复制数据类型理论研究
(CCF 2018 第九届优博论坛)

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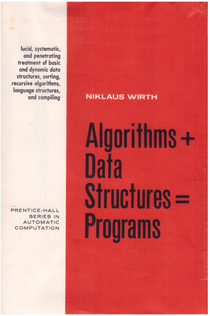
南京大学软件所

2018 年 08 月 08 日

Abstract Data Types [Liskov and Zilles, 1974]

(单线程; 顺序语义)



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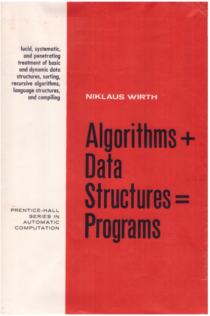
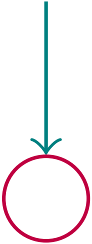
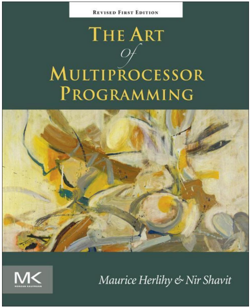




Concurrent Data Types [Herlihy and Wing, 1990]

(多线程; 并发语义)



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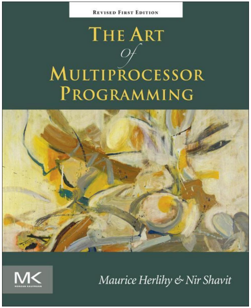

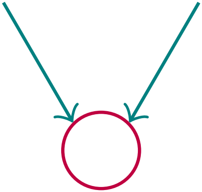

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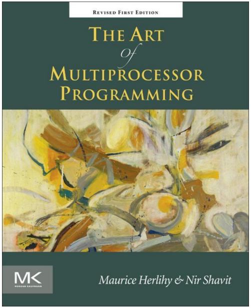

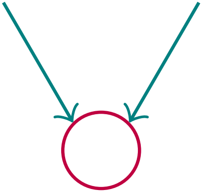

PL (Programming Language)

Replicated Data Types [Shapiro et al., 2011a] [Burckhardt et al., 2014]

(多副本; 复制语义)

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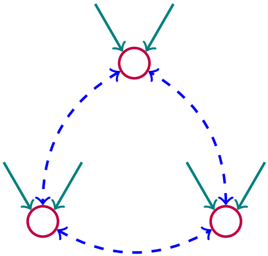

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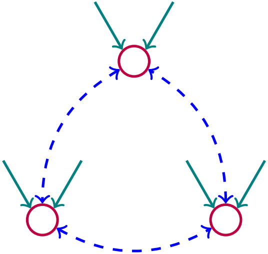

DC (Distributed Computing)

新平台：大规模分布式系统

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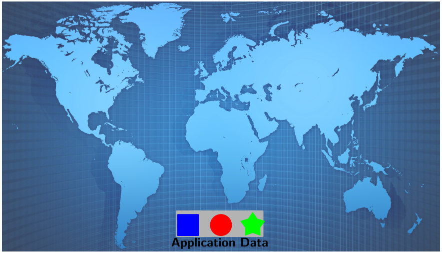
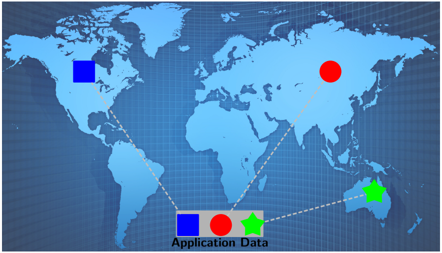
低延迟　高可用性 (4 个 9)　高容错性　高可扩展性

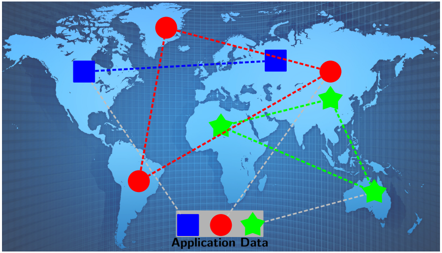


分布数据 (distributed data):



分布数据 (distributed data):

1. 分区 (partition): 水平扩展



分布数据 (distributed data):

1. 分区 (partition): 水平扩展
2. 副本 (replication): 就近访问, 容灾备份

复制数据类型 [Shapiro et al., 2011a]

- Read/Write Register
- Counter
- Set
- List
- HashMap
- Disjoint Set
- Graph
- ...

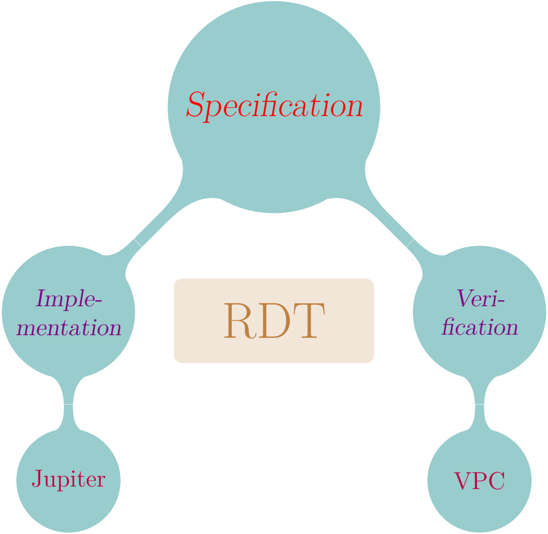
新问题，新挑战

新问题，新挑战

Replicated Data Types: Specification, Verification, Optimality

Sebastian Burckhardt　Alexey Gotsman　Hongseok Yang　Marek Zawirski

[Burckhardt et al., 2014]


Specification
Imple-
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RDT
Veri-
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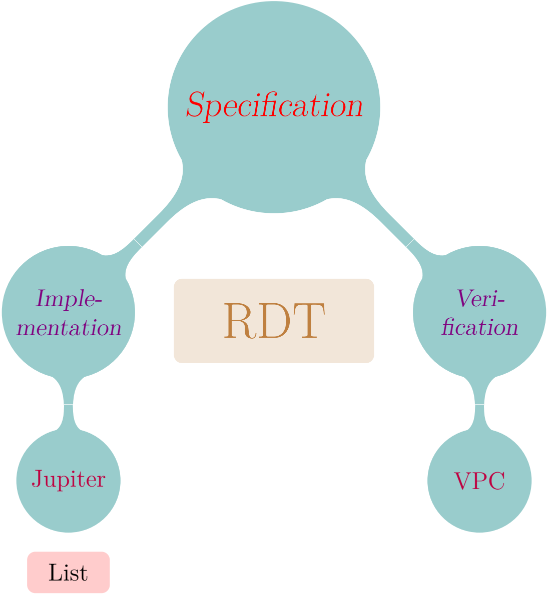
Jupiter
VPC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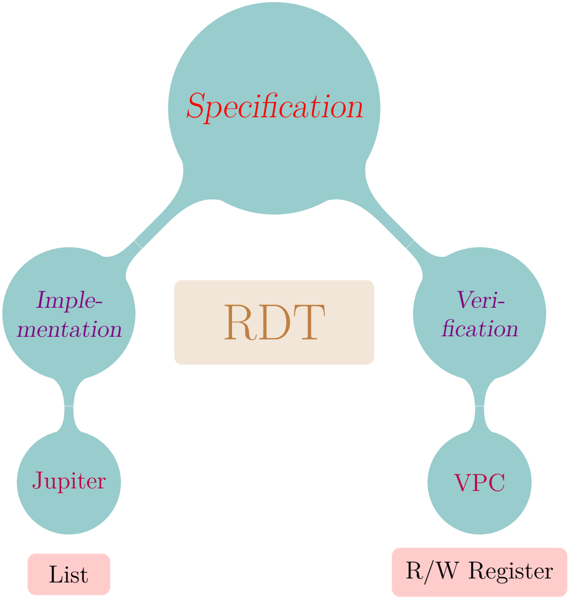
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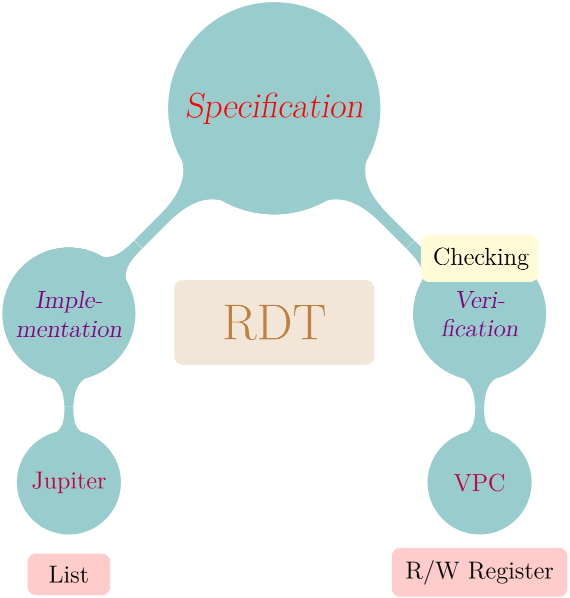
R/W Register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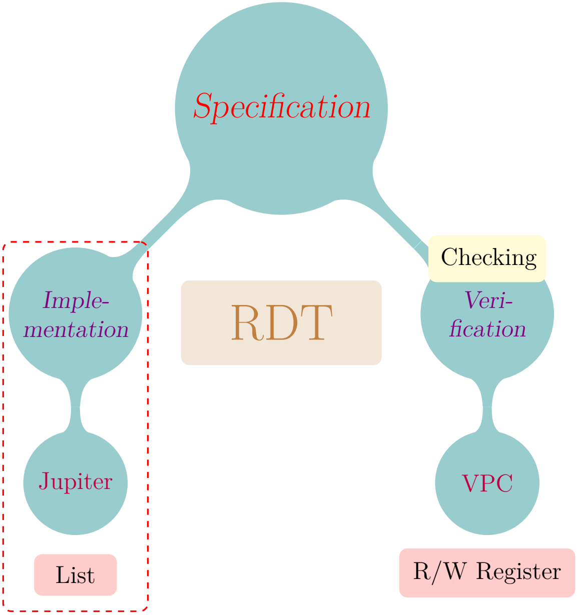
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Brief Announcement @ PODC'2018 [1]

