

PROJECT – 1
Comparison-based Sorting Algorithms

UNIVERSITY OF NORTH CAROLINA AT CHARLOTTE

DEPARTMENT OF INFORMATION TECHNOLOGY

ITCS-6114

Algorithms and Data Structures

Submitted By

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1. INSERTION SORT:

Insertion sort is a simple sorting algorithm that works similar to the way you sort playing cards in your hands. The array is virtually split into a sorted and an unsorted part. Values from the unsorted part are picked and placed at the correct position in the sorted part.

Algorithm:

To sort an array of size n using Insertion Sort:

- 1: Iterate from $\text{array}[0]$ to $\text{array}[n]$ over the array.
- 2: Compare the current element to its most recent left element.
- 3: If the current element is smaller than its immediate left element then compare it with the elements from left to 0. Move the greater elements one position to the right.

Time Complexity: $O(n^2)$ for large arrays and $O(n)$ for sorted arrays.

Boundary Cases: Insertion sort takes maximum time to sort if elements are sorted in reverse order. And it takes minimum time $O(n)$ when elements are already sorted.

Uses: Insertion sort is used when number of elements is small. It can also be useful when input array is almost sorted, only few elements are misplaced in complete big array.

CODE:

```
def insertionsort(array):  
    # Loop from arr[1] to arr[n]  
    for i in range(1, len(array)):  
        current_element = array[i]  
        # Now we iterate backwards from i to 0 and compare the values with the current element  
        j = i-1  
        while j >= 0 and current_element < array[j]:  
            array[j + 1] = array[j]  
            j -= 1  
        array[j + 1] = current_element  
    return array
```

2. MERGE SORT

Merge sort is the sorting technique which based on the divide and conquer technique. Merge sort first divides the array into equal halves. Each sub array of size $n/2$ is sorted in a recursive manner in $O(\log(n))$ time and then combines them in a sorted manner.

Algorithm:

- 1: Find the middle point to divide the array into two halves
- 2: Call merge Sort function for first half: Call merge Sort (array, 1, m)

3: Call merge Sort for second half: Call merge Sort (array, m+1, r)

4: Merge two halves sorted in step 2 and 3: Call merge (array, l, m, r)

Time Complexity: $O(n \cdot \log(n))$ for best, average and worst cases.

Uses: Merge Sort is useful for sorting linked lists in $O(n \cdot \log(n))$ time.

CODE:

```
def mergeSort(array):
    if len(array)>1:
        mid_pos = len(array)//2
        left = array[:mid_pos]
        right = array[mid_pos:]

        sorted_left = mergeSort(left)
        sorted_right = mergeSort(right)
        array = []
        while len(sorted_left)>0 and len(sorted_right)>0:
            if sorted_left[0] < sorted_right[0]:
                array.append(sorted_left[0])
                sorted_left.pop(0)
            else:
                array.append(sorted_right[0])
                sorted_right.pop(0)
        for i in sorted_left:
            array.append(i)
        for i in sorted_right:
            array.append(i)
        #print("Time taken to execute Merge Sort is : ",datetime.now()-starttime,"\n")
    return array
```

3. HEAP SORT

Heap sort is a comparison-based sorting technique whose structure is of Tree type. The tree is based on Binary heap data structure. This sorting technique is similar to selection sort where I first find the maximum element and place the maximum element at the end. I repeat the heapify process for remaining elements.

Algorithm:

- 1: Build a max heap from the input data.
- 2: At this point, the largest item is stored at the root of the heap. Replace it with the last item of the heap followed by reducing the size of heap by 1. Finally, heapify the root of the tree.
- 3: Repeat step 2 while size of heap is greater than 1.

Time Complexity: The Running time for heapify operation takes $O(\log(n))$ which is the height of the binary tree. After that, a sorted array is created by repeatedly removing the largest element from the heap and adding the element to the array. In this way, Heapsort takes $O(n \cdot \log(n))$ time.

