

Main Formulation

$$\min \sum_t \sum_{i \in \mathcal{N}} \sum_{j \in \Delta_i} \sum_{\omega \in \Omega} \left\{ \frac{1}{2|\Omega|} (r_{ij}(\omega, t) I_{ij}(\omega, t)) \right\} \quad (1)$$

- We have current in both directions, so we divide by 2 for loss.

Where:

$\mathcal{N} = \{1, 2, \dots, i, \dots, N\}$ set of nodes,

$\Delta_i = \{j = l, \dots, k\}$ set of neighboring nodes of node i ,

$\Omega = \{1, 2, \dots, \omega, \dots, \bar{\Omega}\}$ set of scenarios,

$r_{ij}(\omega, t)$ = resistance of the line between node i and j in scenario ω and time t ,

$I_{ij}(\omega, t)$ = squared current between nodes i and j in scenario ω and time t .

Subject to: DistFlow Constraints for specific time t

$$\sum_{k \in \Delta_i} P_{ik}(\omega, t) = P_i(\omega, t) - u_i P_{Gi}(\omega, t) \quad \forall i \in \mathcal{N}, \omega \in \Omega \quad (2a)$$

$$\sum_{k \in \Delta_i} Q_{ik}(\omega, t) = Q_i(\omega, t) - u_i Q_{Gi}(\omega, t) \quad \forall i \in \mathcal{N}, \omega \in \Omega \quad (2b)$$

$$I_{ij}(\omega, t) V_i(\omega, t) \geq (P_{ij}(\omega, t))^2 + (Q_{ij}(\omega, t))^2 \quad \forall i \in \mathcal{N}, \forall j \in \Delta_i, \forall \omega \in \Omega \quad (2c)$$

$$V_i(\omega, t) - V_j(\omega, t) - 2(r_{ij} P_{ij}(\omega, t) + x_{ij} Q_{ij}(\omega, t)) + (r_{ij}^2 + x_{ij}^2) I_{ij}(\omega, t) = 0 \quad \forall i \in \mathcal{N}, \forall j \in \Delta_i, \forall \omega \in \Omega \quad (2d)$$

Where:

- u_i : defines if node i is HV/MV substation (parameter)
- $P_i(\omega, t), Q_i(\omega, t)$: power generated by external grid
- $P_{ik}(\omega, t)$: active power flow from node i to node k at scenario ω

$$\begin{aligned} P_{ik}(\omega, t) \geq 0 &\Rightarrow \text{flow from } i \text{ to } k \\ P_{ki}(\omega, t) \leq 0 &\Rightarrow \text{flow from } k \text{ to } i \end{aligned}$$

- $Q_{ik}(\omega, t)$: reactive power from node i to node k
- $V_i(\omega, t)$: square of voltage at node i
- x_{ij} : reactance of the line from node i to node j
- $P_i(\omega, t)$: net consumption at node i

Radiality of Distribution Network

Index sets.

N : set of buses (nodes),
 L : set of undirected lines (edges),
 $S \subseteq L$: subset of lines equipped with switches,
 $nS = L \setminus S$: lines without switches,
 $C \subseteq L \times N \times N : \{(\ell, i, j)\}$ directed candidate arcs for each line.

Parameters.

$\text{slack_node} \in N$: index of the root (slack) bus,
 $\varepsilon > 0$: small constant (default 1) for numerical stability.

Decision variables.

$\delta_\ell \in \{0, 1\}$ (or $[0, 1]$ if relaxed) $\forall \ell \in S$: switch-status (1=closed),
 $f_{\ell, i, j} \in \mathbb{R}$ $\forall (\ell, i, j) \in C$: artificial “flow” on arc (ℓ, i, j) ,
 $\theta \in \mathbb{R}$: Benders-cut auxiliary variable.

