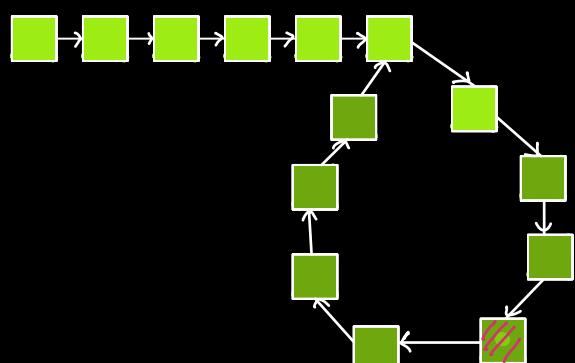
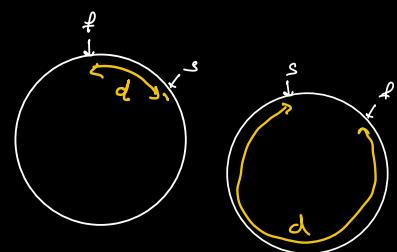
To do M iterations from the meeting point:

start one forth from head \longrightarrow
the other from meeting point

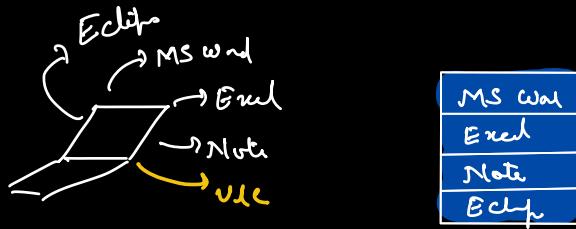


$TC : O(n)$

d
 \downarrow
 $d-1$
 \downarrow
 $d-2$
 \downarrow
 $d-3$
 \vdots
 \circ



Cache



MS Word
Email
Note
Eclips

Cache Eviction Remove from the cache.

LRU : Least Recently Used
(Cache eviction policy)

Music

MRU \Rightarrow

Munni
Rock
Enemey
Believers

LRU \Rightarrow

Songs

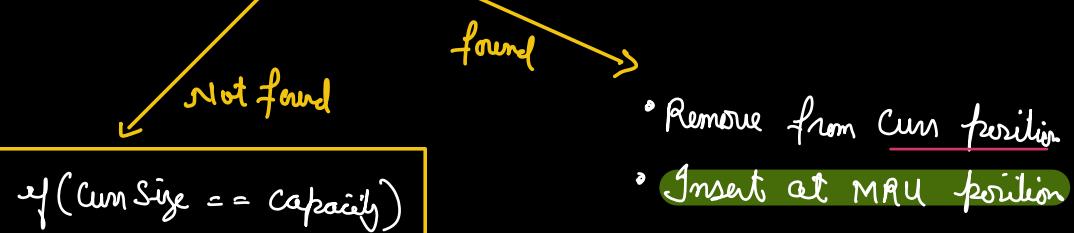
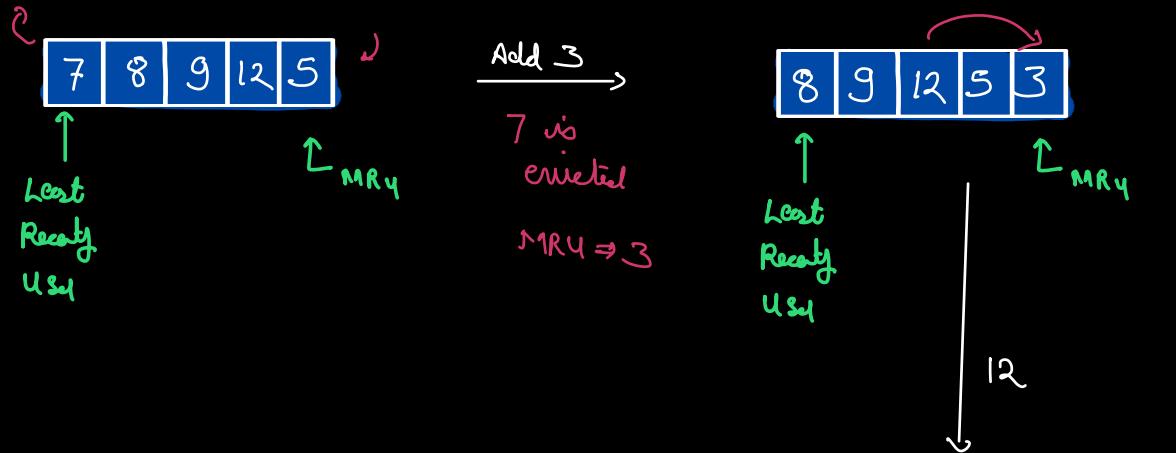
MRU \Rightarrow

On 3
On 1
On 2
On 4

LRU \Rightarrow

Browsers history

7, 8, 9, 12, 5, 3, 12



- Remove from LRU
- Insert at **MRU** position

- Insert at **MRU** position

- Search(x)
- delete
- Insert at one end

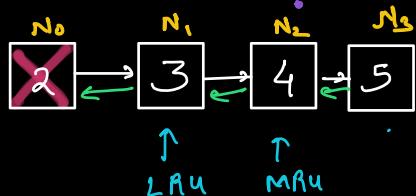
	Arrays	L L	DL + HM
Search	$O(N)$	$O(N) \rightarrow O(1)$ (HM)	$O(1)$
delete	$O(N)$	$O(1)$ (clone after search)	$O(1)$
Insert (at one end)	$O(1)$	$O(1)$	$O(1)$

To optimize search \longrightarrow Hash Map

HashMap < Integer, ListNode > map;

C = 3

2, 3, 4, -5, 4



size: $\emptyset \neq 3$

Map

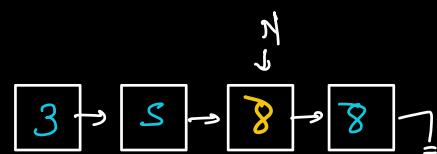
id : Node

2 : N₀

3 : N₁

4 : N₂

5 : N₃



MFU, LFU

TTL + LRU
↳ time to live.

class DLL {

int date;
DLL next;
DLL prev;

}

