Basale datastrukturer, DMA uge 4

- Hvad er en datastruktur?
- Stakke
- Køer
- Hægtede lister
- Prioritetskøer
- Hobe

Mikkel Abrahamsen

Hvad er en datastruktur?

En måde at organisere data på.

Mål: Søge/tilgå/ændre effektivt, dvs. hurtigt og pladsbesparende (kompakt).

Eksempler vi har set: Array og sorteret array.

Hvad er en datastruktur?

En måde at organisere data på.

Mål: Søge/tilgå/ændre effektivt, dvs. hurtigt og pladsbesparende (kompakt).

Eksempler vi har set: Array og sorteret array.

Husk: Alle datastrukturer skal gemmes i hukommelsen, dvs. i et stort array.

	0	1	2	3									
H:													

Push(S,28)

 $\begin{array}{r}
3 \\
4 \\
18 \\
16 \\
S
\end{array}$

 $egin{array}{c} 28 \\ \hline 3 \\ \hline 4 \\ \hline 18 \\ \hline 16 \\ S \\ \hline \end{array}$

Push(S,28)

Push(S,28)Pop(S)

Push(S,28)

 $\mathsf{Pop}(S)$

3	Push(S,28)	
4	Pop(S)	returnér 28
18	Pop(S)	
$\begin{bmatrix} 16 \end{bmatrix}$		

Push(S,28)

18

16

 $\mathsf{Pop}(S)$ $\mathsf{Pop}(S)$ returnér 28

Push(S,28)

 $\mathsf{Pop}(S)$

returnér 28

 $\mathsf{Pop}(S)$

returnér 3

 $\frac{16}{S}$

18

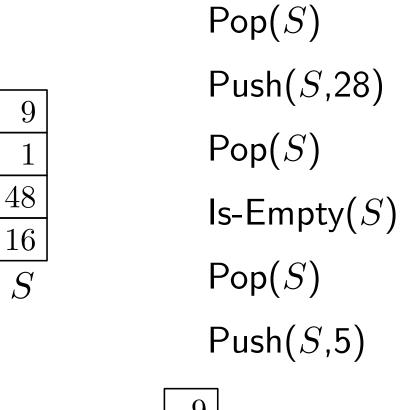
 $\mathsf{Push}(S,7)$

7	Push(S,28)	
4	Pop(S)	returnér 28
18 16	Pop(S)	returnér 3
S	Push(S,7)	

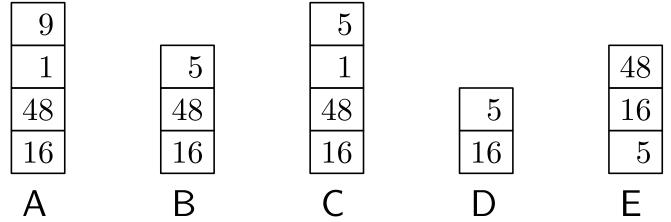
7	Push(S,28)	
4	Pop(S)	returnér 28
18 16	Pop(S)	returnér 3
S	Push(S,7)	
	Is-Empty(S)	

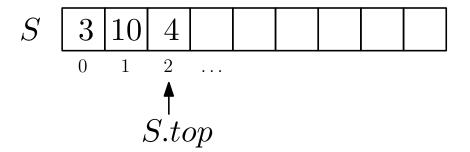
7	Push(S,28)	
4	Pop(S)	returnér 28
18 16	Pop(S)	returnér 3
S	Push(S,7)	
	Is-Empty(S)	returnér false

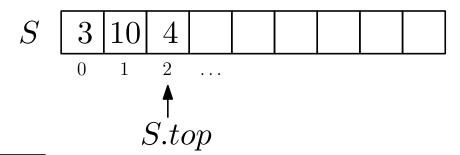
Hvordan ser stakken ud til sidst?



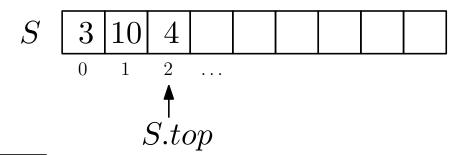
socrative.com → Student login, Room name: ABRAHAMSEN3464







