

# Iteration Functions

## 1 Taylor Polynomial

Given a polynomial  $p(z)$  of degree  $n$  and a complex number  $a$  we have

$$p(z) = p(a) + p'(a)(z - a) + \frac{p''(a)}{2!}(z - a)^2 + \cdots + \frac{p^{(n)}(a)}{n!}(z - a)^n. \quad (1)$$

Suppose  $\theta$  is a root of  $p(z)$ , i.e.  $p(\theta) = 0$ . Let  $z = \theta$ ,  $a = z$  in the above we get

$$0 = p(\theta) = p(z) + p'(z)(\theta - z) + \frac{p''(z)}{2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (2)$$

The above holds for any root of  $\theta$  and for any  $z$ . Adding  $z$  to both sides of the above, define

$$B_1(z) \equiv z - p(z) = z + p'(z)(\theta - z) + \frac{p''(z)}{2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (3)$$

Equivalently,

$$B_1(z) = z - p(z) = \theta + (p'(z) - 1)(\theta - z) + \frac{p''(z)}{2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (4)$$

Note that  $B_1(\theta) = \theta$ . So  $\theta$  is a fixed point of  $B_1(z)$ .

Given  $z_0$ , define the fixed point iteration

$$z_{k+1} = B_1(z_k), \quad k \geq 0. \quad (5)$$

What can we say about the fixed point iteration? Will it converge when  $z_0$  is close to  $\theta$ ?

From (4) we can write

$$B_1(z) - \theta = (p'(z) - 1)(\theta - z) + \frac{p''(z)}{2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (6)$$

Or,

$$B_1(z) - \theta = (\theta - z) \left( (p'(z) - 1) + \frac{p''(z)}{2!}(\theta - z) + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^{n-1} \right). \quad (7)$$

Notice that we can write the above after factoring  $(\theta - z)$  as

$$B_1(z) - \theta = (\theta - z) \left( (p'(z) - 1) + (\theta - z)G(z) \right), \quad (8)$$

where  $G(z)$  is a sum of terms. What is important is that when  $(\theta - z)$  is small,  $(\theta - z)G(z)$  is small. So if  $|p'(\theta) - 1| < 1$ . Then there is a neighborhood of the root  $\theta$  so that for any  $z_0$  in this neighborhood fixed point iteration converges to  $\theta$ . A neighborhood of  $\theta$  means the disc of some radius  $r > 0$  centered at  $\theta$ :

$$D_r(\theta) = \{z : |z - \theta| < r\}. \quad (9)$$

More formally, using the triangle inequality we can write,

$$|(p'(\theta) - 1) + (\theta - z)G(z)| \leq |(p'(\theta) - 1)| + |(\theta - z)G(z)|. \quad (10)$$

So for example if  $|(p'(\theta) - 1)| < .9$ , there will be a neighborhood where  $|(p'(z) - 1)| < .95$ , and  $|(\theta - z)G(z)| < .1$  for every  $z$  in this neighborhood.

## 2 Newton Method

From (2) we also get

$$zp'(z) - p(z) = \theta p'(z) + \frac{p''(z)}{2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (11)$$

Dividing by  $p'(z)$  we get Newton's iteration function

$$B_2(z) \equiv z - \frac{p(z)}{p'(z)} = \theta + \frac{p''(z)}{p'(z)2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{p'(z)n!}(\theta - z)^n. \quad (12)$$

This means

$$B_2(z) \equiv z - \frac{p(z)}{p'(z)} - \theta \approx \frac{p''(z)}{p'(z)2!}(\theta - z)^2. \quad (13)$$

If  $p'(\theta) \neq 0$ , i.e.  $\theta$  is a simple roots of  $p(z)$ , then  $B_2(\theta) = \theta$ , i.e.  $\theta$  is a fixed point of  $B_2(z)$ . In fact multiple roots are also fixed points. And any fixed point of  $B_2(z)$  is a root of  $p(z)$ .

