

# COMP2046

## Problem Solving Using Object Oriented Approach

### Additional - Recursion

not part of the exam

# What is Recursion?

- A **recursion** (aka recursive method, recursive function) is a method calling itself directly or indirectly.
- Prior to this chapter, we have shown some cases of accidental recursion which ends up as an infinite loop, e.g.

- Lecture 9 - Professor - if we remove `super` (directly call itself)

```
class Person {  
    private final String name;  
    Person(String name) {this.name = name;}  
    protected String getName() {return name; }  
}  
class Professor extends Person {  
    public Professor(String name) {  
        super(name);  
    }  
    public String getName() {  
        //if super is missing  
        return "Dr. " + getName();  
    }  
}
```

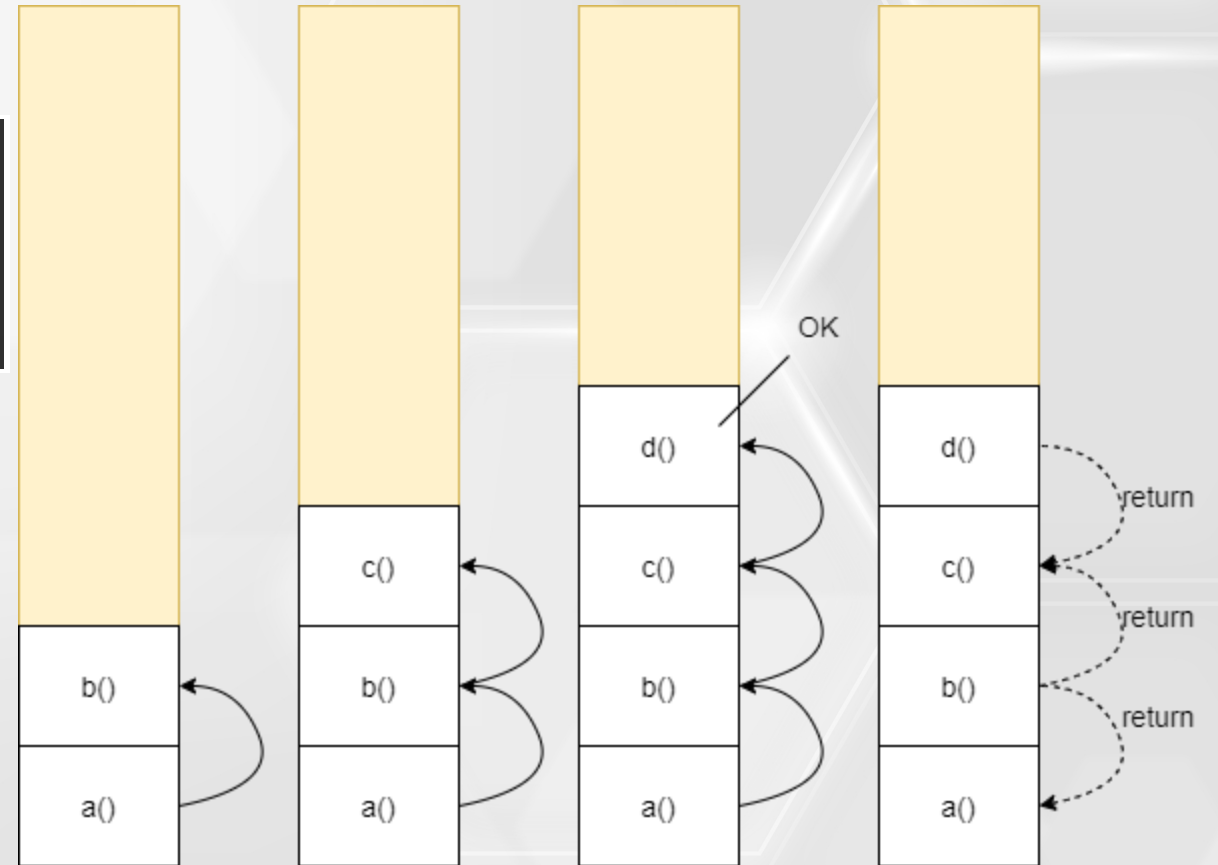
- Lecture 6 - Make Friends (indirectly call itself)

```
public void makeFriend(Person a) {  
    if (numOfFriend < 5) {  
        friends[numOfFriend++] = a;  
        a.makeFriend(this);  
    }  
}
```