Uses: Priority Queues, Order Statistics

CODE:

```
def heapify(array, length, index):
    max_element = index
    left = 2 * index + 1
    right = 2 * index + 2
    # See if left child of root exists and is greater than root
    if left < length and array[index] < array[left]:
        max_element = left
    # See if right child of root exists and is greater than root
    if right < length and array[max_element] < array[right]:
        max_element = right
    if max_element != index:
        array[index], array[max_element] = array[max_element], array[index]
        heapify(array, length, max_element)

def heapSort(array):
    length = len(array)
    for i in range(length//2 - 1, -1, -1):
        heapify(array, length, i)
    for i in range(length-1, 0, -1):
        array[i], array[0] = array[0], array[i]
        heapify(array, i, 0)

    return array
```

4. QUICK SORT

Quicksort is a Divide and Conquer algorithm. It picks an element as pivot and partitions the given array around the picked pivot. There are many different versions of quicksort that pick pivot in different ways such as Always picking first element as pivot, always pick last element as pivot, Pick a random element as pivot, Pick median as pivot. The key process in quicksort is partition (). Target of partitions is, given an array and an element x of array as pivot, put x at its correct position in sorted array and put all smaller elements (smaller than x) before x, and put all greater elements (greater than x) after x. All this should be done in linear time.

Algorithm:

/* low --> Starting index, high --> Ending index */

quicksort(array[], low, high){if (low < high) {/* pi is partitioning index, array[pi] is now at right place */ pi = partition(array, low, high); quicksort(array, low, pi - 1); // Before pi quicksort (array, pi + 1, high); // After pi}}

Time Complexity: $O(n^2)$ is the worst case and $O(N \log(n))$ is the average case complexity.

Uses: Quick Sort is a cache friendly sorting algorithm as it has good locality of reference when used for arrays.

CODE:

```
def partition(array, first, last):
    if last - first > 0:
        pivot, left, right = array[first], first, last
        while left <= right:
            while array[left] < pivot:
                left += 1
            while array[right] > pivot:
                right -= 1
            if left <= right:
                array[left], array[right] = array[right], array[left]
                left += 1
                right -= 1
        partition(array, first, right)
        partition(array, left, last)

def quicksort(array):
    partition(array, 0, len(array) - 1)
    return array
```

5. MODIFIED QUICK SORT

I used median-of-three as pivot in Modified Quick sort. Which means, in this algorithm, I can pick the pivot by taking the median of left most, middle and the right most elements from the array. Compare three elements and swap these elements if necessary. If the input size is less than or equal to 15, Insertion sort algorithm is used.

CODE:

```

def MedianOfThree(arr, left, right):
    mid = (left + right)//2
    if arr[right] < arr[left]:
        array[left],array[right] = array[right],array[left]
    if arr[mid] < arr[left]:
        array[left],array[mid] = array[mid],array[left]
    if arr[right] < arr[mid]:
        array[right],array[mid] = array[mid],array[right]
    return arr[mid]

def modifiedpartition(array, first, last):
    if last - first > 0:
        left, right = first, last
        pivot = MedianOfThree(array,first,last)
        while left <= right:
            while array[left] < pivot:
                left += 1
            while array[right] > pivot:
                right -= 1
            if left <= right:
                array[left], array[right] = array[right], array[left]
                left += 1
                right -= 1
        modifiedpartition(array, first, right)
        modifiedpartition(array, left, last)

def modified_quicksort(array):
    if len(array) <= 15:
        for i in range(1, len(array)):
            current_element = array[i]
            j = i-1
            while j >= 0 and current_element < array[j]:
                array[j + 1] = array[j]
                j -= 1
            array[j + 1] = current_element
    else:
        modifiedpartition(array, 0, len(array) - 1)

    return array

```

EXECUTION CODE:

Function to generate array of random numbers for a given range “n”.

```
from time import time
import random

def random_number_generator(n):
    array = set()
    while len(array) < n:
        array.add(random.randint(1,n))
    array = list(array)
    random.shuffle(array)
    return array
```

