

V8中的指针压缩及其源码分析

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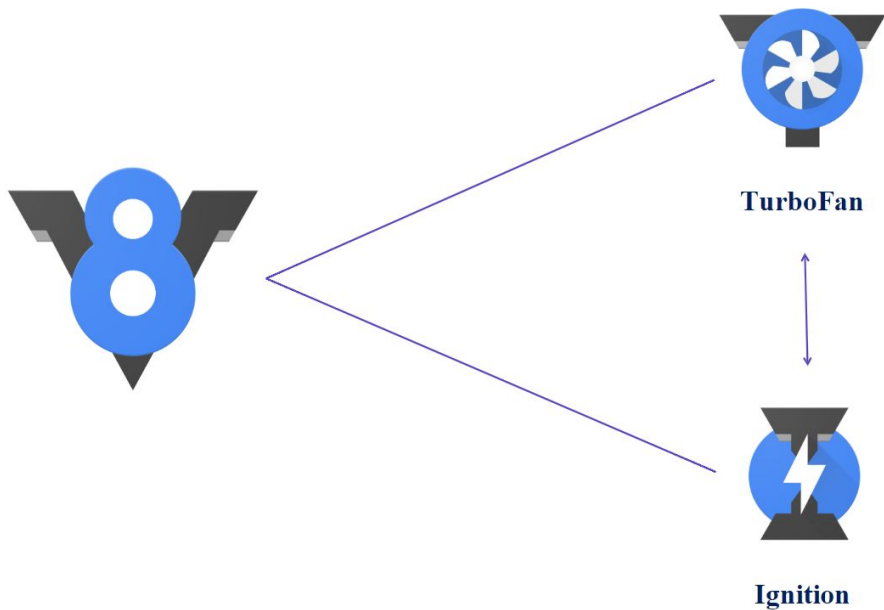
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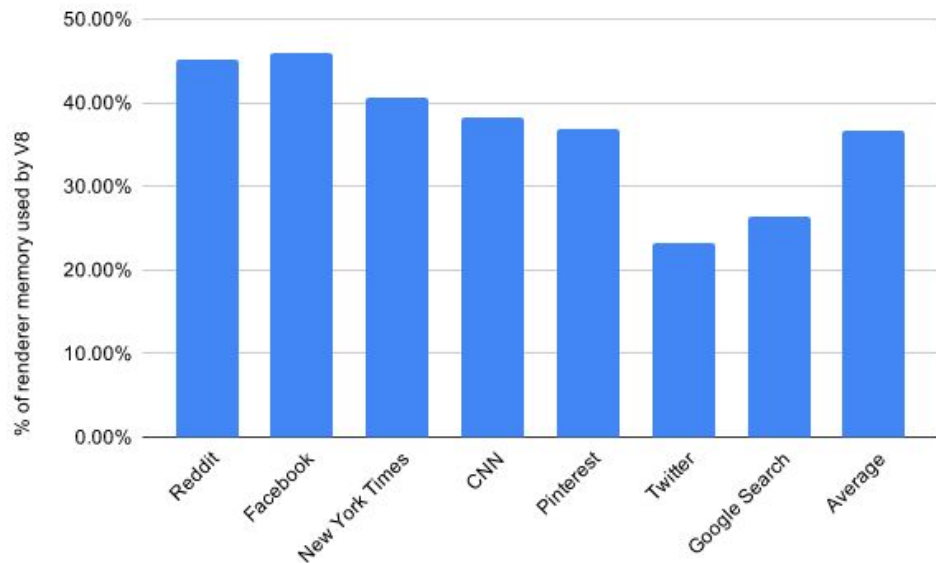
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关于V8



V8 引擎是用 C ++编写的开源高性能 JavaScript 和 WebAssembly 引擎, 它已被用于 Chrome 和 Node.js 等.