实现复制列表的 Jupiter 协议 [Nichols et al., 1995][a] 满足 weak list specification [Attiya et al., 2016][b].

[a] David A. Nichols et al. (1995). "High-latency, Low-bandwidth Windowing in the Jupiter Collaboration System". In: *Proceedings of the 8th Annual ACM Symposium on User Interface and Software Technology*. UIST '95. ACM, pp. 111–120.

[b] Hagit Attiya et al. (2016). "Specification and complexity of collaborative text editing". In: *Proceedings of the 2016 ACM Symposium on Principles of Distributed Computing*. PODC '16. ACM, pp. 259–268.

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Weak List Specification

基于副本的协同文本编辑系统

(a) Google Docs

(b) Apache Wave

(c) Wikipedia

(d) LaTeX Editor

复制列表对象：建模编辑系统的核心功能

$\textsc{Ins}(a, p)$: 在 p 位置插入元素 a

$\textsc{Del}(p)$: 删除 p 位置上的元素

$\textsc{Read}$: 返回该列表

Specification and Complexity of Collaborative Text Editing

Hagit Attiya
Technion

Sebastian Burckhardt
Microsoft Research

Alexey Gotsman
IMDEA Software Institute

Adam Morrison
Technion

Hongseok Yang
University of Oxford

Marek Zawirski*
Inria & Sorbonne Universités,
UPMC Univ Paris 06, LIP6

定义 (Weak List Specification $\mathcal{A}_{\text{weak}}$ [Attiya et al., 2016])

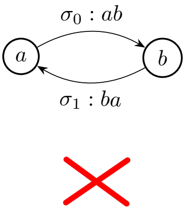
Informally, $\mathcal{A}_{weak}$ requires the ordering between elements that are not deleted to be consistent across the system.

定义在系统所有列表状态上的全局性质

定义（状态对兼容性 (Pairwise State Compatibility Property)）

任给两个列表状态 σ_0、σ_1，若它们含有两个共同元素 a、b，则 a、b 在 σ_0 与 σ_1 中的相对顺序保持一致。

$\boxed{\sigma_0 : ab}$ $\boxed{\sigma_1 : ba}$



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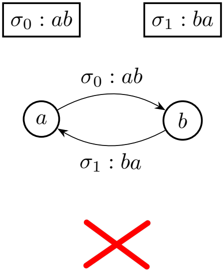
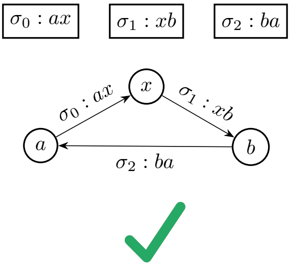
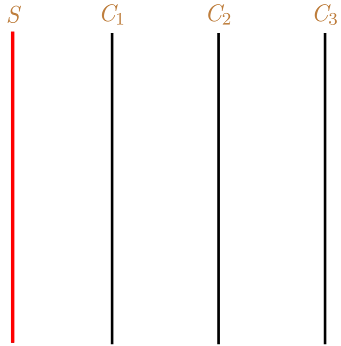




Jupiter

$(n+1)$ replicas $\triangleq$ (n) Client + (1) Server [Nichols et al., 1995]



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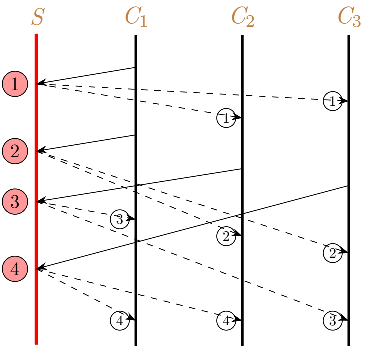

Server 负责将所有操作序列化

$$\text{用户操作} \xrightarrow[\boxed{\text{立即返回}}]{\text{提交}} \text{Client} \xrightarrow[\text{更新操作}]{\text{FIFO}} \text{Server} \xrightarrow[\text{更新操作}]{\text{FIFO}} \text{其它 Clients}$$

操作转换 (Operational Transformation; OT) [Ellis and Gibbs, 1989]

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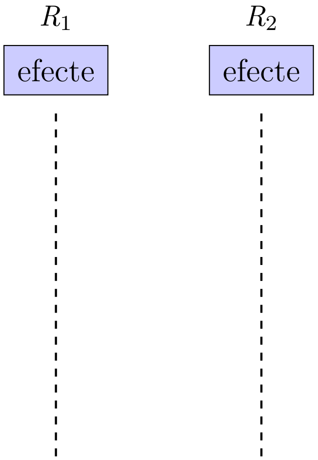

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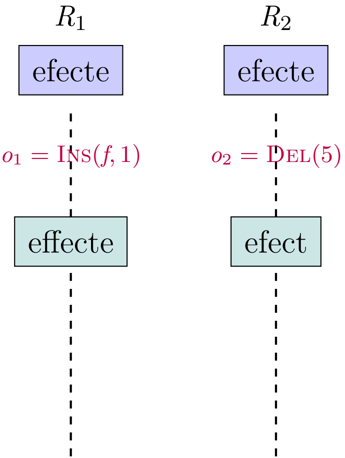

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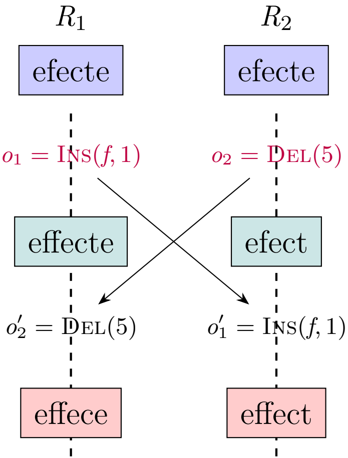

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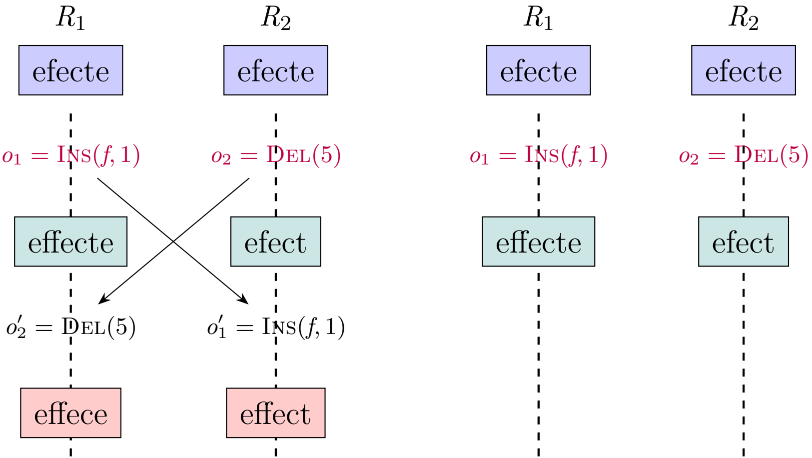

