Z3 Python API 入门

SAT SMT by example (Basics and Equations)

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在线观看幻灯片

本地Jupyter观看幻灯片需要安装RISE

Z3 介绍

Z3是微软开发的一个跨平台的高性能定理证明器,用于许多应用程序,例如:软件/硬件验证和测试,约束求解,混合系统分析,安全性验证,生物学和几何问题等.

Z3提供多种语言的API,使得调用Z3变得十分方便. 我们使用Python API.

主页: https://github.com/Z3Prover/z3

SAT和SMT的关系

SMT solver是SAT solver的前端,它将输入的SMT表达式转换为CNF,再调用SAT solver.

SMT就像高级编程语言, SAT像汇编语言.

一些SMT solver使用第三方SAT solver作为后端, Z3有自己的SAT solver后端.

常见SMT solver

- MathSAT (http://mathsat.fbk.eu/)
- Z3 (https://github.com/Z3Prover/z3/)
- CVC4 (http://cvc4.cs.nyu.edu/web/)
- Yices 2 (http://yices.csl.sri.com/)
- CUDD (http://vlsi.colorado.edu/~fabio/CUDD/)
- PicoSAT (http://fmv.jku.at/picosat/)
- Boolector (http://fmv.jku.at/boolector/)
- PySMT (https://github.com/pysmt/pysmt)
- JavaSMT (https://github.com/sosy-lab/java-smt)
- 作者自己开发的玩具 MK85 (https://yurichev.com/MK85/)

CNF表达式

举例: $(\neg A \lor B) \land (C \lor \neg D)$

SAT solver专用于解CNF形式的布尔表达式

常见SAT solver

- CaDiCaL (https://github.com/arminbiere/cadical)
- Glucose (http://www.labri.fr/perso/lsimon/glucose/)
- Lingeling (http://fmv.jku.at/lingeling/)
- MapleCM (http://sat2018.forsyte.tuwien.ac.at/solvers/main_and_glucose_hack/)
- Maplesat (https://sites.google.com/a/gsd.uwaterloo.ca/maplesat/)
- Minicard (https://github.com/liffiton/minicard)
- Minisat (http://minisat.se/MiniSat.html)
- PySAT (https://github.com/pysathq/pysat)

NP问题的另一种解释

许多算法 快 或 慢 依赖于输入规模.

- 常数时间 O(1)
- 线性时间 O(n)
- 对数时间 O(log

(n))

• 指数时间 $O(2^{p(n)})$

NP问题没有大O表示法.

CNF表达式的求解是NP问题,可以简单使用暴力枚举的方法. SAT solver实际上也是bruteforce,但它使用很多优化的方法,然而没有任何的SAT/SMT solver能预测求解需要的时间.

SAT/SMT solver到底是什么?

可以理解为解NP问题的API,因为所有NP都可以规约到SAT问题.

资源链接

幻灯片中的资料主要来源:

- SAT SMT by example (GIT)
- SAT SMT by example (PDF)
- MK85 toy-level bit-blaster (SRC)
- Z3Py Guide (Jupyter)
- Z3 GitHub
- Z3Examples

环境搭建

最新的发布版: https://github.com/Z3Prover/z3/releases

使用PyCharm的配置步骤: https://github.com/hengxin/sat-smt-satisfying/tree/master/z3-config

或通过Python PIP安装(只包含Python API, 没有Z3可执行程序):

n []: !pip3 install z3-solver

第一件事

import z3 或 from z3 import *

用多了发现还是 from z3 import * 方便

In [6]: from z3 import *

一个简单的例子

```
egin{aligned} x > 2 \ y < 10 \ x + 2 * y = 7 \end{aligned}
```

```
In [7]: x = Int('x')

y = Int('y')

solve(x > 2, y < 10, x + 2*y == 7)

[y = 0, x = 7]
```

- Int('x') 在Z3中创建了整型变量x,
- solve()函数可以求解一系列规约是否可满足,这里有三个规约,
- <, <=, >, >=, ==, !=是用于比较的操作符,类似C++中的重载操作符.

