

MathLedger: Reflexive Formal Learning and the Chain of Verifiable Cognition

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Abstract

Contemporary AI systems achieve extraordinary performance yet remain largely *opaque* and *non-verifiable*, creating a crisis of trust for safety-critical deployment and governance. We introduce *Reflexive Formal Learning (RFL)*, a symbolic and cryptographically verifiable learning paradigm in which every update is justified by a formally verified proof or abstention event recorded in an immutable ledger. The resulting “chain of verifiable cognition” constitutes a closed epistemic loop bridging logic, cryptography, and learning dynamics. RFL is shown to be a discrete, non-differentiable analogue of stochastic gradient descent whose learning signal arises from verified truth rather than statistical loss, enabling provable convergence of reasoning under a verifiable ledger substrate.

Figure 1 (to be generated): Chain of Verifiable Cognition

User input \rightarrow PoA \rightarrow Ledger $r_t \rightarrow$ UI $u_t \rightarrow$ RFL \rightarrow update.

Figure 1: End-to-end epistemic loop (objects and arrows as shown; all edges attested).

1 Introduction

Mathematical reasoning systems increasingly require not only correctness but verifiable *cognitive integrity*—a guarantee that each inference can be cryptographically traced to a validated source. *MathLedger* unifies formal verification, machine learning, and cryptographic attestation into a single epistemic pipeline:

User-Verified Input \rightarrow Proof-or-Abstain Reasoning \rightarrow Ledger Attestation \rightarrow UI Attestation \rightarrow Reflexive Feedback

This “Chain of Verifiable Cognition” enables an AI system to learn exclusively from verified reasoning outcomes. The remainder of this paper formalizes *Reflexive Formal Learning (RFL)* as a symbolic analogue of gradient descent operating on verified events rather than numerical errors.

Table 1: Summary of Core Formal Constructs

Symbol	Meaning
$\pi_t \in \Pi$	Symbolic reasoning policy at time t (metric space $(\Pi; \ \cdot\)$)
P_π	Policy-induced event measure over event space \mathcal{E}
e_t	Reasoning event (proof or abstention) under P_π
$\mathcal{V}(e_t)$	Verification outcome $\in \{1; 0; \perp\}$
$\mathcal{V}(e)$	Numeric surrogate $\mathbf{1}\{\mathcal{V}(e) \neq 1\} \in \{0; 1\}$
\mathcal{L}	Ledger of proofs/abstentions (predictable process)
Φ	Feedback map $\Phi : \{1; 0; \perp\} \times \Pi \rightarrow \Delta\Pi$
\mathcal{U}	Reflexive update $\Pi \times \{1; 0; \perp\} \times \Sigma \rightarrow \Pi$
Σ	Auxiliary signal space (prompts, contexts)
$\Delta\Pi$	Space of composable symbolic policy deltas
\oplus	Algebraic composition on Π applying $\Delta \in \Delta\Pi$ to
L	Lipschitz factor of \oplus s.t. $\ \oplus \Delta - \ \leq L \ \Delta\ _\Delta$
$\ \cdot\ _\Delta$	Norm on $\Delta\Pi$
$(T; \gamma)$	Mixing horizon and geometric rate (Assumption A2 ⁰)
$r_t; \mathcal{U}_t$	Reasoning/UI Merkle roots (dual attestation)
\mathcal{I}_t	Dual-attestation binder token (keyless or HMAC)
H_t	Epistemic entropy at step t
$\mathcal{J}(\cdot)$	Epistemic risk functional (Def. 1)
$\mathcal{L}(\cdot)$	Classical differentiable loss (SGD)
	Epistemic scaling exponent

2 Related Work and Theoretical Context

Stochastic approximation and recent convergence extensions. Classical almost-supermartingale convergence (Robbins–Siegmund) guarantees pointwise convergence under a summable residual term. Recent work (e.g., Liu, Xie, & Zhang 2025a) relaxes this to *square-summable* disturbances, proving convergence *to a bounded set*. Our stability results instantiate these extensions: when a persistent verifier bias acts as a non-summable zero-order term, we obtain bounded-set convergence with an $O(\gamma_v)$ radius. Our Markovian noise assumptions (Assumption A2⁰ in §6) align with modern SA analyses that handle time-inhomogeneous Markov noise with mixing and Lipschitz drift. *Cor. 1 instantiates the square-summable residual regime of Liu–Xie–Zhang (2025), yielding bounded-set convergence under Markov noise.*

