Free All The Things

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Free All The Things

- · well known: free monads
- maybe known: free applicatives
- free monoids
- free <you name it>

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This Talk

- explain the technique behind "Free X"
- apply the technique to different examples
- takeaway: it's easier than you thought
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/lambdaconf

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Let's Get Started

Ready?

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What's The Problem

A free functor is left adjoint to a forgetful functor



What's The Problem

A free functor is left adjoint to a forgetful functor

what's the problem?



What Is Free

A free "thing" **FreeA** on a type(class) A is a A and a function

def inject(x: A): FreeA

such that for any other "thing" B and a function

val $f: A \Rightarrow B$

there exists a unique homomorphism g such that

What Is Free

- still complicated!
- rejoice: there is a pretty mechanical recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - Check laws

Why Free

- nice API using typeclass
- use Free X as if it was X
- program reified into datastructure
- structure can be analyzed/optimized
- one program many interpretations

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Disclaimer Before We Start

- deep embeddings / initial encoding / data structure representation
- not: finally tagless, optimization

Freeing The Monad

The Monad Typeclass

```
trait Monad[F[_]] {
   def pure[A](x: A): F[A]

def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

Give Me The Laws

```
// Left identity
   pure(a).flatMap(f) === f(a)
3
  // Right identity
   fa.flatMap(pure) === fa
6
  // Associativity
   fa.flatMap(f).flatMap(g) ===
8
     fa.flatMap(a \Rightarrow f(a).flatMap(g))
```

```
trait Monad[F[ ]] {
1
       def pure [A](x:A):F[A]
3
       def flatMap[A, B](fa: F[A])(f: A \Rightarrow F[B]): F[B]
4
5
```

- remember our recipe
 - Define AST
 - Add Inject
 - Write Interpreter
 - · Check laws

Freeing The Monad

```
sealed abstract class Free[F[ ], A]
1
2
      final case class Pure[F[ ], A](a: A)
3
          extends Free[F, A]
4
5
      final case class FlatMap[F[], A, B](
6
          fa: Free[F, A].
7
          f: A \Rightarrow Free[F, B]
8
          extends Free[F, B]
9
10
      final case class Inject[F[_], A](fa: F[A])
11
          extends Free[F, A]
12
```

Freeing The Monad

```
implicit def freeMonad[F[_], A]: Monad[Free[F, ?]] =
new Monad[Free[F, ?]] {
    def pure[A](x: A): Free[F, A] = Pure(x)

def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] =
FlatMap(fa, f)
}
```

Interpreter

```
def runFree[F[_], M[_]: Monad, A](nat: F ~> M)(
    free: Free[F, A]): M[A] = free match {
    case Pure(x) => Monad[M].pure(x)
    case Inject(fa) => nat(fa)
    case FlatMap(fa, f) =>
        Monad[M].flatMap(runFree(nat)(fa))(x =>
        runFree(nat)(f(x)))
    }
}
```

What about the laws?

```
// The associativity law
  fa.flatMap(f).flatMap(g) ===
     fa.flatMap(fa, a \Rightarrow f(a).flatMap(q))
  val exp1 = FlatMap(FlatMap(fa, f), g)
  val exp2 = FlatMap(fa, (a: Int) => FlatMap(f(a), q))
3
  exp1 != exp2
```

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What about the laws?



The Laws

- actually, we don't satisfy them
- programmer: after interpretation it's no longer visible
- mathematician: that's not the free monad!
- tradeoff: during construction vs during interpretation

The Right Free Monad

- common transformation: associate flatMap's to the right
- avoids having to rebuild the tree repeatedly during construction
- how: during construction time

Transforming Free Monads

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

Transforming Free Monads

```
def flatMap[A, B](fa: Free[F, A])(
    f: A => Free[F, B]): Free[F, B] = fa match {
    case Pure(x) => f(x) // Left identity
    case Inject(fa) => FlatMap(Inject(fa), f)
    case FlatMap(ga, g) => // Associativity
    FlatMap(ga, (a: Any) => FlatMap(g(a), f))
}
```

We Freed Monads

- DSL with monadic expressiveness
- context sensitive, branching, loops, fancy control flow
- familiarity with monadic style for DSL
- big drawback: interpreter has limited possibilities

Freeing The Applicative

Freeing The Applicative

- free monads are great, but also limited
- we can't analyze the programs
- how about a smaller abstraction?

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Recall

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Applicative Typeclass

