

DMA: Logic

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Warm-up: Which of the following are true propositions?

1. Every $n \in \mathbb{Z}^+$ can be expressed as a product of primes and $n \log_2(n) \in O(n^2)$.
2. If $1 + 4 = 7$ then 4 is a prime.
3. Please, take a seat.
4. If someone in this class has a birthday today, then $\text{LCM}(12, 8) = 24$.

Plan for this class

- ▶ Properties of logical operations
- ▶ Predicates
- ▶ Logical statements with quantifiers
- ▶ Negation of statements with quantifiers

Properties of logical operations

Recall: $p \equiv q$ means that p , q always have the same truth value.

Example: $(r \Rightarrow s) \equiv ((\sim s) \Rightarrow (\sim r))$

Properties of AND, OR, NOT [Thm 2.2.1]

$$(1) \quad p \vee q \equiv q \vee p$$

$$(2) \quad p \wedge q \equiv q \wedge p$$

$$(3) \quad p \vee (q \vee r) \equiv (p \vee q) \vee r$$

$$(4) \quad p \wedge (q \wedge r) \equiv (p \wedge q) \wedge r$$

$$(7) \quad p \vee p \equiv p$$

$$(8) \quad p \wedge p \equiv p$$

$$(9) \quad \sim(\sim p) \equiv p$$

$$(10) \quad \sim(p \vee q) \equiv (\sim p) \wedge (\sim q)$$

$$(11) \quad \sim(p \wedge q) \equiv (\sim p) \vee (\sim q)$$

We can prove the above by comparing the truth tables of LHS and RHS.

Properties involving \Rightarrow and \Leftrightarrow [Thm. 2.2.2]

- (a) $(p \Rightarrow q) \equiv (\sim p \vee q)$
- (b) $(p \Rightarrow q) \equiv ((\sim q) \Rightarrow (\sim p))$
- (c) $(p \Leftrightarrow q) \equiv (p \Rightarrow q) \wedge (q \Rightarrow p)$
- (d) $\sim(p \Rightarrow q) \equiv (p \wedge (\sim q))$
- (e) $\sim(p \Leftrightarrow q) \equiv (p \wedge (\sim q)) \vee (q \wedge (\sim p))$

We proved (a) and (b) on Monday.

Items (d) and (e) can be proven the previous properties.

Thm 2.2.4 gives more properties...

Predicates

Predicates

Consider the sentence: " $x^2 + y^2 = 4$ "

It has a truth value only after we fix x and y .

Def. A **predicate** is a declarative sentence whose truth value depends on one or more (free) variables.

Notation: $P(x)$, $Q(x)$, $P(x, y)$ etc.

Examples of predicates

- ▶ $Q(x, y) : x^2 + y^2 = 4$
- ▶ $P(n) : 3^n + 6n - 1$ is divisible by 4

Defining sets using predicates

$S' = \{x \in S \mid P(x)\}$ denotes the collection of all $x \in S$ for which the statement $P(x)$ is true.

- ▶ Let $S = \mathbb{Z}$ and $P(x) : x > 0$. We have

$$\{x \in \mathbb{Z} \mid x > 0\} = \{1, 2, 3, 4, \dots\} = \mathbb{Z}^+$$

- ▶ We can define the positive reals, \mathbb{R}^+ , as $\{x \in \mathbb{R} \mid x > 0\}$
- ▶ Let $Q(x, y) : x^2 + y^2 = 4$. What is the set

$$\{(x, y) \in \mathbb{R}^2 \mid Q(x, y)\}?$$

Quantifiers

Quantifiers

Given a predicate $P(x)$ where $x \in S$, we can form two propositions:

1. Proposition

$$\forall x \in S \ P(x)$$

is true if $P(x)$ is true **for all** $x \in S$.

Pronounced: "for all elements x of S , $P(x)$ is true"

2. Proposition

$$\exists x \in S \ P(x)$$

is true if $P(x)$ is true for **some** $x \in S$.

Pronounced: "there exists $x \in S$ such that $P(x)$ is true"

Quantifiers: examples

1. Let $P(n) : n^2 + 1 \geq 2$. We can write the statement “*For all positive integers n , we have $n^2 + 1 \geq 2$* ” as

$$\forall n \in \mathbb{Z}^+ \quad P(n)$$

or as

$$\forall n \in \mathbb{Z}^+ \quad n^2 + 1 \geq 2$$

2. We can also specify sets using inequalities and write

$$\forall x \in \mathbb{R}^+ \quad x^2 + 1 \geq 2$$

as

$$\forall \text{real } x > 0 \quad x^2 + 1 \geq 2$$

Note: The statements in (1) are true, but the ones in (2) are false.