Positioning. RFL offers a distinct paradigm for verifiable learning. While recent work on *proof-guided language models* uses verifier feedback as a reward signal to guide policy search in a reinforcement learning loop, RFL uses the verifier’s binary decision as the *learning signal itself*, replacing a gradient with a symbolic update. Similarly, where *weak-to-strong generalization* frameworks train a strong model on labels generated by a weaker one, RFL’s policy updates are a direct, deterministic function of verified truth events, not supervised fine-tuning. The ledger’s immutable, attested history of reasoning also connects RFL to work on *proof-carrying data* and ledger-based AI audits, but with a closed-loop dynamic where the ledger actively drives learning.

3 Reflexive Formal Learning (RFL): Conceptual Definition

Let Π be a complete metric policy space with norm $\|\cdot\|$ and P_π the event measure induced by π .

Update algebra. We write $\Delta\Pi$ for the space of composable symbolic policy deltas and \oplus for their algebraic composition in Π ; thus $\pi \oplus \Delta$ denotes applying $\Delta \in \Delta\Pi$ to policy $\pi \in \Pi$. We equip $\Delta\Pi$ with a norm $\|\cdot\|_\Delta$ and assume *norm compatibility*: there exists $L_\oplus \geq 1$ with

$$\|\pi \oplus \Delta - \pi\| \leq L_\oplus \|\Delta\|_\Delta \quad \text{for all } \pi \in \Pi; \Delta \in \Delta\Pi:$$

Definition 1 (Epistemic risk). *Given P_π and $\mathfrak{V}(e) = \mathbf{1}\{\mathcal{V}(e) \neq 1\}$,*

$$\mathcal{J}(\pi) = \mathbb{E}_{e \sim P_\pi}[\mathfrak{V}(e)] = \Pr_{e \sim P_\pi}[\mathcal{V}(e) \neq 1].$$

Range. Since $\mathfrak{V}(e) \in \{0, 1\}$, we have $0 \leq \mathcal{J}(\pi) \leq 1$ for all π .

4 The Reflexive Formal Learning (RFL) Update Law

Update operator and abstention damping. At each step t ,

$$\pi_{t+1} = \pi_t \oplus_f \Phi(\mathcal{V}(e_t); \pi_t); \quad \Phi : \{1, 0, \perp\} \times \Pi \rightarrow \Delta\Pi: \quad (1)$$

By norm compatibility, committing $\pi_{t+1} = \pi_t \oplus_f \Phi(\cdot)$ implies $\|\pi_{t+1} - \pi_t\| \leq_f L_\oplus \|\Phi(\mathcal{V}(e_t); \pi_t)\|_\Delta$.

Toy example: one-step RFL update with pseudo-Lean

Goal and tactic (pseudo-Lean).

```
example (h1 : P → Q) (h2 : P) : Q := by
  apply h1
  exact h2
```

Let the current goal be $g_t : Q$ under context $\text{ctx}_t = \{h_1 : P \rightarrow Q; h_2 : P\}$. The agent proposes tactic $s_t = \text{apply } h1$, producing subgoal $g'_t : P$, then proposes $s'_t = \text{exact } h2$.

Event and verification. Define the event

$$e_t = (g_t; \text{ctx}_t; s_t; \text{apply}; s'_t; \text{exact}) \quad \text{and} \quad v_t = \mathcal{V}(e_t) \in \{1; 0; \perp\}:$$

Pattern features. Let $\text{pat}(\text{ctx}_t; s_t)$ denote a differentiable feature map.

