The following MIPS instruction adds the values inside of registers \$t2 and \$t3 and stores the result in \$t1:

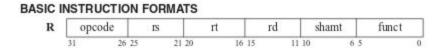
add \$t1 \$t2 \$t3

One thing you'll be asked to do in this class is to convert this instruction to machine code, which is what the assembler does for us. To do this you'll need the MIPS Green Sheet to decode the instructions.

The first thing you'll notice if you lookup the 'add' instruction is that it is an R-type instruction:

$$\begin{array}{c|ccccc} \textbf{CORE INSTRUCTION SET} & \textbf{OPCODE} \\ & & & & & & & / \text{FUNCT} \\ \hline NAME, MNEMONIC & MAT & OPERATION (in Verilog) & (Hex) \\ Add & add & R & R[rd] = R[rs] + R[rt] & (1) & 0 / 20_{hex} \\ \end{array}$$

Looking further down the sheet you will see the instruction format for an R-type instruction:



The first thing we'll do is look up the opcode, which 6-bits long, for the 'add' instruction: $0_{10} = \frac{000000_{2}}{10}$.

Next we'll decode the registers. Before we do, let me point out an important subtelty between MIPS syntax and the instruction format. Notice that the 'add' instruction has the three registers in the order: rd, rs, rt. But notice the R-type instruction format has them in a slightly different order: rs, rt, rd. Applying this to our example, rs = \$t2, rt = \$t3, and rd = \$t1. I hope this distinction is clear! Ok, moving on! Referring to our Green Sheet again, we see the following mapping of registers to integers:

REGISTER NAME, NUMBER, USE, CALL CONVENTION

NUMBER	USE	PRESERVED ACROSS A CALL?		
0	The Constant Value 0	N.A.		
1	Assembler Temporary	No		
2-3	Values for Function Results and Expression Evaluation	No		
4-7	Arguments	No		
8-15	Temporaries	No		
16-23	Saved Temporaries	Yes		
24-25	Temporaries	No		
26-27	Reserved for OS Kernel	No		
28	Global Pointer	Yes		
29	Stack Pointer	Yes		
30	Frame Pointer	Yes		
31	Return Address	No		
	0 1 2-3 4-7 8-15 16-23 24-25 26-27 28 29 30	0 The Constant Value 0 1 Assembler Temporary Values for Function Results and Expression Evaluation 4-7 Arguments 8-15 Temporaries 16-23 Saved Temporaries 24-25 Temporaries 26-27 Reserved for OS Kernel 28 Global Pointer 29 Stack Pointer 30 Frame Pointer		

Since \$t0 = 8, that means \$t1 = 9_{10} , \$t2 = 10_{10} , and \$t2 = 11_{10} . Converting these to 5-bit binary values (as the instruction format tells us to do), we get: \$t1 = $\frac{01001_2}{2}$, \$t2 = $\frac{01010_2}{2}$, and \$t3 = $\frac{01011_2}{2}$. Great!

The next step is simple. In nearly all cases, the 5-bit shamt (shift amount) will be 0_{10} , which is 00000_{2} . Done.

Finally, we need to decode the 6-bit function code. For this we just consult the Green Sheet for the 'add' instruction again and we see that the function code is $20_{16} = 32_{10} = \frac{100000_2}{2}$.

Now that we've decode all of the fields (highlighted in yellow), let's assemble them according to the R-type instruction format:

000000	01010	01011	01001	00000	100000
opcode	rs	rt	rd	shamt	func

The 32-bit binary string on the top row is the machine code equivalent of the MIPS instruction!

You will sometimes be asked to give the hex representation of an instruction. Since 1 hex digit = 4 binary bits, let's break the 32-bit string up into 8 4-bit chunks and then convert to hex:

0000	0001	0100	1011	0100	1000	0010	0000
0	1	4	В	4	8	2	0

And the bottom row of the above table is the hex representation of the MIPS 'add' instruction!