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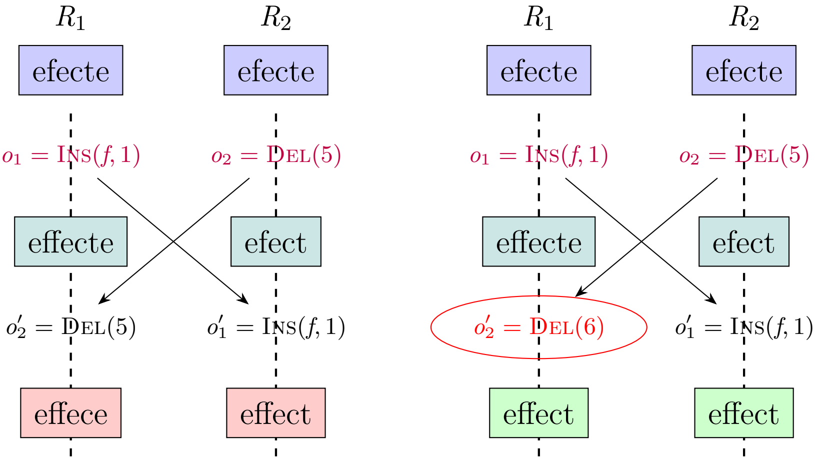
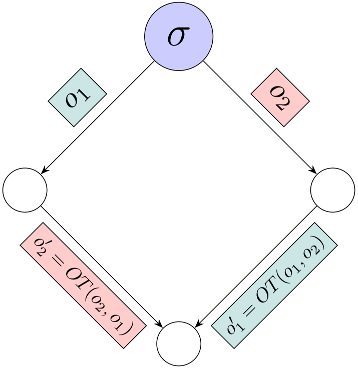
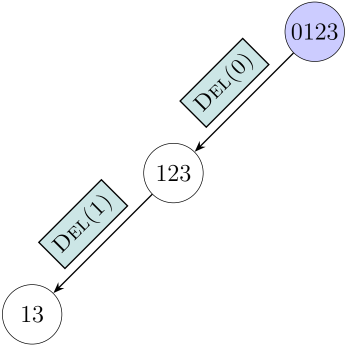




交换律 $\sigma; o_1; o'_2 \equiv \sigma; o_2; o'_1$

[Ellis and Gibbs, 1989]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DEL(0)
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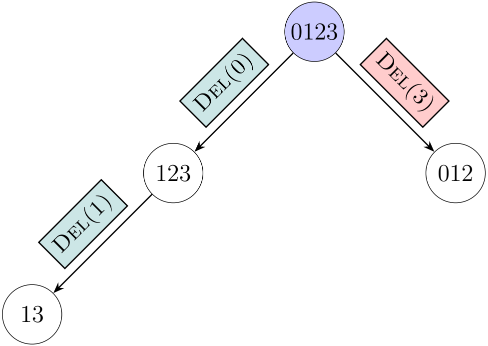
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012
DEL(1)
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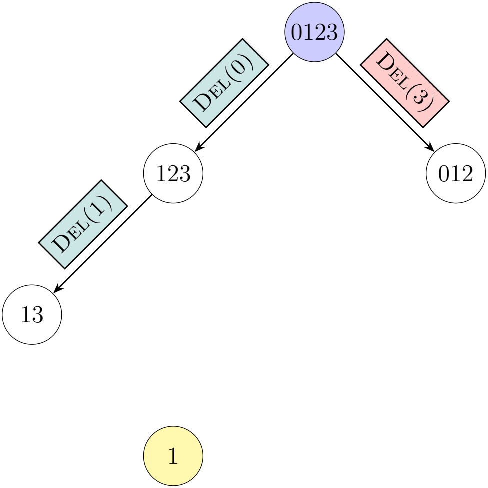
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DEL(1)
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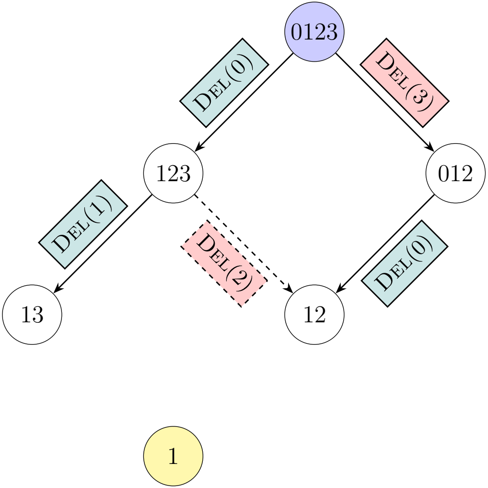
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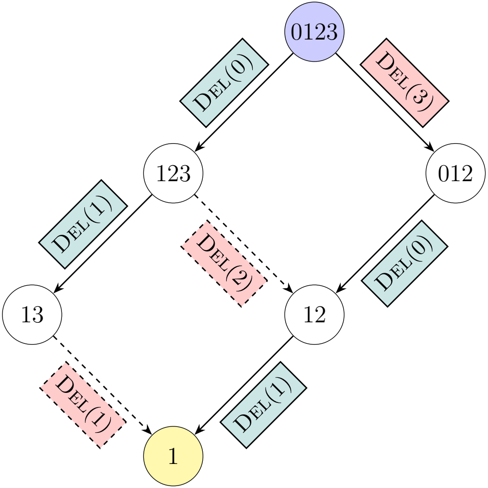
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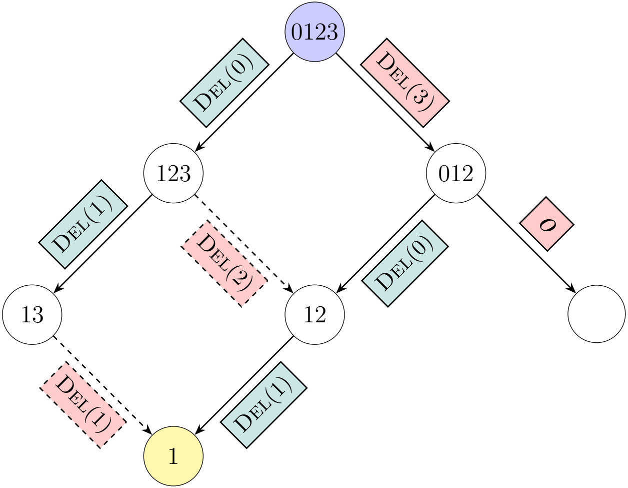
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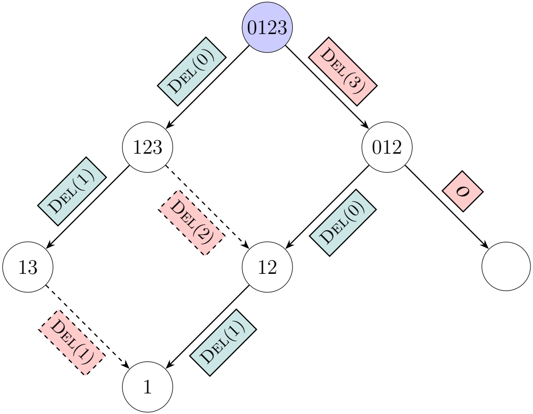
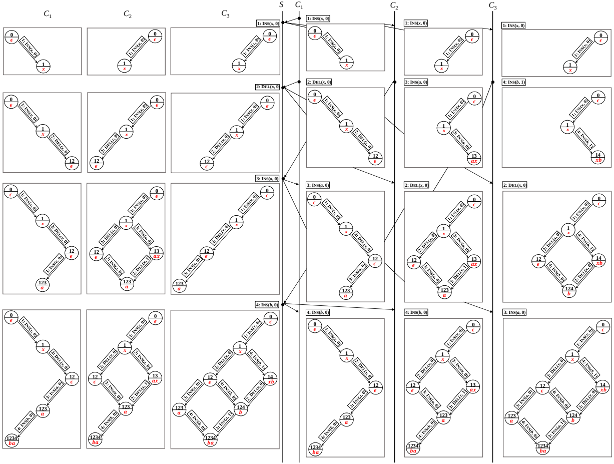
利用数据结构 2D 状态空间 [Xu, Sun, and Li, 2014] 控制何时以及如何执行"操作转换"



$2D$: LOCAL *vs.* GLOBAL

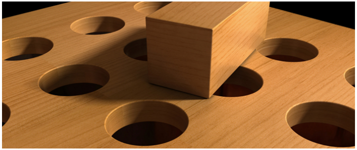
每个 Client 维护一个 $2D$ 状态空间



Server 维护 n 个 $2D$ 状态空间, 与 n 个 Clients 对应

Mismatch!

