MAXIMUM LIKELIHOOD ESTIMATION OF FACTORED REGULAR DETERMINISTIC STOCHASTIC LANGUAGES

Chihiro Shibata and Jeffrey Heinz





University of Toronto July 19, 2019

We thank JSPS KAKENHI #JP18K11449 (CS) and NIH #R01HD87133-01 (JH)

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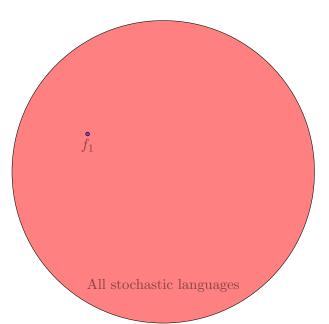
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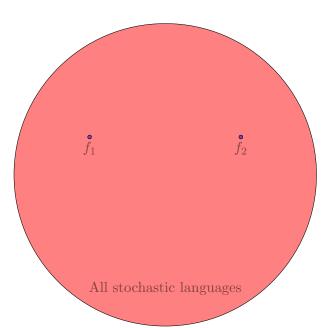
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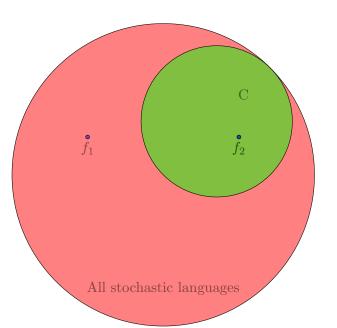
For a class of stochastic languages C, is there an algorithm which reliably returns a Maximum-Likelihood Estimate (MLE) of an observed data sample D?

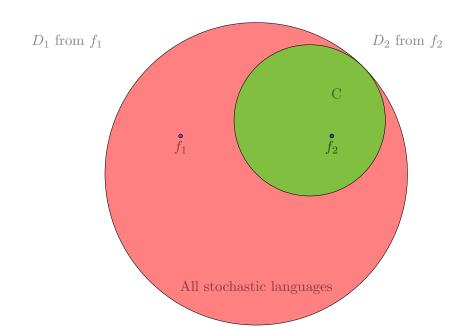
IN PICTURES All stochastic languages

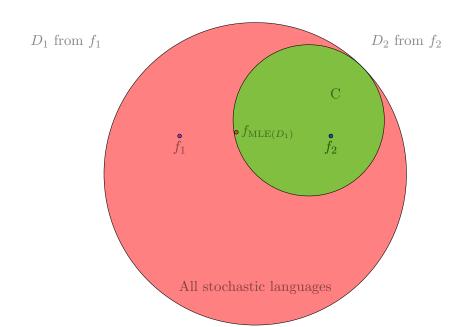
In Pictures

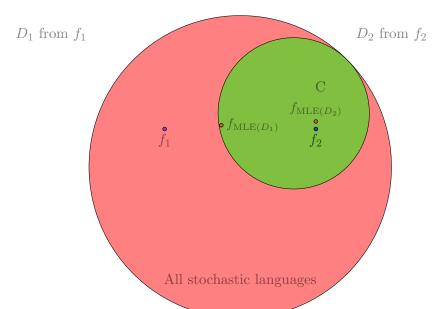






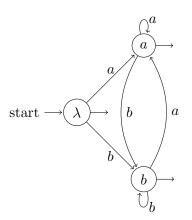






CLASSES DEFINED BY SINGLE DFAS

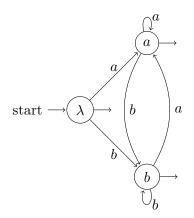
Example: Bigram model



Parameters
$ heta_{ times a}$
$ heta_{ times b}$
$ heta_{ times dash}$
$\overline{ heta_{aa}}$
$ heta_{ab}$
θ_{a}
θ_{ba}
$ heta_{bb}$
$ heta_{b\ltimes}$

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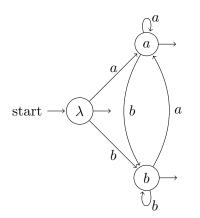
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Parameters
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$$D = \langle ab, aabb \rangle$$