1. Flow–Balance (one parent per node).

$$\sum_{\substack{(\ell, i, j) \in C \\ j=n}} f_{\ell, i, j} = \begin{cases} -\varepsilon, & n \in N \setminus \{\text{slack_node}\}, \\ \geq \varepsilon (|N| - 1), & n = \text{slack_node} \end{cases} \quad \forall n \in N \quad (1)$$

2. Orientation (zero net on each physical line).

$$\sum_{(\ell, i, j) \in C} f_{\ell, i, j} = 0 \quad \forall \ell \in L \quad (2)$$

3. Switch–Propagation (flow only if switch closed).

$$f_{\ell, i, j} \leq \varepsilon |N| \delta_\ell, \quad (3)$$

$$f_{\ell, i, j} \geq -\varepsilon |N| \delta_\ell \quad \forall (\ell, i, j) \in C, \ell \in S \quad (4)$$

4. Number of Closed Switches (spanning-tree size).

$$\sum_{\ell \in S} \delta_\ell = |N| - |nS| - 1 \quad (5)$$

Power & Current Relation with State of Switches

$$-M\delta_{ij} \leq P_{ij}(\omega, t) \leq M\delta_{ij} \quad \forall (i, j) \in \text{Lines With Switches} \quad (4a)$$

$$-M\delta_{ij} \leq Q_{ij}(\omega, t) \leq M\delta_{ij} \quad \forall (i, j) \in \text{Lines With Switches} \quad (4b)$$

$$0 \leq I_{ij}(\omega, t) \leq M\delta_{ij} \quad \forall (i, j) \in \text{Lines With Switches} \quad (4c)$$

$$V_i(\omega, t) - V_j(\omega, t) - 2(r_{ij}P_{ij}(\omega, t) + x_{ij}Q_{ij}(\omega, t)) + (r_{ij}^2 + x_{ij}^2)I_{ij}(\omega, t) + M(1 - \delta_{ij}) \geq 0 \quad (4d)$$

$$V_i(\omega, t) - V_j(\omega, t) - 2(r_{ij}P_{ij}(\omega, t) + x_{ij}Q_{ij}(\omega, t)) + (r_{ij}^2 + x_{ij}^2)I_{ij}(\omega, t) - M(1 - \delta_{ij}) \leq 0 \quad (4f)$$

$$\forall (i, j) \in \text{Lines With Switches}$$

Constraints Limiting:

$$V_i^{min} \leq V_i(\omega, t) \leq V_i^{max} \quad \forall i \in \mathcal{N} \quad (6)$$

$$0 \leq I_{ij}(\omega, t) \leq I_{ij}^{max} \quad \forall i \in \mathcal{N}, \forall j \in \Delta_i \quad (7)$$

$$V_i(\omega, t) = \tilde{V}_i(\omega, t) \quad \forall i \in \text{Substation Nodes}, \forall \omega \in \Omega \quad (8)$$

Where:

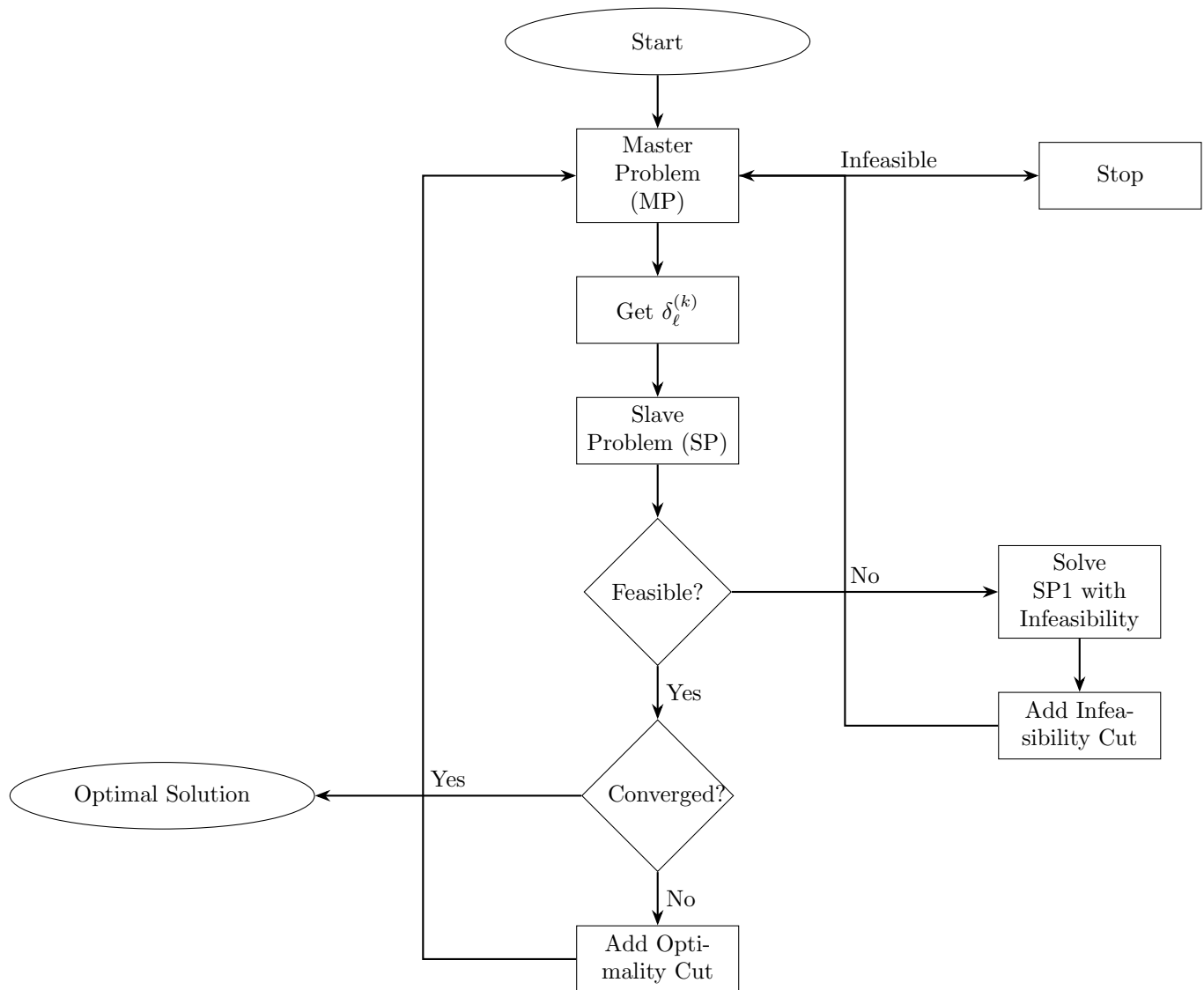
- $\underline{V}_i, \bar{V}_i$ are lower/upper levels of voltage.
- $\tilde{V}_i(\omega, t)$ is the voltage of substation at scenario ω (imposed by external grid).

Sets:

- \mathcal{N} : set of nodes
- Δ_i : set of neighboring nodes of node i
- Ω : set of scenarios
- \mathcal{LS} : set of lines with switches
- \mathcal{SN} : set of substation nodes

Decision Variables and Parameters:

Symbol	Description	Index Domain
$I_{ij}(\omega, t)$	Squared current from node i to j in scenario ω	$\forall i \in \mathcal{N}, j \in \Delta_i, \omega \in \Omega$
$P_{ij}(\omega, t)$	Active power flow from i to j in scenario ω	Same as above
$Q_{ij}(\omega, t)$	Reactive power flow from i to j in scenario ω	Same as above
$V_i(\omega, t)$	Squared voltage at node i in scenario ω	$\forall i \in \mathcal{N}, \omega \in \Omega$
$P_i(\omega, t)$	Net consumption at node i (parameter)	$\forall i \in \mathcal{N}, \omega \in \Omega$
$Q_i(\omega, t)$	Net reactive power at node i (parameter)	Same as above
$P_{Gi}(\omega, t)$	Power generated by external grid at node i	$\forall i \in \text{ExternalPoints}, \omega \in \Omega$
$Q_{Gi}(\omega, t)$	Reactive power from external grid	Same as above
d_{ij}	Binary variable for direction between i and j	$\forall (i, j) \in \text{LinesWithSwitches}$
δ_{ij}	Binary variable for line/switch status	Same as above



Slave Subproblem SP^ω for each scenario $\omega \in W$

Index sets.

N : set of buses, L : set of undirected lines,
 $S \subseteq L$: lines equipped with switches, $C \subseteq L \times N \times N$: directed candidate arcs,
 W : set of scenarios.

