

# Parameterized and Runtime-tunable Snapshot Isolation in Distributed Transactional Key-value Stores

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# Parameterized and Runtime-tunable Snapshot Isolation

## RVSI: Relaxed Version Snapshot Isolation

### 1 Definition of RVSI

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## RVSI: Relaxed Version Snapshot Isolation

### 1 Definition of RVSI

# Principle of RVSI

- ▶ Parameters  $(k_1, k_2, k_3)$  to control the severity of the anomalies w.r.t SI

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<sup>1</sup>RC: Read Committed

# Principle of RVSI

- ▶ Parameters  $(k_1, k_2, k_3)$  to control the severity of the anomalies w.r.t SI
- ▶  $RC^1 \supset RVSI(k_1, k_2, k_3) \supset SI$
- ▶  $RVSI(\infty, \infty, \infty) = RC$        $RVSI(1, 0, *) = SI$

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# Principle of RVSI

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– “*Snapshot Read*” property of SI

## 1. “stale” data versions

# Principle of RVSI

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– “*Snapshot Read*” property of SI

1. “stale” data versions
2. “concurrent” data versions

# Principle of RVSI

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– “*Snapshot Read*” property of SI

1. “stale” data versions
2. “concurrent” data versions
3. “non-snapshot” data versions



# Principle of RVSI

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– “*Snapshot Read*” property of SI

1. “stale” data versions
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bounded staleness

# Principle of RVSI

...

– “*Snapshot Read*” property of SI

1. “stale” data versions bounded staleness
2. “concurrent” data versions bounded concurrency level
3. “non-snapshot” data versions

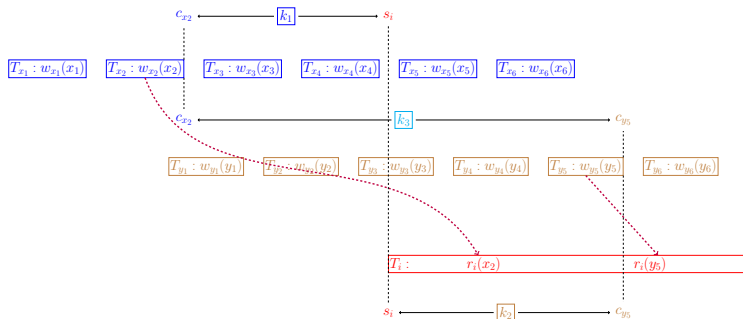
# Principle of RVSI

...

– “*Snapshot Read*” property of SI

- |                                 |                           |
|---------------------------------|---------------------------|
| 1. “stale” data versions        | bounded staleness         |
| 2. “concurrent” data versions   | bounded concurrency level |
| 3. “non-snapshot” data versions | bounded distance          |

# Illustration of RVSI



# Definition of RVSI

$(k_1\text{-BV})$

$$\forall r_i(x_j), w_k(x_k), c_k \in h : \left( c_j \in h \wedge \bigwedge_{k=1}^m (c_j \prec_h c_k \prec_h s_i) \right) \Rightarrow m < k_1,$$

$(k_2\text{-FV})$

$$\forall r_i(x_j), w_k(x_k), c_k \in h : \left( c_j \in h \wedge \bigwedge_{k=1}^m (s_i \prec_h c_k \prec_h c_j) \right) \Rightarrow m \leq k_2,$$

$(k_3\text{-SV})$

$$\forall r_i(x_j), r_i(y_l), w_k(x_k), c_k \in h : \left( \bigwedge_{k=1}^m (c_j \prec_h c_k \prec_h c_l) \right) \Rightarrow m \leq k_3.$$

# Definition of RVSI

$$h \in \text{RVSI} \iff h \in k_1\text{-BV} \cap k_2\text{-FV} \cap k_3\text{-SV} \cap \text{WCF}.$$

## Theorem

$$\text{RVSI}(1, 0, *) = \text{SI}.$$