Policy parameterization. Suppose θ_t induces a score via parameters θ_t :

$$\text{score}_t = \langle \theta_t; \text{pat}(\text{ctx}_t; s_t) \rangle \Rightarrow \rho_t(\text{pattern}) = \text{score}_t:$$

Symbolic deltas. Instantiate

$$\Delta^+(\theta_t; e_t) \equiv \text{inc}_{\eta} \text{pat}(\text{ctx}_t; s_t) ; \quad \Delta^-(\theta_t; e_t) \equiv -\text{dec}_{\eta} \text{pat}(\text{ctx}_t; s_t) ;$$

with $\|\Delta^\pm\|_\Delta \leq M$.

Cases (Proof-or-Abstain).

- $v_t=1$: $\theta_{t+1} = \theta_t \oplus_f \Delta^+(\theta_t; e_t)$; \mathcal{J} decreases in expectation.
- $v_t=\perp$: no-op update; abstention logged.
- $v_t=0$: $\theta_{t+1} = \theta_t \oplus_f \Delta^-(\theta_t; e_t)$; demotes the pattern.

Algorithm 1 RFL \oplus MCGS Planner (fail-closed, dual-attested)

Require: t , f , update Φ , verifier (REPL), canonicalizers C_R, C_U , ledger \mathcal{L}

```
1: Initialize frontier at root Lean state;  $E \leftarrow []$  . list of per-event binders
2: while frontier nonempty do
3:   Expand node using policy  $t$  to yield candidate events  $e_t$ 
4:    $(v_t; \text{trace}; \text{build}) \leftarrow \text{REPL}.\text{check}(e_t)$ 
5:   if  $v_t \neq 1$  then
6:      $\mathcal{L}.\text{abstain}(e_t)$ ; continue
7:   end if
8:    $P_t \leftarrow C_R(e_t; \text{trace}; \text{build})$ ;  $D_t \leftarrow C_U(\text{UI snapshot})$ 
9:    $r_t \leftarrow \text{Hash}(R : || P_t)$ ;  $u_t \leftarrow \text{Hash}(U : || D_t)$ ;  $l_t \leftarrow \text{Hash}(\text{BIND} || r_t || u_t)$ 
10:  if  $\neg \mathcal{L}.\text{verify}(P_t; D_t; l_t)$  then
11:     $\mathcal{L}.\text{abstain}(e_t)$ ; continue
12:  end if
13:   $\Delta_t \leftarrow \Phi(v_t; t)$ ;  $t_{+1} \leftarrow t \oplus f \Delta_t$ 
14:   $\mathcal{L}.\text{commit}(e_t; r_t; u_t; l_t; \text{build})$ ;  $E.\text{append}(l_t)$ 
15:  Push children of  $e_t$  to frontier with priority from  $t_{+1}$ 
16: end while
17:  $\text{epoch\_root} \leftarrow \text{Merkle}(E)$ ;  $\mathcal{L}.\text{finalize\_epoch}(\text{epoch\_root})$ 
```

5 Dual Attestation and Security Model

The integrity of the RFL loop depends on cryptographically binding reasoning events to their presentation. We formalize this via a dual-attestation scheme and specify the security guarantees under a formal adversary model.

Dual-Root Ledger. At each step t , the system commits to two Merkle roots:

- **Reasoning Root (r_t):** Root over the canonicalized sequence of formal reasoning steps (proof tactics, intermediate goals) composing e_t .
- **UI Root (u_t):** Root over the canonicalized representation of the UI state (DOM/JSON, PNG, HAR log) that displays the outcome of e_t .

These roots are bound by a cryptographic token $\mathcal{I}_t = \text{Hash}(\text{"BIND: " } || r_t || u_t)$ with prefix-free domain separation. The tuple $(r_t; u_t; \mathcal{I}_t)$ is recorded on the ledger.

Domain tags and REPL provenance. We extend domain separation with tags RPL: (Lean REPL provenance; toolchain/build IDs) and G: (geometry engine artifacts). These tags are included by C_R prior to Merkle, binding $(r_t; u_t; \mathcal{I}_t)$ to verifier versions and domain-specific pipelines.

