



CS6109 – GRAPH THEORY

Module – 5

Presented By

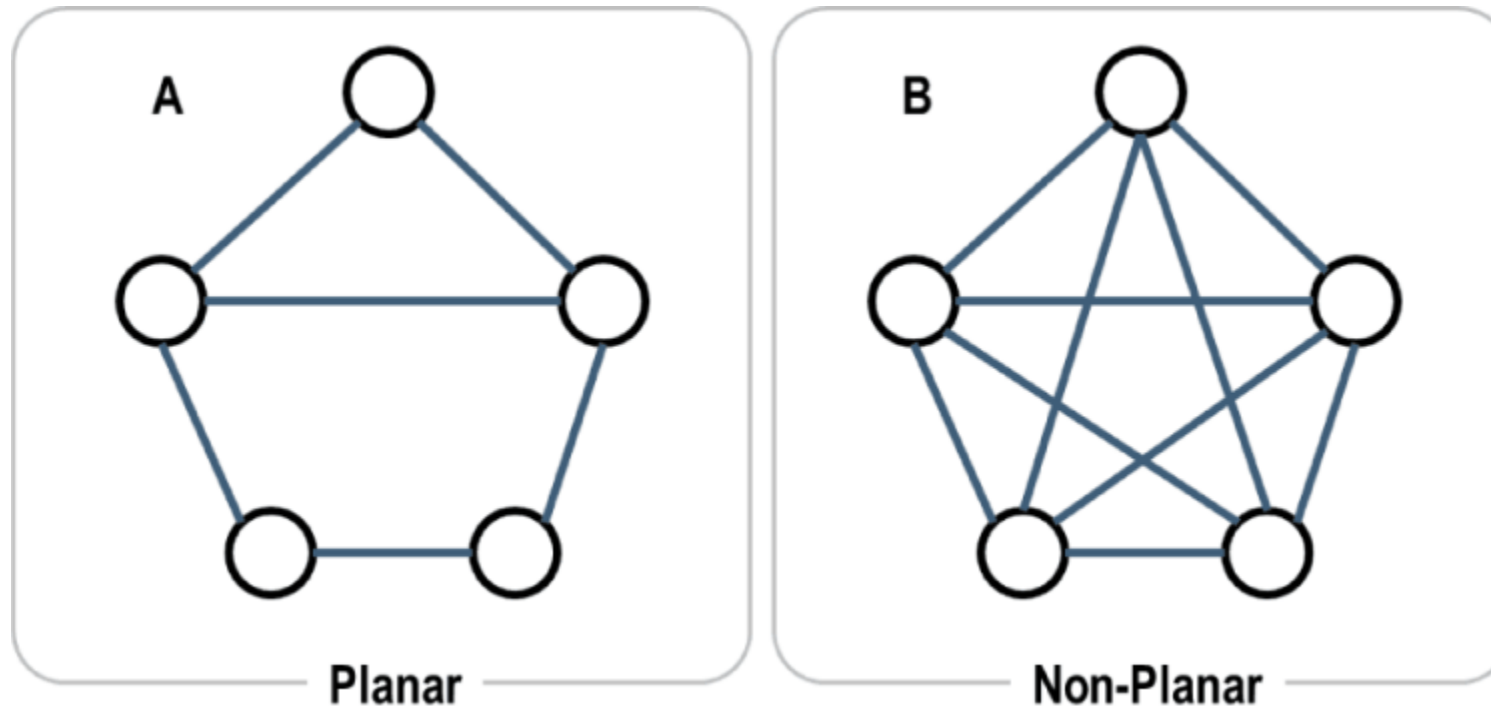
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Module - 5

- Planar Graph
- Representation
- Detection of planarity
- Dual Graph
- Related Theorems