Given from the master: $\delta_\ell^* \in \{0, 1\}$ for each $\ell \in S$.

Variables (per ω):

$v_n^2(\omega, t)$, $p_{ij}(\omega, t)$, $q_{ij}(\omega, t)$, $I_{ij}(\omega, t)$, (and if infeasible: slack variables $s_{vn}(\omega, t)$, $s_{Iij}(\omega, t)$).

Objective.

$$\min_{v^2, p, q, I} \sum_{(i,j) \in C} \frac{1}{2} r_{ij} I_{ij}(\omega, t) \longrightarrow \text{optimality cost } \text{obj}^{\text{feas}, (k)}(\omega, t),$$

or when infeasible

$$\min \sum_{n \in N} (s_{vn}^+(\omega, t) + s_{vn}^-(\omega, t)) + \sum_{(i,j) \in C} s_{Iij}(\omega, t) \longrightarrow \text{obj}^{\text{infeas}, (k)}(\omega, t).$$

Constraints (for each ω):

1. Switch-status propagation:

$$\delta_{\ell\omega} = \delta_\ell^*, \quad \forall \ell \in S.$$

2. Slack-bus voltage:

$$V(\omega, t) = V_n(\omega, t), \quad \text{if } n = \text{slack_node}, \quad \text{else skipped.}$$

3. Nodal active-power balance:

$$\sum_{(i,j) \in C: i=n} p_{ij}(\omega, t) = \begin{cases} -P^{\text{slack}}(\omega, t), & n = \text{slack_node}, \\ -P_n(\omega, t), & \text{otherwise.} \end{cases}$$

4. Nodal reactive-power balance:

$$\sum_{(i,j) \in C: i=n} \left(q_{ij}(\omega, t) - \frac{1}{2} b_{ij} V_i(\omega, t) \right) = \begin{cases} -Q^{\text{slack}}(\omega, t), & n = \text{slack_node}, \\ -Q_n(\omega, t), & \text{otherwise.} \end{cases}$$

5. Branch voltage-drop (DistFlow):

$$V_i(\omega, t) - V_j(\omega, t) - 2(r_{ij} p_{ij}(\omega, t) + x_{ij} q_{ij}(\omega, t)) + (r_{ij}^2 + x_{ij}^2) I_{ij}(\omega, t) \begin{cases} = 0, & \ell \notin S, \\ \geq -M_{ij}(1 - \delta_\ell(\omega, t)), \\ \leq +M_{ij}(1 - \delta_\ell(\omega, t)), & \ell \in S. \end{cases}$$

6. Rotated-cone current constraint (no switch):

$$p_{ij}^2(\omega, t) + q_{ij}^2(\omega, t) \leq I_{ij}(\omega, t) \frac{V_i(\omega, t)}{n_{ij}^2}, \quad \forall \ell \notin S.$$

If $\ell \in S$, simply $I_{ij}(\omega, t) = 0$.

7. Branch power-balance (Kirchhoff):

$$\sum_{(i,j) \in C: \ell=ij} (p_{ij}(\omega, t) - \frac{1}{2} r_{ij} I_{ij}) = 0, \quad \sum_{(i,j) \in C: \ell=ij} (q_{ij}(\omega, t) - \frac{1}{2} x_{ij} I_{ij}(\omega, t)) = 0.$$

8. Switch P/Q bounds ($\ell \in S$):

$$-M_{ij} \delta_\ell(\omega, t) \leq p_{ij}(\omega, t) \leq M_{ij} \delta_\ell(\omega, t), \quad -M_{ij} \delta_\ell(\omega, t) \leq q_{ij}(\omega, t) \leq M_{ij} \delta_\ell(\omega, t).$$

9. Physical limits (only in optimal subproblem):

$$0 \leq I_{ij}(\omega, t) \leq I_{ij}^{\max}, \quad V_n^{\min}(\omega, t) \leq v_n^2(\omega, t) \leq V_n^{\max}(\omega, t).$$

Infeasible Subproblem $\text{SP}_{\text{infeas}}(\omega, t)$

Objective (infeasibility penalties):

$$\min_{V, p, q, I, s_v^\pm, s_I} \sum_{n \in N} (s_{vn}^+(\omega, t) + s_{vn}^-(\omega, t)) + \sum_{(i,j) \in C} s_{Iij}(\omega, t) \longrightarrow \text{obj}^{\text{infeas}, (k)}(\omega, t) \quad (9)$$