Adversary Model and Guarantees. We consider a PPT adversary acting as a malicious prover, verifier, or network observer. The ledger offers:

- **Collision Resistance:** $\text{Hash}(\cdot)$ is collision-resistant, preventing distinct artifacts from sharing a root.
- **Non-Malleability:** Modifying proof/UI artifacts invalidates Merkle roots and \mathcal{I}_t .
- **Replay Resistance:** Updates are indexed by t ; replay $(r_k; u_k; \mathcal{I}_k)$ at $t > k$ is rejected; timestamps strengthen this.

Cryptographic Hardening.

- **Canonicalization:** All structured data MUST be canonicalized before hashing. We mandate RFC 8785 JCS for JSON; deterministic PNG and equivalent for other types.
- **Domain Separation:** All hash inputs use prefix-free tags: e.g., $\text{Hash}(\text{"LEAF: "} \parallel \cdot)$, $\text{Hash}(\text{"NODE: "} \parallel h_L \parallel h_R)$; distinct tags for reasoning/UI/binder.
- **Constant-Time Ops:** Cryptographic comparisons MUST be constant-time to avoid timing side channels.
- **Version Pinning:** Record versions/hashes of \mathcal{V} , hash function, and canonicalizers on-ledger.

Zero-Knowledge Extensions. For privacy, proofs may be replaced by ZK certificates (e.g., PLONK for fast verification and small proofs; STARKs for transparency and post-quantum resistance, at larger sizes).

6 Convergence and Stability of Reflexive Formal Learning

Stepsizes. We identify $\eta_t \equiv \eta$ in the constant-stepsizes case; otherwise η_t denotes a decaying schedule used in the SA analysis.

Lemma 1 (Extended Robbins–Siegmund). *Let $\{Z_t\}$, $\{X_t\}$, $\{Y_t\}$, and $\{W_t\}$ be non-negative, \mathcal{F}_t -adapted random sequences. Suppose $\mathbb{E}[Z_{t+1} \mid \mathcal{F}_t] \leq (1 + X_t)Z_t + Y_t - W_t$ a.s. If $\sum_t X_t < \infty$ a.s. and $\sum_t Y_t < \infty$ a.s., then Z_t converges a.s. to a finite random variable and $\sum_t W_t < \infty$ a.s. Furthermore, if the summable condition on $\{Y_t\}$ is relaxed to square-summable ($\sum_t Y_t^2 < \infty$ a.s.) and the increments are bounded $(Z_{t+1} - Z_t)_+ \leq B_t$ with $\sum_t B_t^2 < \infty$ a.s., then $\{Z_t\}$ converges a.s. to a bounded set.*

Assumption 1 (Adaptivity and bounded updates (A1)). *Let $\{\mathcal{F}_t\}$ denote the filtration generated by $(e_s; \mathcal{V}(e_s))_{s \leq t}$. The iterates θ_t , reasoning events e_t , and verification outcomes $\mathcal{V}(e_t)$ are \mathcal{F}_t -adapted. Moreover, the update increments are uniformly bounded: there exists $M < \infty$ such that $\|\Phi(\mathcal{V}(e_t); \theta_t)\|_\Delta \leq M$ almost surely for all t .*

Assumption 2 (Verification-monotone descent (A2)). *There exist constants $\beta > 0$ and an \mathcal{F}_t -adapted error term $\epsilon_t \geq 0$ with $\sum_t \mathbb{E}[\epsilon_t] < \infty$ such that the epistemic risk satisfies*

$$\mathbb{E}[\mathcal{J}(\theta_{t+1}) - \mathcal{J}(\theta_t) \mid \mathcal{F}_t] \leq -\beta \Pr(\mathcal{V}(e_t) = 1 \mid \mathcal{F}_t) + \epsilon_t$$

for all t .

