

# MathLedger: Reflexive Formal Learning and the Chain of Verifiable Cognition

—  
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## Abstract

Contemporary AI systems achieve extraordinary performance yet remain largely *opaque* and *non-verifiable*, creating a crisis of trust for safety-critical deployment and governance. We introduce *Reflexive Formal Learning (RFL)*, a symbolic and cryptographically verifiable learning paradigm in which every update is justified by a formally verified proof or abstention event recorded in an immutable ledger. The resulting “chain of verifiable cognition” constitutes a closed epistemic loop bridging logic, cryptography, and learning dynamics. RFL is shown to be a discrete, non-differentiable analogue of stochastic gradient descent whose learning signal arises from verified truth rather than statistical loss, enabling provable convergence of reasoning under a verifiable ledger substrate.

**Figure 1 (to be generated): Chain of Verifiable Cognition**  
User input → PoA → Ledger  $r_t$  → UI  $u_t$  → RFL → update.

Figure 1: End-to-end epistemic loop (objects and arrows as shown; all edges attested).

## 1 Introduction

Mathematical reasoning systems increasingly require not only correctness but verifiable *cognitive integrity*—a guarantee that each inference can be cryptographically traced to a validated source. *MathLedger* unifies formal verification, machine learning, and cryptographic attestation into a single epistemic pipeline:

User-Verified Input → Proof-or-Abstain Reasoning → Ledger Attestation → UI Attestation → Reflexive Feedback

This ‘‘Chain of Verifiable Cognition’’ enables an AI system to learn exclusively from verified reasoning outcomes. The remainder of this paper formalizes *Reflexive Formal Learning (RFL)* as a symbolic analogue of gradient descent operating on verified events rather than numerical errors.

Table 1: Summary of Core Formal Constructs

Symbol	Meaning
$t \in \Pi$	Symbolic reasoning policy at time $t$ (metric space $(\Pi; \ \cdot\ )$ )
$P_\pi$	Policy-induced event measure over event space $\mathcal{E}$
$e_t$	Reasoning event (proof or abstention) under $P_\pi$
$\mathcal{V}(e_t)$	Verification outcome $\in \{1; 0; \perp\}$
$\mathbb{V}(e)$	Numeric surrogate $\mathbf{1}\{\mathcal{V}(e) \neq 1\} \in \{0; 1\}$
$\mathcal{L}$	Ledger of proofs/abstentions (predictable process)
$\Phi$	Feedback map $\Phi : \{1; 0; \perp\} \times \Pi \rightarrow \Delta\Pi$
$\mathcal{U}$	Reflexive update $\Pi \times \{1; 0; \perp\} \times \Sigma \rightarrow \Pi$
$\Sigma$	Auxiliary signal space (prompts, contexts)
$\Delta\Pi$	Space of composable symbolic policy deltas
$\oplus$	Algebraic composition on $\Pi$ applying $\Delta \in \Delta\Pi$ to
$L$	Lipschitz factor of $\oplus$ s.t. $\ \oplus\Delta - \ \leq L\ \Delta\ _\Delta$
$\ \cdot\ _\Delta$	Norm on $\Delta\Pi$
$(T; )$	Mixing horizon and geometric rate (Assumption A2 <sup>0</sup> )
$r_t, u_t$	Reasoning/UI Merkle roots (dual attestation)
$\mathcal{I}_t$	Dual-attestation binder token (keyless or HMAC)
$H_t$	Epistemic entropy at step $t$
$\mathcal{J}(\cdot)$	Epistemic risk functional (Def. 1)
$\mathsf{L}(\cdot)$	Classical differentiable loss (SGD)
	Epistemic scaling exponent

## 2 Related Work and Theoretical Context

**Stochastic approximation and recent convergence extensions.** Classical almost-supermartingale convergence (Robbins–Siegmund) guarantees pointwise convergence under a summable residual term. Recent work (e.g., Liu, Xie, & Zhang 2025a) relaxes this to *square-summable* disturbances, proving convergence to a bounded set. Our stability results instantiate these extensions: when a persistent verifier bias acts as a non-summable zero-order term, we obtain bounded-set convergence with an  $O(\cdot v)$  radius. Our Markovian noise assumptions (Assumption A2' in §6) align with modern SA analyses that handle time-inhomogeneous Markov noise with mixing and Lipschitz drift. *Cor. 1 instantiates the square-summable residual regime of Liu–Xie–Zhang (2025), yielding bounded-set convergence under Markov noise.*

