

Unit-1

Number System, Boolean Algebra & Logic Gates

Number system:-

Number system is a base for counting various items. The modern computers (or) Digital system communicate, operates with Binary numbers which use only digits 0 & 1.

To define any Number system, the following specifications are considered.

⇒ Total no:of digits or symbols

⇒ Based on number system which such as 2,8,16

⇒ Here base indicates the total no:of digits available in number system.

Classification:-

There are 4 types. They are,

* Decimal Number System (base 10)

* Binary Number system (base 2)

* Octal Number system (base 8)

* Hexa Number system (base 16)

Decimal Number Systems:-

The total no:of digits (or) symbol present in a decimal number system are 10.

Since the base (or) Radix of the decimal number system is 10. The digits (or) symbols in decimal number system starts with (10).

Ex:- $567_{(10)}$

From the above example we can say that the decimal number 567 can also be written as $(567)_{10}$, where 10 is base. i.e $(567)_{10} = 5 \times 10^2 + 6 \times 10^1 + 7 \times 10^0$.

∴ The decimal number system is the positional weighted system and its weight is expressing in power of 10.

In heating point representation, the left side digits have positive powers and right side digits have negative powers.

Ex:- 567.87

$$\begin{array}{r} 5 \ 6 \ 7 \ . \ 8 \ 7 \\ \hline 10^2 \ 10^1 \ 10^0 \ 10^{-1} \ 10^{-2} \end{array}$$

$$5 \times 10^2 + 6 \times 10^1 + 7 \times 10^0 + 8 \times 10^{-1} + 7 \times 10^{-2}$$

Binary Number system:-

The total no:of digits or symbols are represented in binary numbers (or) are 2.

Since the base (or) Radix of binary number system is 2.

The elements starts from zero to one.

Ex:- 1001

$$\begin{array}{r} 1 \ 0 \ 0 \ 1 \\ \hline 2^3 \ 2^2 \ 2^1 \ 2^0 \end{array}$$

$$\text{i.e;} 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0$$

$$= 8 + 0 + 0 + 1$$

$$= (9)_2$$

* 1001.01

$$\Rightarrow 1 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 + 0 \times 2^{-1} + 1 \times 2^{-2}$$

$$= 8 + 0 + 0 + 1 + 0 + 0.25$$

$$= 9.25$$

- ⇒ In a Binary^{number} system each binary digit is known as bit.
- ⇒ In Binary number system the weight is expressed as powers of 2; for an n-bit number of combinations are 2^n .

Octal Number system:-

Total no:of digits or symbols present in a octal number system is 8. Since the base/radix of the octal number system is 8 i.e; (0-7).

$$\text{Ex:- } \begin{array}{r} 567 \\ \downarrow \quad \downarrow \quad \downarrow \\ 8^2 \quad 8^1 \quad 8^0 \end{array}$$

$$\Rightarrow 5 \times 8^2 + 6 \times 8^1 + 7 \times 8^0$$

In floating point:- 567.12

$$\Rightarrow 5 \times 8^2 + 6 \times 8^1 + 7 \times 8^0 + 1 \times 8^{-1} + 2 \times 8^{-2}$$

In octal Number system the weight is expressed in powers of 8.

Hexa Decimal Number system:-

The total number of digits or symbols present in a hexa decimal number systems are 16.

Since the base/radix of hexa decimal number system is 16.

$$\text{Ex:- } \begin{array}{r} 4 A 5 \\ \downarrow \quad \downarrow \quad \downarrow \\ 16^2 \quad 16^1 \quad 16^0 \end{array}$$

$$\text{i.e;} 4 \times 16^2 + A \times 16^1 + 5 \times 16^0$$

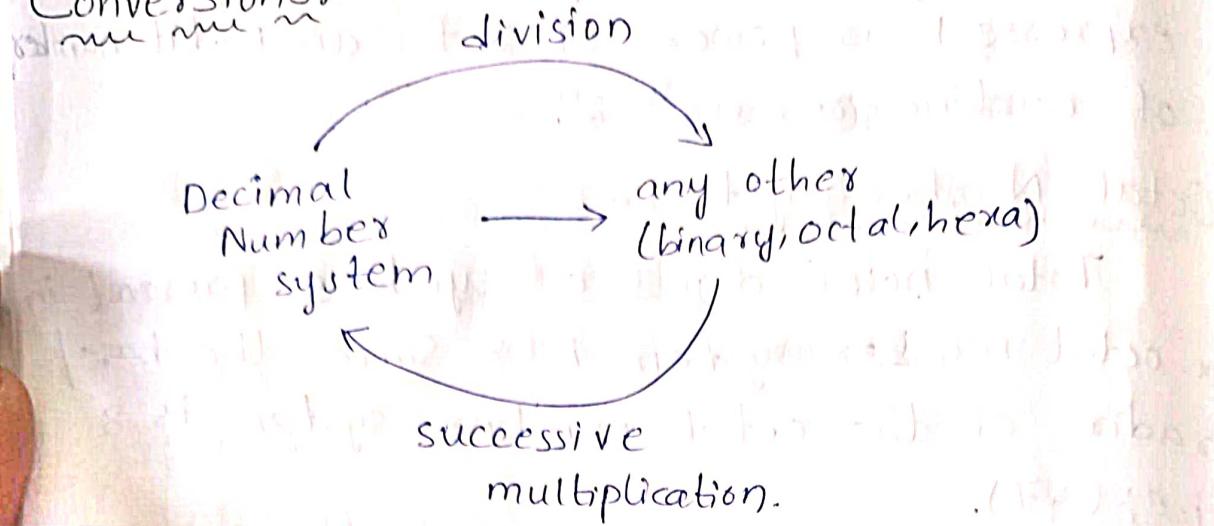
In floating point:-

$$\begin{array}{r} 4 A 5.8 \\ \downarrow \quad \downarrow \quad \downarrow \\ 16^2 \quad 16^1 \quad 16^0 \end{array}$$

$$\text{i.e;} 4 \times 16^2 + A \times 16^1 + 5 \times 16^0 + 8 \times 16^{-1}$$

In hexa decimal number system the weight is expressed in powers of 16.

Conversions:-



* Decimal Number system to any other number system.

$$(i) (40)_{10} \quad (ii) (18)_{10} \quad (iii) (56)_{10}$$

Decimal to Binary system:-

$$(i) \begin{array}{r} 2 | 40 \\ 2 | 20 - 0 \\ 2 | 10 - 0 \\ 2 | 5 - 0 \\ 2 | 2 - 1 \\ \hline 1 - 0 \end{array}$$

$$(ii) \begin{array}{r} 2 | 18 \\ 2 | 9 - 0 \\ 2 | 4 - 1 \\ 2 | 2 - 0 \\ \hline 1 - 0 \end{array}$$

$$(iii) \begin{array}{r} 2 | 56 \\ 2 | 28 - 0 \\ 2 | 14 - 0 \\ 2 | 7 - 0 \\ 2 | 3 - 1 \\ \hline 1 - 1 \end{array}$$

$$(40)_{10} = (101000)_2 \Rightarrow (10010)_2 \Rightarrow (111000)_2$$

Decimal to octal number system:-

$$(i) (40)_{10} \quad \begin{array}{r} 8 | 40 \\ 8 | 5 - 0 \end{array}$$

$$\Rightarrow (50)_8$$

$$(ii) (214)_{10} \quad \begin{array}{r} 8 | 214 \\ 8 | 26 - 6 \\ 8 | 3 - 2 \end{array}$$

$$\Rightarrow (326)_8$$

$$(iii) (3509)_{10} \quad \begin{array}{r} 8 | 3509 \\ 8 | 438 - 5 \\ 8 | 54 - 6 \\ \hline 6 - 6 \end{array}$$

$$\Rightarrow (6665)_8$$

Decimal to Hexa decimal number system:-

$$(i) (40)_{10} \quad \begin{array}{r} 16 | 40 \\ 16 | 2 - 8 \end{array}$$

$$(28)_{16}$$

$$(ii) (214)_{10} \quad \begin{array}{r} 16 | 214 \\ 16 | 13 - 6 \\ \hline 0 = 13 \end{array}$$

$$(D6)_{16}$$

$$(iii) (3509)_{10} \quad \begin{array}{r} 16 | 3509 \\ 16 | 219 - 5 \\ \hline 3 = 11 \end{array}$$

$$\Rightarrow (DB5)_{16}$$

Binary to Decimal number system:-

(i) $(1010)_2$

$\begin{array}{r} 1010 \\ \downarrow \downarrow \downarrow \downarrow \\ 2^3 2^2 2^1 2^0 \end{array}$