# What was happening?

- Stack layout of methods calling

```
void a() { b(); }  
void b() { c(); }  
void c() { d(); }  
void d() { System.out.println("OK"); }
```

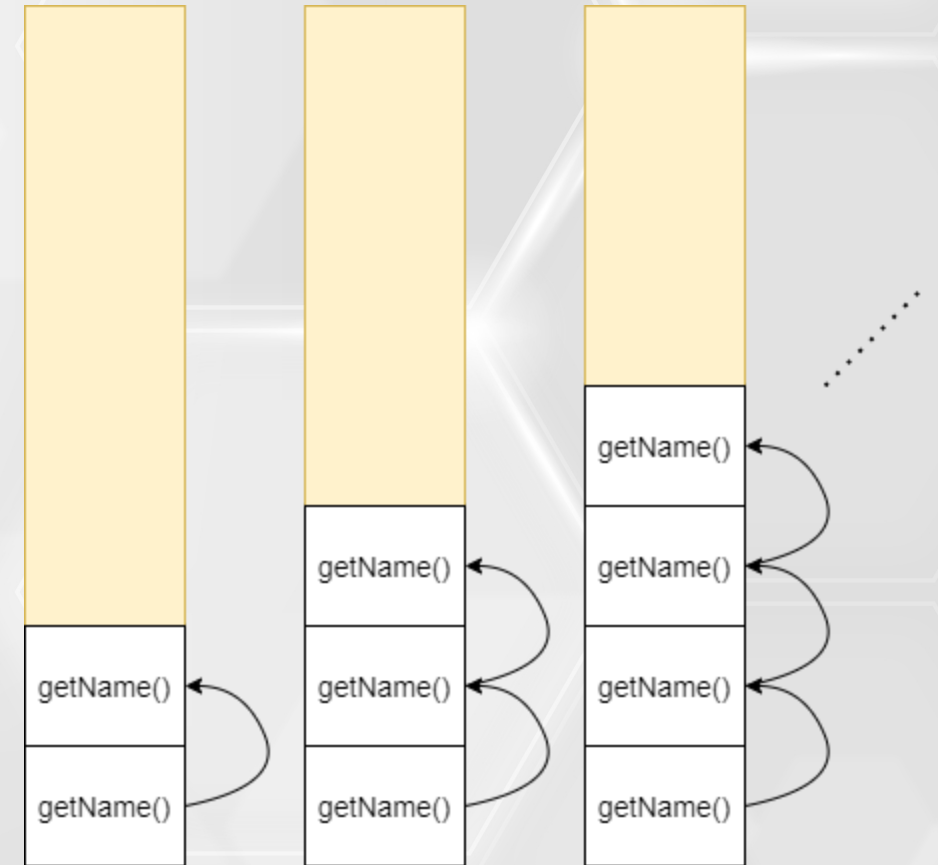


# What was happening?

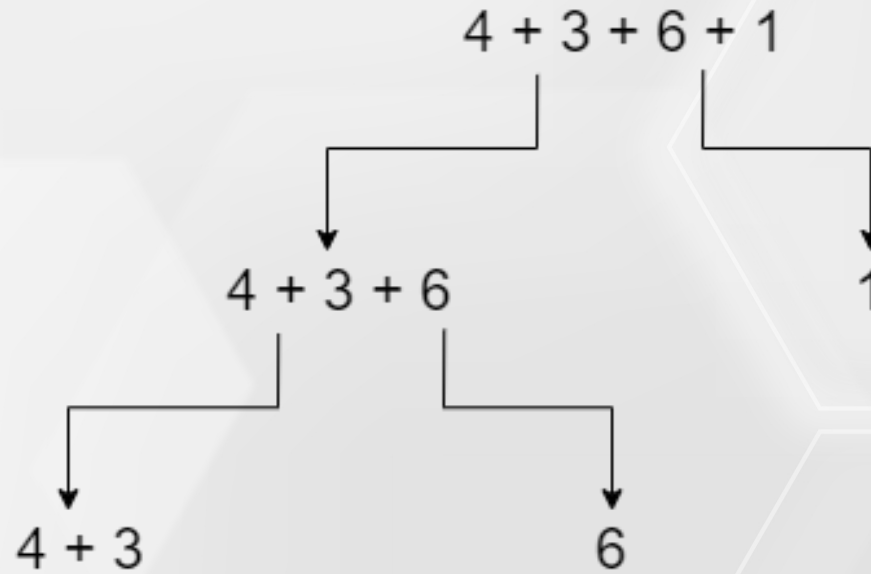
- Stack layout of methods calling

```
class Professor extends Person {  
    ...  
    public String getName() {  
        return "Dr. " + getName();  
    }  
}
```

- This cause **Stack Overflow**.



- A recursion algorithm can be useful to solve some special problems (e.g. searching problem).
- A recursion breaks down a bigger problem into a smaller problem instance.
  - Divide-and-conquer
- Like mathematical induction (if you know what is it).



# A first glance of a working recursion

```
void printNumber(int n) {  
    if (n == 0)  
        return;  
    System.out.print(n + " ");  
    printNumber(n-1);  
}  
void run() {  
    printNumber(5);  
}
```

5 4 3 2 1

We loop without loop! 🤯

# Key-steps of Recursion

1. Always starts with a base case
  2. Call yourself at certain point, with a smaller input
  3. Return the value to your caller
- **Base case:** some very simple situation that you can actually hardcode it
  - **Call yourself:** the recursion. Use yourself to solve the problem
  - **Return value:** return value is not a must for recursion. Yet, many problems use return values in their recursion.

# Recursion example - Factorial

**Factorial:**  $n!$  is defined as  $n! = n \times (n - 1) \times \dots \times 2 \times 1$  with  $0! = 1! = 1$

An ordinary factorial method (without recursion)

