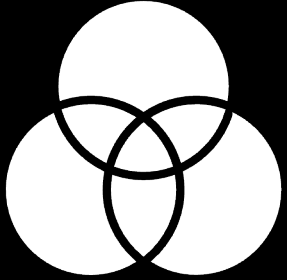


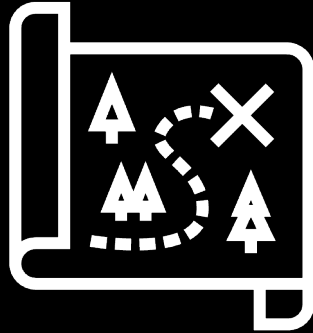
# Venn Diagrams

Week 6 . Deeper dive.

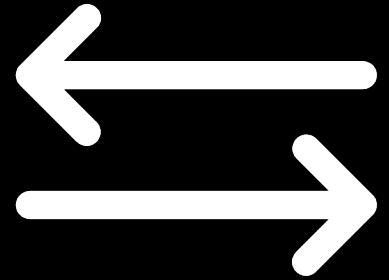
Types of models we will consider:



Venn diagrams

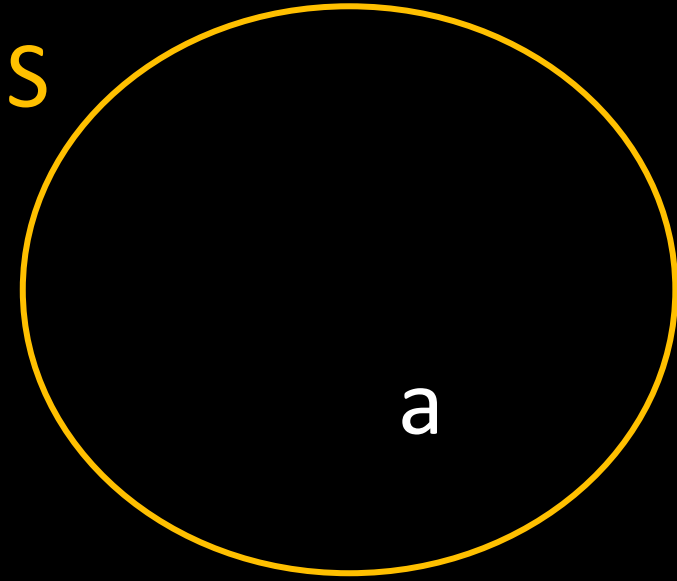


Map models



Arrow models

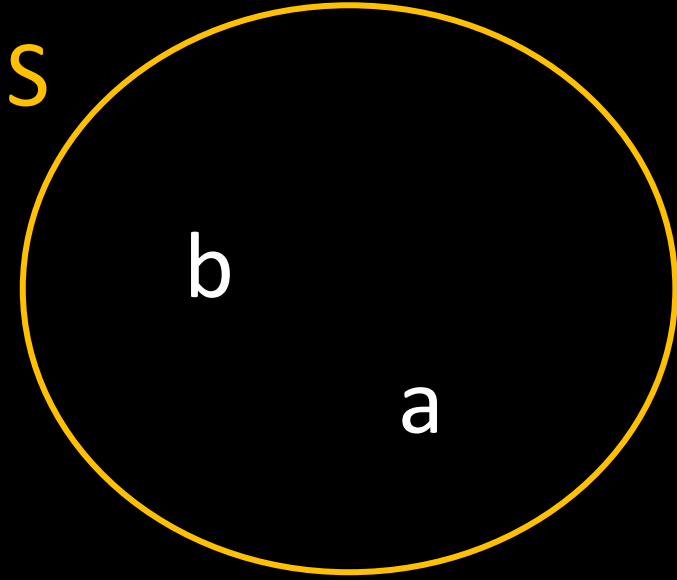
# Reading Venn diagram models



$Sa$  is true in the model

$Sa$

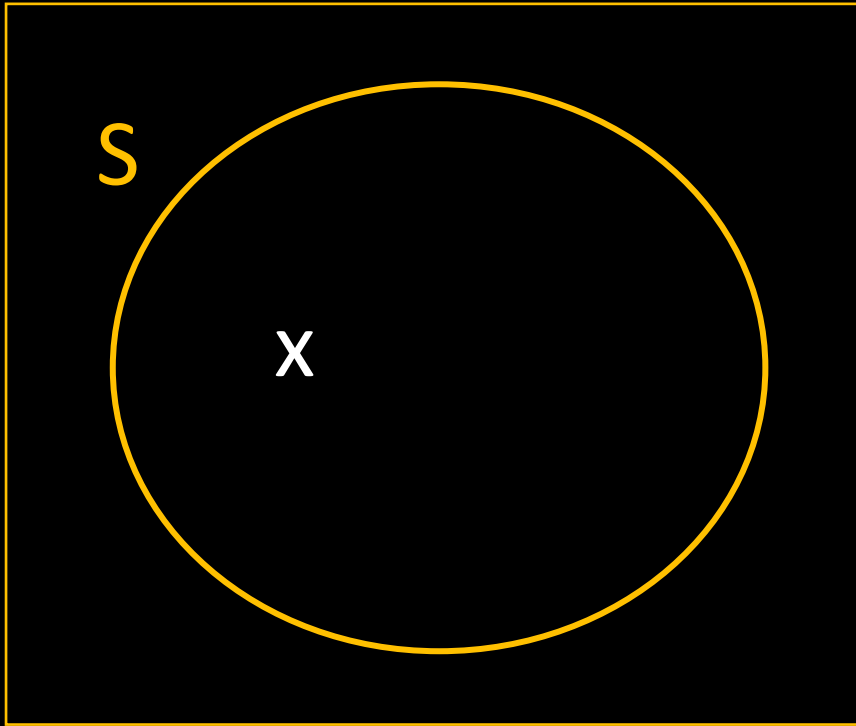
# Reading Venn diagram models



Sa and Sb are both true in the model

$$Sa \wedge Sb$$

# Reading Venn diagram models



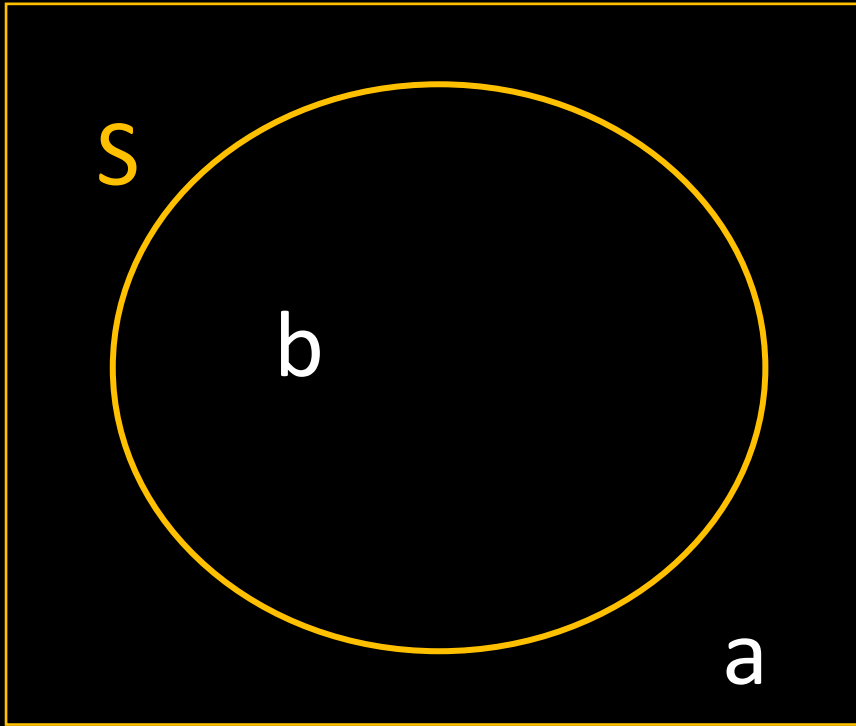
There is at least one object with property S.

$$\exists x Sx$$

Is there anything that isn't an S? unclear!