$$\Rightarrow 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$$

$$\Rightarrow 8 + 0 + 2 + 0$$

$$\Rightarrow (10)_{10}$$

(ii) $(1010110000)_2$

$$\Rightarrow 0 \times 2^8 + 1 \times 2^7 + 0 \times 2^6 + 1 \times 2^5 +$$

$$1 \times 2^4 + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 +$$

$$\Rightarrow 2^7 + 2^5 + 2^4 \Rightarrow 128 + 32 + 16$$

$$\Rightarrow (176)_{10}.$$

Binary to Octal number system:-

(i) $(010110001)_2$

$$(010\cancel{11}0001)_2$$

$$(261)_8$$

(ii) $(010011010)_2$

$$(0100\cancel{11}010)_2$$

$$(232)_8$$

Binary to Hexa decimal number system:-

(i) $(000010110001)_2$

$$(0000101\cancel{1}0001)_2$$

$$0 \quad 11 \quad 1$$

$$\Rightarrow (0B1)_{16}$$

(ii) $(1000101)_2$

$$(1000\cancel{1}01)_2$$

$$4 \quad 5$$

$$\Rightarrow (45)_{16}$$

| | | |
|-------|-------|-------|
| 2^2 | 2^1 | 2^0 |
| 4 | 2 | 1 |

| | | | |
|---|---|---|---|
| 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 1 |

| | | | |
|---|---|---|---|
| 2 | 0 | 1 | 0 |
| 3 | 0 | 1 | 1 |

| | | | |
|---|---|---|---|
| 4 | 1 | 0 | 0 |
| 5 | 1 | 0 | 1 |

| | | | |
|---|---|---|---|
| 6 | 1 | 1 | 0 |
| 7 | 1 | 1 | 1 |

| | | | |
|---|---|---|---|
| 8 | 1 | 1 | 1 |
| 9 | 1 | 1 | 0 |

| | | | |
|----|---|---|---|
| 10 | 1 | 0 | 0 |
| 11 | 1 | 0 | 1 |

| | | | |
|----|---|---|---|
| 12 | 0 | 1 | 1 |
| 13 | 1 | 1 | 0 |

| | | | |
|----|---|---|---|
| 14 | 1 | 1 | 1 |
| 15 | 1 | 1 | 1 |

* Octal to Decimal number system:-

(i) $(324)_8$ (ii) $(473)_8$

$$\begin{aligned} i.e. & 3 \times 8^2 + 2 \times 8^1 + 4 \times 8^0 \\ \Rightarrow & 192 + 16 + 4 \\ \Rightarrow & (212)_{10} \end{aligned} \qquad \begin{aligned} \Rightarrow & 4 \times 8^2 + 7 \times 8^1 + 3 \times 8^0 \\ \Rightarrow & 4 \times 64 + 56 + 3 \\ \Rightarrow & (315)_{10} \end{aligned}$$

* Octal to Binary number system:-

(i) $(473)_8$

$$\Rightarrow (473)_8$$

$$\Rightarrow (100111011)_2$$

(ii) $(273)_8$

$$\Rightarrow (273)_8$$

$$\Rightarrow (010111011)_2$$

* Octal to hexa decimal system:-

(i) $(324)_8$

(ii) $(526)_8$

(iii) $(273)_8$

(iv) $(5327)_8$

$$(i) (324)_8 = 011010100 \Rightarrow (D4)_{16}$$

$$(ii) (526)_8 = 101010110 \Rightarrow (156)_{16}$$

$$(iii) (273)_8 = 010111011 \Rightarrow (8BB)_{16}$$

$$(iv) (5327)_8 = 101011010111 \Rightarrow (AD7)_{16}$$

* Hexa decimal to decimal system:-

(i) $(56)_{16}$

$$\Rightarrow 5 \times 16^1 + 6 \times 16^0$$

$$= 80 + 6$$

$$= (86)_{10}$$

(ii) $(3AD)_{16}$

$$\Rightarrow 3 \times 16^2 + A \times 16^1 + D \times 16^0$$

$$= 3 \times 256 + 10 \times 16 + 13 \times 1$$

$$= 768 + 160 + 13$$

$$= (941)_{10}$$

Hexa decimal to Binary system:-

(i) $(FBC)_6 \Rightarrow (111110111100)_2$

(ii) $(AF25)_{16} = (1010111100100101)_2$

* Hexa decimal to Octal number system:-

(i) $(FBC)_{16} = 111110111100 \Rightarrow (7674)_8$

(ii) $(ABC)_{16} = 101010111100 \Rightarrow (5274)_8$

* Fractional Point Conversions:-

* Decimal to Binary conversion:-

$$\underline{\text{Ex:}} (0.8125)_{10}$$

$$\xrightarrow[\text{Division (LCM)}]{\quad} 8125 \times 2 \Rightarrow 1.625$$

$$0.625 \times 2 \Rightarrow 1.25$$

$$0.25 \times 2 \Rightarrow 0.5$$

$$0.5 \times 2 \Rightarrow 0.0$$

$$(0.1101)_2$$

* Decimal to Octal :-

$$\underline{\text{Ex:}} (0.640625)_{10}$$

$$0.640625 \times 8 \Rightarrow 5.125$$

$$0.125 \times 8 \Rightarrow 1.00$$

$$\Rightarrow (0.51)_8$$

* Decimal to Hexa-decimal:-

$$\underline{\text{Ex:}} (22.64)_{10}$$

$$16 \overline{)22} \qquad \qquad 0.64 \times 16 \Rightarrow 10.24$$

$$\boxed{1-6} \qquad \qquad 0.24 \times 16 \Rightarrow 3.84$$

$$(16) \qquad 100111 \qquad 0.84 \times 16 \Rightarrow 13.44$$

$$0.44 \times 16 \Rightarrow 7.04$$

$$0.04 \times 16 \Rightarrow 0.64$$

$$\Rightarrow (16. A3070)_{16}$$

Binary conversions:-

* Binary conversion addition

* Binary subtraction

* Binary multiplication

* Binary division.

* Binary Addition:-

| A | B | $A+B$ | carry |
|---|---|-------|-------|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |

Ex:-

$$\begin{array}{r} 1010 \\ 1110 \\ \hline 11000 \end{array}$$

$$\begin{array}{r} 101011 \\ 111000 \\ \hline 1100011 \end{array}$$

* Binary subtraction:-

| A | B | $A-B$ | barrow |
|----|---|-------|--------|
| 10 | 0 | 0 | 0 |
| 0 | 1 | -1 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |

$$\begin{array}{r} 1110 \\ 0011 \\ \hline 1011 \end{array}$$

$$\begin{array}{r} 111001 \\ 001010 \\ \hline 110111 \end{array}$$

* Binary multiplication:-

| A | B | $A \cdot B$ |
|---|---|-------------|
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

$$\text{Ex:- (i) } 1101 \times 110 \quad \text{(ii) } 1110 \times 101$$

$$\begin{array}{r}
 1101 \times 110 \\
 \hline
 0000 \\
 1101 \times \\
 1101 \times x \\
 \hline
 1001110
 \end{array}$$

$$\begin{array}{r}
 1110 \times 101 \\
 \hline
 1110 \\
 0000 \times \\
 1110 \times x \\
 \hline
 1000110
 \end{array}$$

* Binary division:-

$$\text{Ex:- } 101101 \div 110$$

$$\begin{array}{r}
 \text{neglect} \\
 110) 001101 (111.1 \\
 \hline
 110 \\
 \hline
 0010 \\
 110 \\
 \hline
 001 \\
 110 \\
 \hline
 001 \\
 110 \\
 \hline
 0
 \end{array}$$

$$101110 \div 101$$

$$\begin{array}{r}
 101) 001110 (110 \\
 \hline
 101 \\
 \hline
 010 \\
 101 \\
 \hline
 00 \\
 0 \\
 \hline
 0
 \end{array}$$