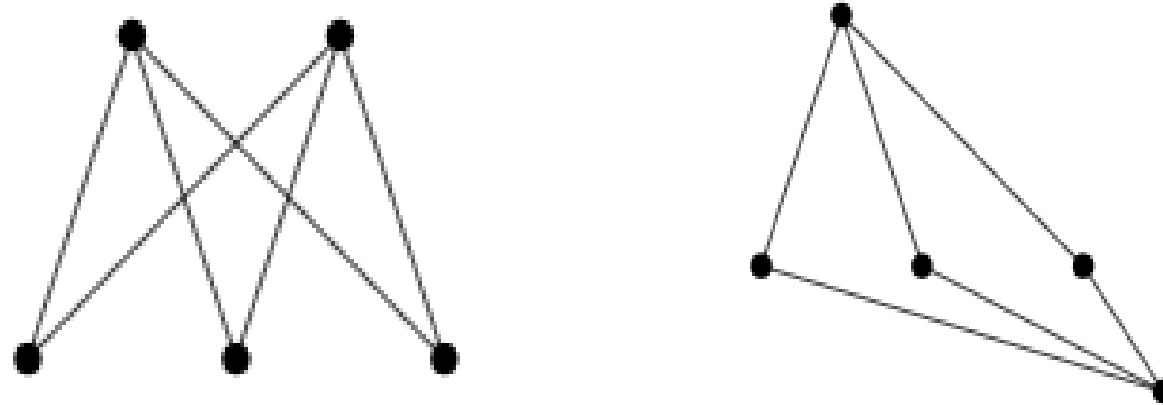
Planar Graph

- A graph is said to be planar if it can be drawn in a plane so that no edge cross.
- When a connected graph can be drawn without any edges crossing, it is called ***planar***. When a planar graph is drawn in this way, it divides the plane into regions called ***faces***.

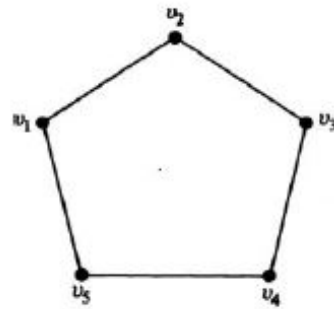


Planar Graph

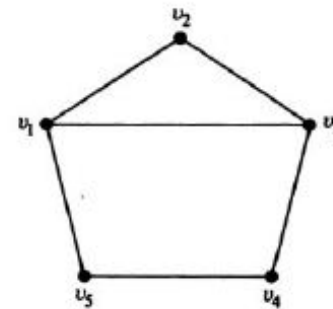
- A drawing of a geometric representation of a graph on any surface such that no edges intersect is called *embedding*.
- To declare that a graph G is nonplanar, we have to show that of all possible geometric representations of G none can be embedded in a plane.
- An embedding of a planar graph G on a plane is called a *plane representation* of G .



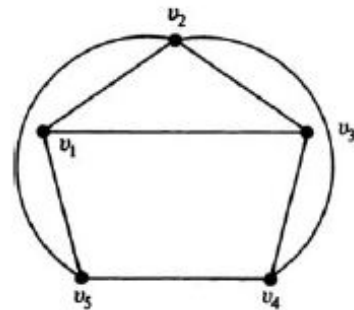
THEOREM 5-1: The complete graph of five vertices is nonplanar.



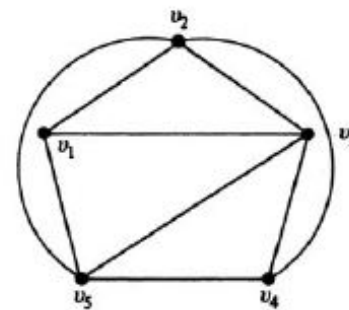
(a)



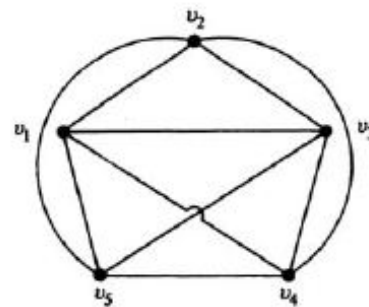
(b)



(c)



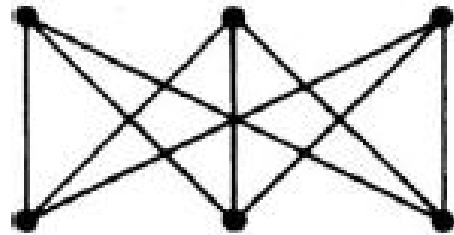
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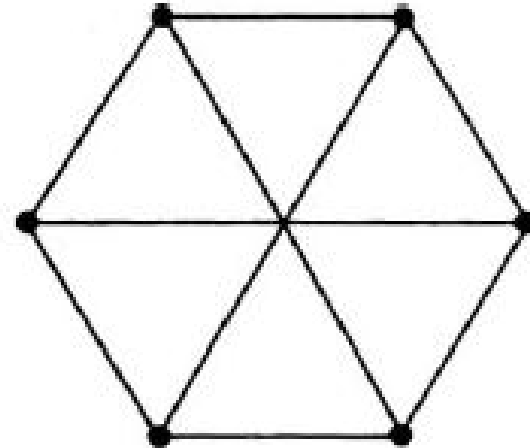
(e)

THEOREM 5-2: Kuratowski's second graph is also nonplanar.