If an area is empty,  
we must not assume  
that the  
corresponding set  
has members.

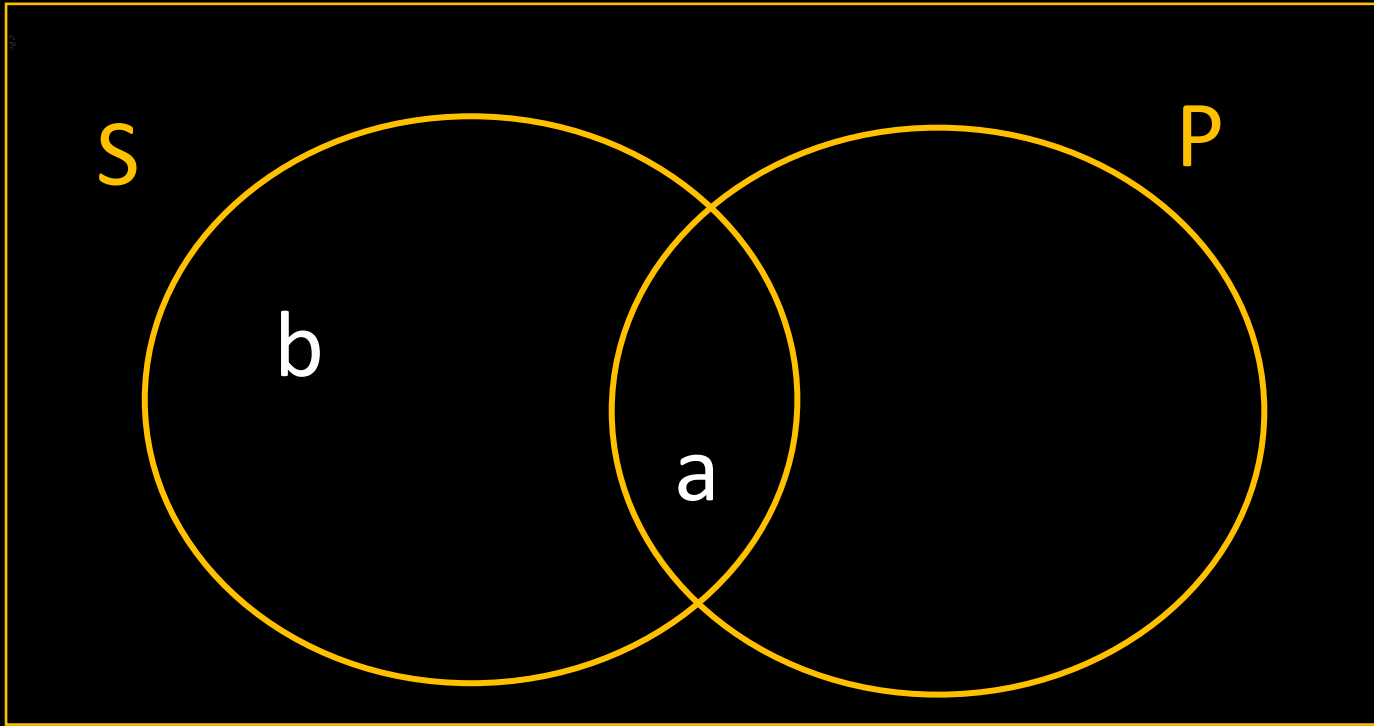
# Reading Venn diagram models



$Sb$  is true in the model but  $Sa$  is not.

$$Sb \wedge \neg Sa$$

# Reading Venn diagram models

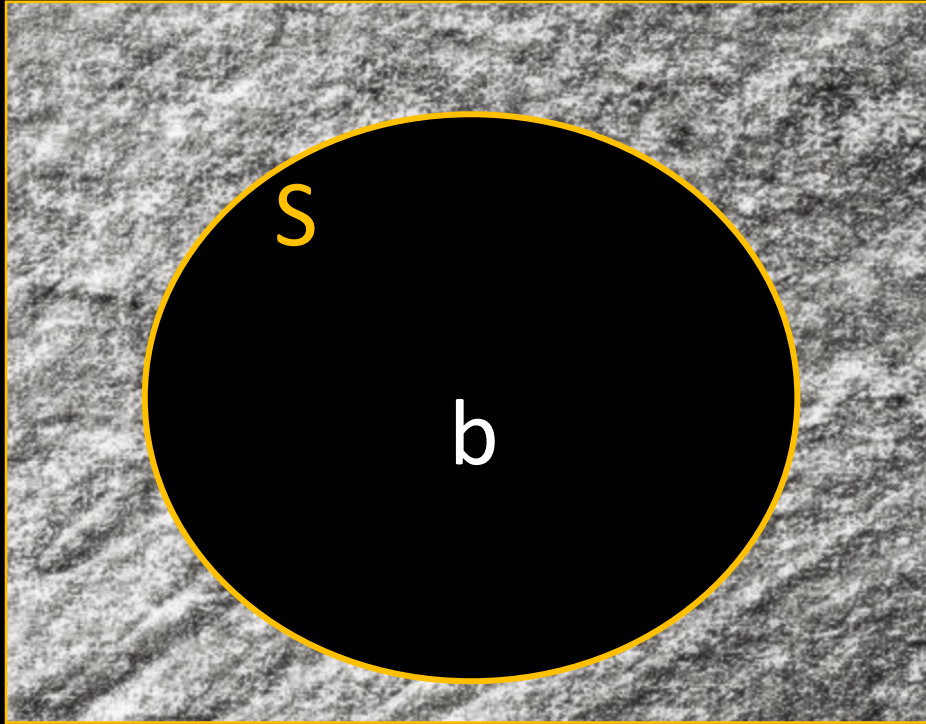


$Sb$  is true in the model.

$Sa$  and  $Pa$  are both true in the model.

$$Sb \wedge (Sa \wedge Pa)$$

# Reading Venn diagram models



$Sb$  is true in the model and there is nothing that is not an  $S$  in the model (we've "blocked out" everything else by shading it in).

