



Reidemeister Moves on Alternating Knots and their Signed Planar Graphs

Internship Report

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1 Introduction

“We learn the rope of life by untying its knots.”

— Jean Toomer

In the nineteenth century, Lord Kelvin proposed the *vortex theory of the atom* to try and explain why chemists found only a few different types of atoms, but there are a lot of each type. The theory stated “atoms were knotted vortices in the ether” [3]. This resulted in efforts to try and find different kinds of knots. And although, Lord Kelvin’s theory was wrong, *knot theory* thrived anyway. Many tabulations of knots showed up, mathematicians defined terms such as *achiral*, *amphicheiral*, *prime*, *alternating* to describe different knots.

The question of equivalence lies of course in the core of many of these terms, it is natural to come to the question; when are two knots equivalent? And what seems to be even a simpler question turns out to be not as straight forward; how to know if a knot is not the unknot? A knot as a structure is very unique, two knots could be equal although they look nothing alike. So what is equal in the eyes of knot theory? Knot theorists started to provide such definitions of equivalence and knot invariants.

Although the knot lives in three dimensional space, we study it in two dimensional space. And Translating the properties of the knot to 2D space is done by knot diagrams and other tools. The definition of equivalences of course transitions (although with some care!) to the 2D world.

Another mathematical structure that lives in 2D is planar graphs. Insights regarding the connection between knots and planar graphs were investigated. A representation of knots as planar graphs helped in tabulating knots for example[3]. It is also possible to derive the Jones polynomial, a knot invariant, from the representative planar graphs of knots[4]. An exciting connection exists between those two mathematical concepts, and graph theory has been around at least a century more before any keen mathematical investigations in knots. It is surely beneficial for both fields to aim to develop, create tools for and investigate this connection[3].

The focus of this project was to dapple further with the relation between planar graphs and knots. As we stated, the question of whether two knots are the same is a unique question. Under special constraints, certain actions on the knot does not change the knot. We will discuss such actions in more details in our report. For now, one of these actions fall under the name of *Reidemeister* moves. It has been shown[3][4][7] that these *Reidemeister* moves have an equivalent correspondence when carried to the planar graphs’ world.

In this report, we present a program that takes in as input a knot diagram represented in *Gauss* Notation, and outputs the related planar graphs. The program also performs the aforementioned *Reidemeister* moves on the planar graphs.

The rest of report is divided as follows; Section 2 will discuss related background and theory in order to help us build an overview of our objective. Section 3, we state our problem statement and our end goal for the program. Section 4 is the Methodology, we highlight our software tools, class structure, external sources used and the Organization of our program. We then dive into the program’s code, we highlight the different aspects and roles of each class, how every piece of the program works with the rest. Section 5, we provide the results and also highlight two examples for better clarity on the program’s procedure. In section 6, we discuss our program regarding performance and efficiency, we also highlight major draw backs. we then continue to provide what paths and aspects should be regarded in future work and what protocols might highly impact the success of future projects regarding our problem statement. Section 7, we conclude our report, summarizing everything so far. Lastly, we have provided an Appendix that contains all of our results so far, a tabulation of the knot diagrams the program worked on and the entire code of the program.

2 Background & Theoretical Concepts

“If I have seen further than others, it is by standing upon the shoulders of giants.”

— Isaac Newton

2.1 The Knot in \mathbb{R}^3

There are many formal definitions for a knot, we choose the following one[2] for its ease to build on further concepts.

Definition 1. A *knot* is a simple closed polygonal curve in \mathbb{R}^3 .



Figure 1: The trefoil as a smooth curve and a polygonal curve.

This definition gives us the benefit to exclude such notions as *wild knots*[2][5].

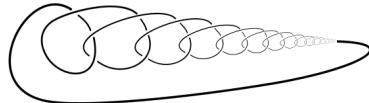


Figure 2: A wild wild knot.

We would then like to define what is a *vertex* of a knot.

Definition 2. Let the ordered set of points (p_1, p_2, \dots, p_n) define a knot, and no proper ordered subset defines the same knot, then the points of the set $\{p_i\}$ are called the **vertices** of the knot.

It is only natural to then utilize the notion of vertices of a knot and bask in its merits, and we do so in the next definition.

Definition 3. Let k be a knot defined by the sequence of points (p_1, p_2, \dots, p_n) and let j be a knot defined by the sequence $(p_0, p_1, p_2, \dots, p_n)$, where:

1. The point p_0 is not collinear with p_1 and p_n .
2. The triangle spanned by (p_0, p_1, p_n) intersects the knot k only in the segment $[p_1, p_n]$.

Then, the knot j is called an **elementary deformation** of the knot k .

But we are not done yet, we can now state the upcoming powerful result[2].

Definition 4. For any two knots k and j , if there exists a sequence of knots $k = k_0, k_1, \dots, k_n = j$, where for each $k_i + 1$ an elementary deformation of k_i , for $i < 0$, then the knots k and j are **equivalent**.

Since a knot is considered a closed polygonal curve, it has no starting point and no end point. Therefore, we can assign an orientation to it[5]. In this report, we will use relay on the notion of orientation in our program. So let's delve into it.

Definition 5. An **oriented** knot consists of a knot and an ordering of its vertices. The ordering must determine the original knot. Two orderings are considered equal if they differ by a cyclic permutation.

The equivalence relation between knots is also carried through to oriented knots.

Definition 6. Let there be a sequence of elementary deformations changing one oriented knot to another, then these two knots are called **oriented equivalent**.

The necessity of displaying knots on paper, on boards and in books led to the formalization of knot projections and diagrams in the plane, or in other words, \mathbb{R}^2 .

2.2 Knot Projections and Knot Diagrams in \mathbb{R}^2

Let k be a knot, then we denote the projection of k by $p(k) = \hat{k}$. If k has an orientation then naturally the projection \hat{k} inherits the orientation. Note that k lives in \mathbb{R}^3 space while \hat{k} lives in \mathbb{R}^2 [4][5]. Where k crosses under or over itself, there will exist an intersection in its projection \hat{k} . We would like now to reference the next definition[5].

Definition 7. Let \hat{k} have the following properties:

1. \hat{k} has finite points of intersection.
2. If Q is a point of intersection of \hat{k} , then the inverse image of Q in k has exactly 2 points.
3. A vertex in the knot k is never mapped onto a double point in \hat{k} .

Then, we call k' a regular projection.

Figure 3 show a point that does not abide by property number 2. Similarly, figures 4 and 5 show points that do not abide by property number 3.

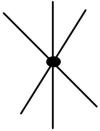


Figure 3: Violation of Property 2.

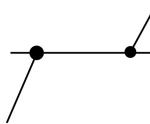


Figure 4: Violation of property 3.

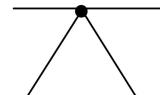
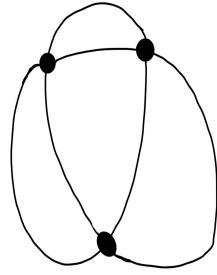
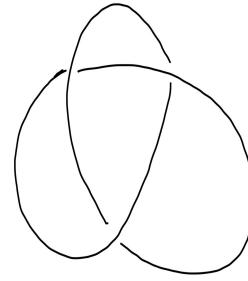


Figure 5: Violation of property 3.

In order to incorporate the “under-and-over” crossing information in the projection, we draw the projection to *appear* as if it cuts itself. And we call this altered projection a *regular diagram*[5]. From here on out, we will denote the diagram of the knot k as k' (e.g. if the knot is called k_1 then its diagram is k'_1).



(a) A knot projection.



(b) A knot diagram.

Figure 6: The trefoil as a projection and as a diagram.

Definition 8. A regular diagram k' with no **reducible** crossings is called a **reduced regular diagram**.

For example, figures 7 and 8 are reduced diagrams while figure 9 is not.

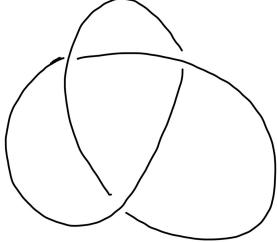


Figure 7: The trefoil.

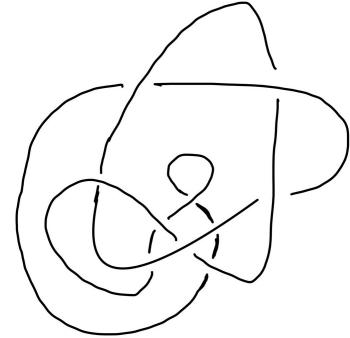
Figure 8: 7_1 knot.

Figure 9: The star with reducible crossings.

Definition 9. An alternating knot diagram is a knot diagram that has the crossings alternate between over and under as we go any orientation[3].

Definition 10. An **alternating knot** is a knot that has at least one alternating diagram.

2.2.1 Equivalence with respect to Knot Diagrams and Projections

We would like to highlight the importance of knot projections and diagrams. After all, through our report we will be relying so much on them in our program.

In [2], there are two theorems regarding projections that set up the field for us. The first theorem states that even if a knot does not have a regular projection, there is an equivalent *nearby* knot with a regular projection. The second theorem states that if a knot has a regular projection, then all *nearby* knots also have a regular projection. The term *nearby* relates to the distance between vertices as we make elementary deformations. We now come to one of major result we wanted to introduce[2].

Theorem 1. Let k and j be knots that have regular projections and identical diagrams, then k and j are **equivalent**.

2.3 Reidemeister moves

Reidemeister moves are ways to alter a knot diagram that will change the relation between the crossings[3]. There are three Reidemeister moves, each paired with its inverse. The *first Reidemeis-*

ter move allows us to add or take a twist away. The second Reidemeister move allows to add or remove two crossings, under certain conditions. The third Reidemeister move allows us to slide a segment of the knot from one side of a crossing to the other end. Figure ?? depict those moves.

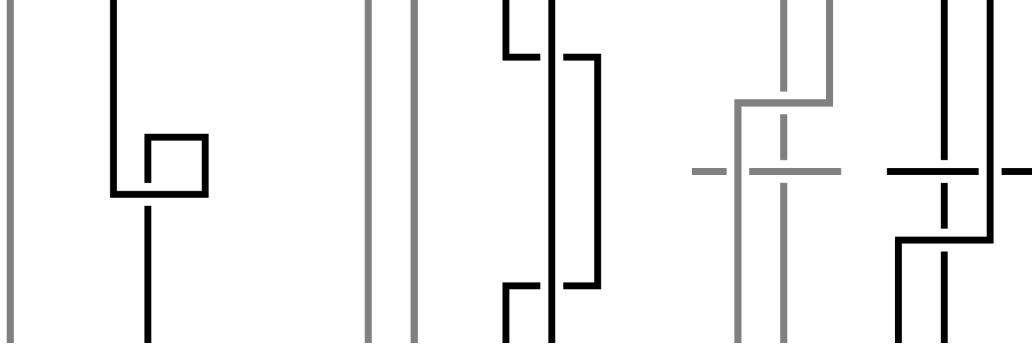


Figure 10: *RI* aka “a twist”. Figure 11: *RII* aka “a poke”. Figure 12: *RIII* aka “a slide”.

Figure 13: The Reidemeister Moves

It is not unreasonable to think that the alterations done by these Reidemeister moves on the knot diagram will result in representing a totally different knot. However, we have the following observation by Alexander and Briggs[2]:

Theorem 2. *If Two knots are equivalent, then their diagrams are related by a sequence of Reidemeister moves.*

We also include the next theorem from [11]:

Theorem 3. *Two knot diagrams are equivalent, if and only if, they are related by Reidemeister moves.*

We will sometimes refer to Reidemeister moves by their nicknames; a “twist” for the first, a “poke” for the second and a “slide” for the third.

2.4 Notation for Knots

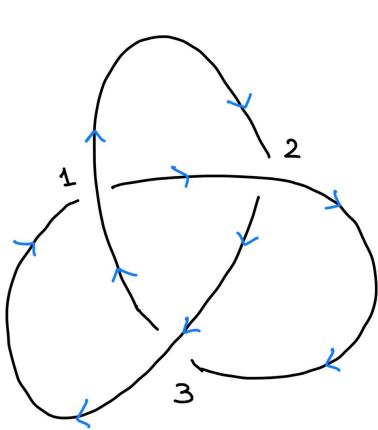
In the process of tabulating knots, codes and notations have been created to represent them in a compact and concise manner. We would like to discuss one specific notation as it relates closely to our program.

Gauss Code

It is amazing that even the great Gauss dabbled with knot theory. The notation is straight forward and has some similarities with another notation, Dowker–Thistlethwaite notation.

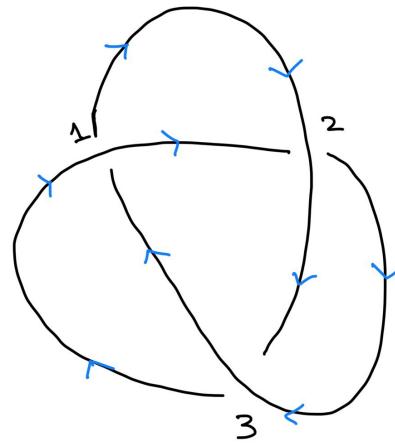
If we have some knot diagram, we give it an orientation, we then pick an arbitrary point and walk along the strands of the knot, when we meet a crossing for the first time we label it, starting with 1 and so on. We then start again at the same point, and if we are passing the knot at over-crossing, we give the label a positive sign, if it was an under-crossing we give the label a negative sign[14].

This method has the downside that it does not differentiate between a knot and its mirror if they are different. For example the Gauss code for the trefoil and the left trefoil is the same: $1, -2, 3, -1, 2, -3$ [14].



Gauss Code: 1, -2, 3, -1, 2, -3

(a) The trefoil.



Gauss Code: 1, -2, 3, -1, 2, -3

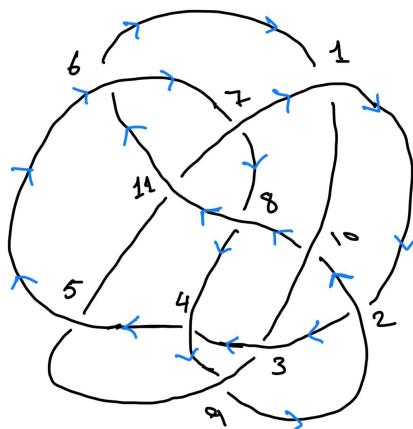
(b) Its mirror.

To address this, the *extended Gauss code* was proposed which accounts for the crossings local orientation. To define local orientation, we define first the difference between right-handedness and left-handedness[14]:

- A right-handed crossing means that if our thumb points in the direction of the orientation of the over-strand then the rest of our finger will point in the direction of the orientation in the under-strand.
- A left-handed crossing means the same thing but using our left hand.

And now, the *extended Gauss code* changes the Gauss code as follows[14]:

We do the same thing in our first iteration and label the crossing. However, in the second iteration over the knot the first occurrence of the label has a sign to depict whether it was an over or under crossing and the second occurrence has a sign to depict the local orientation of the crossings. If the local orientation is right-handed, we give the label a positive sign, if it is left-handed, we give it a negative sign.



Gauss Code:

1 -2 3 -4 5 6 -7 -8 4 -9 2
10 8 11 -6 -1 10 -3 9 -5 -11 7

Extended Gauss Code:

1 -2 3 -4 5 6 -7 -8 -4 -9
-2 -10 8 11 6 -1 -10 3 9 -5
-11 -7

Figure 15: A knot and its Gauss Codes. The red numbers indicate the difference between the Extended Gauss Code and the normal notation.

2.5 Knots and their Signed Planar Graphs

Most of the work our program does, it will be doing it with the aforementioned *signed planar graphs*. In the upcoming subsection, we will show how they are constructed, their equivalence with their respective knots and more.

The Knot Projection and The Plane Graph

Let's again reconsider the notion of the knot projection. Assume we have a knot k , its projection \hat{k} and its diagram k' . If k has at least one crossing, then its projection can be described as a 4-regular plane graph, where each crossing is a vertex. The diagram k' can also be described by the same 4-regular plane graph, but we will need to account for the information of the over and under crossing at each vertex[4].

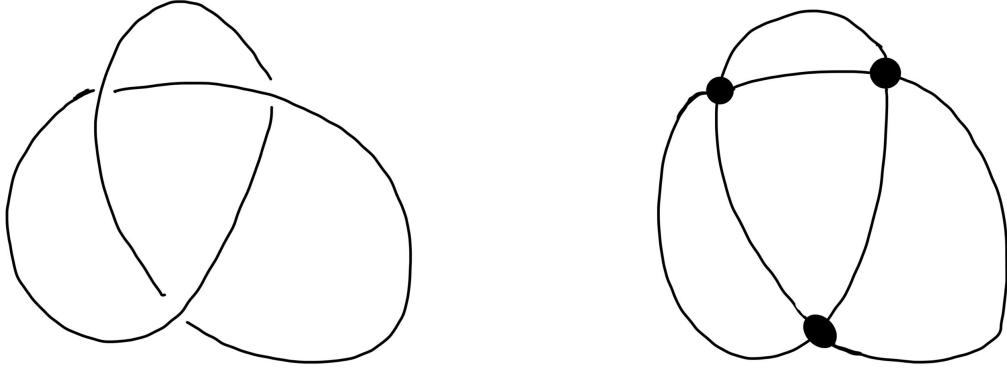


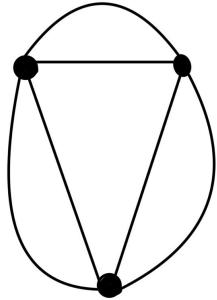
Figure 16: The trefoil and its *plane graph*.

From the Plane Graph to the Signed Planar Graphs

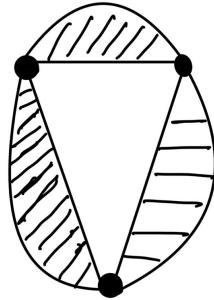
In the context of the following, we can consider \hat{k} as the shadow of k' . Since the knot projection \hat{k} is a 4-regular plane graph, then its dual graph is bipartite[4].

We now can colour the faces of the shadow of the knot projection, or in other words the faces of the plane graph, using two colours, usually they are black and white. We start by colouring the exterior face white, and now each face that share a boundary with the exterior face, we colour black, and then each uncoloured face that shares a boundary with a black face we colour white, and so on we alternate until we have coloured all faces. For its respective colour, each face now represent a vertex for that graph, and we draw an edge between the vertices for each vertex the faces shared in the plane graph. In the case that a vertex of the plane graph lies only in one coloured face, then we draw loop at the respective vertex that represents that face.

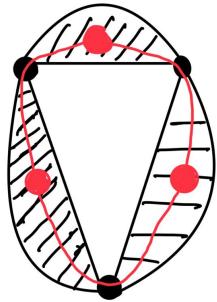
To sum up, we have two planar graphs, one black and one white. let's denote the black graph with b and the white graph with w . We would like to highlight again up to this point what do the graphs actually represent. Let's take b as an example, every vertex in b represents a black face, and every edge of b represents a crossing that these faces share. Figure 17 help to visualise the process.



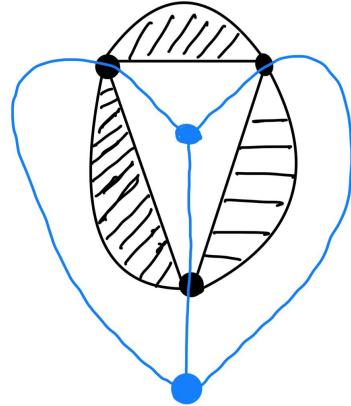
(a) The plane graph for the trefoil.



(b) Shaded plane graph.



(c) Red representing the black graph.



(d) Blue representing the white graph.

Figure 17: Construction of the planar graphs.

The graphs b and w are always connected graphs. Moreover, b and w are the graph duals of each other and from either b or w , we can go backwards and construct the plane graph (the knot projection \hat{k} , the shadow of the knot diagram)[4].