Complements:-

$$1's - 2's, \quad 9's - 10's.$$

$$\text{Ex:- } (10010)_2$$

$$\begin{array}{r}
 10010 \\
 1's \rightarrow 01101 \xrightarrow{\text{complement of}} \\
 \hline
 2's \rightarrow 01110
 \end{array}$$

$$\begin{array}{r}
 562 \\
 \hline
 2^5 \rightarrow 563
 \end{array}$$

2's complement
is done with any
~~one can add~~
~~to~~

1's complement
is only with
0 and 1

$$\text{Ex:- } 9's - 10's$$

$$(i) 468$$

$$\begin{array}{r}
 999 \\
 -468 \\
 \hline
 531
 \end{array}$$

$$\begin{array}{r}
 531 \\
 1 - \text{add 1} \\
 \hline
 532
 \end{array}$$

$$(ii) 3524$$

$$\begin{array}{r}
 9999 \\
 -3524 \\
 \hline
 6475
 \end{array}$$

$$\begin{array}{r}
 6475 \\
 1 \\
 \hline
 6476
 \end{array}$$

$$(iii) 6234$$

$$\begin{array}{r}
 9999 \\
 -6234 \\
 \hline
 3765
 \end{array}$$

$$\begin{array}{r}
 3765 \\
 1 \\
 \hline
 3766
 \end{array}$$

$$\begin{array}{r}
 3765 \\
 1 \\
 \hline
 3766
 \end{array}$$

Other binary codes:-

(1) BCD code / 8421 code

(2) Excess-3 code ($X_5 - 3$)

(3) Gray code

(1) BCD code

| | 8 4 2 1 |
|----|---------|
| 0 | 0 0 0 0 |
| 1 | 0 0 0 1 |
| 2 | 0 0 1 0 |
| 3 | 0 0 1 1 |
| 4 | 0 1 0 0 |
| 5 | 0 1 0 1 |
| 6 | 0 1 1 0 |
| 7 | 0 1 1 1 |
| 8 | 1 0 0 0 |
| 9 | 1 0 0 1 |
| 10 | 1 0 1 0 |
| 11 | 1 0 1 1 |
| 12 | 1 1 0 0 |
| 13 | 1 1 0 1 |
| 14 | 1 1 1 0 |
| 15 | 1 1 1 1 |

(2) $X_5 - 3$ code { decimal no. (0-9) }

| | 8 4 2 1 |
|--------|-----------|
| 0+3=3 | 0 0 1 1 |
| 0+3=4 | 0 1 0 0 |
| 2+3=5 | 0 1 0 1 |
| 3+3=6 | 0 1 1 0 |
| 4+3=7 | 0 0 1 1 1 |
| 5+3=8 | 1 0 0 0 |
| 6+3=9 | 1 0 0 1 |
| 7+3=10 | 1 0 1 0 |
| 8+3=11 | 1 0 1 1 |
| 9+3=12 | 1 1 0 0 |

Gray Code:-

(1) Binary to Gray code

(2) Gray code to Binary

if we get different \rightarrow '1' is easy to convert binary form.

same \rightarrow '0'.

* The gray code is a non-weighted code and it is

not suitable for Arithmetic operations.

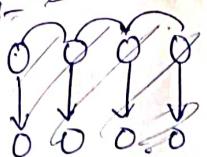
* The only reason for the pa-

larity of gray code is if

different \rightarrow '1' is easy to convert binary form.

* Gray code is also called as "Unit distance code".

* There is only 1 bit change b/w the two consecutive numbers.

Q:- D = (0-09) Binary to Gray conversion :-

 In this the first bit is constant, and adjacent bits were added compared simultaneously.

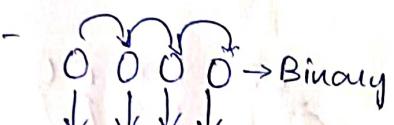
⇒ If we have different bits then '1' is the answer.

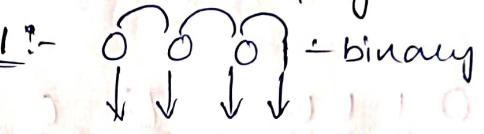
⇒ If we have same bits the '0' is the answer.

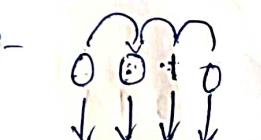
different - 1

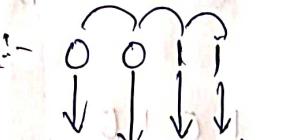
same - 0

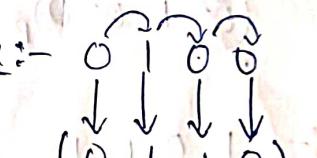
Convert the decimal number to Gray code (Q9)

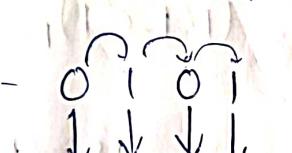
Q1:- 
 (0 0 0 0)_{Gray}

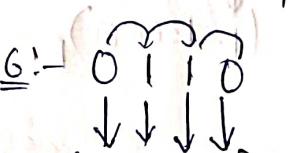
Q2:- 
 (0 0 0 1)_G

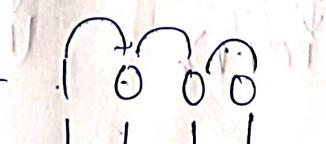
Q3:- 
 (0 0 1 1)_G

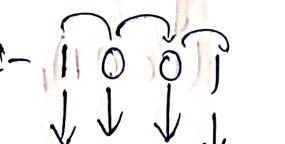
Q4:- 
 (0 0 1 0)_G

Q5:- 
 (0 1 1 0)_G

Q6:- 
 (0 1 1 1)_G

Q7:- 
 (0 1 0 0)_G

Q8:- 
 (1 1 0 0)_G

Q9:- 
 (1 1 0 1)_G

Convert the gray to binary (Q-15) :-

⇒ In this the first bit is constant, the resultant answer added compared to the next bit (second bit).

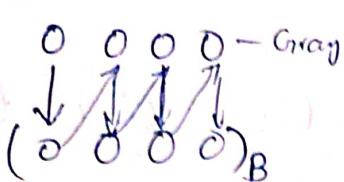
⇒ If we have different bits then '1' is the answer.

→ If we have same bits '0' is the answer.

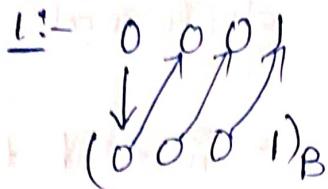
different - 1

same - 0.

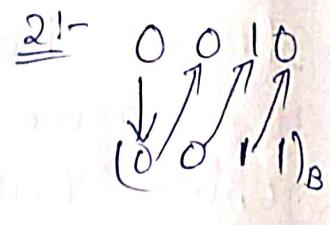
0 :-



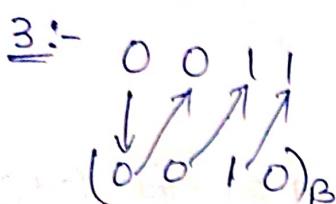
1 :-



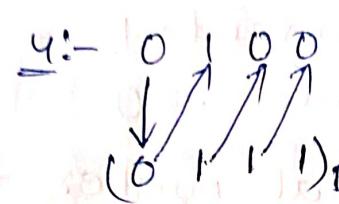
2 :-



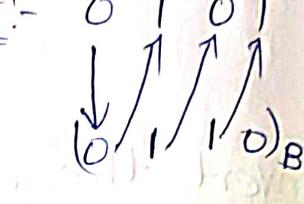
3 :-



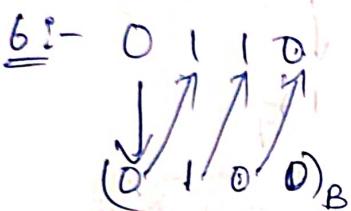
4 :-



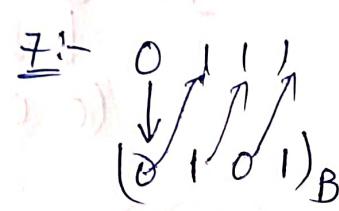
5 :-



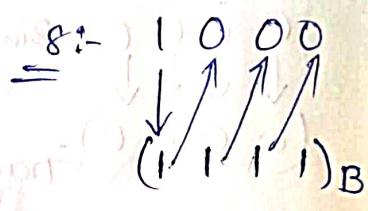
6 :-



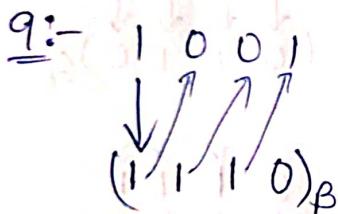
7 :-



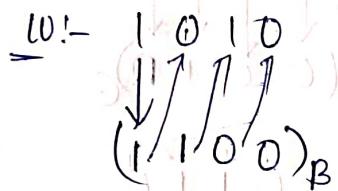
8 :-



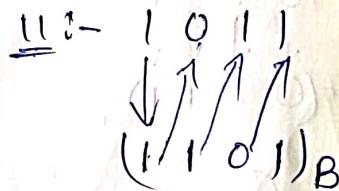
9 :-



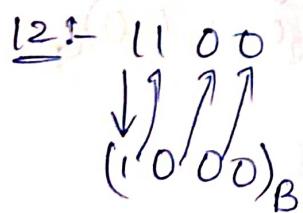
10 :-



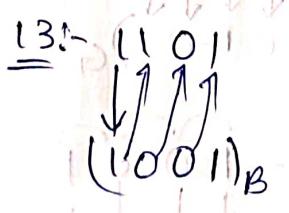
11 :-



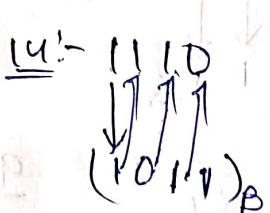
12 :-



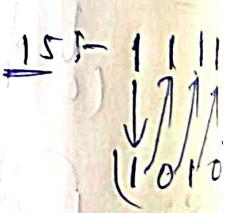
13 :-



14 :-



15 :-



Binary to Gray code:

$$\star (11011001)_B \quad * (00100110)_B \quad * (111000111)_B$$

$$(11011001)_B$$
$$00100110$$
$$111000111$$

$$10110011$$
$$00100110$$
$$111000111$$

$$\star (00001111)_B \quad * (10111000)_B$$

$$(00001111)_B$$
$$10111000$$

$$10110010$$
$$10111000$$

Gray = binary :-

$$\begin{array}{r} \text{* } (10110101)_G \\ \text{↓ } 1 \text{ 1 } 1 \text{ 1 } 1 \\ \text{↓ } (11010001)_B \end{array}$$

$$\begin{array}{r} \text{* } (1100111)_G \\ \text{↓ } 1 \text{ 1 } 1 \text{ 0 } 1 \text{ 1 } 1 \\ \text{↓ } (10110110)_B \end{array} \quad \begin{array}{r} \text{* } (10001000)_G \\ \text{↓ } 1 \text{ 0 } 0 \text{ 0 } 1 \text{ 0 } 0 \text{ 0 } \\ \text{↓ } (1110000)_B \end{array}$$

$$\begin{array}{r} \text{* } (1111010)_G \\ \text{↓ } 1 \text{ 1 } 1 \text{ 1 } 0 \text{ 1 } 0 \\ \text{↓ } (10101100)_B \end{array} \quad \begin{array}{r} \text{* } (00001110)_G \\ \text{↓ } 0 \text{ 0 } 0 \text{ 0 } 1 \text{ 1 } 1 0 \\ \text{↓ } (00001011)_B \end{array}$$

Binary one's compliment for subtraction:-

⇒ Find out the one's compliment of subtrahend and added to minuend (largest no.). If there is a carry bring that carry and add to the LSB (lost significant bit).

⇒ Look at the msb (most significant bit). If it is zero the result is positive and it is in true binary form.

⇒ If the msb is 1 (whether irrespective of carry the result is negative and it is in 1's compliment form)

⇒ Take one's compliment of result to get binary value.

Ex:- Add 27.50 from 68.75

Here the minuend is 68.75 / greater number
subtrahend (lesser number.) = 27.50.

$$\begin{array}{r} 2 | 27 \\ 2 | 12-1 \\ 2 | 6-1 \\ 2 | 3-0 \\ \boxed{1-1} \\ (11011) \end{array}$$

$$\begin{array}{r} 2 | 68 \\ 2 | 34-0 \\ 2 | 17-0 \\ 2 | 8-1 \\ 2 | 4-0 \\ 2 | 2-0 \\ \hline 100100 \end{array}$$

Add 27.50 from 68.75.

$$68.75 - 01000100.1100$$
$$27.50 - 00011.011000$$

complement

$$27.50 - 11100100.0111$$

$$68.75 - 01000100.1100$$

$$\begin{array}{r} & & 1 & 1 \\ & \overline{000701001.0011} \\ \text{carries} & \nearrow & & \searrow \\ \hline 00101001.0100 \end{array}$$

Ex: 52 - 20.

$$\begin{array}{r} 52 \\ 2 | 26 - 0 \\ 2 | 13 - 0 \\ 2 | 6 - 1 \\ 2 | 3 - 0 \\ 2 | 1 - 1 \end{array}$$

$$\begin{array}{r} 20 \\ 2 | 10 - 0 \\ 2 | 5 - 0 \\ 2 | 2 - 1 \\ 1 - 0 \end{array}$$

$$(110100) \quad (10100)$$

$$20 - 110100$$
$$52 - \underline{110100}$$

$$110$$

$$001001$$

- * Binary subtraction for 2's complement.
 - Find out the 2's complement of subtrahend and add it with the minuend.
 - If there is a carry ignore it. And look at the msB of the sum.
 - If msB is zero, the result is positive and it is the true binary number.
 - If the msB is "one" (whether there is a carry or no carry), the result is negative and it is in 2's complement form.
- ⇒ Take its 2's complement find out the true binary value.

Ex:- subtract 14 from 46 using 8 bit 2's compliment method. ($(46)_{10} - (14)_{10}$)

$$\begin{array}{r} 2 \Big| 14 \\ 2 \Big| 7 - 0 \\ 2 \Big| 3 - 1 \\ \hline 1 - 1 \end{array}$$

$$(1110)_B$$

$$\begin{array}{r} 2 \Big| 46 \\ 2 \Big| 23 - 0 \\ 2 \Big| 11 - 1 \\ \hline 5 - 1 \\ 2 \Big| 2 - 1 \\ \hline 1 - 0 \end{array}$$

$$(10110)_B$$

$$14 - 00001110$$

$$46 - 00101110$$

$$\Rightarrow 14 - 11110001$$

$$\begin{array}{r} & 11 \\ \hline 11110010 - 00101110 \\ \hline 11010000 \end{array}$$

BCD Addition:-

- Convert a given decimal number into their equivalent BCD codes.
- Add the BCD numbers using the rules of binary addition.
- Check the result, if it is valid (≤ 9) no correction is needed.
- If the result is invalid add 0110 i.e '6' to the four bits sum to get the correct result.
- Ex:- perform the BCD addition for 25 and 13.

$$\begin{array}{r}
 \text{BCD} + 25 \\
 \text{---} \\
 \begin{array}{r}
 \begin{array}{c} 0010 & 0101 \\ \swarrow & \searrow \\ 13 - 0001 & 0011 \end{array} \\
 \begin{array}{r}
 \begin{array}{r}
 \begin{array}{r}
 0011 & 1000 \\
 \hline
 \downarrow & \downarrow \\
 3 & 8
 \end{array}
 \end{array}
 \end{array}
 \end{array}$$

$\therefore 3 \leq 9$, it is valid.

$\therefore 8 \leq 9$, it is valid.

\therefore The total sum is '38'.

* 91 and 81

$$\begin{array}{r}
 \text{BCD} - 91 \\
 \text{---} \\
 \begin{array}{r}
 \begin{array}{c} 1001 & 0001 \\ \swarrow & \searrow \\ 81 - 1000 & 0001 \end{array} \\
 \begin{array}{r}
 \begin{array}{r}
 \begin{array}{r}
 10001 & 0010 \\
 \hline
 \downarrow & \downarrow \\
 1 & 2
 \end{array}
 \end{array}
 \end{array}
 \end{array}$$

add '6' $\leftarrow 0110$

$\begin{array}{r} 00010011 \\ \hline 17 \end{array}$

, 2) is ≤ 9 , it is valid

∴ The total sum is 172.

BCD subtraction:-

→ The subtraction of BCD numbers can be performed by using q's & 10's complement methods.

→ find the q's & 10's compliment of the given -ve numbers.

→ Add 2 numbers to the BCD addition.

→ If the carry is generated add that carry to the result for correction.

Ex:- perform the subtraction in BCD by the q's compliment method. $305.5 - 168.8$

q's compliment of $168.8 = 999.9$

$$\begin{array}{r} 1110 \quad 0110 \quad 0001 \\ + 1110 \quad 0110 \quad 0001 \\ \hline 1000 \quad 1100 \quad 0000 \end{array} \xrightarrow{-} \begin{array}{r} 168.8 \\ 831.1 \end{array}$$

$\Rightarrow 305.5$

$$\begin{array}{r} 1110 \quad 0110 \quad 0001 \\ + 1110 \quad 0110 \quad 0001 \\ \hline 0000 \quad 1000 \quad 0000 \end{array}$$

BCD:-

$305.5 \Rightarrow 0011 \quad 0000 \quad 0101 \cdot 0101$

$831.1 \Rightarrow 1000 \quad 0011 \quad 0001 \cdot 0001$

$$\begin{array}{r} 1011 \quad 0011 \quad 0110 \cdot 0110 \\ \downarrow \quad \downarrow \quad \downarrow \quad \downarrow \\ 11 \quad 3 \quad 6 \quad 6 \end{array}$$

Here $11 > 9$; 1011

add 8 $\leftarrow 0110$

$$\begin{array}{r} 11 \\ 0001 \quad 0011 \quad 0110 \cdot 0110 \\ + 0001 \quad 0011 \quad 0110 \cdot 0110 \\ \hline 0001 \quad 0011 \quad 0110 \cdot 0111 \end{array}$$

* Note:-

While doing subtraction by using BCD if there is a borrow from the next group then it is subtracted from the difference.