Relaxed physical-limit constraints (per ω):

$$V_n(\omega, t) \leq V_n^{\max} + s_{vn}^+(\omega, t), \quad \forall n \in N, \quad (10)$$

$$V_n(\omega, t) \geq V_n^{\min} - s_{vn}^-(\omega, t), \quad \forall n \in N, \quad (11)$$

$$I_{ij}(\omega, t) \leq \bar{I}_{ij}^{\max} + s_{Iij}(\omega, t), \quad \forall (i, j) \in C, \quad (12)$$

$$s_{vn}^+(\omega, t), s_{vn}^-(\omega, t), s_{Iij}(\omega, t) \geq 0, \quad \forall n \in N, \forall (i, j) \in C. \quad (13)$$

All of the network constraints

{Slack-voltage, Power-balances, Voltage-drops, Conic-current, Switch-bounds, ...}

are assembled in `slave_model_constraints(model)` and thus appear in both the *optimal* and *infeasible* subproblem models. The only differences between them are considering the slack variable in the physical limitation of current and voltage

Master Problem at Iteration k

$$\begin{aligned}
& \min_{\{\theta_\omega\}_{\omega \in \Omega}, \{\delta_\ell\}_{\ell \in S}, \{f_{\ell,i,j}\}_{(\ell,i,j) \in C}} \sum_{\omega \in \Omega} \theta_\omega \\
& \text{s.t.} \quad \underbrace{\sum_{\substack{(\ell,i,j) \in C \\ j=n}} f_{\ell,i,j} = \begin{cases} -\varepsilon, & n \in N \setminus \{\texttt{slack_node}\}, \\ \geq \varepsilon(|N| - 1), & n = \texttt{slack_node}, \end{cases}}_{(1) \text{ Flow-Balance}} \quad \forall n \in N, \\
& \quad \underbrace{\sum_{\substack{(\ell,i,j) \in C \\ \ell=\ell'}} f_{\ell,i,j} = 0}_{(2) \text{ Orientation}} \quad \forall \ell' \in L, \\
& \quad \underbrace{-\varepsilon|N|\delta_\ell \leq f_{\ell,i,j} \leq \varepsilon|N|\delta_\ell}_{(3) \text{ Switch-Propagation}} \quad \forall (\ell,i,j) \in C, \ell \in S, \\
& \quad \underbrace{\sum_{\ell \in S} \delta_\ell = |N| - |nS| - 1}_{(4) \text{ Spanning-Tree Size}} \\
& \quad \underbrace{\theta_\omega \geq \varphi_\omega^{(m)} + \sum_{\ell \in S} \pi_{\ell\omega}^{(m)} (\delta_\ell - \delta_\ell^{(m)})}_{\substack{\text{Optimality cuts} \\ m=1,\dots,k-1}} \quad \forall \omega \in \Omega, \\
& \quad \underbrace{0 \geq \psi_\omega^{(m)} + \sum_{\ell \in S} \rho_{\ell\omega}^{(m)} (\delta_\ell - \delta_\ell^{(m)})}_{\substack{\text{Infeasibility cuts} \\ m=1,\dots,k-1}} \quad \forall \omega \in \Omega, \\
& \quad \delta_\ell \in \{0, 1\}, \forall \ell \in S, \quad f_{\ell,i,j} \in \mathbb{R}, \forall (\ell,i,j) \in C, \quad \theta_\omega \in \mathbb{R}, \forall \omega \in \Omega.
\end{aligned}$$

Notes on the Benders-cut notation: - $\varphi_\omega^{(m)}, \pi_{\ell\omega}^{(m)}$ are the dual-derived intercept and slope coefficients from the *optimality* subproblem for scenario ω at previous iteration m . - $\psi_\omega^{(m)}, \rho_{\ell\omega}^{(m)}$ likewise come from an *infeasibility* cut for scenario ω , iteration m . - $\delta_\ell^{(m)}$ is the value of δ_ℓ at iteration m .