

Secure Computation - II

CS 601.642/442 Modern Cryptography

Fall 2022

Securely Computing *any* Function

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Solution: Using Yao's *garbled circuits* with OT

Circuit Representation

Function $f(x, y)$ can be written as a boolean circuit C :

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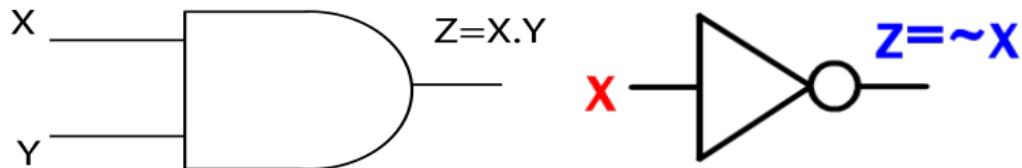
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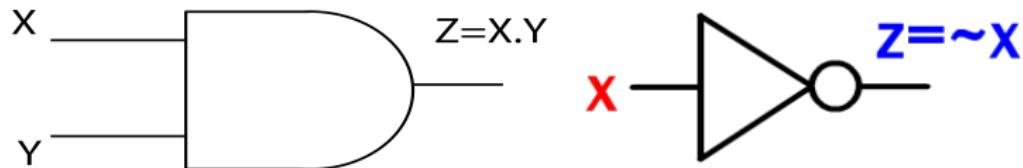
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- *Output*: Output wires of C correspond to output of $f(x, y)$

Garbled Circuits

A Garbling Scheme consists of two procedures (**Garble**, **Eval**):

- **Garble**(C): Takes as input a circuit C and outputs a collection of garbled gates \hat{G} and garbled input wires \hat{In} where

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- **Eval**(\hat{G}, \hat{In}_x): Takes as input a garbled circuit \hat{G} and garbled input wires \hat{In}_x corresponding to an input x and outputs $z = C(x)$

Garbled Circuits: Ideas

- Each wire i in the circuit C is associated with two keys (k_0^i, k_1^i) of a secret-key encryption scheme, one corresponding to the wire value being 0 and other for wire value being 1

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- We want the evaluator to use the input wire keys and the encrypted truth tables to “uncover” a single key k_v^i for every internal wire i corresponding to the value v of that wire. However, k_{1-v}^i should remain hidden from the evaluator

Special Encryption Scheme

Special Encryption Scheme: We need a secret-key encryption scheme $(\text{Gen}, \text{Enc}, \text{Dec})$ with an extra property: there exists a negligible function $\nu(\cdot)$ s.t. for every n and every message $m \in \{0, 1\}^n$,

$$\Pr[k \leftarrow \text{Gen}(1^n), k' \leftarrow \text{Gen}(1^n), \text{Dec}_{k'}(\text{Enc}_k(m)) = \perp] > 1 - \nu(n)$$

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Construction: Modify the secret-key encryption scheme discussed earlier in the class s.t. instead of encrypting m , we encrypt $0^n \| m$. Upon decrypting, check if the first n bits of the message are all 0’s; if not, then output \perp .

Garbled Circuits: Construction

Let $(\text{Gen}, \text{Enc}, \text{Dec})$ be a special encryption scheme. Assign an index to each wire in C s.t. the input wires have indices $1, \dots, n$.

Garble(C):

- For every non-output wire i in C , sample $k_0^i \leftarrow \text{Gen}(1^n)$, $k_1^i \leftarrow \text{Gen}(1^n)$. For every output wire i in C , set $k_0^i = 0$, $k_1^i = 1$.
- For every $i \in [n]$, set $\text{in}_i = (k_0^i, k_1^i)$. Set $\text{In} = (\text{in}_1, \dots, \text{in}_n)$
- For every gate g in C with input wires (i, j) , output wire ℓ :

First Input	Second Input	Output
k_0^i	k_0^j	$z_1 = \text{Enc}_{k_0^i}(\text{Enc}_{k_0^j}(k_{g(0,0)}^\ell))$
k_0^i	k_1^j	$z_2 = \text{Enc}_{k_0^i}(\text{Enc}_{k_1^j}(k_{g(0,1)}^\ell))$
k_1^i	k_0^j	$z_3 = \text{Enc}_{k_1^i}(\text{Enc}_{k_0^j}(k_{g(1,0)}^\ell))$
k_1^i	k_1^j	$z_4 = \text{Enc}_{k_1^i}(\text{Enc}_{k_1^j}(k_{g(1,1)}^\ell))$

Set $\hat{g} = \text{RandomShuffle}(z_1, z_2, z_3, z_4)$. Output $(\hat{\mathbf{G}} = (\hat{g}_1, \dots, \hat{g}_{|C|}), \hat{\text{In}})$

Garbled Circuits: Construction (contd.)

Think: Why is RandomShuffle necessary?

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$\text{Eval}(\hat{\mathbf{G}}, \hat{\mathbf{In}}_x)$:

- Parse $\hat{\mathbf{G}} = (\hat{g}_1, \dots, \hat{g}_{|C|})$, $\hat{\mathbf{In}}_x = (k^1, \dots, k^n)$
- Parse $\hat{g}_i = (\hat{g}_i^1, \dots, \hat{g}_i^4)$
- Decrypt each garbled gate \hat{g}_i one-by-one, in a canonical order:
 - Let k^i and k^j be the input wire keys for gate g .
 - Repeat the following for every $p \in [4]$:

$$\alpha_p = \text{Dec}_{k^i}(\text{Dec}_{k^j}(\hat{g}_i^p))$$

If $\exists \alpha_p \neq \perp$, set $k^\ell = \alpha_p$

- Let out_i be the value obtained for each output wire i . Output $\text{out} = (\text{out}_1, \dots, \text{out}_n)$

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- **Solution:** A will transmit the garbled wire keys corresponding to B 's input using Oblivious Transfer!

Secure Computation from Garbled Circuits: Details

Ingredients: Garbling scheme (**Garble**, **Eval**), 1-out-of-2 OT scheme
 $\text{OT} = (S, R)$

Common Input: Circuit C for $f(\cdot, \cdot)$

A 's input: $x = x_1, \dots, x_n$, **B 's input:** $y = y_1, \dots, y_n$

Protocol $\Pi = (A, B)$:

$A \rightarrow B$: A computes $(\hat{G}, \hat{\text{In}}) \leftarrow \text{Garble}(C)$. Parse $\hat{\text{In}} = (\hat{\text{in}}_1, \dots, \hat{\text{in}}_{2n})$ where $\hat{\text{in}}_i = (k_0^i, k_1^i)$. Set $\hat{\text{In}}_x = (k_{x_1}^1, \dots, k_{x_n}^n)$. Send $(\hat{G}, \hat{\text{In}}_x)$ to B .

$A \leftrightarrow B$: For every $i \in [n]$, A and B run $\text{OT} = (S, R)$ where A plays sender S with input (k_0^{n+i}, k_1^{n+i}) and B plays receiver R with input y_i . Let $\hat{\text{In}}_y = (k_{y_1}^{n+1}, \dots, k_{y_n}^{2n})$ be the outputs of the n OT executions received by B .

B : B outputs $\text{Eval}(\hat{G}, \hat{\text{In}}_x, \hat{\text{In}}_y)$

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Additional Reading: Read security proof from [Lindell-Pinkas'04]