Computer Network Laboratory

Assignment 7

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Course: CSN-361

Github Link: https://github.com/hemant84/CSN-361

Problem Statements:

Problem 1:

Transmit a binary message (from a sender to a receiver) using socket programming in C and report whether the received msg is correct or not; using the following error detection algorithms:

Algorithms used:

- socket_fd and sockaddr_in for socket creation
- bind function to bin the socket
- listen function to wait for the client to approach the server
- some basic data structures like int and arrays

Screenshot : Single Parity Check

ALGO: Number of ones are added and a bit with total%2 is added at the end

```
File Edit View Search Terminal Help
Choose which algorithm to check:

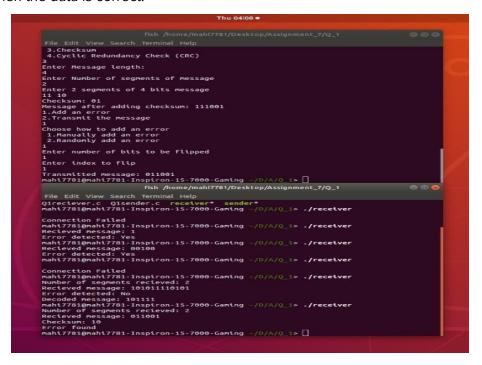
1. Single Parity Check
2. Iwo-dinensional Parity Check
3. Checksum
3. Checksum
4. Cyclic Redundancy Check (CRC)
1. Enter Message length:
4. Enter Message length:
4. Enter the message to send
Bessage after adding parity: 10100
1. Add an error
2. Transmit the message
1. Choose how to add an error
1. Hanualty add an error
1. Hanualty add an error
1. Hanualty add an error
1. Flamonly add an error
1. Fl
```

Two-dimensional Parity Check

ALGO: Single parity is calculated for each row and each column and an extra bit row and column are added respectively

Checksum

ALGO: Data is divided into segments and its checked sum is calculated using the wrapped sum concept. It is checked at the receiver's end and if the total comes to be zero then the data is correct.



Cyclic Redundancy Check (CRC)

ALGO: Based on binary-xor division CRC bits are calculated and are appended at the end and the receivers end its division is performed again. If the result comes to be zero, then the data is correct.

Problem 2:

Transmit a binary message (from a sender to a receiver) using socket programming in C. Using Hamming code detect and correct errors in the transmitted message, if any.

Algorithm used:

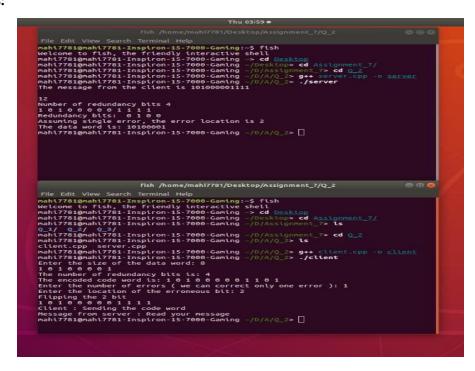
- 1. Write the bit positions starting from 1 in binary form (1, 10, 11, 100, etc).
- 2. All the bit positions that are a power of 2 are marked as parity bits (1, 2, 4, 8, etc).
- 3. All the other bit positions are marked as data bits.
- 4. Each data bit is included in a unique set of parity bits, as determined by its bit position in binary form.
 - a. Parity bit 1 covers all the bits positions whose binary representation includes a 1 in the least significant position (1, 3, 5, 7, 9, 11, etc).
 - b. Parity bit 2 covers all the bits positions whose binary representation includes a 1 in the second position from the least significant bit (2, 3, 6, 7, 10, 11, etc).

- c. Parity bit 4 covers all the bits positions whose binary representation includes a 1 in the third position from the least significant bit (4-7, 12-15, 20-23, etc).
- d. Parity bit 8 covers all the bits positions whose binary representation includes a 1 in the fourth position from the least significant bit bits (8–15, 24–31, 40–47, etc).
- e. In general, each parity bit covers all bits where the bitwise AND of the parity position and the bit position is non-zero.
- 5. Since we check for even parity set a parity bit to 1 if the total number of ones in the positions it checks is odd.
- 6. Set a parity bit to 0 if the total number of ones in the positions it checks is even.

Data Structures:

- bind function binds the socket to the address and port number specified in addr and socket creation as mentioned in the above question
- accept function to extract the first connection request on the queue of pending connections for the listening socket, sockfd, creates a new connected socket, and returns a new file descriptor referring to that socket.
- Socket connection: same as that of the server's socket creation.
- connect function to establish a connection.

Screenshots:



Problem 3:

Write a C++ program to compress a message (non-binary, can be anything like a text message or a code like hexadecimal, etc.) using the following data compression algorithm:

a) Huffman b) Shannon-Fano

Data Structures Used:

- Node- a struct tree node denotes a char and its properties
- comp-Comparison object to be used to order the heap
- encode-traverse the Huffman Tree and store Huffman Codes in a map.
- inbuilt sort function to sort according to probabilities
- decode-traverse the Huffman Tree and decode the encoded string
- Huffman Tree Builder
- Probability table storing probabilities
- ifstream and ofstream : Stream class to read from files
- fstream: Stream class to both read and write from/to files.

Screenshots:

Huffman Algorithm:

There are two major parts in Huffman Coding

1) Build a Huffman Tree from input characters. Steps to build Huffman Tree

Input is an array of unique characters along with their frequency of occurrences and output is Huffman Tree.

Create a leaf node for each unique character and build a min-heap of all leaf nodes (Min Heap is used as a priority queue. The value of frequency field is used to compare two nodes in min heap. Initially, the least frequent character is at root)

Extract two nodes with the minimum frequency from the min-heap.

Create a new internal node with a frequency equal to the sum of the two nodes frequencies.

Make the first extracted node as its left child and the other extracted node as its right child. Add this node to the min-heap.

Repeat steps#2 and #3 until the heap contains only one node. The remaining node is the root node and the tree is complete.

2) Traverse the Huffman Tree and assign codes to characters.

```
fish /home/mahi7781/Desktop/old1
File Edit View Search Terminal Help
Q3.cpp:379:10: warning: invalid suffix on literal; C++11 requires a space betwee
n literal and string macro [-Wliteral-suffix]
printf(" shannon -d encoded.txt"NL);
mahi7781@mahi7781-Inspiron-15-7000-Gaming -/D/old1> ./Q3 input.txt
Select Encoding
1. Huffman
2. Shannon-Fano
Huffman Codes are :
e:11111
i:11110
u:1110
:110
n:1011
m:0010
h:000
a:01
k:100
:0011
p:1010
```

Shannon-Fano Algorithm:

The steps of the algorithm are as follows:

- Create a list of probabilities or frequency counts for the given set of symbols so that the relative frequency of occurrence of each symbol is known.
- Sort the list of symbols in decreasing order of probability, the most probable ones to the left and least probable to the right.
- Split the list into two parts, with the total probability of both the parts being as close to each other as possible.
- Assign the value 0 to the left part and 1 to the right part.
- Repeat the steps 3 and 4 for each part, until all the symbols are split into individual subgroups.