


Correct

$T(n)$
★
 $O(n)$
 $O(\log n)$
 $O(n \log n)$
 $a * m$ $O(n^2)$

$$T(m, n) = O(m \log n)$$

$$\Rightarrow O(m \log n) \leq C * m \log n$$

$$\Rightarrow C * k \log \frac{n}{2} + C * (m - k) \log \frac{n}{2} \leq C * m \log n$$

$$\Rightarrow C * m (\log n - \log 2) + a * m \leq C * m \log n$$

$\hookrightarrow 1$

$$\Rightarrow C * m \log n - C * m + a * m \leq C * m \log n$$

$$\Rightarrow -C * m + a * m \leq 0$$

$$\Rightarrow \boxed{a \leq C}$$

$O(\sqrt{n})$

C exists ?

$$a_n = 6a_{n-2} - a_{n-1}$$

$$\Rightarrow a_n x^n = 6a_{n-2} x^n - a_{n-1} x^n$$

$$\sum_{n=2}^{\infty} a_n x^n = \sum_{n=2}^{\infty} 6a_{n-2} x^n - \sum_{n=2}^{\infty} a_{n-1} x^n \quad (1)$$

$$G(x) = \sum_{n=0}^{\infty} a_n x^n$$

$$\sum_{n=2}^{\infty} \quad \sum_{n=2}^{\infty}$$

$$G(x) = \sum_{n=2}^{\infty} a_n x^n + a_1 x + a_0 \quad (2)$$

$$\begin{cases} a_0 = 1 \\ a_1 = 2 \end{cases}$$

$$\text{CO2: } G(x) - 2x - a_0 = 6x^2 \cdot G(x) - x \cdot G(x) + a_0 x$$

$$G(x) = \frac{3x+1}{-6x^2+x+1}$$

$$G(x) = \frac{1}{-2x+1} = \sum_{n=0}^{\infty} a_n x^n$$

$$= \frac{1}{1-2x} = \sum_{n=0}^{\infty} a_n x^n$$

$$a_n = 2^n$$

$$\begin{cases} T(n) = 3T\left(\frac{n}{3}\right) + \underbrace{a \cdot (\log n)^3}_{\downarrow} & n > 2 \\ T(n) = b. & n \leq 2 \end{cases}$$

Master theorem $a \neq n^{\frac{1}{2}}$

first rule of Master theorem.

Divide & Conquer

First, find the middle row j of matrix $M(n, m)$

Linearly scan row j to find leftmost minimum element, x , and corresponding index, k , the leftmost min element for all rows having index less than j can only be exist at position $< k$.

Hence, we can recursively call the top-left part of $M[1 \dots (j-1), 1 \dots k]$

$M[(j+1) \dots n, k \dots m]$

Pseudo code :

find Leftmost Min val (M) :

if $n == 1$,

linear search for min value x
return x

else :

$$j = \frac{n}{2}$$

linear search for row j of M ,
which is $x(j, k)$

$$M_1 = M[\underbrace{1 \dots (j-1)}, \underbrace{1 \dots k}]$$

$$M_2 = M[\underbrace{(j+1) \dots n}, \underbrace{k \dots m}]$$

$$I_1 = \text{findLeftmost Min val} (M_1)$$

$$I_2 = \text{findLeftmost Min val} (M_2)$$

return x, I_1, I_2

$$\underline{O(m \log n)}$$