Finite-state morphology/phonology tutorial, UD, Jan 23, 2013

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Resources

 The Finite-State Morphology Book [Draft version] (Beesley and Karttunen, 2003)

https://victorio.uit.no/langtech/tags/DIVVUN-NO 1 0 RELEASE/gt/doc/book.pdf 1.pdf

• Foma: program, documentation, etc.

http://foma.googlecode.com

Morphological analysis tutorial (with foma)

https://code.google.com/p/foma/wiki/MorphologicalAnalysisTutorial

Tutorial slides

https://foma.googlecode.com/files/lrec2010 slides.tar.gz

Foma frequently used commands

Starting foma

foma

foma -l <filename> Start foma and execute script on startup

Interface commands

help X Get help on X

apropos X Short help on X

clear Clear stack

define X Y; define X with regular expression Y (Ex.: define V [a|e|i|o|u])

down Enter apply down mode (CTRL-D exits this mode)

net Print top FSM with text

pop Pop the top FSM off of stack

up Enter apply up mode (CTRL-D exits this mode)

random-lower Print random selection of output words from top FSM

random-upper Print random selection of input words from top FSM

regex X; Compile regular expression X

test equivalent Test top two FSMs for equality

tseq Test top FSM for sequentiality

view View top FSM visually

words Print all words of top FSM (or subset if cyclic)

Regular expressions

A | B Union

A* Kleene Star

A+ Kleene Plus

A B Concatenation

? Any symbol

0 Epsilon symbol

\$A "Contains" A (equiv. to ?* A ?*)

A - B Subtraction

A & B Intersection

~A Complement of A

(A) "Optionally" A: equivalent to A | 0

[] Grouping brackets

A .o. B Composition

A.r Reversal

A -> B Unconditional rewriting

A (->) B	Optional rewriting

$A \rightarrow B \parallel L R$	Context-restricted	rewriting (.	\\. //. \\ `) are all valid conte	xt specifiers
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[..] -> B Epenthesis rules are specified like this (do not use
$$0 \rightarrow B$$
)

.#. Word-boundary symbol in context specifications

A.u Extract input projection of transducer A

A.1 Extract output projection of transducer A

A < B All instances of A precede B

A > B All instances of A follow B

A => B _ C Context restriction: the language where all instances of A occur only

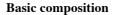
between C and D

A -> B, C -> D, ... Multiple simultanous rewrites with same context

 $A \rightarrow B \parallel C \perp D$,, ... Multiple simultanous rewrites with different contexts

"^" Use quotes to escape special characters

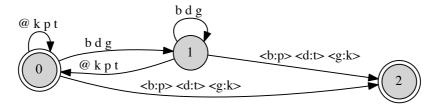
Some examples covered



regex a -> b .o. b -> c ;



End of word stop devoicing:



Spreading with // rules vs. regular ||-type rules:

Illustration of differences between rule types $(\parallel, //, \parallel, \vee)$

a b -> x \\ a b _ a [L holds on input side, R on output side]

abababa (INPUT)
abab x a (OUTPUT)

a b -> x // a b _ a [L holds on output side, R on input side]
a b a b a b a (INPUT)
a b x a b a (OUTPUT)

 ${\tt a} {\tt b} {\tt -> x} \ {\tt //} {\tt a} {\tt b} {\tt _a}$ [L holds on output side, R on output side, gives two valid pairings]

abababa ababa (INPUT)
ab x aba abab x a (OUTPUT)

Sibilant harmony example

#We condition the sibilant harmony on the sibilant immediately to $\#the\ left$, with possibly intervening nonsibilants, encoded by #[?-SIB]*

```
define SIB [s|\[];
define Rule1 s -> \[ || \[ [?-SIB]* \_ ;
define Rule2 s -> \[ // \[ [?-SIB]* \];
```

• With Rule1, we get no spreading:

```
foma[0]: regex Rule1;

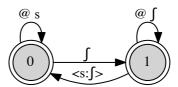
423 bytes. 2 states, 6 arcs, Cyclic.

foma[1]: down

apply down> fisisi

fifisi

apply down>
```



• With Rule2, spreading occurs, since we condition on output side

```
foma[0]: regex Rule2;

423 bytes. 2 states, 6 arcs, Cyclic.

foma[1]: down

apply down> fisisi

fififi
```