What remains now to fully reflect the original knot is the under-and-over information. In order to represent this in the graph, we go to each crossing and we imagine the over-strand rotating *anticlockwise* until it lies on the under-strand, we take note of the face the over-strand traversed¹ and keep track of the face's colour. Back to the graphs b and w , we go to each edge and since each edge symbols a crossing, we check the colour we took note of previously, if it is the same as the graph, the edge gets a positive sign, if it is opposite, then the edge gets a negative sign.

¹The over-strand passes simultaneously through two faces as it rotates. However, they will always be of the same colour (figure 25 gives a visual representation).

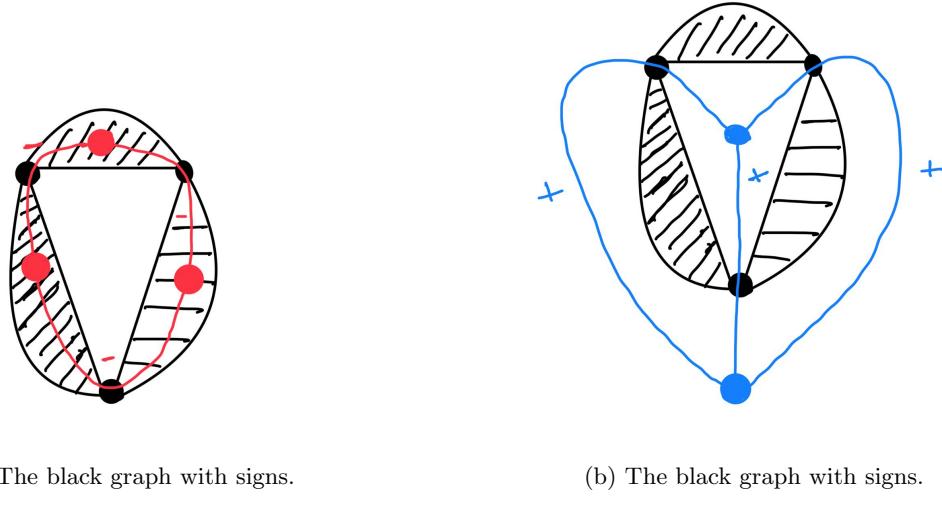


Figure 18: The planar signed graphs for the trefoil.

After adding the signs, we now call b and w the *signed planar graphs* of the knot.

We note that from the *signed planar graphs*², we can obtain the original knot diagram. The process is described in [3] and [5].

2.6 Reidemeister Moves on the Signed Planar Graphs

We carry the same concepts of Reidemeister moves to *signed planar graphs*.

The *first Reidemeister move* translates to deleting a loop in one planar graph and a deleting spike in the other.

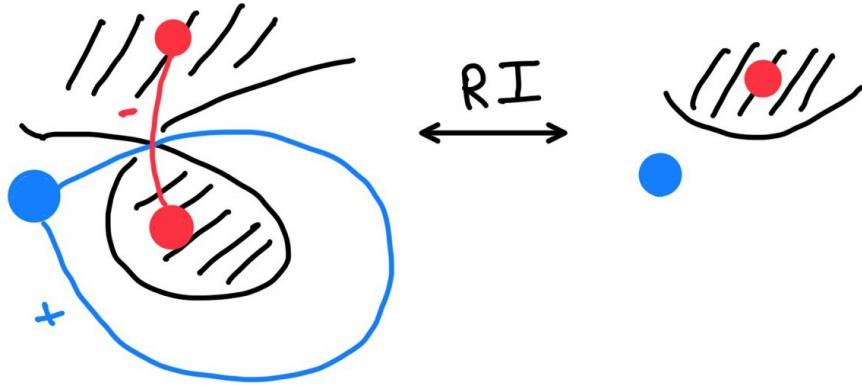


Figure 19: The first Reidemeister move.

²Actually, from only just one!!!

The *second Reidemeister move* translates to deleting two parallel edges in one graph of different signs and contracting a 2-path (with edges of different signs) in the other graph.

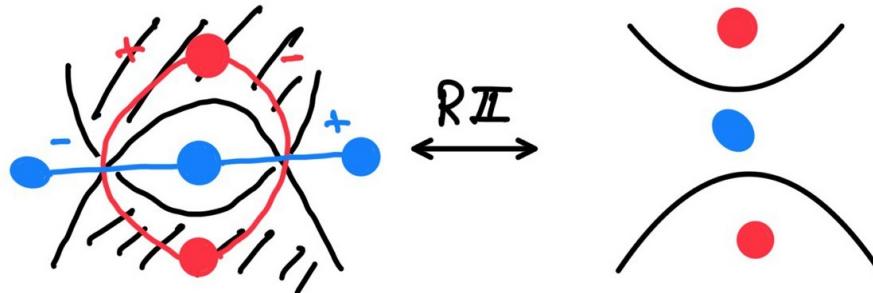


Figure 20: The second Reidemeister move.

The *third Reidemeister move* translates to changing a cycle C_3 in one graph to a star $K_{1,3}$ and doing the reverse in the other.

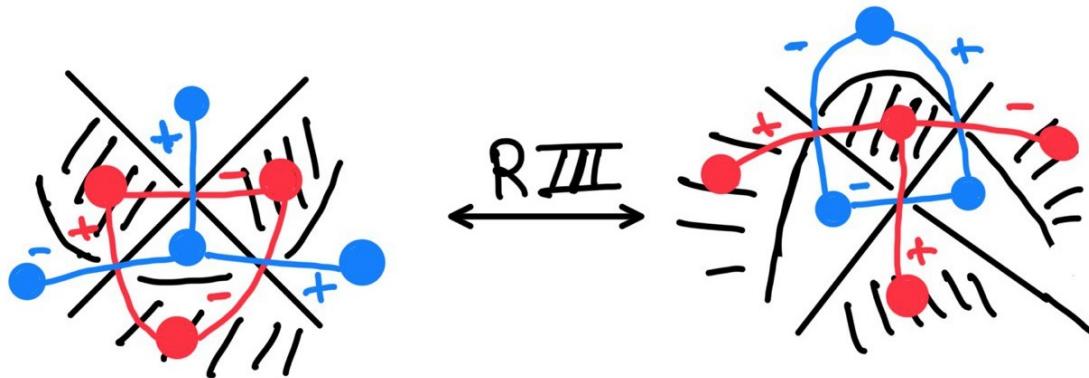


Figure 21: The third Reidemeister move.

More technical detail will be shown in the Methodology(4). The Reidemeister moves here are shown on a more of a local level and we will delve in the details of how our programs handles each *Reidemeister move* and how it affects the whole graph.

3 Problem Statement

“No problem can withstand the assault of sustained thinking.”

— Voltaire

The main inspiration for our internship is the following:

- Let k be a knot and k' be its projection.
- We perform a sequence of *Reidemeister moves* on k' to obtain the knot diagram j' .
- Let b_k, w_k, b_j, w_j be the signed planar graphs of k' and j' respectively.

Now, can we do the same on graphs b_k, w_k to obtain b_j, w_j ?

The question as it stands is hard³, so we have reformulated our starting point.

let k be a knot, k' its knot diagram and j' its **reduced regular diagram**. Let (R_1, R_2, \dots, R_n) be a sequence of Reidemeister moves taking j' to k' , $(k'_0, k'_1, \dots, k'_n)$ be a sequence of knot diagrams, where $i \in \{1, \dots, n\}$, $k'_0 = j'$, $k'_n = k'$ and k'_i is the knot obtained from applying Reidemeister move R_i on k_{i-1} . we constrain that $\forall j, l \in \{1, \dots, n\}$:

- k'_j does not have more crossings than k'_l , for $j < l$.
- $\#R_j, R_l$ such that R_j is a *third Reidemeister move* and R_l is its inverse.

These restrictions are mainly to avoid the situation where we have to perform a *third Reidemeister move* and increase the crossings in order to reduce the knot diagram. An example of such can be found in [3].

And now with the above notation, let b_k, w_k be the signed planar graphs of k' , our question becomes:

Can we find a sequence of Reidemeister moves on b_k, w_k to obtain the signed graphs b_j, w_j of $j'?$

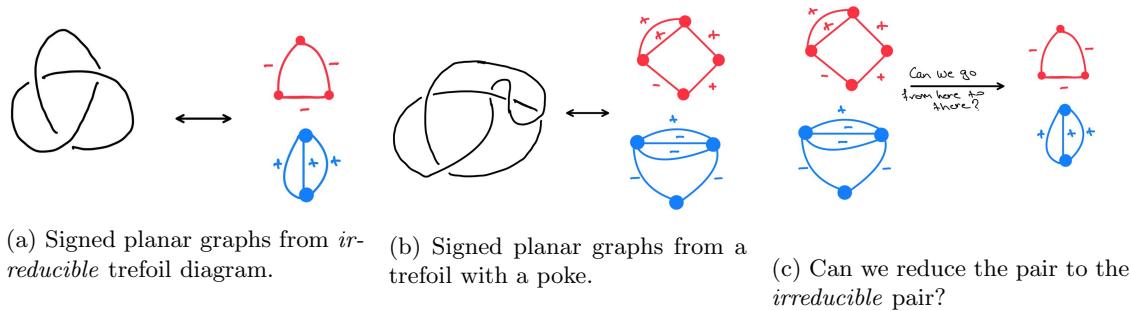


Figure 22: A visual overview of our problem statement.

³Assuming we know nothing else except the planar graphs, it is difficult to find a suitable sequence of Reidemeister moves.

4 Methodology

“It makes no difference how they are loosed.”

— Alexander the Great,
referring to the Gordian Knot.

We have laid out the background, given our motivation and stated our problem. Now, it's time to get to work!

We would like to begin first by detailing the organization, class structure and different software tools.

4.1 Software and Technicalities

4.1.1 Programing Language and Main External Packages

we would like to state the overall structure of our project and the different classes. However, to avoid ambiguity, we will state first the different tools used in our program.

Our program is written in *Python* and nothing else. In order to achieve our objectives we utilized the following two libraries:

- NetworkX.
- pyknotid.

These two packages have been essential to our program. NetworkX[6] deals with graphs, while pyknotid[8] deals with knots and related concepts. More can be found about these two libraries in their respective docs.

Library	Link for Docs
NetworkX	https://networkx.github.io/
pyknotid	https://pyknotid.readthedocs.io/en/latest/

Table 1: Python Libraries and Documentation Links

When we go through the code, we will highlight the crucial methods and classes that we utilized from these sources.

4.1.2 Class Structure and Organization

In our program, there are two classes, one utility class, one helper module, two main scripts and one test script. let's go through them quickly.

HardCodedKnots In this class, we have encoded Gauss codes to represent different knot diagrams. The main functionality of this class is to call instances of it and pass it to the **SignedPlanarGraph** class.

ReidemeisterMoves The name says it all, this *utility* class is responsible for performing the *Reidemeister moves* on the signed planar graphs.

SignedPlanarGraph This class is the main player. It is responsible for taking a Gauss Code representing a knot diagram and turning it into a pair of signed planar graphs.

main This is one of the main scripts, where we could lauch the whole program.

main_iterative The other main script, the main difference is that it is set up such that the program will run over all knot diagrams encoded in **HardCodedKnots**.

check_gauss_code This test script simply takes a Gauss code and identifies the knot up to reflection. It is useful in checking results, debugging and verifying inputted Gauss codes to the program.

my_helpers A module with many methods we use throughout the code.

File Name	Type	Description
HardCodedKnots	class	Encoded Gauss codes.
ReidemeisterMoves	class	Performing Reidemeister moves.
SignedPlanarGraphs	class	Signed Planar Graph construction and manipulation.
check_gauss_code	test script	verifying input and output through Gauss codes.
main	main script	Launch the program.
main_iterative	main script	Launch the program on everything we got!
my_helpers	module	helper methods used throughout the program.

Table 2: Brief overview of different files.

```

project/
├── .vscode/
│   ├── launch.json
│   └── settings.json
└── pyknotid_helali/
    └── my_code_1/
        ├── images
        ├── check_gauss_code.py
        ├── main.py
        ├── main_iterative.py
        ├── my_helpers.py
        ├── HardCodedKnots.py
        ├── ReidemeisterMoves.py
        └── SignedPlanarGraph.py

```

Figure 23: Organization of Project

4.1.3 Major Modules from Imported Libraries

In the upcoming, we will state the major functionalities we used. More details is available in the package respective docs.

NetworkX

Graph This is a class to represent an un-directed graph structure. It allows loops, but no parallel edges are allowed.

MultiGraph This is a class to represent an un-directed graph structure. It allows both loops and parallel edges. This is the main class we will relay to represent our *signed planar graphs*.

PlanarEmbedding This is a class that represents a planar graph with its embedding.

is_isomorphic This is a method that takes two graphs and checks if they are isomorphic.

check_planarity Checks if a graph is planar and returns a embedding for it.

PlanarEmbedding.traverse_face This method is called on an instance of PlanarEmbedding. The method takes an edge and returns the nodes of the face that belong to the edge.

cycle_basis This method takes a graph and returns a list of cycles.

drawing.nx_graph.to_agraph This method is used to draw graphs and export to images.

pyknotid

catalouge pyknotid has a prebuilt databases for knots up to 15 crossings. This package provides a knot lookup by name and by some invariants. However, the download of the database is required (This is provided in pyknotid's docs).

spacecurves This module includes multiple classes to perform topological analysis on knots and work with them in \mathbb{R}^3 .

representations This module provides different ways to represent knots from invariants, diagrams and notations.

representations.gausscode.GaussCode This is a Class that represents a knot as a Gauss code. We will provide more details momentarily([4.1.3](#)) as the notation here is slightly different from the usual theoretical notation and from the one we introduced in our background section([2.4](#)).

catalogue.identify.get_knot This method takes a string and returns a an object of type “ $\langle \text{Model: Knot} \rangle$ ”. Examples of strings that can be passed are “3_1” for the trefoil and “4_1” for the figure eight knot.

catalogue.database.space_curve This method takes the object returned from get_knot and returns a Knot Object. The method space_curve is defined in the module mentioned before “space-curves”.

representation.space_curve This method creates a Knot object from an representation. It also relays on the module “spacecurves”.

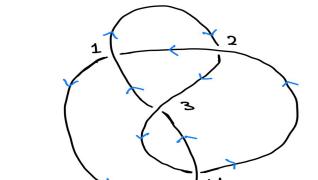
In our code, it will become clear how we utilize the different functionalities of pyknotid. Although we mentioned many modules and classes, pyknotid have them connected and integrated in a way that using pyknotid does not actually require deep analysis of the library itself, and going from one type of representation and/or class becomes straightforward after some experimentation.

Gauss Code in pyknotid

The major concepts are the same as the extended Gauss code we provided in our background section. However, pyknotid uses a different kind of notation. As we traverse the knot, each occurrence of a crossing is represented by three characters; the label number, “-” or “+” to indicate “under-and-over” information, and “a” or “c” to indicate local orientation. And the Gauss code is made up of these three characters separated by a comma. What follows elaborates more.

Symbol	Description
+	An over-crossing.
-	An under-crossing.
a	Anticlockwise/right-handed local orientation.
c	Clockwise/left-handed local orientation.

Table 3: Description of Gauss code components in pyknotid.



Pyknotid's Gauss code:
 $1+a, 2-a, 3+c, 4-c$
 $2+a, 1-a, 4+c, 3-c$.

Figure 24: an example of pyknotid's Gauss code notation.

For the diagram of the figure eight shown, with the given labeling and given orientation, the Gauss code according to pyknotid will be “1+a,2-a,3+c,4-c,2+a,1-a,4+c,3-c”.

Other Packages

We would be remiss if we don't mention the following packages/libraries.

copy Used to create copies of objects.

collections Used through the construction process of the *signed planar graphs*

planarity Through NetworkX, we use crucial modules from planarity. They are related to planar graphs and their embeddings[20].

4.2 The Algorithm and the Program

4.2.1 Overview of the Algorithm

Keeping things straight forward, we start by stating the steps of our approach:

1. We have a knot k .
2. We create two knot diagrams k' and j' . Keeping to our notation, j' is the reduced diagram.
3. k' has been obtained from a sequence of Reidemeister moves made on j' , abiding by the constraints we stated in the Problem Statement(3).
4. We calculate the Gauss code for both knot diagrams.
5. We input the Gauss codes to the program.
6. We obtain the signed graphs b_k, w_k, b_j, w_j .
7. We input b_k, w_k to the program.
8. The program performs all valid Reidemeister moves possible and outputs the signed graphs \hat{b}_k, \hat{w}_k .

9. We compare b_j , w_j and \hat{b}_k , \hat{w}_k .

Simple enough!

4.2.2 The Program

main

Our strategy in this subsection is to start with the main script and as we go through it, we will cover all aspects of our program. And without further ado, here it is:

```

1  from SignedPlanarGraph import SignedPlanarGraph
2  from HardCodedKnots import HardCodedKnots
3  from ReidemeisterMoves import ReidemeisterMoves
4
5  hcg = HardCodedKnots(code_no=3113)
6
7  k_gauss = hcg.get_code_gauss()
8
9  sp_graph = SignedPlanarGraph(code_gauss = k_gauss)
10 sp_graph.create_and_draw_sp()
11
12 reid_sp = ReidemeisterMoves.reidemeister_main(sp_graph)
13
14 reid_sp = SignedPlanarGraph(code_gauss=reid_sp.code_gauss)
15 reid_sp.create_and_draw_sp()
16
17 #print(reid_sp.code_gauss)

```

The first three lines are simply importing our three classes we mentioned before. let's look at line number 5 with more detail.

```
5  hcg = HardCodedKnots(code_no=3113)
```

In line 5, we create the variable `hcg` and we call the constructor of the class **HardCodedKnots**. In the constructor we pass an identifier to the knot diagram, here it is the variable `code_no` and the value is “3113”. We will highlight what does this value mean in the **HardCodedKnots** class. In line 6, we pass the Gauss code to the variable `k_gauss`.

```
9  sp_graph = SignedPlanarGraph(code_gauss = k_gauss)
10 sp_graph.create_and_draw_sp()
```

In line 9, we create an instance of the class **SignedPlanarGraph** and this is saved in the variable `sp_graph`. In line 10, we call the method `create_and_draw_sp()`. This method takes the knot diagram as a Gauss code, produces the two signed planar graphs saved as attributes in the instance `sp_graph`. The method also produces four graph images, two graphs for black and two graphs for white. for each, One graph with no sign labels and the other is with.

```
13  reid_sp = ReidemeisterMoves.reidemeister_main(sp_graph)
```

In line 13, we use the utility class **ReidemeisterMoves**. The method `reidemeister_main()` takes as input the signed planar graph instance `sp_graph` and produces the graph `reid_sp` after performing all possible Reidemeister Moves.

```
15  reid_sp = SignedPlanarGraph(code_gauss=reid_sp.code_gauss)
16  reid_sp.create_and_draw_sp()
```

In line 15, we recreate the instance `reid_sp` this to ensure that all attributes have been cleared and set correctly. In line 16, we again produce the two signed graphs and export the graph images.