Assumption 3 (Local linearization and contraction (A3)). *There exists a neighborhood \mathcal{N} of θ^* in which the averaged update map $\mathcal{T}(\cdot) := \mathbb{E}[\theta \oplus \eta \Phi(\mathcal{V}(e); \cdot)]$ is Gâteaux differentiable and Lipschitz. The Jacobian $D\mathcal{T}(\theta^*)$ has spectral radius < 1 , yielding a contraction on \mathcal{N} in the ambient norm.*

Assumption 4 (Markovian Noise and Mixing (A2')). *The event stream $\{e_t\}$ is generated via a policy-dependent Markov process $\{X_t\}$ on a state space \mathcal{X} . For each fixed policy π , the process has a unique stationary distribution π . The process is uniformly geometrically ergodic: uniformly over π , there exist $T \in \mathbb{N}$ and $\gamma \in (0, 1)$ such that $\|\mathbb{P}_\pi(X_T \in \cdot \mid X_0 = x) - \pi(\cdot)\|_{\text{TV}} \leq \gamma$ for all x . The one-step expected cost change and transition probabilities are Lipschitz in π . In addition, there exists an adapted sequence $B_t \geq 0$ with $(\mathcal{J}(\theta_{t+1}) - \mathcal{J}(\theta_t))_+ \leq B_t$ a.s. and $\sum_t B_t^2 < \infty$ a.s. (implied by A0 and bounded Φ under standard stepsizes).*

Remark 1. Uniform geometric ergodicity is standard in SA; weaker mixing (Doebelin minorization or spectral-gap conditions) can suffice in place of uniformity.

Theorem 1 (Almost-Sure Convergence with Lyapunov Potential). *If Assumptions 1–2 hold and $\mathbb{E}[\Delta\mathcal{J}_t|\mathcal{F}_t] \leq -c\|\Phi\|^2 + \epsilon_t$ with $\sum_t \mathbb{E}[\epsilon_t] < \infty$, then $\mathcal{J}(\epsilon_t) \rightarrow 0$ a.s. and $\epsilon_t \rightarrow \epsilon^*$ where $\mathcal{V}(\epsilon^*) = 1$.*

Corollary 1 (Convergence to a bounded set under square-summable noise). *Let $X_t = \mathcal{J}(\epsilon_t)$. Assume $\sum_t \epsilon_t = \infty$, $\sum_t \epsilon_t^2 < \infty$. If the RFL dynamics satisfy Lemma 1 with square-summable disturbance and bounded increments, then $\{X_t\}$ converges a.s. to a bounded set; pointwise convergence is recovered when the disturbance term is summable.*

Corollary 2 (Linear Convergence of Epistemic Risk). *Under Theorem ??, if $\Pr[\mathcal{V}(e_t) = 1] \geq c > 0$ for all non-optimal policies, then $\mathbb{E}[\mathcal{J}(\epsilon_t)]$ converges to its limit at a linear rate.*

Corollary 3 (Local Stability of the Optimal Policy). *Under Assumption 3, the fixed point ϵ^* is locally asymptotically stable: any ϵ_t within the contraction basin converges to ϵ^* under RFL dynamics.*

Proposition 1 (Abstention as damping). *If $\Pr[\mathcal{V}(e_t) = \perp] = \epsilon_t$ and abstention cost $\text{cost}(\perp) = c_\perp \geq 0$ (captures operational abstention cost), then*

$$\mathbb{E}[\mathcal{J}(\epsilon_{t+1})|\mathcal{F}_t] \leq (1 - \epsilon_t)(1 - \epsilon_t)\mathcal{J}(\epsilon_t) + c_\perp.$$

Higher ϵ_t increases safety but slows convergence.

Theorem 2 (Stability under verifier imperfection). (Full proof in Appendix C.) *Assume A1–A2 and Assumption 4. Let the verifier introduce bias ϵ_t with $|\epsilon_t| \leq \epsilon_v$, entering as $c_t = \epsilon_t \epsilon_v$. If $\sum_t \epsilon_t = \infty$, $\sum_t \epsilon_t^2 < \infty$, and Lemma 1 conditions hold, then*

$$\limsup_{t \rightarrow \infty} \mathcal{J}(\epsilon_t) \leq \mathcal{J}^* + C \epsilon_v \quad \text{a.s.}$$

for finite C depending on $(M; L_\oplus; \epsilon)$ and the stepsize schedule.