**Positioning.** RFL offers a distinct paradigm for verifiable learning. While recent work on *proof-guided language models* uses verifier feedback as a reward signal to guide policy search in a reinforcement learning loop , RFL uses the verifier’s binary decision as the *learning signal itself*, replacing a gradient with a symbolic update. Similarly, where *weak-to-strong generalization* frameworks train a strong model on labels generated by a weaker one , RFL’s policy updates are a direct, deterministic function of verified truth events, not supervised fine-tuning. The ledger’s immutable, attested history of reasoning also connects RFL to work on *proof-carrying data* and ledger-based AI audits , but with a closed-loop dynamic where the ledger actively drives learning.

### 3 Reflexive Formal Learning (RFL): Conceptual Definition

Let  $\Pi$  be a complete metric policy space with norm  $\|\cdot\|$  and  $P_\pi$  the event measure induced by  $\cdot$ .

**Update algebra.** We write  $\Delta\Pi$  for the space of composable symbolic policy deltas and  $\oplus$  for their algebraic composition in  $\Pi$ ; thus  $\oplus \Delta$  denotes applying  $\Delta \in \Delta\Pi$  to policy  $\cdot \in \Pi$ . We equip  $\Delta\Pi$  with a norm  $\|\cdot\|_\Delta$  and assume *norm compatibility*: there exists  $L_\oplus \geq 1$  with

$$\|\cdot \oplus \Delta - \cdot\| \leq L_\oplus \|\Delta\|_\Delta \quad \text{for all } \cdot \in \Pi; \Delta \in \Delta\Pi;$$

**Definition 1** (Epistemic risk). Given  $P_\pi$  and  $\mathfrak{P}(e) = \mathbf{1}\{\mathcal{V}(e) \neq 1\}$ ,

$$\mathcal{J}(\cdot) = \mathbb{E}_{e \sim P_\pi} [\mathfrak{P}(e)] = \Pr_{e \sim P_\pi} [\mathcal{V}(e) \neq 1];$$

*Range.* Since  $\mathfrak{P}(e) \in \{0; 1\}$ , we have  $0 \leq \mathcal{J}(\cdot) \leq 1$  for all  $\cdot$ .

### 4 The Reflexive Formal Learning (RFL) Update Law

**Update operator and abstention damping.** At each step  $t$ ,

$${}_{t+1} = {}_t \oplus {}_f \Phi(\mathcal{V}(e_t); {}_t); \quad \Phi : \{1; 0; \perp\} \times \Pi \rightarrow \Delta\Pi; \quad (1)$$

By norm compatibility, committing  ${}_{t+1} = {}_t \oplus {}_f \Phi(\cdot)$  implies  $\|{}_{t+1} - {}_t\| \leq {}_f L_\oplus \|\Phi(\mathcal{V}(e_t); {}_t)\|_\Delta$ .

## Toy example: one-step RFL update with pseudo-Lean

**Goal and tactic (pseudo-Lean).**

```
example (h1 : P → Q) (h2 : P) : Q := by
apply h1
exact h2
```

Let the current goal be  $g_t : Q$  under context  $\text{ctx}_t = \{h_1 : P \rightarrow Q; h_2 : P\}$ . The agent proposes tactic  $s_t = \text{apply } h_1$ , producing subgoal  $g'_t : P$ , then proposes  $s'_t = \text{exact } h_2$ .

**Event and verification.** Define the event

$$e_t = (g_t; \text{ctx}_t; s_t; \text{apply}; s'_t; \text{exact}) \quad \text{and} \quad v_t = \mathcal{V}(e_t) \in \{1; 0; \perp\}:$$

**Pattern features.** Let  $\text{pat}(\text{ctx}_t; s_t)$  denote a differentiable feature map.

**Policy parameterization.** Suppose  $\pi_t$  induces a score via parameters  $\theta_t$ :

$$\text{score}_t = \langle \pi_t; \text{pat}(\text{ctx}_t; s_t) \rangle \Rightarrow p_t(\text{pattern}) = (\text{score}_t):$$

**Symbolic deltas.** Instantiate

$$\Delta^+(\pi_t; e_t) \equiv \text{inc}_{\eta} \text{pat}(\text{ctx}_t; s_t); \quad \Delta^-(\pi_t; e_t) \equiv -\text{dec}_{\eta} \text{pat}(\text{ctx}_t; s_t);$$

with  $\|\Delta^{\pm}\|_{\Delta} \leq M$ .