```
int factorial(int n) {  
    int result = 1;  
    if (n < 0)  
        throw new ArithmeticException("n cannot be negative");  
    for (int i = 1; i <= n; i++)  
        result *= i;  
    return result;  
}
```



# Recursion example - Factorial #1

A recursion of factorial, started with the base case. The simplest case you know how to do it

```
int factorial(int n) {  
    if (n < 0)  
        throw new ArithmeticException("n cannot be negative");  
    if (n == 1 || n == 0)  
        return 1;  
    // rest of the code...  
}
```



# Recursion example - Factorial #2

- Now imagine Kevin is a very kind person. He is kind enough to write a method `int X()` for you. This method can help you compute factorial of a smaller instance, say  $(n - 1)!$ 
  - You don't challenge Kevin how he implements this method. All you know is `int X()` can really compute  $(n - 1)!$ .
- Now, can you use this `X()` to finish your program?

```
int factorial(int n) {  
    if (n < 0)  
        throw new ArithmeticException("n cannot be negative");  
    if (n == 1 || n == 0)  
        return 1;  
    // rest of the code...  
}
```

# Recursion example - Factorial #2 (con't)

```
int factorial(int n) {  
    if (n < 0)  
        throw new ArithmeticException("n cannot be negative");  
    if (n == 1 || n == 0)  
        return 1;  
    int result = X() * n;  
    return result;  
}
```

- Why `result = X() * n`?
- If `X()` really can compute  $(n - 1)!$  ➡  $(n - 1)! * n = n!$
- The only problem is
  - Kevin is not that kind to write me `int X()...`