Given  $z_0$ , define the fixed point iteration

$$z_{k+1} = B_2(z_k), \quad k \geq 0. \quad (14)$$

We have,

$$B_2(z) - \theta = \frac{p''(z)}{p'(z)2!}(\theta - z)^2 + \cdots + \frac{p^{(n)}(z)}{p'(z)n!}(\theta - z)^n. \quad (15)$$

This means if  $p'(\theta) \neq 0$ , there is a neighborhood around  $\theta$  so that starting with any  $z_0$  in this neighborhood the fixed point iteration converges. Furthermore,

$$\lim_{k \rightarrow \infty} \frac{z_{k+1} - \theta}{(\theta - z_k)^2} = \frac{p''(\theta)}{p'(\theta)2!}. \quad (16)$$

We say the rate of convergence is quadratic. This, roughly speaking, says when  $z_k$  is close enough to a root  $\theta$  the error in each iteration doubles. Thus if it is  $10^{-1}$ , you expect the next errors to be approximately  $10^{-2}$ ,  $10^{-4}$ ,  $10^{-8}$ , etc.

## 3 Newton Method for Multiple Roots

Given a polynomial  $p(z)$  and a root  $\theta$  we say it is a root of multiplicity  $m$  if

$$p(z) = (z - \theta)^m q(z),$$

where  $m \geq 1$  and  $q(\theta) \neq 0$ . When  $m = 1$  we say  $\theta$  is a simple root.

It is easy to show that in Newton's method,

$$B_2'(\theta) = 1 - \frac{1}{m}.$$

This implies

$$\lim_{k \rightarrow \infty} \frac{z_{k+1} - \theta}{(z_k - \theta)} = \frac{m - 1}{m}. \quad (17)$$

## 4 Some Facts on Dynamics of Newton's Method

The basin of attraction of a root  $\theta$  of a polynomial is the set of all input  $z_0$  so that orbit of  $z_0$ , denoted by  $O(z_0) = \{z_1, z_2, \dots\}$ , converges to  $\theta$ . The basin of attraction is denoted by  $A(\theta)$ .

What kind of a set it is?

A subset  $O$  of the Euclidean plane is said to be open if given any point  $z_0$  in  $O$ , there is an open disk  $D_r(z_0) = \{z : |z - z_0| < r\}$  that is contained in  $O$ .

We claim  $A(\theta)$  is an open set. First, there is a neighborhood, open disk at  $\theta$ , say  $D_r(\theta)$  for which any point in it converges to  $\theta$ .

Fact: Newton's iteration function is continuous at  $\theta$ .

Fact: Under continuity, inverse image of an open set is an open set.

$B_2^{-1}(D_r(\theta)) = \{z : B_2(z) \in D_r(\theta)\}$ .

Now if  $O(z_0)$  lies in  $D_r(\theta)$ , that is if the orbit at  $z_0$  converges to  $\theta$ , we claim that there is an open neighborhood of  $z_0$ , say  $D_t(z_0)$ , for some  $t > 0$ , such that any point  $z'$  in  $D_t(z_0)$ , the orbit of  $z'$  converges to  $\theta$ . To prove the existence of  $D_t(z_0)$ , we use the fact that since  $z_0$  converge to  $\theta$  this means after so many Newton iterations, say  $N$  iterations, all the subsequent iterates stay in  $D_r(\theta)$ . Now take the inverse of  $D_r(\theta)$ . This is an open set, say  $O_1$ . Take the inverse image of  $O_1$ , say  $O_2$ . This is an open set. So after  $N$  inverse images we get an open set  $O_N$  which contains  $z_0$  and every point in  $O_N$  converges to  $\theta$ .

Immediate basin of attraction: The largest connected component of  $A(\theta)$  that contains  $\theta$ .

## 5 Halley Method

We start again with the equation

$$0 = p(z) + p'(z)(\theta - z) + \frac{p''(z)}{2!}(\theta - z)^2 + \dots + \frac{p^{(n)}(z)}{n!}(\theta - z)^n. \quad (18)$$

Using the above and a mixture of  $B_1(z)$  and  $B_2(z)$  we would like to make a new iteration function that its order of convergence is cubic.