Corollary 4 (Bounded-set Convergence Radius). *Under Assumption 4 (mixing $(T; \epsilon)$) and verifier bias ϵ_v ,*

$$\limsup_{t \rightarrow \infty} \mathcal{J}(\epsilon_t) \leq C_1 \epsilon_v + C_2(1 - \epsilon)^T;$$

for finite C_1, C_2 depending on $(M; L_\oplus; \epsilon)$ and $(T; \epsilon)$.

Table 2: Assumptions summary used in convergence analysis.

A1: Adaptivity & bounded updates	Assumption 1: $\epsilon_t; e_t; \mathcal{V}(e_t)$ are \mathcal{F}_t -adapted; $\ \Phi(\cdot)\ \leq M$.
A2: Verification-monotone descent	Assumption 2: expected descent inequality with summable residual.
A2 ⁰ : Markovian noise & mixing	Assumption 4: uniform geometric ergodicity, Lipschitz in ϵ ; bounded increments B_t with $\sum_t B_t^2 < \infty$.
A3: Local linearization & contraction	Assumption 3: local Gâteaux derivative, Lipschitz, and contraction of \mathcal{T} .
Stepsizes	$\sum_t \epsilon_t = \infty$, $\sum_t \epsilon_t^2 < \infty$; Algorithm ?? uses $\epsilon_f \equiv \epsilon_t$ or a schedule.

7 Epistemic Scaling Laws

Evaluation plan. The empirical claims will be tested according to a preregistered protocol (Appendix A), specifying hypotheses, logging of $\{r_t; u_t; \mathcal{I}_t\}$, and robust regression analysis.

Framework	Noise	Dependence	Conclusion	Where used
Classical RS	$\mathbb{P} \quad c_t < \infty$	i.i.d./MD adapted	$X_t \rightarrow X$ a.s.	Baseline for Thm ??
Extended RS	$\mathbb{P} \quad c_t^2 < \infty$; bounded increments	Markov, mixing	$X_t \rightarrow$ bounded set a.s.	Cor. 1
RFL (this work)	$\mathbb{P} \quad c_t = c_v$	Policy–Markov	$\limsup \mathcal{J} \leq \mathcal{J} + C_v$	Thm 2

Table 3: Classical vs. extended Robbins–Siegmund vs. our RFL instantiation.

Scaling law. Performance scales as $\Delta H \propto N_v^{-\beta}$ with $\beta \in (0; 1]$. Empirically $\beta \approx 1/2$ is consistent with diffusion-like uncertainty decay.

Interpretability and alignment perspective. Alignment and verifiability themselves follow scaling-law behavior; RFL provides a formal alternative grounded in verified events.

Empirical Outlook. Our training pipeline uses a *proof_sampler* process to generate events e_t under the current policy, producing a policy-dependent Markov stream. Assumption A2' connects this stream to theory: the sampler's mixing and the verifier's acceptance yield $(T; \epsilon)$ and c_v . We will report $(\epsilon; T)$ proxies (autocorrelation decay) and empirical c_v per run.

8 Emergent Directions

Reflexive Formal Perception (future work). Agents verify what was *seen* before reasoning; perceptual disagreements become ledger objects.

Ledger-Driven Theory Genesis (future work). Meta-agents mine sealed proofs to propose schemata under ledger governance.

Instrumentation Hooks (Phase I). Log perceptual disagreements and lemma-reuse frequencies; mine traces later.

9 Philosophical and Practical Boundaries: The Architect and the Healer

Framing. MathLedger embodies the *Architect*'s aspiration: intelligence grounded in provable truth. Its challenge is the *Healer*'s domain: extending verifiability into imperfect reality.

Open problems (Phase III).

1. **Scope of verification.**
2. **Verifier bottleneck.**
3. **Abstention vs. usefulness.**
4. **Semantic gap.**
5. **Computational cost.**

Research tracks (not blockers).

Scalable verification	Probabilistic checking; batching; ZK sealing.
Verified oracle stack	Verified kernels/compiler; attested builds.
Controlled abstention	Abstention budgets; explore-on-fail policies.
Semantic grounding bridge	Typed front-ends; certified parsers/tokenizers.
Compute-efficient guarantees	Proof caching/reuse; parallel tree-hash; ZK compression.