⇒ If there is no borrow from the next higher group then no correction is required.

Ex:- $679.6 - 885.9$

9's compliment of $885.9 \Rightarrow 999.9$

$$\begin{array}{r} 885.9 \\ \underline{-} \\ 114.0 \end{array}$$

BCD:

$$679.6 \Rightarrow 0110 \ 0111 \ 1001.0110$$

$$885.9 \Rightarrow 0000 \ 1000 \ 0100.0000$$

$$\begin{array}{r} & & 1 & 1 & 1 \\ & 0 & 1 & 1 & 1 & 1000 & 1 & 101 & 0 & 110 \\ & \downarrow & & \downarrow & & \downarrow & & \downarrow & & \downarrow \\ & 7 & & 8 & & 13 & & 10 & & 6 \\ \hline & 0 & 1 & 1 & 1 & 1000 & 0110 & 0110 & & \\ & \cancel{0} & \cancel{0} & \cancel{0} & \cancel{1} & \cancel{001} & \cancel{1} & \cancel{1} & & \end{array}$$

Signed binary numbers:-
The sum is -6 .

⇒ If the msB bit is '0' then the resulting value is 'tve', and for remaining have to perform binary addition.

⇒ If the msB bit is '1' then the resulting value is 'lve', and remaining have to perform binary addition.

Ex:- ① $\begin{array}{r} 1 & 0 & 1 & 0 & 0 & 1 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 \end{array}$

$$32 + 0 + 8 + 1 \Rightarrow 41$$

Ex:- $\begin{array}{r} 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 2^5 & 2^4 & 2^3 & 2^2 & 2^1 & 2^0 \end{array}$

$$\Rightarrow -41$$

0 1 0 1 1 0

$$\begin{aligned} &\Rightarrow 1 \times 2^5 + 0 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 \\ &\Rightarrow 32 + 0 + 8 + 4 + 2 + 0 \\ &\Rightarrow 46. \end{aligned}$$

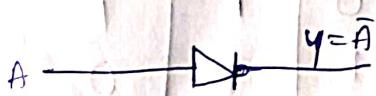
Digital logic Gates:-

- ⇒ Logic gate is a digital circuit which has one or more inputs & only 1 output. The inputs or outputs are logic '0' or '1'.
- ⇒ These logic gates are basic building blocks for digital system.

Types of logic gates:-

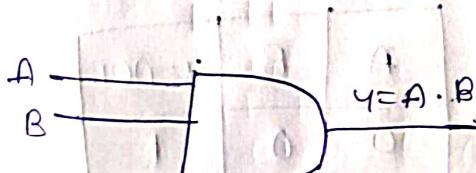
- | | | |
|-----------|---|-----------------|
| (1) NOT | } | basic gates |
| (2) AND | | |
| (3) OR | | |
| (4) NAND | } | universal gates |
| (5) NOR | | |
| (6) X-NOR | } | special gates |
| (7) X-NOR | | |

NOT:-



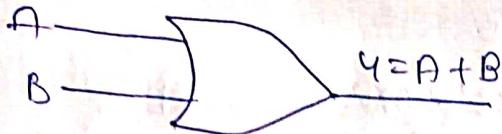
| A | $y = \bar{A}$ |
|---|---------------|
| 0 | 1 |
| 1 | 0 |

AND:-



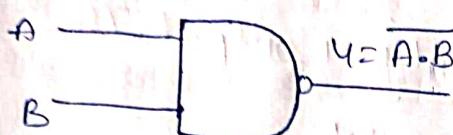
| A | B | $y = A \cdot B$ |
|---|---|-----------------|
| 0 | 0 | 0 |
| 1 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 1 | 1 |

OR :-



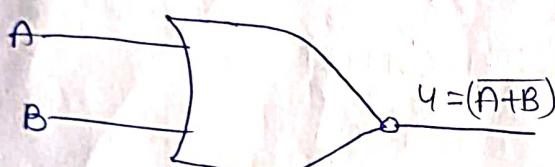
| A | B | $Y = A + B$ |
|---|---|-------------|
| 0 | 0 | 0 |
| 1 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 1 | 1 |

NAND :-



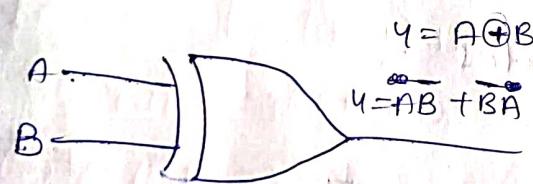
| A | B | $Y = \overline{A} \cdot \overline{B}$ |
|---|---|---------------------------------------|
| 0 | 0 | 1 |
| 1 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 1 | 0 |

NOR :-



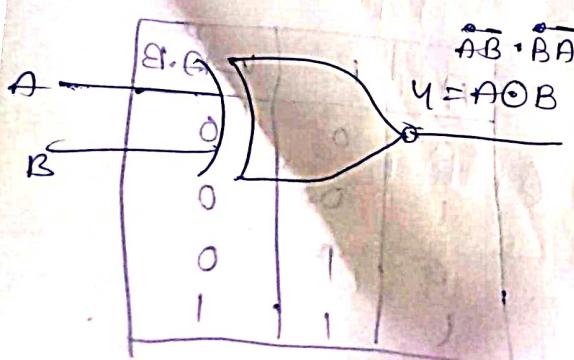
| A | B | $Y = \overline{A} + \overline{B}$ |
|---|---|-----------------------------------|
| 0 | 0 | 1 |
| 1 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 1 | 0 |

X-OR :- (odd is out)



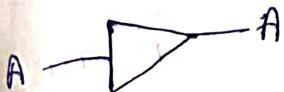
| A | B | $Y = A \oplus B$ |
|---|---|------------------|
| 0 | 0 | 0 |
| 1 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 1 | 0 |

X-NOR :- (even is out)



| A | B | $Y = A \ominus B$ |
|---|---|-------------------|
| 0 | 0 | 1 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

Buffer:



$$0 \rightarrow 0$$

$$1 \rightarrow 1$$

Boolean Algebra:-

Rules of Boolean Algebra.

$$(1) A + 0 = A$$

$$(2) A + 1 = 1$$

$$(3) A \cdot 0 = 0$$

$$(4) A \cdot 1 = A$$

$$(5) A + A = A$$

$$(6) A + \bar{A} = 1$$

$$(7) A \cdot A = A$$

$$(8) A \cdot \bar{A} = 0$$

$$(9) \bar{\bar{A}} = A$$

$$(10) A + AB = A$$

$$(11) A + \bar{A}B = A + B$$

$$(12) (A+B)(A+C) = A+B C$$

Laws and Properties of Boolean Algebra:-

| Name | <u>AND form</u> | <u>OR form</u> |
|--------------------------------------|-----------------------|---------------------|
| (1) Identity law | $1 \cdot A = A$ | $0 + A = A$ |
| (2) Null law | $0 \cdot A = 0$ | $1 + A = 1$ |
| (3) Idempotent Idempotent | $A \cdot A = A$ | $A + A = A$ |
| (4) Inverse law | $A \cdot \bar{A} = 0$ | $A + \bar{A} = 1$ |
| (5) Commutative law | $AB = BA$ | $A + B = B + A$ |
| (6) Associative law | $(AB)C = A(BC)$ | $(A+B)+C = A+(B+C)$ |

$$(7) \text{ Distributive law } A+BC = (A+B)(A+C) \quad A(B+C) = AB + AC$$

$$(8) \text{ Absorption law } A(A+B) = A \quad A+A\bar{B} = A$$

$$(9) \text{ DeMorgan's law } \overline{AB} = \overline{A} + \overline{B} \quad \overline{A+B} = \overline{A} \cdot \overline{B}$$

