

# ‘modal logic’

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October 28, 2022

*Math. Struct. in Comp. Science* (2001), vol. 11, pp. 511–540. Printed in the United Kingdom  
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## **A judgmental reconstruction of modal logic**

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*Received 3 December 1999; revised 3 May 2000*

# Outline

- 1 Pre paper
  - Pre Pre 1960
  - Pre 1960
  - Kripke Semantics
- 2 Pfenning, Davies
  - Judgment, in passing

- prior

- hilbert proofs

## Definition

$$\mathcal{M} = \langle \mathcal{W} : Set, R : \mathcal{W} \times \mathcal{W} \rightarrow 2, V : \omega \rightarrow \mathcal{W} \rightarrow 2 \rangle$$

## Definition

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} P_n \iff V(n)(\alpha)$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} \perp \iff \text{false}$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} A \wedge B \iff \models_{\alpha}^{\mathcal{M}} A \text{ and } \models_{\alpha}^{\mathcal{M}} B$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} A \vee B \iff \models_{\alpha}^{\mathcal{M}} A \text{ or } \models_{\alpha}^{\mathcal{M}} B$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} A \rightarrow B \iff \models_{\alpha}^{\mathcal{M}} A \implies \models_{\alpha}^{\mathcal{M}} B$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} \Box A \iff \forall \beta \in \mathcal{W}. \alpha R \beta \implies \models_{\beta}^{\mathcal{W}} A$$

$$\models_{\alpha \in \mathcal{W}}^{\mathcal{M}} \Diamond A \iff \exists \beta \in \mathcal{W}. \alpha R \beta \wedge \models_{\beta}^{\mathcal{W}} A$$

## Definition

$$\mathcal{M} = \langle \mathcal{W} : Set, \leq, V : \omega \rightarrow \mathcal{W} \rightarrow 2 \rangle$$

$$\alpha \leq \beta \implies \models_{\alpha}^{\mathcal{M}} P_n \implies \models_{\beta}^{\mathcal{M}} P_n$$

## Definition

$$\models_{\alpha}^{\mathcal{M}} A \rightarrow B \iff \forall \beta. \alpha \leq \beta \implies \models_{\beta}^{\mathcal{M}} A \implies \models_{\beta}^{\mathcal{M}} B$$

$$\Box(A \rightarrow B)$$

# Judgment

- “Of course, there would be needed here an analysis of what is understood by an expression, but that is a comparatively trivial matter, as compared with explaining the notion of proposition and judgement. An expression in the most general sense of the word is nothing but a form, that is, something that we can passively recognize as the same in its manifold occurrences and actively reproduce in many copies. But I think that I shall have to rely here upon an agreement that we have such a general notion of expression, which is formal in character, so that the rule can now count as a formal rule.”



# Local Soundness/Completeness

# References I

[1] TODO: ask me