化简函数 simplify() 和 网页格式化设置 set_option(html_mode=True)

```
In [8]: print(simplify(And(x + 1 >= 3, x**2 + x**2 + y**2 + 2 >= 5)))

And(x >= 2, 2*x**2 + y**2 >= 3)

In [9]: set_option(html_mode=True) print(simplify(And(x + 1 >= 3, x**2 + x**2 + y**2 + 2 >= 5))) set_option(html_mode=False)

x ≥ 2 ∧ 2·x<sup>2</sup> + y<sup>2</sup> &ge; 3
```

2 2 上面的html代码显示效果: x ≥ 2 ∧ 2·x + y ≥ 3

实数

Z3中能表示任意大的整数,有理数和代数数

有理数也可以用十进制小数表示,最后的?表示输入被截断到指定的精度:

布尔逻辑

```
Z3支持的布尔逻辑操作: And(), Or(), Not(), Implies() (implication), If() (ifthen-else). == (Bi-implications).
```

举例: $p \Rightarrow q, r \Leftrightarrow \neg q, \neg p \lor r$

[p = False, q = True, r = False]

```
In [14]:  p = Bool('p') 
 q = Bool('q') 
 r = Bool('r') 
 solve(Implies(p, q), r == Not(q), Or(Not(p), r))
```

Solver()

如果规约复杂, solve() 难以直接使用.

```
In [15]: x = Int('x')

y = Int('y')

s = Solver()

print(s)

s.add(x > 10, y == x + 2)

print(s)

print(s.check())

if s.check() == sat:

print(s.model())

[]

[x > 10, y == x + 2]

sat

[x = 11, y = 13]
```

Z3无法解非线性多项式时间的规约,如求解 $2^x=3$ 将会输出 unknown:

```
In [16]: x = Real('x')

s = Solver()

s.add(2**x == 3)

print(s.check())
```

unknown

约束无法满足时, solve() 输出 no solution, s.check() 输出 unsat

```
In [17]: x = \text{Real}(\mathbf{'x'}) solve(x > 4, x < 0)

no solution

In [18]: x = \text{Real}(\mathbf{'x'}) s = Solver() s.add(x > 4, x < 0) print(s.check())

unsat
```

一次创建多个变量

```
In [19]: a, b, c = Ints('a b c')
          d, e = Reals('d e')
          solve(a > b + 2, a == 2*c + 10, c + b <= 1000, d >= e)
          [c = 0, b = 0, e = 0, d = 0, a = 10]
In [20]: X = [Int('x\%s'\% i) \text{ for i in range}(5)]
          print(X)
          X = IntVector('x', 5)
          Y = RealVector('y', 5)
          P = BoolVector('p', 5)
          print(X, Y, P, sep='\n')
          [x0, x1, x2, x3, x4]
          [x_0, x_1, x_2, x_3, x_4]
          [y_0, y_1, y_2, y_3, y_4]
[p_0, p_1, p_2, p_3, p_4]
```

机器数表示

```
In [21]: x, y = BitVec('x', 16), BitVec('y', 16)
print(x + 2)

# -1 is equal to 65535 for 16-bit integers
print(simplify(x + y - 1))

# Creating bit-vector constants

a, b = BitVecVal(-1, 16), BitVecVal(65535, 16)
print(simplify(a == b))

a, b = BitVecVal(-1, 32), BitVecVal(65535, 32)
# -1 is not equal to 65535 for 32-bit integers
print(simplify(a == b))
```

```
x + 2
65535 + x + y
True
False
```

Z3并不区分有符号或无符号,对于无符号的操作,有特定的操作符: ULT, ULE, UGT, UGE, UDiv, URem and LShR 对应 <, <=, >, >=, /, % and >>

```
In [22]: solve(x & y == \simy)

solve(x < 0)

# using unsigned version of <

solve(ULT(x, 0))

[y = 65535, x = 0]

[x = 65535]

no solution
```

为什么使用 Z3 Python API?

接下来将对比使用Python和直接使用Z3解下面的三元一次方程组:

$$3x + 2y - z = 1$$

 $2x - 2y + 4z = -2$
 $-x + \frac{1}{2}y - z = 0$

使用Z3 Python API:

```
In [23]: x, y, z = \text{Reals}('x y z')

s = \text{Solver}()

s.\text{add}(3 * x + 2 * y - z == 1)

s.\text{add}(2 * x - 2 * y + 4 * z == -2)

s.\text{add}(-x + 0.5 * y - z == 0)

print(s.\text{check}())

print(s.\text{model}())
```