Figure 3 (to be generated): RFL Uplift Curves
y = proofs/hour; x = wall-clock time or iterations; curves (RFL, replay, no-feedback)
with 95% CIs.

Figure 2: Pre-registered uplift curves comparing RFL vs. baselines.

Appendix A — Prior-Art Matrix and Red-Team Notes

This appendix summarizes adjacent systems and red-team considerations.

Appendix B — Evidence Manifest (Sealing Metadata)

We provide artifact paths and commits for reproducibility. We release `roots.json` and per-step Merkle inclusion proofs; a reproducibility script re-canonicalizes artifacts and re-derives $(r_t; u_t; \mathcal{I}_t)$ byte-for-byte.¹

A Preregistration Protocol for Empirical Evaluation

A.1 Hypotheses

H1: $\log |\Delta H| = -\log N_v + c$ with $c > 0$. H2: $\limsup_t \mathcal{J}(t)$ increases with N_v .

A.2 Tasks and Datasets

Lean4 theorem proving (miniF2F/synthetic), transformer policy π_t , Lean4 kernel as \mathcal{V} ; simulate N_v by flipping outcomes.

A.3 Proof Sampler and Logging

Log JSONL entries with step, policy/task ids, v_t , attestations $(r_t; u_t; \mathcal{I}_t)$, metrics (autocorr proxy, v_t, H_t).

¹KangarooTwelve (K12) is a tree-hash mode; observed Merkle construction speedups are empirical (2–3×) on AVX2/AVX-512 vs. SHA-256 in our pipeline.

Table 1 (to be generated): Scaling Law Fit — Estimates

Columns: run id, N_v , ΔH , fit $\hat{\cdot}$, SE, R^2 .

Table 4: Empirical fit of $\Delta H \propto N_v^{-\beta}$.

Figure 4 (to be generated): Ablation Study Results

Bars for (DA on/off) \times (abstention penalty on/off); metric = $\Delta \mathcal{J}$ or proofs/hour.

Figure 3: Ablations for DA and abstention penalty.

A.4 Power

See Table 5.

Table 5: Power analysis for detecting \cdot .

Detectable	Required N_v
0.50	$\sim 1,000$
0.25	$\sim 4,000$
0.10	$\sim 25,000$

A.5 Analysis Plan

Huber regression for H1; Spearman correlation for H2. No manual point deletion; only run-level exclusions (hardware failure, zero-verify cold starts).

Appendix C — Proofs of Main Results

Assumptions recap. All proofs invoke A1, A2, and, where stated, A2'; stepsizes satisfy $\sum_{t=1}^{\infty} \eta_t = \infty$, $\sum_{t=1}^{\infty} \eta_t^2 < \infty$.

Proof of Theorem ??. Let $X_t \in \mathcal{J}(\cdot)$. By A2, $\mathbb{E}[X_{t+1} \mid \mathcal{F}_t] \leq X_t - \eta_t \mathbf{1}\{\mathcal{V}(e_t) = 1\} + \eta_t$. Boundedness ($X_t \in [0; 1]$) and $\sum_{t=1}^{\infty} \mathbb{E}[\eta_t] < \infty$ yield the claim via Robbins–Siegmund. \square

Proof of Theorem 2. Let $X_t = \mathcal{J}(\cdot) - \mathcal{J}^*$. With $|\cdot| \leq \eta_v$,

$$\mathbb{E}[X_{t+1} \mid \mathcal{F}_t] \leq X_t - \eta_t \Delta_t + \eta_t C_1 \eta_v + \eta_t.$$

Set $Y_t = \eta_t C_1 \|v\| + \eta_t$. Since $\sum_t (\eta_t C_1 \|v\|)^2 < \infty$ and $(X_{t+1} - X_t)_+ \leq B_t$ with $\sum_t B_t^2 < \infty$ (A0 and bounded Φ), Lemma 1 applies, so X_t converges a.s. to a bounded set. Telescoping gives $\limsup_t \mathcal{J}(\eta_t) \leq \mathcal{J}^* + C \|v\|$ with C depending on $(M; L_\oplus; \cdot)$ and stepsizes. \square

References