The DPLL algorithm

Combinatorial Problem Solving (CPS)

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Overview of the session

- Designing an efficient SAT solver
- DPLL: A Bit of History
- Abstract DPLL:
 - ◆ Rules
 - **♦** Examples
 - ◆ Theoretical Results

Designing an efficient SAT solver

INPUT: formula F in **CNF**

OUTPUT:

- If F is SAT: YES + model
- If F is UNSAT: NO + refutation (proof of unsatisfiability)

Two possible methods:

- resolution-based:
 - not direct to obtain model
 - straightforward to give refutation
- DPLL-based:
 - + straightforward to obtain model
 - not direct to give refutation

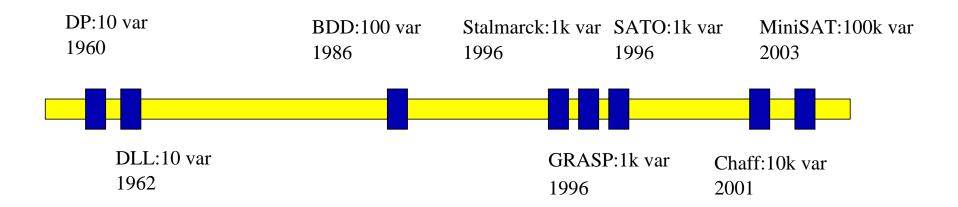
Due to their efficiency, DPLL-based solvers are the method of choice

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DPLL - A Bit of History

- Original DPLL was incomplete method for FOL satisfiability
- First paper (Davis and Putnam) in 1960: memory problems
- Second paper (Davis, Logemann and Loveland) in 1962: Depth-first-search with backtracking
- Late 90's and early 00's improvements make DPLL efficient:
 - Break-through systems: GRASP, SATO, Chaff, MiniSAT



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Our Abstraction of DPLL

- lacksquare Given F in CNF, DPLL tries to build assignment M s.t. $M \models F$
- Assignments M are represented as sequences of literals (those to be true): **EXAMPLE:** sequence $p\overline{q}r$ is M(p)=1, M(q)=0, M(r)=1 (overlining bar $\bar{}$ may be used to represent negation, like $\bar{}$)
 - lacktriangle Order in M matters
 - lacktriangle No literal appears twice in M
 - lacktriangle No contradictory literals in M
- \blacksquare Sequences may have decision literals, denoted l^{d} .
- We will introduce a transition system modelling DPLL
- States in the transition system are pairs $M \parallel F$, where M is a (partial) assignment and F is a CNF
- The algorithm starts with an empty assignment (and the original formula)
- The transition rules in the transition system indicate which steps $M \parallel F \Longrightarrow M' \parallel F'$ are allowed.

Abstract DPLL - Rules

Extending the assignment:

Decide

$$M \parallel F \implies M l^{\mathsf{d}} \parallel F \text{ if } \left\{ egin{array}{l} l \text{ or } \overline{l} \text{ occurs in } F \\ l \text{ is undefined in } M \end{array} \right.$$

UnitProp

$$M \parallel F, C \lor l \implies M \ l \parallel F, C \lor l$$
 if $\left\{ \begin{array}{l} M \models \neg C \\ l \text{ is undefined in } M \end{array} \right.$

Abstract DPLL - Rules (2)

Repairing the assignment:

Fail

Backtrack

$$M \ l^{\operatorname{d}} \ N \parallel F, \ C \implies M \ \overline{l} \parallel F, \ C \ \ \mathbf{if} \ \left\{ \begin{array}{l} M \ l^{\operatorname{d}} \ N \models \neg C \\ N \ \ \text{contains no decision lits} \end{array} \right.$$

 $\emptyset \parallel \overline{1} \vee 2, \overline{3} \vee 4, \overline{5} \vee \overline{6}, 6 \vee \overline{5} \vee \overline{2} \Longrightarrow$

 $\emptyset \parallel \overline{1} \vee 2, \overline{3} \vee 4, \overline{5} \vee \overline{6}, 6 \vee \overline{5} \vee \overline{2} \Longrightarrow (\mathsf{Decide})$

 $\emptyset \parallel \overline{1} \vee 2 \vee 3, \ 1, \ \overline{2} \vee 3, \ \overline{2} \vee \overline{3}, \ 2 \vee 3, \ 2 \vee \overline{3} \Longrightarrow$

 $\emptyset \parallel \overline{1} \vee 2 \vee 3, \ 1, \ \overline{2} \vee 3, \ \overline{2} \vee \overline{3}, \ 2 \vee 3, \ 2 \vee \overline{3} \qquad \Longrightarrow \qquad (\mathsf{UnitProp})$

Abstract DPLL

- lacktriangle There are no infinite sequences of the form $\emptyset \parallel F \implies \dots$
- \blacksquare If $\emptyset \parallel F \Longrightarrow^* M \parallel F$ with state $M \parallel F$ final, then
 - lacktriangleright F is satisfiable
 - lacktriangledown M is a model of F
- If $\emptyset \parallel F \implies^* fail$ then F is unsatisfiable

Hence the transition system gives a decision procedure for SAT

Bibliography - Some further reading

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