HardCodedKnots

```

1  from pyknotid.catalogue import get_knot
2  import pyknotid.representations as rep_
3
4  class HardCodedKnots:
5
6      code_numbers = [31, 311, 3111, 312, 3113, 41, 411, 4112, 412,
7          51, 513, 5131, 52, 522, 61, 62, 63,
8          71, 72, 73, 74, 75, 76, 77]

```

In the beginning of the file we import the modules relates to pyknotid. At line 4, the class begins. The class has a global attribute `code_numbers`, we will highlight there use later. Let's move forward to the `__init__` function.

```

10 def __init__(self, code_no = None, code_string = None):
11
12     if code_no == 31 or code_string == "trefoil" or code_string
13         == "3_1":
14         k = get_knot('3_1').space_curve()
15         self.code_gauss = k.gauss_code()
16         return
17     if code_no == 311 or code_string == "trefoil_twist" or
18         code_string == "3_1_1":
19         self.code_gauss = rep_.GaussCode("1+c,2-c,3+c,1-c,2+c,4+
20             c,4-c,3-c")
21         return

```

the `__init__` function can take two parameters, `code_no` or `code_string`. We see that we have `if` blocks, one after the other. There are multiple `if` blocks after the one we shown here, In each `if` block a knot diagram is represented by a Gauss Code. Let's zoom in on the first one of them.

```

12     if code_no == 31 or code_string == "trefoil" or code_string
13         == "3_1":
14         k = get_knot('3_1').space_curve()
15         self.code_gauss = k.gauss_code()

```

The parameters that could be passed to the `__init__` function where `code_no` and/or `code_string`. So this `if` block checks if `code_no` is equal to 31 or `code_string` is equal to “trefoil” or to “3_1”. If one of these conditions hold, the program enters the `if` block.

In line 13, we use pyknotid's `get_knot()` method to obtain the trefoil and we use pyknotid's `space_curve()` to flesh out the knot and calculate certain invariants and notations that were not there yet just from the `get_knot()` method. This information is now saved in the variable `k`. Line 14 sets the attribute `code_gauss` of the instance of the class **HardCodedKnots** to the **Gauss Code** related to `k`. Line 15 simply exits the `__init__` function.

We hope it was clear what was the overall objective of the `__init__` function. recall that in “main” we had the following:

```

5 hcg = HardCodedKnots(code_no=3113)

```

if `code_no` was equal to 31, `hcg` would be initialized with the attribute `code_gauss` set to the Gauss code of the trefoil.

All reduced regular diagrams of prime knots up to seven crossings are set up as the trefoil. Here is another example for the knot 6_3 .

```

80     if code_no == 63 or code_string == "6_3":
81         k = get_knot("6_3").space_curve()
82         self.code_gauss = k.gauss_code()
83         return

```

let's check the `if` block for the `code_no` of 3113.

```

23     if (code_no == 3113 or code_string == "trefoil_twist_slide"
24         or code_string == "3_1_1_3"):
25         self.code_gauss = rep_.GaussCode("1+c,2-c,3+c,4+c,5+a,6-
a,7-c,3-c,8-a,5-a,6+a,1-c,2+c,7+c,4-c,8+a")

```

In this case, we rely on `pyknotid`'s module “representation”. In line 30, we use it to create a representation of the knot using the Gauss code notation of `pyknotid`.

We would like also to point out, that there is a certain convention we followed while creating the different `code_no` and `code_string` for different Gauss codes. for `code_no`, The name of the knot is followed by numbers representing the Reidemeister moves we performed to obtain the knot. For example, from `code_no` 3113, we deduce that this is the knot 3_1 and we performed a first Reidemeister move then a third Reidemeister move. For `code_string`, we have two options; if the knot has a famous name, we chain the name and then the Reidemeister moves performed, separated by underscore. If there is no famous name, we simply chain the numbers in `code_no` also separated by underscore. That's why in the previous `if` block, “trefoil_twist_slide” and “3_1_1_3” show up.

In table (4), there is a full tabulation showing each knot diagram, its relevant Gauss code, `code_no` and `code_string`.

And to the final detail in `HardCodedKnots`, the getter method `get_code_gauss()`.

```

119     def get_code_gauss(self):
120         return self.code_gauss

```

SignedPlanarGraph

In the land of **SignedPlanarGraph**, knot diagrams find themselves transformed, gaining magical powers to shape-shift from diagrams to graphs and back again.

let's start by stating the imports and global attributes.

```

1 import networkx as nx
2 import my_helpers as help
3 import copy
4
5
6 class SignedPlanarGraph:
7
8     drawing_number = 0

```

NetworkX makes its entrance, we also import other relevant packages and our helpers. the global attribute `drawing_number` will be used when we start producing images for the signed planar graphs.

The following can be passed as parameters to the `__init__` method.

- `black_graph_data`.

- *white_graph_data*.
- *knot*.
- *code_gauss*.
- *signed_black*.
- *signed_white*.

We would like to point out that our approach changed thorough out the project, and that in the program only *code_gauss* will be passed as a parameter.

The class **SignedPlanarGraph** has the following attributes:

- *black_graph*; this has data regarding the black graphs' nodes and edges.
- *white_graph*; The same as above but for the white graph.
- *black_loops*; Loops need special attention and this attribute is responsible to keep track of loops in the black graph.
- *white_loops*; Same as before but for the white graph.
- *code_gauss*; This is initial Gauss code of the knot diagram.
- *signed_black*; this attributes saved the signing labels for the edges of the black graph.
- *signed_white*; the same thing but for the white graph.

```

9  sp_graph = SignedPlanarGraph(code_gauss = k_gauss)
10 sp_graph.create_and_draw_sp()

```

From the main script, in line 9 we created an instance of **SignedPlanarGraph** and saved it in the variable *sp_graph*. As stated we only passed the Gauss code saved in *hcg*. So as things stand, the instance of **SignedPlanarGraph** has the Gauss code saved in the attribute *code_gauss*. Now let's direct our attention to line 10 in the main script and the method *create_and_draw_sp()*.

```

547     def create_and_draw_sp(self):
548         graph = help.create_planar_graph_from_gauss_code(self.
549             code_gauss)
550         self.create_nodes_from_faces(graph)
551         self.create_edges_after_nodes()
552         self.sign_sp()
553         self.draw_sp_graph()
554         self.draw_sp_graph_sign()

```

create_planar_graph_from_gauss_code() takes a Gauss code and returns a graph. Although, the name of the method has “planar” in it, we rely on NetworkX and planarity to handle the creation of a planar embedding.

create_nodes_from_faces() is responsible for the creation of the vertices/nodes of the signed planar graphs.

create_edges_after_nodes() is responsible for adding edges with respect to the crossings shared between faces.

sign_sp() adds the signs on the edges.

`draw_sp_graph()` draws the planar graphs and exports them as images.

`draw_sp_graph_sign()` draws the graphs with the signs on the edges.

Nodes/Vertices Creation from the Faces of the Shadow

The method `create_nodes_from_faces()` takes variable the plane graph (The shadow of the knot diagram) as input. The following is what occurs;

`nx.check_planarity()` This method creates an instance of **PlanarEmbedding** from the passed variable `graph`. The method is defined in NetworkX.

`get_faces_list()` This method takes an instance of **PlanarEmbedding** and produces the faces of the plane graph. we utilize the method `PlanarEmbedding.traverse_face()` to create the faces. However, this method does not account for loops or parallel edges.

`add_loops_face_list()` This methods adds the faces with respect to loops and parallel edges.

`get_longest_region()` This method chooses the exterior face based on length.

`get_outer_regions()` This method starts colouring the exterior face as white and its neighbours as black.

`shade_faces()` This methods completes the colouring of the faces and returns two lists of faces one for each colour.

We then iterate over the list of faces, add the nodes to their respective graph, i.e. black faces are added to the attribute `black_graph` and in respective fashion for `white_graph`. We also save loops in either `black_loops` or `white_loops` depending on which graph they belong too.

Edge Creation

The method `create_edges_after_nodes()` uses the new initialized attributes to get the job done for edges. for each graph we utilize the node name to check if two nodes share a crossing or not. We also utilize the attributes `black_loops` and `white_loops` to account for loops in the graphs.

Sigining the Edges From the Gauss Code

We now have the nodes and edges for each graph, what remains is to sign the edges and our construction is complete. This is the work of the method `sign_sp()`. We rely on the attribute `code_gauss` to help us sign each edge. We also utilize the attributes `black_loops` and `white_loops` as loops create corner cases for signing.

For each edge in one of the signed graph, we check what crossing it represented, then using the Gauss Code, we obtain the face that the over-strand and under-strand bound. Then, depending on the local orientation of the crossing, the colour of the bounded face and which graph we are signing, we know what sign to give to the edge.

The information is saved in `signed_black` for the black graph and `signed_white` for the white.

The following figures display different scenarios for signing.

Assuming that
we are considering
the black graph

Legend:

- △ face with same colour as graph
- △ face with different colour

↷ Rotating over strand
anticlockwise

↗ Orientation

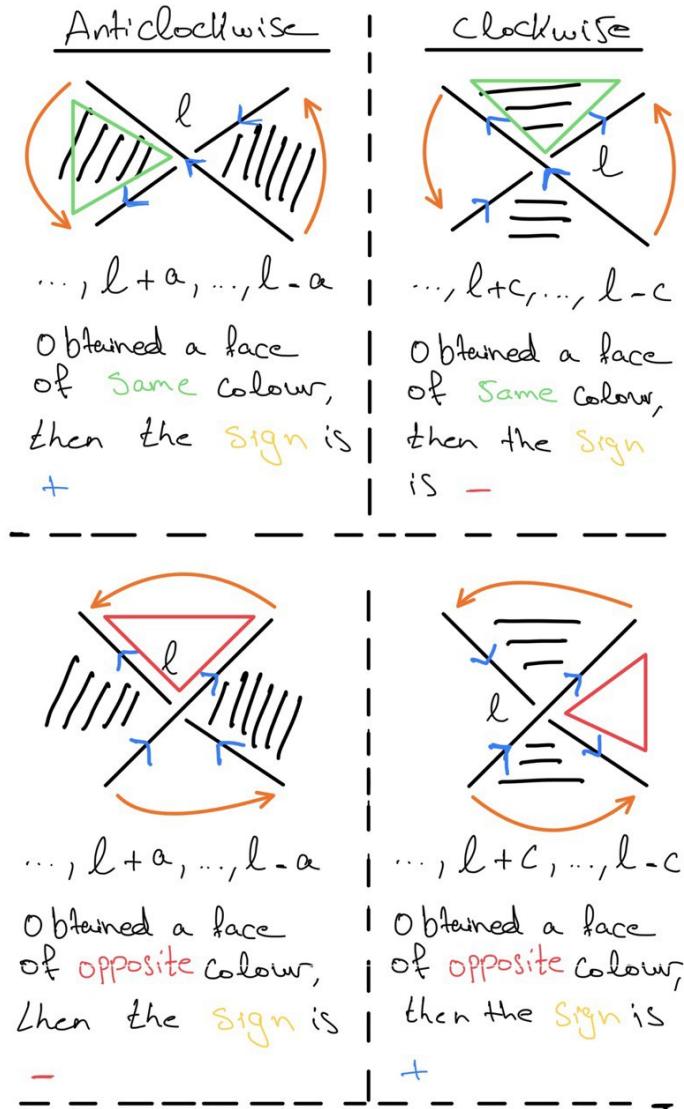


Figure 25: Procedure of how we sign.

Drawing and Exporting the Signed Planar Graphs

What we thought would be one of the main challenges in this program, turned out to be an easy task thanks to the external libraries. The methods `draw_sp_graph()` and `draw_sp_graph.sign()` perform this responsibility. From NetworkX, We utilize the methods `to_agraph()` and `draw()` for producing images.

`draw_sp_graph()` draws the graphs without signs. Then, we add the signs from `signed_black` and `signed_white` and `draw_sp_graph_sign()` draws them as labels on the edges.

Methods for ReidemeisterMoves

We would like to highlight some methods in this class that are used in the utility class **ReidemeisterMoves**.

`create_and_sign_sp()` In the process of performing Reidemeister moves, we obtain new signed graphs, this method ensures that all attributes of the new signed graphs are set correctly.

find_cycles_of_length() This method is used for the *third Reidemeister move*, where we find cycles of length three in one of the graphs. We utilize NetworkX's method *cycle_basis()*.

ReidemeisterMoves

The class **ReidemeisterMoves** is a utility class, which is to say it contains only methods and we never create instances of it.

In the main script, we had the following.

```

13  reid_sp = ReidemeisterMoves.reidemeister_main(sp_graph)
14
15  reid_sp = SignedPlanarGraph(code_gauss=reid_sp.code_gauss)
16  reid_sp.create_and_draw_sp()

```

Performing All valid Reidemeister moves

The method *reidemeister_main()* is responsible for performing all possible Reidemeister moves. It contains the following methods.

reidemeister_I() This method performs Reidemeister moves of the first type.

reidemeister_II() This method performs Reidemeister moves of the second type.

reidemeister_III() This method performs Reidemeister moves of the third type.

reidemeister_main() contains a while loop with the above methods in order. It keeps track of whether any of the above methods changed our signed planar graphs. When an iteration of the while loop occurs and no changes were performed on the signed planar graphs, it breaks out of the loop and returns the resulted signed planar graphs.

Reidemeister I

The method *reidemeister_I()* is responsible to perform the first Reidemeister move. In similar fashion as *reidemeister_main()*, it keeps performing the *first Reidemeister move* until no more changes occur. It contains the following.

reidemeister_I_helper() This method contains two smaller helper methods:

```

reidemeister_I_helper_black()
reidemeister_I_helper_white()

```

Each one of the smaller helper method is related to one of the signed planar graph. let's consider *reidemeister_I_helper_black()*, as *reidemeister_I_helper_white()* does the exact same steps but for white. *reidemeister_I_helper_black()* starts by searching for loops in the black graph. If we find one, we check if the white graph has a related spike node. If both conditions are met, we modify the Gauss Code to perform the Reidemeister move, and recreate the signed planar graphs.

Reidemeister II

The method *reidemeister_II()* is responsible to perform the second Reidemeister move. It also keeps performing the *second Reidemeister move* until no more changes occur. It contains the following.

reidemeister_II_helper() This method contains two smaller helper methods:

```
reidemeister_II_helper_black()
reidemeister_II_helper_white()
```

Each one of the smaller helper method is related to one of the signed planar graph, we again consider just one of them. *reidemeister_II_helper_black()* starts by saving all available pairs of parallel edges. It then filters them keeping only pairs of edges with different signs. Then it searches for the related 2-path edges in the white graph. If all conditions are met, we again modify the Gauss code and create the new signed graphs.

Reidemeister III

The method *reidemeister_III()* is for the *third Reidemeister move*. It performs the *third Reidemeister move* until no more changes occur. It contains the following.

reidemeister_III_helper() This method contains two smaller helper methods:

```
reidemeister_III_helper_black()
reidemeister_III_helper_white()
```

Each one of the smaller helper method is related to one of the signed planar graph. for *reidemeister_III_helper_black()*, we use the method *get_cycles_of_length()* to obtain cycles of length three. We iterate over the cycles, checking if one of them have a related star graph in the white graph. If the condition is met for one of the cycles. We modify the Gauss code and recreate the signed planar graphs.

my_helpers

Let's highlight the methods in our helper module.

is_cyclic_list() Used to check whether two lists are equivalent with respect to cyclic order. This method is mainly used to check that we do not have duplicate faces When creating signed planar graphs.

sublistExists() Used to check if a list is a subset of the other, this is used in edge creation.

edges_set_creation() Used in the colouring of the faces, it creates the edges that bounded a face in the plane graph (shadow of the knot diagram).

get_bounded_face() Used in signing to obtain the face bounded by the over-strand and the under-strand at a crossing.

is_subset_both_ways() Used in finding star graphs when performing *third Reidemeister moves*.

create_planar_graph_from_gauss_code() creates a plane graph from the Gauss code. Parallel edges and loops are omitted from the graph.

main_iterartive

Performs the same utility of the main script, with the added feature that it iterates over the global attribute *code_numbers* in the class **HardCodedKnots**.

check_gauss_code

For convenience and testing, we provided a module that identifies the Gauss code associated with a pair of signed planar graphs.

identify() It takes a Gauss code and identifies it, giving a knot tabulation (e.g. “4_1”).

Additional Notes

The whole code is provided in the Appendix([C](#)) with comments.

5 Results & Examples

“Mathematics is the art of giving the same name to different things.”

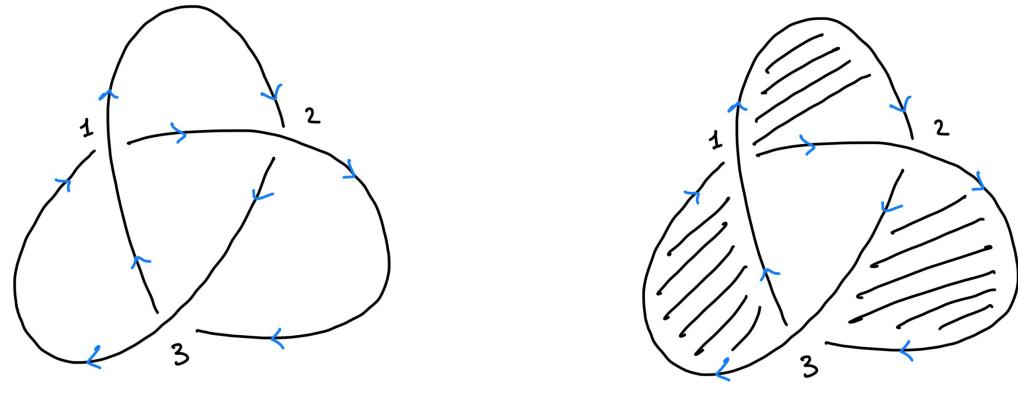
— Henri Poincaré

Our program worked with 24 knot diagrams, constructing their signed planar graphs, executing valid Reidemeister moves and comparing. We obtained the expected results from our program.

We will provide two examples to display the outputs of the program. And we start with the easiest one.

5.1 Example 1: The Trefoil

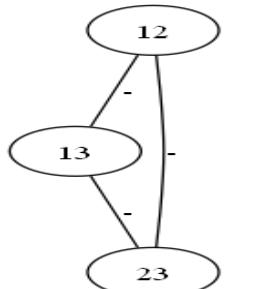
We start with the Trefoil, the Gauss code is: “ $1+c, 2-c, 3+c, 1-c, 2+c, 3-c$ ”. The next figure is its knot diagram and our own colouring work.



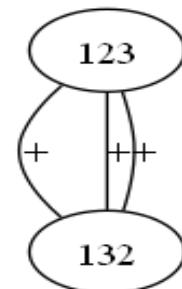
Gauss Code (pyKnotid notation):
 $1+c, 2-c, 3+c, 1-c, 2+c, 3-c$

Figure 26: Our example trefoil with its Gauss code and shaded.

What follows is the signed planar graphs construction.



(a) Black Graph.



(b) White Graph.

Figure 27: The Signed Planar Graphs for the Trefoil.

We like to note that the above figure 26 is a *reduced* knot diagram. And we obtained its signed planar graphs in figure 27. In what follows, we will have a knot diagram with *reducible* crossings, obtain its signed planar graphs and reduce and compare our results.

5.2 Example 2: The Trefoil With three valid Reidemeister moves

Let's move to the case where there are three valid Reidemeister moves, one for each type. Here is the knot diagram and its respective Gauss Code: “1+c,2-c,3-c,6a,7-a,8-c,6+a,9+c,2+c,0+c,0-c,4-c,5-a,1-c,8+c,7+a,9-c,3+c,4+c,5+a”

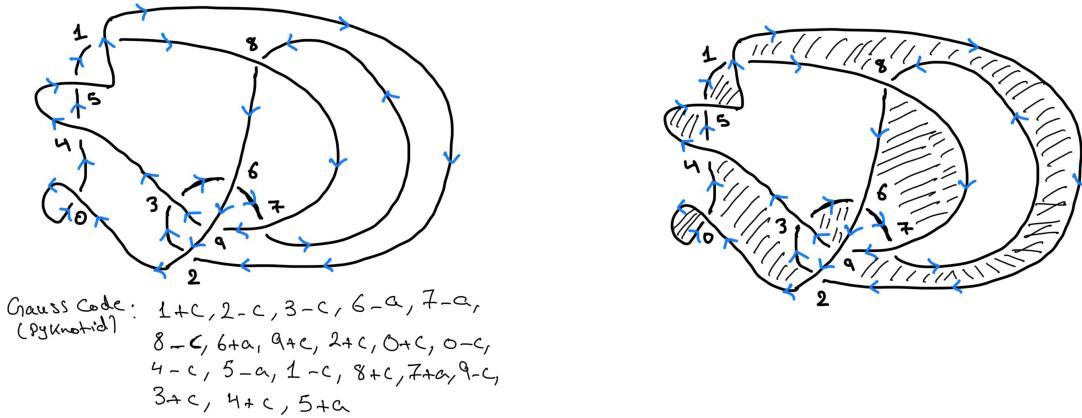
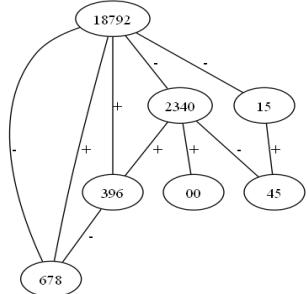
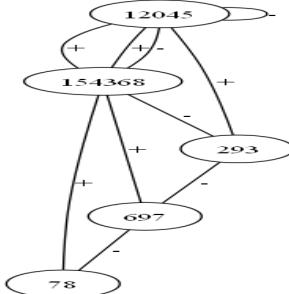


Figure 28: Our second example, the trefoil with three valid Reidemeister moves.

