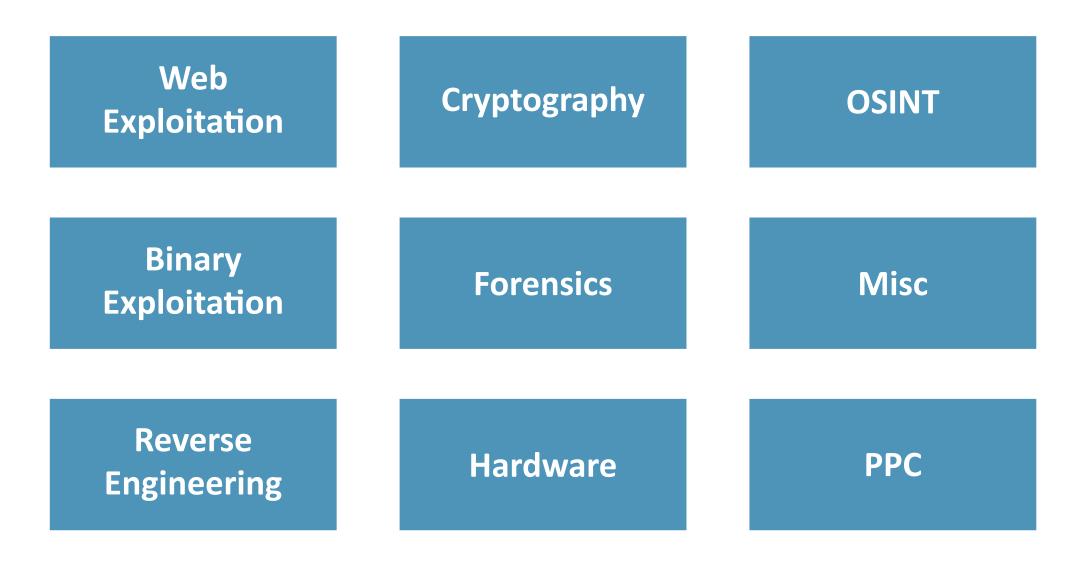
Pedagogy behind crypto CTF challenges and a way to improve them

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What's CTF?

- Capture The Flag is a competition for cybersecurity field.
- Problems are solved obtaining *flags* (words with specific format) and they are given when we broke a weak protocol, corrupted file or system for it.
- Problem categories:



Importance

- Practical experience and exposure to real-world vulnerabilities or attacks.
- Teamwork and community across social websites, jobs and universities. A way to combine all the worlds of learning (self learning, academy, industry).
- Worldwide competitions: *DEF CON, ICC, ECSC, Pwn20wn, Google CTF*.
- Gamification way to learn and teach:
 - To improve playing in cybersecurity.
 - Discovering new attacks or ways to think problems.
 - Mixing theory, practice and tricks.
- Used in a lot of universities for competitions, teach and generate more interest in cybersecurity field.

Complexity for a problemsetter

- Find interesting topics for problems.
- Being creative and make it useful to learn.
- Make it neither too easy nor too difficult.
- Feasible to solve in the way we thought and, also, ensure it as unique to test specific objective skills for evaluation.
- Avoid being solved only with pre-existent automated tools or LLMs (as with ico problem of DEFCON finals 2025).

How would you create a chall in your study field?

- CTF challs are a useful way to learn and teach each area of study.
- Do you want to try creating one for teaching or share with partners?
- Tell me more! I'll be doing feedback until 31/10 (form closes on 10).