直接使用Z3命令行工具(需要 z3 所在目录在 PATH 中):

```
In [24]:
          !echo \
          (declare-const x Real) \
          (declare-const y Real) \
          (declare-const z Real) \
          (assert (=(-(+(* 3 x) (* 2 y)) z) 1)) \setminus
          (assert (=(+(-(* 2 x) (* 2 y)) (* 4 z)) -2)) \
          (assert (=(-(+ (- 0 x) (* 0.5 y)) z) 0)) \
          (check-sat) \
          (get-model) | z3 -in
          sat
          (model
           (define-fun z () Real
            (-2.0))
           (define-fun y () Real
            (-2.0))
           (define-fun x () Real
             1.0)
```

另一个解方程的例子

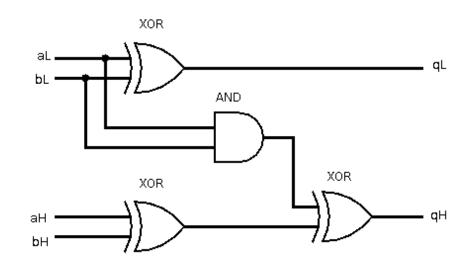
 $\nabla = \frac{1}{3}$

```
In [25]: circle, square, triangle = Ints('circle square triangle')
s = Solver()
s.add(circle + circle == 10)
s.add(circle * square + square == 12)
s.add(circle * square - triangle * circle == circle)
print(s.check())
if s.check() == sat:
    print(s.model())
```

sat

[triangle = 1, square = 2, circle = 5]

直接使用SAT的例子



上图中的2-bit 加法器电路没有进位,求使得输出为3(qL和qH均为1)的输入

容易想到有四种可能的输入: 0+3=3, 1+2=3, 2+1=3, 3+0=3

将该问题改写为CNF: $(\neg a \lor \neg c) \land (a \lor c) \land (\neg b \lor \neg d) \land (b \lor d)$

使用MiniSat解该问题:

p cnf 4 4 -1 -3 0 1 3 0 -2 -4 0 2 4 0

输出:

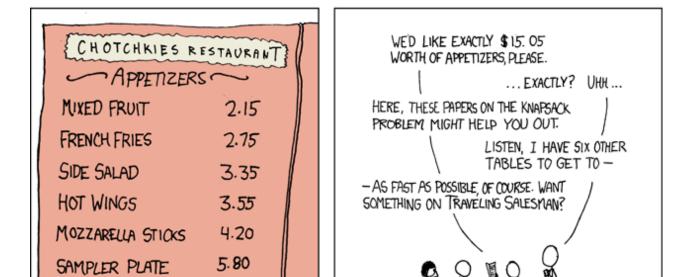
SAT与SMT相比,就像汇编语言与C语言.

更多示例 (第三章 Equations)

解XKCD 287

https://www.xkcd.com/287/

MY HOBBY: EMBEDDING NP-COMPLETE PROBLEMS IN RESTAURANT ORDERS



等价于 2.15a + 2.75b + 3.35c + 3.55d + 4.20e + 5.80f == 15.05 ,其中 $a \dots f$ 是整数.

这是一个丢番图方程(Diophantine Equation, 又称不定方程, 整系数多项式方程).

下面使用Z3进行求解(原文中使用了他自己开发的MK85):

```
In [27]: var = [BitVec(chr(i+ord('a')), 16) for i in range(6)]

s = Solver()
for i in var:
    s.add(ULE(i, 10))
s.add(var[0] * 215 + var[1] * 275 + var[2] * 335 + var[3] * 355 + var[4] * 420 + var[5] * 580 == 1505)

while s.check() == sat:
    m = s.model()
    print(m)
    # block current solution and solve again:

fmt = And(*[i == m.evaluate(i) for i in var])
    s.add(Not(fmt))
```

$$[f = 1, b = 0, a = 1, c = 0, d = 2, e = 0]$$

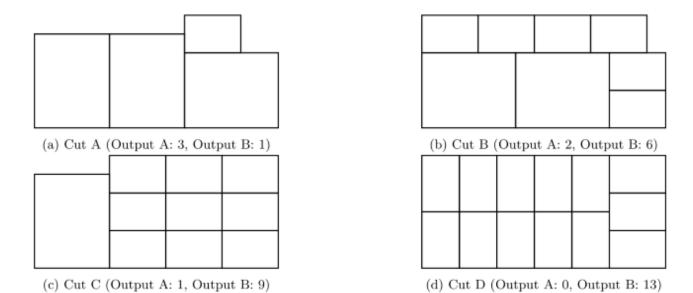
 $[f = 0, b = 0, a = 7, c = 0, d = 0, e = 0]$

