

Lecture 5

CS 510

Linear systems

Want to solve

$$AX = b$$

A  $n \times n$

invertible

b  $n \times 1$

x  $n \times 1$  unknown

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \ddots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{pmatrix}$$

$$A = \begin{pmatrix} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} \\ a_{21}^{(1)} & a_{22}^{(1)} & \ddots & a_{2n}^{(1)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1}^{(1)} & a_{n2}^{(1)} & \cdots & a_{nn}^{(1)} \end{pmatrix}$$

assume  $a_{11}^{(1)} \neq 0$

$$A = \begin{pmatrix} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} \\ a_{21}^{(1)} & a_{22}^{(1)} & \ddots & a_{2n}^{(1)} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1}^{(1)} & a_{n2}^{(1)} & \cdots & a_{nn}^{(1)} \end{pmatrix}$$

Set

$$m_{ii} = \frac{a_{ii}^{(1)}}{a_{ii}^{(1)}}, i=1, \dots, n$$

Multiply row  $i$  by  
-  $m_{ii}$  and add to row  
 $i$ . This gives:

$$\left( \begin{array}{cccc} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & \cdots & a_{2n}^{(2)} \\ \vdots & \ddots & & \\ 0 & a_{nn}^{(2)} & \cdots & a_{nn}^{(2)} \end{array} \right)$$

$$a_{ij}^{(2)} = a_{ij}^{(1)} - m_{ij} a_{1j}^{(1)}$$

Set

$$m_{i2} = \frac{a_{i2}^{(2)}}{a_{22}^{(2)}}, i=3, \dots, n$$

Multiply row 2 by

- $m_{i2}$  and add to row

$i$ . This gives :

$$\left( \begin{array}{cccc} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & \cdots & a_{2n}^{(2)} \\ \vdots & 0 & \ddots & a_{3n}^{(3)} \\ 0 & 0 & \cdots & a_{nn}^{(2)} \end{array} \right)$$

$$a_{ij}^{(3)} = a_{ij}^{(2)} - m_{ij} a_{2j}^{(2)}$$

Eventually we get

$$\begin{pmatrix} a_{11}^{(1)} & \cdots & a_{1n}^{(1)} \\ 0 & a_{22}^{(2)} & \cdots & a_{2n}^{(2)} \\ \vdots & & \ddots & \\ 0 & & & a_{nn}^{(n)} \end{pmatrix} = U$$

Let

$$L = \begin{pmatrix} 1 & & & & \\ m_{21} & 1 & & & \\ m_{31} & m_{32} & 1 & & \\ \vdots & \vdots & & \ddots & \\ m_{n1} & m_{n2} & \dots & -m_{n,n-1} & 1 \end{pmatrix}$$

Then. If  $a_{ii}^{(i)} \neq 0$ ,  $i = 1, \dots, n-1$

This produces  $L, U$  &

$$A = L U$$

Operation Count  
Will Count multiplication/div

First step:

$$(n-1)^2$$

Second step

$$(n-2)^2 \text{ etc}$$

over n!!

i | v

$$1^2 + 2^2 + \dots + (n-1)^2 =$$

$$\frac{(n-1)n(2(n+1)+3)}{6}$$

$$\approx \frac{n^3}{3}$$

Solving  $Ax = b$ :

$$\underline{L}Ux = b$$

Let  $y = Ux$ ,  $Ly = b$   
Solve for  $y$ , then solve for  
 $x$ .

These are done by  
forward + backward  
substitution respectively.

Each costs

$$1+2+\dots+(n-1) = \frac{n(n-1)}{2} \approx \frac{n}{2}$$

∴ +<sub>0</sub> solve

$$Ax = b$$

(or + s)

$$\frac{n^3}{3} + \frac{n^2}{2} + \frac{n^2}{2}$$

operations

LU factorization with  
partial pivoting

Pick pivot row  $i_0$  to be row  
with largest  $a_{i_0 i}$  entry,  $i=1, \dots, n$   
This makes absolute value of  
multipliers to be  $\leq 1$

Do this for each  
column.