Executing for shuffled Input:

```

for i in input_size_of_array:
    n = i
    array = random_number_generator(n)

    starttime1 = time()
    res = quicksort(array)
    executiontime1 = (time()-starttime1)
    quicksorttime[n] = executiontime1

    random.shuffle(array)
    starttime2 = time()
    res = insertionsort(array)
    executiontime2 = (time()-starttime2)
    insertionsorttime[n] = executiontime2

    random.shuffle(array)
    starttime3 = time()
    res = mergeSort(array)
    executiontime3 = (time()-starttime3)
    mergesorttime[n] = executiontime3

    random.shuffle(array)
    starttime4 = time()
    res = heapSort(array)
    executiontime4 = (time()-starttime4)
    heapsorttime[n] = executiontime4

    random.shuffle(array)
    starttime5 = time()
    res = modified_quicksort(array)
    executiontime5 = (time()-starttime5)
    modifiedquicksorttime[n] = executiontime5

print("Insertion Sort runtimes = \n", list(insertionsorttime.values()), "\n")
print("Heap Sort runtimes = \n", list(heapsorttime.values()), "\n")
print("Merge Sort runtimes = \n", list(mergesorttime.values()), "\n")
print("Quick Sort runtimes = \n", list(quicksorttime.values()), "\n")
print("Modified Quick Sort runtimes = \n", list(modifiedquicksorttime.values()))

```

Executing for Sorted input:


```

input_sorted_array = [10,50,100,500,1000,1500,2000,2500,3000]
for i in input_sorted_array:
    n = i
    array = random_number_generator(n)
    array.sort()

    starttime1 = time()
    res = quicksort(array)
    executiontime1 = (time()-starttime1)
    sortedquicksorttime[n] = executiontime1

    starttime2 = time()
    res = insertionsort(array)
    executiontime2 = (time()-starttime2)
    sortedinsertionsorttime[n] = executiontime2

    starttime3 = time()
    res = mergeSort(array)
    executiontime3 = (time()-starttime3)
    sortedmergesorttime[n] = executiontime3

    starttime4 = time()
    res = heapSort(array)
    executiontime4 = (time()-starttime4)
    sortedheapsorttime[n] = executiontime4

    starttime5 = time()
    res = modified_quicksort(array)
    executiontime5 = (time()-starttime5)
    sortedmodifiedquicksorttime[n] = executiontime5

```

Executing for Reversely Sorted input:

```

input_sorted_array = [10,50,100,500,1000,1500,2000,2500,3000]
for i in input_sorted_array:
    n= i
    array = random_number_generator(n)
    array.sort(reverse = True)

    starttime1 = time()
    res = quicksort(array)
    executiontime1 = (time()-starttime1)
    reversesortedquicksorttime[n] = executiontime1

    starttime2 = time()
    res = insertionsort(array)
    executiontime2 = (time()-starttime2)
    reversesortedinsertionsorttime[n] = executiontime2

    starttime3 = time()
    res = mergeSort(array)
    executiontime3 = (time()-starttime3)
    reversesortedmergesorttime[n] = executiontime3

    starttime4 = time()
    res = heapSort(array)
    executiontime4 = (time()-starttime4)
    reversesortedheapsorttime[n] = executiontime4

    starttime5 = time()
    res = modified_quicksort(array)
    executiontime5 = (time()-starttime5)
    reversesortedmodifiedquicksorttime[n] = executiontime5

```

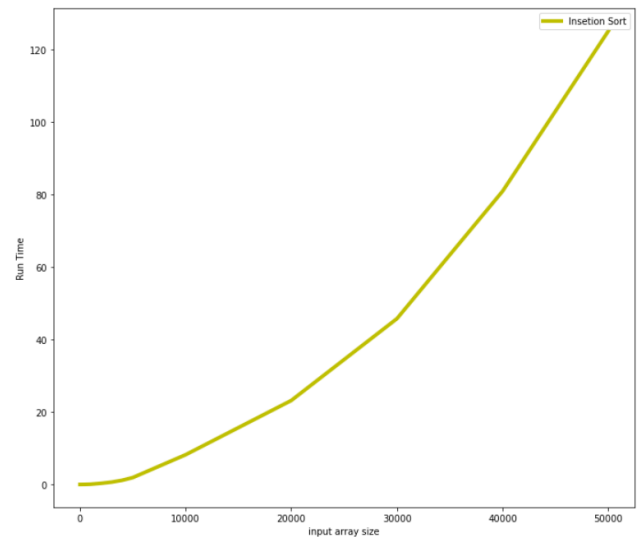
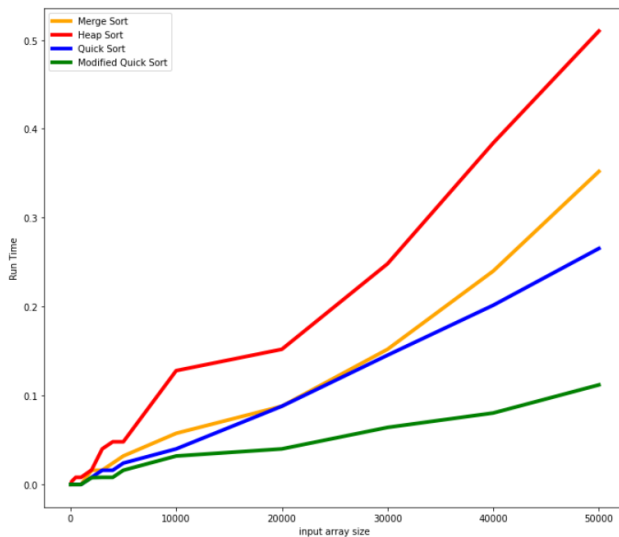
ANALYSIS FOR SHUFFLED INPUT: (Time in Seconds)