$\mathcal{A}_{\text{weak}}$ 所规定的全局性质



Jupiter 协议中，每个 replica 所维护的局部视图

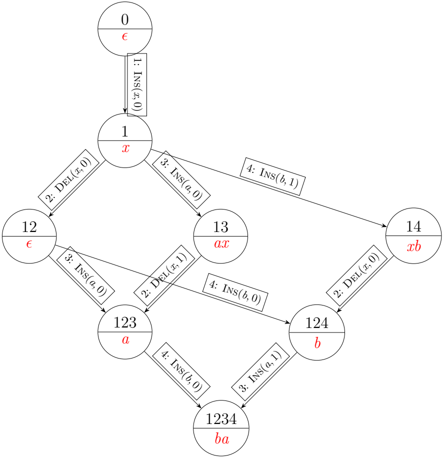
CJupiter (Compact Jupiter)

CJupiter (Compact Jupiter)

定理 (等价性)

在相同的操作调度下，CJupiter 与 Jupiter 中的对应 *replica* 的行为 *(*状态序列*)* 是相同的。

CJupiter 为每个 replica 维护一个 n-ary 有序状态空间



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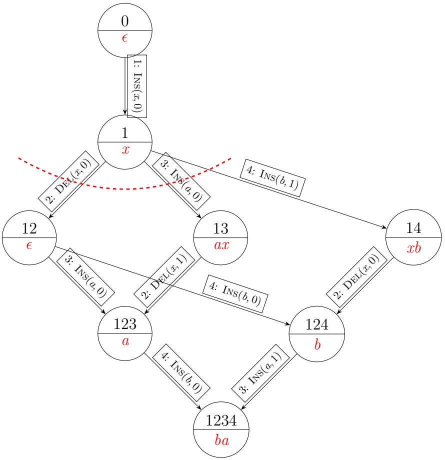
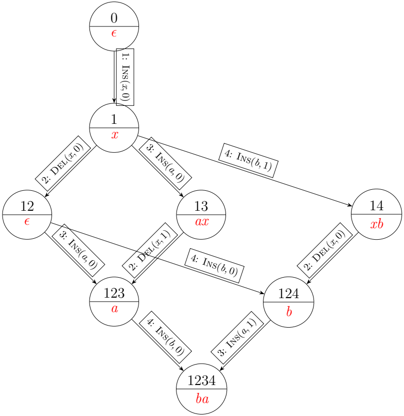


命题 (Compactness of CJupiter)

CJupiter 所维护的 $(n+1)$ 个 n-*ary* 有序状态空间是相同的。



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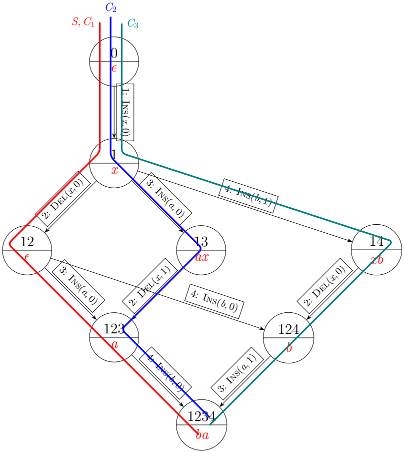
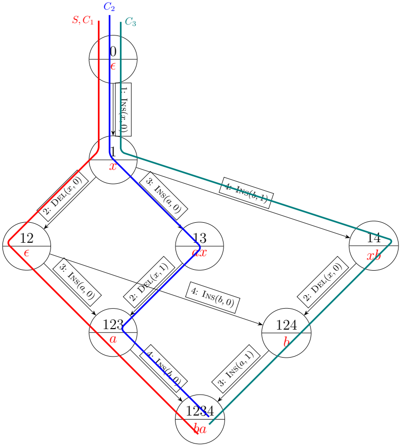


每个 replica 的行为对应于该状态空间中的一条**路径**

CJupiter 满足 Weak List Specification

关注某个 n-ary 有序状态空间，三步骤证明“状态对兼容性”



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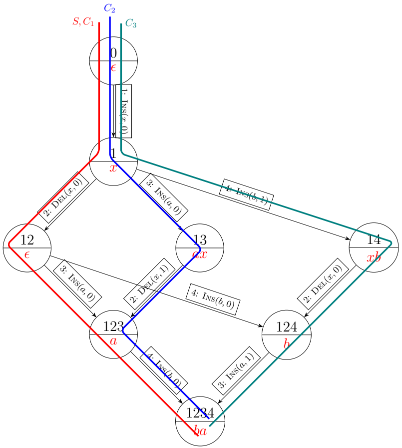
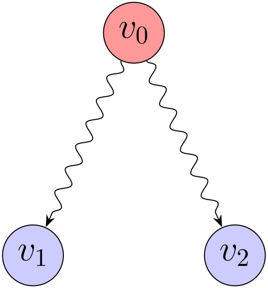


反证法、数学归纳法、分情形分析法

1 任取两个状态节点 v_1 和 v_2

引理 (LCA (Lowest Common Ancestor))

n-ary 有序状态空间中的任意一对状态节点都有**唯一的**最近公共祖先。

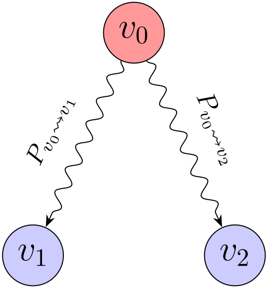


$$v_0 = \mathrm{LCA}(v_1, v_2)$$

② 考虑从 $v_0 = \mathrm{LCA}(v_1, v_2)$ 到 v_1 和 v_2 的两条路径

引理 (Disjoint Paths)

路径 $P_{v_0 \rightsquigarrow v_1}$ 上包含的操作集 $O_{v_0 \rightsquigarrow v_1}$ 与路径 $P_{v_0 \rightsquigarrow v_2}$ 上包含的操作集 $O_{v_0 \rightsquigarrow v_2}$ 不相交。

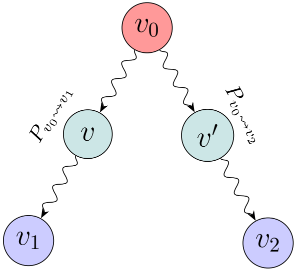


$v_0 = \mathrm{LCA}(v_1, v_2)$

③ 考虑两条路径上的状态

引理 (Compatible Paths)

$P_{v_0 \rightsquigarrow v_1}$ 上的任一状态 v 与 $P_{v_0 \rightsquigarrow v_2}$ 上的任一状态 v' 是兼容的。

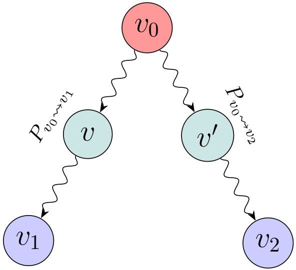


$$v_0 = \mathrm{LCA}(v_1, v_2)$$

3 考虑两条路径上的状态

引理 (Compatible Paths)

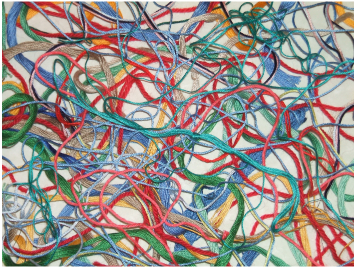
$P_{v_0 \rightsquigarrow v_1}$ 上的任一状态 v 与 $P_{v_0 \rightsquigarrow v_2}$ 上的任一状态 v' 是兼容的。



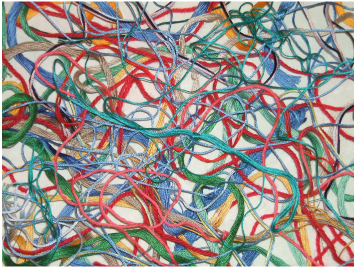
$$v_0 = \mathrm{LCA}(v_1, v_2)$$

$\therefore$ v_1 和 v_2 是兼容的

个人体会: 基于 OT 思想的协议晦涩难懂

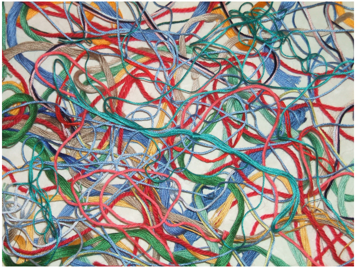


个人体会: 基于 OT 思想的协议晦涩难懂



- 协议多种多样
- 经常不加证明
- 证明是错误的

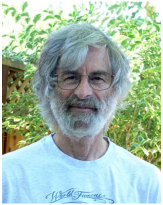
个人体会: 基于 OT 思想的协议晦涩难懂



- 协议多种多样
- 经常不加证明
- 证明是错误的
- 勘误也是错的

Model Checking: 使用 TLAPLUS

jupiter-tlaplus@github



Model Checking: 使用 TLAPLUS

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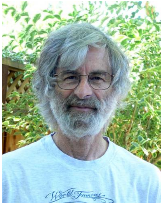

