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- module CC
        TLA+ specification of Causal Consistency variants, including CC, CM, and CCv.
        See the paper "On Verifying Causal Consistency" (POPL'2017).
     EXTENDS TLC, Sequences, FiniteSets, FunctionUtils
      Constants Keys, Vals
      InitVal \stackrel{\triangle}{=} CHOOSE \ v : v \notin (Keys \cup Vals)
      Operation \triangleq [op : \{ \text{"read"}, \text{"write"} \}, key : Keys, val : Vals]
      R(k, v) \stackrel{\triangle}{=} [op \mapsto \text{``read''}, key \mapsto k, val \mapsto v]
      W(k, v) \triangleq [op \mapsto \text{"write"}, key \mapsto k, val \mapsto v]
      Session \stackrel{\triangle}{=} Seg(Operation) A session s \in Session is a sequence of operations.
     History \stackrel{\triangle}{=} SUBSET Session A history <math>h \in History is a set of sessions.
        Test case: The following histories are from Figure 2 of the POPL'2017 paper.
        Naming:
         -ha: history of Figure 2(a)
        - hasa: session a of history ha
     hasa \stackrel{\triangle}{=} \langle W("x", 1), R("x", 2) \rangle
      hasb \triangleq \langle W(\text{"x"}, 2), R(\text{"x"}, 1) \rangle
      ha \stackrel{\Delta}{=} \{hasa, hasb\} CM but not CCv
     hbsa \stackrel{\Delta}{=} \langle W(\text{``z''}, 1), W(\text{``x''}, 1), W(\text{``y''}, 1) \rangle
     hbsb \triangleq \langle W(\text{"x"}, 2), R(\text{"z"}, 0), R(\text{"y"}, 1), R(\text{"x"}, 2) \rangle
     hb \stackrel{\Delta}{=} \{hbsa, hbsb\} CCv but not CM
     hcsa \stackrel{\triangle}{=} \langle W(\text{"x"}, 1) \rangle
     hcsb \triangleq \langle W(\text{"x"}, 2), R(\text{"x"}, 1), R(\text{"x"}, 2) \rangle
      hc \stackrel{\Delta}{=} \{hcsa, hcsb\}\ CC \text{ but not } CM \text{ nor } CCv
     hdsa \triangleq \langle W("x", 1), R("y", 0), W("y", 1), R("x", 1) \rangle
     hdsb \stackrel{\triangle}{=} \langle W(\text{"x"}, 2), R(\text{"y"}, 0), W(\text{"y"}, 2), R(\text{"x"}, 2) \rangle
      hd \stackrel{\Delta}{=} \{hdsa, hdsb\}\ CC, CM, \text{ and } CCv \text{ but no } SC
     \begin{array}{l} hesa \; \stackrel{\triangle}{=} \; \langle W(\text{``x''},1), \; W(\text{``y''},1) \rangle \\ hesb \; \stackrel{\triangle}{=} \; \langle R(\text{``y''},1), \; W(\text{``x''},2) \rangle \\ hesc \; \stackrel{\triangle}{=} \; \langle R(\text{``x''},2), \; R(\text{``x''},1) \rangle \end{array}
      he \stackrel{\triangle}{=} \{hesa, hesb, hesc\} \text{ not } CC \text{ (nor } CM, \text{ nor } CCv)
50 F
       Program order: a union of total orders among operations in the same session.
      Program Order(h) \stackrel{\triangle}{=}
54
         LET RECURSIVE SessionProgramOrder(_)
55
                 SessionProgramOrder(s) \triangleq
56
                    IF s = \langle \rangle THEN \{\}
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\begin{array}{ccc} \text{ELSE} & \text{LET} & sh & \stackrel{\triangle}{=} & Head(s) \\ & st & \stackrel{\triangle}{=} & Tail(s) \end{array}
58
59
                                   \{\langle sh, t \rangle : t \in Range(st)\} \cup SessionProgramOrder(st)
60
             UNION \{SessionProgramOrder(s) : s \in h\}
61
       Test case: TODO: Cardinality testing
       CardOfProgramOrderOf(h) \stackrel{\triangle}{=}
66
       Theorem CardOfProgramOrderTheorem \stackrel{\Delta}{=}
67
         \forall h \in \{ha, hb, hc, hd, he\}:
68
           Cardinality(ProgramOrder(h)) = CardOfProgramOrderOf(h)
69
70 F
      Sequential semantics of read-write registers.
74 F
      Utilities.
       FIXME: there may be multiple same operations in one and more sessions.
     Ops(h) \stackrel{\triangle}{=} Return the set of all operations in history <math>h \in History.
       UNION \{Range(s): s \in h\}
81
82 |-
      Specification of Causal Consistency: CC, CCv, and CM
     CCv(h) \stackrel{\triangle}{=} Check whether h \in History satisfies CCv (Causal Convergence)
86
        \wedge LET ops \stackrel{\triangle}{=} Ops(h)
87
                 \land \exists co \in \text{SUBSET} (ops \times ops) :
88
                        \exists arb \in SUBSET (ops \times ops) :
89
                           \forall op \in ops : TRUE
90
         \wedge FALSE
91
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