Pointer-compressed 背景



Chrome浏览不同页面V8消耗的内存占比
<https://v8.dev/blog/pointer-compression>

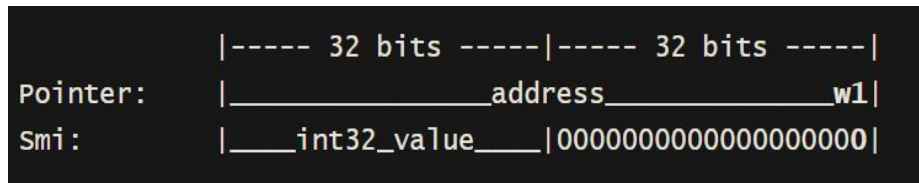
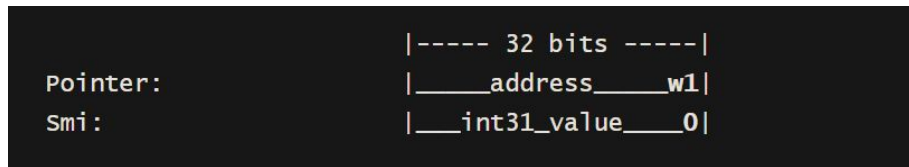
2014年,Chrome浏览器由32位切换为64位,这给浏览器带来更好的安全性、稳定性和性能,但与此同时也带来了内存消耗的增加:指针占用的内存由原来的4字节变为8字节。

而V8中所有的Object都是分配在Heap中,故每个Object都有一个指针指向该Object.

Value tagging

V8 中任何变量(objects, arrays, numbers or strings),都是存储在Heap中的,因此即使是int型变量,也需要一个指针. 为了降低指针带来的内存消耗,V8采用了 Tagged pointer技术.

Tagged pointer利用指针无论是32位架构还是64位架构,低2bit的值总是0的特点,将一些tag存储在低2bit中.V8利用tag来区别是否是Pointer还是Smi,Pointer是Weak还是Strong.



上图: 32位处理器架构的Value tagging

下图: 64位处理器架构的Value tagging

来源:<https://v8.dev/blog/pointer-compression>

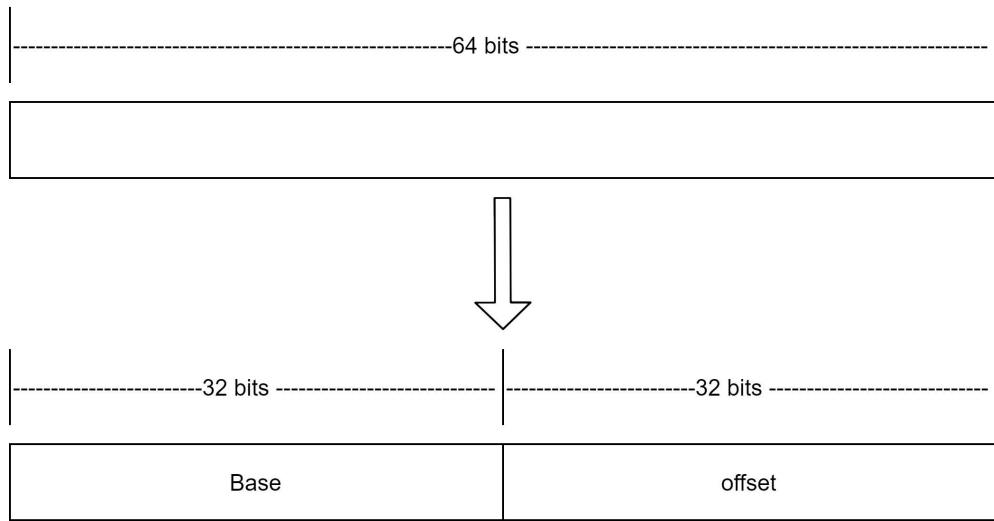
Pointer-compressed 基本原理

Pointer-Compressed 是谷歌用来减少内存占用的方法之一, 原理:

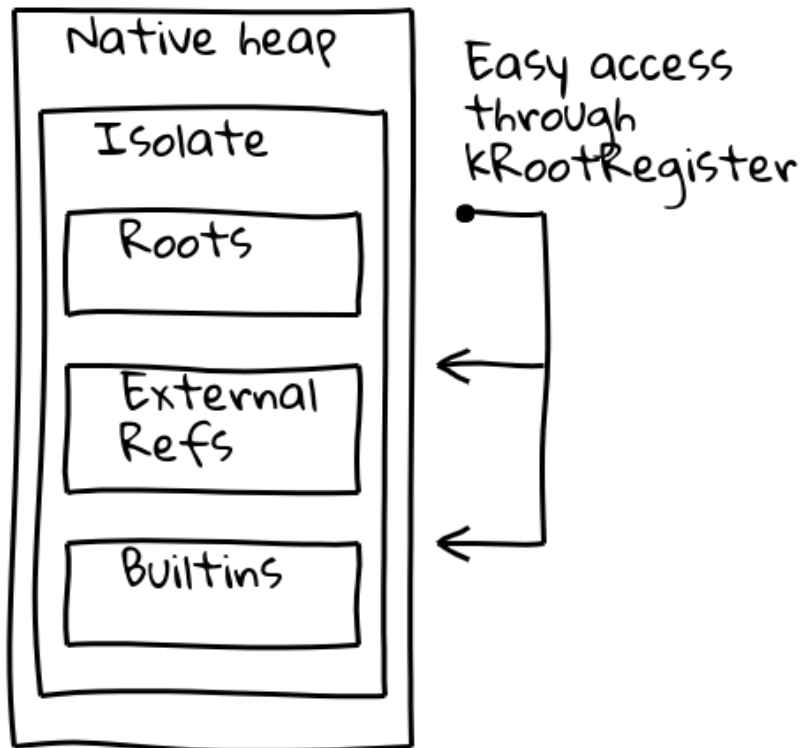
- 采用32位的offset和32bit base来代替64位的指针
- base由一个全局变量持有, 指针只需存储offset到内存中