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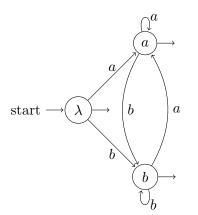
Parameters		
$\overline{\theta_{ times a}}$	1	
$ heta_{ times b}$		
$\theta_{\rtimes \ltimes}$		
θ_{aa}		
$ heta_{ab}$	1	
θ_{a} ×		
θ_{ba}		
$ heta_{bb}$		
$\theta_{b\bowtie}$	1	

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MLE is obtained by passing D through DFA and normalizing. (Vidal et al. 2005)

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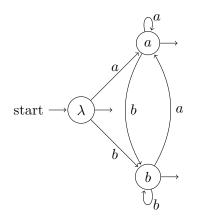


Paran	neters
$\theta_{ times a}$	2
$\theta_{\rtimes b}$	
$\theta_{ m MK}$	
θ_{aa}	1
$ heta_{ab}$	2
θ_{a}	
θ_{ba}	-
$ heta_{bb}$	1
$\theta_{b\ltimes}$	2

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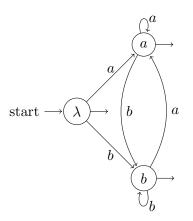
Parameters			
$ heta_{ times a}$	1		
$\theta_{\rtimes b}$	0		
$ heta_{ m imes imes}$	0		
θ_{aa}	1/3		
$ heta_{ab}$	2/3		
$\theta_{a\ltimes}$	0		
θ_{ba}	0		
$ heta_{bb}$	1/3		
θ_{b}	2/3		

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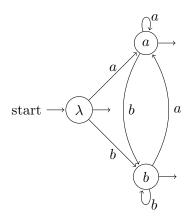
Example: Strictly 2-Local Languages



Parameters
$\overline{ heta_{ times a}}$
$ heta_{ times b}$
$\theta_{\rtimes \ltimes}$
$\overline{\theta_{aa}}$
$ heta_{ab}$
$ heta_{a \ltimes}$
θ_{ba}
$ heta_{bb}$
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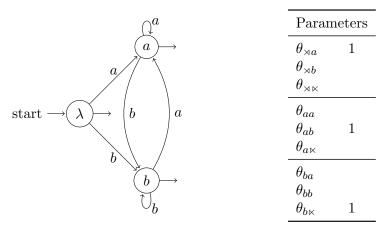
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$ heta_{ times oldsymbol{ee}}$
θ_{aa}
$ heta_{ab}$
$\theta_{a \ltimes}$
θ_{ba}
$ heta_{bb}$
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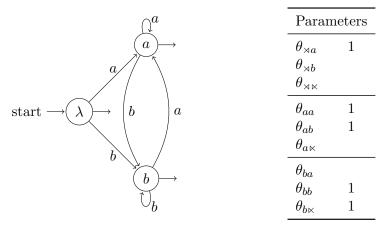
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Smallest language consistent with D in C is obtained by passing D through DFA and 'activating' parsed transitions. (Heinz and Rogers 2013)

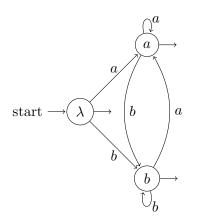
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θ_{\bowtie}	0	
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Class C defined with

	single DFA	finitely many DFA
$f: \Sigma^* \to \{0, 1\}$ $f: \Sigma^* \to [0, 1]$		

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Overview of Related Results (part 2)

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- 6 Maximization-Expectation techniques are used to learn the class of PNFAs, but there is no guarantee to find a global $\underset{\text{U. Toronto}}{\text{optimum}} \ (\text{Rabiner 1989}).$