```
trait Applicative[F[_]] {
   def pure[A](x: A): F[A]

def ap[A, B](fab: F[A => B], fa: F[A]): F[B]
}
```

```
sealed abstract class FreeAp[F[ ]. A]
2
      final case class Pure [F], A (a: A)
3
          extends FreeAp[F. A]
4
5
      final case class Ap[F[ ], A, B](
6
          fab: FreeAp[F, A \Rightarrow B],
7
          fa: FreeAp[F, A])
8
          extends FreeAp[F, B]
9
10
      final case class Inject[F[\ ], A](fa: F[A])
11
          extends FreeAp[F, A]
12
```

Laws

```
1 // identity
   Ap(Pure(identity), v) === v
3
   // composition
   Ap(Ap(Ap(Pure(.compose), u), v), w) ===
     Ap(u, Ap(v, w))
  // homomorphism
   Ap(Pure(f), Pure(x)) === Pure(f(x))
10
   // interchange
   Ap(u, Pure(y)) === Ap(Pure((v)), u)
12
```

Don't Forget The Laws

```
def ap[A, B](fab: FreeAp[F, A \Rightarrow B],
1
                    fa: FreeAp[F, A]): FreeAp[F, B] =
2
        (fab, fa) match {
3
          case (Pure(f), Pure(x)) =>
4
            Pure(f(x)) // homomorphism
5
          case (u. Pure(v)) =>
6
            Ap(Pure((f: A \Rightarrow B) \Rightarrow f(y)), u) // interchange)
7
          case ( , ) => Ap(fab, fa)
8
```

Running FreeApplicatives

```
def runFreeAp[F[\_], M[\_]: Applicative, A](
1
         nat: F \sim M)(free: FreeAp[F, A]): M[A] =
2
       free match {
3
         case Pure(x) => Applicative[M].pure(x)
4
         case Inject(fa) => nat(fa)
5
         case Ap(fab, fa) =>
6
           Applicative[M]
7
             .ap(runFreeAp(nat)(fab), runFreeAp(nat)(fa))
8
```

We Freed Applicatives

- DSL with applicative expressiveness
- context insensitive
- pure computation over effectful arguments
- more freedom during interpretation

Freeing The Functor

And Once Again

- Define AST
- Add Inject
- Write Interpreter
- Check laws

The Functor Typeclass

```
trait Functor[F[_]] {
   def map[A, B](fa: F[A])(f: A => B): F[B]
}
```

Freeing The Functor

```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

case class Inject[F[_], A](fa: F[A])
    extends FreeFunctor[F, A]
```

Freeing The Functor

```
sealed abstract class FreeFunctor[F[_], A]

case class Fmap[F[_], X, A](fa: F[X])(f: X => A)
    extends FreeFunctor[F, A]

def inject[F[_], A](value: F[A]) =
    Fmap(value)(identity)
```

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Clean Code Police



only one subclass?

Freeing The Functor

```
sealed abstract class Fmap[F[_], A] {
1
       type X
2
       def fa: F[X]
       def f: X => A
4
6
     def inject[F[_], A](v: F[A]) = new Fmap[F, A] {
7
       type X = A
8
       def fa = v
9
       def f = identity
10
11
```

Freeing The Functor

```
sealed abstract class Coyoneda [F[], A]
1
       type X
2
       def fa: F[X]
       def f: X => A
4
6
     def inject[F[], A](v: F[A]) = new Coyoneda[F, A] {
7
       type X = A
8
       def fa = v
9
       def f = identity
10
11
```

Free Functor Instance

```
implicit def coyoFun[F[_]]: Functor[Coyoneda[F, ?]] =
1
        new Functor[Covoneda[F, ?]] {
2
          def map[A, B](coyo: Coyoneda[F, A])(
3
              q: A \Rightarrow B: Covoneda[F, B] =
4
            new Coyoneda[F, B] {
5
              type X = coyo.X
6
              def fa = coyo.fa
7
              def f = q.compose(coyo.f)
8
9
10
```

Free Functor Interpreter

```
def runCoyo[F[_]: Functor, A](
coyo: Coyoneda[F, A]): F[A] =
Functor[F].map(coyo.fa)(coyo.f)
```

We Freed Functors

- DSL with hmm functorial expressiveness?
- map fusion! (functor law)
- boring interpreter, though
- still fun!

Freeing The Monoid

The Monoid Typeclass

```
trait Monoid[A] {
def empty: A
def combine(x: A, y: A): A
}
```

The Free Monoid — First Try

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Inject[A](x: A) extends FreeMonoid[A]

case class Combine[A](x: FreeMonoid[A],

y: FreeMonoid[A])

extends FreeMonoid[A]
```