为了适配指针压缩, 需要满足以下两个条件:

- 所有 V8 objects 都要分配在4GB范围中
- 将指针用offset重新解释, 需要一个全局变量保存base



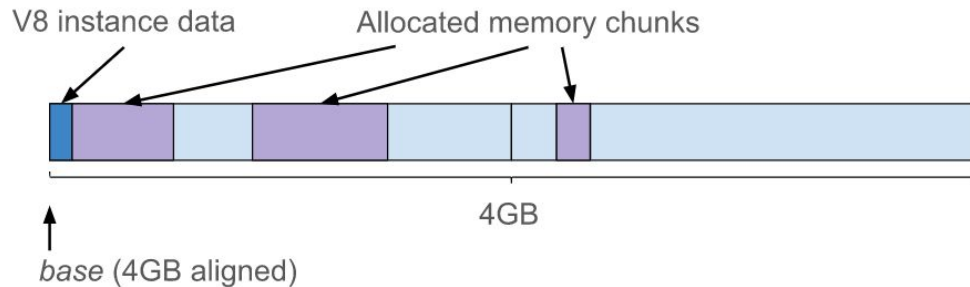
V8的先天优势



V8的isolate布局

来源:<https://v8.dev/blog/pointer-compression>

Pointer-compress下内存布局



Base对齐到4GB内存的起始位置
<https://v8.dev/blog/pointer-compression>

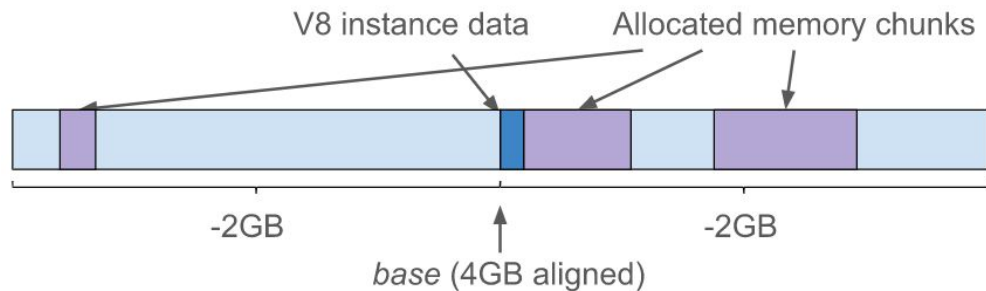
```
uint64_t uncompressed_tagged;  
uint32_t compressed_tagged = uint32_t(uncompressed_tagged);
```

compression

```
uint32_t compressed_tagged;  
  
uint64_t uncompressed_tagged;  
if (compressed_tagged & 1) {  
    // pointer case  
    uncompressed_tagged = base + uint64_t(compressed_tagged);  
} else {  
    // Smi case  
    uncompressed_tagged = int64_t(compressed_tagged);  
}
```