Is-Empty(S)
return S.top == -1

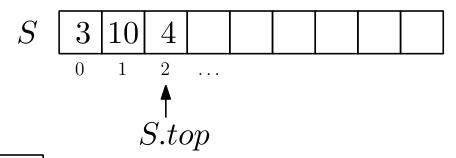


Is-Empty(
$$S$$
)
return $S.top == -1$

$$Push(S, x)$$

$$S.top = S.top + 1$$

$$S[S.top] = x$$



Is-Empty(
$$S$$
)
return $S.top == -1$

$$Push(S, x)$$

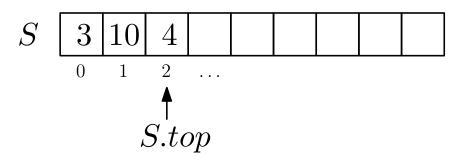
$$S.top = S.top + 1$$

$$S[S.top] = x$$

$$\mathsf{Pop}(S)$$

$$S.top = S.top - 1$$

$$\mathsf{return}\ S[S.top + 1]$$



Is-Empty(
$$S$$
)
return $S.top == -1$

$$Push(S, x)$$

$$S.top = S.top + 1$$

$$S[S.top] = x$$

Køretid: $\Theta(1)$

$$Pop(S)$$
 $S.top = S.top - 1$
 $return \ S[S.top + 1]$

Κø

 $Q \ \boxed{3 \ 15 \ 4}$



 $Q \ \boxed{3 \ 15 \ 4}$

 $\mathsf{Enqueue}(Q,1)$



$$Q \ \boxed{3} \ \boxed{15} \ \boxed{4} \ \boxed{1}$$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $Q \qquad \boxed{15 \mid 4 \mid 1 \mid 22}$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $Q \qquad \boxed{15 \mid 4 \mid 1 \mid 22}$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $\mathsf{Dequeue}(Q)$

 $Q \qquad \boxed{4 \mid 1 \mid 22}$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $\mathsf{Dequeue}(Q)$

returnér 3

 $Q \qquad \boxed{4 \mid 1 \mid 22}$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $\mathsf{Dequeue}(Q)$

Enqueue(Q, 6)

returnér 3

 $Q \qquad \boxed{4 \mid 1 \mid 22 \mid 6}$

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $\mathsf{Dequeue}(Q)$

Enqueue(Q, 6)

returnér 3

Q 4 1 22 6

 $\mathsf{Enqueue}(Q,1)$

 $\mathsf{Enqueue}(Q,22)$

 $\mathsf{Dequeue}(Q)$

 $\mathsf{Dequeue}(Q)$

Enqueue(Q, 6)

 $\mathsf{Is}\text{-}\mathsf{Empty}(Q)$

returnér 3

Q 4 1 22 6

 $\mathsf{Enqueue}(Q,1)$

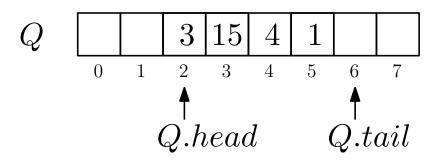
 $\mathsf{Enqueue}(Q,22)$

Dequeue(Q) returnér 3

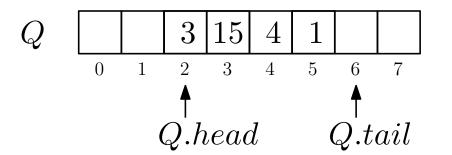
Dequeue(Q) returnér 15

Enqueue(Q, 6)

 $\mathsf{Is}\text{-}\mathsf{Empty}(Q)$ returnér false

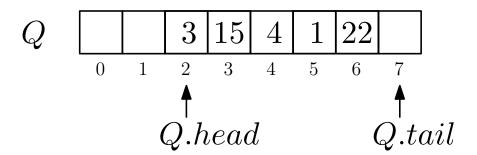


Q.N: størrelse af array



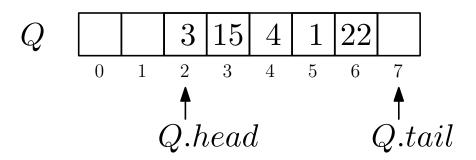
Q.N: størrelse af array

Enqueue(Q, 22)



Q.N: størrelse af array

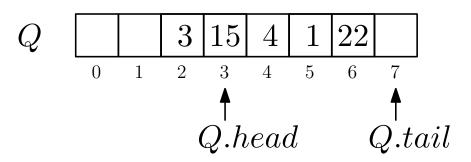
Enqueue(Q, 22)



Q.N: størrelse af array

Enqueue(Q, 22)

 $\mathsf{Dequeue}(Q)$

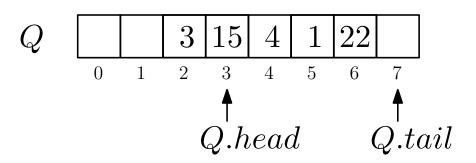


Q.N: størrelse af array

Enqueue(Q, 22)

 $\mathsf{Dequeue}(Q)$

returnér 3



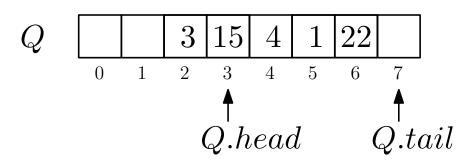
Q.N: størrelse af array

Enqueue(Q, 22)