**Cases (Proof-or-Abstain).**

- $v_t=1$ :  $\pi_{t+1} = \pi_t \oplus f \Delta^+(\pi_t; e_t)$ ;  $\mathcal{J}$  decreases in expectation.
- $v_t=\perp$ : no-op update; abstention logged.
- $v_t=0$ :  $\pi_{t+1} = \pi_t \oplus f \Delta^-(\pi_t; e_t)$ ; demotes the pattern.

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**Algorithm 1** RFL $\oplus$ MCGS Planner (fail-closed, dual-attested)

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**Require:**  $t, f$ , update  $\Phi$ , verifier (REPL), canonicalizers  $C_R, C_U$ , ledger  $\mathcal{L}$

- 1: Initialize frontier at root Lean state;  $E \leftarrow []$  . list of per-event binders
- 2: **while** frontier nonempty **do**
- 3:   Expand node using policy  $\tau_t$  to yield candidate events  $e_t$
- 4:    $(v_t; \text{trace}; \text{build}) \leftarrow \text{REPL}.\text{check}(e_t)$
- 5:   **if**  $v_t \neq 1$  **then**
- 6:      $\mathcal{L}.\text{abstain}(e_t)$ ; **continue**
- 7:   **end if**
- 8:    $P_t \leftarrow C_R(e_t; \text{trace}; \text{build})$ ;  $D_t \leftarrow C_U(\text{UI snapshot})$
- 9:    $r_t \leftarrow \text{Hash}(\text{R} : \| P_t)$ ;  $u_t \leftarrow \text{Hash}(\text{U} : \| D_t)$ ;  $I_t \leftarrow \text{Hash}(\text{BIND} \| r_t \| u_t)$
- 10:   **if**  $\neg \mathcal{L}.\text{verify}(P_t; D_t; I_t)$  **then**
- 11:      $\mathcal{L}.\text{abstain}(e_t)$ ; **continue**
- 12:   **end if**
- 13:    $\Delta_t \leftarrow \Phi(v_t; \tau_t)$ ;  $\tau_{t+1} \leftarrow \tau_t \oplus f\Delta_t$
- 14:    $\mathcal{L}.\text{commit}(e_t; r_t; u_t; I_t; \text{build})$ ;  $E.\text{append}(I_t)$
- 15:   Push children of  $e_t$  to frontier with priority from  $\tau_{t+1}$
- 16: **end while**
- 17:  $\text{epoch\_root} \leftarrow \text{Merkle}(E)$ ;  $\mathcal{L}.\text{finalize\_epoch}(\text{epoch\_root})$

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## 5 Dual Attestation and Security Model

The integrity of the RFL loop depends on cryptographically binding reasoning events to their presentation. We formalize this via a dual-attestation scheme and specify the security guarantees under a formal adversary model.

**Dual-Root Ledger.** At each step  $t$ , the system commits to two Merkle roots:

- **Reasoning Root ( $r_t$ ):** Root over the canonicalized sequence of formal reasoning steps (proof tactics, intermediate goals) composing  $e_t$ .
- **UI Root ( $u_t$ ):** Root over the canonicalized representation of the UI state (DOM/JSON, PNG, HAR log) that displays the outcome of  $e_t$ .

These roots are bound by a cryptographic token  $\mathcal{I}_t = \text{Hash}(\text{"BIND": } \| r_t \| u_t)$  with prefix-free domain separation. The tuple  $(r_t; u_t; \mathcal{I}_t)$  is recorded on the ledger.

**Domain tags and REPL provenance.** We extend domain separation with tags RPL: (Lean REPL provenance; toolchain/build IDs) and G: (geometry engine artifacts). These tags are included by  $C_R$  prior to Merkle, binding  $(r_t; u_t; \mathcal{I}_t)$  to verifier versions and domain-specific pipelines.