,

This results in

$$\underline{L} U = P A$$

where  $P$  is a permutation matrix

Ex.

$$x_1 - x_2 + 2x_3 = 2$$

$$-x_1 + 2x_2 + x_3 = 2$$

$$2x_1 - 4x_2 + x_3 = -1$$

$$A = \begin{bmatrix} 1 & -1 & 2 \\ -1 & 2 & 1 \\ 2 & -4 & 1 \end{bmatrix}, P = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$$

$$P = \begin{bmatrix} 3 \\ 2 \\ 1 \end{bmatrix}, \quad \begin{bmatrix} 2 & -4 & 1 \\ -1 & 2 & 1 \\ 1 & -1 & 2 \end{bmatrix}$$

$$m_{21} = -\frac{1}{2}$$

$$m_{31} = \frac{1}{2}$$

$$\begin{bmatrix} 2 & -4 & 1 \\ 0 & 0 & \frac{3}{2} \\ 0 & 1 & \frac{3}{2} \end{bmatrix}$$

$$, P = \begin{bmatrix} 3 \\ 1 \\ 2 \end{bmatrix}$$

$$\begin{bmatrix} 2 & -4 & 1 \\ 0 & 0 & \frac{3}{2} \\ 0 & 1 & \frac{3}{2} \end{bmatrix} = U$$

$$, P = \begin{bmatrix} 3 \\ 2 \\ 1 \end{bmatrix}$$

$$\begin{bmatrix} 2 & -4 & 1 \\ m_{21} & 0 & \frac{3}{2} \\ m_{31} & 1 & \frac{3}{2} \end{bmatrix}$$

$$\begin{bmatrix} 2 & -4 & 1 \\ m_{31} & 1 & \frac{3}{2} \\ m_{21} & 0 & \frac{3}{2} \end{bmatrix} = U$$

$$L = \begin{bmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ -\frac{1}{2} & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ -\frac{1}{2} & 1 & 0 \\ -\frac{1}{2} & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & -1 & 1 \\ 0 & 1 & 3 \\ 0 & 0 & t \end{bmatrix} = \begin{bmatrix} 2 & -4 & 1 \\ 1 & -1 & 2 \\ -1 & 2 & 1 \end{bmatrix}$$

To solve  $AX = b$ , we solve

$$Ly = Pb = \begin{bmatrix} -1 \\ 2 \\ 2 \end{bmatrix}$$

This gives  $y_1 = -1$

$$\frac{1}{2}y_1 + y_2 = 2$$

$$y_2 = 2 + \frac{1}{2} = \frac{5}{2}$$

$$-\frac{1}{2}y_1 + y_3 = 2, \quad y_3 = \frac{3}{2}$$

Next we solve

$$UX = Y$$

$$\begin{bmatrix} 2 & -4 & 1 \\ 0 & 1 & 0 \\ 0 & 0 & 3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} -1 \\ \frac{5}{2} \\ \frac{3}{2} \end{bmatrix}$$

$$x_3 = 1, \quad x_2 + \frac{3}{2}x_3 = \frac{5}{2}, \quad x_2 = -1, \quad x_1 = 1$$
$$2x_1 - 4x_2 + x_3 = -1$$

In summary :  
in LU factorization with partial  
pivoting we use the matrix  
A to store  $L + U$  and when  
necessarily we permute rows  
but keep track using an array  
 $P$ .

Computing  $A^{-1}$

First Compute LU

Then Compute with  
column of  $A^{-1}$ , say  $x^{(i)}$

by solving  $\mathcal{G}$

$$L U X^{(i)} = e^{(c)}$$

$$e^{(i)} = \begin{bmatrix} 0 \\ \vdots \\ i \\ 0 \end{bmatrix} \leftarrow | \quad \text{in } \underset{\cancel{i}}{i-th} \text{ location}$$

Complexity:

$$\frac{n^3}{3} + n \cdot n^2 = \frac{4n^3}{3}$$

Iterative Method to solve

$$A X = b$$

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & & & \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{bmatrix}, \quad b = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

$$A = D - L - U$$

$$D = \begin{bmatrix} a_{11} & & & \\ & \ddots & & \\ & & a_{ii} & \\ & 0 & \cdots & a_{nn} \end{bmatrix}, \quad L = \begin{bmatrix} 0 & & & \\ -a_{21} & \ddots & & \\ \vdots & & \ddots & \\ -a_{n1} & \cdots & 0 & \end{bmatrix}, \quad U = \begin{bmatrix} 0 & -a_{12} & \cdots & \\ \ddots & \ddots & -a_{ij} & \\ & & \ddots & \\ & & & 0 \end{bmatrix}$$