Len(array)	Insertion Sort	Merge Sort	Heap Sort	Quick Sort	Modified Quick Sort
1000	0.08133721351623535	0.008001327514648438	0.008000612258911133	0.0030663013458251953	0.008000612258911133
2000	0.2960543632507324	0.01600170135498047	0.01598334312438965	0.008000373840332031	0.01602005958557129

3000	0.75160455703 73535	0.02401876 449584961	0.0399858 951568603 5	0.0079965591 43066406	0.00798368453979 4922
4000	1.25752973556 51855	0.03200364 112854004	0.0399854 183197021 5	0.0159995555 87768555	0.01600217819213 8672
5000	1.85445833206 17676	0.04802370 071411133	0.0479865 074157714 84	0.0159828662 87231445	0.0160098609924 3164
10000	6.50246453285 2173	0.06249356 269836426	0.0781359 672546386 7	0.0480051040 64941406	0.01561379432678 2227
20000	18.0062973499 2981	0.14401602 745056152	0.1680183 410644531 2	0.0468790531 1584473	0.04800534248352 051
30000	40.6528863906 86035	0.24803924 560546875	0.2480275 630950927 7	0.0937485694 8852539	0.08000850677490 234
40000	79.3864862918 8538	0.41416168 212890625	0.3279118 537902832	0.1280131340 0268555	0.16954922676086 426
50000	136.041074752 80762	0.54304718 97125244	0.4724144 93560791	0.1600291728 9733887	0.15645718574523 926

GRAPH:

Shuffled inputs



Observations:

- The quick sort algorithm running time depends on the selection of pivot in the given random data. If pivot is around the median value, less time is taken, or else algorithm takes more time.
- In our case modified quick sort took more time as the size of the input increases.
- Time taken by Insertion Sort also increased as the input size increases.
- Time taken by Merge Sort and heap sort is less if the input is random data.

ANALYSIS FOR SORTED INPUT: (Time in Seconds)

```
print("Insertion Sort runtimes =", list(sortedinsertionsorttime.values()), "\n")
print("Heap Sort runtimes = \n", list(sortedheapsorttime.values()), "\n")
print("Merge Sort runtimes = \n", list(sortedmergesorttime.values()), "\n")
print("Quick Sort runtimes = \n", list(sortedquicksorttime.values()), "\n")
print("Modified Quick Sort runtimes = \n", list(sortedmodifiedquicksorttime.values()))
```

```
Insertion Sort runtimes =
[0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.007995128631591797, 0.0]
```