And now we construct the signed planar graphs.



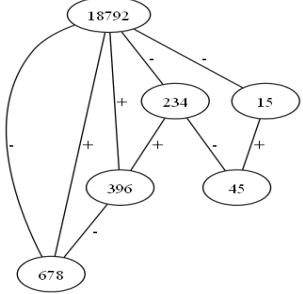
(a) The Black Graph.



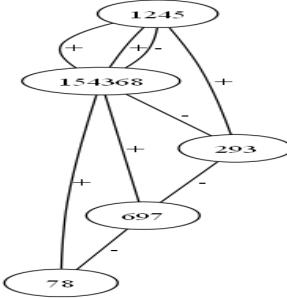
(b) The White Graph.

Figure 29: The Signed Planar Graphs for the Trefoil with three valid moves.

If we look closely at the figures above 29, we notice that in the black graph there is a spike node, namely “00”. And in the white graph there is a loop in the node “12045”. That pair of nodes constitute a Reidemeister move of type one. And our program (thankfully!!) performs the Reidemeister move and reduces the graph.



(a) Black Graph after performing RI.



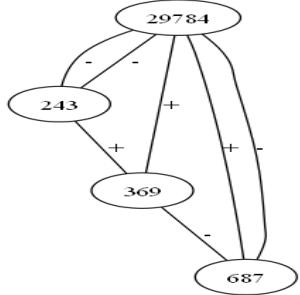
(b) The White Graph after Performing RI

Figure 30: The Signed Planar Graphs after RI.

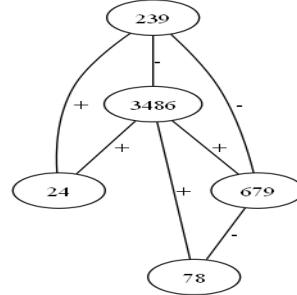
The above graphs are the result after performing one Reidemeister move of type one. We observe that both the spike and the loop have vanished. The new name nodes reflect that the crossing “0” is no longer there. The current Gauss Code is: “1+c,2-c,3-c,6-a,7-a,8-c,6+a,9+c,2+c,4-c,5-a,1-c,8+c,7+a,9-c,3+c,4+c,5+a”.

Since, there are no more valid Reidemeister moves of type one. The program then searches for a valid Reidemeister move of type two. We were surprised (albeit pleasantly!) that the program choose the crossings “1” and “5” to perform the second Reidemeister move instead of our intention of the crossings of “4” and “5”. Still, this is a valid move. The next figures 31 show the resulting operation. If we look closely at the black graph, we will find that the 2-path from node “18792” to “15” to “45” has been replaces by an edge between the nodes “29784” and “243”. In the white graph, a pair of parallel edges between “1245” and “154368” have been replaced by one edge in the new graph between “3486”. We also note that the naming of the nodes have changed to reflect that the crossings “1” and “5” are no more. The current Gauss Code is: “2-c,3-c,6-a,7-a,8-c,6+a,9+c,2+c,4-c,8+c,7+a,9-c,3+c,4+c”.

There are still parallel edges in the black graph. However, the parallel edges between “29784” and “243” have the same sign. As for the parallel edges between “29784” and “687”, although the are of different signs, the relevant faces in the white graph are “3486”, “679” and “78”. The relevant faces do not pass one of our conditions for the second Reidemeister move, which that they only share two crossings in common (“6”, “7” and “8” are common crossings between the faces.). The program concludes there are no more valid Reidemeister moves of type two and moves onward.



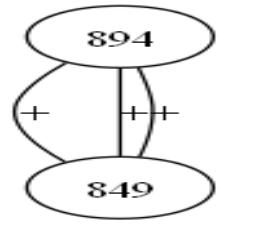
(a) Black Graph after RII.



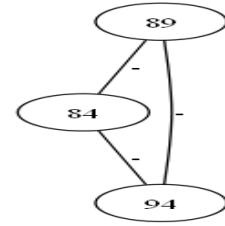
(b) White Graph after RII.

Figure 31: The Signed Planar Graphs after RII.

The cycle (“29784”, “369”, “243”) in the black graph 31a and the star (“3486”, “239”, “24”, “679”) in white 31b fulfill the conditions for the Reidemeister move of type three. Below in figure 32 are the resulting graphs with the Gauss code: “8-c,9+c,4-c,8+c,9-c,4+c”. We note that crossings “2”, “3”, “6”, “7” disappear.



(a) Black Graph after RIII.



(b) White Graph after RIII.

Figure 32: The Signed Planar Graphs after RIII.

Although, black and white have swapped graphs. The pair in the above result is isomorphic to our previous signed planar graphs construction of the Trefoil in figure 27.

In the appendix(A), we display our other results.

6 Discussion & Future Work

“The thing that doesn’t fit is the thing that’s the most interesting: the part that doesn’t go according to what you expected.”

— Richard Feynman

6.1 Discussion

When we approached our problem, we build up blocks as we tested our implementation on different knot diagrams, accommodating for the different anomalies/corner-cases as they showed up. We believe this approach was heavily flawed. Deep insights are required for each process in the signed planar graph construction. What we have initially assumed would be relatively easy turned out to be major challenges. Loops and parallel edges specifically created a lot of problems regarding the derivation of the faces, their colouring, the edges and their signs.

While the program does achieve the expected results given the knot diagrams in table(4), we cannot claim that it is foolproof, it still has its limitations. We could summarize them as following:

- Works up to ten crossings.
- Approach lacks abstraction.
- More verification and tests needed.
- Dealing with less constraints on Reidemeister moves.
- Program construction could deviate from our construction.

Let’s talk about each of these points.

Limit on the Number of Crossings

This limit comes from the nature in our approach to our program. We have used the data type “string” to utilize and manipulate the Gauss Code, naming nodes and crossings. A major issue occurs when we hit the 11th crossing and we realized this issue very late in our coding phase. If we start counting from “0” up to “9”, we face the issue when we want to add the next crossing which will take the label “10”. Currently, the program is not able to differentiate between the labels “0”, “1” and “10”.

Approach lacks Abstraction

The design of our program could benefit from abstracting our implementation. In building up our code, we handled different knot diagrams and modified our code as we went along the way to identify errors and faults in our Algorithm’s construction. Our design relies heavily on testing, validating our results and spotting anomalies. A more general approach that builds on deeper insights is required.

More Verification and Testing required

There are many Gauss Codes that we have not tested. The program *definitely* will fail in its constructions in many of them. We highlight that parallel edges and loops provide difficult and intriguing corner cases. More tests are needed to validate and update our code.

Less Constraints on Reidemeister Moves

We have altered our initial problem statement. It would be interesting to tackle the problem where the program tries to reduce the knot diagram by adding crossings and not just reduce them. It would also be interesting to take the graphs of knot diagram to another, even if both are not *reduced* knot diagrams.

Program's Construction deviates

While not necessarily a limitation, it does cause its own problems. In many cases, our program would choose a different exterior face/region while constructing the signed planar graphs from our own manual construction. Also, the labeling of the crossing from the Gauss Code heavily influence the program's choices, again causing deviations from our manual construction. This makes it harder to verify results, and keep track of what the program is doing. While this issue is of small significance as the number of crossings is not yet large enough. Issues will definitely rise later, and keeping this in the back of our minds as we recreate and modify our program will be of great benefit.

In the next subsection, we give our humble opinion on how future projects should handle our problem statement, and what could benefit upcoming undertakers (we also are definitely one of them!) as they create their own programs.

6.2 Future Work

The major feature that all of our proposed future works share, is the need for generalization. We believe the major aims of future contributions should be geared towards a more abstract approach. We highly recommend utilizing the more algebraic features of both knots and graphs. We would like to discuss the following ideas:

- Abstraction of processes.
- An Object Oriented Approach.

Abstraction of Processes

Many parts of our construction deal with other major areas in Mathematics. For example, the derivation of the faces from the shadow of the knot diagram could be considered from the perspective of planar graphs and their embedding. We highly recommend (mainly for ourselves!) a deeper study of this area in Graph Theory before tackling our problem statement. The labeling of edges with “+” or “-” could benefit from a deeper understanding of the Gauss Notation and similar algebraic structure. Each major step in the construction should be handled on a deeper level, shifting our perspective so that we underline all the nuisances and trickery that hide within the objective of the considered step.

An Object Oriented Approach

As discussed in the previous subsection, one of the major limitations is the number of crossings the program can work with. An OOP approach can easily eliminate such a drawback. Also, we envision objects specifically designed to handle the complex situations regarding parallel edges and loops as the interact with different elements in the construction. Two objects/classes that come to the forefront mind are; faces, crossings.

We also would recommend the review of the different software and solutions already available regarding graphs and knots to draw inspiration, such as *KnotAtlas*[13] and *Sage*[16]. However, up to the moment of writing this report and to our knowledge, there was no software that constructed the signed planar graphs of a knot diagram.

7 Conclusion

“It ain’t over till it’s over.”

— Yogi Berra

In this report, we present our program dealing with Reidemeister moves on alternating knot diagrams and their signed planar graphs. We highlight theorems (1) and (2) as they provide the basis for our approach. We developed our program to take as input knot diagrams represented as a Gauss codes, one of them is a *reducible* knot diagram, the program then derives their signed planar graphs. It takes the signed planar graph of the **unreduced** knot diagram, checks if there are any possible Reidemeister moves, performs them and modify its signed planar graphs. In the end, the program compares the pairs of signed planar graphs.

Our program was written using *Python*. We have also utilized external modules and libraries, two important mentions are “NetworkX” and “pyknotid”. More details are found throughout the report and also in the References.

The details of the program were shown. We gave an overview the overall structure and algorithm. We then displayed the different classes and their roles. We also introduced the main script and other crucial modules. We walked through the process of the construction of the signed planar graphs and execution Reidemeister moves. We provided detailed examples of such constructions, and in the appendix(A), we display the different outputs on our knot diagram collection.

The program still has many areas that require development. We highlighted the approach in our development phase, and showed how our program might face difficulties with anomalies such loops or parallel edges in the signed planar graphs. More tests and different knot diagrams should be run on the program as validation tests, this would definitely optimize our results and discover bugs or errors. Also, a major drawback currently is the limit on knot diagrams to have a maximum of 10 crossings.

We provided insights regarding future work. Firstly, The adaptation of a more general approach than our trial-and-error methodology should be used. Secondly, utilizing OOP to overcome certain challenges, we add that OOP is an easy and straightforward solution for the ten crossing drawback.

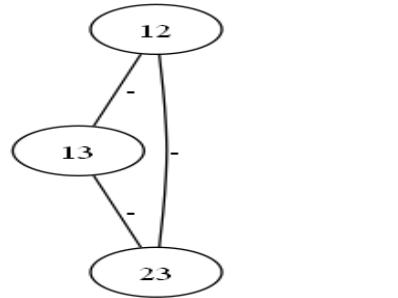
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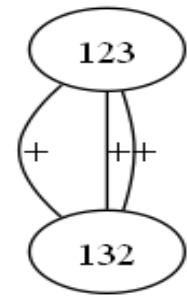
A Other Results

In the upcoming, We consider first the *irreducible* knot diagram and its signed planar graphs. These are the special outputs that we compare with the rest of our graph constructions, checking each knot diagram with its respective *irreducible* diagram.

A.1 The Trefoil



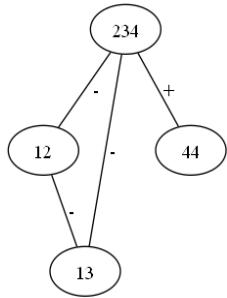
(a) Black Graph.



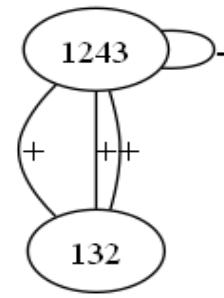
(b) White Graph.

Figure 33: The Signed Planar Graphs for the Trefoil.

A.1.1 The Trefoil with a Twist

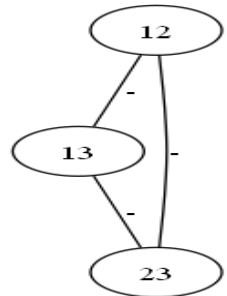


(a) Black Graph.

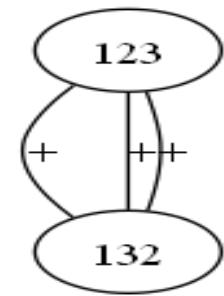


(b) White Graph.

Figure 34: initial graphs for the trefoil with a twist.



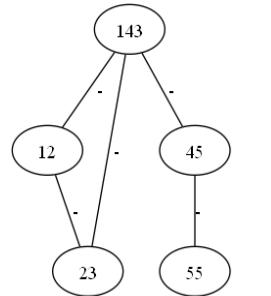
(a) Black Graph.



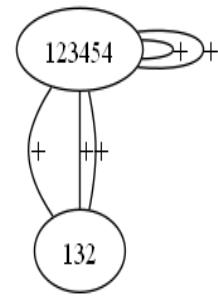
(b) White Graph.

Figure 35: final result for the trefoil with a twist.

A.1.2 The Trefoil with 2 Twists

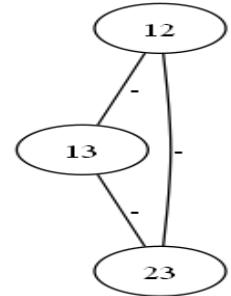


(a) Black Graph.

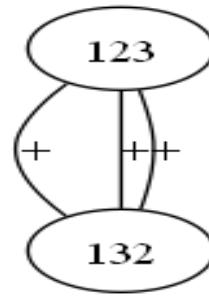


(b) White Graph.

Figure 36: initial graphs for the Trefoil with 2 twist.



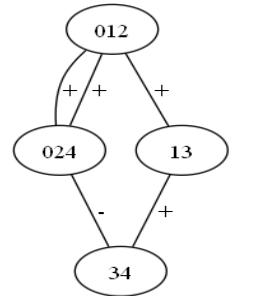
(a) Black Graph.



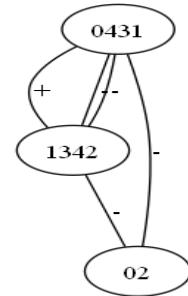
(b) White Graph.

Figure 37: final results for the Trefoil with 2 twist.

A.1.3 The Trefoil with a Poke

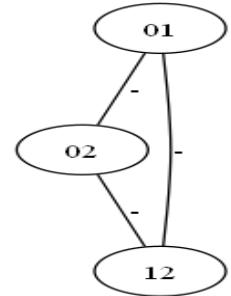


(a) Black Graph.

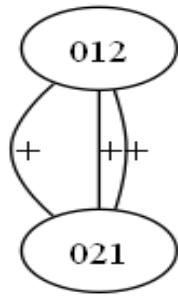


(b) White Graph.

Figure 38: initial graphs for the trefoil with a poke



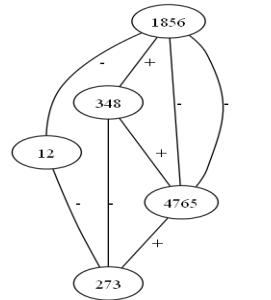
(a) Black Graph.



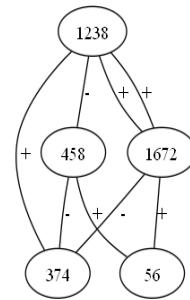
(b) White Graph.

Figure 39: final results for the trefoil with a poke

A.1.4 The Trefoil with a Twist then a Slide

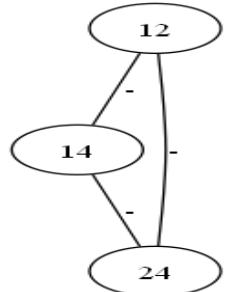


(a) Black Graph.

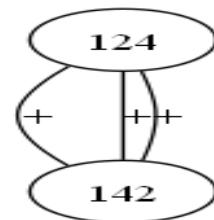


(b) White Graph.

Figure 40: initial graphs for the trefoil with a twist then a slide



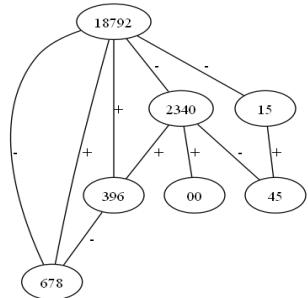
(a) Black Graph.



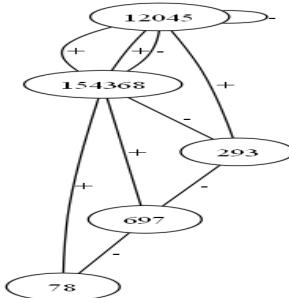
(b) White Graph.

Figure 41: final results for the trefoil with a twist then a slide

A.1.5 The Trefoil with a Twist then a Poke then a Slide

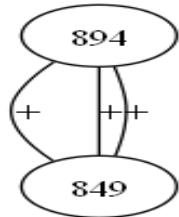


(a) The Black Graph.

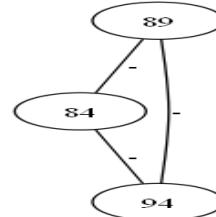


(b) The White Graph.

Figure 42: initial graphs for the trefoil with three valid moves.



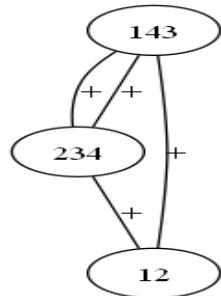
(a) Black Graph.



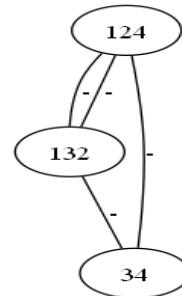
(b) White Graph.

Figure 43: final results of the trefoil with three valid moves.

A.2 The Figure Eight



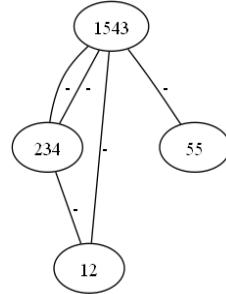
(a) Black Graph.



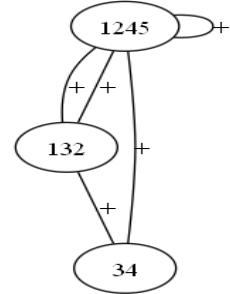
(b) White Graph.

Figure 44: Graphs for the Trefoil with a poke

A.2.1 Figure Eight with a Twist

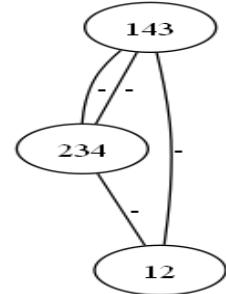


(a) Black Graph.