decompression

Pointer-compress下几种内存布局



Base对齐到4GB内存的中间位置
<https://v8.dev/blog/pointer-compression>

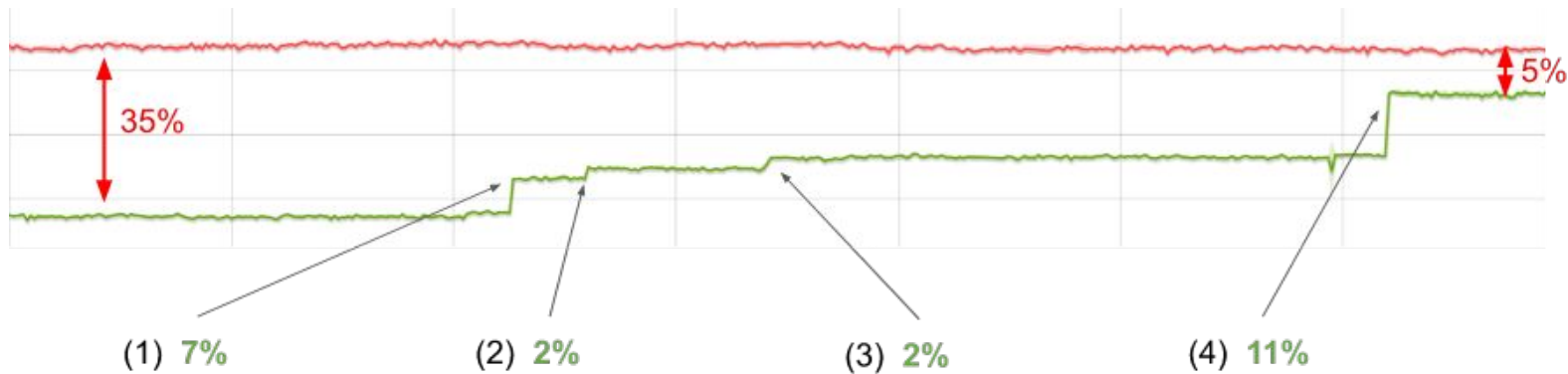
```
int32_t compressed_tagged;  
  
// Common code for both pointer and Smi cases  
int64_t uncompressed_tagged = int64_t(compressed_tagged);  
if (uncompressed_tagged & 1) {  
    // pointer case  
    uncompressed_tagged += base;  
}
```

decompression

```
int32_t compressed_tagged;  
  
// Same code for both pointer and Smi cases  
int64_t sign_extended_tagged = int64_t(compressed_tagged);  
int64_t selector_mask = -(sign_extended_tagged & 1);  
// Mask is 0 in case of Smi or all 1s in case of pointer  
int64_t uncompressed_tagged =  
    sign_extended_tagged + (base & selector_mask);
```

无分支的decompression

Pointer-compress带来的性能损失及其优化



Octane's score on x64 architecture
<https://v8.dev/blog/pointer-compression>

优化一 Branchful version was 7% faster on x64

Decompression	Branchless	Branchful
Code	<pre>movsxlq r11, [...] movl r10, r11 andl r10, 0x1 negq r10 andq r10, r13 addq r11, r10</pre>	<pre>movsxlq r11, [...] testb r11, 0x1 jz done addq r11, r13 done:</pre>
Summary	20 bytes	13 bytes
	6 instructions executed	3 or 4 instructions executed
	no branches	1 branch
	1 additional register	

X64下decompression汇编代码对比
<https://v8.dev/blog/pointer-compression>

Decompression	Branchless	Branchful
Code	<pre>ldur w6, [...] sbfx x16, x6, #0, #1 and x16, x16, x26 add x6, x16, w6, sxtw</pre>	<pre>ldur w6, [...] sxtw x6, w6 tbz w6, #0, #done add x6, x26, x6 done:</pre>
Summary	16 bytes	16 bytes
	4 instructions executed	3 or 4 instructions executed
	no branches	1 branch
	1 additional register	

ARM64下decompression汇编代码对比
<https://v8.dev/blog/pointer-compression>

同样Arm64下有分支的代码同样比无分支代码快

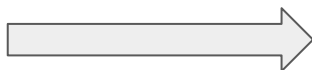
现代CPU分支预测技术已经十分强大, 影响类似这样的代码执行的效率主要取决于代码的执行指令数量或code size。

优化二: Eliminate decompressions directly followed by compressions 2%

```
Reduction DecompressionElimination::ReduceCompress(Node* node) {  
    DCHECK(IrOpcode::IsCompressOpcode(node->opcode()));  
  
    DCHECK_EQ(node->InputCount(), 1);  
    Node* input_node = node->InputAt(0);  
    IrOpcode::Value input_opcode = input_node->opcode();  
    if (IrOpcode::IsDecompressOpcode(input_opcode)) {  
        DCHECK(IsValidDecompress(node->opcode(), input_opcode));  
        DCHECK_EQ(input_node->InputCount(), 1);  
        return Replace(input_node->InputAt(0));  
    } else if (IsReducibleConstantOpcode(input_opcode)) {  
        return Replace(GetCompressedConstant(input_node));  
    } else {  
        return NoChange();  
    }  
}
```