**Adversary Model and Guarantees.** We consider a PPT adversary acting as a malicious prover, verifier, or network observer. The ledger offers:

- **Collision Resistance:**  $\text{Hash}(\cdot)$  is collision-resistant, preventing distinct artifacts from sharing a root.
- **Non-Malleability:** Modifying proof/UI artifacts invalidates Merkle roots and  $\mathcal{I}_t$ .
- **Replay Resistance:** Updates are indexed by  $t$ ; replay  $(r_k; u_k; \mathcal{I}_k)$  at  $t > k$  is rejected; timestamps strengthen this.

## Cryptographic Hardening.

- **Canonicalization:** All structured data MUST be canonicalized before hashing. We mandate RFC 8785 JCS for JSON; deterministic PNG and equivalent for other types.
- **Domain Separation:** All hash inputs use prefix-free tags: e.g.,  $\text{Hash}(\text{"LEAF": " } \parallel \cdot)$ ,  $\text{Hash}(\text{"NODE": " } \parallel h_L \parallel h_R)$ ; distinct tags for reasoning/UI/binder.
- **Constant-Time Ops:** Cryptographic comparisons MUST be constant-time to avoid timing side channels.
- **Version Pinning:** Record versions/hashes of  $\mathcal{V}$ , hash function, and canonicalizers on-ledger.

**Zero-Knowledge Extensions.** For privacy, proofs may be replaced by ZK certificates (e.g., PLONK for fast verification and small proofs; STARKs for transparency and post-quantum resistance, at larger sizes).

## 6 Convergence and Stability of Reflexive Formal Learning

**Stepsizes.** We identify  $\gamma_t \equiv \gamma_f$  in the constant-stepsize case; otherwise  $\gamma_t$  denotes a decaying schedule used in the SA analysis.

**Lemma 1** (Extended Robbins–Siegmund). *Let  $\{Z_t\}$ ,  $\{X_t\}$ ,  $\{Y_t\}$ , and  $\{W_t\}$  be non-negative,  $\mathcal{F}_t$ -adapted random sequences. Suppose  $\mathbb{E}[Z_{t+1} | \mathcal{F}_t] \leq (1 + X_t)Z_t + Y_t - W_t$  a.s. If  $\sum_t X_t < \infty$  a.s. and  $\sum_t Y_t < \infty$  a.s., then  $Z_t$  converges a.s. to a finite random variable and  $\sum_t W_t < \infty$  a.s. Furthermore, if the summable condition on  $\{Y_t\}$  is relaxed to square-summable ( $\sum_t Y_t^2 < \infty$  a.s.) and the increments are bounded  $(Z_{t+1} - Z_t)_+ \leq B_t$  with  $\sum_t B_t^2 < \infty$  a.s., then  $\{Z_t\}$  converges a.s. to a bounded set.*

**Assumption 1** (Adaptivity and bounded updates (A1)). *Let  $\{\mathcal{F}_t\}$  denote the filtration generated by  $(s; e_s; \mathcal{V}(e_s))_{s \leq t}$ . The iterates  $x_t$ , reasoning events  $e_t$ , and verification outcomes  $\mathcal{V}(e_t)$  are  $\mathcal{F}_t$ -adapted. Moreover, the update increments are uniformly bounded: there exists  $M < \infty$  such that  $\|\Phi(\mathcal{V}(e_t); x_t)\|_\Delta \leq M$  almost surely for all  $t$ .*

**Assumption 2** (Verification-monotone descent (A2)). *There exist constants  $\gamma > 0$  and an  $\mathcal{F}_t$ -adapted error term  $\epsilon_t \geq 0$  with  $\sum_t \mathbb{E}[\epsilon_t] < \infty$  such that the epistemic risk satisfies*

$$\mathbb{E}[\mathcal{J}(x_{t+1}) - \mathcal{J}(x_t) | \mathcal{F}_t] \leq -\gamma \Pr(\mathcal{V}(e_t) = 1 | \mathcal{F}_t) + \epsilon_t$$

for all  $t$ .