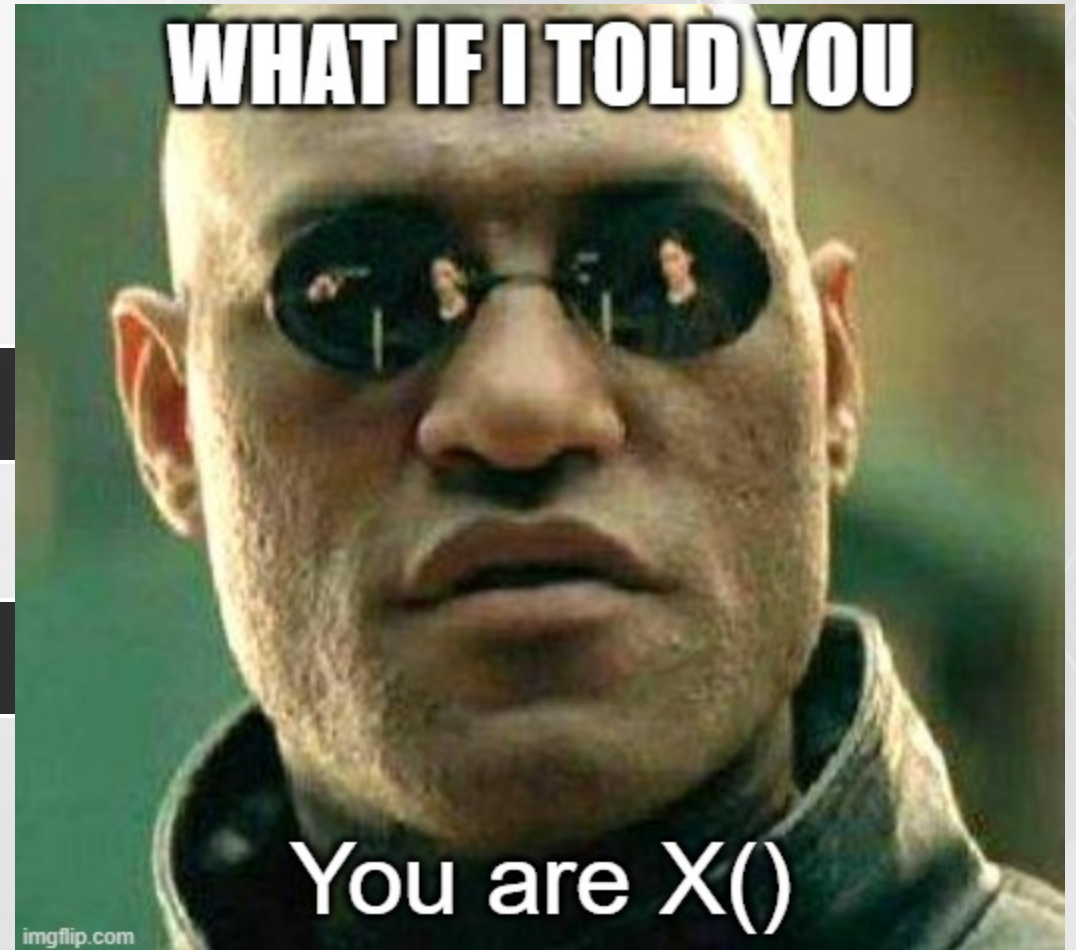
# Recursion example

- Revisit what is the ***purpose*** of your method `int factorial(int n)`
  - To compute  $n$  factorial!
- So `X()` is indeed `factorial(n-1)`.

```
int result = X() * n;
```