- A regular connected graph with six vertices and nine edges is also nonplanar.



(a)



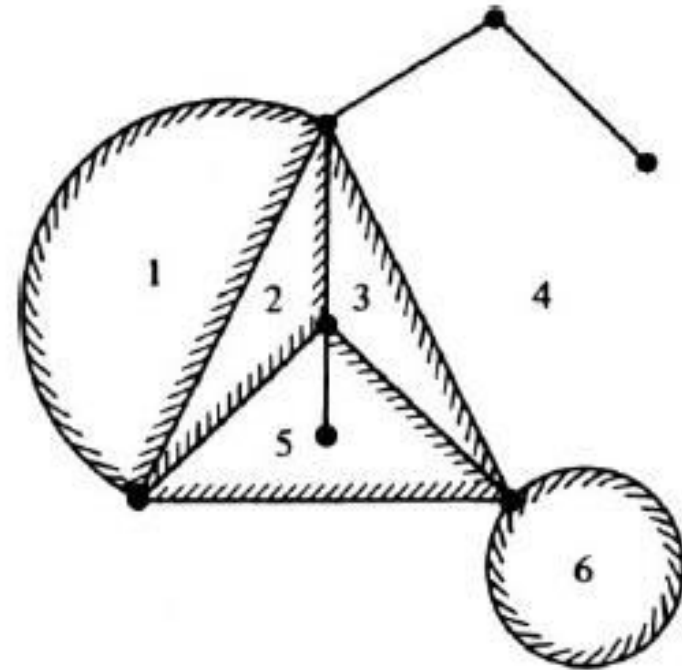
(b)

Representation of planar graph

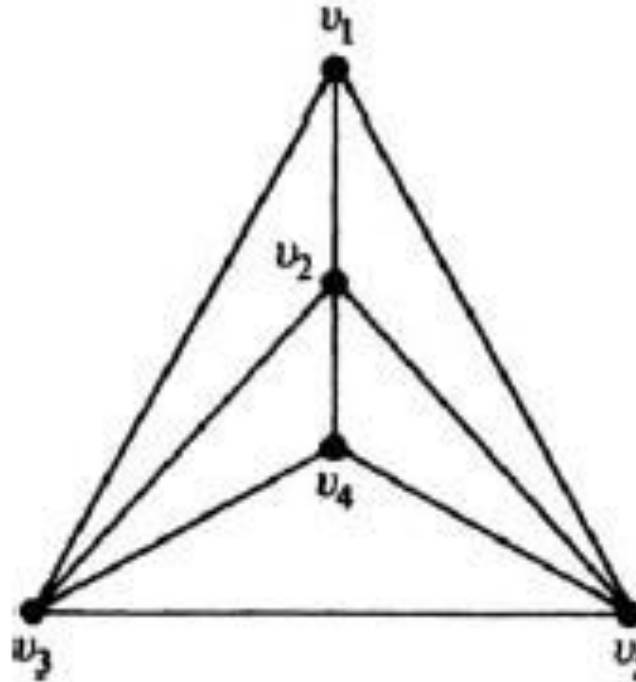
Region of a Graph: Consider a planar graph $G=(V,E)$. A region is defined to be an area of the plane that is bounded by edges and cannot be further subdivided. A planar graph divides the plane into one or more regions. One of these regions will be infinite.

Finite Region: If the area of the region is finite, then that region is called a finite region.

Infinite Region: If the area of the region is infinite, that region is called a infinite region. A planar graph has only one infinite region.



THEOREM 5-3: Any simple planar graph can be embedded in a plane such that every edge is drawn as a straight line segment.



Straight-line representation of the graph

THEOREM 5-6: A connected planar graph with n vertices and e edges has $e - n + 2$ regions.

In a connected planar graph with n vertices, m edges and f regions (faces), $n - m + f = 2$.

Proof:

Without loss of generality, assume that the planar graph is simple. Since any simple planar graph can have a plane representation such that each edge is a straight line, any planar graph can be drawn such that each region is a polygon (a polygon net).

Let the polygon net representing the given graph consist of f regions. Let k_p be the number of p -sided regions.