**Assumption 3** (Local linearization and contraction (A3)). *There exists a neighborhood  $\mathcal{N}$  of  $x^*$  in which the averaged update map  $\mathcal{T}(\cdot) := \mathbb{E}[\cdot \oplus \gamma_f \Phi(\mathcal{V}(e_t); \cdot)]$  is Gâteaux differentiable and Lipschitz. The Jacobian  $D\mathcal{T}(x^*)$  has spectral radius  $< 1$ , yielding a contraction on  $\mathcal{N}$  in the ambient norm.*

**Assumption 4** (Markovian Noise and Mixing (A2')). *The event stream  $\{e_t\}$  is generated via a policy-dependent Markov process  $\{X_t\}$  on a state space  $\mathcal{X}$ . For each fixed policy  $\pi$ , the process has a unique stationary distribution  $\pi$ . The process is uniformly geometrically ergodic: uniformly over  $x$ , there exist  $T \in \mathbb{N}$  and  $\gamma \in (0; 1)$  such that  $\|\mathbb{P}_\pi(X_T \in \cdot | X_0 = x) - \pi(\cdot)\|_{\text{TV}} \leq \gamma$  for all  $x$ . The one-step expected cost change and transition probabilities are Lipschitz in  $\cdot$ . In addition, there exists an adapted sequence  $B_t \geq 0$  with  $(\mathcal{J}(x_{t+1}) - \mathcal{J}(x_t))_+ \leq B_t$  a.s. and  $\sum_t B_t^2 < \infty$  a.s. (implied by A0 and bounded  $\Phi$  under standard stepsizes).*

*Remark 1.* Uniform geometric ergodicity is standard in SA; weaker mixing (Doeblin minorization or spectral-gap conditions) can suffice in place of uniformity.

**Theorem 1** (Almost-Sure Convergence with Lyapunov Potential). *If Assumptions 1–2 hold and  $\mathbb{E}[\Delta\mathcal{J}_t|\mathcal{F}_t] \leq -c\|\Phi\|^2 + \gamma_t \mathbb{E}[\gamma_t] < \infty$ , then  $\mathcal{J}(\gamma_t) \rightarrow 0$  a.s. and  $\gamma_t \rightarrow \gamma^*$  where  $\mathcal{V}(\gamma^*) = 1$ .*

**Corollary 1** (Convergence to a bounded set under square-summable noise). *Let  $X_t = \mathcal{J}(\gamma_t)$ . Assume  $\gamma_t \gamma_t = \infty$ ,  $\gamma_t \gamma_t^2 < \infty$ . If the RFL dynamics satisfy Lemma 1 with square-summable disturbance and bounded increments, then  $\{X_t\}$  converges a.s. to a bounded set; pointwise convergence is recovered when the disturbance term is summable.*

**Corollary 2** (Linear Convergence of Epistemic Risk). *Under Theorem ??, if  $\Pr[\mathcal{V}(e_t) = 1] \geq c > 0$  for all non-optimal policies, then  $\mathbb{E}[\mathcal{J}(\gamma_t)]$  converges to its limit at a linear rate.*

**Corollary 3** (Local Stability of the Optimal Policy). *Under Assumption 3, the fixed point  $\gamma^*$  is locally asymptotically stable: any  $\gamma$  within the contraction basin converges to  $\gamma^*$  under RFL dynamics.*

**Proposition 1** (Abstention as damping). *If  $\Pr[\mathcal{V}(e_t) = \perp] = \gamma_t$  and abstention cost  $\text{cost}(\perp) = c_\perp \geq 0$  (captures operational abstention cost), then*

$$\mathbb{E}[\mathcal{J}(\gamma_{t+1})|\mathcal{F}_t] \leq (1 - \gamma_t)(1 - \gamma_t)\mathcal{J}(\gamma_t) + c_\perp:$$

Higher  $\gamma_t$  increases safety but slows convergence.

**Theorem 2** (Stability under verifier imperfection). *(Full proof in Appendix C.) Assume A1–A2 and Assumption 4. Let the verifier introduce bias  $\gamma_t$  with  $|\gamma_t| \leq \gamma_v$ , entering as  $c_t = \gamma_t \gamma_v$ . If  $\gamma_t \gamma_t = \infty$ ,  $\gamma_t \gamma_t^2 < \infty$ , and Lemma 1 conditions hold, then*

$$\limsup_{t \rightarrow \infty} \mathcal{J}(\gamma_t) \leq \mathcal{J}^* + C \gamma_v \quad \text{a.s.}$$

for finite  $C$  depending on  $(M; L_\oplus; \gamma)$  and the stepsize schedule.

