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Extracting Success from IBM's 20-Qubit Machines Using Error-Aware Compilation

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Outline

- Problem to solve
- Background
- Experiment
- Results and Analysis

The problem

"NISQ (Noisy, Intermediate-Scale Quantum) computing requires error mitigation to achieve meaningful computation."

The authors needed to develop a compilation tool in order to maximize the success probability of real-world subroutines such as an adder circuit.

The problem

- Quantum computers exist, and can surpass classical computers on a range of important problems.
- Full realization of quantum error correction remains out of reach.
- Need error-aware compilers for NISQ computing.

"Not All Qubits Are Created Equal: A Case for Variability-Aware Policies for NISQ-Era Quantum Computers" by Tannu and Qureshi

"Full-Stack, Real-System Quantum Computer Studies: Architectural Comparisons and Design Insights" by Murali et al.

Background

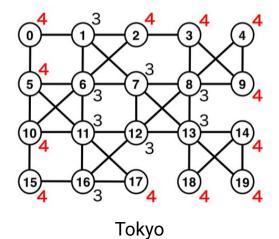
- 1. Machine-level compilation
- 2. IBM's 20-qubits quantum processor
- 3. Error model
- 4. Tomography
- 5. Randomized benchmarking
- 6. Architecture-aware compilation

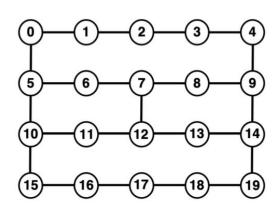
Machine-level compilation

 Decomposition of high-level language constructs into a series of one- and two-qubit operations that can be executed on the target system;

- Mapping of the variables defined by the programmer to locations in the system, in tandem with generation of appropriate execution of gates between qubits unfortunately placed far apart;
- 3. Generation of low-level control for the hardware itself.

IBM's 20-qubits quantum processor





Poughkeepsie

Each vertex represents a qubit, and each edge indicates that a CNOT gate can be executed between the two qubits.

The number outside the circles indicate the eccentricity, or maximum distance to another qubit.

Error model

- Single-Qubit Gate errors -> G
 - a. Unitary bit flips
 - b. Unitary phase flips
- 2. Bi-Qubit Gate errors(CNOT error) -> B
 - a. One or both qubits value flips
 - b. One or both qubits phase flips
 - c. IMPORTANT, can propagate errors from one qubit to another
- 3. State Preparation and measurement(SPAM) errors -> S
 - a. Can appear only once per qubit
 - b. Less cumulative impact

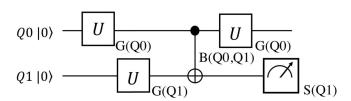


Figure 4: Errors and the gates with which they occur

Tomography

Two primary forms: state tomography & process tomography

Tells us how well we have done at creating our desired state, or how well a particular process (gate or set of gates) works, respectively.

The number of possible states naturally grows exponentially with the number of qubits. Need to test not just for bit flip errors, but also for phase flip errors.

Randomized benchmarking

- 1. Randomly select m gates from the *Clifford* group and arrange them in any order;
- Select the Clifford gate (or gates, if performing RB on more than one qubit) that will reverse the operation of the entire preceding sequence of m Clifford gates, C_{m+1} execute this as the m+1th gate;
- 3. Measure the qubit(s). If the output state is not equal to the input state, then an error has occurred somewhere in the whole circuit;
- 4. Change the number m and perform steps 1 to 3 again.

Clifford gates are the elements of the Clifford group, a set of mathematical transformations which effect permutations of the Pauli operators.

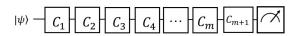


Figure 5: Randomized Benchmarking

$$C_{m+1} = \left(\prod_{i=1}^m C_i\right)^{\dagger}$$

In this way, it's possible to estimate the average value of errors per gate included in the gate set.

Architecture-aware compilation

Previously:

Focus on the impact on execution time.

Recently:

Pay attention to fidelity improvements.

Experiment

- 1. Long-distance CNOTs and making choices
 - a. Path selection for remote CNOT
 - b. Circuit selection for remote CNOT
- 2. Compiling complete programs
 - a. Search space
 - b. The optimization algorithm
 - c. Subroutines for compilation algorithm
 - d. Experimental evaluation of compilation

Long-distance CNOTs and making choices

WHY?

