20171099 Parallel Prefix algorithm on an array of 8 elements parallel- frefix algo for (i=1 to n/2) { l'in farallel bi = a2i-1 + a2i find be recursively and store in Co for (i=1 ton) { // in parallel if (i is even) { Si = Ci/2 cosif (i=1) S1=C1 if (i is odd) f Si = COOD C:1/2 + ai end return S. end upward traversal Time Com: 0 (logn) downward work: o(n) retain 10 Prefixsum parents unde+self retain

Harelita Sharme (2) Parallel search algo P= Processors = 4, no. to be searched = 10 = 2 1 n=16 0 11 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 1 0 1. divide array into p parts [l=0; r=n+1) [R=17] 1 A3 A4 [13 14]5 16 1 5678 eneck & with first element of each 1 0 part parallely 1 P1 211 2 2 1 --farallel 1 $P_2 \quad \mathcal{U} \rightarrow 5 ; \quad \mathcal{X} > 5$ 4. checks P3 279;2>97_ we know 279 k x < 13 2 + 13; X<13 Py So, x lies in Az. here Divide (Recuesive Step) A3 9 10 11 12 in p parts 2 12 11) 10 91 Check uz parallely each part now that each part has only one 4 checks elemen x - 9 --- MATCH!. 2-10 $\chi - 11$ u-12 time: O(logpn) = 2 p processors -> Work = 4+4=8

 $\boxed{3} \rightarrow \boxed{4} \rightarrow \boxed{5} \rightarrow \boxed{9} \rightarrow \boxed{6} \rightarrow \boxed{1} \rightarrow \boxed{7}$ 1 & i & n do in parallel R[i] = 1 R[i] = 0 if node i is last while S[i] + S[S[i]] do R[i] = R[i] + R[S[i]] S[i] = 5[s[i]] end Iteration 1 Time: O(logn) W: O(n logn) Iteration 2 Iteration 3:

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