The Laws

```
1  // left identity
2  empty |+| x === x
3
4  // right identity
5  x |+| empty === x
6
7  // associativity
8  1 |+| (2 |+| 3) === (1 |+| 2) |+| 3
```

The Laws and Free Monoid

- let's try to enforce those laws in our structure
- goal: correct by construction
- arbitrary decision: associate left vs right

Fixing Associativity

```
sealed trait NotCombine[+A]
1
2
      sealed abstract class FreeMonoid[+A]
3
4
      case object Empty
5
          extends FreeMonoid[Nothing]
6
          with NotCombine[Nothing]
7
8
      case class Inject[A](x: A)
9
          extends FreeMonoid[A]
10
          with NotCombine[A]
11
12
      case class Combine[A](x: NotCombine[A],
13
                             v: FreeMonoid[A])
14
          extends FreeMonoid[A]
15
```

The Problem With Neutral Elements

- get rid completely? not possible
- limit ourselves to a single element
- restrict Combine to have only real values on the left side
- goal: minimal canonical structure

Minimizing Structure — Extract Inject

```
case class Inject[A](x: A)

sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: Inject[A], y: FreeMonoid[A])

extends FreeMonoid[A]
```

Minimizing Structure — Remove Inject

```
sealed abstract class FreeMonoid[+A]

case object Empty extends FreeMonoid[Nothing]

case class Combine[A](x: A, y: FreeMonoid[A])
extends FreeMonoid[A]
```

The Monoid Instance

```
implicit def monoid[A]: Monoid[FreeMonoid[A]] =
1
       new Monoid[FreeMonoid[A]] {
2
          override def empty = Empty
3
          override def combine(
4
              x: FreeMonoid[A],
5
              y: FreeMonoid[A]): FreeMonoid[A] = x match {
6
            case Empty
7
                                => V
            case Combine(h, t) => Combine(h, combine(t, y))
8
9
10
```

Minimizing Structure — List

```
case object Nil extends List[Nothing]
case class Cons[A](head: A, tail: List[A])
extends List[A]
```

sealed abstract class List[+A]

1

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We Freed Monoids

- DSL for combining things
- works beautifully with folds
- interpretation can be in parallel (associativity)

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Now That We Can Free Anything



What should we free?

Credit Where It's Due

- Once upon a time: https://engineering.wingify.com/posts/Free-objects/
- use free boolean algebra to define DSL for event predicates
- credits to Chris Stucchio (@stucchio)

Let's Free A Boolean Algebra

- Define AST
- Add Inject
- Write Interpreter
- Check laws

Boolean Algebras

```
trait BoolAlgebra[A] {
   def tru: A
   def fls: A

def not(value: A): A

def and(lhs: A, rhs: A): A
   def or(lhs: A, rhs: A): A
}
```

```
sealed abstract class FreeBool[+A]
1
2
     case object Tru extends FreeBool[Nothing]
3
     case object Fls extends FreeBool[Nothing]
4
5
     case class Not[A](value: FreeBool[A])
6
          extends FreeBool[A]
7
     case class And[A](lhs: FreeBool[A], rhs: FreeBool[A])
8
          extends FreeBool[A]
9
     case class Or[A](lhs: FreeBool[A], rhs: FreeBool[A])
10
          extends FreeBool[A]
11
     case class Inject[A](value: A) extends FreeBool[A]
12
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A => B)(
1
         implicit B: BoolAlgebra[B]): B = {
2
       fb match {
3
         case Tru
                      => B.tru
4
         case Fls
                       => B.fls
5
         case Inject(v) \Rightarrow f(v)
6
         case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                      => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                       => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                       => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

```
def runFreeBool[A, B](fb: FreeBool[A])(f: A \Rightarrow B)(
1
          implicit B: BoolAlgebra[B]): B = {
2
        fb match {
3
          case Tru
                       => B.tru
4
          case Fls
                       => B.fls
5
          case Inject(v) \Rightarrow f(v)
6
          case Not(v) => B.not(runFreeBool(v)(f))
7
          case Or(lhs, rhs) =>
8
            B.or(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
9
          case And(lhs, rhs) =>
10
            B.and(runFreeBool(lhs)(f), runFreeBool(rhs)(f))
11
12
13
```

Using Free Bool

- that was simple (though boilerplate-y)
- what can we do with our new discovered structure
- reminder: boolean operators
 - true, false
 - and, or
 - · xor, implies, nand, nor, nxor

Free Bool Example: Search

sealed trait Search

1