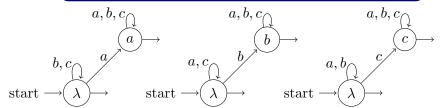
Defining C with finitely many DFA

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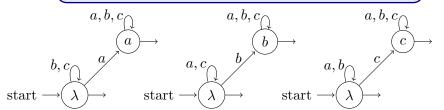
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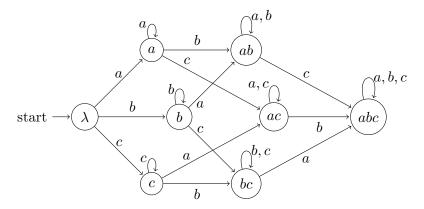
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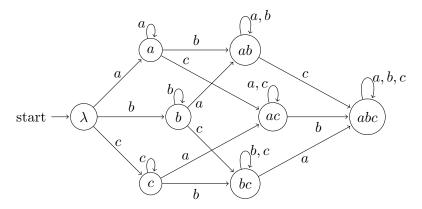
Product Operations

- 1 For Boolean languages, use **acceptor product** (yields intersection)
- 2 For Stochastic languages, use co-emission product (yields joint distribution)

THE PRODUCT OF THOSE THREE ACCEPTORS



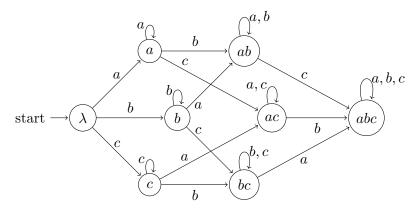
THE PRODUCT OF THOSE THREE ACCEPTORS



(exit/accepting arrow at each state is not shown)

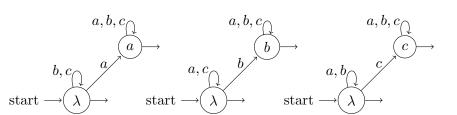
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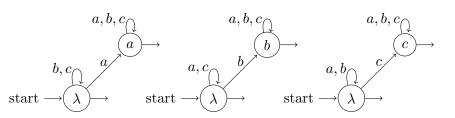


- If C is defined by this DFA, then C = Piecewise 2-Testable.
- If C is defined by the 3 atomic DFAs, then C = Strictly 2-Piecewise.

Cause ...

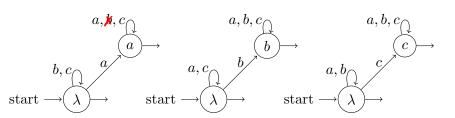


Cause ...

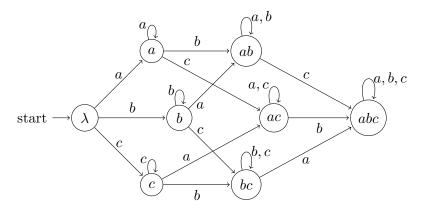


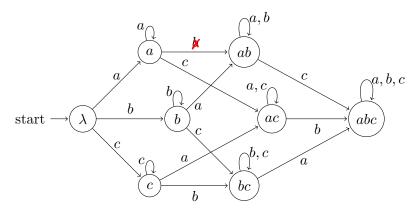
The parameters of the model are set at the level of the individual DFA.

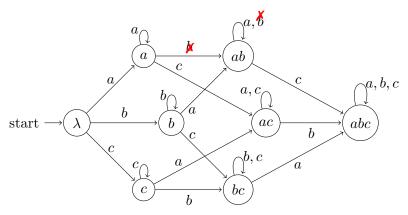
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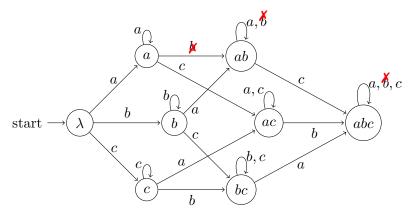


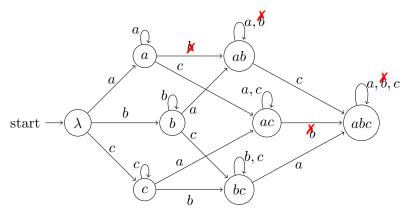
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- 2 The transitions in the product are NOT independent.



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- They show it always returns the smallest Boolean language in the class consistent with the data, and thus identifies the class in the limit from positive data.