First note

$$B_1(z) - B_2(z) = -p(z) + \frac{p(z)}{p'(z)} = (p'(z) - 1)(\theta - z) + \sum_{i=2}^n \frac{(p'(z) - 1)p^{(i)}(z)}{i!p'(z)}(\theta - z)^i. \quad (19)$$

Multiply the above by  $p(z)$  and (18) by  $-(p'(z) - 1)(\theta - z)$  and adding, we get

$$p(z)(B_1(z) - B_2(z)) = p^2(z) \frac{(1 - p'(z))}{p'(z)} = \sum_{i=2}^n u_i(z)(\theta - z)^i, \quad (20)$$

where

$$u_i(z) = (p'(z) - 1) \left( \frac{p(z)p^{(i)}(z)}{i!p'(z)} - \frac{p^{(i-1)}(z)}{(i-1)!} \right). \quad (21)$$

Multiplying (20) by

$$\frac{-p''(z)}{2p'(z)u_2(z)} \quad (22)$$

and adding it to the expansion of  $B_2(z)$  we get

$$B_3(z) \equiv z - p(z) \frac{p'(z)}{p'(z)^2 - p(z)p''(z)/2} = \theta + \sum_{i=3}^n v_i(z)(\theta - z)^i, \quad (23)$$

where

$$v_i(z) = \left( \frac{p^{(i)}(z)}{i!p'(z)} - \frac{p''(z)}{2p'(z)} \frac{u_i(z)}{u_2(z)} \right). \quad (24)$$

The above iteration function is called Halley's method and has cubic order of convergence. This roughly means near a simple root if the current error is  $10^{-1}$ , the next one roughly  $10^{-3}$ ,  $10^{-6}$ , etc.

## 6 Horner's Method

Consider a polynomial

$$p_n(z) = a_n z^n + a_{n-1} z^{n-1} + \cdots + a_1 z + a_0.$$

To efficiently evaluate the polynomial at a point  $z_0$  we use nested multiplication:

$$p(z_0) = (\cdots ((a_n z_0 + a_{n-1}) z_0 + a_{n-2}) z_0 + \cdots + a_1) z_0 + a_0.$$

Let  $b_n = a_n$  and recursively define

$$b_{n-m} = b_{n-m+1} z_0 + a_{n-m}, \quad m = 1, \dots, n.$$

Performing  $n$  multiplications and  $n$  additions we evaluate  $p(z)$  at  $z_0$  to get

$$p(z_0) = b_0.$$

Let

$$p_1(z) = b_n z^{n-1} + b_{n-2} z^{n-2} + \cdots + b_2 z + b_1.$$

We have

$$p_1(z)(z - z_0) + b_0 = p(z).$$

From the above, differentiating we get

$$p'(z) = p_1'(z)(z - z_0) + p_1(z). \quad (25)$$

Thus

$$p'(z_0) = p_1(z_0).$$

So from Horner's recursion we get  $p(z_0)$  and  $p'(z_0)$ . By repeating this process we can compute all the normalized derivatives of  $p(z)$  at  $z_0$ . By induction from (25) we get

$$p^{(i)}(z) = p_1^{(i)}(z)(z - z_0) + i p_1^{(i-1)}(z). \quad (26)$$

Substituting  $z = z_0$  gives

$$p^{(i)}(z_0) = i p_1^{(i-1)}(z_0), \quad i = 1, \dots, n. \quad (27)$$

To summarize, given a polynomial

$$p_n(z) = a_n z^n + a_{n-1} z^{n-1} + \cdots + a_1 z + a_0,$$

and a particular  $z_0$ , we can compute all the normalized derivatives of  $p(z)$  at  $z_0$ :

$$b_i^{(i)} = \frac{p^i(z_0)}{i!}, \quad i = 0, \dots, n,$$

Set

$$b_m^{(-1)} = a_m, \quad m = 0, \dots, n.$$

For  $i = 0, \dots, n$ , do

$$b_n^{(i)} = b_n^{(i-1)}.$$

For  $m = 1, \dots, n - i$ , do

$$b_{n-m}^{(i)} = z_0 b_{n-m+1}^{(i)} + b_{n-m}^{(i-1)}.$$