is actually

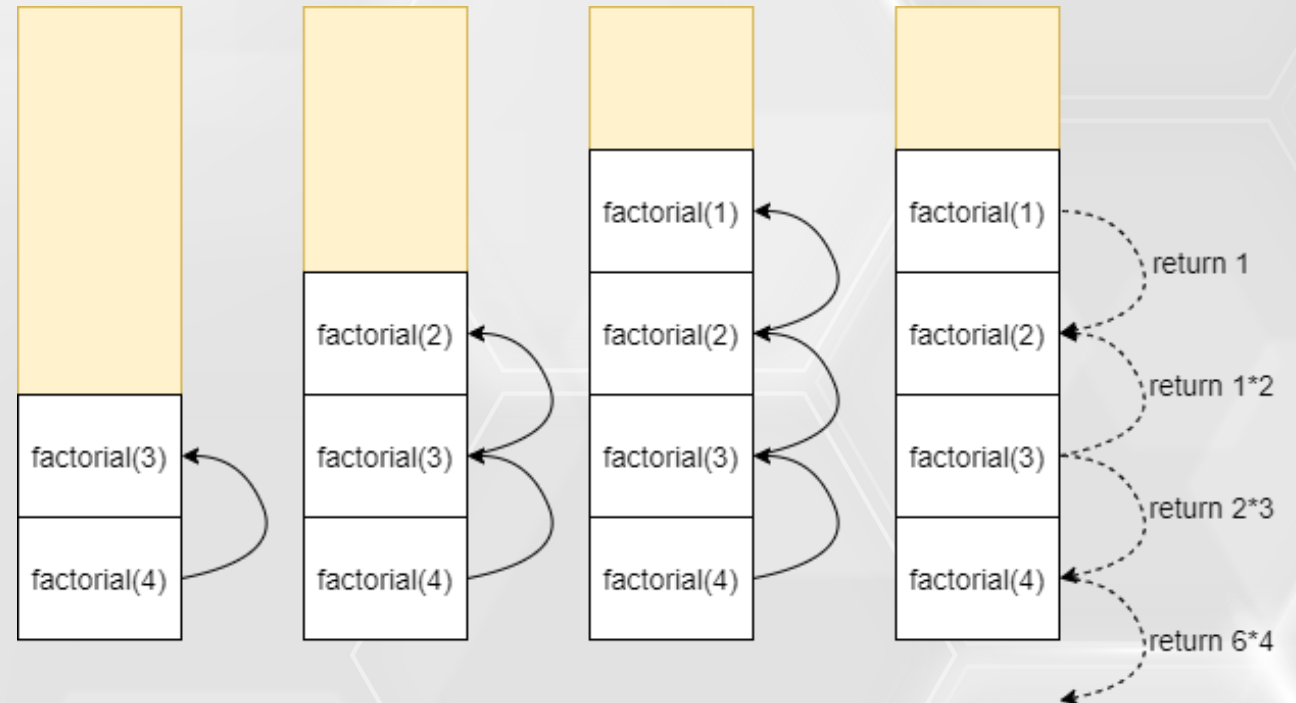
```
int result = factorial(n-1) * n;
```