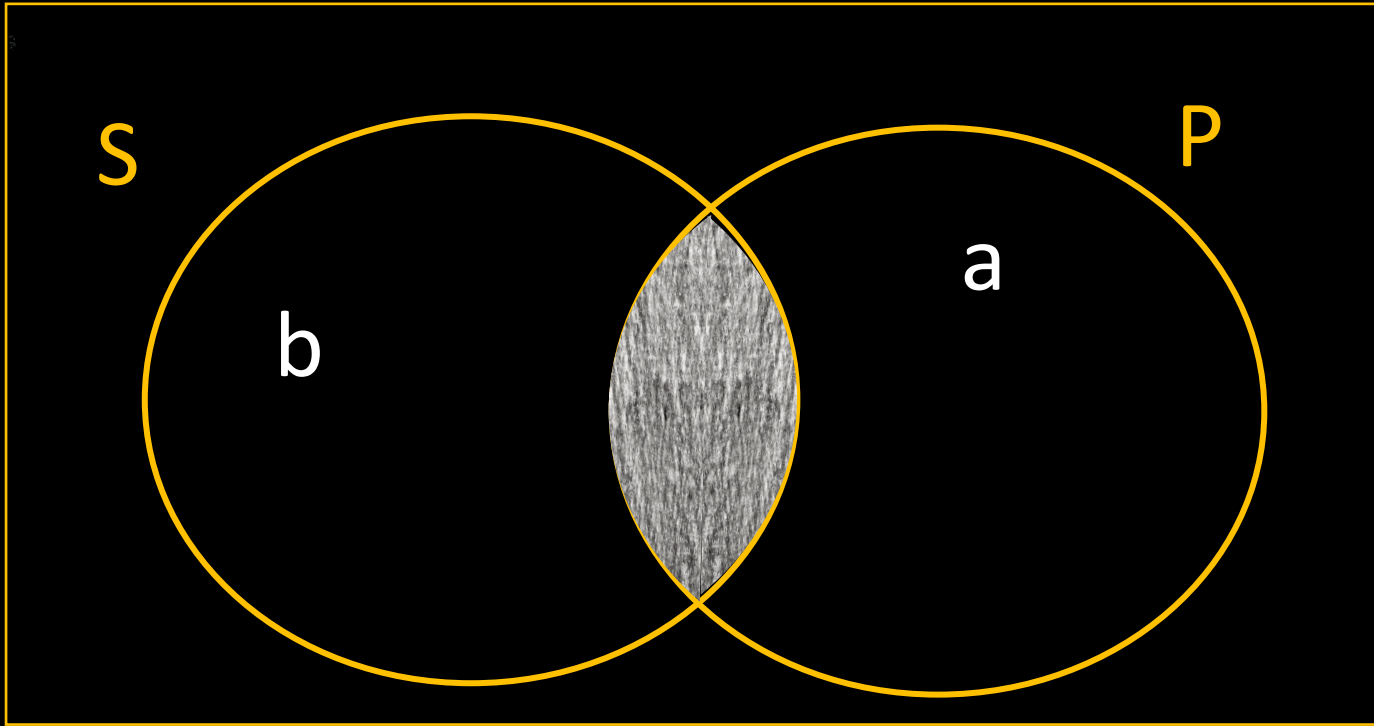
$$Sb \wedge \neg \exists x \neg Sx$$

or equivalently

$$Sb \wedge \forall x Sx$$



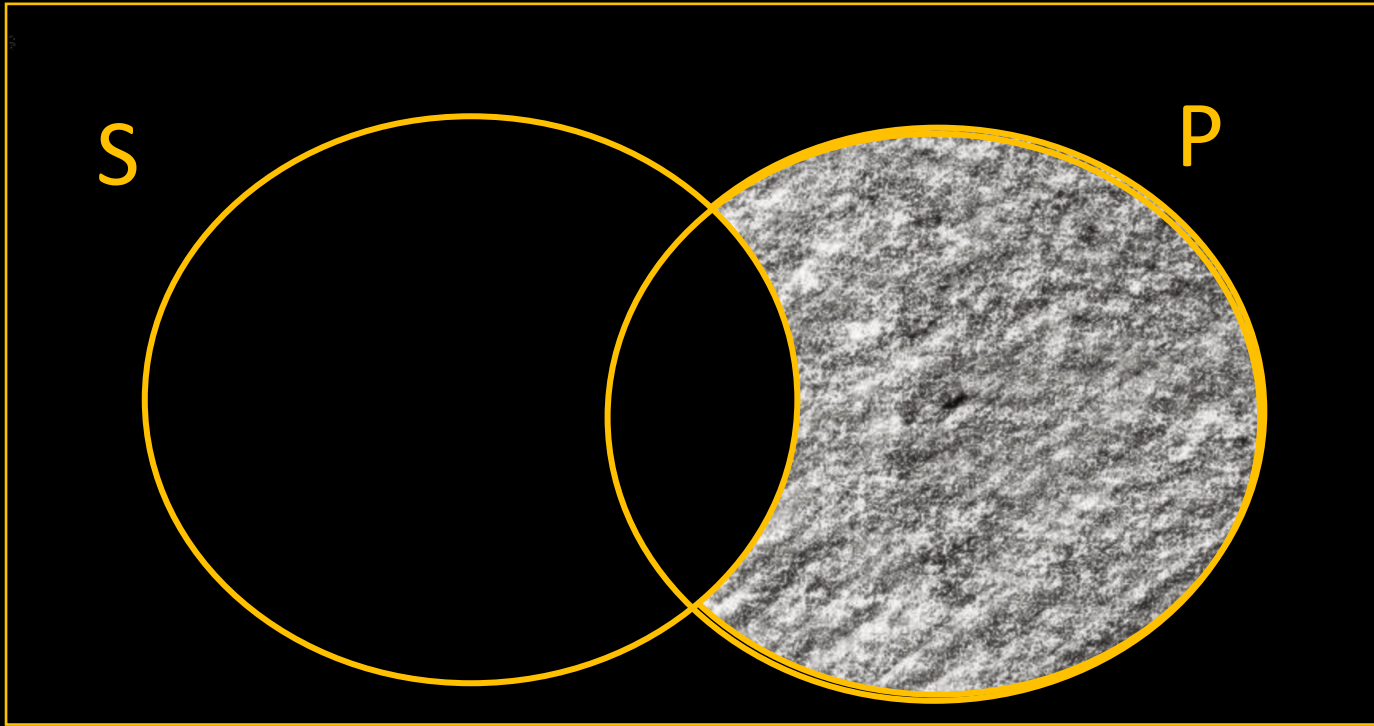
# Reading Venn diagram models



$Sb$  is true in the model,  
 $Pb$  is true in the model,  
and there is nothing that  
has both properties.