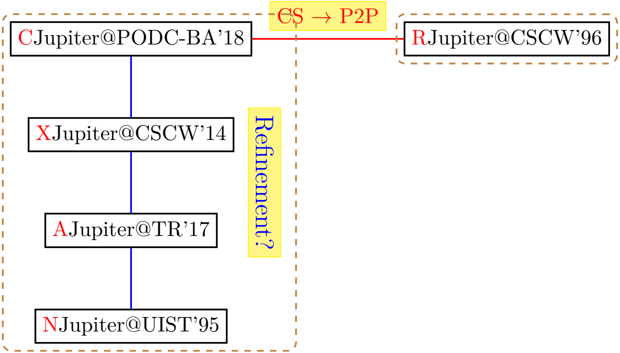

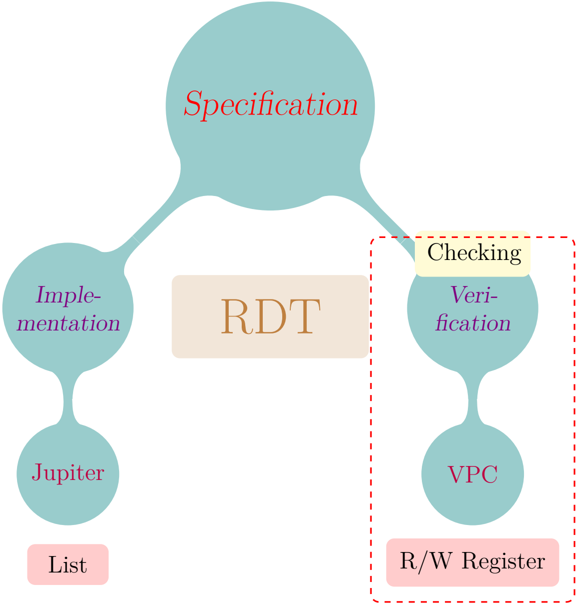
Specification
Imple-
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RDT
Checking
Veri-
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Jupiter
VPC
List
R/W Register

协议验证 (Verification of a Protocol)

[Bouajjani, Enea, and Hamza, 2014] [Bouajjani et al., 2017]

执行验证 (Verification of an Execution)

协议验证 (Verification of a Protocol)

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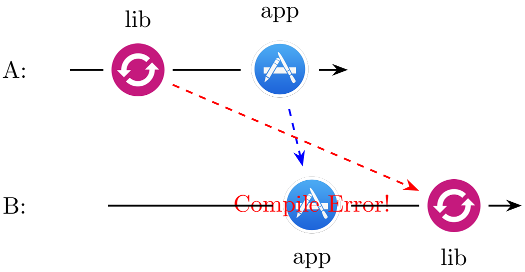
执行验证 (Verification of an Execution)

黑盒测试/确认系统是否提供了其所声称的数据一致性

[DeCandia et al., 2007] [Golab, Li, and Shah, 2011]

PRAM: 包含存储系统常提供的最基本的 “会话” (session) 一致性 [Terry et al., 1994] [Brzezinski, Sobaniec, and Wawrzyniak, 2004]



PRAM 保证 “单调写” 性质

定义 (VPC (Verifying PRAM Consistency) 判定问题)

实例: 系统执行 *(execution e)*

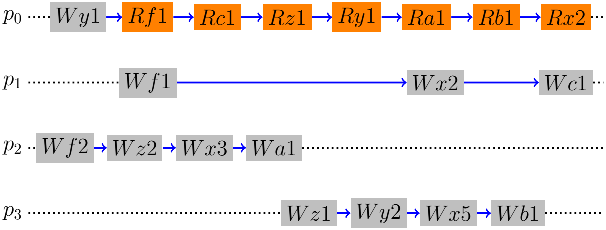
问题: 该执行 e 是否满足 PRAM 一致性模型 $(\mathcal{C})$?

$$e \in \mathcal{C} \Rightarrow \{0, 1\}?$$

定义 (系统执行)

系统执行 $e \triangleq \{h_p \mid h_p : \text{进程 } p \text{ 上的读写操作序列}\}$

规模 n: 系统执行中读写操作的总数

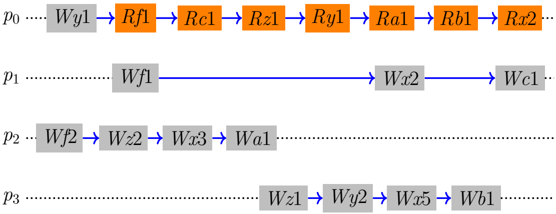


定义 (PRAM 一致性模型)

系统执行 e 满足 PRAM 一致性

$$\Longleftrightarrow$$

$\forall p : p$ 上所有操作与其它进程上所有写操作存在合法调度



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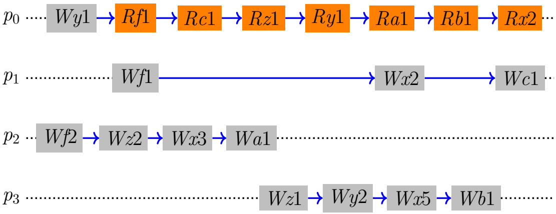

$$p_0 : Wf2\ Wf1\ Wz2\ Wz1\ Wy2\ Wy1\ Rf1\ Wx5\ Wx3\ Wx2\ Wc1\ Rc1\ Rz1\ Ry1\ Wa1\ Ra1\ Wb1\ Rb1\ Rx2$$

定义 (PRAM 一致性模型)

系统执行 e 满足 PRAM 一致性

$$\iff$$

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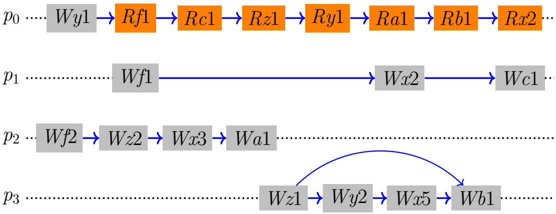

$p_0 :$ $Wf2$ $Wf1$ $Wz2$ $Wz1$ $Wy2$ $Wy1$ $Rf1$ $Wx5$ $Wx3$ $Wx2$ $Wc1$ $Rc1$ $Rz1$ $Ry1$ $Wa1$ $Ra1$ $Wb1$ $Rb1$ $Rx2$

定义 (PRAM 一致性模型)

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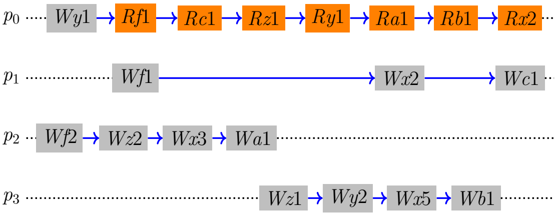

$$p_0 : Wf2\ Wf1\ Wz2\ Wz1\ Wy2\ Wy1\ Rf1\ Wx5\ Wx3\ Wx2\ Wc1\ Rc1\ Rz1\ Ry1\ Wa1\ Ra1\ Wb1\ Rb1\ Rx2$$

VPC 问题的四种变体 (按 “执行” 的类型) 及复杂度
([*] : 本文工作)

	(S)ingle variable	*(M)ultiple variables*
write (D)uplicate values	VPC-SD	VPC-MD
write (U)nique value	VPC-SU	VPC-MU

VPC 问题的四种变体 (按 “执行” 的类型) 及复杂度
([*]：本文工作)

	(S)ingle variable	*(M)ultiple variables*
write (D)uplicate values	VPC-SD (NP-complete) [*]	VPC-MD (NP-complete) [*]
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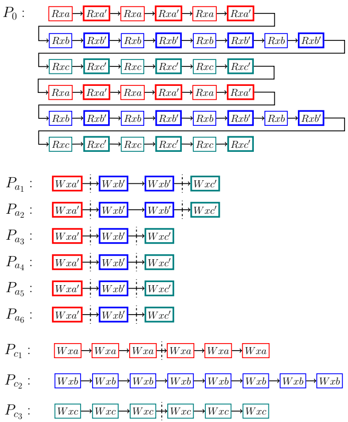
VPC 问题的四种变体 (按 “执行” 的类型) 及复杂度
([∗] : 本文工作)

	(S)ingle variable	*(M)ultiple variables*
write (D)uplicate values	VPC-SD (NP-complete) [∗]	VPC-MD (NP-complete) [∗]
write (U)nique value	VPC-SU (P) [Golab, Li, and Shah, 2011]	VPC-MU (P) [∗]

Read-mapping [Gibbons and Korach, 1997]: $\forall r,\ \exists! w,\ f(r) = w.$

VPC-SD (VPC-MD) 是 NP-complete 问题

VPC-SD (VPC-MD) 是 NP-complete 问题



Unary 3-Partition 实例 $A = \{2,2,1,1,1,1\}, m = 2, B = 4$ 对应的 VPC-SD 执行


Marzio De Biasi
30k 5 76 231

"Basically I'm a programmer :-)

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I enjoy reading papers about the complexity of puzzle games, and I'm writing (amateur) proofs on the complexity of a few puzzle games."