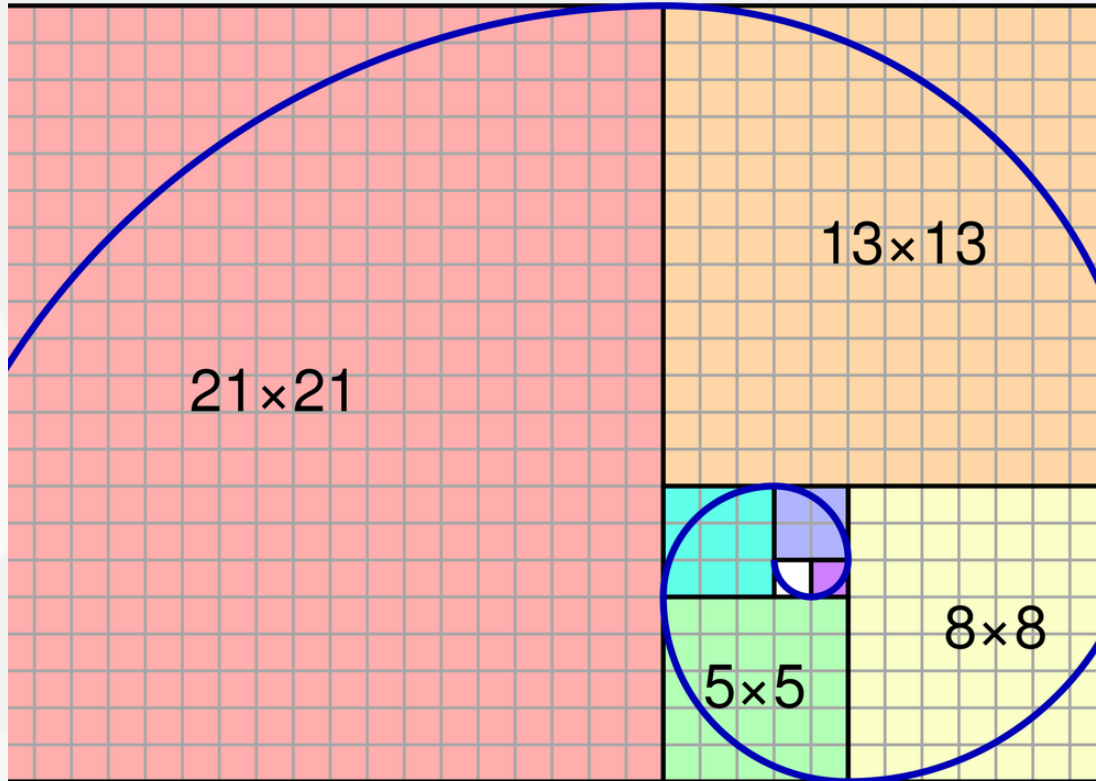
# Factorial Recursion

```
int factorial(int n) {  
    if (n < 0)  
        throw new ArithmeticException("n cannot be negative");  
    if (n == 1 || n == 0)  
        return 1;  
    return factorial(n-1) * n;  
}
```

1. Base case:  $n == 1$
2. Call yourself with input =  $n-1$
3. Return value to caller



# Fibonacci Sequence - the Golden Ratio



## The Fibonacci Sequence

1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, 233, 377...

$$1+1=2$$

$$1+2=3$$

$$2+3=5$$

$$3+5=8$$

$$5+8=13$$

$$8+13=21$$

$$13+21=34$$

$$21+34=55$$

$$34+55=89$$

$$55+89=144$$

$$89+144=233$$

$$144+233=377$$

$$f(n) = f(n-1) + f(n-2), \forall n \geq 3$$
$$f(1) = f(2) = 1$$

# Fibonacci Sequence Do together

1. What is your base case?
2. How would Kevin help you this time?
3. Return value?

```
int fib(int n) {  
    ...  
}
```