Order of quantifiers

Example: Let

$$S = \{\text{Living and deceased people}\}$$

Consider the following statements:

1. $\forall p \in S \exists m \in S$ m is the mother of p
2. $\exists m \in S \forall p \in S$ m is the mother of p

Are these different? Yes!

1. Every p has a mother
2. There is an m that is the mother for **all** p

The order matters! In general $\forall x \exists y P(x, y)$
is different from $\exists y \forall x P(x, y)$

Negation of statements with quantifiers

How can we rewrite $\sim(\forall x \in S \ P(x))$ mean?

Example.

q : **All** students in this class have read KBR Chapter 2.

What is $\sim q$?

$\sim q$: **There is** a student in this class
who has not read KBR Chapter 2.

In general, we have

$$\sim(\forall x \in S \ P(x)) \equiv (\exists x \in S \ \sim P(x)) \quad \text{Thm 2.2.3 (a)}$$

$$\sim(\exists x \in S \ P(x)) \equiv (\forall x \in S \ \sim P(x)) \quad \text{Thm 2.2.3 (b)}$$

Negation of Big- Θ using quantifiers

Big- Θ : Let $f, g: \mathbb{R}^+ \rightarrow \mathbb{R}$ be asymptotically positive. Then f is $\Theta(g)$ if there exists $c_1, c_2, x_0 > 0$ such that $c_1 g(x) \leq f(x) \leq c_2 g(x)$ for all $x \geq x_0$.

With quantifiers:

$$\exists c_1 > 0 \exists c_2 > 0 \exists x_0 > 0 \forall x \geq x_0 \ c_1 g(x) \leq f(x) \leq c_2 g(x).$$

The negation:

$$\begin{aligned} & \sim [\exists c_1 > 0 (\exists c_2 > 0 \exists x_0 > 0 \forall x \geq x_0 \ c_1 g(x) \leq f(x) \leq c_2 g(x))] \\ \equiv & \forall c_1 > 0 \sim [\exists c_2 > 0 (\exists x_0 > 0 \forall x \geq x_0 \ c_1 g(x) \leq f(x) \leq c_2 g(x))] \\ \equiv & \forall c_1 > 0 \forall c_2 > 0 \sim [\exists x_0 > 0 (\forall x \geq x_0 \ c_1 g(x) \leq f(x) \leq c_2 g(x))] \\ \equiv & \forall c_1 > 0 \forall c_2 > 0 \forall x_0 > 0 \sim [\forall x \geq x_0 (c_1 g(x) \leq f(x) \leq c_2 g(x))] \\ \equiv & \forall c_1 > 0 \forall c_2 > 0 \forall x_0 > 0 \exists x \geq x_0 \sim [c_1 g(x) \leq f(x) \leq c_2 g(x)] \\ \equiv & \forall c_1 > 0 \forall c_2 > 0 \forall x_0 > 0 \exists x \geq x_0 \\ & [c_1 g(x) > f(x) \vee f(x) > c_2 g(x)] \end{aligned}$$

What is going on here?

Three mathematicians walk into a bar.
The bartender asks: “Do all three of you want a beer?”
The first mathematician says: “I don’t know.”
The second one says: “I don’t know.”
The third one says: “Yes!”

Let S be the set of three mathematicians.

Consider the predicate $P(x) : x \text{ wants to have a beer.}$

The bartender asks if the following is true

$$\forall x \in S \ P(x) \tag{1}$$

Why don’t the first two mathematicians know the truth value of (1) but the last one does?

Implications involving predicates

Let $P(x)$ and $Q(x)$ be predicates concerning $x \in S$.

It is a **convention** that with $P(x) \Rightarrow Q(x)$ we mean the statement

“For all $x \in S$ we have that $P(x) \Rightarrow Q(x)$.”

Example. Statement

$$x > 1 \Rightarrow x^2 > 1$$

means

“For all x we have that if $x > 1$ then $x^2 > 1$.”

You should be able to:

- ▶ Rewrite sentences using logical connectives and quantifiers.
- ▶ Translate logical statements into everyday language.
- ▶ Evaluate whether a given proposition is true or false.
- ▶ Rewrite and simplify logical statements using properties of logical operations and quantifiers.

Puzzle: The blue-eyed islanders

The islanders: 60 with blue eyes and 70 with brown eyes.

They know: everyone on the island is an excellent logician.

They don't know: their own eye color, the total number of blue-eyed people. Also they cannot talk to each other.

The pirate captain: visits every night and will free any *blue-eyed* person who can tell him their own eye color. (He will “dispose of” anyone who guesses incorrectly.)

The guru: proclaims on Day 1 that someone on the island has blue eyes.

The question: Who gets off the island and when?