XKCD 287的其他解法(各种奇奇怪怪的语言):

https://stackoverflow.com/questions/141779/solving-the-np-complete-problem-in-xkcd

优化问题: Wood workshop

有很多大小为6x13的木块,需加工为800个4x5的小木块和400个2x3的小木块. 有四种加工方式:



怎样切使用最少的木块?

```
In [28]: s = Optimize()
         workpieces_total = Int('workpieces_total')
         cut_A, cut_B, cut_C, cut_D = Ints('cut_A cut_B cut_C cut_D')
         out_A, out_B = Ints('out_A out_B')
         s.add(workpieces_total == cut_A + cut_B + cut_C + cut_D)
         s.add(cut_A >= 0)
         s.add(cut_B >= 0)
         s.add(cut C >= 0)
         s.add(cut_D >= 0)
         s.add(out_A == cut_A * 3 + cut_B * 2 + cut_C * 1)
         s.add(out_B == cut_A * 1 + cut_B * 6 + cut_C * 9 + cut_D * 13)
         s.add(out\_A == 800)
         s.add(out_B == 400)
         s.minimize(workpieces_total)
         print(s.check(), s.model())
         sat [cut_B = 25,
         cut_D = 0,
         cut A = 250,
```

out_B = 400, out_A = 800,

cut C = 0

workpieces_total = 275,

子集和问题(Subset sum)

给一个整数集合,问是否存在某个非空子集,使得子集内中的数字和为某个特定数值。

例:给定集合{-7,-3,-2,5,8},是否存在子集和为0的集合?答案是YES,因为子集{-3,-2,5}的数字和是0。

```
In [29]: set = [-7, -3, -2, 5, 8]
    var = [Int('var_%d' % i) for i in range(len(set))]
    s = Solver()
    rt = []
    for i, j in zip(var, set):
        rt.append(i * j)
        s.add(Or(i == 0, i == 1)) # like bools
        s.add(sum(rt) == 0)

s.add(sum(var) >= 1) # subset must not be empty
    if s.check() == sat:
        m = s.model()
        print([j for i, j in zip(var, set) if m[i].as_long() == 1])
    else:
        print("unsat")
```

[-3, -2, 5]

2017 AMC 12A Problems/Problem 2

The sum of two nonzero real numbers is 4 times their product. What is the sum of the reciprocals of the two numbers?

$$x + y = 4xy$$

$$\frac{1}{x} + \frac{1}{y} = ?$$

```
In [32]:  x, y = \text{Reals}('x y') 
 s = \text{Solver}() 
 s.add(x > 0) 
 s.add(x + y == 4 * x * y) 
 print(s.check()) 
 m = s.model() 
 print("the model: ", m) 
 print("the answer:", m.evaluate(1 / x + 1 / y)) 
 sat 
 the model: [x = 1, y = 1/3]
```

the answer: 4

模逆元

验证: 0x1234567

给定一个机器数(能被3整除),它乘以什么数的结果就像它除以3一样?

```
In [33]: m = BitVec('m', 32)
s = Solver()
divisor = 3 # wouldn't work for 10, etc
random_const = 0x1234567 # random constant, must be divisible by divisor
const = (random_const * divisor)
s.add(const * m == const / divisor)
print(s.check())
magic_number = s.model()[m].as_long()
print("Ox%x" % magic_number)

sat
Oxaaaaaaaab

In [34]: print('验证: Ox%x' % (const * magic_number & Oxffffffff))
```