Jacobi Method

$$AX = DX - (L+U)X = b$$

$$DX = (L+U)X + b$$

Now given  $X^{(k)}$  as an  
approximation to solution of  
 $AX = b$ , we compute  $X^{(k+1)}$   
from

$$D X^{(k+1)} = (L + U) X^{(k)} + b$$

or

$$X^{(k+1)} = D^{-1}(L + U) X^{(k)} + D^{-1}b$$

Let  $T = D^{-1}(L + U)$ ,  $C = D^{-1}b$

so  $X^{(k+1)} = T X^{(k)} + C$

This is Jacobi method

Clearly  $a_{ii}$  must be non zero.

$$D^{-1} = \begin{bmatrix} \frac{1}{a_{11}} & & & \\ & \ddots & & 0 \\ & & \ddots & \frac{1}{a_{nn}} \\ 0 & \ddots & 0 & \end{bmatrix}$$

Let's write

$$T_j = D^{-1}(L + U), \quad c_j = D^{-1}b$$

Each iteration takes  $O(n^2)$  operations so if we don't need to iterate many times we gain over Gaussian Method

Will it converge?

Suppose  $X_x$  is solution to

$$AX = b$$

Then

$$X_x = T_j X_x + C \quad (1)$$

Also

$$X^{(k)} = T_j X^{(k-1)} + C \quad (2)$$

Subtracting (1) from (2)

$$X^{(k)} - X_x = T_j(X^{(k-1)} - X_x)$$

If we let

$$\ell^{(k)} = x^{(k)} - x^* \text{ (error)}$$

then

$$\ell^{(k)} = T_j \ell^{(k-1)}$$

$$so \quad \|\ell^{(k)}\| \leq \|T_j\| \cdot \|\ell^{(k-1)}\|$$

or

$$\|e^{(k)}\| \leq \|T_j\|^2 \cdot \|e^{(k-2)}\|$$

$$\|e^{(k)}\| \leq \|T_j\|^k \|e^{(0)}\|$$

If  $\rho(T_j) < 1$ , then

We can choose a matrix norm

so that  $\|T_j\| < 1 \Rightarrow$  convergence

We say A is diagonally dominant if

$$|a_{ii}| > \sum_{\substack{j=1 \\ j \neq i}}^n |a_{ij}|.$$

Thm. If A is diagonally dominant then Jacobi method converges.

Pf.

$$T_j = D^{-1}(L + U)$$

$$\|T_j\|_\infty < 1$$

Recall  $\|\cdot\|_\infty$  norm of a matrix  $X$  is  $\max_{\text{its rows}} \|\cdot\|_1$ , norm of

Gauss-Seidel

We start with

$$A x = b$$

$$A = D - L - U$$

$$(D - L - U) x = b$$

We write

$$(D - L)X = UX + b$$

-iteratively:

$$(D - L)X^{(k)} = UX^{(k-1)} + b$$

$$X^{(k)} = (D - L)^{-1}UX^{(k-1)} + (D - L)^{-1}b$$

$$\cancel{x}^{(k)} = T_g x^{(k-1)} + c_g$$

$$T_g = (D - L)^{-1} U$$

$$c_g = (D - L)^{-1} b$$

AS in Jacobi, if  
 $x_*$  is solution to  $Ax = b$

$$x^{(k)} - x_* = T_g(x^{(k-1)} - x_*)$$
$$e^{(k)} = T_g e^{(k-1)} = T_g e^{(0)}$$

Successive Over-relaxation method

SOR

$$A x = b$$

$$A = D - L - U$$

We have

$$Lx = (-D + U)x + b$$

let  $w$  be any scalar

$$wLx = w(-D + U)x + wb$$

add

$$+ Dx = Dx$$

to both sides

we set

$$(D - wL)X = [(1-w)D + wU]X_{wb}^+$$

write  $X^{(k)}$

$$(D - wL)X^{(k)} = [(1-w)D + wU]X^{(k-1)} + wb$$

Then

$$x^{(k)} = T_w x^{(k-1)} + c_w$$

where

$$T_w = (D - wL)^{-1} [(1-w)D + wU]$$

$$c_w = (D - wL)^{-1} wb$$

SOR  $\therefore$  is

$$x^{(k)} = \bar{T}_w x^{(k-1)} + c_w$$

We remark that to solve  
 $(D - wL)^{-1} u = b'$  time.  
can be done in  $O(n^2)$

In general Convergence of  
SOR & the selection of  
 $w$  is not so easy and  
may not converse.