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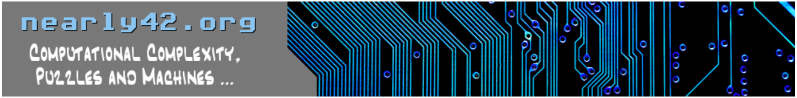
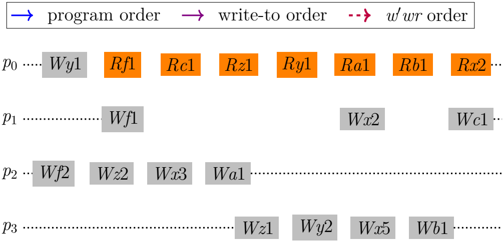
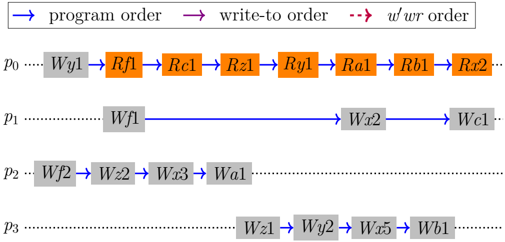


VPC-MU 的多项式算法 RW-CLOSURE



RW-CLOSURE 算法示例: 在传递闭包之上迭代应用 $w' wr$ 规则

VPC-MU 的多项式算法 RW-Closure



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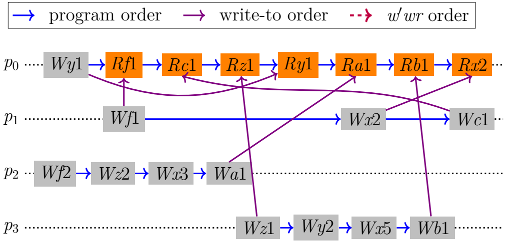
VPC-MU 的多项式算法 RW-CLOSURE



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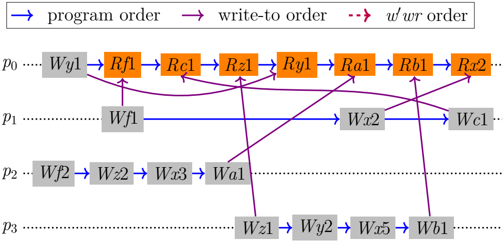
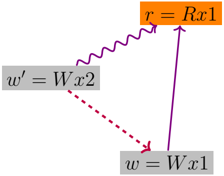
VPC-MU 的多项式算法 RW-CLOSURE



RW-CLOSURE 算法示例：在传递闭包之上迭代应用 $w'\,wr$ 规则



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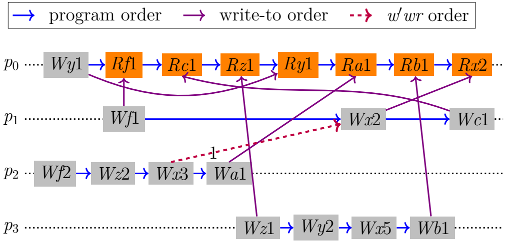

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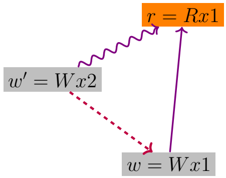

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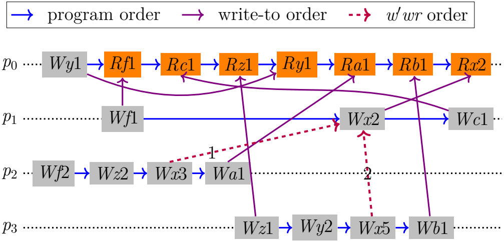

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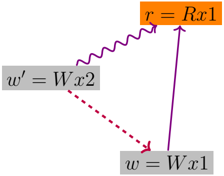

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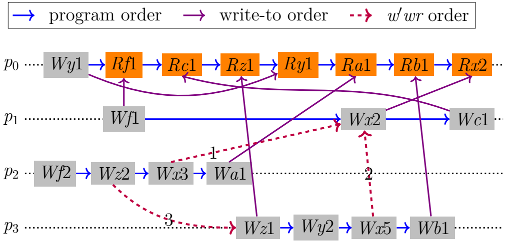

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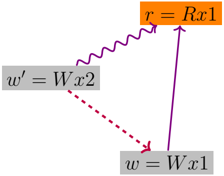
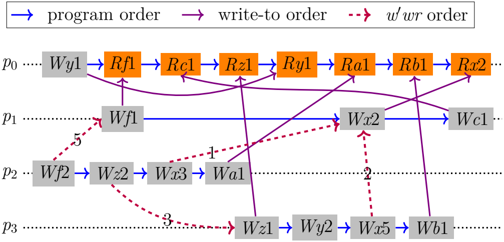


VPC-MU 的多项式算法 RW-Closure



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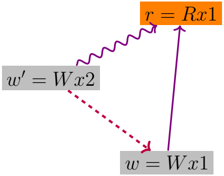

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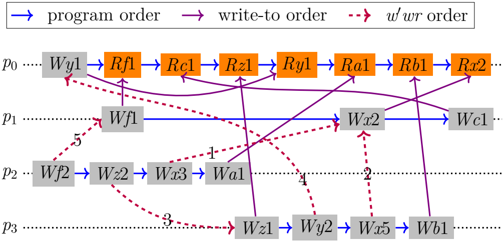

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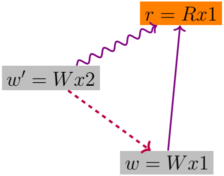

定理 (RW-Closure 算法正确性)

VPC-MU 实例满足 PRAM 一致性

$$\Longleftrightarrow$$

RW-Closure 算法所得图是 DAG 图

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证明

"$\Longrightarrow$" 反证法

"$\Longleftarrow$" 对读操作作数学归纳, 构造合法调度

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$$\iff$$

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证明

"$\Longrightarrow$" 反证法

"$\Longleftarrow$" 对读操作作数学归纳, 构造合法调度

RW-Closure 算法复杂度:

$$\underbrace{O(n^2)}_{\#\text{loops}} \cdot \underbrace{O(n^3)}_{\text{transitive closure}} = O(n^5)$$

RW-CLOSURE 算法的缺点:

- 在全图上应用 $w'wr$ 规则
- 应用 $w'wr$ 规则无特定顺序

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VPC-MU 的多项式算法 Read-Centric 要点:

- 增量式调度每个读操作
- 在读操作诱导的局部子图上按逆拓扑序应用 $w' wr$ 规则

定理 (Read-Centric 算法正确性)

VPC-MU 实例满足 PRAM 一致性

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Read-Centric 算法所得图是 DAG 图

证明

$$\text{Read-Centric} \xleftrightarrow{\textit{Reachability}} \text{RW-Closure}$$

定理 (READ-CENTRIC 算法正确性)

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$$\Longleftrightarrow$$

READ-CENTRIC 算法所得图是 DAG 图

证明

$$\text{READ-CENTRIC} \xleftrightarrow{\textit{Reachability}} \text{RW-CLOSURE}$$

READ-CENTRIC 算法复杂度:

$$\underbrace{O(n)}_{\#\text{reads}} \cdot \underbrace{O(n \cdot n^2)}_{\text{TOPO-SCHEDULE}} = O(n^4)$$

VPC 在相关工作中的意义

较早关注（分布式系统领域）“弱一致性模型验证”问题（2013∼）:

- **强一致性:** [Gibbons and Korach, 1997] [Cantin, Lipasti, and Smith, 2005] [Golab, Li, and Shah, 2011]
- **弱一致性:** [Furbach et al., 2014] [Bouajjani et al., 2017] [Emmi and Enea, 2018]

VSC (Verifying Sequential Consistency) 与 **VL** (Verifying Linearizability) 问题的复杂度 [Gibbons and Korach, 1997]

Variants	**VSC**	**VL**
General	NP-complete	NP-complete
2 Operations/Process	NP-complete	NP-complete
2 Variables	NP-complete	NP-complete
3 Processes	NP-complete	$O(n \log n)$
Read-mapping	NP-complete	$O(n \log n)$
Write-order	NP-complete	$O(n \log n)$
`read&write` only	NP-complete	NP-complete
Conflict-order	$O(n \log n)$	$O(n \log n)$

VMC (Verifying Memory Coherence) 问题的复杂度 [Cantin, Lipasti, and Smith, 2005]

Variants	**Read/Write**	**Read-Modify-Write**
1 Operation/Process	$O(n \lg n)$	$O(n^2)$
2 Operations/Process	?	NP-complete
3+ Operations/Process	NP-complete	NP-complete
Constant k processes	$O(n^k)$	$O(n^k)$
1 Write/Value (Read-mapping)	$O(n)$	$O(n \lg n)$
2 Writes/Value	NP-complete	?
3+ Writes/Value	NP-complete	NP-complete
Write-order	$O(n^2)$	$O(n)$

Atomicity [2] 相关一致性模型验证问题复杂度 (假设: 不允许写重复值)

	Safety	**Regularity**	**Atomicity**	**Sequential**
Offline [Anderson et al., 2010]	$O(n^2)$	$O(n^2)$	$O(n^3)$	*not studied*
Online [3] [Golab, Li, and Shah, 2011]	$O(n)$	$O(n)$	$O(n \log n)$	$\mathsf{Poly}(n)$