扫雷游戏



WIDTH= 9 HEIGHT= 9

```
In [36]:
         def chk_bomb(row, col):
            s = Solver()
            cells = [[Int('r\%d\_c\%d' \% (r, c)) for c in range(WIDTH + 2)] for r in range(HEIGHT + 2)]
            for c in range(WIDTH + 2): # make border
               s.add(cells[0][c] == 0)
               s.add(cells[HEIGHT + 1][c] == 0)
            for r in range(HEIGHT + 2):
               s.add(cells[r][0] == 0)
               s.add(cells[r][WIDTH + 1] == 0)
            for r in range(1, HEIGHT + 1):
               for c in range(1, WIDTH + 1):
                  s.add(Or(cells[r][c] == 0, cells[r][c] == 1)) # otherwise -1 is possible, etc
                  t = known[r - 1][c - 1]
                  if t in "012345678":
                     s.add(cells[r][c] == 0)
                     # we need empty border so the following expression would be able to work for all possible cases:
                     expr = cells[r - 1][c - 1] + cells[r - 1][c] + cells[r - 1][c + 1] + cells[r][c - 1] + 
                          cells[r][c + 1] + cells[r + 1][c - 1] + cells[r + 1][c] + cells[r + 1][c + 1] == int(t)
                     s.add(expr)
            s.add(cells[row][col] == 1) # place bomb
            if s.check() == unsat:
               print("row=%d col=%d, unsat!" % (row, col))
```

```
In [37]: # enumerate all hidden cells:
for r in range(1, HEIGHT + 1):
    for c in range(1, WIDTH + 1):
        if known[r - 1][c - 1] == "?":
            chk_bomb(r, c)

row=1 col=3, unsat!
row=6 col=2, unsat!
row=6 col=3, unsat!
```

row=7 col=4, unsat! row=7 col=9, unsat! row=8 col=9, unsat! 输出unsat的坐标是没有地雷的,在游戏中点击后将结果再传给Z3求解



破解线性同余伪随机数生成器(LCG)

LCG非常高效,使用广泛,如C语言的 rand() 函数.

 $\mathsf{LCG定义}: X_{n+1} = (aX_n + c) \mod m$

但它的安全性非常低,我们使用Z3来破解LCG.

Source	modulus <i>m</i>	multiplier <i>a</i>	increment c
Numerical Recipes	232	1664525	1013904223
Borland C/C++	232	22695477	1
III (II CCC)[15]	0.21	4402545245	10045

```
In [39]: import time
         modulus = 2 ** 31
         a = 1103515245
         c = 12345
         def lcg(seed):
            """Linear congruential generator."""
           while True:
              seed = (a * seed + c) % modulus
              yield (seed >> 16) & 0x7fff # 提高破解效率
         generator = lcg(int(time.time()))
         rand_var = []
         rand_state = []
         for g, _ in zip(generator, range(10)):
           rand_var.append(g % 100)
           rand_state.append(g)
         print('random var:', rand_var)
         print('random state:', rand_state)
```

random var: [97, 17, 49, 55, 8, 42, 14, 58, 68, 28] random state: [19497, 8317, 27149, 20155, 21308, 15142, 29814, 22358, 30868, 15328]

```
In [40]:
         state = [BitVec('state%d' % i, 32) for i in range(10)]
          output_prev = BitVec('output_prev', 32)
          output next = BitVec('output next', 32)
          s = Solver()
         for i in range(9):
            s.add(state[i + 1] == ((state[i] * a + c) % modulus))
         for i in range(1, 9):
            s.add(URem((state[i] >> 16) \& 0x7fff, 100) == rand_var[i])
          s.add(output prev == URem((state[0] >> 16) \& 0x7fff, 100))
          s.add(output_next == URem((state[9] >> 16) & 0x7fff, 100))
          if s.check() == sat:
            m = s.model()
            print('cracked state', [m.evaluate((state[i] >> 16) & 0x7fff) for i in range(10)])
            print('predicate first:', m.evaluate(output_prev), 'last:', m.evaluate(output_next))
          else:
            print('unsat')
```

cracked state [19497, 8317, 27149, 20155, 21308, 15142, 29814, 22358, 30868, 15328] predicate first: 97 last: 28

这个LCG能产生4个连续的0吗?(运行较慢)

```
In [41]: s = Solver()
for i in range(4):
    s.add(state[i + 1] == ((state[i] * a + c) % modulus))

for i in range(1, 5):
    s.add(URem((state[i] >> 16) & 0x7fff, 100) == 0)

print(s.check())
if s.check() == sat:
    print(s.model())

sat
[state3 = 3228838068,
    state1 = 2789793222,
    state2 = 3451711879,
    state0 = 3538020257,
    state4 = 2232704221]
```

因式分解(运行较慢)