**Corollary 4** (Bounded-set Convergence Radius). *Under Assumption 4 (mixing  $(T; \gamma)$ ) and verifier bias  $\gamma_v$ ,*

$$\limsup_{t \rightarrow \infty} \mathcal{J}(\gamma_t) \leq C_1 \gamma_v + C_2(1 - \gamma)^T;$$

for finite  $C_1, C_2$  depending on  $(M; L_\oplus; \gamma)$  and  $(T; \gamma)$ .

Table 2: Assumptions summary used in convergence analysis.

A1: Adaptivity & bounded updates	Assumption 1: $\gamma_t, e_t; \mathcal{V}(e_t)$ are $\mathcal{F}_t$ -adapted; $\ \Phi(\cdot)\  \leq M$ .
A2: Verification-monotone descent	Assumption 2: expected descent inequality with summable residual.
A2 <sup>0</sup> : Markovian noise & mixing	Assumption 4: uniform geometric ergodicity, Lipschitz in $\gamma$ ; bounded increments $B_t$ with $\mathbb{E}[B_t^2] < \infty$ .
A3: Local linearization & contraction	Assumption 3: local Gâteaux derivative, Lipschitz, and contraction of $T$ .
Stepsizes	$\gamma_t \gamma_t = \infty$ , $\gamma_t \gamma_t^2 < \infty$ ; Algorithm ?? uses $\gamma_f \equiv \gamma_t$ or a schedule.

## 7 Epistemic Scaling Laws

**Evaluation plan.** The empirical claims will be tested according to a preregistered protocol (Appendix A), specifying hypotheses, logging of  $\{r_t; u_t; \mathcal{I}_t\}$ , and robust regression analysis.

Framework	Noise	Dependence	Conclusion	Where used
Classical RS	$P[c_t < \infty]$	i.i.d./MD adapted	$X_t \rightarrow X$ a.s.	Baseline for Thm ??
Extended RS	$c_t^2 < \infty$ ; bounded increments	Markov, mixing	$X_t \rightarrow$ bounded set a.s.	Cor. 1
RFL (this work)	$c_t = \langle \cdot \rangle_v$	Policy-Markov	$\limsup \mathcal{J} \leq \mathcal{J}_v + C_v$	Thm 2

Table 3: Classical vs. extended Robbins–Siegmund vs. our RFL instantiation.

**Scaling law.** Performance scales as  $\Delta H \propto N_v^{-\beta}$  with  $\beta \in (0; 1]$ . Empirically  $\beta \approx 1/2$  is consistent with diffusion-like uncertainty decay.

**Interpretability and alignment perspective.** Alignment and verifiability themselves follow scaling-law behavior; RFL provides a formal alternative grounded in verified events.

**Empirical Outlook.** Our training pipeline uses a *proof\_sampler* process to generate events  $e_t$  under the current policy, producing a policy-dependent Markov stream. Assumption A2' connects this stream to theory: the sampler’s mixing and the verifier’s acceptance yield ( $\mathcal{T}_v$ ) and  $\langle \cdot \rangle_v$ . We will report  $(\langle \cdot \rangle, \mathcal{T})$  proxies (autocorrelation decay) and empirical  $\langle \cdot \rangle_v$  per run.

## 8 Emergent Directions

**Reflexive Formal Perception (future work).** Agents verify what was *seen* before reasoning; perceptual disagreements become ledger objects.

**Ledger-Driven Theory Genesis (future work).** Meta-agents mine sealed proofs to propose schemata under ledger governance.

**Instrumentation Hooks (Phase I).** Log perceptual disagreements and lemma-reuse frequencies; mine traces later.

## 9 Philosophical and Practical Boundaries: The Architect and the Healer

**Framing.** MathLedger embodies the *Architect*’s aspiration: intelligence grounded in provable truth. Its challenge is the *Healer*’s domain: extending verifiability into imperfect reality.