Since each edge is on the boundary of exactly two regions,

$$3k_3 + 4k_4 + 5k_5 + \dots + rk_r = 2m, \quad \rightarrow \quad (1)$$

where k_r is the number of polygons with r edges.

$$\text{Also, } k_3 + k_4 + k_5 + \dots + k_r = f. \quad \rightarrow \quad (2)$$

The sum of all angles subtended at each vertex in the polygon net is $2\pi n$. $\rightarrow \quad (3)$

Now, the sum of all interior angles of a p -sided polygon is $\pi(p-2)$ and the sum of the exterior angles is $\pi(p+2)$. The expression in (3) is the total sum of all interior angles of $f-1$ finite regions plus the sum of the exterior angles of the polygon defining the infinite region. This sum is

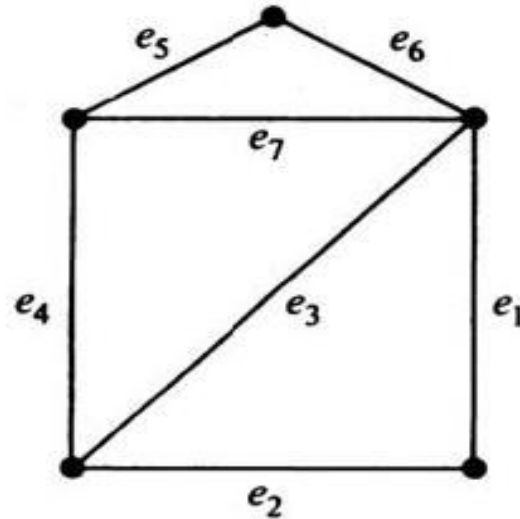
$$\begin{aligned} & \pi(3-2)k_3 + \pi(4-2)k_4 + \dots + \pi(r-2)k_r + 4\pi \\ &= \pi[3k_3 + 4k_4 + \dots + rk_r - 2(k_3 + k_4 + \dots + k_r)] + 4\pi \\ &= \pi(2m - 2f) + 4\pi = 2\pi(m - f + 2). \quad \rightarrow \quad (4) \end{aligned}$$

Equating (3) and (4) we get

$$2\pi(m - f + 2) = 2n\pi, \text{ so that } f = m - n + 2.$$

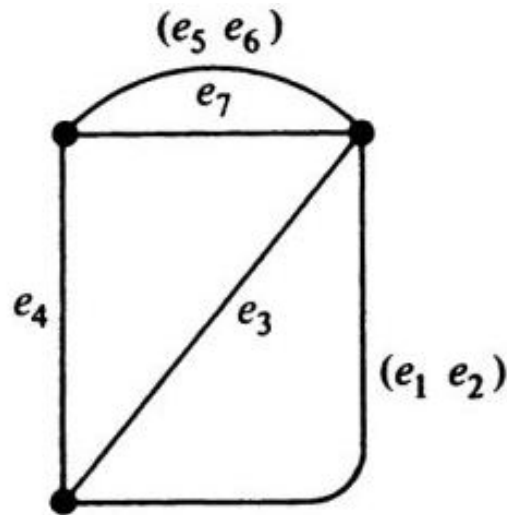
DETECTION OF PLANARITY

Elementary Reduction

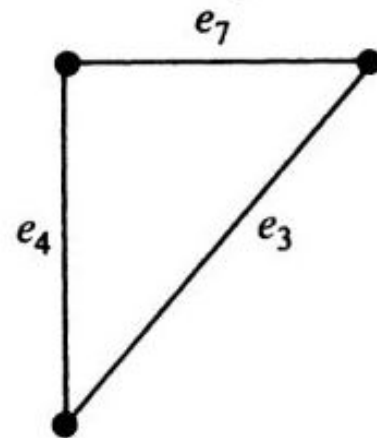


- **Step 1:** Since a disconnected graph is planar if and only if each of its components is planar, we need consider only one component at a time. Also, a separable graph is planar if and only if each of its blocks is planar. Therefore, for the given arbitrary graph G , determine the set $G = \{G_1, G_2, \dots, G_k\}$, where each G_i is a nonseparable block of G . Then we have to test each G_i for planarity.
- **Step 2:** Since addition or removal of self-loops does not affect planarity, remove all self-loops.

