# A Concatenation Operation to Derive Autosegmental Graphs

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### Introduction



- ► What kind of structures are autosegmental phonological representations (APRs; Goldsmith, 1976)?
- ► Current conception: top-down constraints on graph structures
- ► Our result: these fundamental properties emerge from concatenation of an alphabet of primitives (Engelfriet and Vereijken, 1997; Courcelle et al., 2012)
- ► This makes them directly comparable to strings

### Introduction

#### Overview

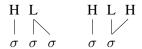
- ▶ Why APRs & what are APRs
- ► Representation of APRs as labeled mixed graphs
- ► Top-down, axiomatic definition of APRs
- Constructive derivation of axioms from concatenation & primitives
- ► Empirical consequences
- ► Discussion

### **APRs**

- ► What are APRs?
- ► Mende (Leben, 1973; Goldsmith, 1976): fewer attested tone combinations than possible F, HL, HLLL, ...
  \*HF, \*HLHL, \*HHLL, \*HLLH, ...
- ► Answer: tone melodies independent from syllables

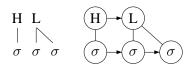


### **APRs**



- ► Fundamentally, APRs are:
  - ► Ordered tiers of like autosegments
  - ► Association lines connecting items on distinct tiers
- ► Set of valid APRs usu. defined *axiomatically* (Goldsmith, 1976; Sagey, 1986; Bird and Klein, 1990; Kornai, 1995)

### **APGs**

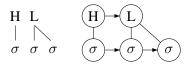


► To represent APRs, we can use (simple) *labeled mixed graphs* 

$$G = \langle V, E, A, \ell \rangle$$

- $ightharpoonup GR(\Sigma)$ 
  - $V = \{0, 1, ..., n\}$
  - $ightharpoonup E \subseteq \mathcal{P}_{=2}(V)$
  - $ightharpoonup A \subseteq V \times V$
  - $\blacktriangleright \ \ell: V \to \Sigma$

### **APGs**

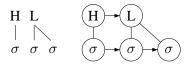


► To represent APRs, we can use (simple) *labeled mixed graphs* 

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- ► Why graphs?
  - ► Graphs explicitly encode all information in APRs
  - ▶ Deep literature on pattern matching, characterizing & learning graph sets, graph transductions, etc. (Courcelle et al., 2012; Engelfriet and Hoogeboom, 2001; López et al., 2012; Ferrera, 2013, inter alia)

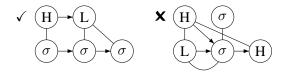
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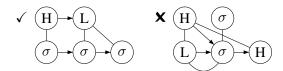
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$$G = \langle V, E, A, \ell \rangle$$

- ► To select *well-formed* autosegmental phonological graphs (APGs), there are two options
  - ► Axiomatically (top-down approach)
  - ► Constructively (bottom-up approach)



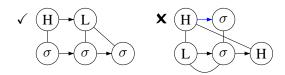
- ▶ We first define the basic autosegmental structure in  $APG(\Sigma, T)$
- ▶ We consider two *tiers*



▶ Let  $\leq$  be the reflexive, transitive closure of *A* 

### Axiom (Tier structure)

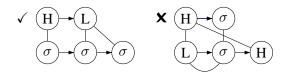
*V* is partitioned into two sets  $V_0$ ,  $V_1$  which are each totally ordered by  $\preccurlyeq$ .  $V_0$  and  $V_1$  are the **tiers** of *G*.



▶ Let *T* be a partition  $\{T_t, T_m\}$  on  $\Sigma$ 

Axiom (Like tier labels)

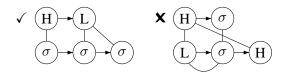
- ▶ Let  $V_m$  correspond to  $T_m$ , likewise  $V_t$ ,  $T_t$
- ▶ Let's set  $T_m = \{H, L\}, T_t = \{\sigma\}$



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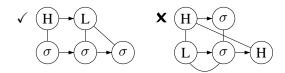
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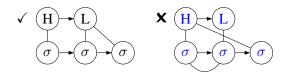
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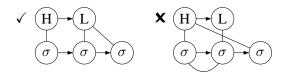
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► Associations can only be between tiers

Axiom (Associations between tiers)

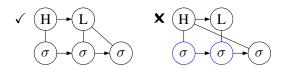
For all  $\{x, y\} \in E$ ,  $x \nleq y$  and  $y \nleq y$ .



