

MTH 3105: Discrete Mathematics Take Home Assignment I

1.

- (a) Define the sequence c_0, c_1, \dots by the equations $c_0 = 0$ and $c_n = c_{\lfloor n/2 \rfloor} + 3$ for all $n > 0$. Prove that $c_n \leq 2n$ for all $n \geq 3$. (5 Marks)

First of all, $c_3 = c_4 = c_5 = 6$, so the hypothesis is true for $n = 3, 4$ and 5 . Now assume the hypothesis is true for $n = 3, 4, \dots, k$, where $k \geq 5$. Consider $n = k + 1$:

$$\begin{aligned} c_{k+1} &= c_{\lfloor (k+1)/2 \rfloor} + 3 \\ &\leq 2 \lfloor (k+1)/2 \rfloor + 3 \quad \text{by assumption} \\ &\leq (k+1) + 3 \\ &\leq 2(k+1) \quad \text{since } k \geq 5 \end{aligned}$$

Hence by induction the hypothesis is true for all $n \geq 3$.

- (b) Prove that any positive integer N is divisible by 11 if and only if the difference between the sum of odd digits and the sum of even digits is divisible by 11. (5 Marks)

First, note that $10 \equiv -1 \pmod{11}$, and $100 \equiv 1 \pmod{11}$.

We observe that $10^{2k+1} \equiv -1 \pmod{11}$ and $10^{2k} \equiv 1 \pmod{11}$.

Let $N = \sum_{i=0}^k 10^i d_i$ so that k is the largest integer satisfying $10^k \leq N$.

So $N \equiv \sum_{i=0}^k 10^i d_i \pmod{11}$.

We set $d_{k+1} = 0$ if k is even.

It follows that we can write the previous modulo equation as

$$N \equiv \sum_{i=0}^{\lfloor \frac{k}{2} \rfloor} (d_{2i} - d_{2i+1}) \pmod{11}$$

The right-hand side shows the difference of sum of odd digits and even digits and the proof completes.

2.

- (a) Assume that $\forall x \exists y P(x, y)$ is false and the domain of them is nonempty. Which of the following must be false? (2 Marks)

- | | | |
|-------|-------------------------------|-------|
| (i) | $\forall x \forall y P(x, y)$ | False |
| (ii) | $\exists x \forall y P(x, y)$ | False |
| (iii) | $\exists x \exists y P(x, y)$ | True |

- (b) Let $P(x)$ denote the statement “ x is an accountant” and let $Q(x)$ denote the statement “ x owns a Porsche”. Write each statement below in first order logic.

- | | | | |
|-------|--|-------------------------------------|-------------|
| (i) | All accountants own Porsches. | $\forall x (P(x) \rightarrow Q(x))$ | (2 Marks) |
| (ii) | Some accountant owns a Porsche. | $\exists x (P(x) \wedge Q(x))$ | (2 Marks) |
| (iii) | All owners of Porsches are accountants. | $\forall x (Q(x) \rightarrow P(x))$ | (2 Marks) |
| (iv) | Someone who owns a Porsche is an accountant. | $\exists x (P(x) \wedge Q(x))$ | (2 Marks) |

3.

- (a) Prove that $\sqrt{6}$ is irrational. (4 Marks)

Suppose $\sqrt{6}$ is rational, so we can write $\sqrt{6} = \frac{m}{n}$ so that m and n do not have common factor other than

1. So, we get $6 = \frac{m^2}{n^2}$

$$6n^2 = m^2$$

Since left hand side is even, m must be even, put $n = 2k$. Then

$$6n^2 = (2k)^2$$

$$3n^2 = 2k^2$$

Now right-hand side is even, n must also be even. This contradicts the assumption that n and m do not have common factor other than 1. Hence $\sqrt{6}$ is irrational.

(b) A detective has interviewed four witnesses to a crime. From their stories, the detective has concluded that:

- (i) If the butler is telling the truth, then so is the cook.
- (ii) The cook and the gardener cannot both be telling the truth.
- (iii) The gardener and the handyman are not both lying.
- (iv) If the handyman is telling the truth then the cook is lying.

Who must be lying? There may be more than one liar. Show your steps. (6 Marks)

Let B , C , G and H be the proposition that the butler, cook, gardener and handyman is telling the truth respectively. Then from the problem description we get

$$i. B \rightarrow C$$

$$ii. \neg (C \wedge G)$$

$$iii. G \vee H$$

$$iv. H \rightarrow \neg C$$

Now suppose C is true, then by (iv) we get H is false, then by (iii) we get G is true. By (ii) we get C is false, which is a contradiction. As a result, C must be false. Furthermore, B must be false by (i). Here we cannot proceed any further, so we know **the cook and butler must be lying**.

4.

(a) Simplify the following statement:

(2 Marks)

$$\neg (\neg q \wedge \neg (\neg q \vee s)) \vee (q \wedge (r \rightarrow r))$$

$$\equiv \neg (\neg q \wedge \neg (\neg q \vee s)) \vee q$$

identity laws

$$\equiv (q \vee (\neg q \vee s)) \vee q$$

$$\equiv (q \vee \neg q) \vee s \vee q$$

$$\equiv \text{true}$$

(b) Determine whether or not the following arguments are valid:

$$(i) \frac{\neg p \rightarrow q, \neg q}{p}$$

(4 Marks)

$$(ii) \frac{\frac{p}{\neg p \rightarrow \neg q}}{p \rightarrow q}$$

(4 Marks)

p	q	$\neg p$
F	F	T
F	T	T
T	F	F
T	T	F

$\neg p \rightarrow q$	$\neg q$	p
F	T	F
T	F	F
T	T	T
T	F	T

Valid

$\neg q$	$\neg p \rightarrow \neg q$	$p \rightarrow q$
T	T	T
F	F	T
T	T	F
F	T	T

Invalid