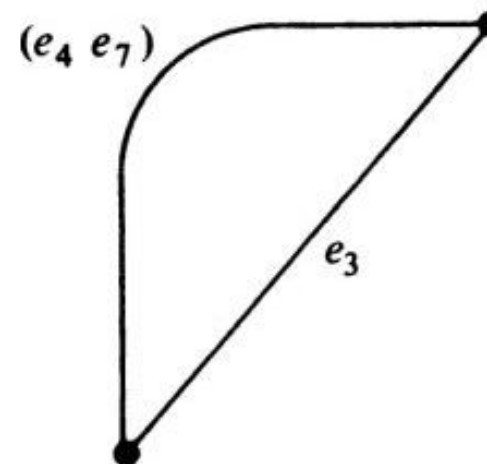
- **Step 3:** Since parallel edges also do not affect planarity, eliminate edges in parallel by removing all but one edge between every pair of vertices.
- **Step 4:** Elimination of a vertex of degree two by merging two edges in series does not affect planarity. Therefore, eliminate all edges in series. Repeated application of steps 3 and 4 will usually reduce a graph drastically.



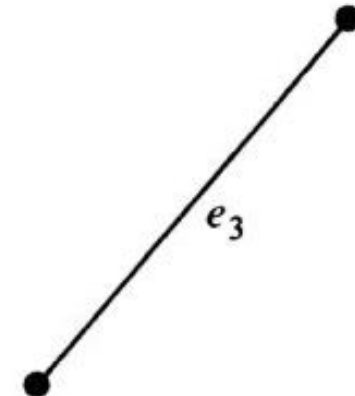
(a) Series Reduced



(b) Parallel Reduced



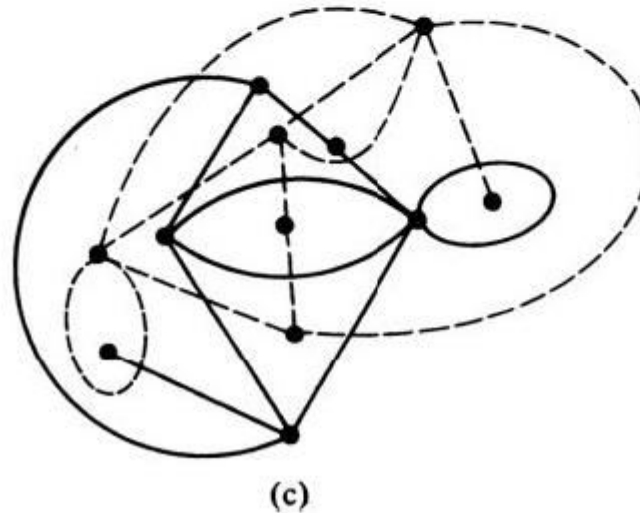
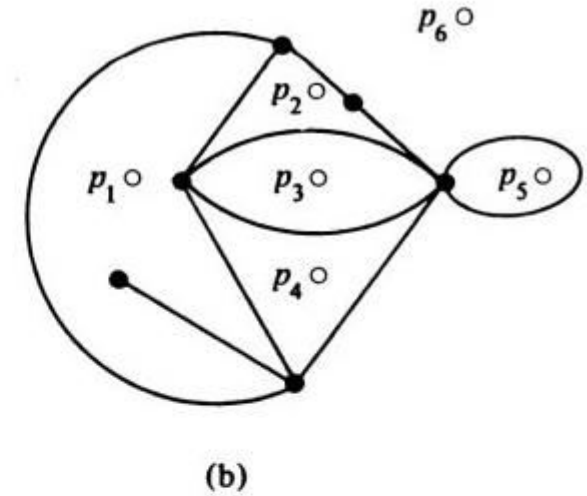
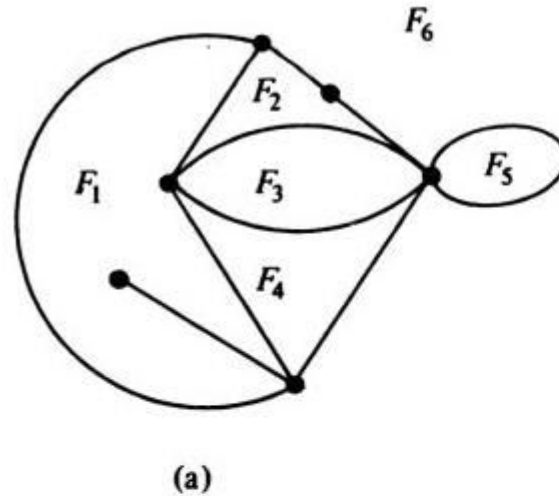
(c) Series Reduced



(d) Parallel Reduced

DUAL GRAPH

1. Divide the graph into region or faces.
2. For each face assign points.
3. Connect the points for each edges separating them.



REFERENCES:

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Thank you.