

Topological sorting, O, and Θ

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Outline

Topological sorting

O as a relation

 Θ as a relation



Recap and intro

A partial order on set A is a relation R which is reflexive, antisymmetric, and transitive.

- A poset (A, ≤)
- Posets can contain incomparable elements.
 Recall: Elements 5 and 3 are incomparable in poset (Z⁺, |).
- If any two elements in a poset (A, ≤) are comparable, then we say that (A, ≤) is totally ordered.
- If (A, ≤) is finite and totally ordered, then its elements can be sorted as

$$a_1, \ldots, a_n$$

where
$$|A| = n$$
 and $a_1 \le a_2 \le ... \le a_n$

Q: How can we sort the elements of a partial order?



Sorting elements of a poset (A, \leq)

Goal: Given an n-element poset (A, \leq) , produce an ordering

$$a_1, \ldots, a_n$$
 (1)

which is consistent with the "≤" partial order:

$$a_i \leqslant a_j \Rightarrow i \leqslant j$$
 (2)

Def. We refer to the ordering in (1) as topological sorting of poset (A, \leq) .

- Contrapositive of (2): $i > j \Rightarrow a_i \leqslant a_j$
- What does $a_i \leqslant a_i$ mean?



Topological sorting of a poset (A, \leq)

... is an ordering of its elements

$$a_1, a_2, \ldots, a_n$$

where $i > j \Rightarrow \alpha_i \leqslant \alpha_j$.

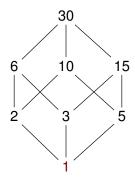
Equivalently:

$$\begin{array}{rcl} i{>}1 & \Rightarrow & \alpha_i{\leqslant}\alpha_1 \\ i{>}2 & \Rightarrow & \alpha_i{\leqslant}\alpha_2 \\ i{>}3 & \Rightarrow & \alpha_i{\leqslant}\alpha_3 \\ & \vdots \\ i{>}(n{-}1) & \Rightarrow & \alpha_i{\leqslant}\alpha_{n{-}1} \end{array}$$

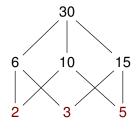


How to sort

Consider poset $(\{1,2,3,5,6,10,15,30\}, |)$

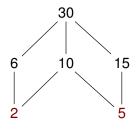






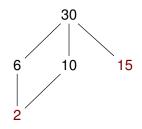
1,





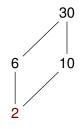
1,3,





1,3,5,





1,3,5,15,





1,3,5,15,2,





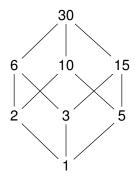
1,3,5,15,2,10



30

1,3,5,15,2,10,6





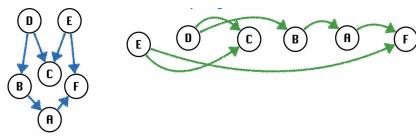
1,3,5,15,2,10,6,30

Note: The ordering produced by topological sorting is not unique.



Topological sorting and digraphs

Informally speaking, a topological sorting of an acyclic digraph is an ordering of the vertices, where all the arrows "point to the right".



$$A = \{A, B, C, D, E, F\}$$

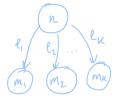
$$R = \{(A, F), (B, A), (D, B), (D, C), (E, C), (E, F)\}$$



Finding the next minimal element efficiently

Kahn's algorithm with O(|V| + |E|) runtime

```
L ← Empty list that will contain the sorted elements
S ← Set of all nodes with no incoming edge
while S is non-empty do
   remove a node n from S
   add n to tail of L
   for each node m with an edge e from n to m do
      remove edge e from the graph
   if m has no other incoming edges then
   insert m into S
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O as a relation



O as a relation

Recall: For asymptotically positive sequences (a_n) , (b_n) , we say that (a_n) is $O((b_n))$ if $\exists C>0$ and $\exists k\geqslant 1$ such that $\forall n\geqslant k$

$$a_n \leq Cb_n$$

How to think of O as a relation on some set A:

- $\bullet \ \ A = \{ \text{all asymptotically positive sequences } (c_{\mathfrak{n}}) \text{ of numbers} \}$
- $\bullet \ (a_n)O(b_n) \Leftrightarrow \exists C > 0 \exists k \geqslant \mathsf{1} \big(n \geqslant k \Rightarrow a_n \leqslant Cb_n \big)$

Question: Is O a partial order?



Is O a partial order? X

Reflexive,
$$(\alpha_n)O(\alpha_n)$$
? \checkmark For $C=1, k=n_0$ we have $(n\geqslant 1\Rightarrow \alpha_n\leqslant 1\cdot \alpha_n)$

Transitive,
$$(a_n)O(b_n) \wedge (b_n)O(c_n) \Rightarrow (a_n)O(c_n)$$
? \checkmark

- $\exists C > 0 \ \exists k \geqslant 1 \ (n \geqslant k \Rightarrow a_n \leqslant Cb_n)$
- $\exists D > 0 \ \exists \ell \geqslant 1 \ (n \geqslant \ell \Rightarrow b_n \leqslant Dc_n)$
- Let E = CD, m = max{k, ℓ }. Assume n \geqslant m. Then $a_n \leqslant Cb_n \leqslant CDc_n = Ec_n$

So
$$n \geqslant m \Rightarrow a_n \leqslant Ec_n$$

Antisymmetric,
$$(a_n)O(b_n) \wedge (b_n)O(a_n) \Rightarrow (a_n) = (b_n)$$
? $(n)O(2n)$ and $(2n)O(n)$, but $(n) \neq (2n)$



Θ as a relation



Θ as a relation

Recall: We say that (a_n) is $\Theta((b_n))$ if (a_n) is $O((b_n))$ and (b_n) is $O((a_n))$

- $\bullet \ \ A = \{ \text{all asymptotically positive sequences } (c_{\mathfrak{n}}) \text{ of numbers} \}$
- $(a_n)\Theta(b_n) \Leftrightarrow (a_n)O(b_n)$ and $(b_n)O(a_n)$
- So $\Theta = O \cap O^{-1}$

Question: Is Θ an equivalence relation¹?



¹reflexive, symmetric, and transitive

Is $\Theta = O \cap O^{-1}$ an equivalence relation?

- O is reflexive and transitive.
- **Thm.** If R is reflexive/transitive then R⁻¹ is also reflexive/transitive.
 - So O^{-1} is reflexive and transitive.
- Recall from last week: Thm. If R and S are reflexive/transitive then R ∩ S is reflexive/transitive.
 - So $\Theta = O \cap O^{-1}$ is reflexive and transitive.
- Final observation: $\Theta = O \cap O^{-1}$ is symmetric.
- Answer: Θ is an equivalence relation.