用这种方法可以破解RSA (规模较小,比如位数少于 2^{40})

```
In [21]: import random
         from functools import reduce
         from operator import mul
         def factor(n):
            print("factoring", n)
            in1, in2, out = Ints('in1 in2 out')
            s = Solver()
            s.add(out == n, in1 * in2 == out, in1 > 1, in2 > 1)
            if s.check() == unsat:
               print(n, "is prime (unsat)")
               return [n]
            if s.check() == unknown:
               print(n, "is probably prime (unknown)")
               return [n]
            m = s.model()
            # get inputs of multiplier:
            in1_n = m[in1].as_long()
            in2_n = m[in2].as_long()
            print("factors of", n, "are", in1_n, "and", in2_n)
            # factor factors recursively:
            rt = sorted(factor(in1_n) + factor(in2_n))
            # self-test:
            assert (reduce(mul, rt, 1) == n)
            return rt
```

In [20]: print(factor(1234567890))

factoring 1234567890 factors of 1234567890 are 11409 and 108210 factoring 11409 factors of 11409 are 3803 and 3 factoring 3803 3803 is prime (unsat) factoring 3 3 is prime (unsat) factoring 108210 factors of 108210 are 15 and 7214 factoring 15 factors of 15 are 5 and 3 factoring 5 5 is prime (unsat) factoring 3 3 is prime (unsat) factoring 7214 factors of 7214 are 3607 and 2 factoring 3607 3607 is prime (unsat) factoring 2 2 is prime (unsat) [2, 3, 3, 5, 3607, 3803]

电子表格公式计算

	0	1	2	3
A	1	0	B0+B2	A0*B0*C0
В	123	10	12	11
С	667	A0+B1	C1*A0*122	A3+C2

s.unsat_core()

如果数据出现循环依赖,则不能求解.而直接给用户返回unsat非常不友好,可以通过unsat_core返回更多信息(退出幻灯片查看下面的例子).

```
In [15]: import sys, re
         # MS Excel or LibreOffice style.
         # first top-left cell is AO, not A1
         def coord_to_name(R, C):
            return "ABCDEFGHIJKLMNOPQRSTUVWXYZ"[R] + str(C)
         ar = [['1', '0', 'B0+B2', 'A0*B0*C0'],
             ['123', '10', '12', '11'],
             ['667', 'A0+B1', 'C1*A0*122', 'A3+C2']]
         \# ar = [['1', '0', 'B0+B2', 'A0*B0*C0'],
         # ['123', '10', '12', '11'],
              ['C1+123', 'C0*123', 'C1*A0*122', 'A3+C2']]
         WIDTH = len(ar[0])
         HEIGHT = Ien(ar)
         s = Solver()
         s.set(unsat_core=True)
         # cells{} is a dictionary with keys like "AO", "B9", etc:
         cells = \{\}
         for R in range(HEIGHT):
            for C in range(WIDTH):
              name = coord_to_name(R, C)
              cells[name] = BitVec(name, 32)
```

```
In [16]: cur_R = 0
          cur C = 0
          for row in ar:
            for c in row:
                # string like "A0+B2" becomes "cells["A0"]+cells["B2"]":
                c = re.sub(r'([A-Z]{1}[0-9]+)', r'cells["\1"]', c)
                st = "cells[\"%s\"]==%s" % (coord_to_name(cur_R, cur_C), c)
                # evaluate string. Z3Py expression is constructed at this step:
                e = eval(st)
                # add constraint:
              s.add(e) # 如果没使用unsat_core
                s.assert_and_track(e, coord_to_name(cur_R, cur_C)) # 如果使用了unsat_core
                \operatorname{cur} C = \operatorname{cur} C + 1
             \operatorname{cur} R = \operatorname{cur} R + 1
             cur_C = 0
          if s.check() == sat:
             m = s.model()
             for r in range(HEIGHT):
                for c in range(WIDTH):
                   print(m.evaluate(cells[coord_to_name(r, c)]), end="\t")
                print()
          else:
             print('unsat', s.unsat_core())
                     135 82041
```

123 10 12 11 667 11 1342

83383

物理问题

一辆车的速度为30m/s,红灯亮了,刹车加速度为 $-8m/s^2$,求开始刹车到停止的位移.