Crypto CTF challenge example: Samsara

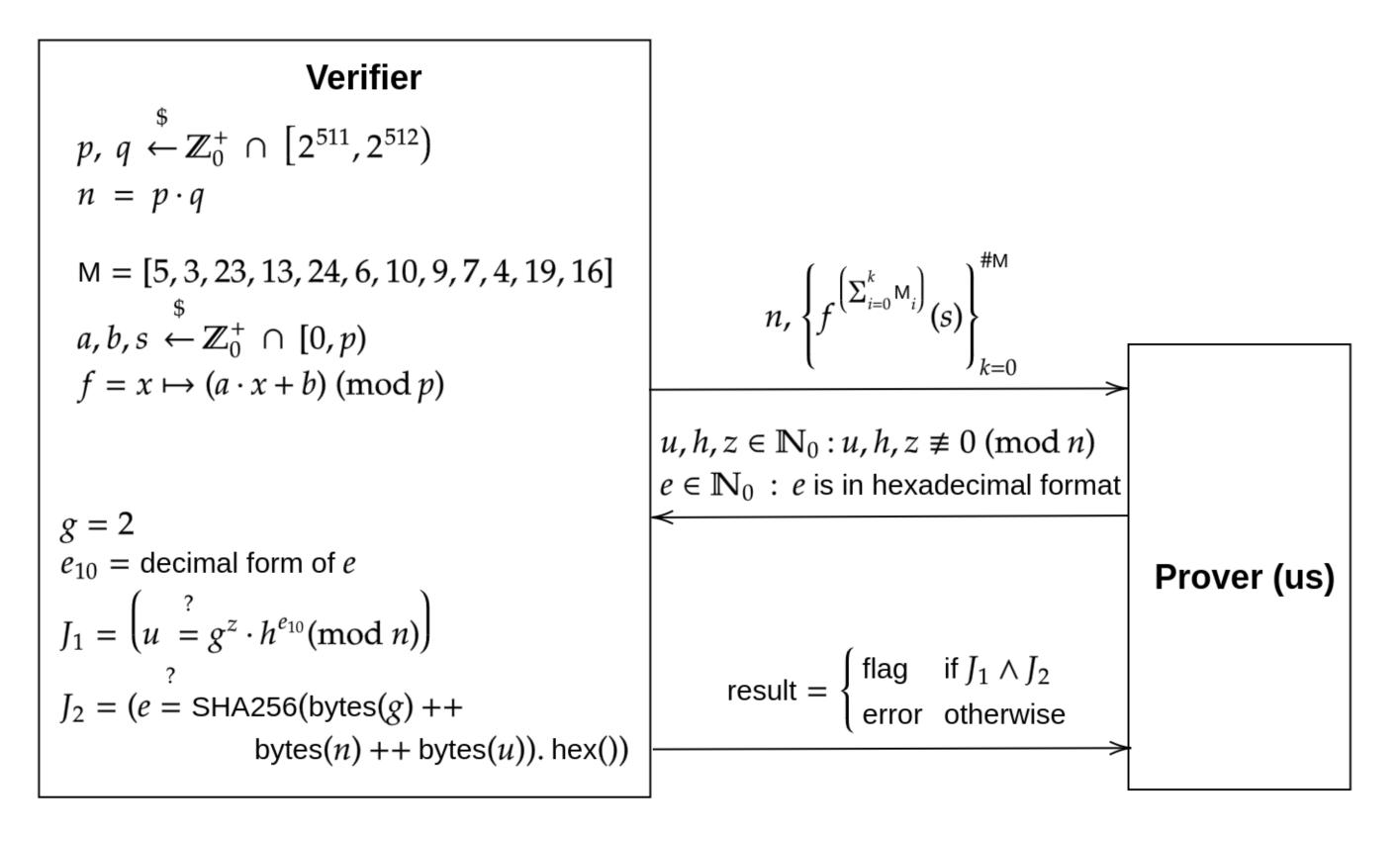
International Cybersecurity Challenge 2024

(Santiago, Chile)

Topics: Linear Congruential Generator, Non-interactive Girault's proof of knowledge (with Fiat-Shamir transformation) and Frozen Heart Vulnerability.

Statement

Break the following protocol to capture the flag:



Solution idea: We've two main parts. Let see each one separately.

First part (Linear Congruential Generator)

- LCG with unknown parameters and are given non consecutive terms of this sequence (with its positions).
- Let be $s_k = f^{(k)}(s)$ the terms we have. An useful property to use is $s_A s_B \equiv (a^A a^B) \left(s + \frac{b}{a-1} \right) \pmod{p}$.
- Let T = pX, S = pY, then $gcd(S, T) = p \iff gcd(X, Y) = 1$. Since the probability that two random numbers are coprime is 61%, we've a good chance to get p with a little quantity of (T, S) tries.
- Aiming $R: R \neq 0 \land R \equiv 0 \pmod{p}$, we'll search for positions values $A, B, C, D, E, F, G, H \in \left\{\sum_{i=0}^k M_i\right\}_{k=0}^{\#M}$ such that:

$$0 \equiv (s_A - s_B)(s_C - s_D) - (s_E - s_F)(s_G - s_H) \equiv B^2 K \pmod{p}$$

where
$$B = s + \frac{b}{a-1}$$
, $K = (a^A - a^B)(a^C - a^D) - (a^E - a^F)(a^G - a^H)$.

■ Since $B \equiv 0 \pmod{p}$ is unlikely to happen, we aim for p|K. For that:

$${A+C, B+D} = {E+G, F+H} \land {A+D, B+C} = {E+H, F+G}$$

Let m=#M+1. The search can be done with bruteforce in $O(m^8)$, but searching intelligently in $O(m^6)$. The total number of useful combinations is 288 and we're able to get p and also $q=\frac{n}{p}$.

Second part (Girault's identification non-interactive protocol)

- h isn't used in hash and it's a weakness (Frozen Heart Vulnerability).
- Choosing random values u, z, we've $e = \operatorname{Hash}(g, N, u)$ and we aim for $h \equiv (u \cdot g^{-z})^{e_{\text{inv}}} \pmod{n}$ where $e_{\text{inv}} \equiv e^{-1} \pmod{\phi(n)}$ because:

$$g^z \cdot h^e \equiv g^z \cdot \left((u \cdot g^{-z})^{e_{\mathsf{inv}}} \right)^e \equiv g^z \cdot u \cdot g^{-z} \equiv u \pmod{n}$$

- We can compute e_{inv} easily with *Extended Euclidean algorithm* because we know $\phi(n)$. If it doesn't exist, try another u, z random values.
- Sending these values to verifier, it should give us the flag message.

Complete solution and implementation



github.com/helcsnewsxd/latincrypt-2025