Each G, B, and S error above will have an error rate ε dependent on type and location.

$$ESP = \prod_{i} (1 - \varepsilon_i).$$

Path selection for remote CNOT

Problem 1: CNOT Path Selection

The programmer wishes to execute the circuit shown in Fig. 6. When the starting point (control qubit) and the end point (target qubit) for the Bi-Qubit gate (CNOT) are not neighbors, which path is the highest fidelity?

Result 1

ESP was able to select the better of two routes with accuracy of 70% at 2 hops, 66.6% at 3 hops and 62.5% at 4 hops.

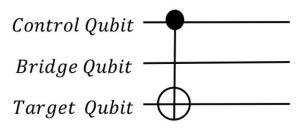
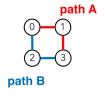
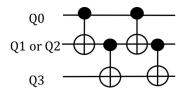


Figure 6: Circuit A, a CNOT gate skipping across another qubit

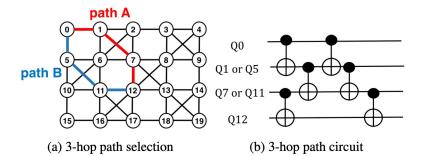
Path selection for remote CNOT

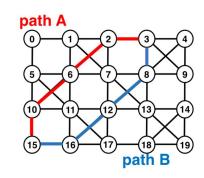


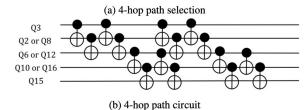
(a) 2-hop path selection



(b) 2-hop path circuit







Problem 2: Circuit Selection

The programmer wishes to execute the circuit shown in Fig. 6. When the CNOT gate is executed with remote qubits, there are multiple theoretically equivalent circuits as shown in Fig. 10.

However, considering errors, these are not necessarily equivalent. Which circuit has the highest fidelity?

Result 2

ESP was able to select the optimal circuit among three candidates with accuracy of 40%.

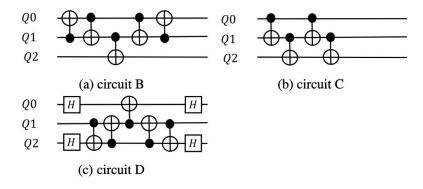


Figure 10: Circuits equivalent to the circuit shown in Fig. 6

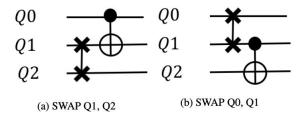


Figure 12: Circuits including SWAP gates

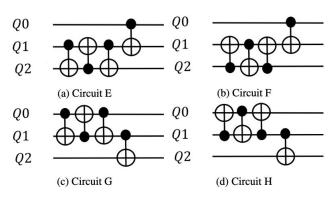


Figure 13: Circuits equivalent to Fig. 12. SWAP gate can be implemented with three CNOT gates.

Problem 2': Circuit Selection (including SWAP)

Perform the same circuit as Problem 2 with additional circuit options, including permitting the relocation of a qubit, using the SWAP gate shown in Fig. 12. Which circuit is the highest fidelity?

Result 2'

ESP was able to select the optimum circuit from among seven candidates 25% of the time. Fig.14 shows the fidelity of selected circuit relative to the other candidate circuits.

Problem 3: Circuit Selection (on Poughkeepsie) -

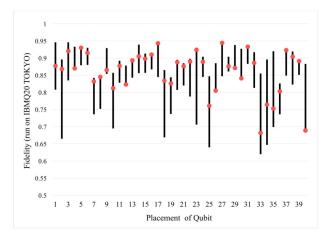
Perform the same circuit as Problem2âĂŹ on the other 20-qubit processor called Poughkeepsie. Is there any change in the reliability of circuit selection by ESP?

Result 3

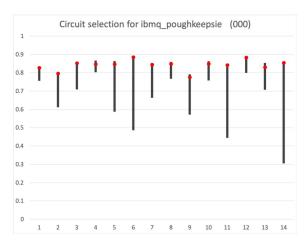
ESP was able to select the optimum circuit from among seven candidates 43% of the time. Fig.15 shows the fidelity of selected circuit relative to the other candidate circuits.