However, some results are  
known.

Theorem (Kahan) If  $a_{ii} \neq 0$

$\forall i$ , then  $P(T_w) \geq$

$|w_i - 1| \cdot S_\delta$  SOR

Cannot converge if

$w \notin (0, 2)$ .

Thm ( Ostrowski-Richter )

If  $A$  is PD (Positive definite)

$w \in (0, 2)$ , then

SOR converges for

any  $x^{(0)}$ .

Thm. If  $A$  is PD & tridiagonal,  
then  $\rho(T_g) = \rho(T_j)^2 < 1$

and optimal  $w \approx$

$$w = \frac{2}{1 + \sqrt{1 - \rho(T_j)^2}}$$

and with this choice  
of  $w$

$$P(T_w) = w - 1.$$

Example,

$$A = \begin{bmatrix} 4 & 3 & 0 \\ 3 & 4 & -1 \\ 0 & -1 & 4 \end{bmatrix}$$

Claim:

A

is Positive  
definite,

i.e.  
 $\forall x \quad x^T A x > 0$

$$\begin{aligned}x^T A x &= \\4x_1^2 + 4x_2^2 + 4x_3^2 \\+ 3x_1x_2 + 3x_2x_1 \\- x_2x_3 - x_2x_3 \\&= 3(x_1 + x_2)^2 + (x_2 - x_3)^2 \\+ x_1^2 &\quad + 3x_3^2\end{aligned}$$

If  $x \neq 0$

$$x^T A x > 0$$

. . .  
Positive definite

This is  $\lambda > 0 + \lambda x$   
So A is Positive  
definite.

So SOR Convers  
for any  $w \in \mathbb{C}^n$

and an  $y$   
initial  $x^{(0)}$ .

What is optimal  
w?

$$T_j = D^{-1}(-L - U)$$

$$= \begin{bmatrix} \frac{1}{4} & 0 & 0 \\ 0 & \frac{1}{4} & 0 \\ 0 & 0 & \frac{1}{4} \end{bmatrix} \begin{bmatrix} 0 & -3 & 0 \\ -3 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$T = \begin{bmatrix} 0 & -\frac{3}{4} & 0 \\ -\frac{3}{4} & 0 & \frac{1}{4} \\ 0 & \frac{1}{4} & 0 \end{bmatrix}$$

What is  $P(\bar{T}_j) = ?$

We need to find  
eigenvalues of  $\bar{T}$

$$\det(\bar{T}_j - \lambda I)$$

$$\begin{array}{cccc} -\lambda & -\frac{3}{4} & 0 & | \\ -\frac{3}{4} & -\lambda & \frac{1}{4} & \\ \hline 0 & \frac{1}{2} & -\lambda & \\ \end{array}$$
$$= -\lambda \left( \lambda^2 - \frac{1}{16} \right)$$

so  $\lambda = \frac{1}{4}$  is target

$P(T_j) = \frac{1}{4}$

opt  $w$  is 2

$$w = \frac{2}{1 + \sqrt{1 - \frac{1}{2}}} = \frac{2}{1 + \frac{\sqrt{3}}{2}} \approx 1.25$$

Accelerated SOR

AOR

$$x^{(k)} = \bar{T}_{\sigma, \omega} x^{(k-1)} + c_{\sigma, \omega}$$

where

$$T_{\sigma, w} =$$

$$(D - \sigma L)^{-1} \left[ [(1-w)D + (w-\sigma)L] + wU \right]$$

$$c_{\sigma, w} = w(D - \sigma L)^{-1} b$$

If  $\sigma = 0$ ,  $w = 1$ , Jacobi

If  $\sigma = w = 1$ , Gauss-Siedel

If  $\sigma = w$ , SOR