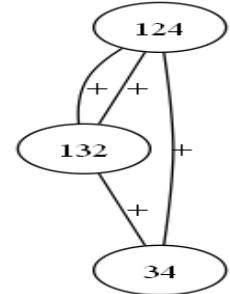


(b) White Graph.

Figure 45: initial graphs for the figure eight with a twist.



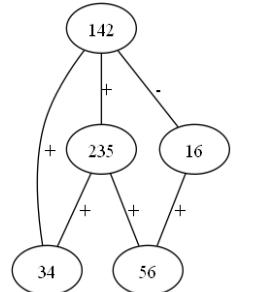
(a) Black Graph.



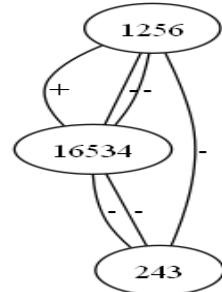
(b) White Graph.

Figure 46: final results for the figure eight with a twist.

A.2.2 Figure Eight with a poke

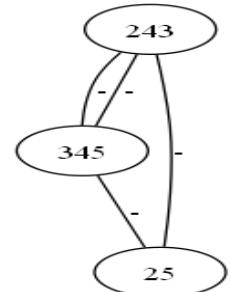


(a) Black Graph.

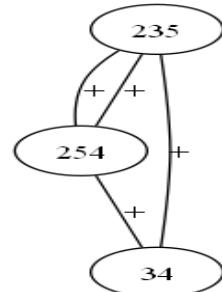


(b) White Graph.

Figure 47: initial graphs for figure eight with a poke.



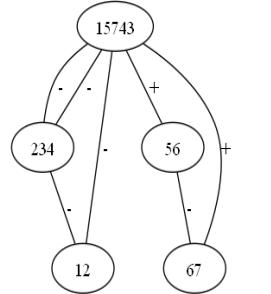
(a) Black Graph.



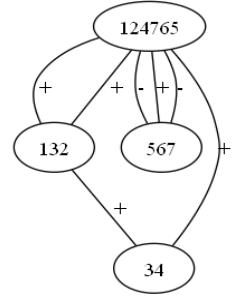
(b) White Graph.

Figure 48: final results for figure eight with a poke.

A.2.3 Figure Eight with a Twist then a Poke

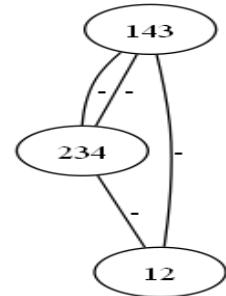


(a) Black Graph.

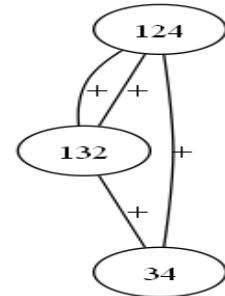


(b) White Graph.

Figure 49: initial graphs for the figure eight with a twist then a poke.



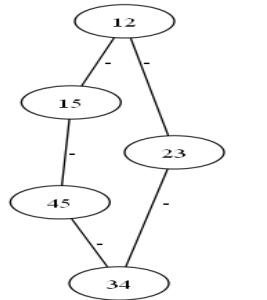
(a) Black Graph.



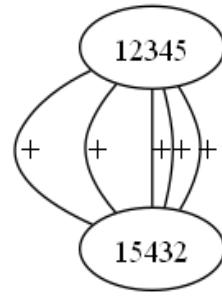
(b) White Graph.

Figure 50: final results for the figure eight with a twist then a poke.

A.3 Star



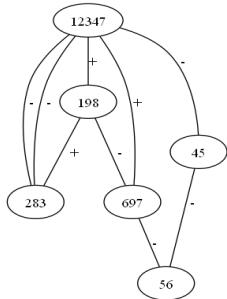
(a) Black Graph.



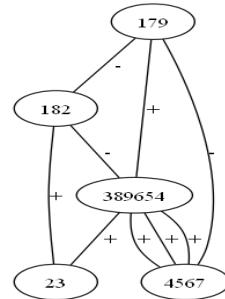
(b) White Graph.

Figure 51: Graphs for the star.

A.3.1 Star with a slide

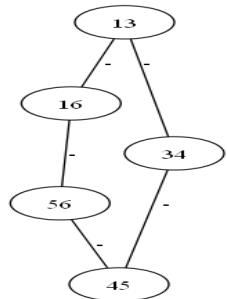


(a) Black Graph.

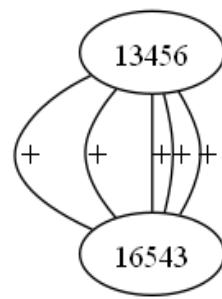


(b) White Graph.

Figure 52: initial graphs for star with a slide.



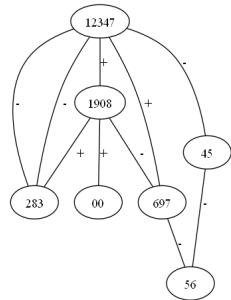
(a) Black Graph.



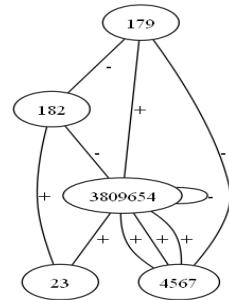
(b) White Graph.

Figure 53: final results for star with a slide.

A.3.2 Star with a Slide then a Twist

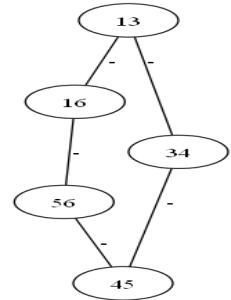


(a) Black Graph.

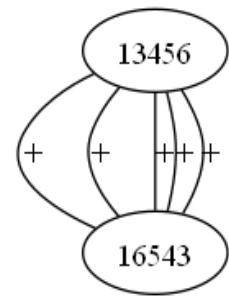


(b) White Graph.

Figure 54: initial graphs for star with a slide then a twist.



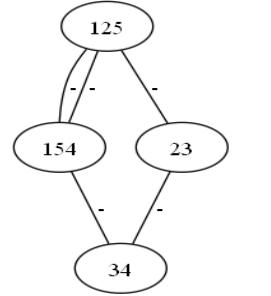
(a) Black Graph.



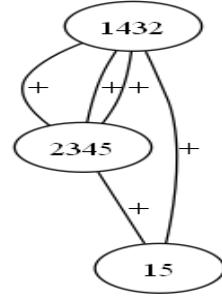
(b) White Graph.

Figure 55: final results for star with a slide then a twist.

A.4 5_2 Knot



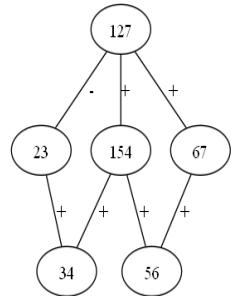
(a) Black Graph.



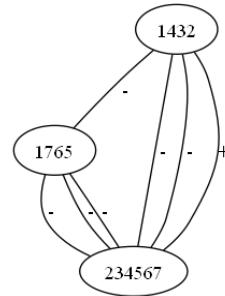
(b) White Graph.

Figure 56: Graphs for 5_2 .

A.4.1 5_2 Knot with a Poke

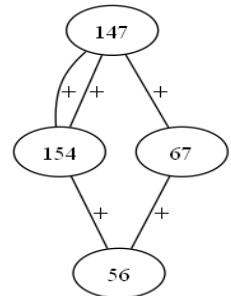


(a) Black Graph.

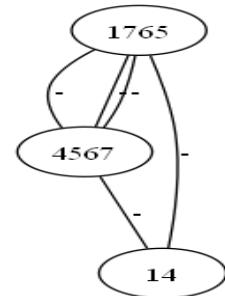


(b) White Graph.

Figure 57: initial graphs for 5_2 with a poke.



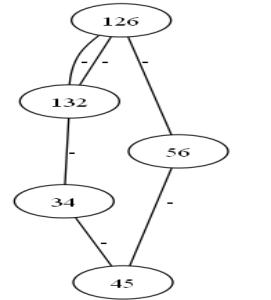
(a) Black Graph.



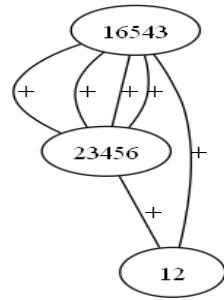
(b) White Graph.

Figure 58: final results for 5_2 with a poke.

A.5 6_1 Knot



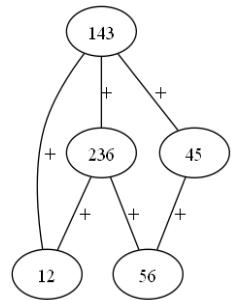
(a) Black Graph.



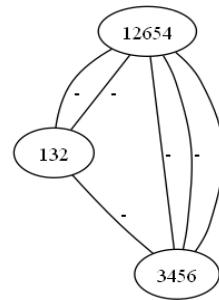
(b) White Graph.

Figure 59: Graphs for 6_1 .

A.6 6_2 Knot



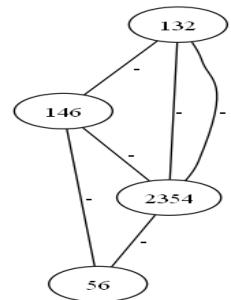
(a) Black Graph.



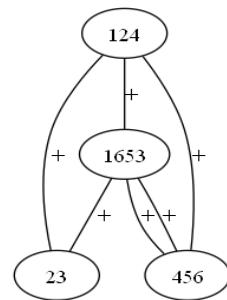
(b) White Graph.

Figure 60: Graphs for 6_2 .

A.7 6_3 Knot



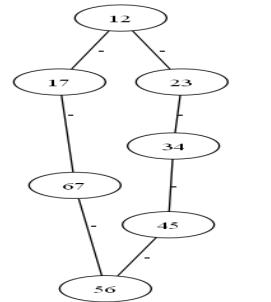
(a) Black Graph.



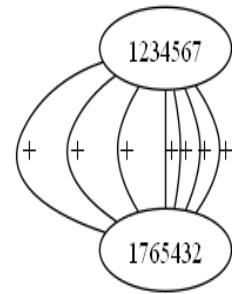
(b) White Graph.

Figure 61: Graphs for 5_3 .

A.8 7_1 Knot



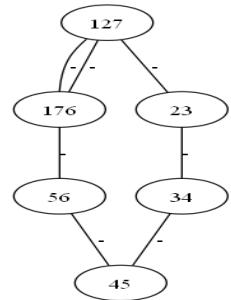
(a) Black Graph.



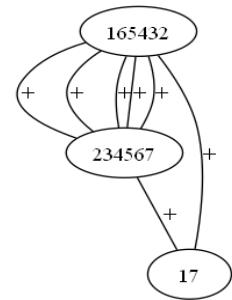
(b) White Graph.

Figure 62: Graphs for 7_1 .

A.9 7_2 Knot



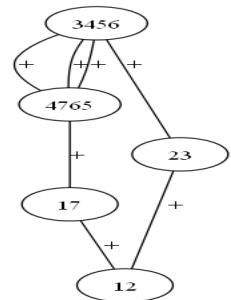
(a) Black Graph.



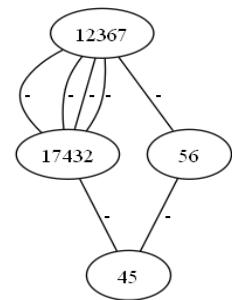
(b) White Graph.

Figure 63: Graphs for 7_2 .

A.10 7_3 Knot



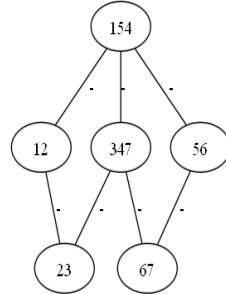
(a) Black Graph.



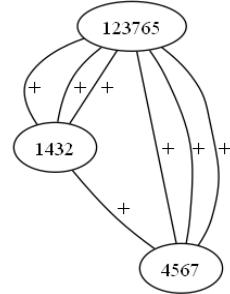
(b) White Graph.

Figure 64: Graphs for 7_3 .

A.11 7_4 Knot



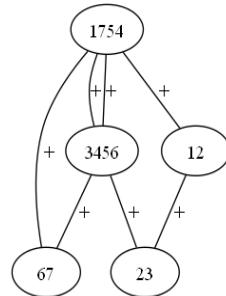
(a) Black Graph.



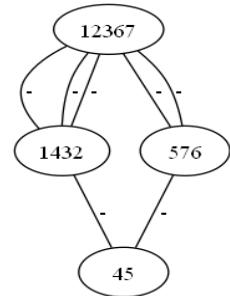
(b) White Graph.

Figure 65: Graphs for 7_4 .

A.12 7_5 Knot



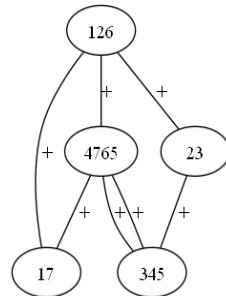
(a) Black Graph.



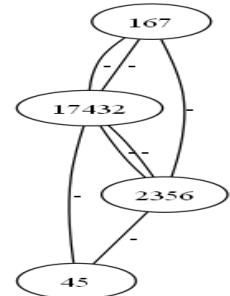
(b) White Graph.

Figure 66: Graphs for 7_5 .

A.13 7_6 Knot



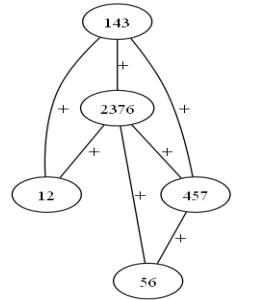
(a) Black Graph.



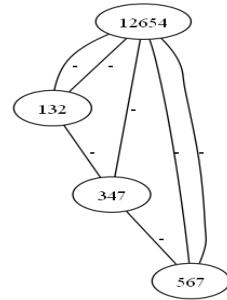
(b) White Graph.

Figure 67: Graphs for 7_6 .

A.14 7_7 Knot



(a) Black Graph.

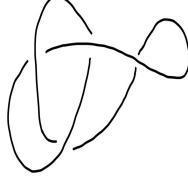
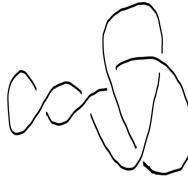
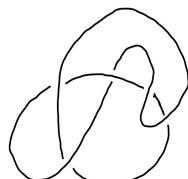
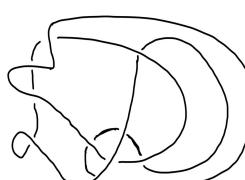


(b) White Graph.

Figure 68: Graphs for 7_7 .

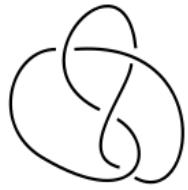
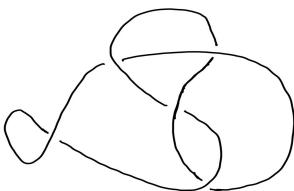
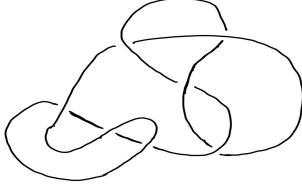
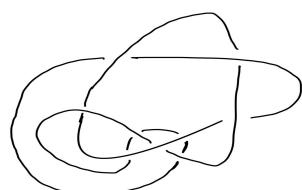
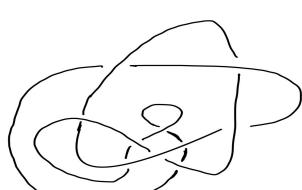
B HardCodedKnots Codes

Table 4: Hard-coded Knots Information

Knot	code_no	code_string	Gauss Code	Description	Knot Diagram
3_1	31	3_1 trefoil	1+c,2-c,3+c, 1-c,2+c,3-c	The trefoil	
3_1	311	3_1_1 trefoil_twist	1+c,2-c,3+c, 1-c,2+c,4+c, 4-c,3-c	The trefoil with one twist	
3_1	3111	3_1_1_1 trefoil_twist _twist	1+c,2-c,3+c, 4-a,5+a,5-a, 4+a,1-c,2+c, 3-c	The trefoil with two twist	
3_1	312	3_1_2 trefoil_poke	0-c,1+c,3-c, 4-a,2-c,0+c, 4+a,3+c,1-c, 2+c	The trefoil with a poke	
3_1	3113	3_1_1_3 trefoil_twist _slide	1+c,2-c,3+c, 4+c,5+a,6-a, 7-c,3-c,8-a, 5-a,6+a,1-c, 2+c,7+c,4-c, 8+a	The trefoil with a twist and a slide	
3_1	31123	3_1_1_2_3 trefoil_twist _poke_slide	1+c,2-c,3-c, 6a,7-a,8-c, 6+a,9+c,2+c, 0+c,0-c,4-c, 5-a,1-c,8+c, 7+a,9-c,3+c, 4+c,5+a	The trefoil with three valid moves.	

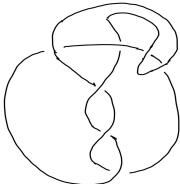
Continued on next page

Table 4 – Hardcoded Knots continued

Knot	code_no	code_string	Gauss Code	Description	Knot Diagram
4 ₁	41	4_1 figure_eight	1-c,2+c,3-a, 4+a,2-c,1+c, 4-a,3+a	The figure eight	
4 ₁	411	4_1_1 figure_eight _twist	1+a,2- a,3+a, 4-c,2+a,1-a, 5+a,5-a,4+c, 3-c	The figure eight with a twist	
4 ₁	4112	4_1_1_2 figure_eight _twist_poke	1+a,2-a,3+c, 4-c,2+a,1-a, 5+a,6+c,7+a, 5-a,6-c,7-a, 4+c,3-c	The figure eight with a twist	
5 ₁	51	5_1 star	1+c,2-c,3+c, 4-c,5+c,1-c, 2+c,3-c,4+c, 5-c	The star	
5 ₁	513	5_1_3 star_slide	1+c,2+c,3-c, 4+c,5-c,6+c, 7+a,1-c,8+c, 3+c,2-c,8-c, 9-a,7-a,4-c, 5+c,6-c,9+a	The star with a slide	
5 ₁	5131	5_1_3_1 star_slide _twist	1+c,2+c,3-c, 4+c,5-c,6+c, 7+a,1-c,8+c, 3+c,2-c,8-c, 0+c,0-c, 9-a,7-a,4-c, 5+c,6-c,9+a	The star with a slide and a twist	

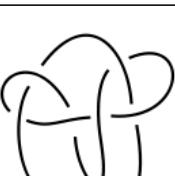
Continued on next page

Table 4 – Hardcoded Knots continued

Knot	code_no	code_string	Gauss Code	Description	Knot Diagram
5 ₂	52	5_2	1+a,2- a,3+a, 4-a,5+a,1-a, 4+a,3- a,2+a, 5-a	The 5 ₂ knot	
5 ₂	522	5_2	1+c,2+a,3+c, 4-c,5+c,6-c, 7+c,1-c,4+c, 3-c,2-a,7-c, 6+c,5-c	The 5 ₂ knot	
6 ₁	61	6_1	1-c,2+c,3-a, 4+a,5- a,6+a, 2-c,1+c,6-a, 5+a,4-a,3+a	The 6 ₁ knot	
6 ₂	62	6_2	1-c,2+c,3-a, 4+a,5- a,6+a, 2-c,1+c,4-a, 5+a,6-a,3+a	The 6 ₂ knot	
6 ₃	63	6_3	1+a,2- a,3+a, 1-a,4+c,5-c, 6+c,4-c,2+a, 3-a,5+c,6-c	The 6 ₃ knot	
7 ₁	71	7_1	1+c,2-c,3+c, 4-c,5+c,6-c, 7+c,1-c,2+c, 3-c,4+c,5-c, 6+c,7-c	The 7 ₁ knot	