Table 1: the performance specifications of two IBMQ 20-qubit systems named Tokyo and Poughkeepsie

	tokyo	poughkeepsie
	(2nd gen 20-qubit system)	(3rd gen 20-qubit system)
Mean of Two-qubit (CNOT) error rates $\times 10^{-2}$	2.84	2.25
best	1.47	1.11
worst	7.12	6.11
Mean of Single-qubit error rates $\times 10^{-3}$	1.99	1.07
best	0.64	0.52
worst	6.09	2.77



Tokyo



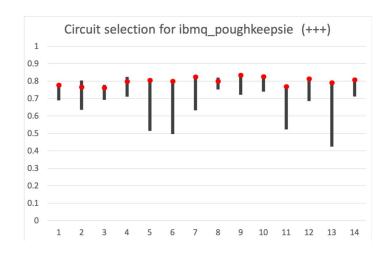
Poughkeepsie

Problem 3': Circuit Selection $(|+++\rangle)$

Perform the same circuit as Problem 3 but with different initial input states on Poughkeepsie. Is there any change in the reliability of circuit selection by ESP?

Result 3' -

ESP was able to select the optimum circuit from among seven candidates 36% of the time. Fig.16 shows the fidelity of selected circuit relative to the other candidate circuits.



Compiling complete programs

"A general purpose compiler, however, is required to combine such realizations in order to deal with much more complex quantum circuits with many CNOT gates."

"The compiler must be able to distinguish more reliable realizations from the others."

$$ESP(C) = \prod_{g \in C} ESP(g) = \prod_{g \in C} (1 - \varepsilon_g)$$

Search space

The search space of this problem can be visualized as a collection of hypercubes. Each hypercube consists of a complete set of execution states, and there is one corresponding hypercube for each qubit mapping. There are two types of transitions between the states: gate execution and swap insertion.

The optimization algorithm

```
Algorithm 1 Compilation Algorithm
 1: function COMPILE(gates, topology)
       N \leftarrow |gates|
                           initialize
       S_0
                                                               with
                                               states
    INITIAL_MAPPING(gates, topology)
       for i = 0 upto N (exclusive) do
          S_{i+1} \leftarrow \{\}
          for all state \in S_i do
             for all g \in gates do
                if dependency of g is satisfied then
 8:
                   qubits \leftarrow state.mapping[g.qubits]
10:
                   swap \leftarrow BEST\_SWAP(qubits, topology)
                   state' \leftarrow state with swap inserted and g ex-
11:
    ecuted
12:
                   state'.esp \leftarrow state.esp \times ESP(swap) \times
    ESP(g)
                   UPDATE_SCORE(state', gates, topology)
13:
                   S_{i+1} \leftarrow S_{i+1} \cup \{state'\}
14:
                end if
15:
             end for
16:
17:
          end for
18:
          S_{i+1} \leftarrow \text{top-}B \text{ states of } S_{i+1}
       end for
       return the best state \in S_N
21: end function
```

Beam-search based heuristic optimization algorithm.

Time Complexity:

$$O(N^2M + N^3B)$$

Subroutines for compilation algorithm

Algorithm 3 Greatest Connecting Edge Mapping

```
1: function INITIAL_MAPPING(gates,topology)
        extract guest graph G_V = (V_V, E_V) from gates
       Mapped = \{\}
        (v_1, v_2) \leftarrow \text{greatest edge} \in E_V
        (q_1,q_2) \leftarrow \text{lowest edge} \in topology
       M \leftarrow \{v_1 \mapsto q_1, v_2 \mapsto q_2\}
       Mapped \leftarrow Mapped \cup \{v_1, v_2\}
       E_V \leftarrow E_V \setminus \{(v_1, v_2)\}
        while not all endpoints in E_V are mapped do
           pick the greatest edge (v, v') \in E_V where v \in S and
    v' \notin S
           if there is a free physical qubit adjacent to M[v] then
11:
12:
              pick the lowest edge (q, q') \in topology where
    q = M[v] and q' is a free qubit
              M \leftarrow M \cup \{v' \mapsto q'\}
13:
              S \leftarrow S \cup \{v'\}
14:
           end if
15:
           E_V \leftarrow E_V \setminus \{(v, v')\}
        end while
        for all v \in V_V \setminus Mapped do
           pick a random free physical qubit q \in topology
           M \leftarrow M \cup \{v \mapsto q\}
        end for
        return M
23: end function
```

Algorithm 4 Find The Best Swap Sequence

```
1: function BEST_SWAP(qubits, topology)
       if qubits.len() == 1 then
          return []
       end if
       ESP_{max} \leftarrow 0
       swap_{max} \leftarrow []
       for all (q_0, q_1) \in topology do
          swap_0 \leftarrow \text{find the swap sequence between } qubits[0]
    and q_0 with greatest ESP(swap_0)
          swap_1 \leftarrow \text{ find the swap sequence between } qubits[1]
    and q_1 with greatest ESP(swap_1)
          ESP \leftarrow ESP(swap_0) \times ESP(swap_1) \times
    ESP(CNOT \text{ over } q_0 \text{ and } q_1)
          if ESP > ESP_{max} then
             ESP_{max} \leftarrow ESP
13:
             swap_{max} \leftarrow swap_0 + swap_1
          end if
14:
       end for
       return swap<sub>max</sub>
17: end function
```

