Specification, Implementation, and Complexity of Replicated Data Types with Composite Operations

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Abstract

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1 Introduction

2 Preliminaries

2.1 Observed-Remove Set (OR-Set)

$$\mathcal{F}_{\texttt{orset}}(\texttt{rd}, E, \mathsf{op}, \mathsf{vis}, \mathsf{ar}) = \{a \mid \exists e \in E.\mathsf{op}(e) = \texttt{add}(a) \tag{1}$$

$$\land \big(\forall f \in E.\mathsf{op}(f) = \texttt{rm}(a) \implies \neg(e \xrightarrow{\mathsf{vis}} f)\big)\}. \tag{2}$$

3 Replicated Data Types with Composite Operations

3.1 Specification

We consider a composite operation of a replicated data type τ in the form of $C = A \oplus B$, where A, B, and C are different objects of type τ.

Q : Generalize to different data types?

Following [1], we specify the semantics of a composite operation $A \oplus B$ of a replicated data type τ by a function $\mathcal{F}_\tau$ that determines the return value of $\oplus$ based on prior operations performed on the two objects involved (i.e., A and B). However, $\mathcal{F}_\tau$ for a composite operation $\oplus$ takes as parameters two, not one as in [1], *operation contexts*, one on each object involved.

Note: Partial operation context [2]

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▶ **Definition 1** (Product of Operation Contexts)**.** Consider two operation contexts for the same replicated data type τ:

$$\mathcal{C}_A = (E_A, \mathsf{op}_A, \mathsf{vis}_A, \mathsf{ar}_A) \tag{3}$$

$$\mathcal{C}_B = (E_B, \mathsf{op}_B, \mathsf{vis}_B, \mathsf{ar}_B) \tag{4}$$

The product $\mathcal{C} = \mathcal{C}_A \times \mathcal{C}_B$ of $\mathcal{C}_A$ and $\mathcal{C}_B$ is also an operation context defined as $\mathcal{C} = (E, \mathsf{op}, \mathsf{vis}, \mathsf{ar})$, where

- $E = E_A \times E_B$
- $\mathsf{op} = \mathsf{op}_A \sqcup \mathsf{op}_B$
- $\mathsf{vis} = \mathsf{vis}_A \times \mathsf{vis}_B$
- $\mathsf{ar} = \mathsf{ar}_A \times \mathsf{ar}_B$

▶ **Definition 2.**

$$\mathcal{F}_\tau(\oplus, \mathcal{C}_A, \mathcal{C}_B) = \mathcal{F}_\tau(\oplus, \mathcal{C}_A \times \mathcal{C}_B) \tag{5}$$

4 Replicated Set with Composite Operations

We consider the replicated set data type with composite operations including union ($\cup$), intersection ($\cap$), and set difference ($\setminus$).

4.1 Specification

$$\mathcal{F}_{\mathtt{orset}}(A \setminus B, \mathcal{C}_A, \mathcal{C}_B) \tag{6}$$

$$= \{a \mid \exists (e, e') \in E_A \times E_B. \Big(\big(\mathtt{op}(e) = \mathtt{add}(a) \land \mathtt{op}(e') = \mathtt{rm}(a)\big) \tag{7}$$

$$\land \forall (f, f') \in E_A \times E_B. \big((\mathtt{op}(f) = \mathtt{add}(a) \land \mathtt{op}(f') = \mathtt{add}(a)) \lor (\mathtt{op}(f) = \mathtt{rm}(a) \land \mathtt{op}(f') = \mathtt{add}(\tag{8}$$

$$\implies \neg\big((e, e') \xrightarrow{\mathsf{vis}} (f, f')\big)\Big)\}. \tag{9}$$

4.2 Protocol

5 Related Work

6 Conclusion and Future Work

References

1 Sebastian Burckhardt, Alexey Gotsman, Hongseok Yang, and Marek Zawirski. Replicated data types: Specification, verification, optimality. In *Proceedings of the 41st ACM SIGPLAN-SIGACT Symposium on Principles of Programming Languages*, POPL '14, pages 271–284, New York, NY, USA, 2014. ACM. URL: http://doi.acm.org/10.1145/2535838.2535848, doi:10.1145/2535838.2535848.

2 Alexey Gotsman and Hongseok Yang. Composite replicated data types. In *European Symposium on Programming Languages and Systems*, pages 585–609. Springer, 2015.