 $\mathsf{Dequeue}(Q)$

returnér 3

```
\begin{aligned} & \mathsf{Enqueue}(Q,x) \\ & Q[Q.tail] = x \\ & \mathsf{if} \ Q.tail < Q.N-1 \\ & Q.tail = Q.tail + 1 \\ & \mathsf{else} \\ & Q.tail = 0 \end{aligned}
```



Q.N: størrelse af array

 $\mathsf{Enqueue}(Q,22)$

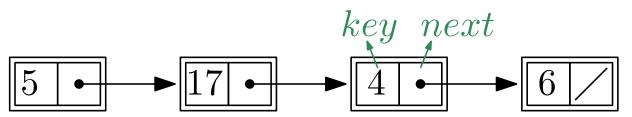
Dequeue(Q)

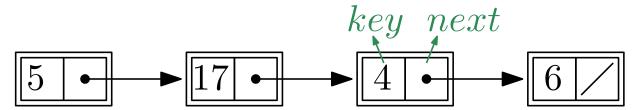
returnér 3

```
\begin{aligned} & \text{Enqueue}(Q, x) \\ & Q[Q.tail] = x \\ & \text{if } Q.tail < Q.N-1 \\ & Q.tail = Q.tail + 1 \\ & \text{else} \\ & Q.tail = 0 \end{aligned}
```

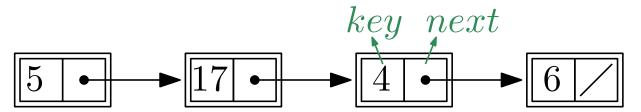
 $\begin{aligned} & \text{Dequeue}(Q) \\ & x = Q[Q.head] \\ & \text{if } Q.head < Q.N-1 \\ & Q.head = Q.head+1 \\ & \text{else} \\ & Q.head = 0 \\ & \text{return } x \end{aligned}$