► Associations can only be between tiers

Axiom (Associations between tiers)

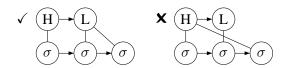
For all  $\{x,y\} \in E$ ,  $x \not\leq y$  and  $y \not\leq y$ .



▶ No line crossing (Goldsmith, 1976; Hammond, 1988; Coleman and Local, 1991)

### Axiom (NCC)

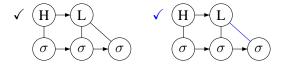
For all  $u, v, x, y \in V$ , if  $\{u, x\}, \{v, y\} \in E$  and  $u \leq v$ , then  $x \leq y$ 

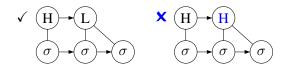


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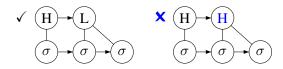




► The Obligatory Contour Principle Leben (1973); Goldsmith (1976); McCarthy (1986); Odden (1986)

Axiom (OCP)

For tier  $V_m$ , for all  $x, y \in V_m$ ,  $(x, y) \in A$  implies  $\ell(x) \neq \ell(y)$ 



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 $ightharpoonup APG(\Sigma,T)$ : graphs in  $GR(\Sigma)$  obeying these axioms

Axioms 1) Tier structure

4) NCC

2) Like tier labels

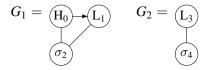
- 5) OCP
- 3) Associations between tiers
- ► These properties can all be derived from concatenation

► Engelfriet and Vereijken (1997); Courcelle et al. (2012): two methods of graph concatenation:

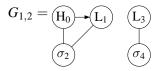
One way is to "glue" them together, by identifying some of their vertices. The other way is to "bridge" them (or rather, "bridge the gap between them"), by adding edges between their vertices. (Courcelle et al., 2012, p. 6)

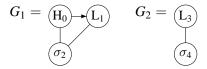
► We define a specialized version of this machinery (using both "gluing" and "bridging") to generate the set of APGs

$$G_1 = \underbrace{H_0}$$
 $C_1 = \underbrace{L_3}$ 
 $G_1 \circ G_2 = \underbrace{H_0}$ 
 $G_2 = \underbrace{H_0}$ 
 $G_2 = \underbrace{G_1}$ 
 $G_3 \circ G_4 = \underbrace{G_2}$ 
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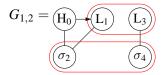


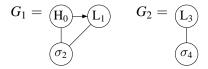
- ► Identify the 'ends' of the tiers
- $\blacktriangleright$  Merge identical end nodes on  $V_m$
- ▶ Draw arcs between any other pair of end nodes



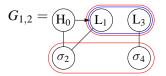


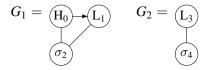
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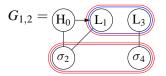


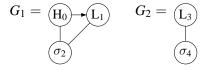
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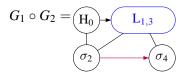




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### The identity graph

▶ Let 
$$G_{\lambda} = \langle \emptyset, \emptyset, \emptyset, \emptyset \rangle$$

Theorem ( $G_{\lambda}$  is identity)

For any 
$$G \in GR(\Sigma)$$
,  $G \circ G_{\lambda} = G_{\lambda} \circ G = G$ .

### **Graph primitives**

- ► Engelfriet and Vereijken (1997): graph sets can interpret string sets
- ► Concatenation of symbols = concatenation of graph primitives

### Definition

An alphabet of graph primitives over  $GR(\Sigma)$  is a finite set  $\Gamma$  of symbols and a naming function  $g:\Gamma\to GR(\Sigma)$ .

### Definition

For  $w \in \Gamma^*$ ,  $g(w) \stackrel{\text{def}}{=}$ 

- $G_{\lambda}$  if  $w = \lambda$
- $g(u) \circ g(\gamma)$  if  $w = u\gamma, u \in \Gamma^*, \gamma \in \Gamma$

### APG graph primitives

Definition (APG graph primitive)

For  $\Sigma$  and  $T = \{T_t, T_m\}$ , an *APG graph primitive* is a graph  $G \in GR(\Sigma)$  for which

- a.  $V_t$  is a singleton set  $\{v_t\}$
- b.  $V_m$  is totally ordered by  $\leq$
- c. All  $e \in E$  are of the form  $\{v_m, v_t\}, v_m \in V_m$

$$\Gamma = \{H, L, F\}$$

$$g(H) = \underbrace{H}_{\sigma} g(L) = \underbrace{L}_{\sigma} g(F) = \underbrace{H}_{\sigma} L$$