Simplify the expressions:-

$$(1) Y(A, B, D) = (A+B)(A+B+D)\bar{D}$$

$$(2) Y(A, B, C) = AB + A\bar{B}(\bar{A}+\bar{C})$$

$$\begin{aligned} (1) Y(A, B, D) &= (\bar{A}+B)(A+B+D)\bar{D} \\ &= (\bar{A}+B)(A\bar{D}+B\bar{D}+D\bar{D}) \quad \{ \because D\bar{D} = 0 \} \\ &= (\bar{A}+B)(A\bar{D}+B\bar{D}+0) \\ &= (\bar{A}+B)(A\bar{D}+B\bar{D}) \\ &= (\bar{A}+B)(A+B)\bar{D} \\ &= (AA + \bar{A}B + BA + B \cdot B)\bar{D} \\ &= (0 + \bar{A}B + BA + B)\bar{D} \\ &= (\bar{A} + A + 1)B\bar{D} \\ &= (1 + 1)B\bar{D} \\ &= (1)B\bar{D} \Rightarrow B\bar{D} \end{aligned}$$

$$(2) Y(A, B, C) = AB + A\bar{B}(\bar{A}+\bar{C})$$

$$\begin{aligned} &= AB + A\bar{A}\bar{B} + A\bar{B}\bar{C} \\ &= AB + 0 + A\bar{B}\bar{C} \\ &= A(B + \bar{B}\bar{C}) \quad \{ \text{by DeMorgan's law} \} \\ &= A(B + \bar{B} + \bar{C}) \\ &= A(1 + \bar{C}) \\ &= A \end{aligned}$$

$$A+B = B+A$$

$$(A+B)+C = A+(B+C)$$

Reduce the expression:-

$$(1) \bar{A}\bar{C} + ABC + \bar{A}C$$

$$\bar{A}[C + \bar{C}] + ABC$$

$$\bar{A}(1) + ABC = \bar{A} + ABC$$

$$(2) f = (\overline{A+B\bar{C}})(A\bar{B} + ABC)$$

$$f = (\overline{A+\bar{B}C})(A\bar{B} + ABC)$$

$$= (\bar{A} \cdot \bar{B}C)(A\bar{B} + ABC)$$

$$= (\bar{A} \cdot BC)(A\bar{B} + ABC)$$

$$= \bar{A} \cdot A \cdot B \cdot \bar{B}C + \bar{A}BC \cdot ABC$$

$$= \bar{A}A \cdot \bar{B}C(\bar{B} + 1) = 0.$$

$$(3) f = (B+BC)(B+\bar{B}C)(B+D)$$

$$f = (B+BC)(B+\bar{B}C)(B+D)$$

$$= B(B+\bar{B}C)(B+D)$$

$$= (BB+B\cdot\bar{B}C)(B+D)$$

$$= (B+0C)(B+D)$$

$$= (B+0)(B+\bar{B}D)$$

$$= (B \cdot B + B \cdot \bar{B}D) \Rightarrow B+0D \Rightarrow B$$

Duality:-

A Dual of Boolean function is obtained by interchanging the logical and operator with logical or operator and interchanging '0' with '1's.

Given expression

$$J = 0$$

$$(1) \bar{0} = 1$$

$$1+0 = 1$$

$$(2) 0 \cdot 1 = 0$$

$$1+1 = 1$$

$$(3) 0 \cdot 0 = 0$$

$$0+0 = 0$$

$$(4) 1 \cdot 1 = 1$$

$$A+A = A$$

$$(5) A \cdot 0 = 0$$

$$A+0 = A$$

$$(6) A \cdot 1 = A$$

$$A+A = A$$

$$(7) A \cdot A = A$$

$$A+\bar{A} = 1$$

$$(8) A \cdot \bar{A} = 0$$

$$A+B = B+A$$

$$(9) A \cdot B = BA$$

$$A+(B+C) = (A+B)+C$$

$$(10) A \cdot (B \cdot C) = (A \cdot B) \cdot C$$

$$A+B C = (A+B)(A+C)$$

$$(11) A(B+C) = AB+AC$$

$$A+AB = A$$

$$(12) A(A+B) = A$$

$$A+A+B = A+B$$

$$(13) A(A \cdot B) = AB$$

$$\overline{A+B} = \overline{A}\bar{B}$$

$$(14) \overline{AB} = \bar{A} + \bar{B}$$

Reduce the expression,

~~$$A(B+\bar{C})(\overline{AB}+\bar{A}\bar{C})$$~~

$$(1) f = A[B+\bar{C}(\overline{AB}+\bar{A}\bar{C})]$$

$$(2) f = (B+BC)(B+\bar{B}C)(B+D)$$

$$(3) \text{ Show that } A\bar{B}C + B + B\bar{D} + AB\bar{D} + \bar{A}C = B+C.$$

$$(4) f = A[B+\bar{C}(\overline{AB}+\bar{A}\bar{C})]$$

$$= A[B+(\bar{A}+\bar{B})\cdot\bar{C}(\bar{A}+\bar{C})]$$

$$= A[B+(\bar{A}+\bar{B})\cdot(\bar{C}\bar{A}+0)]$$

$$= A[B+(\bar{A}+\bar{B})(\bar{C}\bar{A})]$$

$$= A[B+(\bar{A}+\bar{B})(\bar{C}+A)]$$

$$\Rightarrow [AB+(0)(0)] \rightarrow 0$$

$$\begin{aligned}
 (2) f &= (B+BC)(B+\bar{B}C)(B+D) \\
 &= (B+BC) [BB + BD + \bar{B}C \cdot B + \bar{B}C \cdot D] \\
 &= (B+BC) [B + BD + 0 + \bar{B}C \cdot D] \\
 &= (B+BC) [B + BD + \bar{B}CD] \\
 &= BB + BBD + B\bar{B}CD + BBC + BCBD + BCB \cancel{CB} \\
 &= [B+BD] + [\cancel{BC} + \bar{B}CD] \\
 &= B + BC \\
 &= B(1+C) \Rightarrow B(1) \Rightarrow \underline{\underline{B}}
 \end{aligned}$$

Duality:

- (1) $\bar{A}B + \bar{A}\bar{B}\bar{C} + A\bar{B}CD + A\bar{B}\bar{C}\bar{D}$ \Rightarrow if eqn is SOP form it be NAND form
- (2) $\bar{X}YZ + \bar{X}Y\bar{Z} + X\bar{Y}\bar{Z} + XY\bar{Z}$ \Rightarrow if eqn is POS form it be NOR form
- (3) $AB + \bar{A}C + A\bar{B}C$

Implementation by using NAND Gate:-

Boolean functions:-
Boolean functions are also known as logic functions.

⇒ Boolean functions are constructed by connecting the Boolean constants & variables with Boolean operations.

⇒ The boolean expressions are used to describe the boolean functions.
Ex: If the boolean expression $\bar{A}B + \bar{B}C$ is used to describe the function 'f'. Then boolean function can be written as based on structure of Boolean expression. It can be categorized into 2 forms
(1) Sum of product (SOP/normal form)
(2) Product of sum (POS)

(1) Sum of product:- This is a group of product terms and summed together.

Ex:- $f(A, B, C) = \overbrace{\bar{A}B + \bar{B}C}^{\text{product terms}} \quad \downarrow \quad \downarrow \quad \downarrow \quad \text{variable/literals.}$

⇒ It is also called as disjunctive normal form.

(2) Product of sum:- It is a group of sum terms and multiplied together.

Ex:- $f(A, B, C) = (\bar{A} + \bar{B})(A + B)$
 $\qquad\qquad\qquad \underbrace{\quad\quad}_{\text{sum of terms}}$

⇒ It is also called as conjunctive normal form.
 canonical
 Canonical form:-

These are the spl case of SOP & POS forms.

⇒ These are also known as standard SOP form & POS form.

⇒ Standard SOP form is also called as disjunctive canonical form.

⇒ Each product term contains all the variables of the function.

Let us consider a boolean function $f(A, B, C) =$

$$(\bar{A}B + \bar{B}C)$$

$$AB(C+\bar{C}) + \bar{B}C(A+\bar{A})$$

$$\left[\begin{array}{l} A=1 \\ \bar{A}=0 \end{array} \right]$$

$$\bar{A}BC + \bar{A}B\bar{C} + \bar{B}CA + \bar{B}C\bar{A}$$

| <u>Decimal No</u> | <u>A, B, C</u> | <u>min. terms</u> | <u>f(A, B, C)</u> |
|-------------------|----------------|---------------------------|-------------------|
| 0 | 0 0 0 | $\bar{A} \bar{B} \bar{C}$ | 0 |
| 1 | 0 0 1 | $A \bar{B} C$ | 1 |
| 2 | 0 1 0 | $\bar{A} B \bar{C}$ | 1 |
| 3 | 0 1 1 | $\bar{A} B C$ | 1 |
| 4 | 1 0 0 | $A \bar{B} \bar{C}$ | 0 |
| 5 | 1 0 1 | $A \bar{B} C$ | 1 |
| 6 | 1 1 0 | $A B \bar{C}$ | 0 |
| 7 | 1 1 1 | $A B C$ | 0 |

$$f(A, B, C) = m_1 + m_2 + m_3 + m_5$$

$$\{m(1, 2, 3, 5)\}$$

Standard POS form:-
⇒ It is also called as Conjunctive Canonical form.