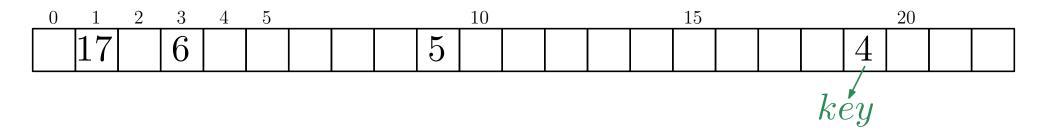


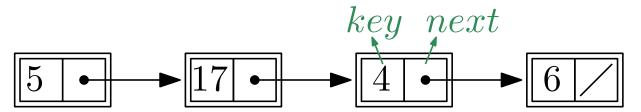


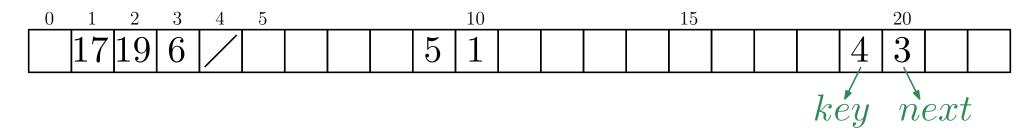


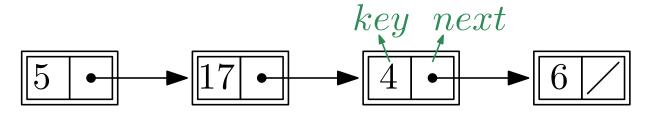
0	1	2	3	4	5		10					15					20				

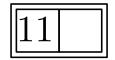


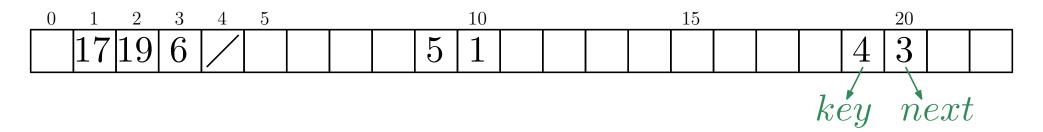


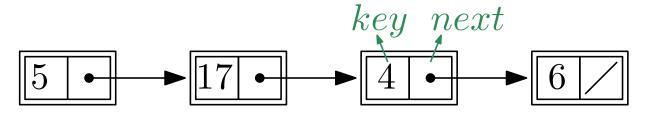


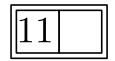


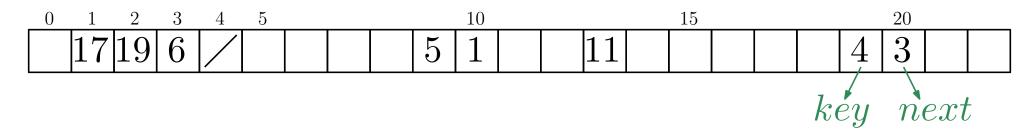


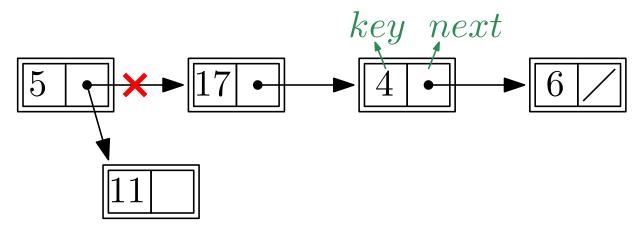