**Open problems (Phase III).**

1. **Scope of verification.**
2. **Verifier bottleneck.**
3. **Abstention vs. usefulness.**
4. **Semantic gap.**
5. **Computational cost.**

**Research tracks (not blockers).**

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Scalable verification	Probabilistic checking; batching; ZK sealing.
Verified oracle stack	Verified kernels/compilers; attested builds.
Controlled abstention	Abstention budgets; explore-on-fail policies.
Semantic grounding bridge	Typed front-ends; certified parsers/tokenizers.
Compute-efficient guarantees	Proof caching/reuse; parallel tree-hash; ZK compression.

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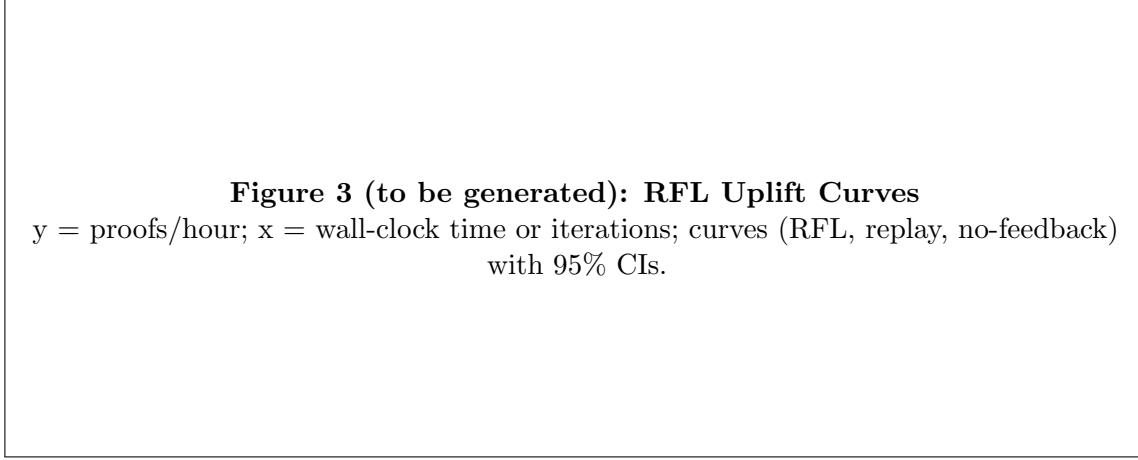


Figure 2: Pre-registered uplift curves comparing RFL vs. baselines.

## Appendix A — Prior-Art Matrix and Red-Team Notes

This appendix summarizes adjacent systems and red-team considerations.

## Appendix B — Evidence Manifest (Sealing Metadata)

We provide artifact paths and commits for reproducibility. We release `roots.json` and per-step Merkle inclusion proofs; a reproducibility script re-canonicalizes artifacts and re-derives  $(r_t; u_t; \mathcal{I}_t)$  byte-for-byte.<sup>1</sup>

## A Preregistration Protocol for Empirical Evaluation

### A.1 Hypotheses

H1:  $\log |\Delta H| = -\log N_v + c$  with  $c > 0$ . H2:  $\limsup_t \mathcal{J}(\cdot_t)$  increases with  $\cdot_v$ .

### A.2 Tasks and Datasets

Lean4 theorem proving (miniF2F/synthetic), transformer policy  $\cdot_t$ , Lean4 kernel as  $\mathcal{V}$ ; simulate  $\cdot_v$  by flipping outcomes.

### A.3 Proof Sampler and Logging

Log JSONL entries with step, policy/task ids,  $v_t$ , attestations  $(r_t; u_t; \mathcal{I}_t)$ , metrics (autocorr proxy,  $\cdot_v, H_t$ ).

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<sup>1</sup>KangarooTwelve (K12) is a tree-hash mode; observed Merkle construction speedups are empirical (2–3×) on AVX2/AVX-512 vs. SHA-256 in our pipeline.

**Table 1 (to be generated): Scaling Law Fit — Estimates**  
 Columns: run id,  $N_v$ ,  $\Delta H$ , fit  $\hat{\beta}$ , SE,  $R^2$ .

Table 4: Empirical fit of  $\Delta H \propto N_v^{-\beta}$ .