$$(Sb \wedge Pa) \wedge \neg \exists x (Sx \wedge Px)$$

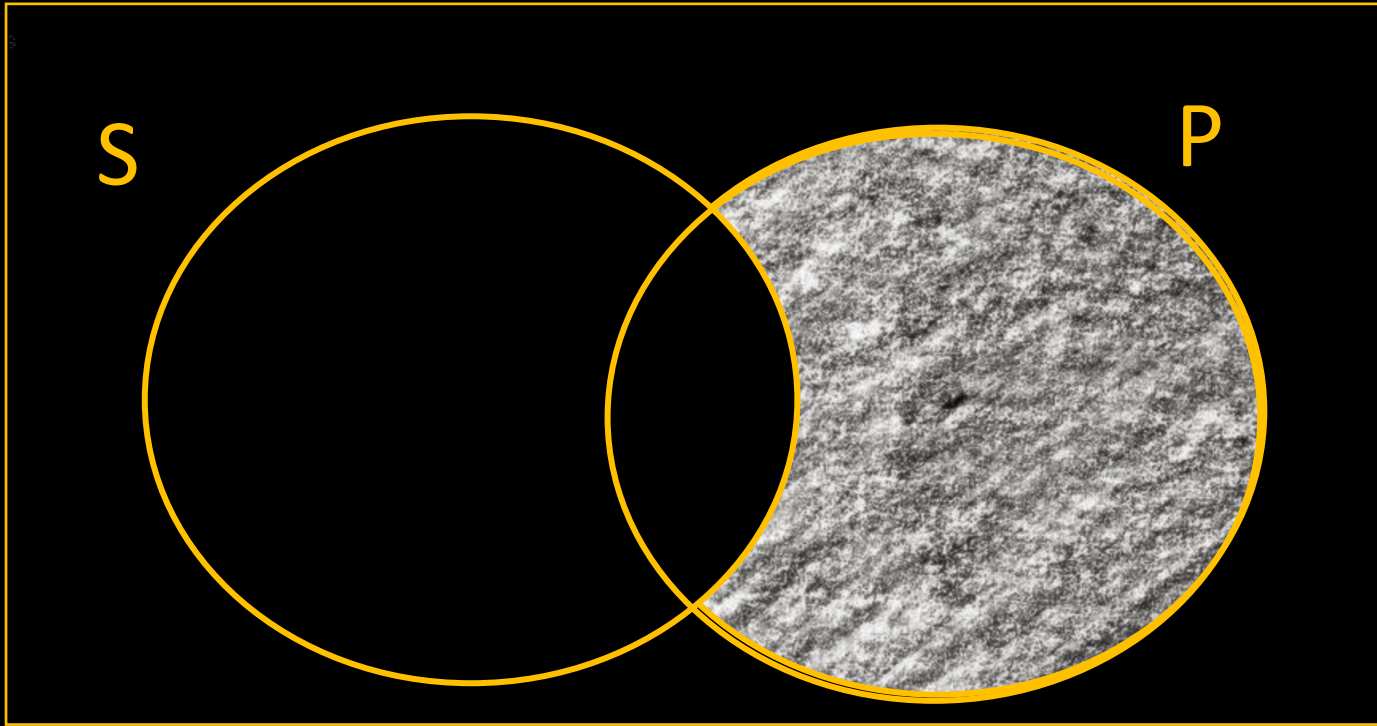
# Reading Venn diagram models



No non-S is a P.  
or, equivalently,  
All Ps are Ss.

$\forall x (Px \rightarrow Sx)$   
or, equivalently,  
 $\neg \exists x (Px \wedge \neg Sx)$

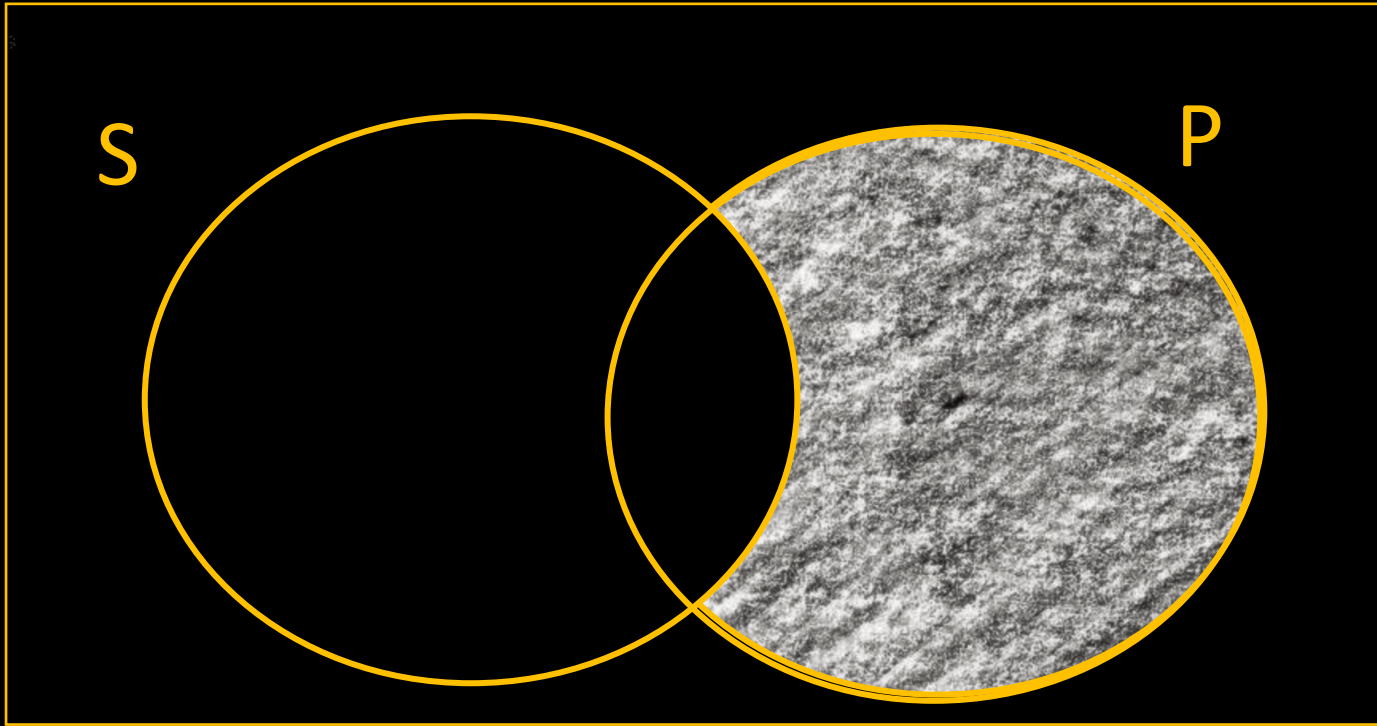
# Reading Venn diagram models



Can I say "All Ps are Ss"?

Yes. This doesn't mean that there is a P necessarily. Just that if there were one, it would also be an S!