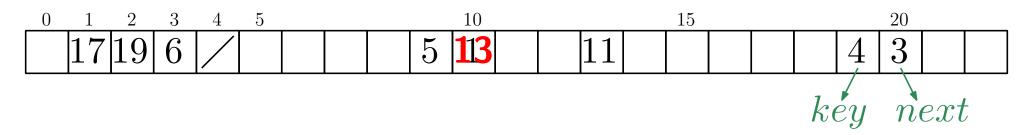


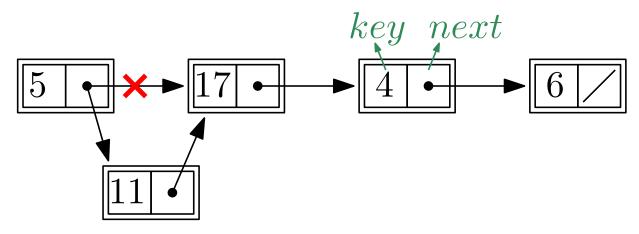


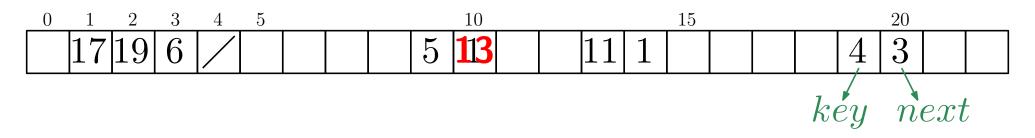


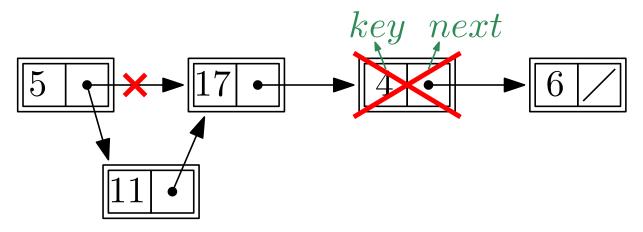


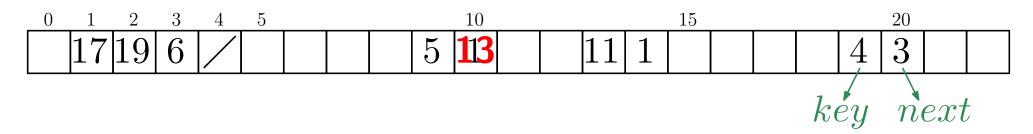


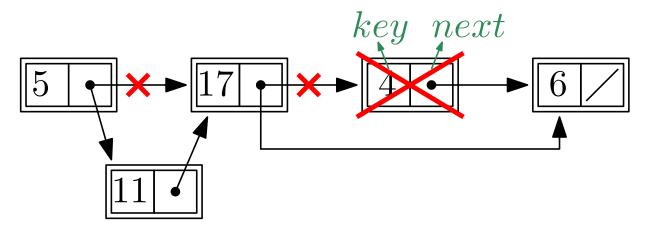


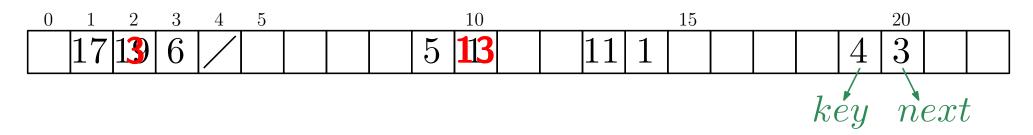


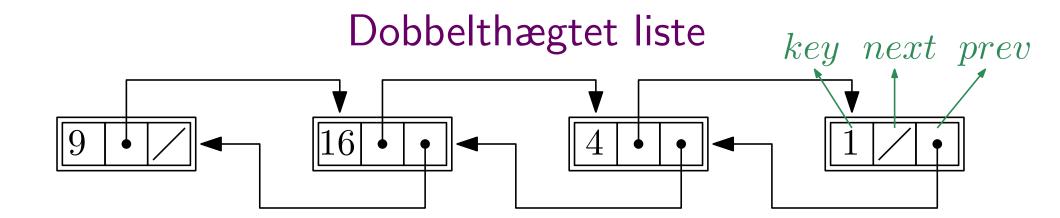


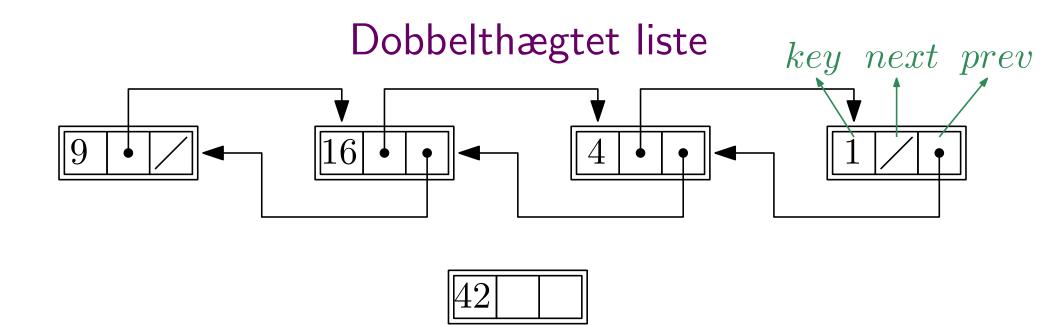


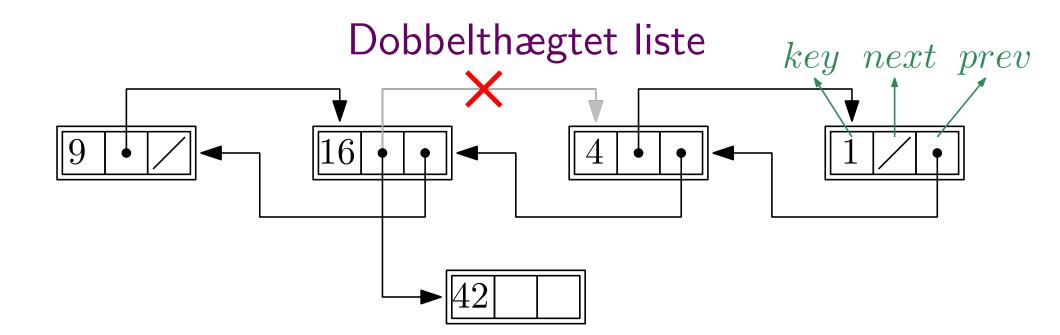