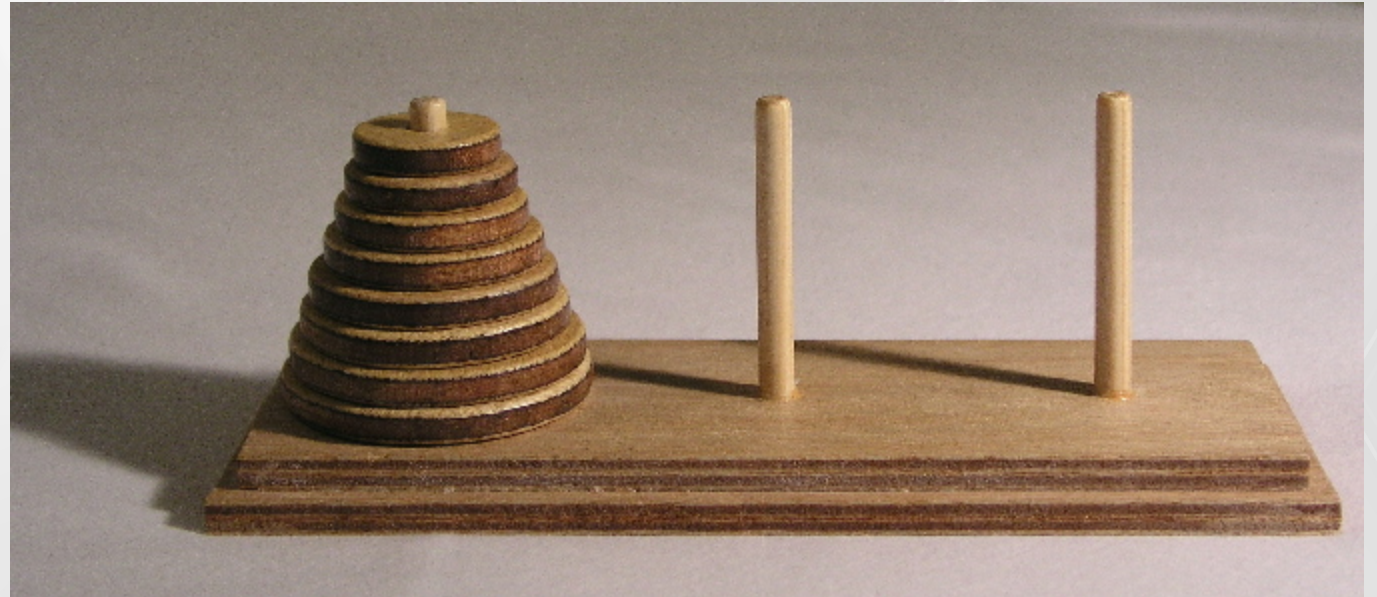




- Towers of Hanoi is one of the classic problems every budding computer scientist must grapple with. Legend. Epic. Classic

## Rules




- Move one disc at a time
- Big disc cannot on top of a small disc
- Goal: Move the disc from tower A to tower B, using tower C as a buffer



Let's work out how two-discs, three discs can be moved.



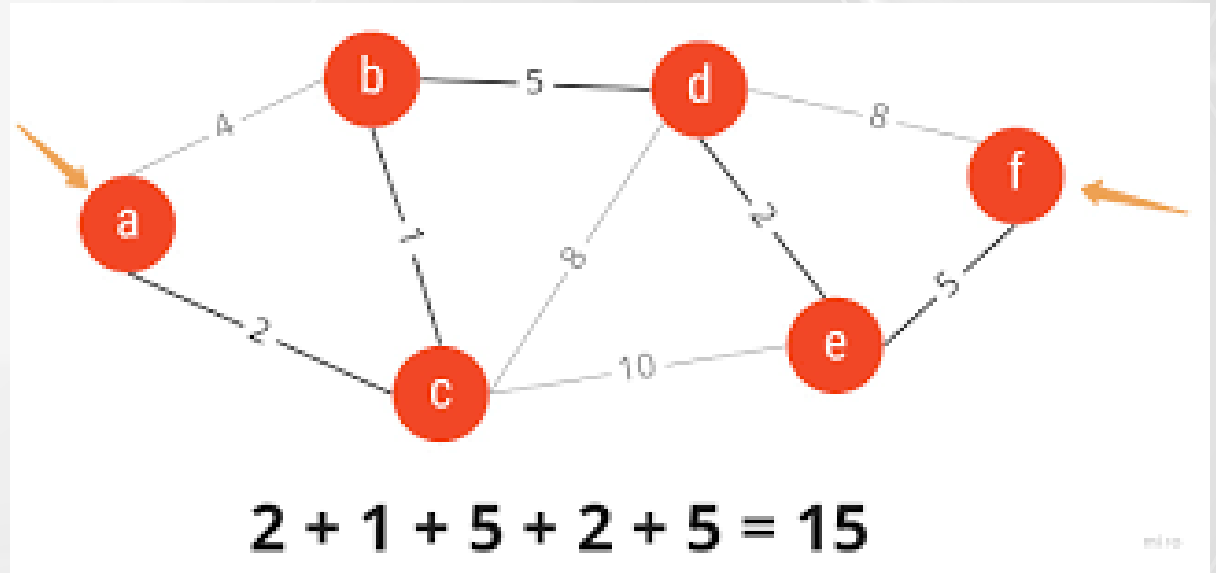
```
public void solveTower(int disks, char src, char dest, char buffer) {  
    if (disks == 1) {  
        System.out.printf("\n Move %c --> %c", src, dest); //base case  
        return;  
    }  
  
    //recursion step -- move discs - 1 from src to buffer  
    solveTower(disks - 1, src, buffer, dest);  
    //move 1 disc from src to dest  
    solveTower(1, src, dest, buffer);  
    //move discs - 1 from buffer back to dest  
    solveTower(disks - 1, buffer, dest, src);  
}
```