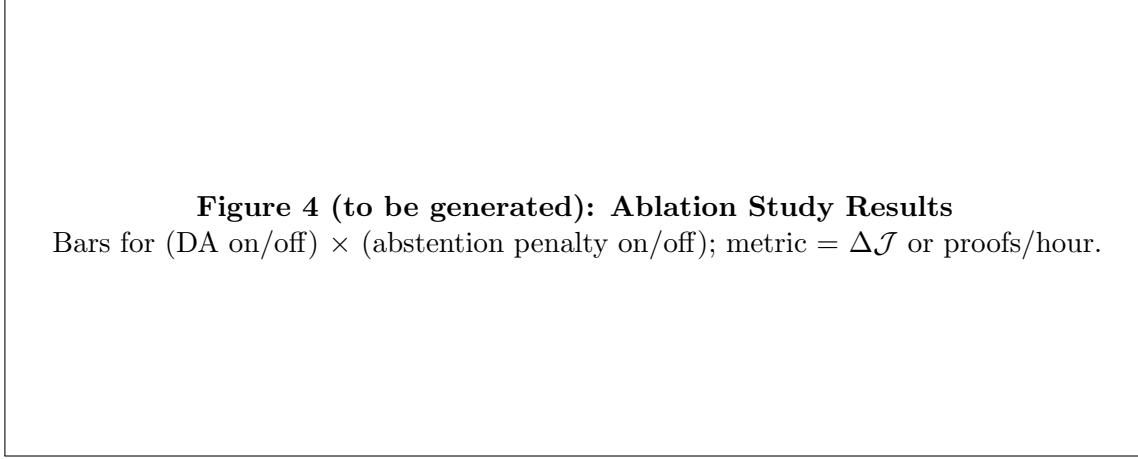


Figure 3: Ablations for DA and abstention penalty.

#### A.4 Power

See Table 5.

Table 5: Power analysis for detecting  $\beta$ .

Detectable	Required $N_v$
0.50	$\sim 1,000$
0.25	$\sim 4,000$
0.10	$\sim 25,000$

#### A.5 Analysis Plan

Huber regression for H1; Spearman correlation for H2. No manual point deletion; only run-level exclusions (hardware failure, zero-verify cold starts).

### Appendix C — Proofs of Main Results

**Assumptions recap.** All proofs invoke A1, A2, and, where stated, A2'; stepsizes satisfy  $\sum_t \frac{\sigma_t^2}{t} < \infty$ ,  $\sum_t \frac{\sigma_t^2}{t} \rightarrow \infty$ .

*Proof of Theorem ??.* Let  $X_t \bar{\rightarrow} \mathcal{J}(e_t)$ . By A2,  $\mathbb{E}[X_{t+1} | \mathcal{F}_t] \leq X_t - \frac{1}{f} \mathbf{1}\{\mathcal{V}(e_t) = 1\} + \frac{1}{t}$ . Boundedness ( $X_t \in [0; 1]$ ) and  $\sum_t \mathbb{E}[|e_t|] < \infty$  yield the claim via Robbins–Siegmund.  $\square$

*Proof of Theorem 2.* Let  $X_t = \mathcal{J}(e_t) - \mathcal{J}^*$ . With  $|e_t| \leq v$ ,

$$\mathbb{E}[X_{t+1} | \mathcal{F}_t] \leq X_t - \frac{1}{t} \Delta_t + \frac{1}{t} C_1 v + \frac{1}{t}:$$

Set  $Y_t = \sum_{\tau=0}^t C_1 \eta_\tau + \epsilon_t$ . Since  $\sum_{\tau=0}^t (\sum_{\tau=0}^\tau C_1 \eta_\tau)^2 < \infty$  and  $(X_{t+1} - X_t)_+ \leq B_t$  with  $\sum_{\tau=0}^t B_t^2 < \infty$  (A0 and bounded  $\Phi$ ), Lemma 1 applies, so  $X_t$  converges a.s. to a bounded set. Telescoping gives  $\limsup_t \mathcal{J}(\cdot_t) \leq \mathcal{J}^* + C \eta_v$  with  $C$  depending on  $(M; L_\oplus; \cdot)$  and stepsizes.  $\square$

## References