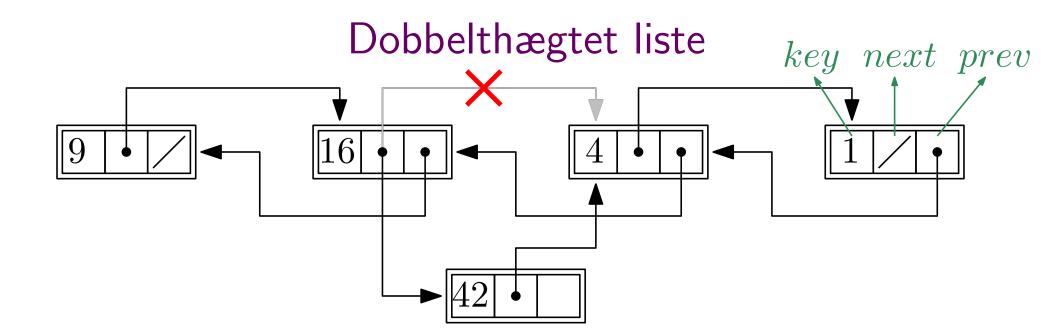


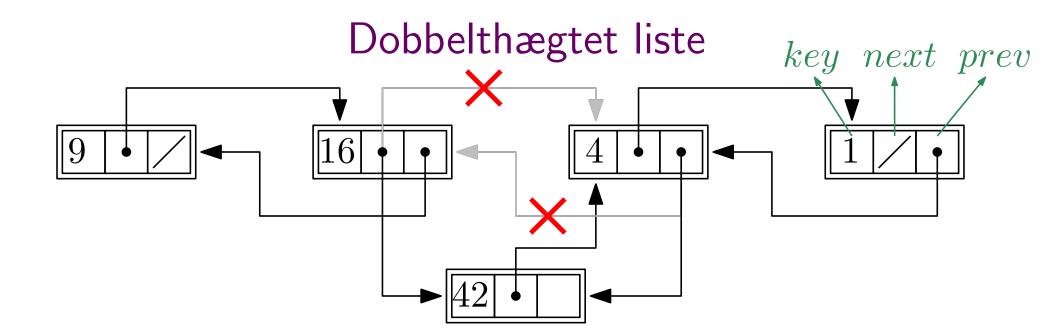


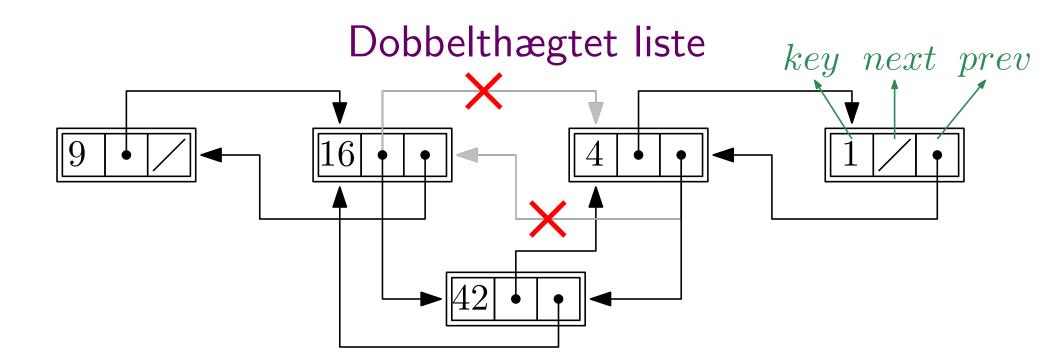


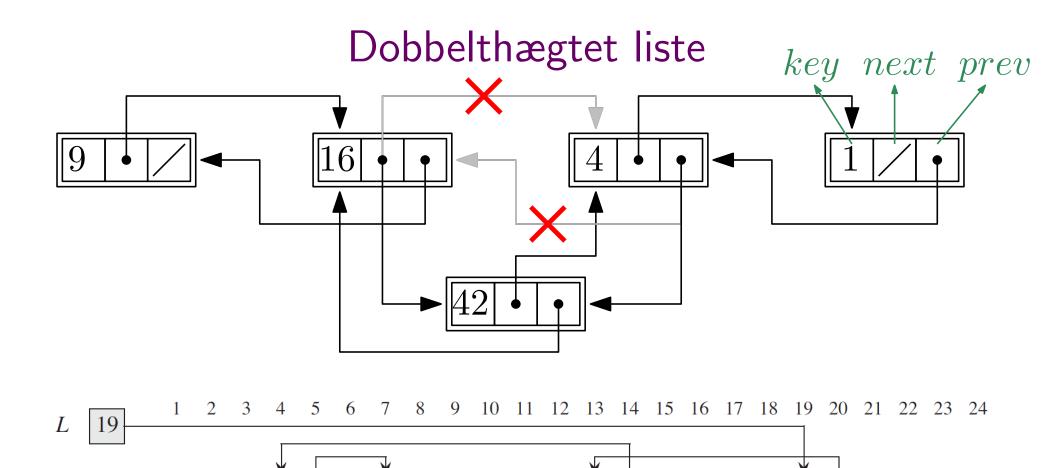












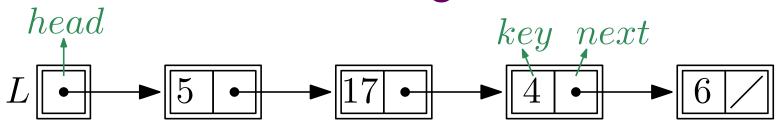
16

key | prev next

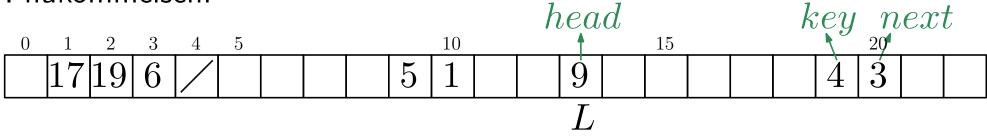
 \boldsymbol{A}

13

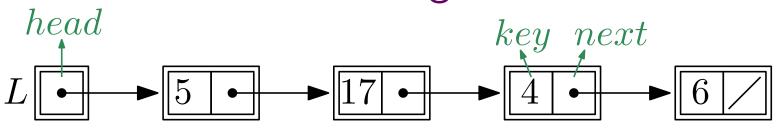
13



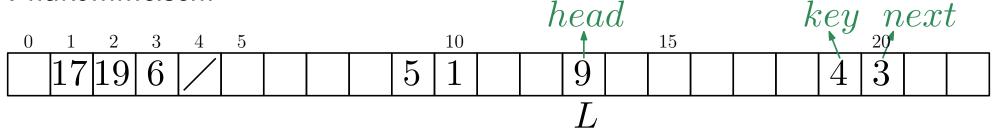
I hukommelsen:



List-Search(L, k): Søg fra begyndelsen efter nøgleværdi k.