Algorithm 2 Update Score of a State

```
1: function UPDATE_SCORE(state, gates, topology)
2: score \leftarrow state.esp
3: for all g \in gates do
4: if g has not been executed then
5: swap \leftarrow BEST\_SWAP(state.mapping, g, topology)
6: score \leftarrow score \times ESP(swap) \times ESP(g)
7: end if
8: end for
9: state'.score \leftarrow score
10: end function
```

Experimental evaluation of compilation

"We cannot decide whether the circuit succeeded or failed from one shot of measurement. Instead, we have to focus on how close the empirical probability distribution sampled by running the compiled circuit on a NISQ machine multiple times and the ideal distribution an imaginary noiseless quantum computer will produce are."

Define KL divergence:

$$D_{KL}(P_{ideal}||P_{empirical}) = \sum_{x} P_{ideal}(x) \log \frac{P_{ideal}(x)}{P_{empirical}(x)}$$

The lower KL divergence is, the better the quality of the compiled circuit is.

Experimental evaluation of compilation

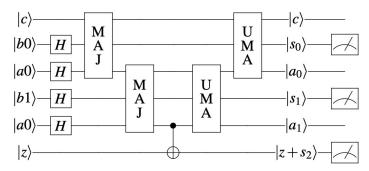
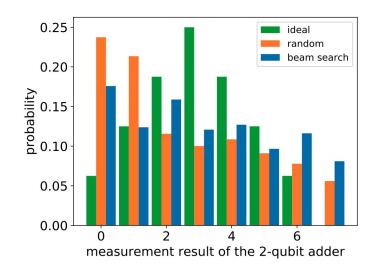
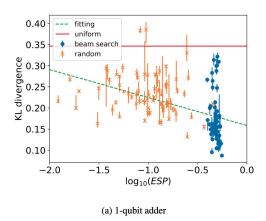
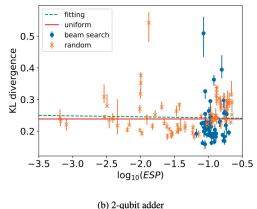


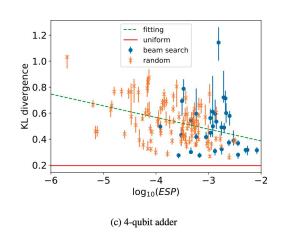
Figure 17: Cucarro's adder circuit of 2-qubit inputs and carryin $(|c\rangle)$ and carry-out $(|z\rangle)$. MAJ (MAJority) gate consists of one Toffoli gate and two CNOT gates and computes carry bit by determining the majority of inputs are $|1\rangle$ state. UMA (UnMajority and Add) is also made of one Toffoli and two CNOTs and computes the addition of this digit on $|b\rangle$, uncomputing the other qubits. We denote the i'th bit of the addition as $|s_i\rangle$ in the figure. Toffoli gates are made of 6 CNOT gates and 9 single qubit gates. Since we give $|0\rangle$ for $|c\rangle$ and $|z\rangle$, the measurement of this circuit gives you the distribution of the added value.



Experimental evaluation of compilation







Linear regression:

1-qubit adder = -0.475

2-qubit adder = -0.0279

4-qubit adder, negative correlation, but absolute value of KL divergence is much worse.

Impacts

Proposed two reliability metrics for quantum gates and circuits: Estimated Success Probability for use during compilation and KL divergence for assessing results.

Designed and implemented an error-aware quantum compiling system.

Could contribute to future quantum architectural design.

WHO should care & WHY

Quantum computer designers Quantum algorithm developers Quantum computing scientists

This paper exhibited a working error-aware compilation, which is crucial for future quantum computer design. The paradigm described in the paper can be incorporated directly into standard release of current quantum systems.

Is it believable?

Yes.

ACM Journal

24 citations for pre-print version(since May.2019)

4 citations for published version(since May.2020)

Limitations

- 1. Even in relatively simple case of selecting a two-hop path, the best success rate is only 70%.
- 2. Relation between ESP and KL divergence vanished on complex circuit.
- 3. Not considered memory errors.

Questions?

Thank you