Continued on next page

Table 4 – Hardcoded Knots continued

Knot	code_no	code_string	Gauss Code	Description	Knot Diagram
7 ₂	72	7_2	1+a,2-a, a,3+a, 4-a,5+a,6-a, 7+a,1-a, a,6+a, 5-a,4+a,3-a, 2+a,7-a	The 7 ₂ knot	
7 ₃	73	7_3	1+a,2-a, 3+a,4-a, a,5+a, 6-a,7+a,1-a, 2+a,3-a, 6+a,5-a, 4+a, 7-a	The 7 ₃ knot	
7 ₄	74	7_4	1-a,2+a,3-a, 4+a,5-a, a,6+a, 7-a,3+a,2-a, 1+a,4-a, 7+a, 6-a,5+a	The 7 ₄ knot	
7 ₅	75	7_5	1+a,2-a, a,3+a, 4-a,5+a,6-a, 7+a,5-a, 4+a, 1-a,2+a,3-a, 6+a,7-a	The 7 ₅ knot	
7 ₆	76	7_6	1-a,2+c,3-c, 4+a,5-a,3+c, 2-c,6+a,7-a, 1+a,6-a, 5+a, 4-a,7+a	The 7 ₆ knot	
7 ₇	77	7_7	1-c,2+c,3-a, 4+a,5-c,6+c, 2-c,1+c,4-a, 7+a,6-c,5+c, 7-a,3+a	The 7 ₇ knot	

C the Whole Code

C.1 main.py

```

1  from SignedPlanarGraph import SignedPlanarGraph
2  from HardCodedKnots import HardCodedKnots
3  from ReidemeisterMoves import ReidemeisterMoves
4
5  hcg = HardCodedKnots(code_no=3113)
6
7  k_gauss = hcg.get_code_gauss()
8
9  sp_graph = SignedPlanarGraph(code_gauss = k_gauss)
10 sp_graph.create_and_draw_sp()
11
12 reid_sp = ReidemeisterMoves.reidemeister_main(sp_graph)
13
14 reid_sp = SignedPlanarGraph(code_gauss=reid_sp.code_gauss)
15 reid_sp.create_and_draw_sp()
16
17 #print(reid_sp.code_gauss)

```

C.2 check_gauss_code.py

```

1  '''
2  We use this file to make sure our creations, our certain results
3  are correct by calculating manually the Gauss Code and checking if
4      it
5  gives us the expected knot
6  '''
7
8  import pyknotid.representations as rep_
9  import numpy as np
10 import sys
11 import pyknotid.make as mk
12
13 gc = rep_.GaussCode("1+c,4-c,5+c,6-c,7+c,1-c,4+c,7-c,6+c,5-c")
14 rep = rep_.Representation(gc)
15 k = rep.space_curve(bool=False)
16 # k.plot()
17
18 print(k.identify())

```

C.3 main_iterative.py

```
1  from pyknotid.catalogue import get_knot
2  import networkx as nx
3  import matplotlib.pyplot as plt
4  from networkx.drawing.nx_agraph import to_agraph
5  import my_helpers as help
6  from SignedPlanarGraph import SignedPlanarGraph
7  from HardCodedKnots import HardCodedKnots
8  from ReidemeisterMoves import ReidemeisterMoves
9  import sys
10
11
12 for code_no in HardCodedKnots.code_numbers:
13
14     hcg = HardCodedKnots(code_no = code_no)
15
16     k_gauss = hcg.get_code_gauss()
17
18
19     sp_graph = SignedPlanarGraph(code_gauss = k_gauss)
20     sp_graph.create_and_draw_sp()
21
22
23     reid_sp = ReidemeisterMoves.reidemeister_main(sp_graph)
24
25     reid_sp = SignedPlanarGraph(code_gauss=reid_sp.code_gauss)
26     reid_sp.create_and_draw_sp()
27
28     # Uncomment following if you would like to check after each run
29     # of the program
30     # on different Gauss Code.
31     # print(reid_sp.code_gauss)
32     # input()
```

C.4 my_helpers.py

```

1      ''
2      This file is for methods and/or modules that will
3      be used across the whole project.
4      For easier access and better structure.
5      ''
6
7
8      import numpy as np
9      import collections as col
10     import pyknotid.representations as rep_
11     import networkx as nx
12
13
14      ''
15      Used to check whether two lists are equivalent w.r.t. Cyclic order.
16      Order is considered from left to right.
17      This method is mainly used to check that we do not have duplicate
18          faces
19      When creating signed planar graphs.
20
21      ex 1:
22      l1 = [ 1, 2 , 3]
23      l2 = [3 , 1, 2]
24      is_cyclic_list(l1, l2) => TRUE
25
26      ex 2:
27      l1 = [ 1, 2 , 3]
28      l2 = [3 , 2, 1]
29      is_cyclic_list(l1, l2) => False
30      ''
31      def is_cyclic_list(l1, l2):
32          if len(l1) != len(l2):
33              return False
34          if not l1:
35              return False
36          elem = l1[0]
37          for i in range(0,len(l1)):
38              if elem == l2[i]:
39                  if l1 == l2[i : ] + l2[ : i]:
40                      return True
41              return False
42      return False
43
44      ''
45      To check if an ordered subset of list1
46      is contained in list2. Order is important!!
47      ''
48      def sublistExists(list1, list2):
49          str_l1 = str(list1).strip("[] ")
50          str_l2 = str(list2).strip("[] ")
51          return str_l1 in str_l2

```

```

52     '''
53     used to create edges list
54     in order to create the edges between the nodes
55     in the dual graph.
56     '''
57     def edges_set_dual_creation(str_value):
58         if not str_value:
59             return
60
61         edges_set = []
62
63         for i in range(0, len(str_value) - 1):
64             edges_set.append(str_value[i: i+2])
65             edges_set.append(str_value[len(str_value) - 1] + str_value[:1])
66
67         return edges_set
68
69     '''
70     To Create all edges as strings
71     each edge consists of two characters
72     the characters represent the nodes.
73     '''
74     def edges_set_creation(str_value):
75         if not str_value:
76             return
77
78         edges_set = []
79
80         for i in range(0, len(str_value) - 1):
81             edges_set.append(str_value[i: i+2])
82             edges_set.append(str_value[len(str_value) - 1] + str_value[:1])
83
84         return edges_set
85
86     '''
87     This method is to take a Gauss Code provided from either
88     The library pyknotid, or from usual conventional mathematical
89     notation and modify to start counting from 0.
90     e.g. 1+c,2+a,... => 0+c,1+a,...
91     '''
92
93     def code_gauss_from_math_to_comp(code_gauss):
94         code_gauss_splitted = str(code_gauss).split(",")
95
96         joker_code_gauss = ""
97
98         for gc in code_gauss_splitted:
99             crossing_no = int(gc[0])
100            crossing_no = crossing_no - 1
101            gc = str(crossing_no) + gc[1] + gc[2]
102            joker_code_gauss = joker_code_gauss + gc + ","
103
104         last_item = joker_code_gauss[len(joker_code_gauss)-1]
105         if last_item == ",":

```

```
106     joker_code_gauss = joker_code_gauss[:len(joker_code_gauss)
107         -1]
108
109     return joker_code_gauss
110
111     '',
112     https://stackoverflow.com/questions/48688494/python-string-search-
113     regardless-of-character-sequence
114     ''',
115     def edge\_contained\(substr, str\):
116         c1 = col.Counter\(str\)
117         c2 = col.Counter\(substr\)
118         return not\(c2-c1\)
119
120     '',
121     Create planar diagram from Gauss Code
122     ''',
123
124     def create\_planar\_graph\_from\_gauss\_code\(code\_gauss\):
125
126         if not code\_gauss:
127             raise Exception\("Gauss\u2022Code\u2022is\u2022Empty"\)
128
129         str\_gc = str\(code\_gauss\)
130         code\_gauss\_splitted = str\_gc.split\(","\)
131
132         if len\(code\_gauss\_splitted\) % 2 != 0:
133             raise Exception\("Gauss\u2022Code\u2022Is\u2022invalid"\)
134
135         # creating a knot from the gauss code
136         # to make sure that the gauss code
137         # is a valid knot
138         joker\_gc = rep\_.GaussCode\(str\(code\_gauss\)\)
139         rep = rep\_.Representation\(joker\_gc\)
140         k = rep.space\_curve\(bool=True\)
141
142         graph = nx.MultiGraph\(\)
143
144         i = 0
145         while i < len\(code\_gauss\_splitted\):
146             gc = code\_gauss\_splitted\[i\]
147             if i == len\(code\_gauss\_splitted\) - 1:
148                 gc\_next = code\_gauss\_splitted\[0\]
149             else:
150                 gc\_next = code\_gauss\_splitted\[i+1\]
151             graph.add\_edge\(int\(gc\[0\]\),int\(gc\_next\[0\]\)\)
152             i+=1
153
154     return graph
155
156     '',
157     To get the face that bounds the over and understrand
158     '''
```

```
158
159 def get_bounded_face(edges, faces):
160     faces_list = []
161     if len(edges[0]) == 2 and edges[0][0] == edges[0][1]:
162         edges[0] = edges[0][0]
163     if len(edges[1]) == 2 and edges[1][0] == edges[1][1]:
164         edges[1] = edges[1][0]
165
166     for face in faces:
167         if edge_contained(edges[0], face):
168             if edge_contained(edges[1], face):
169                 faces_list.append(face)
170
171     if len(faces_list) == 0:
172         return None
173
174     no_of_edges = []
175     for face in faces_list:
176         no_of_edges.append(len(face))
177
178     bounded_face_index = no_of_edges.index(min(no_of_edges))
179     return faces_list[bounded_face_index]
180
181 def is_subset_both_ways(str_1, str_2):
182     chars_1 = set(str_1)
183     chars_2 = set(str_2)
184     return len(str_1) == len(str_2) and chars_1.issubset(chars_2)
185         and chars_2.issubset(chars_1)
```

C.5 HardCodedKnots.py

```

1  from pyknotid.catalogue import get_knot
2  import pyknotid.representations as rep_
3
4  class HardCodedKnots:
5
6      code_numbers = [31, 311, 3111, 312, 3113, 41, 411, 4112, 412,
7          51, 513, 5131, 52, 522, 61, 62, 63,
8          71, 72, 73, 74, 75, 76, 77]
9
10     def __init__(self, code_no = None, code_string = None):
11
12         if code_no == 31 or code_string == "trefoil" or code_string
13             == "3_1":
14             k = get_knot('3_1').space_curve()
15             self.code_gauss = k.gauss_code()
16             return
17
18         if code_no == 311 or code_string == "trefoil_twist" or
19             code_string == "3_1_1":
20             self.code_gauss = rep_.GaussCode("1+c,2-c,3+c,1-c,2+c,4+
21                 c,4-c,3-c")
22             return
23
24         if code_no == 3111 or code_string == "trefoil_twist_twist"
25             or code_string == "3_1_1_1":
26             self.code_gauss = rep_.GaussCode("1+c,2-c,3+c,4-a,5+a,5-
27                 a,4+a,1-c,2+c,3-c")
28             return
29
30         if code_no == 312 or code_string == "trefoil_poke" or
31             code_string == "3_1_2":
32             self.code_gauss = rep_.GaussCode("0-c,1+c,3-c,4-a,2-c,0+
33                 c,4+a,3+c,1-c,2+c")
34             return
35
36         if (code_no == 3113 or code_string == "trefoil_twist_slide"
37             or code_string == "3_1_1_3"):
38             self.code_gauss = rep_.GaussCode("1+c,2-c,3+c,4+c,5+a,6-
39                 a,7-c,3-c,8-a,5-a,6+a,1-c,2+c,7+c,4-c,8+a")

```

```

40     if code_no == 4112 or code_string == "figure_eight_twist_poke" or code_string == "4_1_1_2":
41         self.code_gauss = rep_.GaussCode("1+a,2-a,3+c,4-c,2+a,1-
42                                         a,5+a,6+c,7+a,5-a,6-c,7-a,4+c,3-c")
43         return
44
45     if code_no == 412 or code_string == "figure_eight_poke" or code_string == "4_1_2":
46         self.code_gauss = rep_.GaussCode("1-c,2-a,3+c,4-c,2+a,5-
47                                         a,6+a,1+c,4+c,3-c,5+a,6-a")
48         return
49
50     if code_no == 51 or code_string == "star" or code_string == "5_1":
51         k = get_knot("5_1").space_curve()
52         self.code_gauss = k.gauss_code()
53         return
54
55     if code_no == 513 or code_string == "star_slide" or code_string == "5_1_3":
56         self.code_gauss = rep_.GaussCode("1+c,2+c,3-c,4+c,5-c,6+
57                                         c,7+a,1-c,8+c,3+c,2-c,8-c,9-a,7-a,4-c,5+c,6-c,9+a")
58         return
59
60     if code_no == 5131 or code_string == "star_slide_twist" or code_string == "5_1_3_1":
61         self.code_gauss = rep_.GaussCode("1+c,2+c,3-c,4+c,5-c,6+
62                                         c,7+a,1-c,8+c,3+c,2-c,8-c,0+c,0-c,9-a,7-a,4-c,5+c,6-c,
63                                         ,9+a")
64         return
65
66     if code_no == 52 or code_string == "5_2":
67         k = get_knot("5_2").space_curve()
68         self.code_gauss = k.gauss_code()
69         return
70
71     if code_no == 522 or code_string == "5_2_2":
72         self.code_gauss = rep_.GaussCode("1+c,2+a,3+c,4-c,5+c,6-
73                                         c,7+c,1-c,4+c,3-c,2-a,7-c,6+c,5-c")
74         return
75
76     if code_no == 61 or code_string == "6_1":
77         k = get_knot("6_1").space_curve()
78         self.code_gauss = k.gauss_code()
79         return
80
81     if code_no == 62 or code_string == "6_2":
82         k = get_knot("6_2").space_curve()
83         self.code_gauss = k.gauss_code()
84         return
85
86     if code_no == 63 or code_string == "6_3":
87         k = get_knot("6_3").space_curve()
88         self.code_gauss = k.gauss_code()

```

```
83         return
84
85     if code_no == 71 or code_string == "7_1":
86         k = get_knot("7_1").space_curve()
87         self.code_gauss = k.gauss_code()
88         return
89
90     if code_no == 72 or code_string == "7_2":
91         k = get_knot("7_2").space_curve()
92         self.code_gauss = k.gauss_code()
93         return
94
95     if code_no == 73 or code_string == "7_3":
96         k = get_knot("7_3").space_curve()
97         self.code_gauss = k.gauss_code()
98         return
99
100    if code_no == 74 or code_string == "7_4":
101        k = get_knot("7_4").space_curve()
102        self.code_gauss = k.gauss_code()
103        return
104
105    if code_no == 75 or code_string == "7_5":
106        self.code_gauss = rep_.GaussCode("1+a,2-a,3+a,4-a,5+a,6-
107            a,7+a,5-a,4+a,1-a,2+a,3-a,6+a,7-a")
108        return
109
110    if code_no == 76 or code_string == "7_6":
111        k = get_knot("7_6").space_curve()
112        self.code_gauss = k.gauss_code()
113        return
114
115    if code_no == 77 or code_string == "7_7":
116        k = get_knot("7_7").space_curve()
117        self.code_gauss = k.gauss_code()
118        return
119
120    def get_code_gauss(self):
121        return self.code_gauss
```

C.6 SignedPlanarGraph.py

```
1 import networkx as nx
2 import my_helpers as help
3 import copy
4
5
6 class SignedPlanarGraph:
7
8     drawing_number = 0
9
10    def __init__(self, black_graph_data = None, white_graph_data =
11                  None, knot = None,
12                  code_gauss = None, signed_black = None,
13                  signed_white = None):
14        if black_graph_data:
15            self.black_graph = nx.MultiGraph(black_graph_data)
16        else:
17            self.black_graph = nx.MultiGraph()
18        if white_graph_data:
19            self.white_graph = nx.MultiGraph(white_graph_data)
20        else:
21            self.white_graph = nx.MultiGraph()
22        if knot:
23            self.knot = knot
24        else:
25            self.knot = None
26        if code_gauss:
27            self.code_gauss = code_gauss
28        else:
29            self.code_gauss = None
30        if signed_black:
31            self.signed_black = signed_black
32        else:
33            self.signed_black = []
34        if signed_white:
35            self.signed_white = signed_white
36        else:
37            self.signed_white = []
38
39    def __eq__(self, other):
40
41        if str(self.code_gauss) == str(other.code_gauss):
42            if nx.is_isomorphic(self.black_graph, other.black_graph):
43                :
44                    if nx.is_isomorphic(self.white_graph, other.
45                        white_graph):
46                        if self.signed_black == other.signed_black:
47                            if self.signed_white == other.signed_white:
48                                return True
49
50    return False
51
52
53    def __ne__(self, other):
54
```

```

49         return not self.__eq__(other)
50
51     def create_nodes_from_faces(self, graph):
52
53         (check, planar_graph) = nx.check_planarity(graph)
54         faces_first_run = SignedPlanarGraph.get_faces_list(
55             planar_graph)
56         faces = SignedPlanarGraph.add_loops_face_list(graph,
57             faces_first_run)
58         outer_face = SignedPlanarGraph.get_longest_region(faces)
59         outer_regions = SignedPlanarGraph.get_outer_regions(
60             outer_face, faces)
61         [black, white] = SignedPlanarGraph.shade_faces_new(
62             outer_regions, faces, outer_face)
63
64         black_nodes = []
65         for b in black:
66             black_nodes.append(str(b).strip("[]").replace(",\u2022", ""))
67             self.black_graph.add_nodes_from(black_nodes)
68
69         white_nodes = []
70         for w in white:
71             white_nodes.append(str(w).strip("[]").replace(",\u2022", ""))
72             self.white_graph.add_nodes_from(white_nodes)
73
74         self.black_loops = []
75         for b in self.black_graph.nodes():
76             if len(b) == 2:
77                 if b[0] == b[1]:
78                     self.black_loops.append(b)
79
80         self.white_loops = []
81         for w in self.white_graph.nodes():
82             if len(w) == 2:
83                 if w[0] == w[1]:
84                     self.white_loops.append(w)
85
86     def create_edges_after_nodes(self):
87
88         nodes = self.black_graph.nodes()
89         nodes = list(nodes)
90         edges = []
91
92         i = 0
93         while i < len(nodes):
94             j = i + 1
95             while j < len(nodes):
96                 if len(nodes[j]) == 2:
97                     if nodes[j][0] == nodes[j][1]:
98                         if nodes[j][0] in nodes[i]:

```

```

99                     continue
100                for ch in nodes[j]:
101                    if ch in nodes[i]:
102                        edges.append((nodes[i], nodes[j]))
103                        j += 1
104                white_loops = copy.deepcopy(self.white_loops)
105                for white_loop in white_loops:
106                    if white_loop[0] in nodes[i]:
107                        edges.append((nodes[i], nodes[i]))
108                        the_last = white_loop[0]
109                        the_last_index = nodes[i].find(the_last)
110                        m = (the_last_index - 1) % len(nodes[i])
111                        n = (the_last_index + 1) % len(nodes[i])
112                        while True:
113                            if nodes[i][m] == nodes[i][n]:
114                                edges.append((nodes[i], nodes[i]))
115                                another = the_last + " " + nodes[i][m]
116                                self.white_loops.append(another)
117                                m = (m - 1) % len(nodes[i])
118                                n = (n + 1) % len(nodes[i])
119                                continue
120                            break
121                i += 1
122
123
124        self.black_graph.add_edges_from(edges)
125
126    nodes = self.white_graph.nodes()
127    nodes = list(nodes)
128    edges = []
129
130    i = 0
131    while i < len(nodes):
132        j = i + 1
133        while j < len(nodes):
134            if len(nodes[j]) == 2:
135                if nodes[j][0] == nodes[j][1]:
136                    if nodes[j][0] in nodes[i]:
137                        edges.append((nodes[i], nodes[j]))
138                        j += 1
139                        continue
140                    for ch in nodes[j]:
141                        if ch in nodes[i]:
142                            edges.append((nodes[i], nodes[j]))
143                        j += 1
144                black_loops = copy.deepcopy(self.black_loops)
145                for black_loop in black_loops:
146                    if black_loop[0] in nodes[i]:
147                        edges.append((nodes[i], nodes[i]))
148                        the_last = black_loop[0]
149                        the_last_index = nodes[i].find(the_last)
150                        m = (the_last_index - 1) % len(nodes[i])
151                        n = (the_last_index + 1) % len(nodes[i])
152                        while True:

```

```

153         if nodes[i][m] == nodes[i][n]:
154             edges.append((nodes[i], nodes[i]))
155             another = the_last + " " + nodes[i][m]
156             self.black_loops.append(another)
157             m = (m - 1) % len(nodes[i])
158             n = (n + 1) % len(nodes[i])
159             continue
160         break
161     i += 1
162
163     self.white_graph.add_edges_from(edges)
164
165     @staticmethod
166     def get_faces_list(planar_graph):
167         graph = planar_graph
168         faces_list = []
169         for node_1 in graph.nodes():
170             for node_2 in graph.nodes():
171                 if node_2 not in graph[node_1]:
172                     continue
173                 faces_list.append(graph.traverse_face(node_1, node_2))
174
175         i = 0
176         while i < len(faces_list):
177             j = i + 1
178             while j < len(faces_list):
179                 if help.is_cyclic_list(faces_list[i], faces_list[j]):
180                     :
181                     faces_list.remove(faces_list[j])
182                     j = j - 1
183                     j = j + 1
184                     i = i + 1
185         return faces_list
186
187     @staticmethod
188     def get_longest_region(faces_list):
189         no_edges_in_each_face = []
190
191         for face in faces_list:
192             no_edges_in_each_face.append(len(face))
193
194         outer_region_index = no_edges_in_each_face.index(max(
195             no_edges_in_each_face))
196         return faces_list[outer_region_index]
197
198     #Add loops as faces
199     @staticmethod
200     def add_loops_face_list(orig_graph, faces_list):
201         graph = orig_graph
202         edges = graph.edges()
203         edges_list = list(edges)
204         edge_count = {}

```

```

204     for e in edges:
205         e_str = str(e).strip("()").replace(",\u2022", " ")
206         e_str = e_str.replace("'", " ")
207         rev_e_str = e_str[::-1]
208         if e_str not in edge_count and rev_e_str not in
209             edge_count:
210                 edge_count[e_str] = 1
211                 if e_str[0] == e_str[1]:
212                     edge_count[e_str] += 1
213             elif e_str in edge_count:
214                 edge_count[e_str] += 1
215             else:
216                 edge_count[rev_e_str] += 1
217         for e in edge_count:
218             if edge_count[e] > 1:
219                 face = [int(e[0:1]), int(e[1:])]
220                 for i in range(0, edge_count[e] - 1):
221                     faces_list.append(face)
222     return faces_list
223
224 @staticmethod
225 def get_outer_regions(longest_region, faces_list):
226
227     outer_regions = []
228     loops = []
229
230     for face in faces_list:
231         if face != longest_region:
232             if SignedPlanarGraph.check_same_edge_region(face,
233                 longest_region):
234                 outer_regions.append(face)
235                 check_if_loop = str(face).strip("[]").replace(",",
236                     "\u2022", " ")
237                 if len(check_if_loop) == 2:
238                     loops.append(face)
239
240     for loop in loops:
241         if loop[0] == loop[1]:
242             continue
243         for region in outer_regions:
244             if loop != region:
245                 if SignedPlanarGraph.check_same_edge_region(loop,
246                     region):
247                     outer_regions.remove(region)
248     return outer_regions
249
250 @staticmethod
251 def shade_faces_new(outer_reigons, faces_list, outer_face):
252
253     black_shades = outer_reigons
254     white_shades = [outer_face]
255
256     coloured = black_shades + white_shades

```

```

253     ple_faces = [str(face[0]) + str(face[1]) for face in
254                 faces_list if len(face) == 2]
255
255     nbr_black_turn = True
256     got_coloured= outer_reigons
257
258
259     skipped = False
260     while len(coloured) != len(faces_list):
261         next_got_coloured = []
262         if nbr_black_turn:
263             for face in got_coloured:
264                 for next_face in faces_list:
265                     if next_face not in coloured:
266                         if face != next_face:
267                             if SignedPlanarGraph.
268                                 check_same_edge_region(face,
269                                         next_face):
270                                 if len(next_face) > 2:
271                                     [cond, edge] =
272                                         SignedPlanarGraph.
273                                         check_same_edge_region(
274                                             face, next_face)
275                                         if edge in ple_faces or edge
276                                         [::-1] in ple_faces:
277                                             skipped = True
278                                             continue
279                                         coloured.append(next_face)
280                                         white_shades.append(next_face)
281                                         next_got_coloured.append(
282                                             next_face)
283                                         nbr_black_turn = False
284                                         got_coloured = next_got_coloured
285                                     else:
286                                         for face in got_coloured:
287                                             for next_face in faces_list:
288                                                 if next_face not in coloured:
289                                                     if face != next_face:
290                                                         if SignedPlanarGraph.
291                                                             check_same_edge_region(face,
292                                     next_face):
293                                         if len(next_face) > 2:
294                                             [cond, edge] =
295                                                 SignedPlanarGraph.
296                                                 check_same_edge_region(
297                                                     face, next_face)
298                                                 if edge in ple_faces or edge
299                                                 [::-1] in ple_faces:
300                                                     continue
301                                                 coloured.append(next_face)
302                                                 black_shades.append(next_face)
303                                                 next_got_coloured.append(
304                                                     next_face)
305                                         nbr_black_turn = True

```

```

292             got_coloured = next_got_coloured
293         if skipped and len(got_coloured) == 0 and len(coloured)
294             != len(faces_list):
295             remaining = [face for face in faces_list if face not
296                           in coloured]
297             for face in remaining:
298                 jump = False
299                 for next_face in coloured:
300                     if SignedPlanarGraph.check_same_edge_region(
301                         face, next_face):
302                         if next_face in black_shades:
303                             white_shades.append(face)
304                         elif next_face in white_shades:
305                             black_shades.append(face)
306                             coloured.append(face)
307                             jump = True
308             if jump:
309                 break
310
311     return black_shades, white_shades
312
313 @staticmethod
314 def check_same_edge_region(face_1, face_2):
315     face_str_1 = str(face_1).strip("[]").replace(",□", "")
316     face_str_2 = str(face_2).strip("[]").replace(",□", "")
317     subsets = help.edges_set_creation(face_str_1)
318     for subset in subsets:
319         rev_face_str_2 = face_str_2[::-1]
320         if help sublistExists(subset, face_str_2):
321             return True, subset
322         if help sublistExists(subset, rev_face_str_2):
323             return True, subset
324         last_edge = face_str_2[len(face_str_2)-1:] + face_str_2
325             [:-1]
326         rev_last_edge = last_edge[::-1]
327         if help sublistExists(subset, last_edge):
328             return True, subset
329         if help sublistExists(subset, rev_last_edge):
330             return True, subset
331         if subset[0] == subset[1]:
332             if help sublistExists(subset[0], face_str_2):
333                 return True, subset
334
335     return False
336
337     def draw_sp_graph(self, path=None):
338         if path:
339             pass
340         else:
341             from networkx.drawing.nx_agraph import to_agraph
342
343             A = to_agraph(self.black_graph)
344             path_final = "pyknotid_helali/my_code_1/images/
345                         medial_graph_black" + str(SignedPlanarGraph.
346                                         drawing_number) + ".png"

```

```

340         A.draw( path_final , prog='dot')
341         SignedPlanarGraph.drawing_number += 1
342
343         pos = nx.planar_layout(self.white_graph)
344         A = to_agraph(self.white_graph)
345         path_final = "pyknotid_helali/my_code_1/images/
346             medial_graph_white" + str(SignedPlanarGraph.
347                 drawing_number) + ".png"
348         A.draw( path_final , prog='dot')
349         SignedPlanarGraph.drawing_number += 1
350
351     def draw_sp_graph_sign(self , path=None):
352         if path:
353             pass
354         else:
355             from networkx.drawing.nx_agraph import to_agraph
356
357             joker_graph = nx.MultiGraph()
358             joker_signs = copy.deepcopy(self.signed_black)
359
360             for sign in joker_signs:
361                 joker_graph.add_edge(sign[0][0] , sign[0][1] ,
362                     label = sign[1])
363
364             A = to_agraph(joker_graph)
365             path_final = "pyknotid_helali/my_code_1/images/
366                 medial_graph_black" + str(SignedPlanarGraph.
367                     drawing_number) + ".png"
368             A.draw( path_final , prog='dot')
369             SignedPlanarGraph.drawing_number += 1
370
371             joker_graph = nx.MultiGraph()
372             joker_signs = copy.deepcopy(self.signed_white)
373
374             for sign in joker_signs:
375                 joker_graph.add_edge(sign[0][0] , sign[0][1] ,
376                     label = sign[1])
377
378             A = to_agraph(joker_graph)
379             path_final = "pyknotid_helali/my_code_1/images/
380                 medial_graph_white" + str(SignedPlanarGraph.
381                     drawing_number) + ".png"
382             A.draw( path_final , prog='dot')
383             SignedPlanarGraph.drawing_number += 1
384
385     def sign_sp(self):
386         if self.code_gauss:
387             code_gauss_split = str(self.code_gauss).split(",")

```

```

386     black_loops = copy.deepcopy(self.black_loops)
387     white_loops = copy.deepcopy(self.white_loops)
388
389     already_signed_black = []
390     for b_edge in self.black_graph.edges():
391         if b_edge in already_signed_black:
392             if b_edge[0] != b_edge[1]:
393                 continue
394         already_signed_black.append(b_edge)
395         remove_brackets = str(b_edge).strip("()").replace("",
396                                         " ")
397         nodes = remove_brackets.split(",")
398         shift = False
399         if nodes[0] != nodes[1]:
400             for char in nodes[0]:
401                 if char in nodes[1]:
402                     code_gauss_pair = []
403                     i = 0
404                     while i < len(code_gauss_splitted):
405                         gc = code_gauss_splitted[i]
406                         if gc[0] == char:
407                             the_gc = gc
408                             if i == len(code_gauss_splitted) - 1:
409                                 code_gauss_pair.append(
410                                     code_gauss_splitted[0])
411                             else:
412                                 code_gauss_pair.append(
413                                     code_gauss_splitted[i+1]))
414                         i = i + 1
415             if len(code_gauss_pair) != 2:
416                 raise Exception("Either your
417                                 algorithm doesn't work or the
418                                 Gauss code is wrong")
419
420             edges = []
421             edges.append(char + code_gauss_pair
422                         [0][0])
423             edges.append(char + code_gauss_pair
424                         [1][0])
425
426             if len(edges[0]) == 2 and edges[0][0] ==
427                 edges[0][1]:
428                 shift = True
429             if len(edges[1]) == 2 and edges[1][0] ==
430                 edges[1][1]:
431                 shift = True
432
433             faces= list(self.black_graph.nodes()) +
434                 list(self.white_graph.nodes())
435             the_face = help.get_bounded_face(edges,
436                                              faces)
437             if the_face is not None:

```

```

427         if the_face in self.black_graph.
428             nodes():
429                 if the_gc[2] == "a":
430                     if not shift:
431                         self.signed_black.append
432                             ((b_edge, "+", char))
433                     else:
434                         self.signed_black.append
435                             ((b_edge, "-", char))
436                 elif the_gc[2] == "c":
437                     if not shift:
438                         self.signed_black.append
439                             ((b_edge, "-", char))
440                     else:
441                         self.signed_black.append
442                             ((b_edge, "+", char))
443             elif the_face in self.white_graph.
444                 nodes():
445                     if the_gc[2] == "a":
446                         if not shift:
447                             self.signed_black.append
448                                 ((b_edge, "-", char))
449                         else:
450                             self.signed_black.append
451                                 ((b_edge, "+", char))
452                     elif the_gc[2] == "c":
453                         if not shift:
454                             self.signed_black.append
455                                 ((b_edge, "+", char))
456                         else:
457                             self.signed_black.append
458                                 ((b_edge, "-", char))
459
460             elif nodes[0] == nodes[1]:
461                 crossed_at = None
462                 for white_loop in white_loops:
463                     if white_loop[0] in nodes[0]:
464                         crossed_at = white_loop[0]
465                         white_loops.remove(white_loop)
466                         break
467                 if crossed_at == None:
468                     raise Exception("Loop does not exist!!!"
469                         "Edge is not correct!!!")
470                 for gc in code_gauss splitted:
471                     if gc[0] == crossed_at:
472                         if gc[2] == "a":
473                             self.signed_black.append((b_edge, "+",
474                                         crossed_at))
475                         elif gc[2] == "c":
476                             self.signed_black.append((b_edge, "-",
477                                         crossed_at))
478
479             break

```

```

466
467     already_signed_white = []
468     for w_edge in self.white_graph.edges():
469         if w_edge in already_signed_white:
470             if w_edge[0] != w_edge[1]:
471                 continue
472             already_signed_white.append(w_edge)
473             remove_brackets = str(w_edge).strip("()").replace(",",
474                                         ", ")
475             nodes = remove_brackets.split(", ")
476             shift = False
477             if nodes[0] != nodes[1]:
478                 for char in nodes[0]:
479                     if char in nodes[1]:
480                         code_gauss_pair = []
481                         i = 0
482                         while i < len(code_gauss_splitted):
483                             gc = code_gauss_splitted[i]
484                             if gc[0] == char:
485                                 the_gc = gc
486                                 if i == len(code_gauss_splitted) - 1:
487                                     code_gauss_pair.append(
488                                         code_gauss_splitted[0])
489                                 else:
490                                     code_gauss_pair.append(
491                                         code_gauss_splitted[i+1])
492                                 i = i+1
493             if len(code_gauss_pair) != 2:
494                 raise Exception("Either your
495                               algorithm doesn't work or the
496                               Gauss code is wrong")
497
498             edges = []
499             edges.append(char + code_gauss_pair
500                         [0][0])
501             edges.append(char + code_gauss_pair
502                         [1][0])
503
504             if len(edges[0]) == 2 and edges[0][0] ==
505                 edges[0][1]:
506                 shift = True
507             if len(edges[1]) == 2 and edges[1][0] ==
508                 edges[1][1]:
509                 shift = True
510
511             faces= list(self.black_graph.nodes()) +
512                   list(self.white_graph.nodes())
513             the_face = help.get_bounded_face(edges,
514                                              faces)
515             if the_face is not None:
516                 if the_face in self.white_graph.
517                     nodes():
518                         if the_gc[2] == "a":

```

```

507             if not shift:
508                 self.signed_white.append
509                     ((w_edge, "+", char)
510                         )
511             else:
512                 self.signed_white.append
513                     ((w_edge, "-", char)
514                         )
515             elif the_gc[2] == "c":
516                 if not shift:
517                     self.signed_white.append
518                         ((w_edge, "-", char))
519                 else:
520                     self.signed_white.append
521                         ((w_edge, "+", char))
522             elif the_face in self.black_graph.
523                 nodes():
524                 if the_gc[2] == "a":
525                     if not shift:
526                         self.signed_white.append
527                             ((w_edge, "-", char)
528                                 )
529                     else:
530                         self.signed_white.append
531                             ((w_edge, "+", char)
532                                 )
533             elif the_gc[2] == "c":
534                 if not shift:
535                     self.signed_white.append
536                         ((w_edge, "+", char))
537                 else:
538                     self.signed_white.append
539                         ((w_edge, "-", char))
540             elif nodes[0] == nodes[1]:
541                 crossed_at = None
542                 for black_loop in black_loops:
543                     if black_loop[0] in nodes[0]:
544                         crossed_at = black_loop[0]
545                         black_loops.remove(black_loop)
546                         break
547                 if crossed_at == None:
548                     raise Exception("Loop does not exist!!!"
549                         "Edge is not correct!!!")
550                 for gc in code_gauss splitted:
551                     if gc[0] == crossed_at:
552                         if gc[2] == "a":
553                             self.signed_white.append((w_edge, "+",
554                                 crossed_at))
555                         elif gc[2] == "c":
556                             self.signed_white.append((w_edge, "-",
557                                 crossed_at))
558                         break
559             else:

```

```

545         raise Exception("Gauss_Code is Empty!!! We are flying blind!!!")
546
547     def create_and_draw_sp(self):
548         graph = help.create_planar_graph_from_gauss_code(self.
549             code_gauss)
550         self.create_nodes_from_faces(graph)
551         self.create_edges_after_nodes()
552         self.sign_sp()
553         self.draw_sp_graph()
554         self.draw_sp_graph_sign()
555
556     def create_and_sign_sp(self):
557         graph = help.create_planar_graph_from_gauss_code(self.
558             code_gauss)
559         self.create_nodes_from_faces(graph)
560         self.create_edges_after_nodes()
561         self.sign_sp()
562
563     def label_edges_sp(Agraph, signs):
564         temp_signs = copy.deepcopy(signs)
565         for data in signs:
566             edge = data[0]
567
568     def find_loops_in_sp(self):
569
570         self.black_loops = []
571         for b in self.black_graph.nodes():
572             if len(b) == 2:
573                 if b[0] == b[1]:
574                     self.black_loops.append(b)
575
576         self.white_loops = []
577         for w in self.white_graph.nodes():
578             if len(w) == 2:
579                 if w[0] == w[1]:
580                     self.white_loops.append(w)
581
582     def find_cycles_of_length(graph, desired_length):
583         no_multi_graph = SignedPlanarGraph.multigraph_to_graph(graph
584             )
585         cycles = nx.cycle_basis(no_multi_graph)
586         cycles_of_length = [cycle for cycle in cycles if len(cycle)
587             == desired_length]
588         return cycles_of_length
589
590     def multigraph_to_graph(multi_graph):
591         graph = nx.Graph()
592         for u, v, key, data in multi_graph.edges(data=True, keys=
593             True):
594             graph.add_edge(u, v)
595         return graph

```

C.7 Reidemeistermoves.py

```

1  from SignedPlanarGraph import SignedPlanarGraph
2  import copy
3  import my_helpers as help
4
5  class ReidemeisterMoves:
6
7
8      def reidemeister_main(sp_graph: SignedPlanarGraph) -> SignedPlanarGraph:
9
10         reid_I, reid_II, reid_III = True, True, True
11         joker_sp = copy.deepcopy(sp_graph)
12
13         while reid_I or reid_II or reid_III:
14             if reid_I or reid_II or reid_III:
15                 joker_sp_reid_I = ReidemeisterMoves.reidemeister_I(joker_sp)
16                 if joker_sp != joker_sp_reid_I:
17                     joker_sp = joker_sp_reid_I
18                     reid_I = False
19                     reid_II, reid_III = True, True
20                 else:
21                     reid_I = False
22
23             if reid_I or reid_II or reid_III:
24                 joker_sp_reid_II = ReidemeisterMoves.reidemeister_II(joker_sp)
25                 if joker_sp != joker_sp_reid_II:
26                     joker_sp = joker_sp_reid_II
27                     reid_II = False
28                     reid_I, reid_III = True, True
29                 else:
20                     reid_II = False
31
32             if reid_I or reid_II or reid_III:
33                 joker_sp_III = ReidemeisterMoves.reidemeister_III(joker_sp)
34                 if joker_sp != joker_sp_III:
35                     joker_sp = joker_sp_III
36                     reid_III = False
37                     reid_I, reid_II = True, True
38                 else:
39                     reid_III = False
40
41         joker_sp = SignedPlanarGraph(code_gauss = joker_sp.code_gauss)
42         joker_sp.create_and_sign_sp()
43         return joker_sp
44
45     def reidemeister_I(sp_graph: SignedPlanarGraph) -> SignedPlanarGraph:
46
47         reid_I = copy.deepcopy(sp_graph)
48         while True:
49             reid_I_new = ReidemeisterMoves.reidemeister_I_helper(reid_I)
50             if reid_I == reid_I_new:
51                 return reid_I_new
52             reid_I = reid_I_new
53
54     def reidemeister_II(sp_graph: SignedPlanarGraph) -> SignedPlanarGraph:
55         reid_II = copy.deepcopy(sp_graph)
56         while True:
57             reid_II_new = ReidemeisterMoves.reidemeister_II_helper(reid_II)
58             if reid_II == reid_II_new:
59                 return reid_II_new
60             reid_II = reid_II_new
61
62     def reidemeister_III(sp_graph: SignedPlanarGraph) -> SignedPlanarGraph:
63         reid_III = copy.deepcopy(sp_graph)
64         while True:
65             reid_III_new = ReidemeisterMoves.reidemeister_III_helper(reid_III)