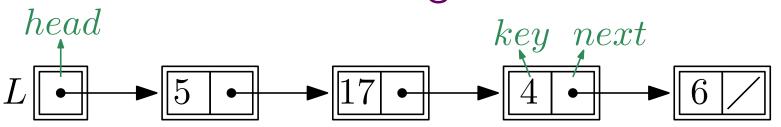


I hukommelsen:

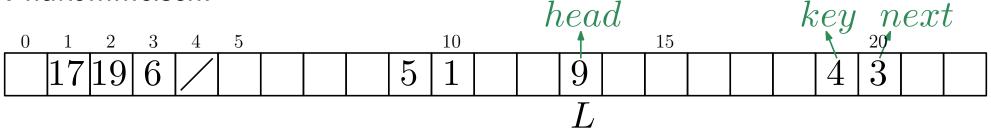


List-Search(L, k): Søg fra begyndelsen efter nøgleværdi k.

$$\begin{aligned} & \text{List-Search}(L,k) \\ & x = L.head \\ & \text{while } x \neq \text{NIL and } x.key \neq k \\ & x = x.next \\ & \text{return } x \end{aligned}$$



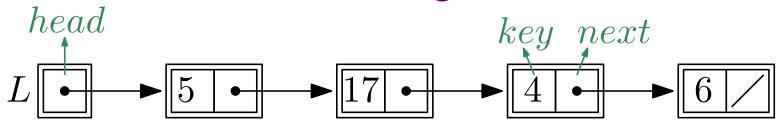
I hukommelsen:



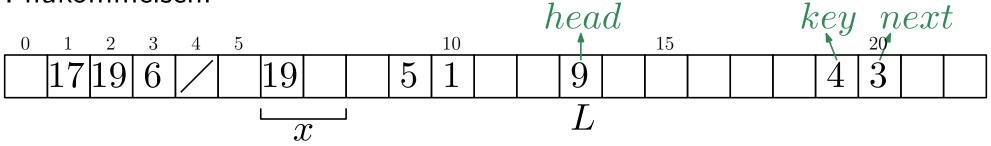
List-Search(L, k): Søg fra begyndelsen efter nøgleværdi k.

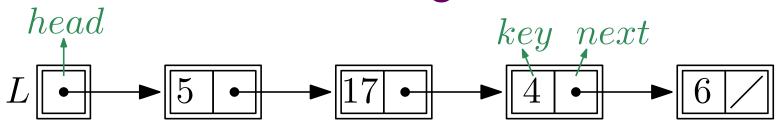
$$\begin{aligned} & \text{List-Search}(L,k) \\ & x = L.head \\ & \text{while } x \neq \text{NIL and } x.key \neq k \\ & x = x.next \\ & \text{return } x \end{aligned}$$

 $\Theta(n)$ tid.

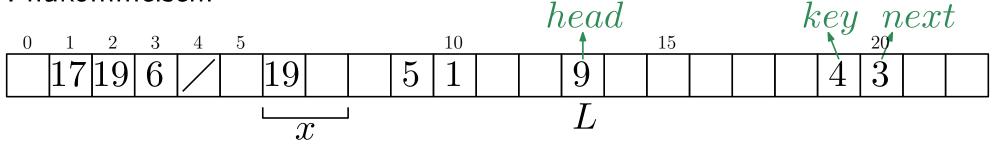


I hukommelsen:



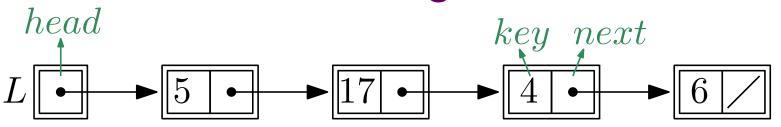


I hukommelsen:

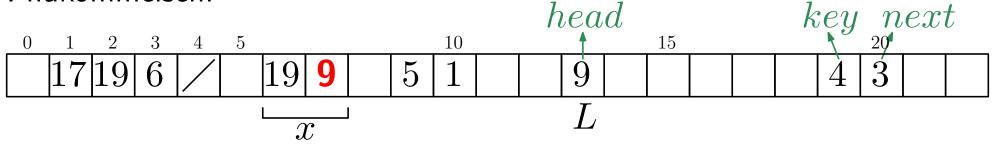


List-Insert
$$(L, x)$$

 $x.next = L.head$
 $L.head = x$

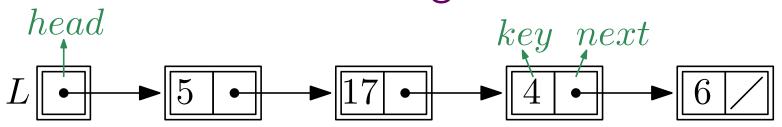


I hukommelsen:

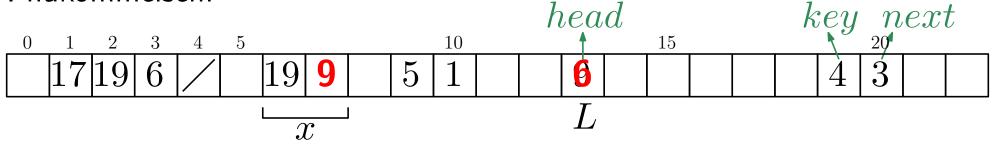


List-Insert
$$(L, x)$$

 $x.next = L.head$
 $L.head = x$



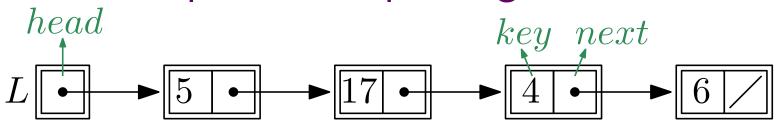
I hukommelsen:



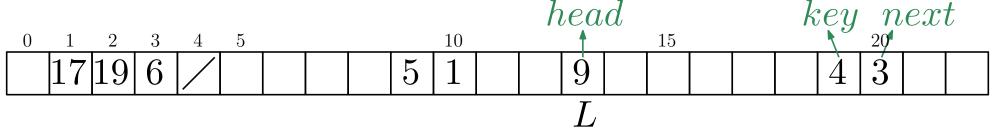
List-Insert
$$(L, x)$$

 $x.next = L.head$
 $L.head = x$

Operationer på hægtede lister



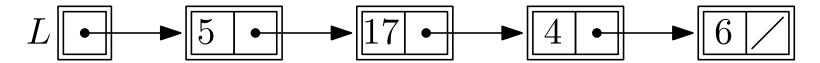
I hukommelsen:



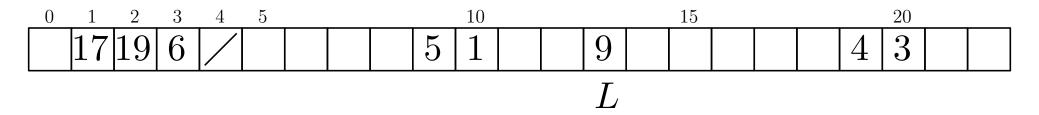
List-Search(L, k): Søg fra begyndelsen efter nøgleværdi k. $\Theta(n)$ tid.

List-Insert(L, x): Indsæt listeelementet x i begyndelsen af L. $\Theta(1)$ tid.

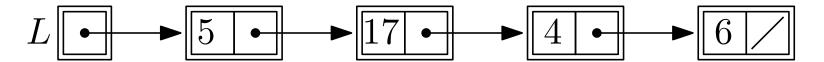
List-Delete(L,x): Slet listeelementet x fra L. For dobbelthægtede lister: $\Theta(1)$ tid.



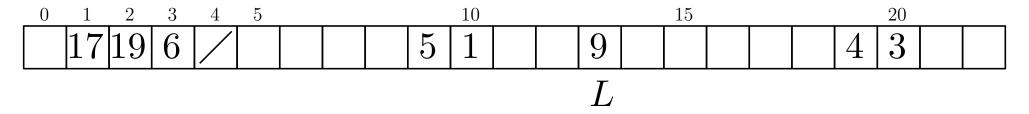
I hukommelsen:



Tilgå felt nummer k i hukommelsen eller i array A tager $\Theta(1)$ tid.

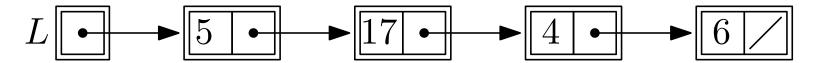


I hukommelsen:

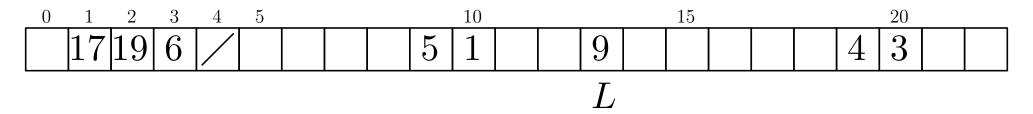


Tilgå felt nummer k i hukommelsen eller i array A tager $\Theta(1)$ tid.

Tilgå element nummer k i en liste tager $\Theta(k)$ tid.



I hukommelsen:



Tilgå felt nummer k i hukommelsen eller i array A tager $\Theta(1)$ tid.

Tilgå element nummer k i en liste tager $\Theta(k)$ tid.

I F#:

mylist.[k] tager $\Theta(k)$ tid!

"Lists in F# are implemented as singly linked lists, which means that operations that access only the head of the list are O(1), and element access is O(n)."