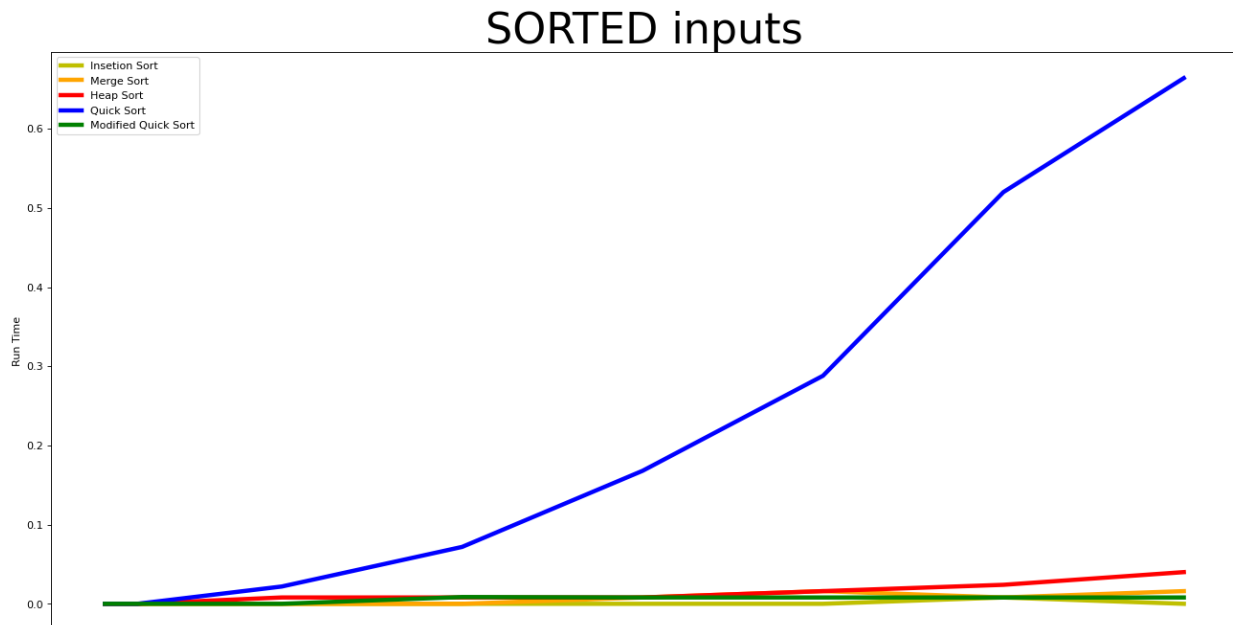
```
Heap Sort runtimes =
[0.0, 0.0, 0.0, 0.008006095886230469, 0.008000850677490234, 0.00799897003173828, 0.016002416610717773, 0.023990869522094727, 0.0400240421295166]
```

```
Merge Sort runtimes =
[0.0, 0.0, 0.0, 0.0, 0.0, 0.008001089096069336, 0.01598525047302246, 0.008000612258911133, 0.015999794006347656]
```

```
Quick Sort runtimes =
[0.0, 0.0, 0.0, 0.021834850311279297, 0.07188296318054199, 0.1679990291595459, 0.28805065155029297, 0.520054817199707, 0.664074182510376]
```

```
Modified Quick Sort runtimes =
[0.0, 0.0, 0.0, 0.008541345596313477, 0.008001327514648438, 0.007997751235961914, 0.008001565933227539, 0.007982254028320312]
```

GRAPH:



ANALYSIS FOR REVERSLY SORTED INPUT: (Time in Seconds)

```

print("Insertion Sort runtimes =\n",list(reversesortedinsertionsorttime.values()),"\n")
print("Heap Sort runtimes = \n", list(reversesortedheapsorttime.values()),"\n")
print("Merge Sort runtimes = \n", list(reversesortedmergesorttime.values()),"\n")
print("Quick Sort runtimes = \n", list(reversesortedquicksorttime.values()),"\n")
print("Modified Quick Sort runtimes = \n", list(reversesortedmodifiedquicksorttime.values()))

Insertion Sort runtimes =
[0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0]

Heap Sort runtimes =
[0.0, 0.0, 0.0, 0.008009195327758789, 0.008199214935302734, 0.016001224517822266, 0.016002893447875977, 0.024003982543945312, 0.039994001388549805]

Merge Sort runtimes =
[0.0, 0.0, 0.0, 0.008006811141967773, 0.008000373840332031, 0.0, 0.015984535217285156, 0.008001089096069336, 0.016013383865356445]

Quick Sort runtimes =
[0.0, 0.0, 0.0, 0.013830423355102539, 0.07200312614440918, 0.16802167892456055, 0.2800486087799072, 0.4480552673339844, 0.7295355796813965]

Modified Quick Sort runtimes =
[0.0, 0.0, 0.0, 0.0, 0.0, 0.008000612258911133, 0.00800013542175293, 0.01135563850402832, 0.007999897003173828]

```

GRAPH:

REVERSELY SORTED inputs