```
case class Term(t: String) extends Search
case class After(date: Date) extends Search
case class InText(t: String) extends Search
case class InUrl(url: String) extends Search
// and the usual smart ctors
```

Free Bool Example: Search

after sneaking in syntactic sugar behind the scenes:

```
val search = term("FP") &
after("20180101") &
 !(term("Java") | inText("spring")) &
inUrl("lambdaconf")
```

The Site Type

```
case class Site(terms: List[String],
url: String,
indexedAt: Date,
text: String)
```

• all predicates are run against a collection of these

Free Bool Example: Search

```
def evalSearch(pred: FreeBool[Search])(
1
         site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                => site.terms.contains(t)
         case Term(t)
4
         case After(d)
                                 => site.indexedAt > d
5
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                 => site.url.contains(w)
8
9
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

Free Bool Example: Search

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                  => site.terms.contains(t)
          case Term(t)
4
         case After(d)
                                 => site.indexedAt > d
5
          case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                  => site.url.contains(w)
8
9
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

Free Bool Example: Search

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                 => site.terms.contains(t)
          case Term(t)
4
          case After(d)
                                  => site.indexedAt > d
5
         case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                  => site.url.contains(w)
8
9
        runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

Free Bool Example: Search

```
def evalSearch(pred: FreeBool[Search])(
1
          site: Site): Boolean = {
2
       def nat(s: Search): Boolean = s match {
3
                                 => site.terms.contains(t)
          case Term(t)
4
          case After(d)
                                 => site.indexedAt > d
5
          case InText(t: String) => site.text.contains(t)
6
         case InUrl(w)
                                  => site.url.contains(w)
8
9
       runFreeBool(pred)(nat)
10
11
12
     val result = Sites.all().filter(evalSearch(search))
13
```

But Wait There's More

- short circuiting and other optimization
- what if you don't have all the information?
 - partially evaluate predicates
 - if evaluates successfully, done
 - · else, send it on
- core language vs extension
 - · Chris also demonstrates extension
 - translate a rich language to base instructions
 - with all the advantages

Optimizing Boolean Algebras

```
def optimize[A](fa: FreeBool[A]): FreeBool[A] =
       fa match {
2
         case Tru => Tru
3
         case Fls => Fls
4
         case Inject(v) => Inject(v)
5
         case Not(Not(v)) => v
6
         case Not(v) => Not(v)
7
         case Or(Tru, ) => Tru
8
         case Or( , Tru) => Tru
9
         case Or(x, y) => Or(x, y)
10
         case And(Fls, ) => Fls
11
         case And(_, Fls) => Fls
12
         case And(x, y) => And(x. y)
13
14
```

Partial Evaluation

- idea: you might have only partial information
- evaluate as much as possible
- optimal: we can already reduce without needing more information
- otherwise: send it on (JSON, Protobuf, ...)

Partial Evaluation

```
def partialEvaluator[A, B](p: FreeBool[A])(
1
          f: A => Option[B])(implicit B: BoolAlgebra[B])
2
        : Either[FreeBool[A], B] = p match {
3
        case Tru => Right(B.tru)
4
        case Inject(v) \Rightarrow f(v).toRight(p)
5
        case Or(lhs, rhs) =>
6
          val(l, r) = (partialEvaluator(lhs)(f),
7
                         partialEvaluator(rhs)(f))
8
         // perform short circuiting
9
          ???
10
        case => ???
11
12
```

Partial Evaluation

```
// fulltext not available
     case class SiteMetadata(terms: List[String],
2
                              url: Strina.
3
                              indexedAt: Date)
4
     def partially(meta: SiteMetadata)(
         p: Search): Option[Boolean] = p match {
6
       case Term(t)
                               => Some(meta.terms.contains(t))
       case After(d)
                               => Some(meta.indexedAt > d)
8
       case InUrl(w)
                               => Some(meta.url.contains(w))
9
       case InText(t: String) => None
10
11
```

We Freed Boolean Algebras

- good example of underused free structure
- partial evaluation
- serialize the AST (JSON, Protobuf, Avro, ...)
- exercise: minimize AST representation

Resources

- Free Boolean Algebra by Chris Stucchio https://engineering.wingify.com/posts/Free-objects/
- Source Code: https://github.com/markus1189/ free-all-the-things/tree/lambdaconf

Go And Free All The Things!

Markus Hauck (@markus1189)

Introduction

Free Monad

Free Applicative

Free Functor

Free Monoid

Free Boolean Algebra

Conclusion

LambdaConf 2018

Going Deeper

- try to encode one of the normal forms for boolean algebras
- try to remove Inject cases from Monad and Applicative
- free Magmas
- define free X using alternative minimal set of ops of the typeclass