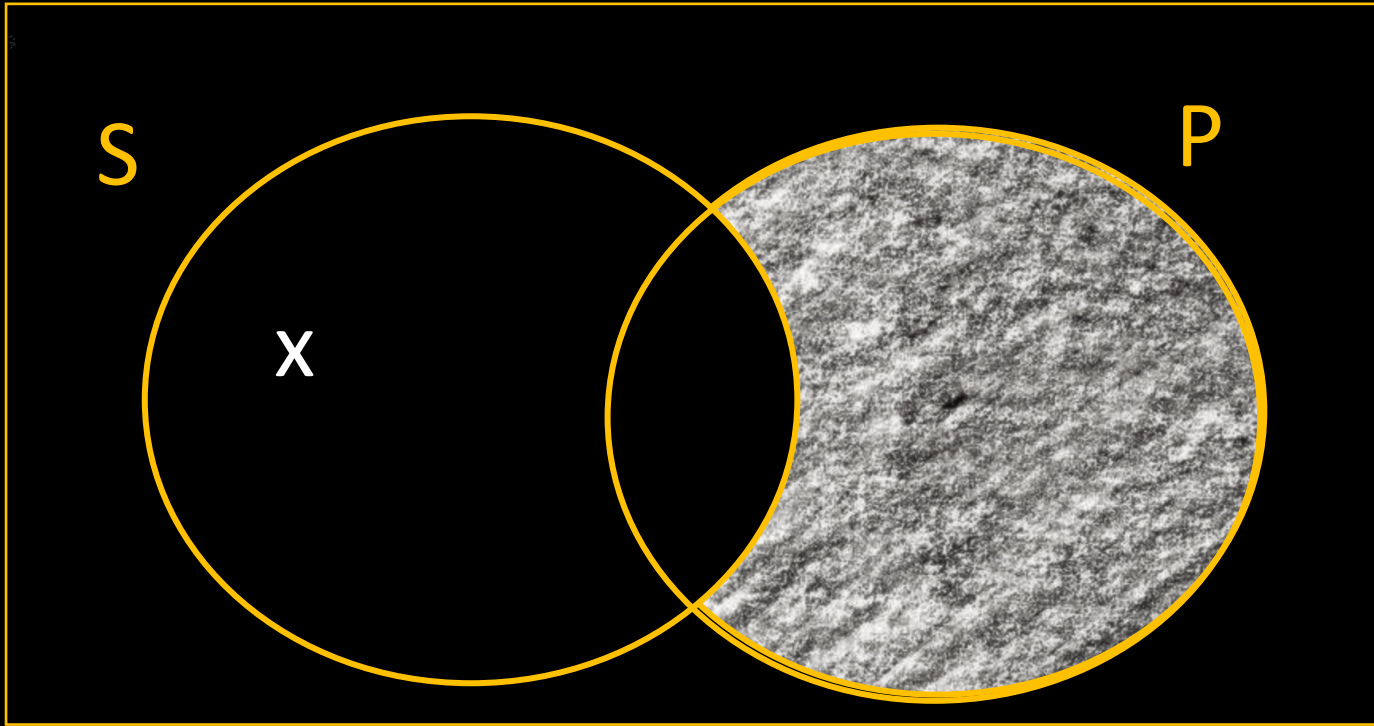
# Reading Venn diagram models



Can I say  
 $\exists x Sx$ ?

No! Not unless  
we're told that  
the set has  
members.

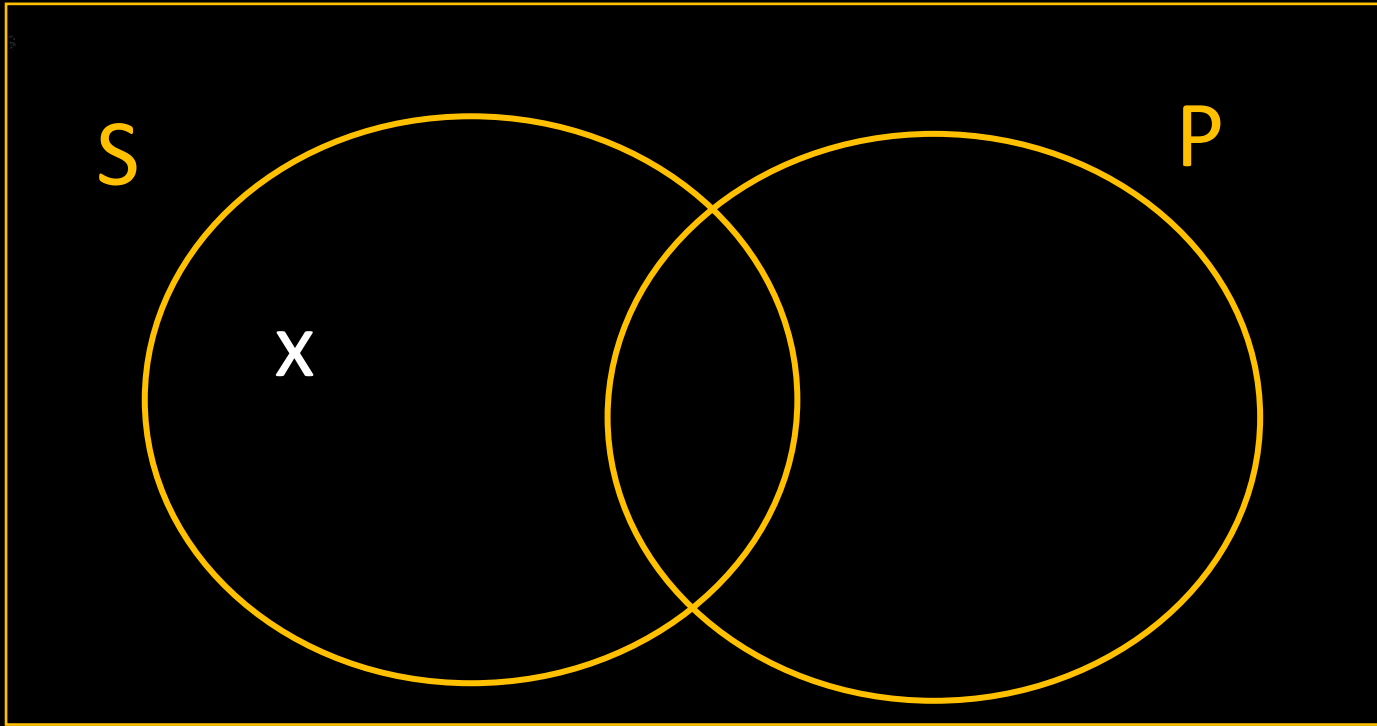
# Reading Venn diagram models



Can I say  
 $\exists x Sx$ ?

Yes! The "x"  
tells us that  
there is at least  
one object with  
property S.

Practice: Are the following true or false?