[2] 也称 Linearizability

[3] 包含其它假设

k-AV (k-Atomicity Verification) 问题复杂度

Problems	Variants	Results	Work
1-AV	General	NP-complete	[Gibbons and Korach, 1997]
1-AV	Write unique value	$O(n \log n)$	[Gibbons and Korach, 1997]
2-AV	Write unique value	$O(n \log n)$	[Golab, Hurwitz, and Li, 2013]
k-AV	Write unique value Bounded concurrency	$O(n^2)$	[Golab et al., 2015] [Golab et al., 2018]
k-AV	Write unique value		

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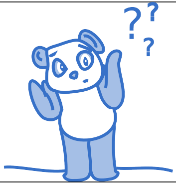
k-AV	Write unique value		

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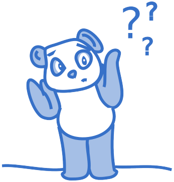
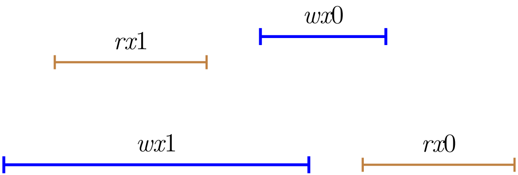
2-AV	Write unique value	$O(n \log n)$	[Golab, Hurwitz, and Li, 2013]
k-AV	Write unique value Bounded concurrency	$O(n^2)$	[Golab et al., 2015] [Golab et al., 2018]
k-AV	Write unique value		

Atomicity = 实时序 + 读写语义

[Lamport, 1986]



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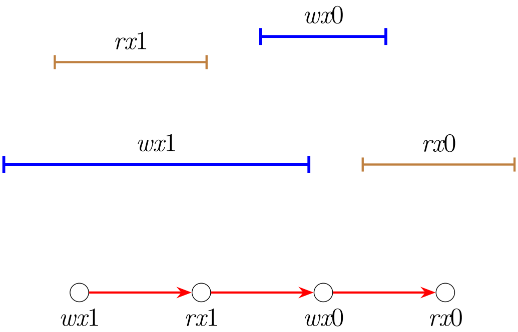

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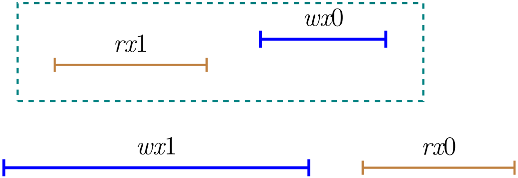

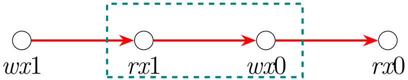

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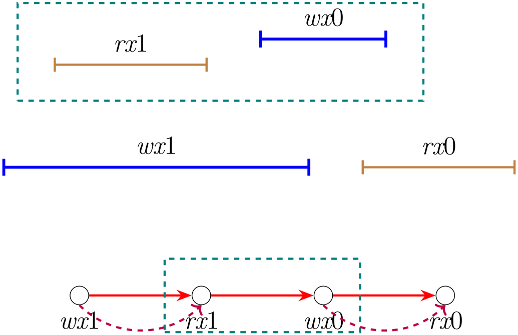
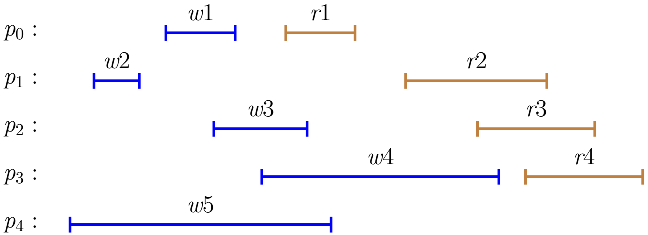


k-Atomicity = 实时序 + k-读写语义

[Aiyer, Alvisi, and Bazzi, 2005] [Taubenfeld, 2013]



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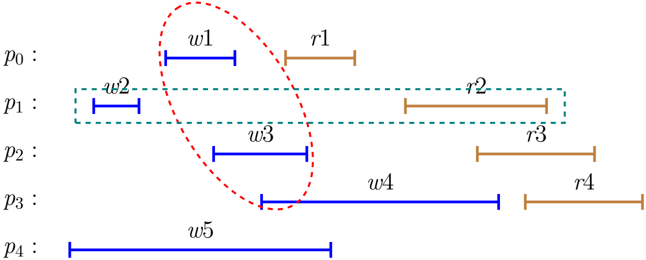

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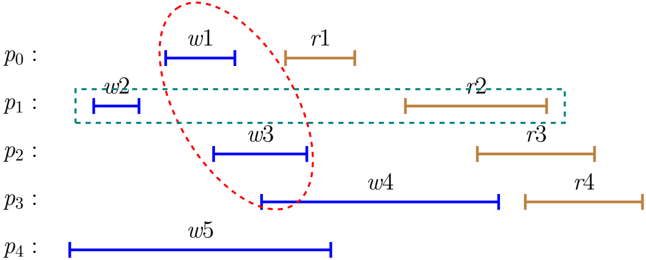
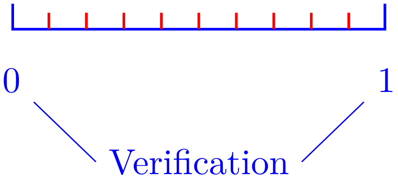


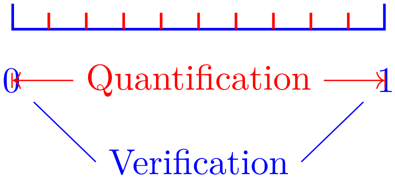
3-Atomicity : $w5$ $w2$ $w1$ $r1$ $w3$ $w4$ $r2$ $r3$ $r4$

定义 (k-AV (k-Atomicity Verification) 判定问题)

实例: 系统执行 e (不允许写重复值)、参数 k

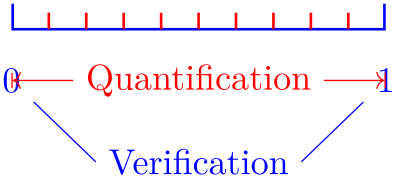
问题: 该执行 e 是否满足 k-Atomicity?


0
1
Verification



协议量化分析

[Lee and Welch, 2005] [Bailis et al., 2012] [Chatterjee and Golab, 2017]

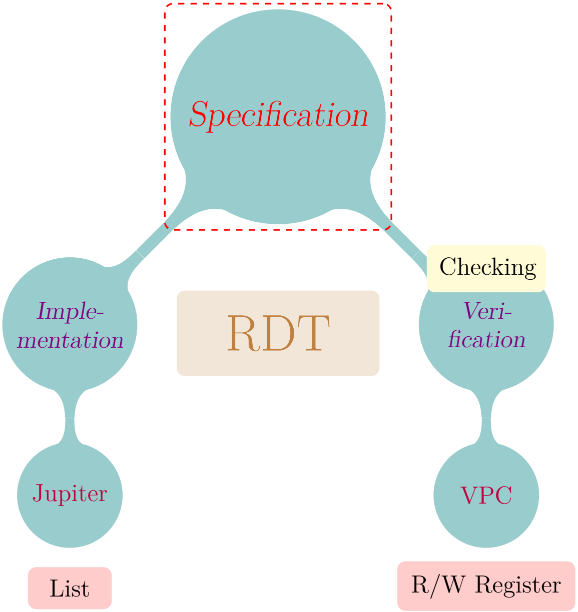


协议量化分析

[Lee and Welch, 2005] [Bailis et al., 2012] [Chatterjee and Golab, 2017]

PA2AM: Probabilistically-Atomic 2-Atomicity


One more thing...