运动方程:

$$d=v_i*t+rac{1}{2}*a*t^2 \ v_f=v_i+a*t$$

```
In [22]:

d, a, t, v_i, v_f = Reals('d a t v_i v_f')

equations = [ d == v_i * t + (a*t**2)/2, v_f == v_i + a*t ]

print("Kinematic equations:")

print(equations)

# Given v_i, v_f and a, find d

problem = [ v_i == 30, v_f == 0, a == -8]

print("Problem:")

print(problem)

print("Solution:")

solve(equations + problem)
```

Kinematic equations:

[d ==
$$v_i*t + (a*t**2)/2$$
, $v_f == v_i + a*t$]
Problem:
[$v_i == 30$, $v_f == 0$, $a == -8$]
Solution:
[$a = -8$, $v_f = 0$, $v_i = 30$, $t = 15/4$, $d = 225/4$]

Sudoku

每一行/每一列/小方框不能出现相同的数字

				9	4		3	
			5	1				7
	8	9					4	
						2		8
	6		2		1		5	
1		2						
	7					5	2	
9				6	5			
	4		9	7				

```
In [25]:
          # 9x9 matrix of integer variables
         X = [ [Int("x_%s_%s" \% (i+1, j+1)) for j in range(9) ] for i in range(9) ]
          # each cell contains a value in {1, ..., 9}
          cells_c = [And(1 \le X[i][j], X[i][j] \le 9) for i in range(9) for j in range(9)]
          # each row contains a digit at most once
          rows_c = [ Distinct(X[i]) for i in range(9) ]
          # each column contains a digit at most once
          cols_c = [ Distinct([ X[i][j] for i in range(9) ]) for j in range(9) ]
          # each 3x3 square contains a digit at most once
          sq_c = [Distinct([X[3*i0 + i][3*j0 + j] for i in range(3) for j in range(3)])
                  for iO in range(3) for jO in range(3) ]
          sudoku c = cells c + rows c + cols c + sq c
          # sudoku instance, we use '0' for empty cells
          instance = ((0,0,0,0,9,4,0,3,0),
                  (0,0,0,5,1,0,0,0,7),
                  (0,8,9,0,0,0,0,4,0),
                  (0,0,0,0,0,0,2,0,8),
                  (0,6,0,2,0,1,0,5,0),
                  (1,0,2,0,0,0,0,0,0),
                  (0,7,0,0,0,0,5,2,0),
                  (9,0,0,0,6,5,0,0,0),
                  (0,4,0,9,7,0,0,0,0)
          instance_c = [If(instance[i][j] == 0, True, X[i][j] == instance[i][j])
                    for i in range(9) for j in range(9) ]
```

```
In [27]: 
s = Solver()
s.add(sudoku_c + instance_c)
if s.check() == sat:
    m = s.model()
    r = [ [ m.evaluate(X[i][j]) for j in range(9) ]
        for i in range(9) ]
        print_matrix(r)
else:
    print("failed to solve")

[[7, 1, 5, 8, 9, 4, 6, 3, 2],
[2, 3, 4, 5, 1, 6, 8, 9, 7],
[6, 8, 9, 7, 2, 3, 1, 4, 5],
[4, 9, 3, 6, 5, 7, 2, 1, 8],
```

