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MODULE CC -
 1 [
      TLA+ specification of Causal Consistency variants, including CC, CM, and CCv.
      See the paper "On Verifying Causal Consistency" (POPL'2017).
 8 EXTENDS Naturals, Sequences, Functions, FiniteSets, FiniteSetsExt, RelationUtils, TLC
    CONSTANTS Keys, Vals
     InitVal \stackrel{\triangle}{=} CHOOSE \ v : v \notin (Keys \cup Vals)
      oid: unique operation identifier
13
     Operation \stackrel{\Delta}{=} [op : \{ \text{"read"}, \text{"write"} \}, key : Keys, val : Vals, oid : Nat] \}
    R(k, v, oid) \stackrel{\square}{=} [op \mapsto \text{``read''}, key \mapsto k, val \mapsto v, oid \mapsto oid]
     W(k, v, oid) \triangleq [op \mapsto "write", key \mapsto k, val \mapsto v, oid \mapsto oid]
     Session \stackrel{\triangle}{=} Seq(Operation) A session s \in Session is a sequence of operations.
     History \stackrel{\triangle}{=} SUBSET Session A history h \in History is a set of sessions.
19
20 F
      Utilities.
     Ops(h) \stackrel{\triangle}{=} Return the set of all operations in history <math>h \in History.
24
       UNION \{Range(s): s \in h\}
25
26 |
      Well-formedness of history h \in History:
       - TODO: type invariants
      - uniqueness of oids
     WellFormed(h) \triangleq
33
      \land h \in History
34
       \wedge Cardinality(Ops(h)) = ReduceSet(LAMBDA s, x : Len(s) + x, h, 0)
35
36 F
      Program order: a union of total orders among operations in the same session.
     ProgramOrder(h) \triangleq
40
       LET RECURSIVE SessionProgramOrder(_)
41
             SessionProgramOrder(s) \triangleq
42
               If s = \langle \rangle then \{\}
43
                 ELSE LET sh \stackrel{\triangle}{=} Head(s)
                                 st \stackrel{\triangle}{=} Tail(s)
45
                                \{\langle sh, t \rangle : t \in Range(st)\} \cup SessionProgramOrder(st)
46
            UNION \{SessionProgramOrder(s) : s \in h\}
47
48
      Sequential semantics of read-write registers.
52 F
      Specification of Causal Consistency: CC, CCv, and CM
     CCv(h) \stackrel{\triangle}{=} Check whether h \in History satisfies CCv (Causal Convergence)
       \land WellFormed(h)
57
        \wedge LET ops \stackrel{\triangle}{=} Ops(h)
58
                \land \exists co \in \text{SUBSET} (ops \times ops) :
59
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\exists arb \in SUBSET (ops \times ops) :
 60
                            \land IsStrictPartialOrder(co, ops)
 61
                            \land IsStrictTotalOrder(arb, ops)
 62
                            \land Respect(co, ProgramOrder(h))
                                                                          AxCausal
 63
                            \land Respect(arb, co)
                                                                           AxArb
 64
                            \land \forall op \in ops : TRUE
                                                                           TODO: AxCausalArb
 65
         \wedge FALSE
 66
 67 H
        Test case: The following histories are from Figure 2 of the POPL'2017 paper.
        Naming Conventions:
        -ha: history of Figure 2(a)
        - hasa: session a of history ha
        TODO: to automatically generate histories
      hasa \triangleq \langle W(\text{``x''}, 1, 1), R(\text{``x''}, 2, 2) \rangle
      hasb \triangleq \langle W("x", 2, 3), R("x", 1, 4) \rangle
      ha \stackrel{\Delta}{=} \{hasa, hasb\} CM but not CCv
      hbsa \triangleq \langle W(\text{"z"}, 1, 1), W(\text{"x"}, 1, 2), W(\text{"y"}, 1, 3) \rangle
      hbsb \triangleq \langle W(\text{"x"}, 2, 4), R(\text{"z"}, 0, 5), R(\text{"y"}, 1, 6), R(\text{"x"}, 2, 7) \rangle
      hb \stackrel{\Delta}{=} \{hbsa, hbsb\} CCv but not CM
     hcsa \triangleq \langle W("x", 1, 1) \rangle
      hcsb \triangleq \langle W("x", 2, 2), R("x", 1, 3), R("x", 2, 4) \rangle
      hc \stackrel{\Delta}{=} \{hcsa, hcsb\} CC but not CM nor CCv
     hdsa \triangleq \langle W("x", 1, 1), R("y", 0, 2), W("y", 1, 3), R("x", 1, 4) \rangle
     hdsb \triangleq \langle W("x", 2, 5), R("y", 0, 6), W("y", 2, 7), R("x", 2, 8) \rangle
      hd \stackrel{\Delta}{=} \{hdsa, hdsb\}\ CC, CM, \text{ and } CCv \text{ but no } SC
     hesa \triangleq \langle W("x", 1, 1), W("y", 1, 2) \rangle
     hesb \stackrel{\triangle}{=} \langle R("y", 1, 3), W("x", 2, 4) \rangle
     hesc \triangleq \langle R("x", 2, 5), R("x", 1, 6) \rangle
      he \stackrel{\triangle}{=} \{hesa, hesb, hesc\} \text{ not } CC \text{ (nor } CM, \text{ nor } CCv)
      THEOREM WellFormedTheorem \stackrel{\triangle}{=}
        \forall h \in \{ha, hb, hc, hd, he\} : WellFormed(h)
100
      CardOfProgramOrderOfHistory(h) \triangleq
102
        LET CardOfProgramOrderOfSession(s) \stackrel{\triangle}{=}
103
           IF Len(s) < 1 THEN 0 ELSE Sum(1 ... Len(s) - 1)
104
                ReduceSet(LAMBDA\ s,\ x: CardOfProgramOrderOfSession(s) + x,\ h,\ 0)
105
     THEOREM Program Order Cardinality Theorem \stackrel{\triangle}{=}
107
        \forall h \in \{ha, hb, hc, hd, he\}:
108
           Cardinality(ProgramOrder(h)) = CardOfProgramOrderOfHistory(h)
109
110
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- $\begin{tabular}{ll} $$ \ $\ $$ Modification $History$ \\ $$ $$ $$ $$ Apr 03 22:21:24 $CST 2021 by $hengxin$ \\ $$ $$ $$ $$ $$ Created $Tue Apr 01 10:24:07 $CST 2021 by $hengxin$ \\ $$ $$$