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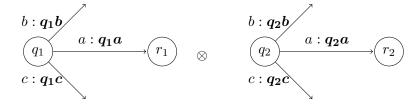
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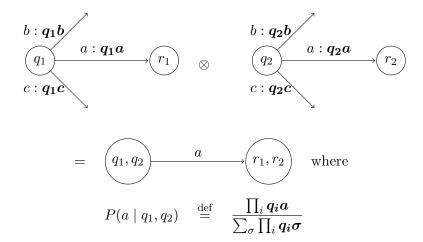
		X			
$Pr(x \mid P_{\leq 1}(y))$		s	$\widehat{\mathrm{ts}}$	ſ	$\widehat{\mathrm{tf}}$
	s	0.0335	0.0051	0.0011	0.0002
	$\widehat{\mathrm{ts}}$	0.0218	0.0113	0.0009	0.
У	ſ	0.0009	0.	0.0671	0.0353
	$\widehat{\mathrm{tf}}$	0.0006	0.	0.0455	0.0313

Table: Results of SP₂ estimation on the Samala corpus. Only sibilants are shown. (Heinz and Rogers 2010, p. 894)

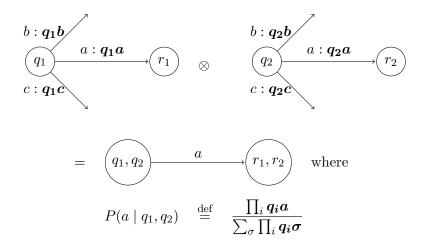
THE CO-EMISSION PRODUCT (DEFINITION)



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For fixed σ , the co-emission product treats the parameters $q_i\sigma$ as independent.

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- 2 Essentially, we prove that parameters which maximize the probability of the data with respect to such models are found by running the corpus through each of the individual factor PDFAs and calculating the relative frequencies.

SOME DETAILS OF THE ANALYSIS

- 1 Probability of Words
- 2 Relative Frequency of Emissions
- 3 Empirical Mean of co-emission probabilities
- 4 Main Theorems

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- Then the probability that σ is emitted after the product machine $\bigotimes_{1 \leq i \leq K} \mathcal{M}_i$ reads the prefix $\sigma_1 \cdots \sigma_{i-1}$ is the following:

$$\operatorname{Coemit}(\sigma, i) = \frac{\prod_{j=1}^{K} T_j(q(j, i), \sigma)}{\sum_{\sigma' \in \Sigma} \prod_{j=1}^{K} T_j(q(j, i), \sigma')}. \tag{1}$$

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• We assume that there is a end marker $\ltimes \in \Sigma$ which uniquely occurs at the end of words.

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- Let q(j,i) denote a state in Q_j that is reached after M_j reads the prefix $\sigma_1 \cdots \sigma_{i-1}$.
- If i = 1 then q(j, i) represents the initial state of M_j .
- Let $T_j(q, \sigma)$ denote a parameter (transitional probabability) in PDFA M_j .
- Then the probability that σ is emitted after the product machine $\bigotimes_{1 \leq j \leq K} \mathcal{M}_j$ reads the prefix $\sigma_1 \cdots \sigma_{i-1}$ is the following:

$$\mathsf{Coemit}(\sigma, i) = \frac{\prod_{j=1}^{K} T_j(q(j, i), \sigma)}{\sum_{\sigma' \in \Sigma} \prod_{j=1}^{K} T_j(q(j, i), \sigma').} \tag{1}$$

• We assume that there is a end marker $\ltimes \in \Sigma$ which uniquely occurs at the end of words.

$$P(w \ltimes) = \prod_{i=1}^{N+1} \mathsf{Coemit}(\sigma_i, i)$$

• Let $m_w(M_i, q, \sigma) \in \mathbb{Z}^+$ denote how many times σ is emitted at the state q while the machine M_i emits w.

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Relative Frequency of Emission

- Let $m_w(M_i, q, \sigma) \in \mathbb{Z}^+$ denote how many times σ is emitted at the state q while the machine M_i emits w.
- Let $n_w(M_i, q) \in \mathbb{Z}^+$ denote how many times the state q is visited while the machine M_i emits w.

Then

$$freq_w(\sigma|M_j, q) = \frac{m_w(M_j, q, \sigma)}{n_w(M_j, q)},$$
(3)

represents the relative frequency that M_i emits σ at q during emission of w.