### **Graph primitives**

Def. of APG prim. a. 
$$V_t = \{v_t\}$$
 b.  $\preccurlyeq$  totally orders  $V_m$  c. All  $e \in E = \{v_m, v_t\}$  for  $v_m \in V_m$  
$$\Gamma = \{H, L, F\} \quad g(H) = H \quad g(L) = L \quad g(F) = H L$$

$$PG(\Gamma) = \{g(w)|w \in \Gamma^*\}$$
 
$$g(\text{HHH}) \quad g(\text{HLL}) \quad g(\text{LF})$$
 
$$(\sigma) + (\sigma) +$$

APG prim. a. 
$$V_t = \{v_t\}$$
 Ax. 1) Tier struct. 4) NCC b.  $\preccurlyeq$  totally orders  $V_m$  2) Tier labels 5) OCP c. All  $e \in E = \{v_m, v_t\}$  3) Assoc. for  $v_m \in V_m$ 

### Theorem

### Every $G \in APG(\Gamma)$ follows Ax. 1-5

- ► APG graph primitives have two tiers V<sub>t</sub> and V<sub>m</sub> and concatenation preserves this
- $\triangleright$  NCC from singleton  $V_t$
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- $\circ$  is associative over  $APG(\Gamma)$ 
  - ▶ Concatenation preserves tier structure; proof similar of associativity over \(\Gamma^\*\)

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### Concatenation and empirical phenomena

- ightharpoonup  $APG(\Gamma)$  also naturally captures facts of phonology
- ► Unbounded spreading, but no 'unbounded contouring'

$$g(\mathbf{H}) = \begin{matrix} \mathbf{H} \\ \begin{matrix} \mathbf{G} \end{matrix} \qquad g(\mathbf{L}) = \begin{matrix} \mathbf{L} \\ \begin{matrix} \mathbf{G} \end{matrix} \qquad g(\mathbf{F}) = \begin{matrix} \mathbf{H} \end{matrix} + \begin{matrix} \mathbf{L} \\ \begin{matrix} \mathbf{G} \end{matrix} \qquad g(\mathbf{F}) = \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \\ \mathbf{G} \end{matrix} + \begin{matrix} \mathbf{G} \end{matrix} +$$

 $g(\mathbf{H}^n) = \underbrace{\mathbf{H}}_{\sigma_1 \bullet \sigma_2 \bullet \cdots \bullet \sigma_n}$ 

► No 
$$w \in \Gamma^*$$
 s.t.  $g(w) = H + L + H + L$ 

▶ (unless added to  $\Gamma$ )

### Concatenation and empirical phenomena

► Most OCP violations at morpheme boundaries or signaled by downstep (Hyman, 2014)

$$g(\mathbf{H}) = \underbrace{\mathbf{H}}_{\boldsymbol{\sigma}} \quad g(^{!}\mathbf{H}) = \underbrace{\mathbf{H}}_{\boldsymbol{\sigma}} \underbrace{\mathbf{H}}_{\boldsymbol{\sigma}} \quad g(\#) = \#$$

► Adjacent Hs across boundaries (Aghem; Hyman, 2014)

$$g(H \# H) = \underbrace{H} + \underbrace{\#} + \underbrace{H}$$

► Adjacent Hs tautomorphemically (Kishambaa; Odden, 1986)

$$g(HH) = H$$
  $g(H^!H) = H + H$ 

### Discussion

- ► g is direct way of relating strings to APRs
- ► Cognitive interpretation: how humans relate linear speech stream to APRs
- ▶  $(APG(\Gamma), \circ)$  as free monoid
- ▶ Multi-tier representations: n-ary partitions on  $\Sigma$
- ► Full feature geometry may require more structure on partition
- ► Less restrictive version w.r.t. OCP could include second concatenation operation without 'gluing'

### Conclusion

- We defined APRs as  $APG(\Gamma)$  parallel to strings
- ► Maintains key properties of APRs
- ► Positive empirical implications (e.g. no unbounded contouring)
- ► Restrictive theory of APRs (w.r.t OCP)

### Acknowledgements

We thank three reviewers for their insightful comments and suggestions. Adam Jardine acknowledges support from a University of Delaware Graduate Research Fellowship.