```

```

66     if Reid_III_new == Reid_III:
67         return Reid_III_new
68     Reid_III = Reid_III_new
69
70     def reidemeister_I_helper(sp_graph: SignedPlanarGraph) -> SignedPlanarGraph:
71         """
72             we start by checking if the black graph has loops
73         """
74         if sp_graph.white_loops != None:
75             [crossing, ch_r_i] = ReidemeisterMoves.reidemeister_I_black(
76                 sp_graph)
77         if ch_r_i:
78             """
79                 We have found an RI pair !!!
80                 Cleaning up
81             """
82             code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
83             code_gauss_final = str(code_gauss_final)
84             code_gauss_final_splitted = code_gauss_final.split(",")
85             code_gauss_final = ""
86             code_gauss_final = []
87             for code in code_gauss_final_splitted:
88                 if crossing not in code:
89                     code_gauss_final += code + ","
90
91             code_gauss_final = code_gauss_final[ : len(code_gauss_final) -
92                                         1]
93             code_gauss_final = ''.join(code_gauss_final)
94
95             final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
96             final_sp.create_and_sign_sp()
97             return final_sp
98
99         if sp_graph.black_loops != None:
100            [crossing, ch_r_i] = ReidemeisterMoves.reidemeister_I_helper_white(
101                sp_graph)
102            if ch_r_i:
103                """
104                    We have found an RI pair !!!
105                    Cleaning up
106                """
107
108                code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
109                code_gauss_final = str(code_gauss_final)
110                code_gauss_final_splitted = code_gauss_final.split(",")
111                code_gauss_final = ""
112                for code in code_gauss_final_splitted:
113                    if crossing not in code:
114                        code_gauss_final += code + ","
115
116                code_gauss_final = code_gauss_final[ : len(code_gauss_final) -
117                                              1]
118                code_gauss_final = ''.join(code_gauss_final)
119
120                final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
121                final_sp.create_and_sign_sp()
122                return final_sp
123
124        return sp_graph
125
126    def reidemeister_I_black(sp_graph):
127        if len(sp_graph.white_loops) != 0:
128            white_loops = copy.deepcopy(sp_graph.white_loops)
129            for white_loop in white_loops:
130                crossing = white_loop[0]
131                for edge in sp_graph.black_graph.edges():
132                    if edge[0] == edge[1]:
133                        if crossing in edge[0]:

```

```

130         for signed_edge in sp-graph.signed_black:
131             if signed_edge[0] == edge:
132                 sign = signed_edge[1]
133                 for w_node in sp-graph.white_graph.nodes():
134                     if len(w_node) == 2:
135                         if w_node[0] == w_node[1]:
136                             if crossing in w_node:
137                                 if len(sp-graph.white_graph[w_node]) == 1:
138                                     for signed_opp_edge in
139                                         sp-graph.signed_white:
140                                         if w_node in
141                                             signed_opp_edge[0]:
142                                                 sign_opp =
143                                                     signed_opp_edge[1]
144                                                 if sign != sign_opp:
145                                                     return crossing,
146                                                     True
147
148             return None, False
149
150     def reidemeister_helper_I_white(sp-graph):
151         if len(sp-graph.black_loops) != 0:
152             black_loops = copy.deepcopy(sp-graph.black_loops)
153             for black_loop in black_loops:
154                 crossing = black_loop[0]
155                 for edge in sp-graph.white_graph.edges():
156                     if edge[0] == edge[1]:
157                         if crossing in edge[0]:
158                             for signed_edge in sp-graph.signed_white:
159                                 if signed_edge[0] == edge:
160                                     sign = signed_edge[1]
161                                     for b_node in sp-graph.black_graph.nodes():
162                                         if len(b_node) == 2:
163                                             if b_node[0] == b_node[1]:
164                                                 if crossing in b_node:
165                                                     if len(sp-graph.black_graph[b_node]) == 1:
166                                                         for signed_opp_edge in
167                                                             sp-graph.signed_black:
168                                                             if b_node in
169                                                                 signed_opp_edge[0]:
170                                                                 sign_opp =
171                                                                     signed_opp_edge[1]
172
173             return None, False
174
175     def reidemeister_II_helper(sp-graph: SignedPlanarGraph) -> SignedPlanarGraph:
176
177         black_reid = ReidemeisterMoves.reidemeister_II_helper_black(sp-graph)
178         if black_reid is not None:
179             return black_reid
180
181         white_reid = ReidemeisterMoves.reidemeister_II_helper_white(sp-graph)
182         if white_reid is not None:
183             return white_reid
184
185
186     def reidemeister_II_helper_black(sp-graph: SignedPlanarGraph) ->
187         SignedPlanarGraph:
188         ''
189
190         We start by checking parallel edges in the black graph.
191         and keeping note of their different signs

```

```

185     """
186     duplicates = {}
187     signs = {}
188     found_parrallel = False
189     for edge in sp_graph.signed_black:
190         if edge[0] in duplicates:
191             duplicates[edge[0]] += 1
192             signs[edge[0]] += edge[1]
193             found_parrallel = True
194         else:
195             duplicates[edge[0]] = 1
196             signs[edge[0]] = edge[1]
197
198     # Parallel Edges exist
199     if found_parrallel:
200         # We filter the edges from the dictionaries,
201         # to only have edges that have count more than 1.
202         filtered_duplicates = {edge: count for edge, count in duplicates.items()
203             if count >= 2}
204         filtered_signs = {edge: sign for edge, sign in signs.items() if edge in
205             filtered_duplicates}
206
207         for edge in filtered_duplicates:
208             sign = filtered_signs[edge]
209             # Skip parallel edges with only one type of sign.
210             if len(set(sign)) == 1:
211                 continue
212
213             parallel_edges = []
214             for signed_edge in sp_graph.signed_black:
215                 if signed_edge[0] == edge:
216                     parallel_edges.append(signed_edge)
217                 i = 0
218                 while i < len(parallel_edges):
219                     j = i + 1
220                     break_outer = False
221                     face = ""
222                     while j < len(parallel_edges):
223                         pe_1 = parallel_edges[i]
224                         pe_2 = parallel_edges[j]
225                         if pe_1[1] != pe_2[1]:
226                             face = pe_1[2] + pe_2[2]
227                             if face in sp_graph.white_graph.nodes():
228                                 face_edges = []
229                                 for w_edge in sp_graph.signed_white:
230                                     if face in w_edge[0]:
231                                         face_edges.append(w_edge)
232                                     if len(face_edges) == 2:
233                                         if face_edges[0][1] != face_edges[1][1]:
234                                             break_outer = True
235                                             break
236                                         if face[::-1] in sp_graph.white_graph.nodes():
237                                             face = face[::-1]
238                                             face_edges = []
239                                             for w_edge in sp_graph.signed_white:
240                                                 if face in w_edge[0]:
241                                                     face_edges.append(w_edge)
242                                                 if len(face_edges) == 2:
243                                                     if face_edges[0][1] != face_edges[1][1]:
244                                                         break_outer = True
245                                                         break
246                                         j += 1
247                                         if break_outer:
248                                             break
249                                         i += 1
250
251             if face is None:
```

```

251             raise Exception("Something is wrong !!")
252     if len(face) == 0:
253         raise Exception("Something is wrong !!")
254
255     """
256     Cleaning up
257     """
258
259     code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
260     code_gauss_final = str(code_gauss_final)
261     code_gauss_final_split = code_gauss_final.split(",")
262
263     code_gauss_final = []
264     for code in code_gauss_final_split:
265         if face[0] not in code and face[1] not in code:
266             code_gauss_final += code + ","
267
268     code_gauss_final = code_gauss_final[:len(code_gauss_final) - 1]
269     code_gauss_final = ''.join(code_gauss_final)
270
271     final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
272     final_sp.create_and_sign_sp()
273     return final_sp
274
275
276 def reidemeister_II_helper_white(sp_graph: SignedPlanarGraph) ->
277     SignedPlanarGraph:
278     """
279     We start by checking parallel edges in the black graph.
280     and keeping note of their different signs
281     """
282     duplicates = {}
283     signs = {}
284     found_parallel = False
285     for edge in sp_graph.signed_white:
286         if edge[0] in duplicates:
287             duplicates[edge[0]] += 1
288             signs[edge[0]] += edge[1]
289             found_parallel = True
290         else:
291             duplicates[edge[0]] = 1
292             signs[edge[0]] = edge[1]
293
294     # Parallel Edges exist
295     if found_parallel:
296         # We filter the edges from the dictionaries, to only have edges that
297         # have count more than 1.
298         filtered_duplicates = {edge: count for edge, count in duplicates.items()
299                               if count >= 2}
300         filtered_signs = {edge: sign for edge, sign in signs.items() if edge in
301                           filtered_duplicates}
302
303         for edge in filtered_duplicates:
304             sign = filtered_signs[edge]
305             # Skip parallel edges with only one type of sign.
306             if len(set(sign)) == 1:
307                 continue
308
309             parallel_edges = []
310             for signed_edge in sp_graph.signed_white:
311                 if signed_edge[0] == edge:
312                     parallel_edges.append(signed_edge)
313
314             i = 0
315             while i < len(parallel_edges):
316                 j = i + 1
317                 break_outer = False
318                 face = ""
319
320                 if parallel_edges[i][1] == parallel_edges[j][1]:
321                     if parallel_edges[i][0] < parallel_edges[j][0]:
322                         face = parallel_edges[i][1]
323                     else:
324                         face = parallel_edges[j][1]
325
326                     if face == "":
327                         break_outer = True
328
329                     if parallel_edges[i][0] < parallel_edges[j][0]:
330                         parallel_edges[i][1] = face
331                     else:
332                         parallel_edges[j][1] = face
333
334                 i += 1
335
336             if break_outer:
337                 break
338
339
340             if len(parallel_edges) == 2:
341                 if parallel_edges[0][1] == parallel_edges[1][1]:
342                     if parallel_edges[0][0] < parallel_edges[1][0]:
343                         parallel_edges[0][1] = parallel_edges[1][1]
344                     else:
345                         parallel_edges[1][1] = parallel_edges[0][1]
346
347             if len(parallel_edges) == 1:
348                 if parallel_edges[0][1] == parallel_edges[0][1]:
349                     parallel_edges[0][1] = ""
350
351             if len(parallel_edges) == 0:
352                 break
353
354             if len(parallel_edges) > 2:
355                 print(f"Warning: Found {len(parallel_edges)} parallel edges for edge {edge} in graph {sp_graph.name}. Only handling pairs of parallel edges at once. Other edges will be ignored.")
```

```

315     while j < len(parrallel_edges):
316         pe_1 = parrallel_edges[i]
317         pe_2 = parrallel_edges[j]
318         if pe_1[1] != pe_2[1]:
319             face = pe_1[2] + pe_2[2]
320             if face in sp_graph.black_graph.nodes():
321                 face_edges = []
322                 for b_edge in sp_graph.signed_black:
323                     if face in b_edge[0]:
324                         face_edges.append(b_edge)
325             if len(face_edges) == 2:
326                 if face_edges[0][1] != face_edges[1][1]:
327                     break_outer = True
328                     break
329             if face[::-1] in sp_graph.black_graph.nodes():
330                 face = face[::-1]
331                 face_edges = []
332                 for b_edge in sp_graph.signed_black:
333                     if face in b_edge[0]:
334                         face_edges.append(b_edge)
335             if len(face_edges) == 2:
336                 if face_edges[0][1] != face_edges[1][1]:
337                     break_outer = True
338                     break
339             j += 1
340             if break_outer:
341                 break
342             i += 1
343
344         if face is None:
345             raise Exception("Something is wrong!!")
346         if len(face) == 0:
347             raise Exception("Something is wrong!!")
348
349     ,
350     Cleaning up
351     ,
352
353     code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
354     code_gauss_final = str(code_gauss_final)
355     code_gauss_final_splitted = code_gauss_final.split(",")
356
357     code_gauss_final = []
358     for code in code_gauss_final_splitted:
359         if face[0] not in code and face[1] not in code:
360             code_gauss_final += code + ","
361
362     code_gauss_final = code_gauss_final[:len(code_gauss_final) - 1]
363     code_gauss_final = ''.join(code_gauss_final)
364
365     final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
366     final_sp.create_and_sign_sp()
367     return final_sp
368
369     return None
370
371     def reidemeister_III_helper(sp_graph:SignedPlanarGraph) -> SignedPlanarGraph:
372         to_remove_gc_wh = ReidemeisterMoves.reidemeister_III_helper_white(sp_graph)
373         if len(to_remove_gc_wh) != 0:
374             code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
375             code_gauss_final = str(code_gauss_final)
376             code_gauss_final_splitted = code_gauss_final.split(",")
377
378             code_gauss_final = []
379             for code in code_gauss_final_splitted:
380                 if code[0] not in to_remove_gc_wh:
381                     code_gauss_final += code + ","
382
383             code_gauss_final = code_gauss_final[:len(code_gauss_final) - 1]

```

```

383         code_gauss_final = ''.join(code_gauss_final)
384
385         final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
386         final_sp.create_and_sign_sp()
387         return final_sp
388
389     to_remove_gc.bl = ReidemeisterMoves.reidemeister_III_helper_black(sp_graph)
390     if len(to_remove_gc.bl) != 0:
391         code_gauss_final = copy.deepcopy(sp_graph.code_gauss)
392         code_gauss_final = str(code_gauss_final)
393         code_gauss_final_split = code_gauss_final.split(',')
394
395         code_gauss_final = []
396         for code in code_gauss_final_split:
397             if code[0] not in to_remove_gc.bl:
398                 code_gauss_final += code + ","
399
400         code_gauss_final = code_gauss_final[:len(code_gauss_final)-1]
401         code_gauss_final = ''.join(code_gauss_final)
402
403         final_sp = SignedPlanarGraph(code_gauss=code_gauss_final)
404         final_sp.create_and_sign_sp()
405         return final_sp
406
407     return sp_graph
408
409 def reidemeister_III_helper_black(sp_graph:SignedPlanarGraph) ->
410     SignedPlanarGraph:
411     black_cycles = ReidemeisterMoves.reidemeister_III_find_cycles(sp_graph, "black")
412     white_stars = []
413
414     for bl_cycle in black_cycles:
415         found = False
416         opp_face = bl_cycle[1][0]
417         for white_face in sp_graph.white_graph.nodes():
418             if help.is_subset_both_ways(white_face, opp_face):
419                 opp_face = white_face
420                 found = True
421                 break
422             if not found:
423                 continue
424             faces_to_consider = []
425             for face in bl_cycle[0]:
426                 if len(face) < 4:
427                     faces_to_consider.append(face)
428
429             for char in opp_face:
430                 if char in faces_to_consider[0] and char in faces_to_consider[1]:
431                     remains = char
432                     break
433             faces_to_consider.append(opp_face)
434             faces_to_consider_str = "".join(faces_to_consider)
435             faces_to_consider_str = faces_to_consider_str.replace(char, "")
436             if len(set(faces_to_consider_str)) != 4:
437                 continue
438             return faces_to_consider_str
439     return ""
440
441 def reidemeister_III_helper_white(sp_graph:SignedPlanarGraph) ->
442     SignedPlanarGraph:
443     white_cycles = ReidemeisterMoves.reidemeister_III_find_cycles(sp_graph, "white")
444     black_stars = []
445
446     for wh_cycle in white_cycles:
447         found = False
448         opp_face = wh_cycle[1][0]

```

```

447     for black_face in sp_graph.black_graph.nodes():
448         if help.is_subset_both_ways(black_face, opp_face):
449             opp_face = black_face
450             found = True
451             break
452         if not found:
453             continue
454         faces_to_consider = []
455         for face in wh_cycle[0]:
456             if len(face) < 4:
457                 faces_to_consider.append(face)
458
459         for char in opp_face:
460             if char in faces_to_consider[0] and char in faces_to_consider[1]:
461                 remains = char
462                 break
463             faces_to_consider.append(opp_face)
464             faces_to_consider_str = "".join(faces_to_consider)
465             faces_to_consider_str = faces_to_consider_str.replace(char, "")
466             if len(set(faces_to_consider_str)) != 4:
467                 continue
468             return faces_to_consider_str
469     return ""
470
471 def reidemeister_III_find_cycles(sp_graph:SignedPlanarGraph, colour) -> list:
472     if colour == "black":
473         return ReidemeisterMoves.reidemeister_III_find_cycles_black(sp_graph)
474     elif colour == "white":
475         return ReidemeisterMoves.reidemeister_III_find_cycles_white(sp_graph)
476
477 def reidemeister_III_find_cycles_black(sp_graph:SignedPlanarGraph):
478     # We start by looking for cycles in the black graph
479     black_cycles = SignedPlanarGraph.find_cycles_of_length(sp_graph.black_graph,
480                 3)
481     filtered_bl_cycles = []
482     # we check the cycles to pick only with edges of different signs.
483     for cycle in black_cycles:
484         i = 0
485         edges_sign_crossing = []
486         while i < len(cycle):
487             j = (i+1) % len(cycle)
488             edges_sign_crossing += [(sgn_edge[1], sgn_edge[2]) for sgn_edge in
489                         sp_graph.signed_black
490                         if sgn_edge[0] == (cycle[i], cycle[j]) or sgn_edge[0] == (
491                             cycle[j], cycle[i])]
492             i += 1
493         crossings = []
494         signs = []
495         for sgn_edge in edges_sign_crossing:
496             signs.append(sgn_edge[0])
497             crossings.append(sgn_edge[1])
498
499         cycle_data = ("").join(crossings), ("").join(signs))
500
501         if len(cycle_data[1]) == 3 and len(set(cycle_data[1])) > 1:
502             filtered_bl_cycles.append((cycle, cycle_data))
503
504     return filtered_bl_cycles
505
506 def reidemeister_III_find_cycles_white(sp_graph:SignedPlanarGraph):
507     # We start by looking for cycles in the black graph
508     white_cycles = SignedPlanarGraph.find_cycles_of_length(sp_graph.white_graph,
509                 3)
510     filtered_wh_cycles = []
511     # we check the cycles to pick only with edges of different signs.
512     for cycle in white_cycles:
513         i = 0
514         edges_sign_crossing = []

```

```
511     while i < len(cycle):
512         j = (i+1) % len(cycle)
513         edges_sign_crossing += [(sgn_edge[1], sgn_edge[2]) for sgn_edge in
514             sp_graph.signed_white
515             if sgn_edge[0] == (cycle[i], cycle[j]) or sgn_edge[0] == (
516                 cycle[j], cycle[i])]
517             i += 1
518
519         crossings = []
520         signs = []
521         for sgn_edge in edges_sign_crossing:
522             signs.append(sgn_edge[0])
523             crossings.append(sgn_edge[1])
524
525         cycle_data = ("").join(crossings), ("").join(signs))
526
527         if len(cycle_data[1]) == 3 and len(set(cycle_data[1])) > 1:
528             filtered_wh_cycles.append((cycle, cycle_data))
529
530     return filtered_wh_cycles
```