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It is straightforward to lift this definition to data sequences $D = \langle w_1 \ltimes, w_2 \ltimes, \dots w_{|D|} \ltimes \rangle$ by letting $w = w_1 \ltimes w_2 \ltimes \dots w_{|D|} \ltimes$.

EMPIRICAL MEAN OF CO-EMISSION PROBABILITIES

Empirical Mean of Co-emission probabilities

$$\mathsf{sumCoemit}_w(\sigma, M_j, q) = \sum_{i \text{ s.t. } q(j, i) = q} \mathsf{Coemit}(\sigma, i).$$

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The *empirical mean of a co-emission probability* is defined as follows:

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This is the sample average of the co-emission probability when $q \in Q_i$ is visited.

Main Theorem

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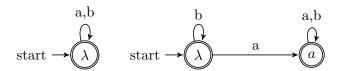
Consider any parameter $T_j(q, \sigma)$ in PDFA M_j .

Theorem

 $\partial P(D)/\partial T_j(q,\sigma)=0$ holds for all j if and only if the following equation is satisfied for all $1\leq j\leq K$:

$$\operatorname{freq}_w(\sigma|M_j,q) = \overline{\operatorname{Coemit}}_w(\sigma|M_j,q)$$
.

EXAMPLE



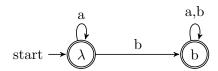


FIGURE: The 2-set of of SD-PDFAs with $\Sigma = \{a, b\}$. There are 15 parameters. Suppose $D = abb \ltimes bbb \ltimes$.

EXAMPLE

$$\begin{split} \operatorname{freq}_D(a|\mathcal{M}_\lambda,\lambda) &= 1/8 & \operatorname{freq}_D(a|\mathcal{M}_a,\lambda) = 1/5 & \operatorname{freq}_D(a|\mathcal{M}_a,a) = 0/3, \\ \operatorname{freq}_D(b|\mathcal{M}_\lambda,\lambda) &= 5/8 & \operatorname{freq}_D(b|\mathcal{M}_a,\lambda) = 3/5 & \operatorname{freq}_D(b|\mathcal{M}_a,a) = 2/3, \\ \operatorname{freq}_D(\ltimes|\mathcal{M}_\lambda,\lambda) &= 2/8 & \operatorname{freq}_D(\ltimes|\mathcal{M}_a,\lambda) = 1/5 & \operatorname{freq}_D(\kappa|\mathcal{M}_a,a) = 1/3, \\ & \operatorname{freq}_D(a|\mathcal{M}_b,\lambda) = 1/3 & \operatorname{freq}_D(a|\mathcal{M}_b,b) = 3/5, \\ & \operatorname{freq}_D(b|\mathcal{M}_b,\lambda) = 2/3 & \operatorname{freq}_D(b|\mathcal{M}_b,b) = 0/5, \\ & \operatorname{freq}_D(\kappa|\mathcal{M}_b,\lambda) = 0/3 & \operatorname{freq}_D(\kappa|\mathcal{M}_b,b) = 2/5, \\ \end{split}$$

Figure: Frequency computations with $D = abb \ltimes bbb \ltimes$ and the 2-set of of SD-PDFAs on previous slide.

Let $\tau_{j,q,\sigma}$ denote $\log T_j(q,\sigma)$; i.e. the log of a parameter of C defined with $\bigotimes_{i} M_{j}$.

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Theorem

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Thus the solution obtained by the previous theorem is a MLE.

At a high level, the problem we considered is a decomposition of complex probability distributions into simpler factors.

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A reviewer points out that this literature may simplify our proofs.

FUTURE WORK

- 1 Language Modeling with various sets of specific factors and various corpora such as . . .
 - $1 \operatorname{SL}_k + \operatorname{SP}_k$
 - 2 $SLP_{k,\ell}$ (Rogers and Lambert 2019, MoL)
 - 3 Atomic PDFA based on phonological features (Chandlee et al. 2019, MoL)
- 2 ... and compare to NNs, ALERGIA, and other algorithms on various benchmarks.
- 3 Connections to probababilistic graphical models
- 4 Extend results to weighted deterministic automata.

THANKS

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Questions?