Specification
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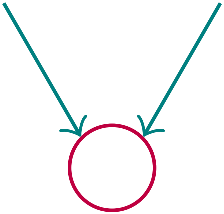
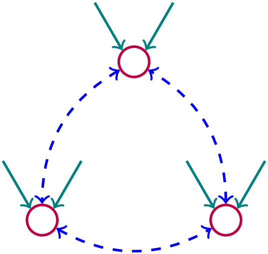
mentation
RDT
Checking
Veri-
fication
Jupiter
VPC
List
R/W Register

多处理器系统中的并发数据类型

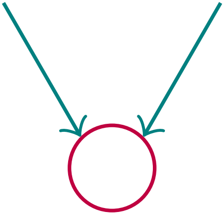


分布式系统中的复制数据类型

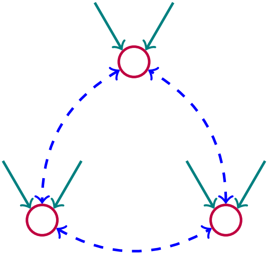


规约：数据一致性模型 (Consistency Model)

多处理器系统中的并发数据类型



分布式系统中的复制数据类型



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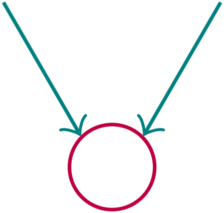
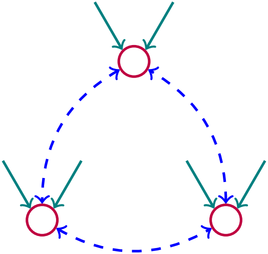


PL + DC

分布式系统中的复制数据类型



[Burckhardt et al., 2014]

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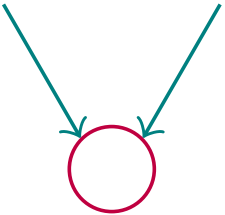
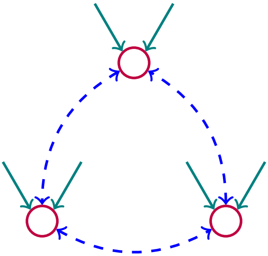
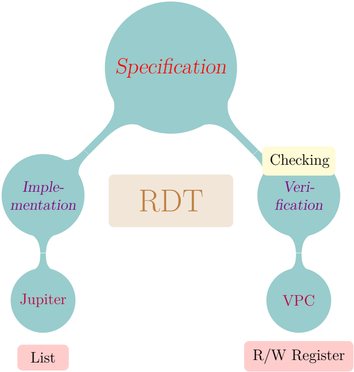
多处理器系统中的并发数据类型



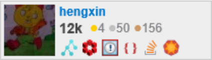
分布式系统中的复制数据类型



PL + DC + FM [Burckhardt et al., 2014]



hfwei@nju.edu.cn



hengxin@homepage hengxin@github hengxin@stackexchange

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当用户不再提交更新操作时, 所有 *replicas* 上的列表是相同的。

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对系统的中间状态缺少足够的约束

针对列表的操作转换函数 [Ellis and Gibbs, 1989]

$$OT\Big(\textsc{Ins}(a_1, p_1, pr_1), \textsc{Ins}(a_2, p_2, pr_2)\Big) = \begin{cases} \textsc{Ins}(a_1, p_1, pr_1) & p_1 < p_2 \\ \textsc{Ins}(a_1, p_1 + 1, pr_1) & p_1 > p_2 \\ \text{NOP} & p_1 = p_2 \wedge a_1 = a_2 \\ \textsc{Ins}(a_1, p_1 + 1, pr_1) & p_1 = p_2 \wedge a_1 \neq a_2 \wedge pr_1 > pr_2 \\ \textsc{Ins}(a_1, p_1, pr_1) & p_1 = p_2 \wedge a_1 \neq a_2 \wedge pr_1 \leq pr_2 \end{cases}$$

$$OT\Big(\textsc{Ins}(a_1, p_1, pr_1), \textsc{Del}(\_, p_2, pr_2)\Big) = \begin{cases} \textsc{Ins}(a_1, p_1, pr_1) & p_1 \leq p_2 \\ \textsc{Ins}(a_1, p_1 - 1, pr_1) & p_1 > p_2 \end{cases}$$

$$OT\Big(\textsc{Del}(\_, p_1, pr_1), \textsc{Ins}(a_2, p_2, pr_2)\Big) = \begin{cases} \textsc{Del}(\_, p_1, pr_1) & p_1 < p_2 \\ \textsc{Del}(\_, p_1 + 1, pr_1) & p_1 \geq p_2 \end{cases}$$

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定理 (RW-CLOSURE 算法正确性)

VPC-MU 实例满足 PRAM 一致性

$$\Longleftrightarrow$$

RW-CLOSURE 算法所得图是 DAG 图

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证明

"$\Longrightarrow$" 反证法

"$\Longleftarrow$" 难点: *DAG* 图蕴含着多个全序

技巧: 对读操作作数学归纳, 构造合法调度

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RW-CLOSURE 算法复杂度:

$$\underbrace{O(n^2)}_{\#\text{loops}} \cdot \underbrace{O(n^3)}_{\text{transitive closure}} = O(n^5)$$

READ-CENTRIC 算法复杂度:

$$\underbrace{O(n)}_{\text{iterations}} \cdot \underbrace{O(n \cdot n^2)}_{\text{TOPO-SCHEDULE}} = O(n^4)$$

引理 (TOPO-SCHEDULE 的非迭代性)

设 TOPO-SCHEDULE 正在处理读操作 r,
则局部子图中的每个写操作最多只有一次机会
在满足规则 $w' wr$ 的三元组中扮演 “w' 角色”。

实验评估

实验目的[1]：

1. 考察 Read-Centric 算法的实际效率 (*vs.* 渐近时间复杂度)
2. 对比 Read-Centric 算法与 RW-Closure 算法的效率

[1] 机器配置: Intel Core i7 3.40GHZ, 4GB RAM.

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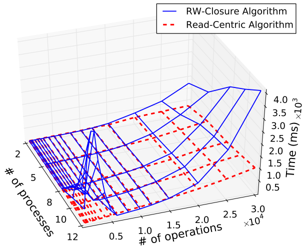
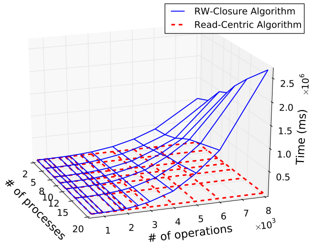
两类负载:

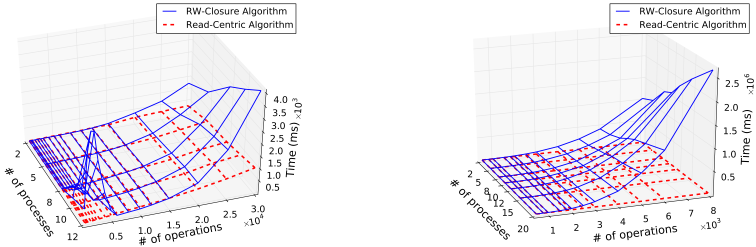
1. 随机生成的系统执行
2. 满足 PRAM 一致性的系统执行 ($\approx$ 最坏情况输入)

[1]机器配置: Intel Core i7 3.40GHZ, 4GB RAM.



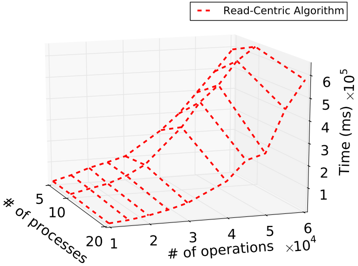


RW-Closure 算法与 Read-Centric 算法在 (左) 随机生成的执行及 (右) 满足 PRAM 一致性的执行上的运行时间。



RW-Closure 算法与 Read-Centric 算法在 (左) 随机生成的执行及 (右) 满足 PRAM 一致性的执行上的运行时间。

(右) 20 个进程、8,000 个操作:
Read-Centric 可获得 694 倍加速.



READ-CENTRIC 算法在满足 PRAM 一致性的执行上的运行时间

READ-CENTRIC: 20 个进程、60,000 个操作 $<$ 600s [1]

RW-CLOSURE: 20 个进程、8,000 个操作 $>$ 3,000s

[1] 用于测试, 规模可用

$\mathcal{I}$-Atomicity = i-实时序 + 读写语义

$\mathcal{I}$: Inversions

$\mathcal{I}$-Atomicity $=$ i-实时序 $+$ 读写语义

$\mathcal{I}$: Inversions

$$f(\{\text{inversions}\}) \leq i$$

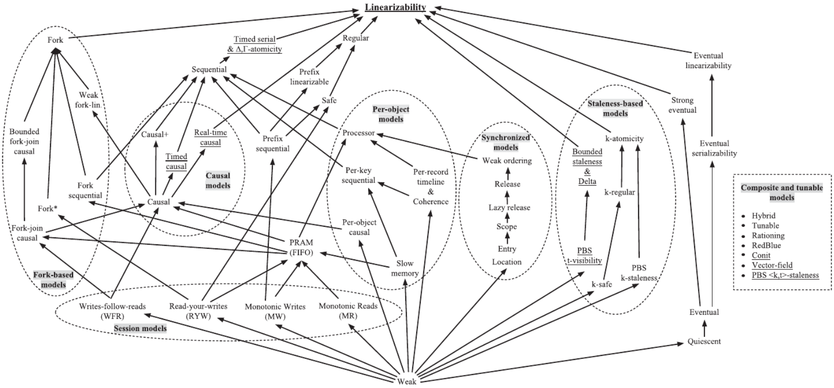
$\boxed{\mathcal{I}\text{-Atomicity} = i\text{-实时序} + \text{读写语义}}$

$\mathcal{I}$: Inversions

$$f(\{\text{inversions}\}) \leq i$$

定义框架

(50 种) 一致性模型关系图 [Viotti and Vukolić, 2016] [Burckhardt, 2014]



建立统一的形式化框架