  Kevin: read a few more times. Yeah, I did not get it for the first time too! 

# Searching problem with Recursion

- How to find the shortest path from a node to another node?

	a	b	c	d	e	f
a		4	2			
b	4		1	5		
c	2	1		8	10	
d		5	8		2	8
e			10	2		5
f				8	5	



# Searching Problem with Recursion

```
boolean search(int steps, int[][] matrix, int src, int dest) {
    if (steps < 0)
        return false;
    if (src == dest) {
        System.out.print((char) ('a' + src));
        return true;
    }
    //find the neighbor that can reach the destination
    for (int i = 0; i < matrix.length; i++) {
        int nextStop = i;
        int cost = matrix[src][nextStop];
        if (search(steps - cost, matrix, nextStop, dest)) {
            System.out.printf(" < %c(%d) ", src + 'a', steps);
            return true;
        }
    }
    return false;
}
```

# Searching Problem with Recursion

- The algorithm presented above only tell if a path can be searched with a given cost.
- How to find the best path (shortest)?

```
int i = 0;
for (; i < 1000000; i++)
    if (search(i, matrix, src, dest))
        break;
System.out.println("Minimum steps required: " + i);
```

- We can use a better algorithm to find the minimum distance.

# Find the minimum distance

```
int minDistance(int[][] matrix, int src,
                int dest, boolean[] visited) {
    if (dest == src)
        return 0;
    if (visited[src])
        return INF;
    visited[src] = true;
    int distance = INF;
    for (int i = 0; i < matrix.length; i++) {
        int d = minDistance(matrix, i, dest, visited)
            + matrix[src][i];
        if (d < distance)
            distance = d; //minimum distance
    }
    visited[src] = false; //back-tracking
    return distance;
}
```

```
visited[src] = false; //back-tracking  
return distance;
```

- Why do we need to set `visited[src] = false;`?
- Because you need to clean this flag for other searching path!
  - $A > B > C > \dots$
  - $A > C > D > \dots$
  - Both path use C, if you do not unflag C, it cannot be used in another path!
- This unset is called back-tracking.

- Refer to your code in Lab 2.
- Build an autopath that reach the destination.

```
*           X
X           X
X X       X   X
        X   X
      X     X X
    X       X X
X         X   X D
```

```
* . . . . X
X   X   X . 
X X   X . . X
        X . X
      X . . X X
    X . . X X
X   . . X . . 
        . . . X D
```

```
/**
 * I want to navigate myself from the position
 * row col to the final destination 'D'.
 * If it is reachable, a path of ...
 * will be written from my current position to
 * the destination & return true
 * If it is not reachable, I will return false
 * & clean up the dots that I have written.
 *
 * @param maze input 2D array
 * @param row my row
 * @param col my col
 * @return true if reachable, false otherwise
 */
boolean autoMaze(char[][] maze, int row, int col) { }
```





# Lab - Bridge Crossing

- This is an optional lab.
- To find the solution that cross the bridge within finite time.
- No marks will be given for this lab.

# *Hall of Fame*

- Year 2021-22
  - Zhuang Ruiqi
- Year 2022-23
  - Cheung Yui Haang
  - Lee Ka Yan
  - UCHE Destiney Nnanna

