

**Problem 1** (6%). Consider the Turán's theorem and its proof, and let  $n$  be a multiple of  $k$ , where  $k \geq 2$  is an integer.

Construct a graph  $G = (V, E)$  with  $n$  vertices that contains no  $(k + 1)$ -clique such that the number of edges attains the upper-bound given in the Turán's theorem, i.e.,

$$|E| = \left(1 - \frac{1}{k}\right) \cdot \frac{n^2}{2}.$$

Justify your answer.

**Problem 2** (7%). Let  $k > 0$  be an integer and let  $p(n)$  be a function of  $n$  with  $p(n) = \Omega((6k \ln n)/n)$  for large  $n$ . Prove that "almost surely" the random graph  $G = G(n, p)$  has no independent set of size  $n/2k$ , i.e., show that

$$\Pr \left[ \alpha(G) \geq \frac{n}{2k} \right] = o(1).$$

**Problem 3** (7%). Let  $K_n$  denote the complete graph with  $n$  vertices. Show that it is possible to color  $K_n$  with at most  $3\sqrt{n}$  colors so that there are no monochromatic triangles.

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Justify your answer.

1. Set a graph  $G = (V, E)$ , with  $n$  vertices and contains no  $(k+1)$ -clique and a maximal # of edges
2. Since above, when we add a new edge to  $G$  will create a  $(k+1)$ -clique
3.  $G$  contains at least one  $k$ -clique, and we let  $A$  be a  $k$ -clique in  $G$ , and let  $B := V \setminus A$ ;  $e_A, e_B, e_{AB}$  denotes edge # in  $A$ , in  $B$ , and between  $A$  and  $B$  respectively.
4. Hence,  $e_A = \binom{k}{2} = \frac{k \cdot k-1}{2}$
5. By induction hypothesis,  $e_B \leq \left(1 - \frac{1}{k}\right) \cdot \frac{(n-k)^2}{2}$
6. Each  $v \in B$  is adjacent to at most  $k-1$  vertices in  $A$ .  
Hence,  $e_{AB} \leq (k-1) \cdot (n-k)$
7. Hence,  $|E| = e_A + e_B + e_{AB} \leq \frac{k^2 - k}{2} + \left(1 - \frac{1}{k}\right) \cdot \frac{(n-k)^2}{2} + (k-1) \cdot (n-k)$ 

$$\begin{aligned} & \frac{k^2 - k}{2} + \frac{\left(n^2 - 2nk + k^2 - \frac{n^2}{k} + 2n - k\right)}{2} + (nk - k^2 - n + k) \\ &= \frac{k^2 - k + n^2 - 2nk + k^2 - \frac{n^2}{k} + 2n - k + nk - k^2 - n + k}{2} \\ &= \left(1 - \frac{1}{k}\right) \cdot \frac{n^2}{2} \# \end{aligned}$$

**Problem 2** (7%). Let  $k > 0$  be an integer and let  $p(n)$  be a function of  $n$  with  $p(n) = \Omega((6k \ln n)/n)$  for large  $n$ . Prove that "almost surely" the random graph  $G = G(n, p)$  has no independent set of size  $n/2k$ , i.e., show that

$$\Pr \left[ \alpha(G) \geq \frac{n}{2k} \right] = o(1).$$

1. Let  $r$  denotes independent set of size at least  $n/2k$

$$\text{i.e. } r = \lfloor \frac{n}{2k} \rfloor$$

2. the probability of choose any subset  $S$  of size  $r$  ( $S \subseteq V$ ), that  $S$  is independent set is

$$\Pr [S \text{ is independent set}] = (1-p)^{\binom{r}{2}}$$

$$\Rightarrow \Pr [S \text{ is independent set}, |S|=r] \leq \binom{n}{r} (1-p)^{\binom{r}{2}}$$

3. 分別對  $\binom{n}{r}$  和  $(1-p)^{\binom{r}{2}}$  做估計的 upper bound

①  $\binom{n}{r}$ :

$$\binom{n}{r} = \frac{n!}{r!(n-r)!} = \frac{n \times (n-1) \times \dots \times (n-r+1)}{r!} \leq \frac{n^r}{r!}, \quad \text{又 } r! \geq \left(\frac{r}{e}\right)^r$$

$$\Rightarrow \binom{n}{r} \leq \frac{n^r}{(r/e)^r} = \left(\frac{en}{r}\right)^r = e^{r \ln(en/r)}$$

②  $(1-p)^{\binom{r}{2}}$ :

$$\text{由 } (1-p) \leq e^{-p}, \quad (1-p)^{\binom{r}{2}} \leq e^{-p \cdot \binom{r}{2}}$$

$$\text{又 } \binom{r}{2} \approx \frac{r^2}{2} = \frac{n^2}{2 \cdot 4k^2} = \frac{n^2}{8k^2}$$

$$P(n) \geq C \cdot \left(\frac{6k \cdot \ln n}{n}\right) \Rightarrow p \cdot \binom{r}{2} \geq C \cdot \left(\frac{6k \cdot \ln n}{n}\right) \cdot \frac{n^2}{8k^2} = \frac{C}{8k} n \ln n$$

$$\therefore \text{因此 } (1-p)^{\binom{r}{2}} \leq e^{\left(-\frac{C}{8k} \cdot n \ln n\right)}$$

$$4. \quad \binom{n}{r} (1-p)^{\binom{r}{2}} \leq e^{r \cdot \ln(\frac{en}{r})} \times e^{\left(-\frac{C}{8k} \cdot n \ln n\right)}$$

當  $n$  很大時  $e^{r \cdot \ln(\frac{en}{r})} \times e^{\left(-\frac{C}{8k} \cdot n \ln n\right)}$  趨近於 0

$$\Rightarrow \Pr \left[ \alpha(G) \geq \frac{n}{2k} \right] = o(1) \#$$

**Problem 3 (7%).** Let  $K_n$  denote the complete graph with  $n$  vertices. Show that it is possible to color  $K_n$  with at most  $3\sqrt{n}$  colors so that there are no monochromatic triangles.

1. For each edge, choose one of  $3\sqrt{n}$  colors uniformly at random to color the edges
2. The probability a given three vertices form a monochromatic triangles is  $\left(\frac{1}{3\sqrt{n}}\right)^2 = \frac{1}{9n}$
3. The # of triangles whose three edge colorings are not independent of a given triangle is  $3n-8$
4. By LLL, it holds that  $e \cdot \frac{1}{9n} (3n-7) \leq 1$  hence, there is a nonzero probability that there are no monochromatic triangles.