$\forall x (Px \rightarrow Sx)$  false

$\forall x (Sx \rightarrow Px)$  false

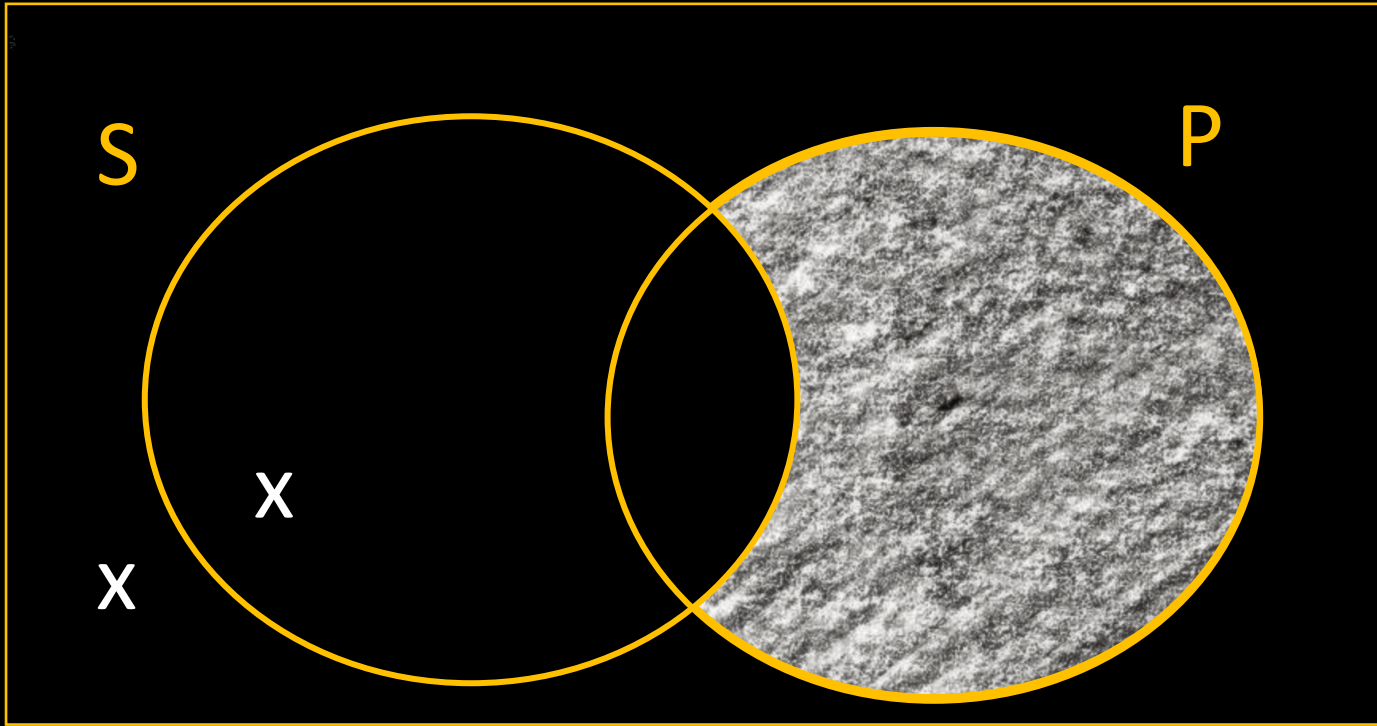
$\exists x Sx$  true

$\exists x \neg Sx$  false

$\exists x Px$  false

$\exists x \neg Px$  true

Practice: Are the following true or false?



$\exists x \neg Sx$  true

$\forall x (Sx \rightarrow Px)$  false

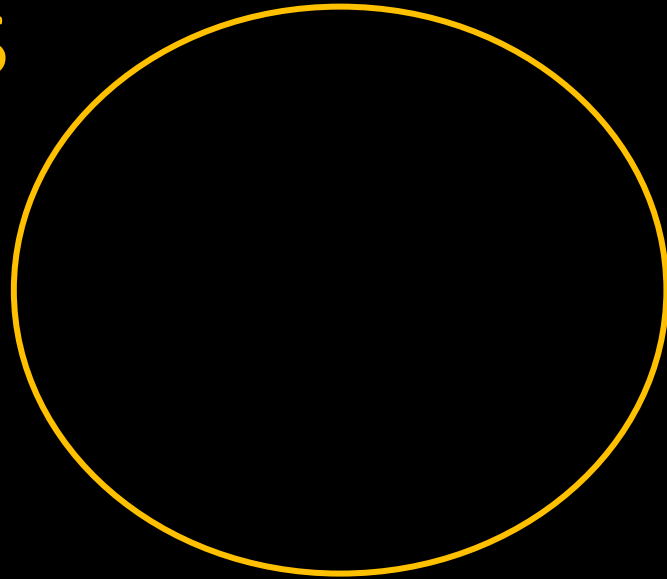
$\forall x (Px \rightarrow Sx)$  true

$\forall x (\neg Px \rightarrow Sx)$  false

$\forall x (\neg Sx \rightarrow \neg Px)$  true

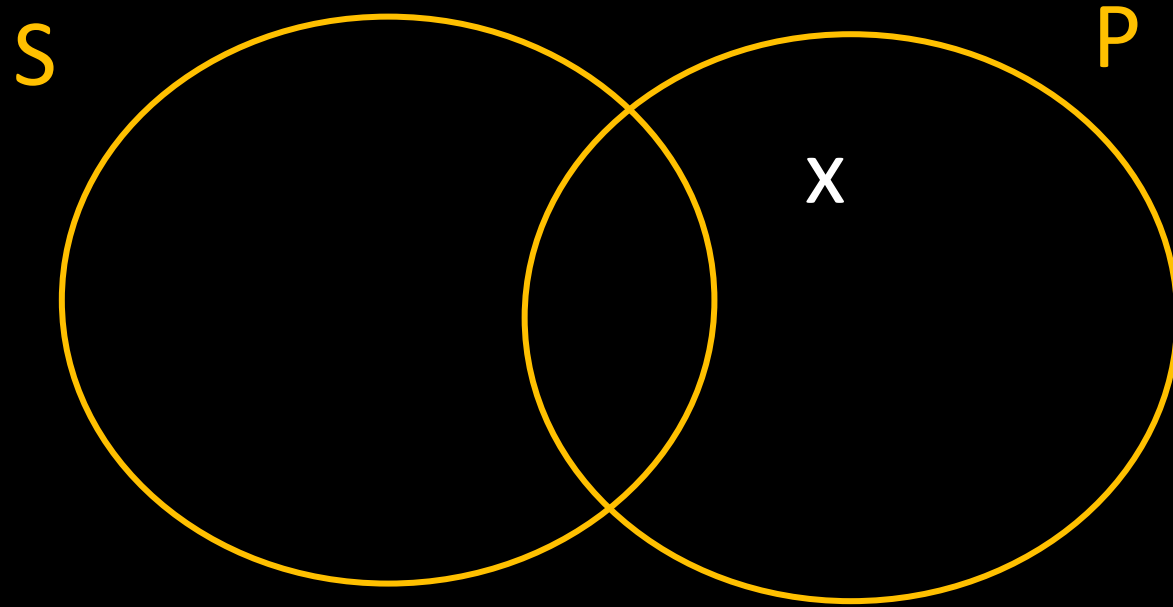
Draw a model that makes  $\exists x (Sx \wedge \neg Px)$  false:

S

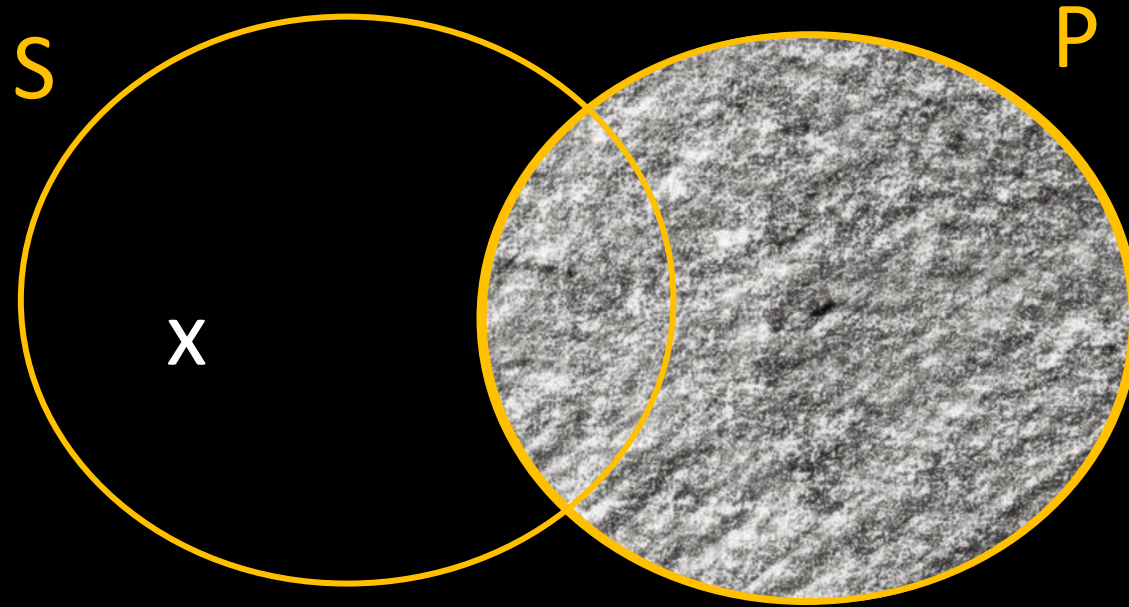




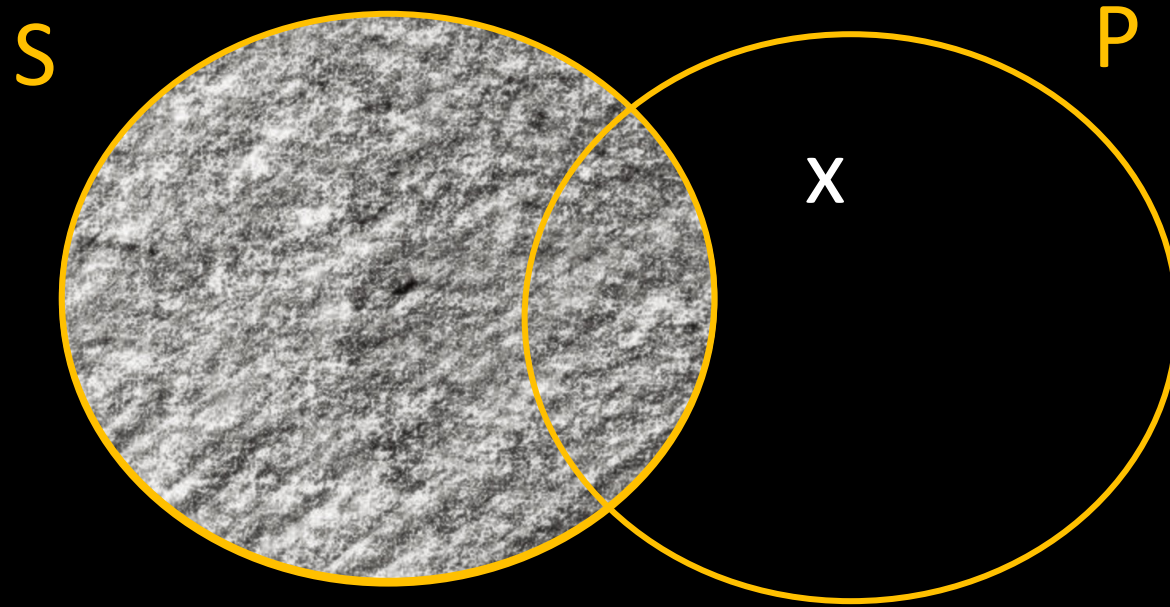
Draw a *non-empty* model that makes  $\exists x (Sx \wedge \neg Px)$  **false**:



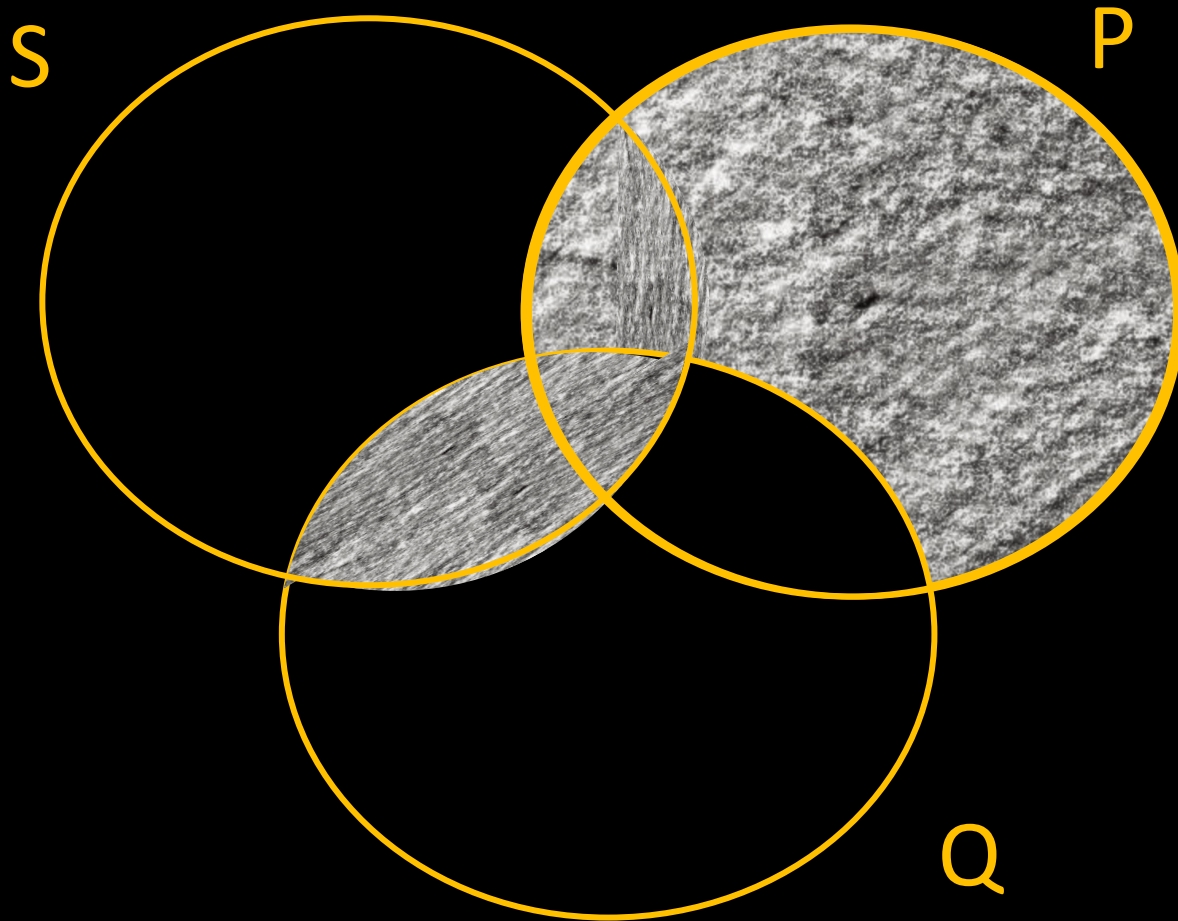
Draw a non-empty model that makes  $\forall x (Sx \wedge \neg Px)$  **true**:



Draw a non-empty model that makes  $\forall x (Sx \wedge \neg Px)$  false:



# Using Venn diagram models for arguments



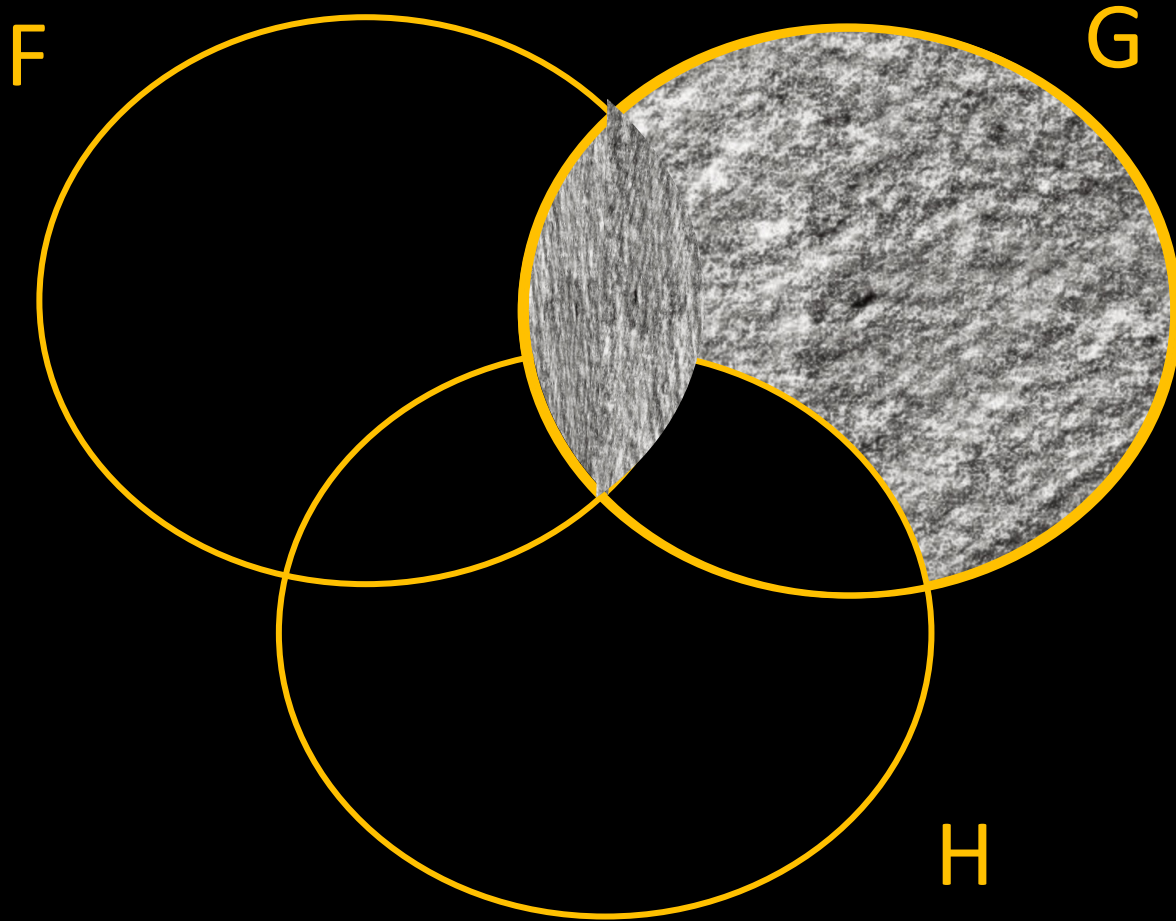
Consider the argument :

All P are Q.  
No Q are S.  
 $\therefore$  No P are S.



valid

# Using Venn diagram models for arguments



Consider the argument :

No F are G

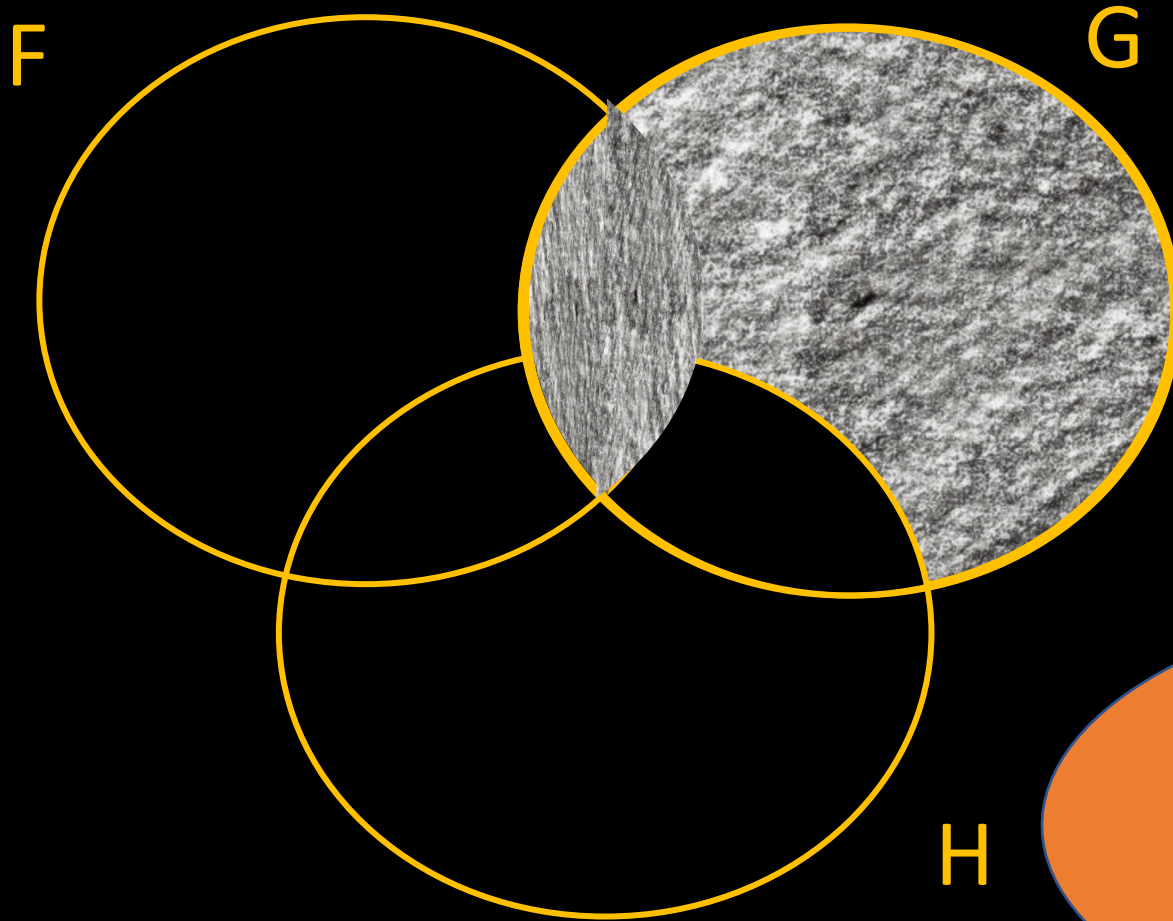
All G are H.

$\therefore$  Some F are H



invalid

# Using Venn diagram models for arguments



Consider the argument :

No F are G

All G are H.

∴ Some F are H



invalid

Why?

Recall, for the argument to be valid, the conclusion **MUST** follow. Here, we don't **have to** have any F that are H!