⇒ Each sum term contains all the variables of the function.

Max terms:-

⇒ Each individual term in standard POS is called max terms represented with 'M'.

⇒ ~~MAX~~ $\pi M(0,1,\dots)$.

Convert the given expression in standard SOP form.

$$f(A,B,C) = \underbrace{A}_{\text{1}} + ABC$$

$$A(B+\bar{B})(C+\bar{C}) + ABC$$

$$\Rightarrow \underbrace{ABC}_{\text{1}} + \underbrace{A\bar{B}C}_{\text{2}} + \underbrace{AB\bar{C}}_{\text{3}} + \underbrace{A\bar{B}\bar{C}}_{\text{4}} + \underbrace{\bar{A}BC}_{\text{5}}$$

$$\Rightarrow \underbrace{ABC}_{\text{1}} + \underbrace{A\bar{B}C}_{\text{2}} + \underbrace{AB\bar{C}}_{\text{3}} + \underbrace{A\bar{B}\bar{C}}_{\text{4}}$$

$$\Rightarrow \pi M(4,5,6,7)$$

* Convert the given expression in standard SOP form. $f(A,B,C) = \underline{AB} + \underline{BC} + \underline{\bar{AC}}$

$$= AB(C+\bar{C}) + BC(A+\bar{A}) + \bar{A}C(B+\bar{B})$$

$$= ABC + A\bar{B}C + AB\bar{C} + \bar{A}BC + \bar{A}C\bar{B}$$

$$= ABC + AB\bar{C} + \bar{A}BC + \bar{A}C\bar{B}$$

Convert the given expression into standard POS form.

$$f(A, B, C) = (A+B)(B+C)(A+C)$$

$$\Rightarrow (A+B+C\bar{C})(B+C+A\bar{A})(A+C+B\bar{B})$$

$$\Rightarrow (A+B+C)(A+B+\bar{C})(B+C+A)(B+C+\bar{A})(A+C+B)(A+C+\bar{B})$$

$$= (A+B+C)(A+B+\bar{C})(B+C+\bar{A})(A+C+\underline{\bar{B}})$$

* Convert convert the given expression into stand. POS form. $f(A, B, C) = A(A+B+C)$

'Compliments of Boolean functions:-'

$$\Rightarrow (A+B)' = A' \cdot B'$$

$$(A \cdot B)' = A' + B'$$

$$\text{Ex:- } 1 : x'y'z + xy + yz'$$

$$= (x'y'z)' + (xy)' + (yz)'$$

$$= x' + y'' + z' + x'y' + y' + z''$$

$$= x' + y + z' + x' + y' + y' + z$$

$$2. ABC' + BC + A'B$$

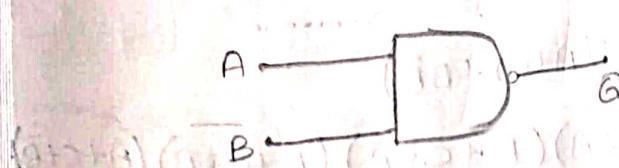
$$= (ABC')' + (BC)' + (A'B)'$$

$$= A' + B' + C'' + B' + C' + A'' + B'$$

$$= A' + B' + C + B' + C' + A + B'$$

NAND GATE implementation

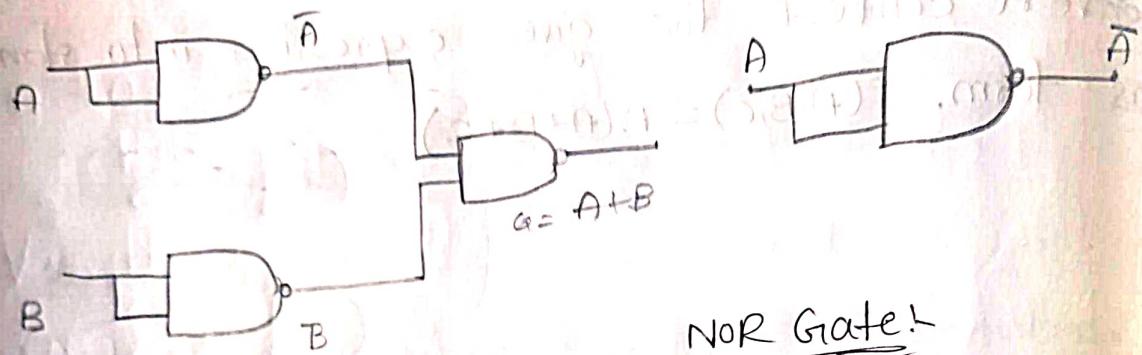
(UNAND Gate symbol!:-



(2) AND Gate:-

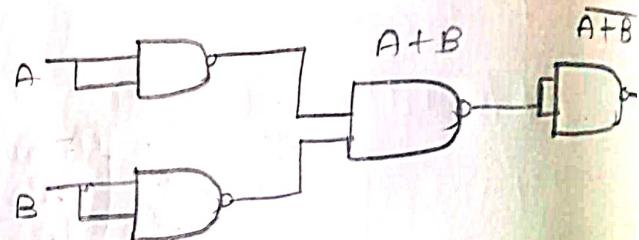
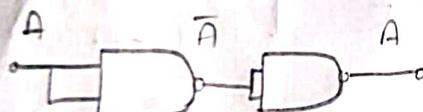