优化三: 去除多余的指令 2%

```
movl rax, <mem> // load  
movlsxlq rax, rax // sign extend
```

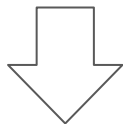


```
movlsxlq rax, <mem>
```

优化四: Updated the pattern matching 与 Decompressed Optimize 11%

优化五: Smi-corrupting 2.5%

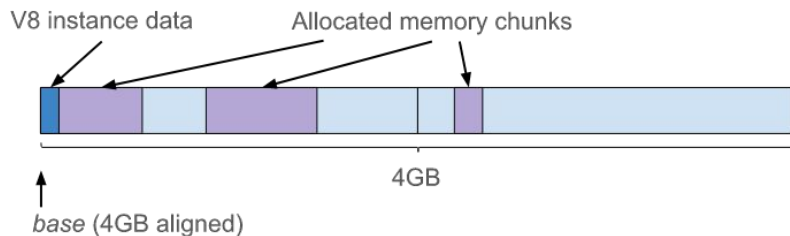
```
int64_t uncompressed_tagged = int64_t(compressed_tagged);  
if (uncompressed_tagged & 1) {  
    // pointer case  
    uncompressed_tagged += base;  
}
```



```
int64_t uncompressed_tagged = base + int64_t(compressed_tagged);
```

若采用的内存布局是 base 对齐到 4GB 内存空间的起始位置时, int64_t 可变为 uint64_t, 将符号扩展变为零扩展, 进一步提高性能

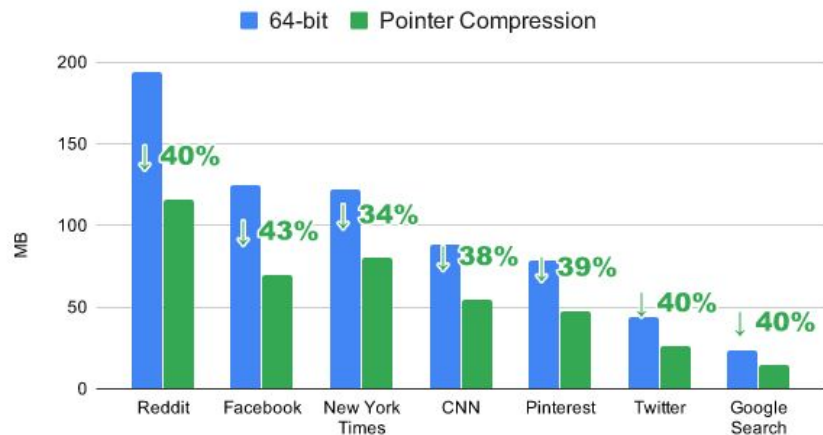
base points to the beginning, 4 GB aligned



Base 对齐到 4GB 内存的起始位置
<https://v8.dev/blog/pointer-compression>

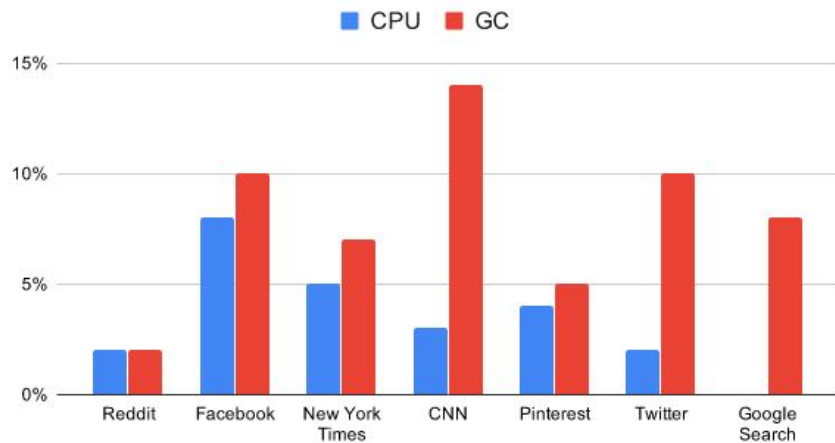
结果

V8 heap memory



采用指针压缩前后内存对比
<https://v8.dev/blog/pointer-compression>

Performance improvements



采用指针压缩后CPU和GC的性能提升
<https://v8.dev/blog/pointer-compression>