[8, 6, 7, 2, 3, 1, 9, 5, 4], [1, 5, 2, 4, 8, 9, 7, 6, 3], [3, 7, 6, 1, 4, 8, 5, 2, 9], [9, 2, 8, 3, 6, 5, 4, 7, 1], [5, 4, 1, 9, 7, 2, 3, 8, 6]]

```
In [29]:
          # Let us remove 9 from the first row and see if there is more than one solution
          instance = ((0,0,0,0,0,4,0,3,0),
                  (0,0,0,5,1,0,0,0,7),
                  (0,8,9,0,0,0,0,4,0),
                  (0,0,0,0,0,0,2,0,8),
                  (0,6,0,2,0,1,0,5,0),
                  (1,0,2,0,0,0,0,0,0),
                  (0,7,0,0,0,0,5,2,0),
                  (9,0,0,0,6,5,0,0,0),
                  (0,4,0,9,7,0,0,0,0))
          instance_c = [If(instance[i][j] == 0, True, X[i][j] == instance[i][j])
                    for i in range(9) for j in range(9) ]
          def n_solutions(n):
            s = Solver()
            s.add(sudoku_c + instance_c)
            i = 0
            while s.check() == sat and i < n:
               m = s.model()
               print([[ m.evaluate(X[i][j]) for j in range(9)] for i in range(9)])
               fml = And([X[i][j] == m.evaluate(X[i][j]) for i in range(9) for j in range(9)])
               s.add(Not(fml))
               i += 1
```

```
6, 7, 2, 9, 1, 4, 5, 3], [1, 5, 2, 4, 8, 3, 7, 9, 6], [6, 7, 1, 3, 4, 8, 5, 2, 9], [9, 2, 8, 1, 6, 5, 3, 7, 4], [3, 4,
5. 9. 7. 2. 8. 6. 1]]
6, 7, 2, 3, 1, 9, 5, 4], [1, 5, 2, 4, 8, 9, 7, 6, 3], [3, 7, 6, 1, 4, 8, 5, 2, 9], [9, 2, 8, 3, 6, 5, 4, 7, 1], [5, 4,
1, 9, 7, 2, 3, 8, 6]]
[[5, 1, 7, 6, 8, 4, 9, 3, 2], [2, 3, 4, 5, 1, 9, 6, 8, 7], [6, 8, 9, 7, 2, 3, 1, 4, 5], [7, 9, 3, 4, 5, 6, 2, 1, 8], [4,
6, 8, 2, 9, 1, 7, 5, 3], [1, 5, 2, 8, 3, 7, 4, 9, 6], [3, 7, 6, 1, 4, 8, 5, 2, 9], [9, 2, 1, 3, 6, 5, 8, 7, 4], [8, 4,
5, 9, 7, 2, 3, 6, 1]]
[[5, 1, 6, 7, 8, 4, 9, 3, 2], [2, 3, 4, 5, 1, 9, 6, 8, 7], [7, 8, 9, 6, 2, 3, 1, 4, 5], [3, 9, 7, 4, 5, 6, 2, 1, 8], [4,
6, 8, 2, 9, 1, 7, 5, 3], [1, 5, 2, 8, 3, 7, 4, 9, 6], [6, 7, 3, 1, 4, 8, 5, 2, 9], [9, 2, 1, 3, 6, 5, 8, 7, 4], [8, 4,
5, 9, 7, 2, 3, 6, 1]]
6, 8, 2, 9, 1, 7, 5, 3], [1, 5, 2, 7, 8, 3, 4, 9, 6], [6, 7, 3, 1, 4, 8, 5, 2, 9], [9, 2, 1, 3, 6, 5, 8, 7, 4], [8, 4,
5. 9. 7. 2. 3. 6. 111
6. 8. 2. 9. 1. 7. 5. 3], [1, 5, 2, 7, 8, 3, 4, 9, 6], [3, 7, 6, 1, 4, 8, 5, 2, 9], [9, 2, 1, 3, 6, 5, 8, 7, 4], [8, 4,
5, 9, 7, 2, 3, 6, 1]]
[[6, 1, 7, 8, 2, 4, 9, 3, 5], [2, 3, 4, 5, 1, 9, 6, 8, 7], [5, 8, 9, 6, 3, 7, 1, 4, 2], [4, 9, 3, 7, 5, 6, 2, 1, 8], [7,
6. 8. 2. 9. 1. 4. 5. 3]. [1. 5. 2. 4. 8. 3. 7. 9. 6]. [3. 7. 6. 1. 4. 8. 5. 2. 9]. [9. 2. 1. 3. 6. 5. 8. 7. 4]. [8. 4.
5. 9. 7. 2. 3. 6. 111
[[6, 1, 7, 8, 2, 4, 9, 3, 5], [2, 3, 4, 5, 1, 9, 6, 8, 7], [5, 8, 9, 7, 3, 6, 1, 4, 2], [4, 9, 3, 6, 5, 7, 2, 1, 8], [7,
6. 8. 2. 9. 1. 4. 5. 3]. [1. 5. 2. 4. 8. 3. 7. 9. 6]. [3. 7. 6. 1. 4. 8. 5. 2. 9]. [9. 2. 1. 3. 6. 5. 8. 7. 4]. [8. 4.
5. 9. 7. 2. 3. 6. 1]]
6, 8, 2, 9, 1, 3, 5, 4], [1, 5, 2, 4, 3, 8, 7, 9, 6], [8, 7, 6, 1, 4, 3, 5, 2, 9], [9, 2, 1, 8, 6, 5, 4, 7, 3], [3, 4,
5. 9. 7. 2. 8. 6. 1]]
6. 8. 2. 9. 1. 7. 5. 3]. [1, 5, 2, 7, 3, 8, 4, 9, 6], [8, 7, 6, 1, 4, 3, 5, 2, 9], [9, 2, 1, 8, 6, 5, 3, 7, 4], [5, 4,
3. 9, 7, 2, 8, 6, 1]]
```

The end

感谢聆听!