OR Gate:-

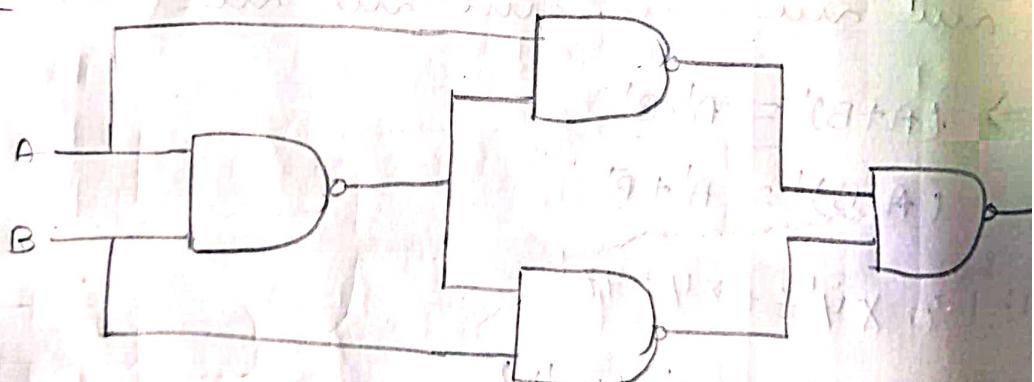


NOR Gate!:-

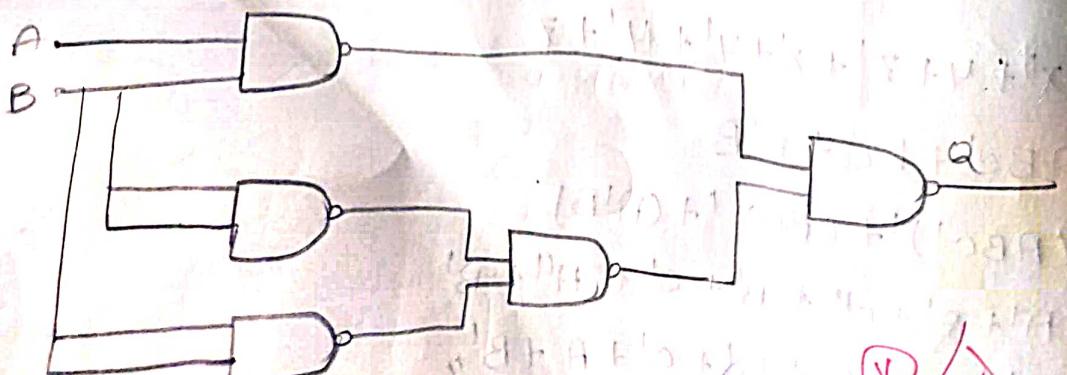
Buffer!:-



EX-OR!:-



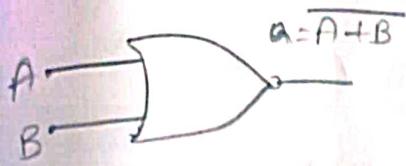
EX-NOR!:-



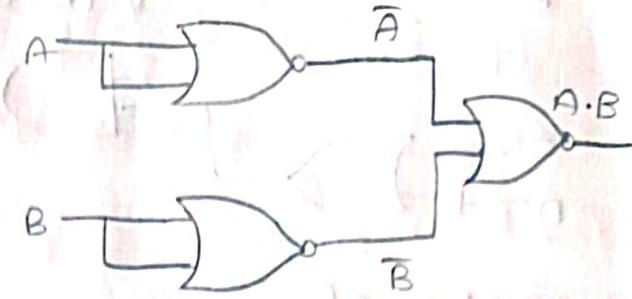
Q/5

NOR Gate implementation by using gates:-

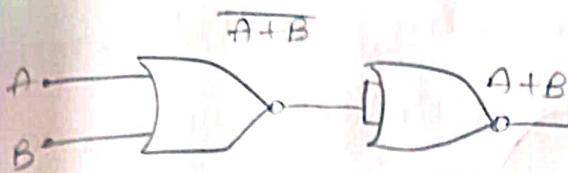
NOR Gate:-



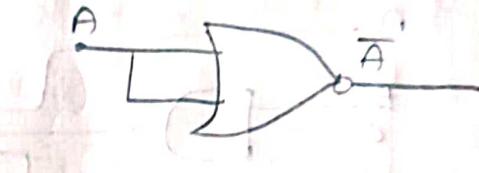
AND Gate:-



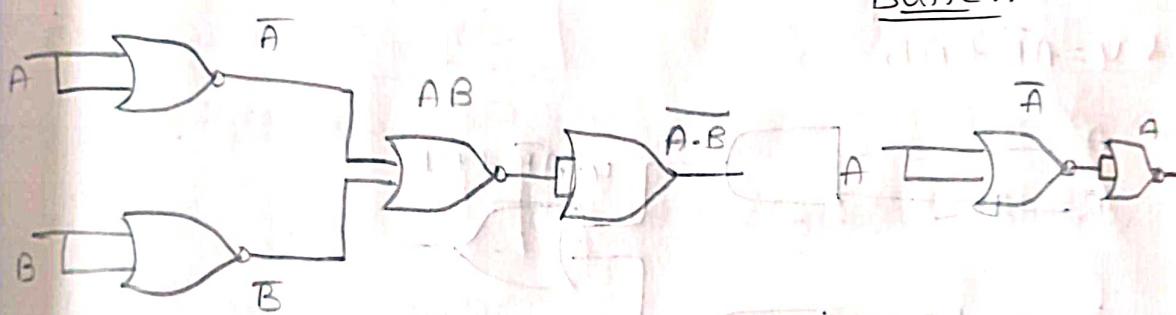
OR Gate:-



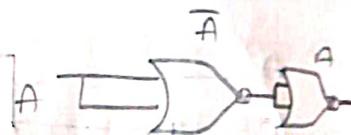
NOT Gate:-



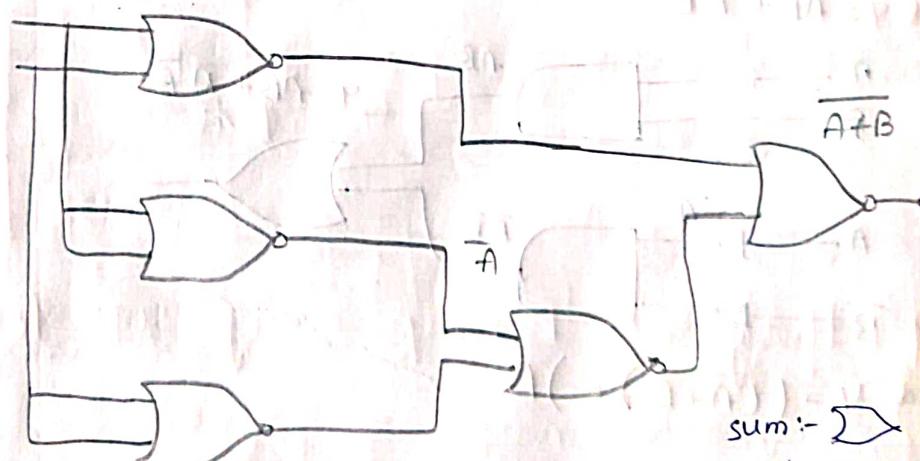
NAND Gate:-



Buffer:-

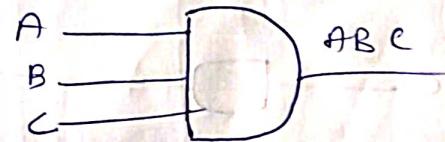
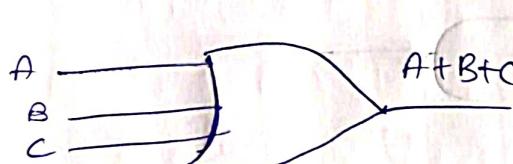


EX-OR Gate:-

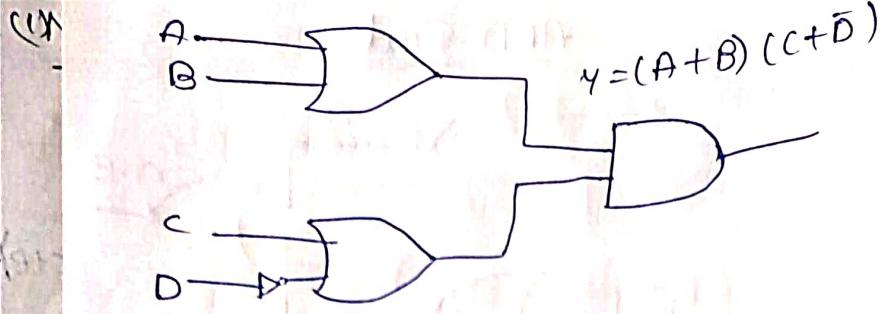


$$\text{sum: } \Sigma$$
$$\text{product: } P$$

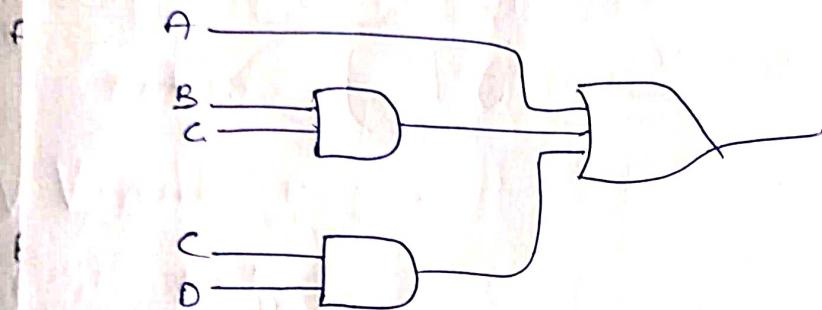
$$* Y = A + B + C$$



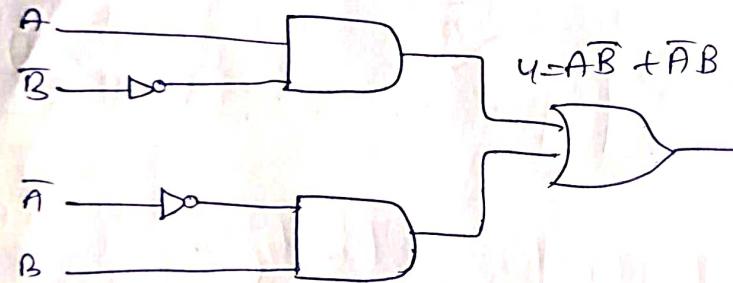
$$Y = (A+B)(C+D)$$



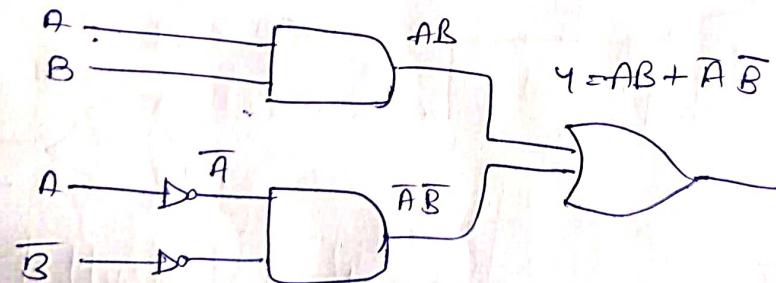
$$* A + BC + CD$$



$$* Y = A\bar{B} + \bar{A}B$$



$$* Y = AB + \bar{A}\bar{B}$$



$$* Y = ((\bar{A}+B)C)D$$

